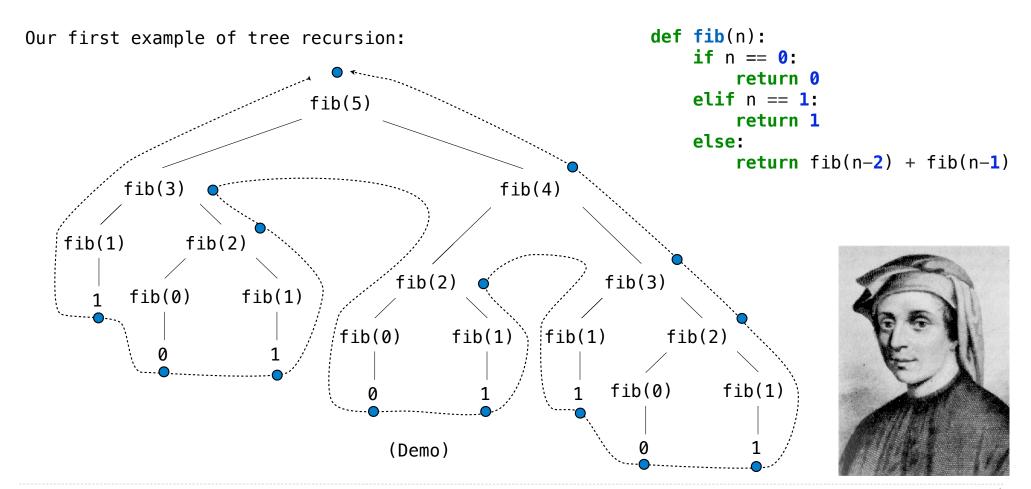
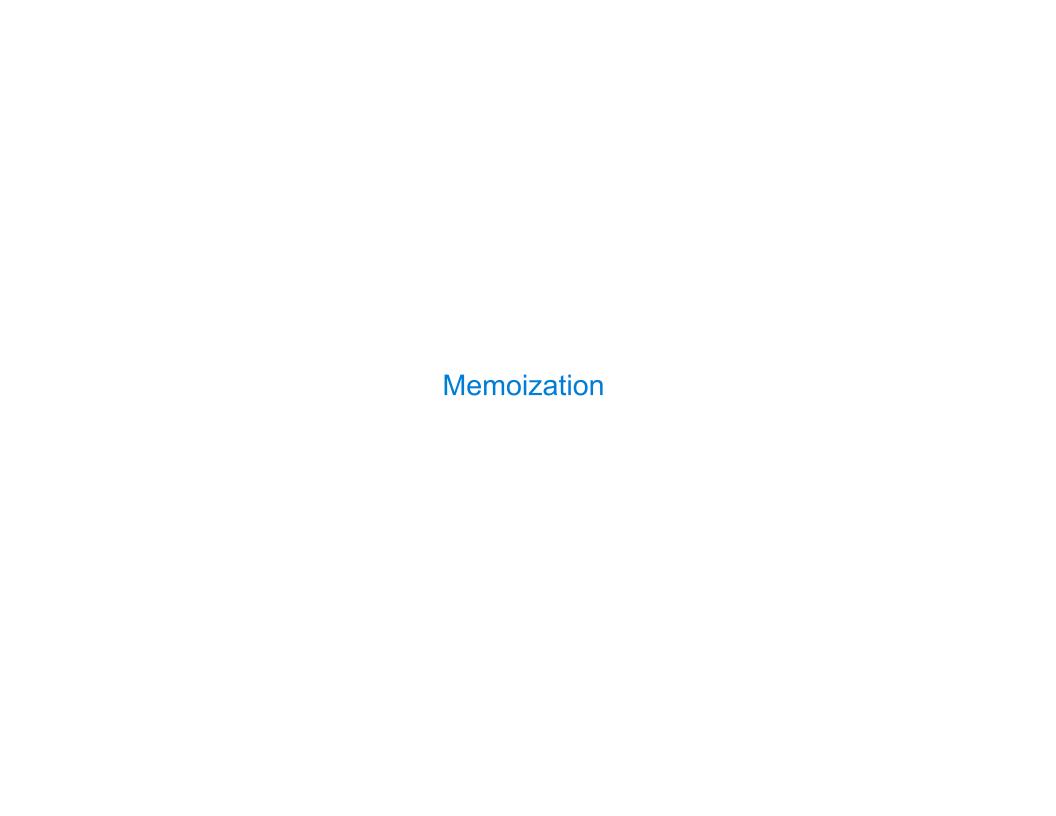


Recursive Computation of the Fibonacci Sequence



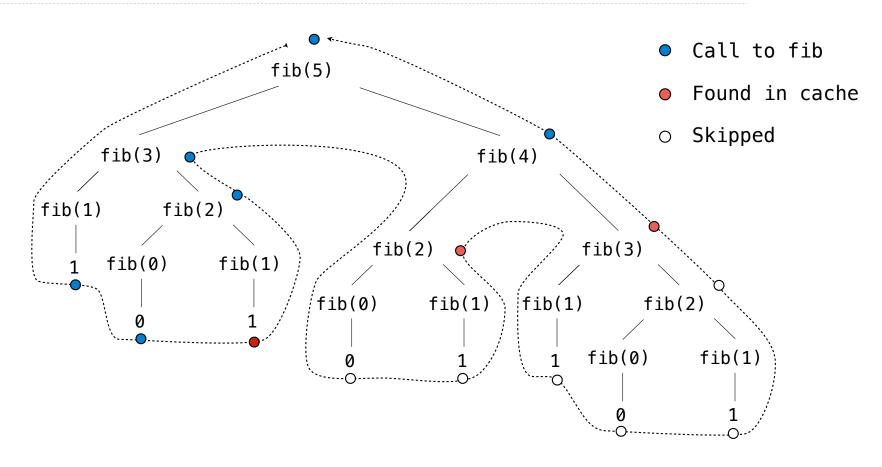


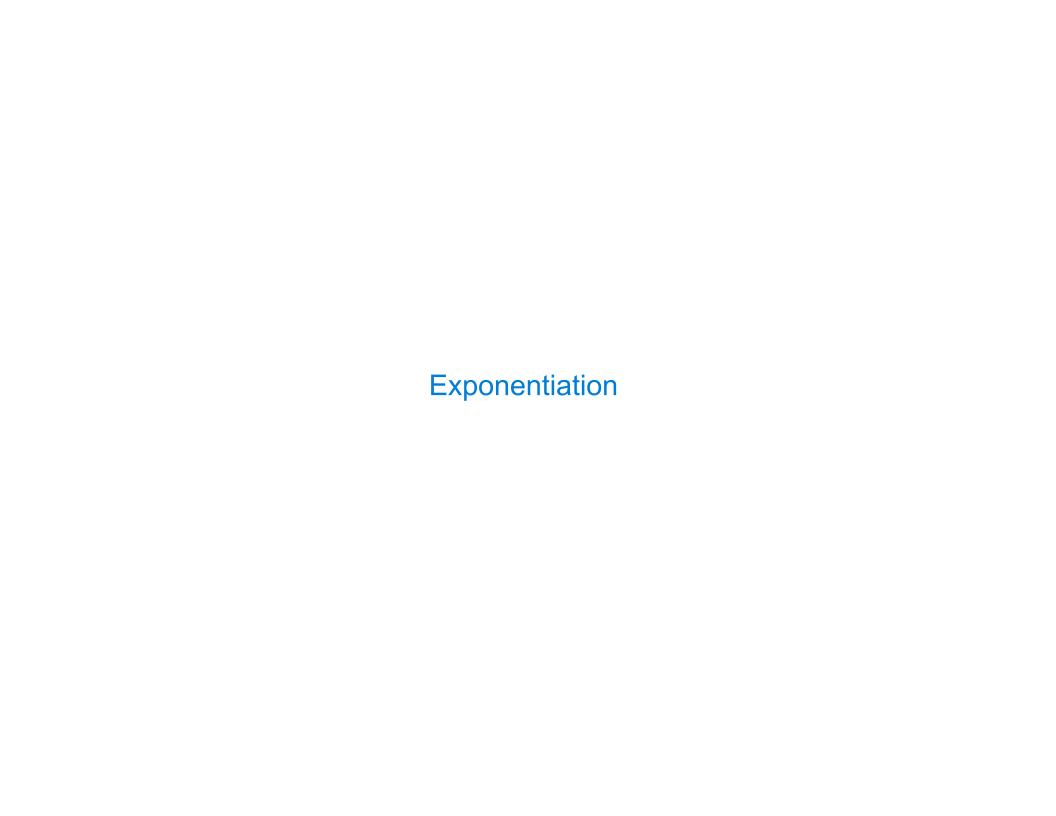
Memoization

Idea: Remember the results that have been computed before

(Demo)

Memoized Tree Recursion





Exponentiation

Goal: one more multiplication lets us double the problem size

```
def exp(b, n):
                                                                                   b^n = \begin{cases} 1 & \text{if } n = 0\\ b \cdot b^{n-1} & \text{otherwise} \end{cases}
       if n == 0:
              return 1
       else:
              return b * exp(b, n-1)
def exp fast(b, n):
       if n == 0:
              return 1
       elif n % 2 == 0:
                                                                                   b^{n} = \begin{cases} 1 & \text{if } n = 0\\ (b^{\frac{1}{2}n})^{2} & \text{if } n \text{ is even}\\ b \cdot b^{n-1} & \text{if } n \text{ is odd} \end{cases}
              return square(exp_fast(b, n//2))
       else:
              return b * exp_fast(b, n-1)
def square(x):
       return x * x
```

(Demo)

Exponentiation

Goal: one more multiplication lets us double the problem size

```
def exp(b, n):
    if n == 0:
        return 1
    else:
        return b * exp(b, n-1)

def exp_fast(b, n):
    if n == 0:
        return 1
    elif n % 2 == 0:
        return square(exp_fast(b, n//2))
    else:
        return b * exp_fast(b, n-1)

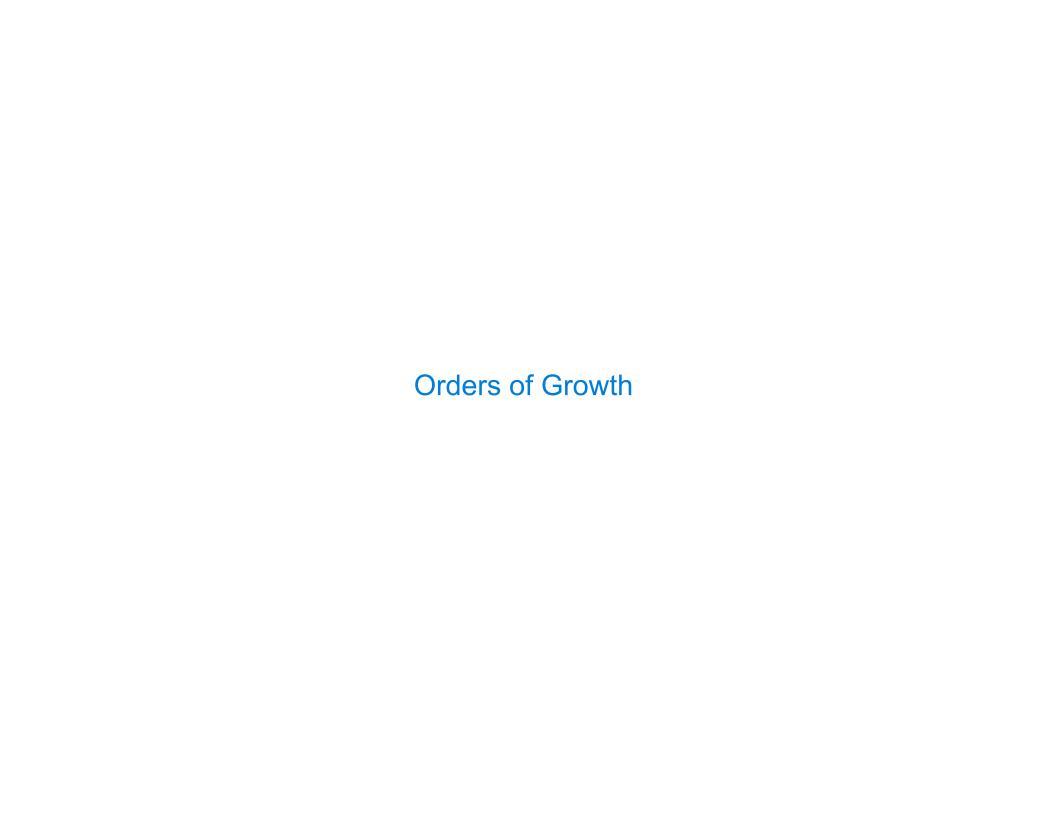
def square(x):
    return x * x
```

Linear time:

- Doubling the input doubles the time
- 1024x the input takes 1024x as much time

Logarithmic time:

- Doubling the input increases the time by one step
- 1024x the input increases the time by only 10 steps



Quadratic Time

Functions that process all pairs of values in a sequence of length n take quadratic time

```
3
def overlap(a, b):
    count = 0
                                                     0
                                                          0
                                               4
    for item in a:
        for other in b:
                                                     0
                                                          1
                                               5
             if item == other:
                 count += 1
    return count
                                                          0
                                                     0
                                               6
overlap([3, 5, 7, 6], [4, 5, 6, 5])
                                                          1
                                                     0
                                               5
```

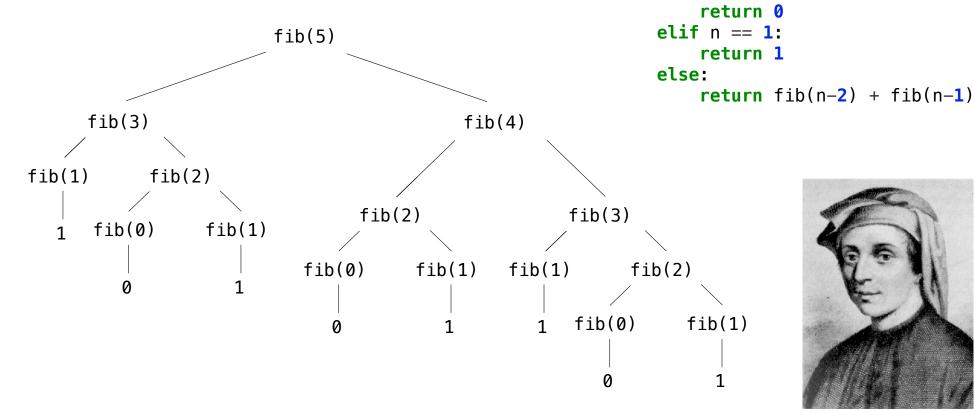
(Demo)

7

6

Exponential Time

Tree-recursive functions can take exponential time





def fib(n):

if n == **0**:

Time for n+n

Time for input n+1

Time for input n

Common Orders of Growth

Exponential growth. E.g., recursive fib

Incrementing n multiplies time by a constant

$$a \cdot b^{n+1} = (a \cdot b^n) \cdot b$$

Quadratic growth. E.g., overlap

Incrementing n increases time by n times a constant

$$a \cdot (n+1)^2 = (a \cdot n^2) + a \cdot (2n+1)$$

Linear growth. E.g., slow exp

Incrementing n increases time by a constant

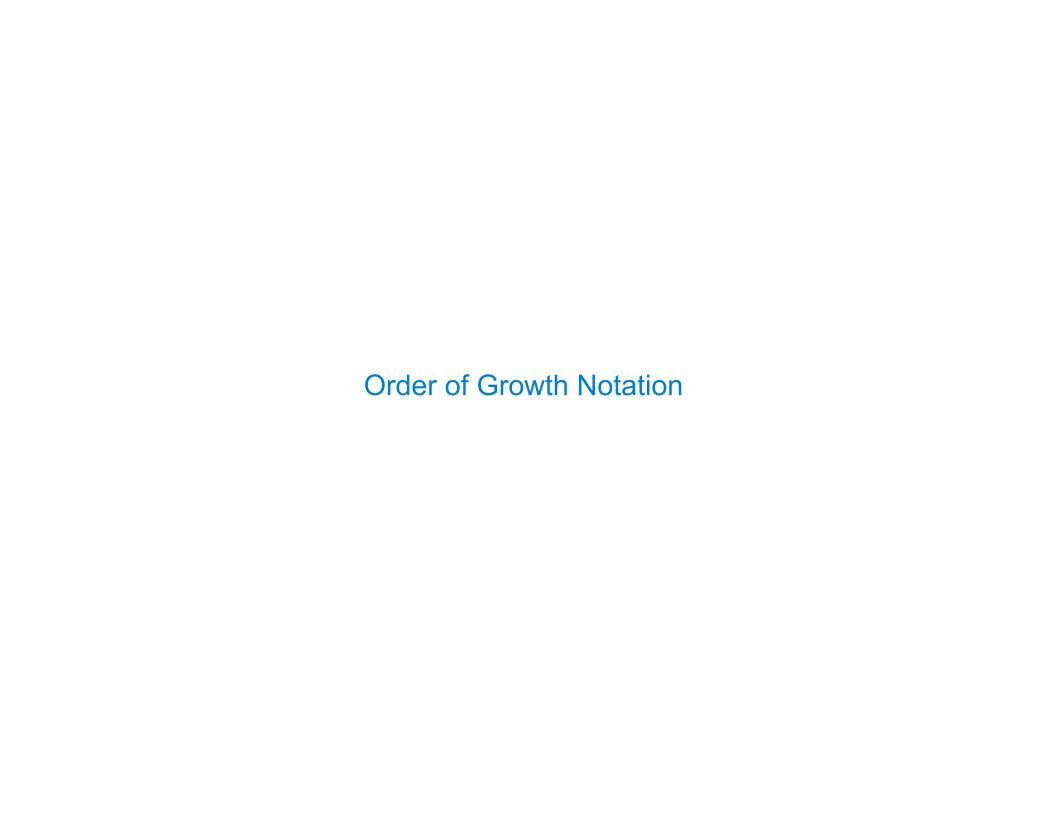
$$a \cdot (n+1) = (a \cdot n) + a$$

Logarithmic growth. E.g., exp_fast

Doubling n only increments time by a constant

$$a \cdot \ln(2 \cdot n) = (a \cdot \ln n) + a \cdot \ln 2$$

Constant growth. Increasing n doesn't affect time



Big Theta and Big O Notation for Orders of Growth

| Exponential growth. E.g., recursive fib Incrementing n multiplies $time$ by a constant | $\Theta(b^n)$ | $O(b^n)$ |
|---|------------------|-------------|
| Quadratic growth. E.g., overlap Incrementing n increases $time$ by n times a constant | $\Theta(n^2)$ | $O(n^2)$ |
| Linear growth. E.g., slow exp Incrementing n increases $time$ by a constant | $\Theta(n)$ | O(n) |
| Logarithmic growth. E.g., exp_fast Doubling n only increments $time$ by a constant | $\Theta(\log n)$ | $O(\log n)$ |
| Constant growth. Increasing n doesn't affect time | $\Theta(1)$ | O(1) |



Space and Environments

Which environment frames do we need to keep during evaluation?

At any moment there is a set of active environments

Values and frames in active environments consume memory

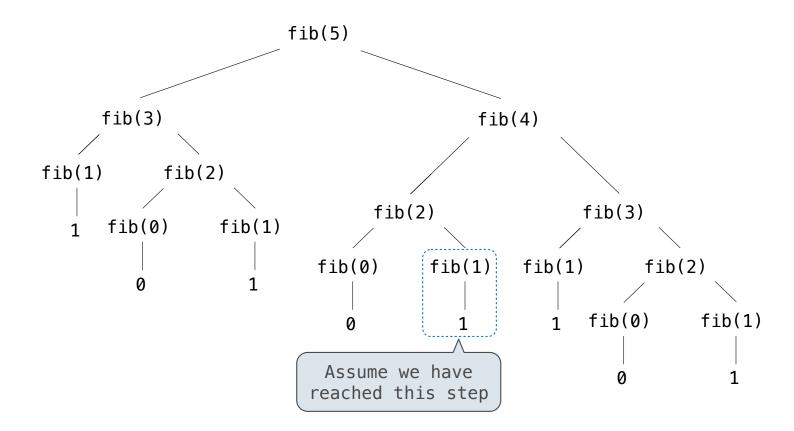
Memory that is used for other values and frames can be recycled

Active environments:

- Environments for any function calls currently being evaluated
- Parent environments of functions named in active environments

(Demo)

Fibonacci Space Consumption



Fibonacci Space Consumption

