CS3334 Data Structures (Suggested Solutions to Extra Exercises)

1. [Total 10 marks]

- 2. [Total 20 marks]
 - (a) [5 marks]
 The functions in non-decreasing order:

$$\sqrt{n}$$
, $n \log n$, $n^2 \log n$, n^3 , 2^n

- (b) [5 marks] $f(n) = O(3n^2) = O(n^2)$
- (c) [10 marks]

$$g(n) = n + \frac{(n-1)(n)}{2} + 10$$
$$= n + \frac{n^2}{2} - \frac{n}{2} + 10$$
$$= O(n^2)$$

(Deduct 2 marks for each error/mistake in the derivation)

3. [Total 25 marks]

(a) [10 marks] Let P be the statement:

$$sum = \sum_{j=0}^{i-1} A[j]^2$$

3 marks: (Base case) When the program execution comes to the loop test for the first time, variable i=0 and sum=0. By the given notation, $\sum_{j=0}^{-1} A[j]^2 = 0$. That is, the statement P claims that sum=0 which is true.

7 marks: (Induction step) Assume P is true when the program execution comes to the loop test for the k-th time for some $k \geq 1$. Then, $sum = \sum_{j=0}^{i-1} A[j]^2$ at that moment.

Suppose the loop test succeeds and the loop body is executed. After executing the statement "sum+=Array[i]*Array[i];", $sum=\sum_{j=0}^{i-1}A[j]^2+A[i]^2=\sum_{j=0}^{i}A[j]^2$. After executing the statement "i++;", $\sum_{j=0}^{i-1}A[j]^2$. Therefore, P is true again when the program execution comes to the loop test for the (k+1)-st time.

By the principle of mathematical induction, P is a loop invariant at the loop test.

(b) [5 marks]

3 marks: Proof of termination: Initially, i = 0. Each execution of the loop body increases i by 1. Eventually, i will become n and at that time, the loop test: i < n will fail. So, the loop will terminate.

2 marks: Proof of total correctness: At the last loop test, i.e., when the loop test fails, i = n. By the loop invariant P, the variable sum contains the value $\sum_{j=0}^{n-1} A[j]^2$.

(c) [10 marks]

2 marks: The initialization of sum and i takes 2 units of time.

2 marks: The loop test is done n+1 times, costing n+1 units of time.

2 marks: The statement i++ is executed n times.

3 marks: The loop body is executed n times. Each execution of the loop body consists of one addition (+), one assignment (=), and one multiplication (*). Thus, each execution takes 3 units of time.

1 mark: Therefore, the worst case time complexity of the function is T(n) = 2 + (n+1) + n + 3n = 5n + 3 = O(n).

2

4. [Total 10 marks]

4 marks: Let T(n) be the worst case time complexity for the function. The function performs at most one recursive call of size n-1. It also takes constant time to perform the local work. Therefore, $T(n) \leq T(n-1) + c$.

3 marks: Solving the recurrence formula:

$$T(n) \leq T(n-1) + c$$

$$T(n-1) \leq T(n-2) + c$$

$$T(n-2) \leq T(n-3) + c$$

$$\cdots$$

$$T(1) \leq T(0) + c.$$

Summing them up, we have

$$T(n) \leq T(0) + nc.$$

where T(0) is a constant.

3 marks: Therefore, T(n) = O(n).

5. [Total 5 marks] There are 4 inversions

