Quiz 03

Name:

Time: Complete and submit to the instructor

Evaluation:

• As described in the syllabus, the Quiz is 20% of the overall grade.

Exercise 1: Find the complexity of the below program:

```
function(int n)
{
   if (n==1)
      return;
   for (int i=1; i<=n; i++)
   {
      for (int j=1; j<=n; j++)
      {
          printf("*");
          break;
      }
}</pre>
```

Exercise 2:

Are each of the following true or false?

```
a) (a) 3 n^2 + 10 n \log n = 0(n \log n)
```

b) (b)
$$3 n^2 + 10 n \log n = Omega(n^2)$$

c) (c)
$$3 n^2 + 10 n \log n = Theta(n^2)$$

d) (d)
$$n \log n + n/2 = O(n)$$

e) (e)
$$10 \text{ SQRT}(n) + \log n = O(n)$$

f) (f)
$$SQRT(n) + log n = O(log n)$$

g) (g)
$$SQRT(n) + log n = Theta(log n)$$

Exercise 1 : Find a theta notation for each expression

h)
$$6n + 1$$

i)
$$3n^2 + 2n \lg n$$

j)
$$2 + 4 + 6 + ... + 2n$$

k)
$$2 + 4 + 8 + 16 + ... + 2^n$$

Exercise 2: Find a theta notation for the number of times the statement x = x + 1 is executed

a) for
$$i = 1$$
 to $2n$

$$x = x + 1$$

b) for
$$i = 1$$
 to $2n$

for
$$j = 1$$
 to n

$$x = x + 1$$

c) for
$$i = 1$$
 to n

for
$$j = 1$$
 to n

for
$$k = 1$$
 to n

$$x = x + 1$$

Exercise 3: Find the greatest common divisor of each pair of integers:

- a) 20, 40
- b) $3^2 \times 7^3$, $3^2 \times 7^3$
- c) 0, 17

Exercise 4 : Find the least common multiple of each pair of integers :

- a) 20, 40
- b) 5, 25

Exercise 5: Express each binary number in decimal.

a) 100000

Exercise 6: Express each decimal number in binary.

a) 43

Exercise 7: Add the binary numbers.

a) 101101 + 11011

 $\underline{Exercise~8:}~Express~the~hexadecimal~number~in~decimal.$

a) B4F

Exercise 9: Add the hexadecimal numbers.

b) F0BA + E9AD

Formula

ANALYSIS OF ALGORITHMS

Analysis an algorithm refers to the process of deriving estimates for the time and space needed to execute the algorithm.

- To derive a theta notation, you must derive both big oh and omega notation.
- An other way to derive big oh, omega and theta estimations is to use known results:

Expression	Name	Estimate
$a_k n^k + a_{k-1} n^{k-1} + \dots + a_1 n + a_0$	Polynomial	Θ(n ^k)
1 + 2 + + n	Arithmetic Sum (Case k = 1 for Next entry)	Θ(n²)
$1^k + 2^k + \dots + n^k$	Sum of Powers	Θ(n ^{k+1})
lg n!	Log n Factorial	Θ(n lg n)

Definition: Let f and g be functions with domain $\{1, 2, 3, ...\}$

We write

$$f(n) = O(g(n))$$

We define f(n) is of order at most g(n) or f(n) is big oh of g(n) if there exists a positive constant C1 such that

$$|f(n)| \le C1 |g(n)|$$

For all but finitely many positive integers n.

We write

$$f(n) = \Omega(g(n))$$

We define f(n) is of order at least g(n) or f(n) is omega of g(n) if there exists a positive constant C2 such that

$$|f(n)| \ge C2 |g(n)|$$

For all but finitely many positive integers n.

We write

$$f(n) = \Theta(g(n))$$

We define f(n) is of order g(n) or f(n) is theta of g(n) if f(n) = O(g(n)) and f(n) = O(g(n)).

By replacing O by Ω and "at most" by "at least", we obtain the definition of what it means for the best-case, worst-case, or average-case time of an algorithm to be of order at least g(n). If the best-case time required by an algorithm is O(g(n)) and $\Omega(g(n))$, we say that the best-case time required by an algorithm is O(g(n)).

Introduction to number of theory

- Let m and n be integers with n and m different to zero. A common divisor of m and n is an integer that divides both m and n. The greatest common divisor, written gcd(m, n)
- Let m and n be positive integers. A *common multiple* of m and n is an integer that divides by both m and n. The *least common divisor*, written lcm(m, n)
- lcd(m, n) is the smallest common multiple of m and n.
 - *The decimal number system*: it represents integers using 10 symbols.
 - *The binary number system*: it represents integers using bits (a bit is a binary digit, that is a 0 or a 1).
 - The hexadecimal number system: it represents integers using 16 symbols
 - The <u>Euclidean algorithm</u> is an odd, famous, and efficient algorithm for finding the greatest common divisor of two integers, The Euclidean algorithm is based on the fact that if $r = a \mod b$, then $\gcd(\mathbf{a}, \mathbf{b}) = \gcd(\mathbf{b}, \mathbf{r})$

Hexadecimal to Decimal Conversion Chart

Hexadecimal	Decimal
0	0
1	1
2	2
3	3
4	4
5	5
6	6
7	7
8	8
9	9
A	10
В	11
С	12
D	13
Е	14
F	15