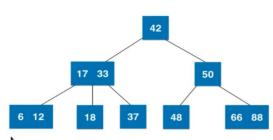
#### 数据结构与算法 14- 红黑树和2-3树

笔记本: 我的笔记

**创建时间:** 2020/10/11 10:52 **更新时间:** 2020/10/11 11:05

作者: liuhouer 标签: 算法

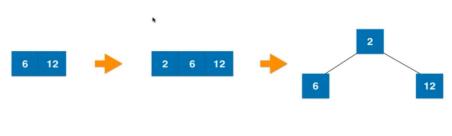


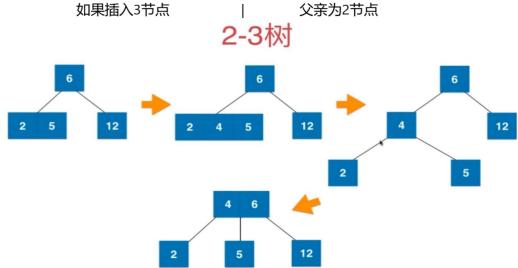


2-3树是一棵绝对平衡的树

#### 2-3树

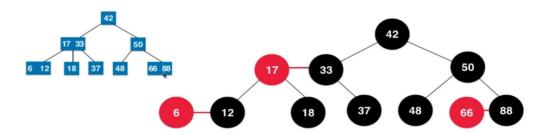
如果插入3-节点





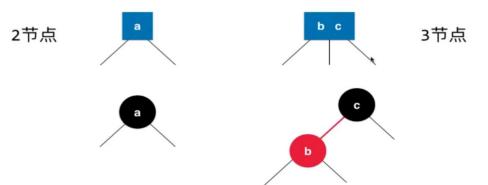
1>>红黑树和二三树在本质上是一样的,只不过用红颜色代表了3节点树根的左侧节点【左倾红黑树】

#### 红黑树和2-3树



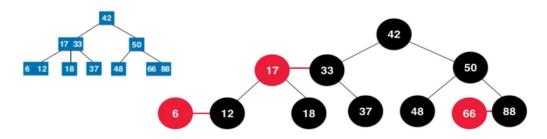
5. 从任意一个节点到叶子节点, 经过的黑色节点是一样的

### 红黑树和2-3树



4. 如果一个节点是红色的, 那么他的孩子节点都是黑色的

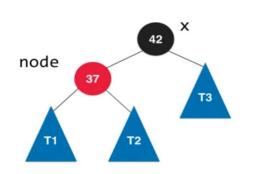
#### 红黑树和2-3树



红黑树是保持"黑平衡"的二叉树

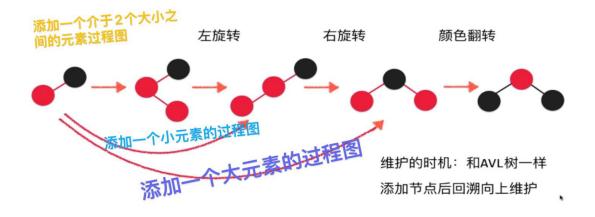
严格意义上, 不是平衡二叉树 最大高度: 21ogn 0(logn)

## 左旋转



node.right = x.left
x.left = node
x.color = node.color
node.color = RED

#### 红黑树添加新元素



# 红黑树的性能总结

对于完全随机的数据,普通的二分搜索树很好用!

缺点: 极端情况退化成链表(或者高度不平衡)

对于查询较多的使用情况, AVL树很好用!

红黑树牺牲了平衡性(2logn的高度)

统计性能更优 (综合增删改查所有的操作)

#### 2.实现红黑树【未实现红黑树的删除】

import java.util.ArrayList;
public class RBTree<K extends Comparable<K>, V> {

```
private static final boolean RED = true;
private static final boolean BLACK = false;
private class Node{
   public K key;
   public V value;
   public Node left, right;
   public boolean color;
   public Node(K key, V value){
       this.key = key;
       this.value = value;
       left = null;
       right = null;
       color = RED;
   }
}
private Node root;
private int size;
public RBTree(){
   root = null;
   size = 0;
}
public int getSize(){
   return size;
public boolean isEmpty(){
   return size == 0;
// 判断节点node的颜色
private boolean isRed(Node node){
   if(node == null)
     return BLACK;
   return node.color;
}
// /\ // T2 T3 T1 T2
                    T1 T2
private Node leftRotate(Node node){
   Node x = node.right;
   // 左旋转
   node.right = x.left;
   x.left = node;
   x.color = node.color;
   node.color = RED;
   return x;
}
   node
```

```
右旋转
         T2
                            node
                            / \
T1 T2
// /\
// y T1
private Node rightRotate(Node node){
   Node x = node.left;
    // 右旋转
   node.left = x.right;
   x.right = node;
   x.color = node.color;
   node.color = RED;
   return x;
}
// 颜色翻转
private void flipColors(Node node){
   node.color = RED;
   node.left.color = BLACK;
   node.right.color = BLACK;
}
// 向红黑树中添加新的元素(key, value)
public void add(K key, V value){
    root = add(root, key, value);
   root.color = BLACK; // 最终根节点为黑色节点
}
// 向以node为根的红黑树中插入元素(key, value), 递归算法
// 返回插入新节点后红黑树的根
private Node add(Node node, K key, V value){
   if(node == null){
       size ++;
       return new Node(key, value); // 默认插入红色节点
   }
    if(key.compareTo(node.key) < 0)</pre>
       node.left = add(node.left, key, value);
    else if(key.compareTo(node.key) > 0)
       node.right = add(node.right, key, value);
    else // key.compareTo(node.key) == 0
       node.value = value;
   if (isRed(node.right) && !isRed(node.left))
       node = leftRotate(node);
    if (isRed(node.left) && isRed(node.left.left))
       node = rightRotate(node);
    if (isRed(node.left) && isRed(node.right))
       flipColors(node);
   return node;
}
```

```
// 返回以node为根节点的二分搜索树中, key所在的节点
private Node getNode(Node node, K key){
    if(node == null)
       return null;
    if(key.equals(node.key))
       return node;
    else if(key.compareTo(node.key) < 0)</pre>
       return getNode(node.left, key);
   else // if(key.compareTo(node.key) > 0)
       return getNode(node.right, key);
}
public boolean contains(K key){
   return getNode(root, key) != null;
public V get(K key){
   Node node = getNode(root, key);
   return node == null ? null : node.value;
public void set(K key, V newValue){
   Node node = getNode(root, key);
   if(node == null)
       throw new IllegalArgumentException(key + " doesn't exist!");
   node.value = newValue;
}
// 返回以node为根的二分搜索树的最小值所在的节点
private Node minimum(Node node){
   if(node.left == null)
       return node;
   return minimum(node.left);
}
// 删除掉以node为根的二分搜索树中的最小节点
// 返回删除节点后新的二分搜索树的根
private Node removeMin(Node node){
    if(node.left == null){
       Node rightNode = node.right;
       node.right = null;
       size --;
       return rightNode;
   node.left = removeMin(node.left);
    return node;
}
// 从二分搜索树中删除键为key的节点
public V remove(K key){
   Node node = getNode(root, key);
   if(node != null){
       root = remove(root, key);
       return node.value;
   }
```

```
return null;
}
private Node remove(Node node, K key){
    if( node == null )
        return null;
    if( key.compareTo(node.key) < 0 ){</pre>
        node.left = remove(node.left , key);
        return node;
    }
    else if(key.compareTo(node.key) > 0 ){
        node.right = remove(node.right, key);
        return node;
          // key.compareTo(node.key) == 0
    else{
        // 待删除节点左子树为空的情况
        if(node.left == null){
            Node rightNode = node.right;
            node.right = null;
            size --;
            return rightNode;
        }
        // 待删除节点右子树为空的情况
        if(node.right == null){
            Node leftNode = node.left;
            node.left = null;
            size --;
            return leftNode;
        }
        // 待删除节点左右子树均不为空的情况
        // 找到比待删除节点大的最小节点,即待删除节点右子树的最小节点
        // 用这个节点顶替待删除节点的位置
        Node successor = minimum(node.right);
        successor.right = removeMin(node.right);
        successor.left = node.left;
        node.left = node.right = null;
        return successor;
    }
}
public static void main(String[] args){
    System.out.println("Pride and Prejudice");
    ArrayList<String> words = new ArrayList<>();
    if(FileOperation.readFile("pride-and-prejudice.txt", words)) {
    System.out.println("Total words: " + words.size());
        RBTree<String, Integer> map = new RBTree<>();
        for (String word : words) {
            if (map.contains(word))
                map.set(word, map.get(word) + 1);
                map.add(word, 1);
```

```
System.out.println("Total different words: " + map.getSize());
System.out.println("Frequency of PRIDE: " + map.get("pride"));
System.out.println("Frequency of PREJUDICE: " +
map.get("prejudice"));
}

System.out.println();
}
```