CS3402: Chapter 6 SQL: Structured Query Language II

The SELECT-FROM-WHERE Structure of Basic SQL Queries

- SELECT statement
 - ◆ The basic statement for retrieving information from a database

```
SELECT <attribute list>
FROM 
WHERE <condition>;

where

SELECT <attribute list>
FROM 
[WHERE <condition>]

[ORDER BY <attribute list>];
```

- <attribute list> is a list of attribute names whose values are to be retrieved by the query.
- is a list of the relation names required to process the query.
- <condition> is a conditional (Boolean) expression that identifies the tuples to be retrieved by the query.

| Specified at | tributes |
|--------------|----------|
|--------------|----------|

Satisfy the conditions

| | | V | |
|---------------|------|--------|-------|
| Title | Year | Length | Туре |
| Star War | 1977 | 124 | Color |
| Mighty Duck | 1991 | 104 | Color |
| Wayne's World | 1992 | 95 | Color |

One possible database state for the COMPANY relational database schema

EMPLOYEE

| Fname | Minit | Lname | Ssn | Bdate | Address | Sex | Salary | Super_ssn | Dno |
|----------|-------|---------|-----------|------------|--------------------------|-----|--------|-----------|-----|
| John | В | Smith | 123456789 | 1965-01-09 | 731 Fondren, Houston, TX | М | 30000 | 333445555 | 5 |
| Franklin | Т | Wong | 333445555 | 1955-12-08 | 638 Voss, Houston, TX | | 40000 | 888665555 | 5 |
| Alicia | J | Zelaya | 999887777 | 1968-01-19 | 3321 Castle, Spring, TX | F | 25000 | 987654321 | 4 |
| Jennifer | S | Wallace | 987654321 | 1941-06-20 | 291 Berry, Bellaire, TX | F | 43000 | 888665555 | 4 |
| Ramesh | K | Narayan | 666884444 | 1962-09-15 | 975 Fire Oak, Humble, TX | М | 38000 | 333445555 | 5 |
| Joyce | Α | English | 453453453 | 1972-07-31 | 5631 Rice, Houston, TX | F | 25000 | 333445555 | 5 |
| Ahmad | V | Jabbar | 987987987 | 1969-03-29 | 980 Dallas, Houston, TX | М | 25000 | 987654321 | 4 |
| James | Е | Borg | 888665555 | 1937-11-10 | 450 Stone, Houston, TX | М | 55000 | NULL | 1 |

DEPARTMENT

| Dname | Dname <u>Dnumber</u> | | Dname <u>Dnumber</u> Mgr_ssn | | | |
|----------------|----------------------|-----------|------------------------------|--|--|--|
| Research | 5 | 333445555 | 1988-05-22 | | | |
| Administration | 4 | 987654321 | 1995-01-01 | | | |
| Headquarters | 1 | 888665555 | 1981-06-19 | | | |

DEPT_LOCATIONS

| Dnumber | Dlocation |
|---------|-----------|
| 1 | Houston |
| 4 | Stafford |
| 5 | Bellaire |
| 5 | Sugarland |
| 5 | Houston |

One possible database state for the COMPANY relational database schema

WORKS ON

| Essn | <u>Pno</u> | Hours |
|-----------|------------|-------|
| 123456789 | 1 | 32.5 |
| 123456789 | 2 | 7.5 |
| 666884444 | 3 | 40.0 |
| 453453453 | 1 | 20.0 |
| 453453453 | 2 | 20.0 |
| 333445555 | 2 | 10.0 |
| 333445555 | 3 | 10.0 |
| 333445555 | 10 | 10.0 |
| 333445555 | 20 | 10.0 |
| 999887777 | 30 | 30.0 |
| 999887777 | 10 | 10.0 |
| 987987987 | 10 | 35.0 |
| 987987987 | 30 | 5.0 |
| 987654321 | 30 | 20.0 |
| 987654321 | 20 | 15.0 |
| 888665555 | 20 | NULL |

PROJECT

| Pname | Pnumber | Plocation | Dnum |
|-----------------|---------|-----------|------|
| ProductX | 1 | Bellaire | 5 |
| ProductY | 2 | Sugarland | 5 |
| ProductZ | 3 | Houston | 5 |
| Computerization | 10 | Stafford | 4 |
| Reorganization | 20 | Houston | 1 |
| Newbenefits | 30 | Stafford | 4 |

DEPENDENT

| Essn | Dependent_name | Sex | Bdate | Relationship |
|-----------|----------------|-----|------------|--------------|
| 333445555 | Alice | F | 1986-04-05 | Daughter |
| 333445555 | Theodore | М | 1983-10-25 | Son |
| 333445555 | Joy | F | 1958-05-03 | Spouse |
| 987654321 | Abner | М | 1942-02-28 | Spouse |
| 123456789 | Michael | М | 1988-01-04 | Son |
| 123456789 | Alice | F | 1988-12-30 | Daughter |
| 123456789 | Elizabeth | F | 1967-05-05 | Spouse |

CS3402

Slide 6- 13

Nested Queries

Nested queries

- ◆ Some queries require that existing values in the database be fetched and then used in a comparison condition.
- ◆ Such queries can be conveniently formulated by using nested queries, which are complete select-from-where blocks within another SQL query.
- ◆ The inner one is called nested query and the outer one is called the outer query.
- ◆ These nested queries can also appear in the WHERE clause or the FROM clause or the SELECT clause or other SQL clauses as needed.

Nested Queries with IN

Comparison operator IN

- Compares a single value v with a set (or multiset) of values V
- Evaluates "v IN V" to TRUE if v is one of the elements in V

e.g. Make a list of all project numbers for projects that involve an employee whose last name is 'Smith', either as a worker or as a manager of the department that controls the project.

```
Q4A:
                  SELECT
                             DISTINCT Pnumber
                  FROM
                             PROJECT
                             Pnumber IN
                  WHERE
                                           Pnumber
                             SELECT
The project no. of projects
                              FROM
                                           PROJECT, DEPARTMENT, EMPLOYEE
that have an manager
                              WHERE
                                           Dnum=Dnumber AND
with last name "Smith"
                                           Mgr_ssn=Ssn AND Lname='Smith')
                             OR
The project no. of projects
                             Pnumber IN
                             ( SELECT
                                           Pno
that have an employee
                              FROM
                                           WORKS_ON, EMPLOYEE
with last name "Smith"
                                           Essn=Ssn AND Lname='Smith');
                              WHERE
```

CS340₂

6

Nested Queries with IN

Comparison operator IN

- ◆ If a nested query returns a single value, it is permissible to use = instead of IN for the comparison operator.
- ◆ In general, the nested query will return a table (relation), which is a set or multiset of tuples. So to be safe IN is recommended.

Nested Queries, IN

Comparison operator IN

- ◆ IN can be used to compare tuples of values by placing them within parentheses.
- v IN V will return true if the tuple v exists in V

e.g. Select the Essn of all employees who work the same (project, hours) as the employee Essn ="123456789"

```
FROM WORKS_ON
WHERE (Pno, Hours) IN (SELECT Pno, Hours
FROM WORKS_ON
The (Pno, Hours) of the employee with Essn
="123456789"

FROM WORKS_ON
WORKS_ON
Essn='123456789');
```

Nested Queries, SOME/ANY, ALL

In addition to the IN operator, a number of other comparison operators can be used to compare a single value v to a set or multiset V by combining the keywords ANY/SOME, ALL.

ANY/SOME

- ◆ = ANY (or = SOME) returns TRUE if the value v is equal to some value in the set V and is hence equivalent to IN
- ◆Other operators that can be combined with ANY (or SOME): >, >=, <, <=, and <>.

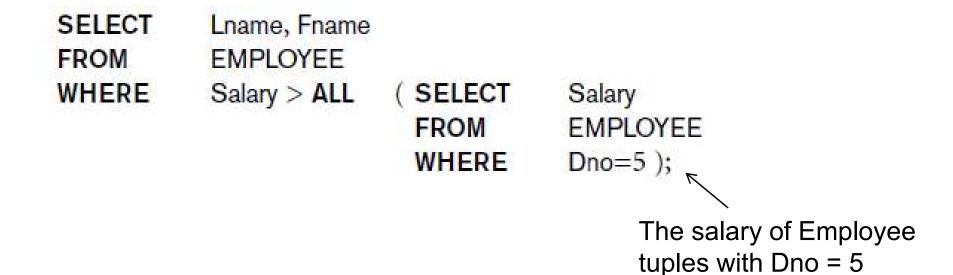
ALL

- ◆ALL can also be combined with each of these operators:=, >, >=, <, <=, and <>.
- ◆E.g. v > ALL V returns TRUE if the value v is greater than all the values in the set V.

Nested Queries, SOME/ANY, ALL

ANY/SOME , ALL

e.g. Select the names of employees whose salary is greater than the salary of all the employees in department 5



What will be returned if we change the ALL to ANY?

Correlated Nested Queries

Correlated nested query

- Whenever a condition in the WHERE clause of a nested query references some attributes of a relation declared in the outer query.
- We can understand a correlated query better by considering that the nested query is evaluated once for each tuple in the outer query.
 - e.g. Retrieve the name of each employee who has a dependent with the same first name and the same sex as the employee.

```
SELECT E.Fname, E.Lname

FROM EMPLOYEE AS E

WHERE E.Ssn IN (SELECT D.Essn

FROM DEPENDENT AS D

WHERE E.Fname=D.Dependent_name AND E.Sex=D.Sex);
```

EXISTS/ NOT EXISTS function

- Boolean functions that return TRUE or FALSE, used in a WHERE clause condition.
- Check whether the result of a nested query is empty or not
- ◆ EXISTS returns TRUE if not empty; NOT EXISTS returns TRUE if empty.
- Typically used in conjunction with a correlated nested query
- e.g. Retrieve the name of each employee who has a dependent with the same first name and the same sex as the employee.

```
SELECT E.Fname, E.Lname

FROM EMPLOYEE AS E

WHERE EXISTS (SELECT *

FROM DEPENDENT AS D

WHERE E.Ssn=D.Essn AND E.Sex=D.Sex

AND E.Fname=D.Dependent_name);
```

EXISTS/ NOT EXISTS function

e.g. retrieve the names of employees who have no dependents:

```
SELECT Fname, Lname

FROM EMPLOYEE

WHERE NOT EXISTS (SELECT *

FROM DEPENDENT

WHERE Ssn=Essn);
```

For each EMPLOYEE tuple, the correlated nested query selects all DEPENDENT tuples whose Essn value matches the EMPLOYEE Ssn; if the result is empty, no dependents are related to the employee, so we select that EMPLOYEE tuple and retrieve its Fname and Lname.

EXISTS/ NOT EXISTS function

e.g. List the names of managers who have at least one dependent :

```
SELECT Fname, Lname
FROM EMPLOYEE
WHERE EXISTS (SELECT *
FROM DEPENDENT
WHERE Ssn=Essn)
AND
EXISTS (SELECT *
FROM DEPARTMENT
WHERE Ssn=Mgr_ssn);
```

The first nested query selects all DEPENDENT tuples related to an EMPLOYEE.
The second selects all DEPARTMENT tuples managed by the EMPLOYEE.
If at least one of the first and at least one of the second exists, we select the EMPLOYEE tuple.

EXISTS/ NOT EXISTS function

e.g. Retrieve the name of each employee who works on all the projects controlled by department number 5:

The first subquery selects all projects controlled by department 5.

The second subquery selects all projects that the particular employee being considered works on.

If the set difference of the first subquery result MINUS (EXCEPT) the second subquery result is empty, it means that the employee works on all the projects and is therefore selected.

- Used to summarize information from multiple tuples into a singletuple
- Built-in aggregate functions
 - ◆ COUNT: returns the number of values / tuples in the set or multiset.
 - SUM, MAX,MIN, and AVG: applied to a set or multiset of numeric values and return, respectively, the sum, maximum value, minimum value, and average (mean) of those values.
- Functions can be used in the SELECT clause or in a HAVING clause

 Used to summarize information from multiple tuples into a singletuple

e.g. Find the sum of the salaries of all employees, the maximum salary, the minimum salary, and the average salary

```
SELECT SUM(Salary), MAX(Salary), MIN(Salary), AVG(Salary)
FROM EMPLOYEE;
```

This query returns a single-row summary of all the rows in the EMPLOYEE table.

| SUM(Salary) | MAX(Salary) | MIN(Salary) | AVG(Salary) |
|-------------|-------------|-------------|-------------|
| 281000 | 55000 | 25000 | 35125 |

- Aggregate Functions combining with where clause
 - Only the tuples satisfying condition will be aggregated.

e.g. Find the sum of the salaries of all employees of the 'Research' department, as well as the maximum salary, the minimum salary, and the average salary in this department

```
SELECT SUM (Salary), MAX (Salary), MIN (Salary), AVG (Salary)
FROM EMPLOYEE, DEPARTMENT
WHERE Dno=Dnumber AND Dname='Research';
```

e.g. Retrieve the number of employees in the 'Research' department.

```
Here the asterisk (*) refers to the rows (tuples), so COUNT (*) returns the number of rows in the result of the query.

WHERE DNO = DNUMBER AND DNAME = 'Research';
```

COUNT:

- will not eliminate duplicated values/tuples
- add keyword "DISTINCT" if we do not want to count duplicated values

e.g. count the number of distinct salary values in the database

```
SELECT COUNT (DISTINCT Salary)
FROM EMPLOYEE;
```

can be used to count tuples with COUNT(*);

- NULL in aggregate functions:
 - In general, NULL values are discarded when aggregate functions are applied to a particular column (attribute); the only exception is for COUNT(*) because tuples instead of values are counted.
 - When an aggregate function is applied to a collection of values, NULLs are removed from the collection before the calculation; if the collection becomes empty because all values are NULL, the aggregate function will return NULL (except in the case of COUNT, where it will return 0 for an empty collection of values).

Aggregate functions in nested query

e.g. retrieve the names of all employees who have two or more dependents:

```
SELECT Lname, Fname

FROM EMPLOYEE

WHERE (SELECT COUNT(*)

FROM DEPENDENT

WHERE Ssn = Essn) >= 2;
```

The correlated nested query counts the number of dependents that each employee has; if this is greater than or equal to two, the employee tuple is selected.

GROUP BY

- Followed by attribute list called the grouping attribute(s).
- partition the relation into nonoverlapping subsets (or groups) of tuples.
- ◆ Each group (partition) will consist of the tuples that have the same value on the grouping attribute(s).
- We can then apply aggregate functions to each such group independently to produce summary information about each group.

■ GROUP BY

◆ The grouping attributes should also appear in the SELECT clause, so that the value resulting from applying each aggregate function to a group of tuples appears along with the value of the grouping attribute(s).

e.g. for each department, retrieve the department number, the number of employees in the department, and their average salary:

```
SELECT Dno, COUNT(*), AVG(Salary)
FROM EMPLOYEE
GROUP BY Dno;
```

Each group having the same value for the GROUP BY attribute Dno.

GROUP BY

Ahmad

James

GROUP BY Dno;

SELECT Dno, COUNT(*), AVG(Salary) FROM EMPLOYEE

> Minit Count (*) Avg (Salary) Fname Lname Salary Super_ssn Dno Dno Ssn B Smith 123456789 30000 333445555 5 33250 John 5 4 T 333445555 888665555 3 Franklin Wong 40000 5 31000 666884444 Ramesh K Narayan 38000 333445555 5 55000 Result of Q24 25000 333445555 A English 453453453 5 Joyce Alicia Zelaya 999887777 25000 987654321 J 4 Jennifer S Wallace 987654321 43000 888665555 4 987654321 V Jabbar 987987987 25000

> > NULL

4

Grouping EMPLOYEE tuples by the value of Dno

888665555

Bong

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55000

■ GROUP BY

- ◆ If NULLs exist in the grouping attribute, then a separate group is created for all tuples with a NULL value in the grouping attribute.
- ◆ GROUP BY can be applied after joining two or more relations.
- ◆ WHERE clause will be performed before GROUP BY, that means only tuples satisfying condition will be grouped.

e.g. for each project, retrieve the project number, the project name, and the number of employees who work on that project:

```
SELECT Pnumber, Pname, COUNT(*)
FROM PROJECT, WORKS_ON
WHERE Pnumber=Pno
GROUP BY Pnumber, Pname;
```

HAVING clause

- Provides a condition to select or reject an entire group. Only the groups that satisfy the condition are retrieved in the result of the query.
- appear in conjunction with a GROUP BY clause.

e.g. For each project on which more than two employees work, retrieve the project number, the project name, and the number of employees who work on the project.

```
SELECT Pnumber, Pname, COUNT(*)
FROM PROJECT, WORKS_ON
WHERE Pnumber=Pno
GROUP BY Pnumber, Pname
HAVING COUNT(*)>2;
```

■ HAVING clause

SELECT Pnumber, Pname, COUNT(*)

FROM PROJECT, WORKS_ON

WHERE Pnumber=Pno

GROUP BY Pnumber, Pname

HAVING COUNT(*)>2;

| Pname | Pnumber | | Essn | Pno | Hours |
|-----------------|---------|-------|-----------|-----|-------|
| ProductX | 1 | | 123456789 | 1 | 32.5 |
| ProductX | 1 | | 453453453 | 1 | 20.0 |
| ProductY | 2 | | 123456789 | 2 | 7.5 |
| ProductY | 2 | | 453453453 | 2 | 20.0 |
| ProductY | 2 | 1 1 | 333445555 | 2 | 10.0 |
| ProductZ | 3 | | 666884444 | 3 | 40.0 |
| ProductZ | 3 | | 333445555 | 3 | 10.0 |
| Computerization | 10 | * * * | 333445555 | 10 | 10.0 |
| Computerization | 10 | | 999887777 | 10 | 10.0 |
| Computerization | 10 | | 987987987 | 10 | 35.0 |
| Reorganization | 20 | | 333445555 | 20 | 10.0 |
| Reorganization | 20 | | 987654321 | 20 | 15.0 |
| Reorganization | 20 | | 888665555 | 20 | NULL |
| Newbenefits | 30 | | 987987987 | 30 | 5.0 |
| Newbenefits | 30 | 1 1 | 987654321 | 30 | 20.0 |
| Newbenefits | 30 | 1 1 | 999887777 | 30 | 30.0 |

These groups are not selected by the HAVING condition of Q26.

| Pname | Pnumber | | Essn | Pno | Hours | | Pname | C | | | | |
|-----------------|---------|-----|-----------|--------|-------|----------------|-------------------|-----|---|--|----------------|--|
| ProductY | 2 | | 123456789 | 2 | 7.5 | □ ⊢ | ProductY | | | | | |
| ProductY | 2 | 1 1 | 453453453 | 2 | 20.0 | 1 J┌ ► | Computerization | | | | | |
| ProductY | 2 | 1 1 | 333445555 | 2 10.0 | 10.0 | | | | ╚ | | Reorganization | |
| Computerization | 10 | | 333445555 | 10 | 10.0 | 17 II- | Newbenefits | | | | | |
| Computerization | 10 | | 999887777 | 10 | 10.0 | 1 _ | Result of Q26 | | | | | |
| Computerization | 10 | | 987987987 | 10 | 35.0 | 1_ | (Pnumber not show | vn) | | | | |
| Reorganization | 20 | 1 | 333445555 | 20 | 10.0 | 17 II | | | | | | |
| Reorganization | 20 | 1 1 | 987654321 | 20 | 15.0 | 1 | | | | | | |
| Reorganization | 20 | 1 1 | 888665555 | 20 | NULL | | | | | | | |
| Newbenefits | 30 | | 987987987 | 30 | 5.0 | 77 | | | | | | |
| Newbenefits | 30 | 1 1 | 987654321 | 30 | 20.0 | | | | | | | |
| Newbenefits | 30 | 1 1 | 999887777 | 30 | 30.0 | 11 | | | | | | |

After applying the HAVING clause condition

After applying the WHERE clause but before applying HAVING

21

Count (*)

3

3

HAVING clause

• the WHERE clause is executed first, to select individual tuples or joined tuples; the HAVING clause is applied later, to select individual groups of tuples.

e.g. count the total number of employees whose salaries exceed \$40,000 in each department, but only for departments where more than five employees work.

SELECT Dno, COUNT(*)
FROM EMPLOYEE
WHERE Salary>40000
GROUP BY Dno
HAVING COUNT(*)>5;

Incorrect!

It will select only departments that have more than five employees who each earn more than \$40,000.

HAVING clause

• the WHERE clause is executed first, to select individual tuples or joined tuples; the HAVING clause is applied later, to select individual groups of tuples.

e.g. count the total number of employees whose salaries exceed \$40,000 in each department, but only for departments where more than five employees work.

```
SELECT Dno, COUNT(*)

FROM EMPLOYEE

WHERE Salary>40000 AND Dno IN

(SELECT Dno
FROM EMPLOYEE

GROUP BY Dno
HAVING COUNT(*)>5)
```

GROUP BY Dno;

EXPANDED Block Structure of SQL Queries

```
SELECT <attribute and function list>
FROM 
[ WHERE <condition> ]
[ GROUP BY <grouping attribute(s)> ]
[ HAVING <group condition> ]
[ ORDER BY <attribute list> ];
```

Summary of SQL Syntax

Table 7.2 Summary of SQL Syntax

```
CREATE TABLE  ( <column name> <column type> [ <attribute constraint> ]
                          {, <column name> <column type> [ <attribute constraint> ]}
                           [  { ,  } ] )
DROP TABLE 
ALTER TABLE  ADD <column name> <column type>
SELECT [ DISTINCT ] <attribute list>
FROM ( { <alias> } | <joined table> ) { , ( { <alias> } | <joined table> ) }
[ WHERE <condition> ]
[GROUP BY <grouping attributes> [HAVING <group selection condition>]]
[ORDER BY <column name> [ <order> ] { , <column name> [ <order> ] } ]
<attribute list> ::= ( * | ( <column name> | <function> ( ( [ DISTINCT ] <column name> | * ) ) )
                   {,(<column name>| <function>(([DISTINCT] <column name>|*))}))
<grouping attributes> ::= <column name> { , <column name> }
<order> ::= ( ASC | DESC )
INSERT INTO  [ ( <column name> { , <column name> } ) ]
(VALUES (<constant value>, {<constant value>}) {, (<constant value>})}
<select statement>)
```

Summary of SQL Syntax

```
Table 7.2 Summary of SQL Syntax

DELETE FROM <a href="mailto:table name">
[WHERE <selection condition">]

UPDATE <a href="mailto:table name">
SET <column name">= <value expression</a> { , <column name">= <value expression</a> }

[WHERE <selection condition</a>]

CREATE [UNIQUE] INDEX <index name</a>
ON <a href="mailto:table name">(<column name</a> [<order</a>) { , <column name</a> [<order</a>) })

[CLUSTER]

DROP INDEX <index name</a>

CREATE VIEW <view name</a> [ (<column name</a> { , <column name</a> })]
```

DROP VIEW <view name>

AS <select statement>

NOTE: The commands for creating and dropping indexes are not part of standard SQL.

References

- **6e**
 - Ch. 4. p. 83 107
 - Ch. 5, p. 111 126

Mid-Term examination

Date: Mar 4th (Mon)

Time: 12:30 - 2:30 pm (2 hours)

Please arrive at least 15 min before the exam starts.

Venue: YEUNG LT-2 OR YEUNG P4703 OR YEUNG Y5-202

Please find your seat and venue from this link:

http://8.210.38.3/midseat/

Format: Open-book exam

Printed version of lecture/tutorial/practice notes, personal notes, and textbook. No electronical devices.

Exam mode: face-to-face, paper writing

Coverage: Lecture 1-6