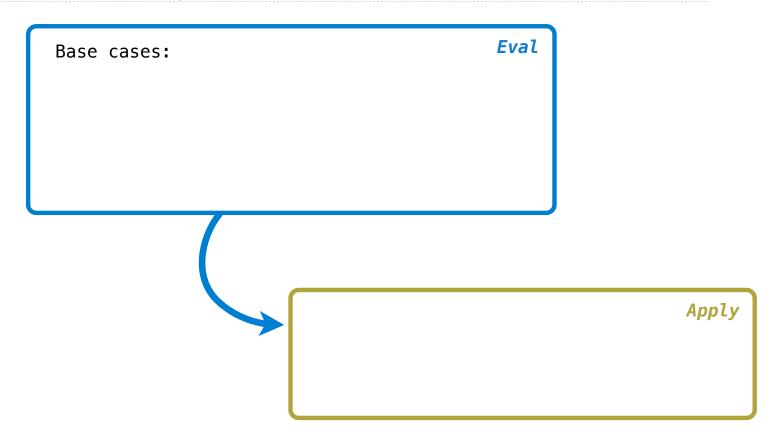
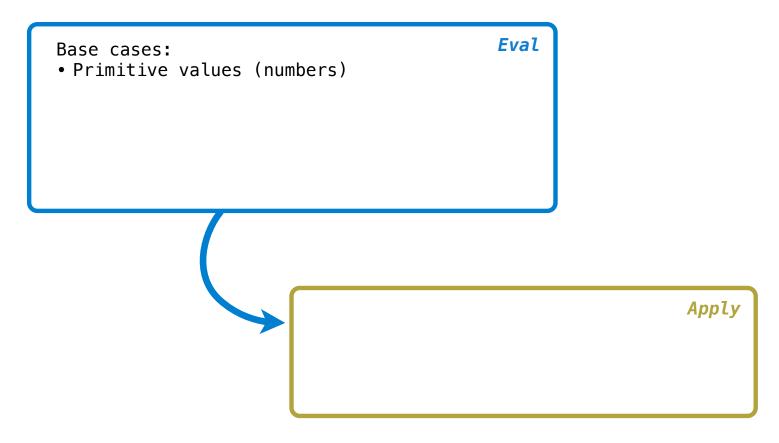
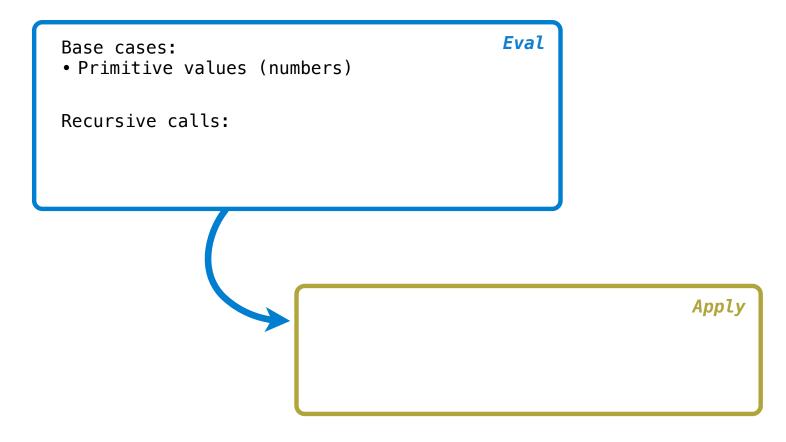
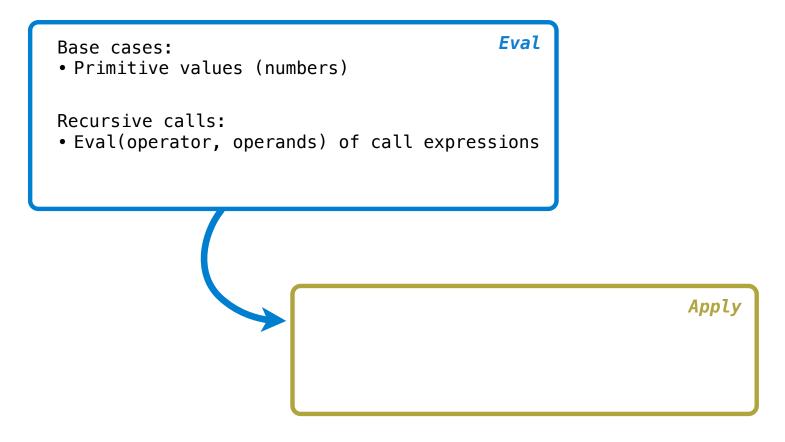


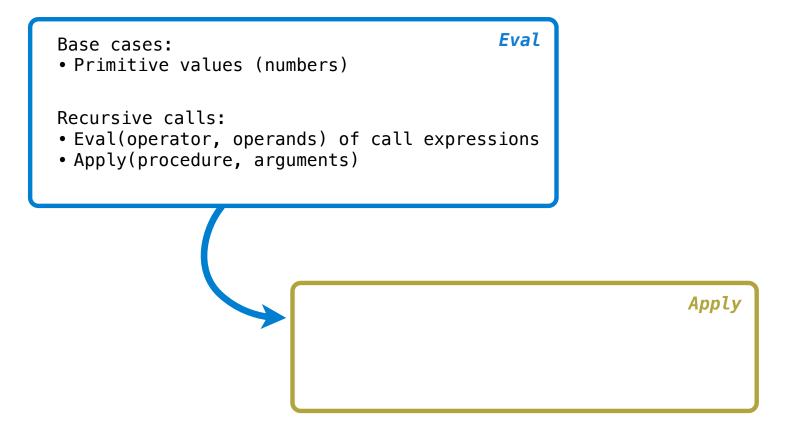
.....











Base cases: • Primitive values (numbers) Recursive calls: • Eval(operator, operands) of call expressions • Apply(procedure, arguments) Base cases: • Built-in primitive procedures

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- Primitive values (numbers)
- Look up values bound to symbols

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Base cases:

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Recursive calls:

• Eval(body) of user-defined procedures

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Recursive calls:

- Eval(operator, operands) of call expressions
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Base cases:

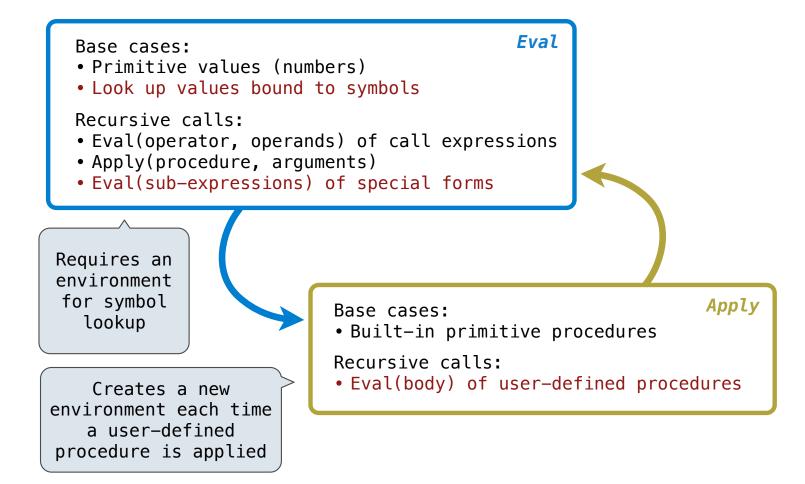
Apply

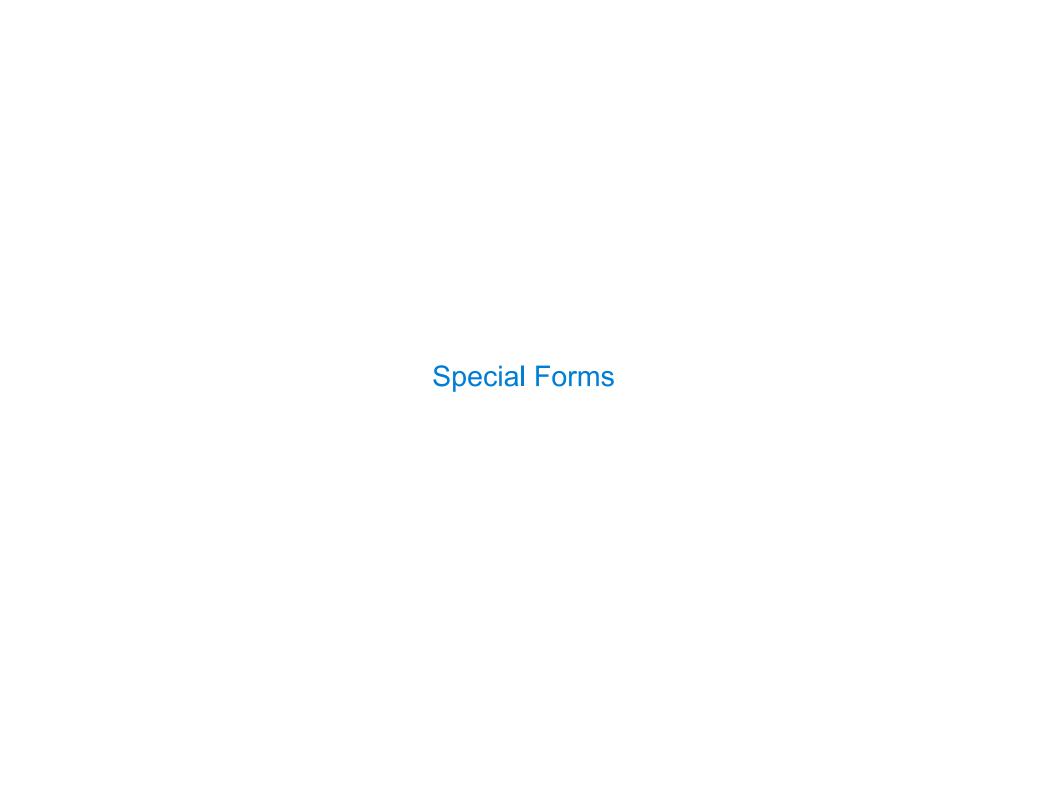
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Recursive calls:

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Eval Base cases: • Primitive values (numbers) • Look up values bound to symbols Recursive calls: • Eval(operator, operands) of call expressions • Apply(procedure, arguments) • Eval(sub-expressions) of special forms Requires an environment for symbol Apply Base cases: lookup • Built-in primitive procedures Recursive calls: • Eval(body) of user-defined procedures





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    (define <name> <expression>)
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```
(if <predicate> <consequent> <alternative>)
  (lambda (<formal-parameters>) <body>)
      (define <name> <expression>)
  (<operator> <operand 0> ... <operand k>)
```

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```
Special forms
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(*coperator* < consequent > < alternative > )

(*coperator* > coperand & > ... < coperand & > )
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(define (demo s) (if (null? s) '(3) (cons (car s) (demo (cdr s))) ))
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(define <name> <expression>)

(define (demo s) (if (null? s) '(3) (cons (car s) (demo (cdr s)))))

(demo (list 1 2))
```



Logical Special Forms

Logical Special Forms

Logical forms may only evaluate some sub-expressions

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```
• If expression: (if f <consequent> <alternative>)
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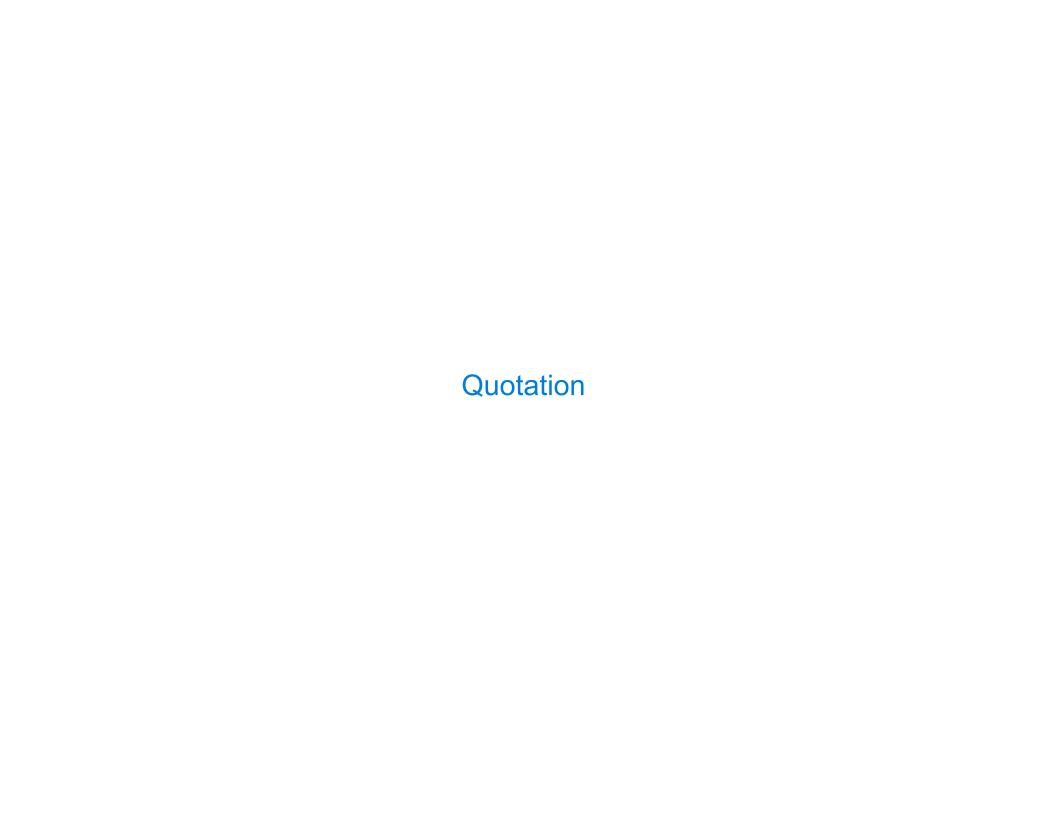
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do_if_form

(Demo)



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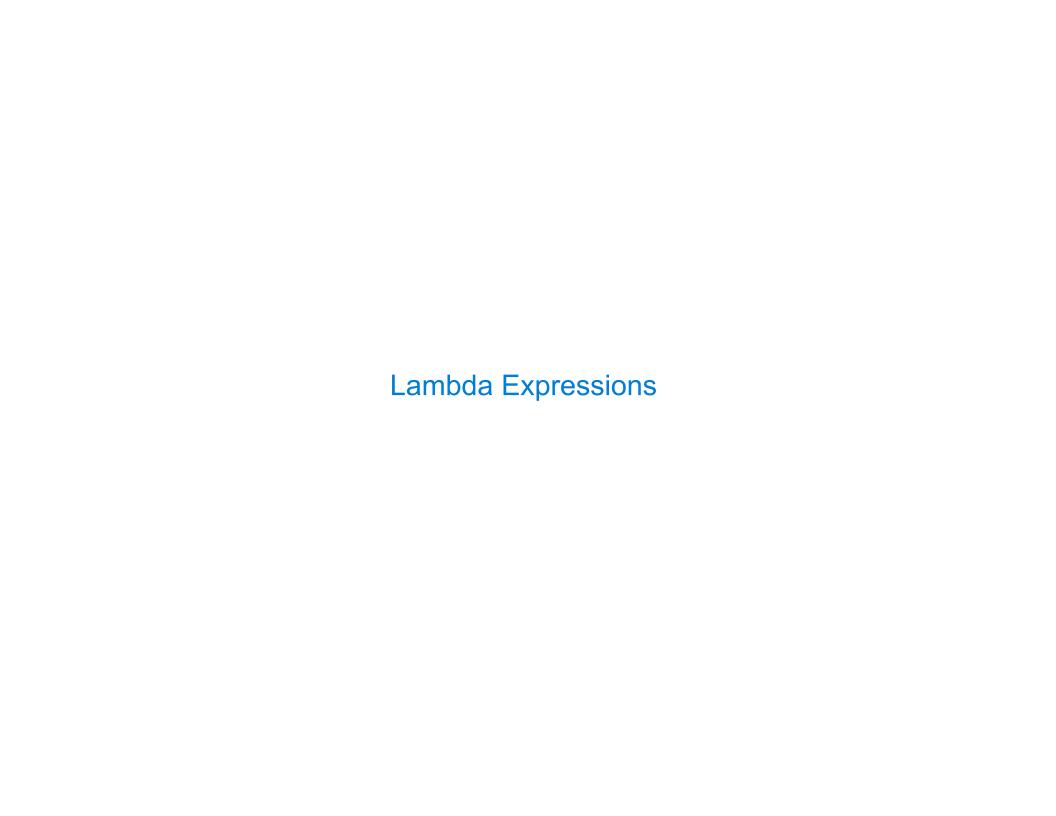
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(Demo)



Lambda expressions evaluate to user-defined procedures

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(lambda (<formal-parameters>) <body>)
(lambda (x) (* x x))
```

self_env = env

```
Lambda expressions evaluate to user-defined procedures

    (lambda (<formal-parameters>) <body>)

    (lambda (x) (* x x))

class LambdaProcedure:
    def __init__(self, formals, body, env):
        self.formals = formals _______ A scheme list of symbols
        self.body = body
```

```
Lambda expressions evaluate to user-defined procedures
```

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(lambda (<formal-parameters>) <body>)
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```
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In Project 4, Frames do not hold return values

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g: Global frame
y 3
z 5

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(Demo)



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```
(define x (+ 1 2))
```

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Procedure definition is shorthand of define with a lambda expression

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(define (<name> <formal parameters>) <body>)
(define <name> (lambda (<formal parameters>) <body>))
```

Applying User-Defined Procedures	
	16

To apply a user-defined procedure, create a new frame in which formal parameters are bound to argument values, whose parent is the **env** attribute of the procedure

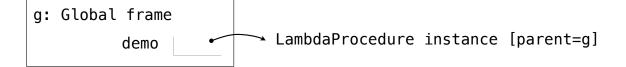
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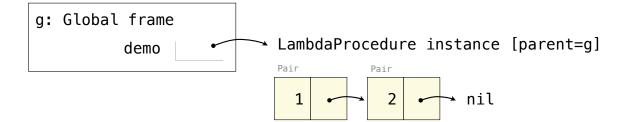


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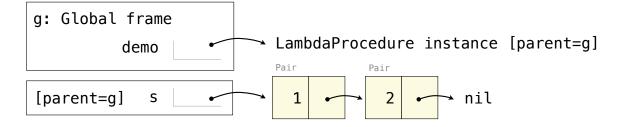
Evaluate the body of the procedure in the environment that starts with this new frame



16

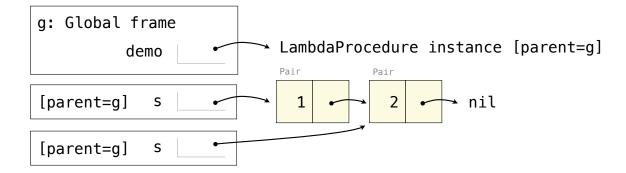
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Evaluate the body of the procedure in the environment that starts with this new frame

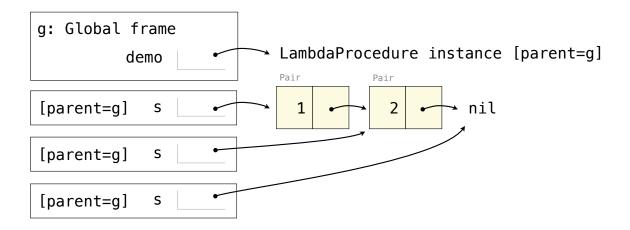


16

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Eval/Ap	ply ir	ı Lisp ´	1.5
---------	--------	----------	-----

Eval/Apply in Lisp 1.5

```
apply[fn;x;a] =
      [atom[fn] \rightarrow [eq[fn;CAR] \rightarrow caar[x];
                     eq[fn;CDR] \rightarrow cdar[x];
                     eq[fn;CONS] \rightarrow cons[car[x];cadr[x]];
                     eq[fn;ATOM] \rightarrow atom[car[x]];
                     eq[fn; EQ] \rightarrow eq[car[x]; cadr[x]];
                     T \rightarrow apply[eval[fn;a];x;a]];
      eq[car[fn]; LAMBDA] \rightarrow eval[caddr[fn]; pairlis[cadr[fn]; x; a]];
      eq[car[fn]; LABEL] - apply[caddr[fn]; x; cons[cons[cadr[fn];
                                                     caddr[fn]];a]]]
eval[e;a] = [atom[e] - cdr[assoc[e;a]];
      atom[car[e]] -
                 [eq[car[e],QUOTE] \rightarrow cadr[e];
                 eq[car[e];COND] - evcon[cdr[e];a];
                 T - apply[car[e];evlis[cdr[e];a];a]];
     T - apply[car[e];evlis[cdr[e];a];a]]
```