

Chapter 8 : Graph Theory

Introduction

Graph Theory

Definition: A graph G consists of two sets V and E , where:

- V is a nonempty set of vertices
- E is a set of edges.

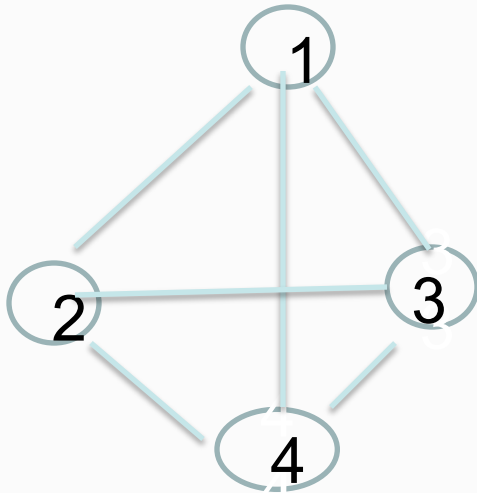
We denote a graph G

$$G = (V, E)$$

Graph Theory

- **Example: undirected Graph G1**

G_1



$$G_1 = (V_1, E_1),$$

$$\text{where } V_1 = \{1, 2, 3, 4\}$$

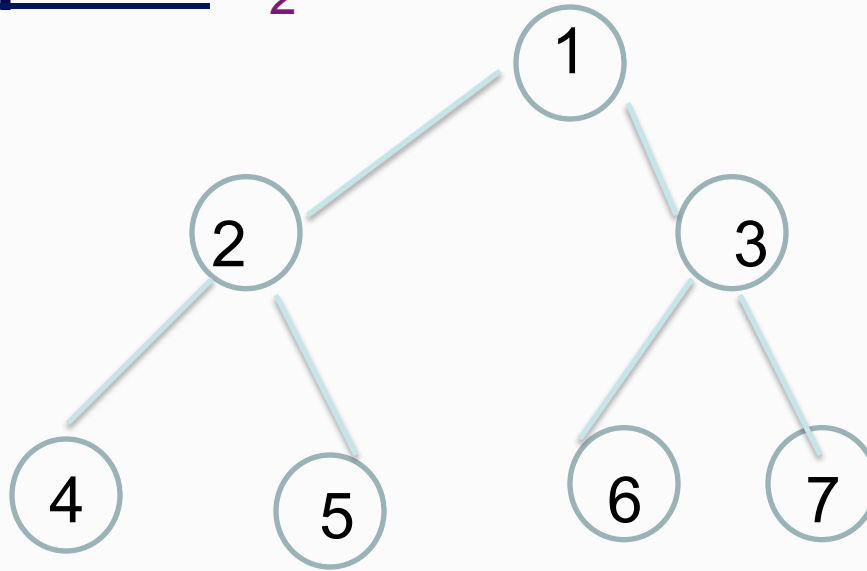
$$E_1 = \{(1,2), (1,3), (1,4), (2,3), (2,4), (3,4)\}$$



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Graph Theory

- Example 2: G_2



$$G_2 = (V_2, E_2)$$

Where $V_2 = \{1, 2, 3, 4, 5, 6, 7\}$

$$E_2 = \{(1, 2), (1, 3), (2, 4), (2, 5), (3, 6), (3, 7)\}$$

- Example 3: Directed Graph G_3

G_3



$$G_3 = (V_3, E_3)$$

$$\text{Where } V_3 = \{1, 2, 3\}$$

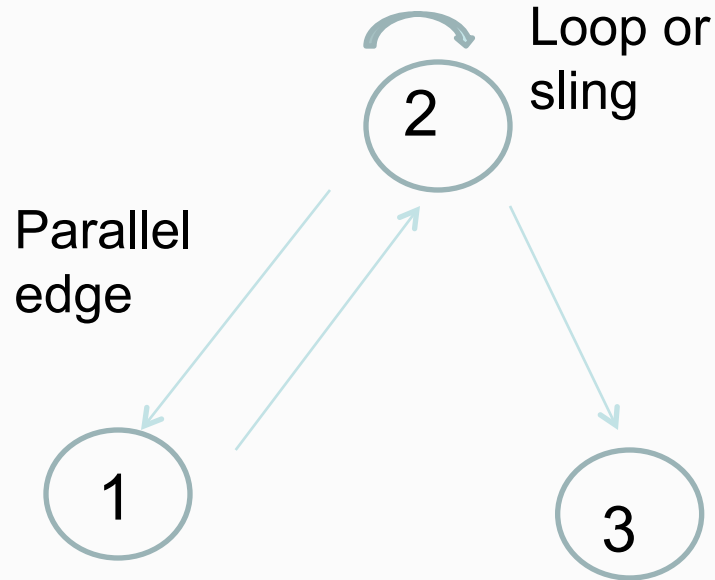
$$E_3 = \{(1, 2), (2, 1), (2, 3)\}$$



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- Example 4: G_4



$$G_4 = (V_4, E_4)$$

Where $V_4 = \{1, 2, 3\}$

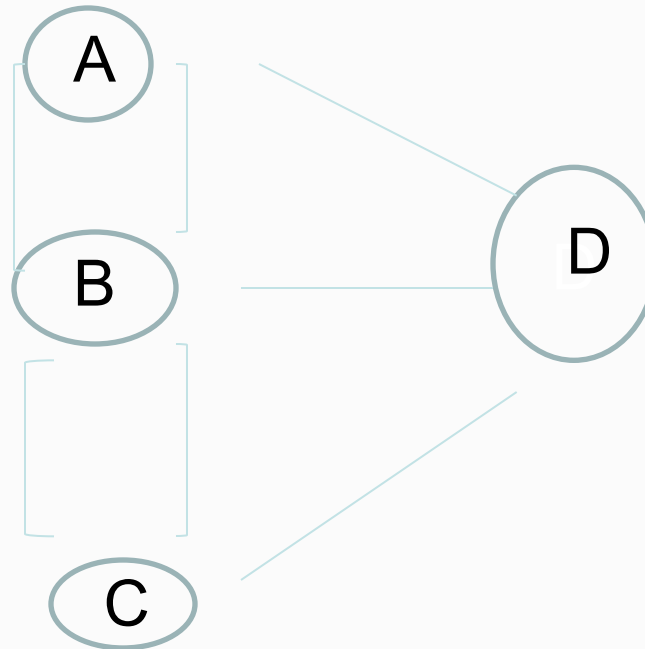
$$E_4 = \{(1, 2), (2, 1), (2, 2), (2, 3)\}$$



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- **Definition:** The **degree** of a vertex v , $\deg(v)$, is the number of edges incident on v .



Deg (A) = 3

Deg (B) = 5

Deg (C) =3

Deg (D) = 3

- **Definition:**

A graph in which the vertices can be partitioned into disjoint sets V_1 and V_2 with every edge incident on one vertex in V_1 and one vertex in V_2 is called a

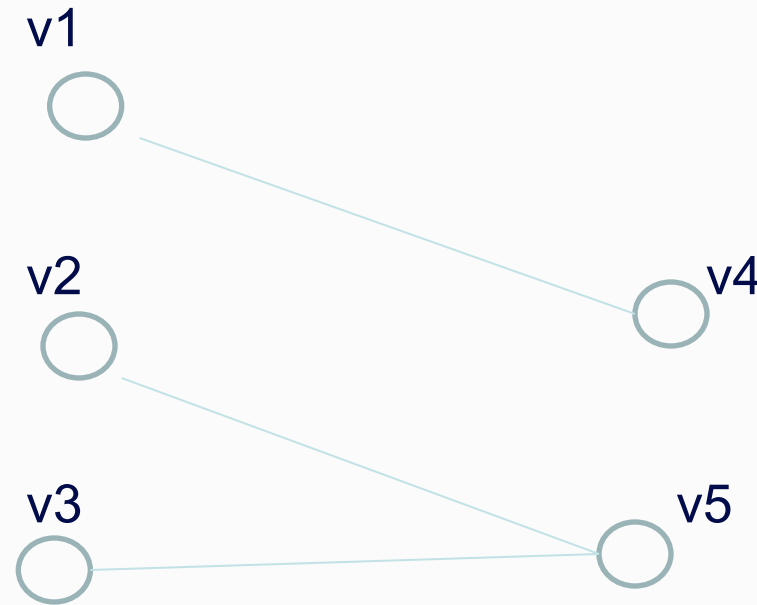
bipartite graph



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- **Example** : This graph is bipartite since if we let
 $V1 = \{v1, v2, v3\}$ and $V2 = \{v4, v5\}$



*Each edge is incident
on one vertex in $V1$ and one vertex in $V2$*



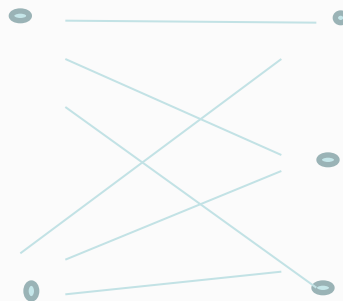
Graph Theory

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- **Definition:**

The graph $K_{m,n}$, called *the complete bipartite graph of m and n vertices*, has disjoint sets V_1 of m vertices and V_2 of n vertices. Every vertex in V_1 is joined to every vertex in V_2 by an edge.

- **Example:**



$K_{2,3}$ is the complete bipartite graph on two and three vertices.

Paths and Cycles

- **Definition:**

Let V_0 and V_n be vertices in a graph. A *path* from V_0 to V_n of length n is an alternating sequence of $n+1$ vertices and n edges beginning with vertex V_0 and ending with vertex V_n .

$$(V_0, e_1, V_1, e_2, V_2, \dots, V_{n-1}, e_n, V_n)$$

In which edge e_i is incident on vertices V_{i-1} and V_i for $i = 1, \dots, n$.



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- Example:

In this graph $(1, e_1, 2, e_2, 3, e_3, 4, e_4, 2)$

Is a path of length 4 from vertex 1 to vertex 2.

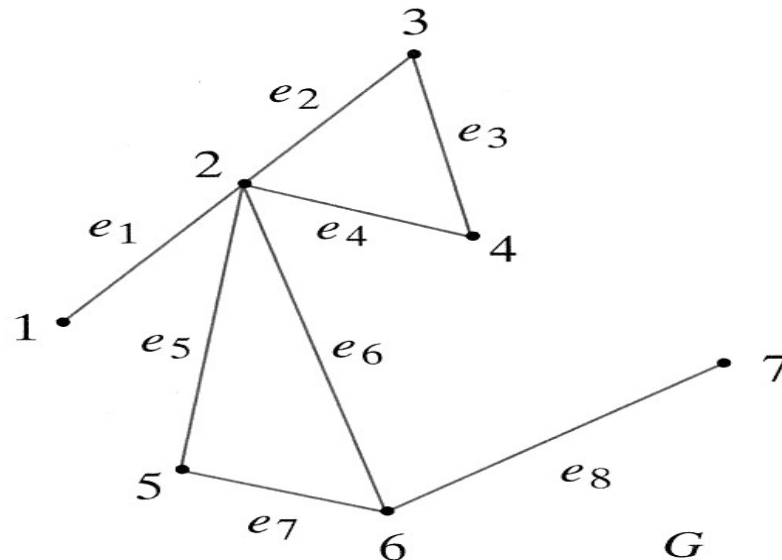


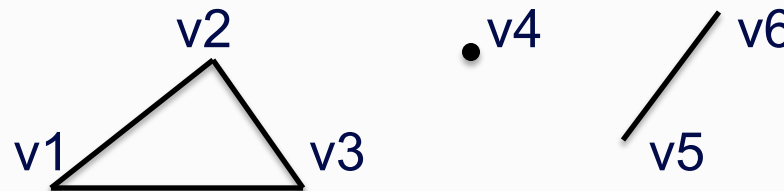
Figure 8.2.1 A connected graph with paths $(1, e_1, 2, e_2, 3, e_3, 4, e_4, 2)$ of length 4 and (6) of length 0.



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- A **simple cycle** is a cycle of the form (v_0, v_1, \dots, v_n) where $v_0 = v_n$ and v_0, v_1, \dots, v_{n-1} are distinct.
- A *graph G is connected* if given any vertices v and w , there is a path from v to w .
- **Example**: the following *graph is NOT connected*:



There is no path from v_2 to v_5

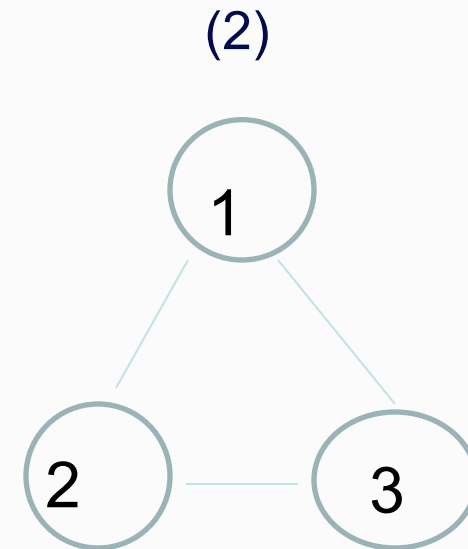
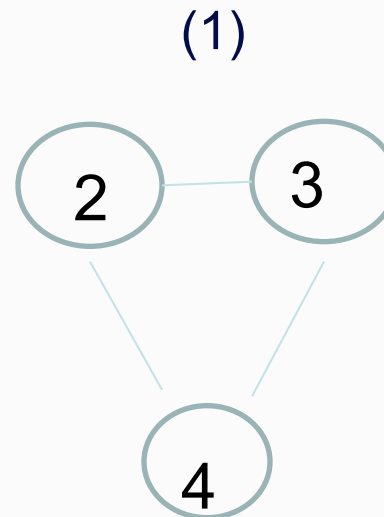
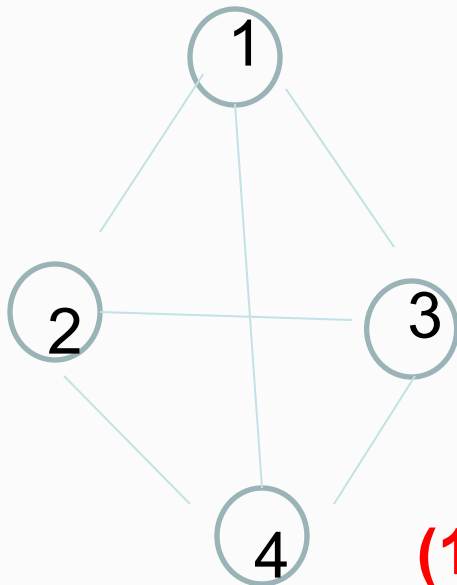


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- **Definition of the Subgraphs:** Let $G = (V, E)$ be a graph. A graph $G' = (V', E')$ is a subgraph of G if $V' \subseteq V$ and $E' \subseteq E$.

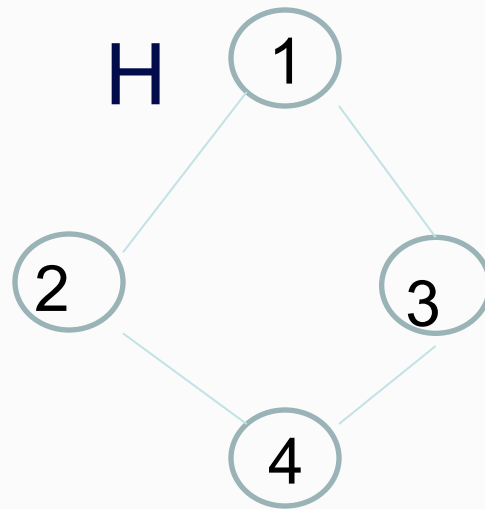
- **Example:** G_1



(1) and (2) are subgraphs of G_1

Graph Theory

- **Definition:** A *connected component* H of an undirected graph G is a maximal connected subgraph.



H is a connected component of G

- **Definition:**

- ❖ A **path** from v to w of length n is an edge sequence from v to w of length n in which the edges are distinct.
- ❖ A **simple path** from v to w of length n is a path of the form $v=v_0, v_1, \dots, v_n=w$, where $v_i \neq v_j$.
- ❖ A **cycle** is a path from v to v .

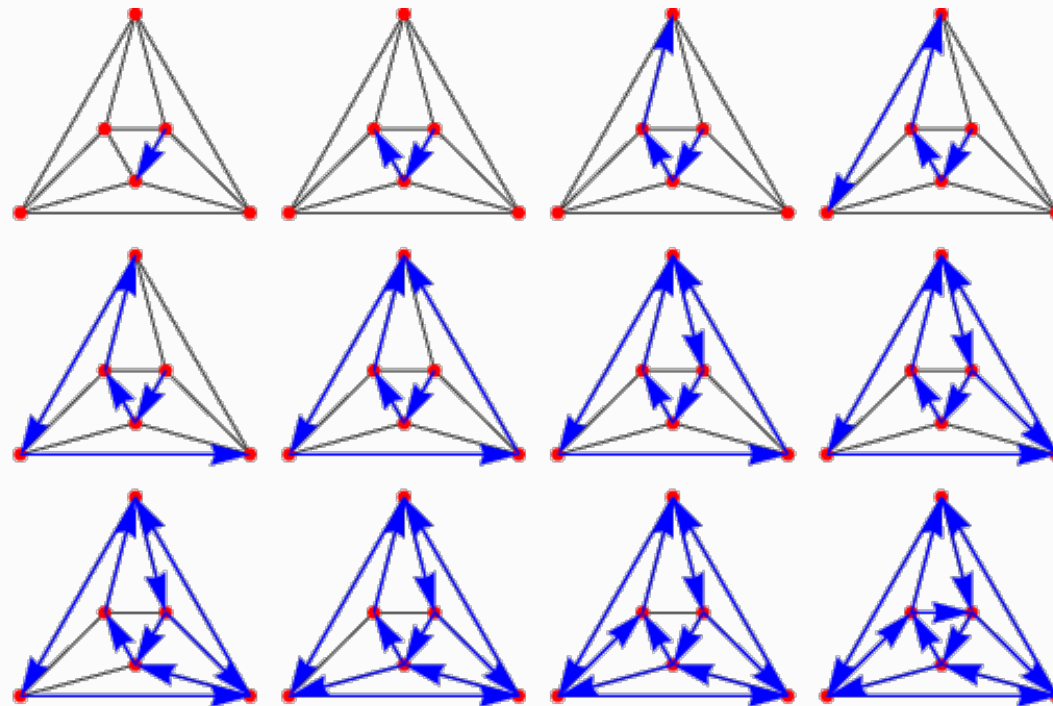


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Graph Theory

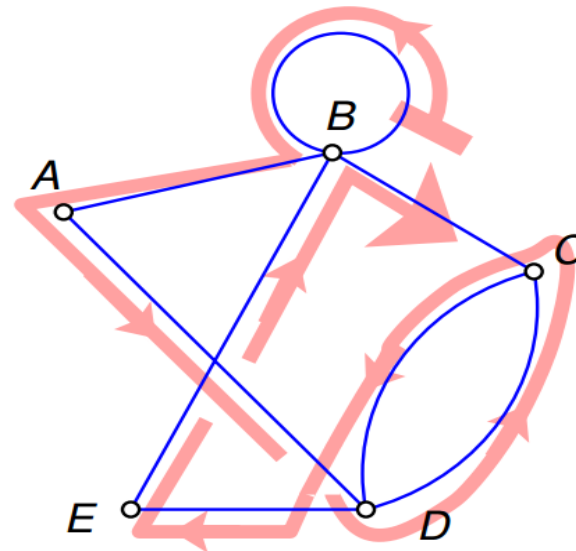
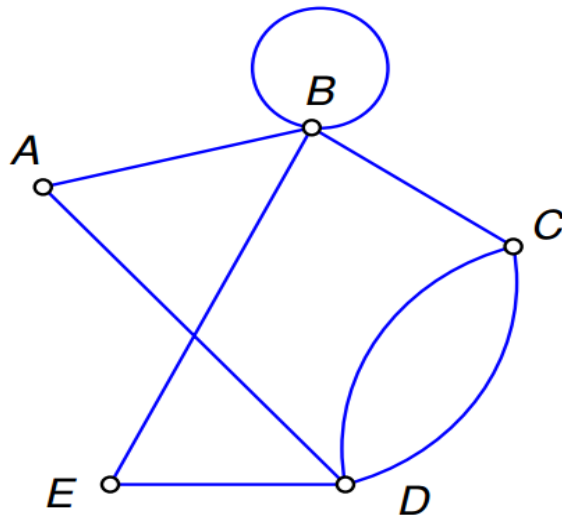
- An Euler Cycle

If a graph G is connected and every vertex has an even degree, *G has an Euler cycle.*



An Euler Path

Euler Paths and Euler Circuits

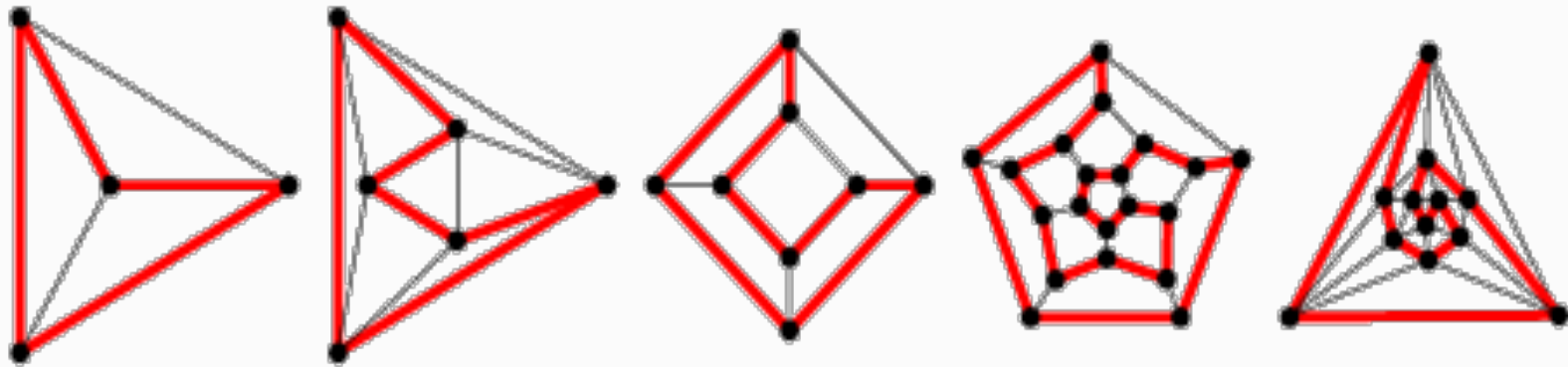


An Euler path: BBADCDEBC

Graph Theory

- *The Hamiltonian Cycle:*

It is a **route** that begins and ends at the same vertex and which passes through each **vertex** of G exactly once.



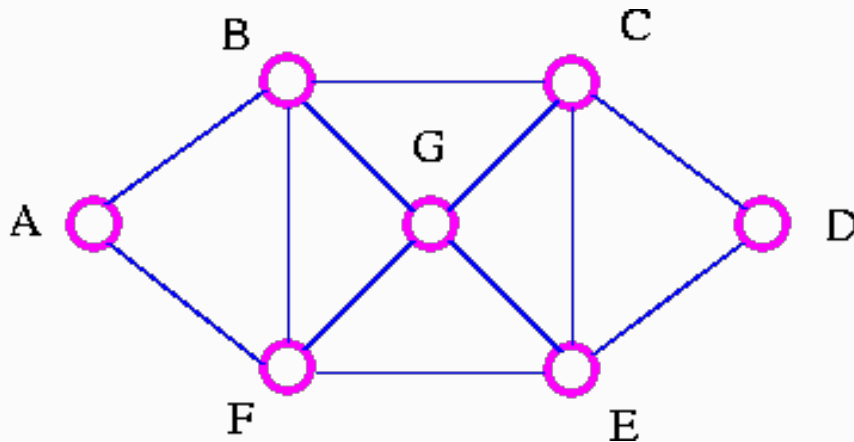


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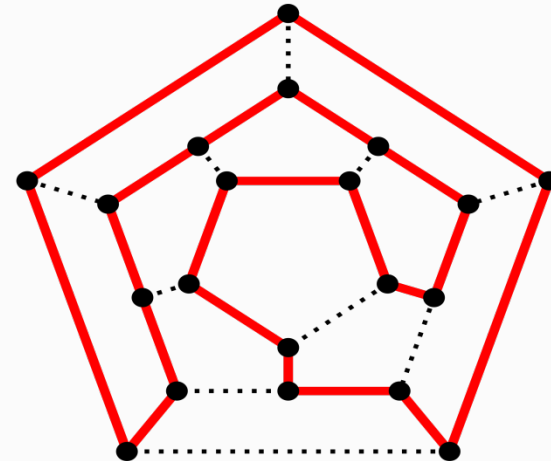
Graph Theory

- An Euler Cycle vs a Hamiltonian Cycle

An Euler cycle visits each **edge** once, whereas a Hamiltonian cycle visits each **vertex** once.



An Euler cycle



A Hamiltonian Cycle

Representations of Graph

Graph Theory

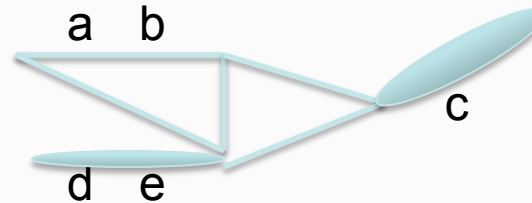
- In the preceding sections we represented a graph by drawing it. Sometimes, as for example in using a computer to analyze a graph, we need a more formal representations.
- Our first method of representing a graph uses *the adjacency matrix*.



Graph Theory

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- Example 1: Consider the graph of the figure:



To obtain the adjacency matrix of this graph:

1. we first select an ordering of the vertices, say a, b, c, d, e.
2. Next, we label the rows and columns of a matrix with the ordered vertices.

The entry in this matrix in row i , column j , $i \neq j$, is the number of edges incident on i and j . If $i = j$, the entry is twice the number of loops incident on i .

The adjacency matrix for this graph is

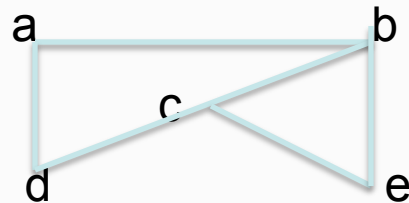
$$\begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{pmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 2 \\ 1 & 1 & 1 & 2 & 0 \end{pmatrix} \end{matrix} = A$$



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- **Example 2:** Consider the graph of the figure:



The adjacency matrix of the simple graph of this figure :

$$\begin{array}{c} \begin{array}{ccccc} & a & b & c & d & e \\ \begin{array}{c} a \\ b \\ c \\ d \\ e \end{array} & \begin{pmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{pmatrix} \end{array} = A\end{array}$$

A, A^2, A^3, \dots

$$A^2 = A * A = ?$$



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A, A^2, A^3, \dots

$A^2 = A * A =$

$$\begin{pmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{pmatrix} = \begin{pmatrix} 2 & 0 & 2 & 0 & 1 \\ 0 & 3 & 1 & 2 & 1 \\ 2 & 1 & 3 & 0 & 1 \\ 0 & 2 & 0 & 2 & 1 \\ 1 & 1 & 1 & 1 & 2 \end{pmatrix}$$

Consider the entry for row a , column c in A^2 , obtained by multiplying pairwise the entries in row a by the entries in column c of the matrix A and summing:

$$\begin{pmatrix} 0 & 1 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 0 \\ 1 \\ 1 \end{pmatrix} = 0 * 0 + 1 * 1 + 0 * 0 + 1 * 1 + 0 * 1 = 2$$



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- **Example 3:** After the previous example, we showed that if A is the matrix of the graph of the previous graph, then

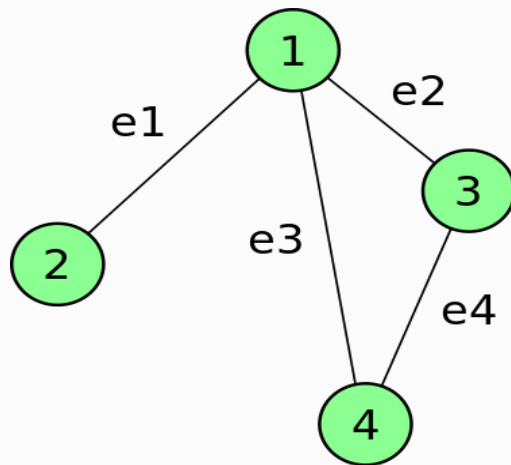
$$A^2 = \begin{pmatrix} 2 & 0 & 2 & 0 & 1 \\ 0 & 3 & 1 & 2 & 1 \\ 2 & 1 & 3 & 0 & 1 \\ 0 & 2 & 0 & 2 & 1 \\ 1 & 1 & 1 & 1 & 2 \end{pmatrix}$$

By multiplying,

$$A^4 = A^2 * A^2 = \begin{pmatrix} 2 & 0 & 2 & 0 & 1 \\ 0 & 3 & 1 & 2 & 1 \\ 2 & 1 & 3 & 0 & 1 \\ 0 & 2 & 0 & 2 & 1 \\ 1 & 1 & 1 & 1 & 2 \end{pmatrix} \begin{pmatrix} 2 & 0 & 2 & 0 & 1 \\ 0 & 3 & 1 & 2 & 1 \\ 2 & 1 & 3 & 0 & 1 \\ 0 & 2 & 0 & 2 & 1 \\ 1 & 1 & 1 & 1 & 1 \end{pmatrix} = \begin{pmatrix} 9 & 3 & 11 & 1 & 6 \\ 3 & 15 & 7 & 11 & 8 \\ 11 & 7 & 15 & 3 & 8 \\ 1 & 11 & 3 & 9 & 6 \\ 6 & 8 & 8 & 6 & 8 \end{pmatrix}$$

Graph Theory

- Definition:** An *Incident Matrix* is a **matrix** that shows the relationship between two classes of objects. If the first class is X and the second is Y, the **matrix** has one row for each element of X and one column for each element of Y.



$$\begin{matrix} & e1 & e2 & e3 & e4 \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix} \end{matrix}$$

Isomorphism of Graphs

Graph Theory

- **Definition:** *Isomorphism of Graphs*

Two **graphs** which contain the same number of **graph** vertices connected in the same way are said to be ***isomorphic***.

Formally, two **graphs** and with **graph** vertices are said to be **isomorphic** if there is a permutation of such that is in the set of **graph** edges.



Graph Theory

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- Example: *Isomorphism of graphs*

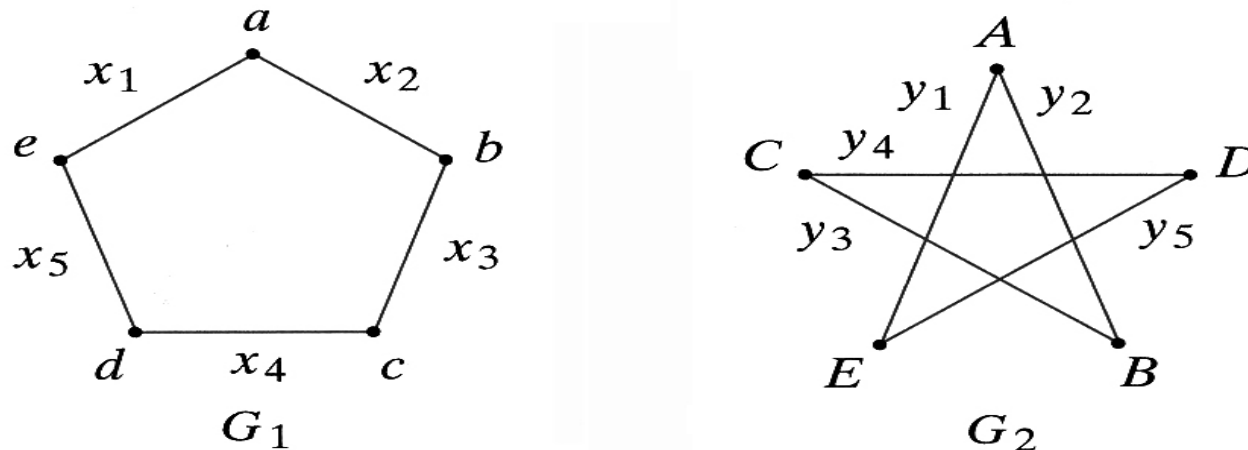


Figure 8.6.1
Isomorphic graphs.

An isomorphism for the graphs G_1 and G_2 of figure 8.6.1 is defined by

$$f(a) = A, \quad f(b) = B, \quad f(c) = C, \quad f(d) = D, \quad f(e) = E$$
$$g(x_i) = y_i, \quad i = 1, \dots, 5$$



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- **Theorem:** Graphs G_1 and G_2 are isomorphic if and only if for some ordering of their vertices, *their adjacency matrices are equal*.
- **Example:**

The adjacency matrix of graph G_1 in Figure 8.6.1 relative to the vertex ordering a, b, c, d, e

$$\begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{pmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{pmatrix} \end{matrix}$$

Is equal to the adjacency matrix of graph G_2 in Figure 8.6.1 relative to the vertex ordering A, B, C, D, E,

$$\begin{matrix} & \begin{matrix} A & B & C & D & E \end{matrix} \\ \begin{matrix} A \\ B \\ C \\ D \\ E \end{matrix} & \begin{pmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{pmatrix} \end{matrix}$$

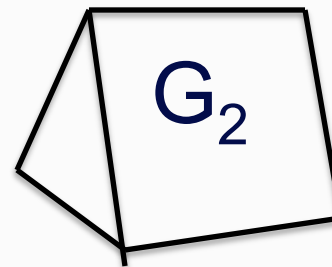
Graphs G_1 and G_2 (fig. 8.6.1) are isomorphic because their adjacency matrices are identical.

Graph Theory

- Example:

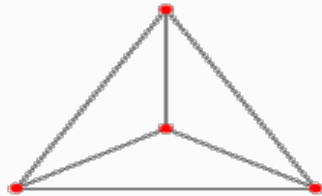
Graphs G_1 and G_2 are NOT isomorphic.

G_1 has 7 edges and G_2 has 6 edges.

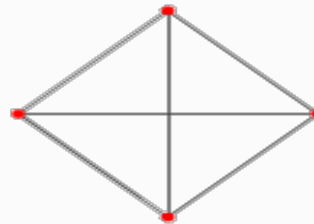


- **Definition:** *Planar Graphs*

A graph is planar if it can be drawn in the plane without its edge crossing.



planar



non-planar