

14. Find grammars for $\Sigma = \{a, b\}$ that generate the sets of

- (a) all strings with exactly two a 's.
- (b) all strings with at least two a 's. $= f(b)$
- (c) all strings with no more than three a 's. $= f(c)$
- (d) all strings with at least three a 's. $= f(d)$
- (e) all strings that start with a and end with b . $= f(e)$

(b) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow S'aS' | S'bS' | \lambda$

(c) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow S'aS'aS' | S'aS'aS' | S'aS' | S', S' \rightarrow bS' | \lambda$

(d) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow S'aS'aS', S' \rightarrow aS' | bS' | \lambda$

(e) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aS'b, S' \rightarrow aS' | bS' | \lambda$

Argument (b)

$\Rightarrow L(G) \subseteq f(b)$

If $n_a(\text{str}) < 2$, then str must not be accepted by $S \rightarrow S'aS'$

$A_1, A_2, A_3 \in L(G)$

W₁₂₃ = A₁ a B₂ a C₃, i, j, k represent the length of A, B, C respectively. base case: aa is equivalent to W₀₀₀ = A₀ a B₀ a C₀. $G' = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aS' | bS' | \lambda$

$A_0 \in L(G), A_1 = aA_0 | A_0a | bA_0 | Ab \in L(G)$. Assume that $A_0 \in L(G)$ for $n \geq 0$.

$A_{n+1} = aA_n | A_n a | bA_n | Ab \in L(G) \forall n \geq 0, A_n \in L(G)$. Similarly, B and C $\in L(G)$.

which implies that $f(b) \subseteq L(G) \Rightarrow L(G) = f(b)$

Argument (c)

$\Rightarrow L(G) \subseteq f(c)$

If $n_a(\text{str}) > 3$, str can't be generated by G

since at least one S' contains a by pigeon-hole principle.

There are only 4 cases:

Argument (d)

$\Rightarrow L(G) \subseteq f(d)$

If $n_a(\text{str}) < 3$, then str must not be accepted by $S \rightarrow S'aS'aS'$

Let W₁₂₃ = A₁ a B₂ a C₃ a D₄, $G' = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aS' | bS' | \lambda$

base case: $A_0 = \lambda \in L(G), A_1 = aA_0 | A_0a | bA_0 | Ab \in L(G)$. Assume that $n \geq 0$.

$A_n \in L(G), A_{n+1} = aA_n | A_n a | bA_n | Ab \in L(G)$, which means that $\forall n \geq 0, A_n \in L(G)$.

Similarly, B_i, C_j, D_k $\in L(G) \forall i, j, k \geq 0$, which implies that $W_{ijkl} \in L(G)$

$\forall i, j, k, l \geq 0 \Rightarrow f(d) \subseteq L(G) \Rightarrow L(G) = f(d)$

Argument (e)

$\Rightarrow L(G) \subseteq f(e)$

If str is not begin by a or end by b, it isn't accepted by

$S \rightarrow aSb$. Let W_k be a string, l is the length. $W_0 = \lambda$

base case: $ab = aW_0 b, G' = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aS' | bS' | \lambda$

$W_0 \in L(G), W_1 = aW_0 | W_0a | bW_0 | bW_0 \in L(G)$. Assume that

for $n \geq 0, W_n \in L(G), W_{n+1} = aW_n | W_n a | bW_n | bW_n \in L(G)$

which implies that $\forall n \geq 0, aW_nb \in L(G) \subseteq f(e) \subseteq L(G)$

$\Rightarrow L(G) = f(e)$

$\Rightarrow f(c) \subseteq L(G) \Rightarrow L(G) = f(c)$

Case 1: (no a \Rightarrow string = A_L)

A_L is a string. l is the length. $A_0 = \lambda$

$L(G) = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow bS' | \lambda$

Base case: $A_0 \in G', A_1 = bA_0 | Ab \in L(G)$

Assume that for $n \geq 0, A_n \in L(G)$

$A_{n+1} = bA_n | Ab \in L(G)$, which means that

$\forall n \geq 0, A_n \in L(G) \Rightarrow f(A_L : bl^0) \subseteq L(G)$

Case 2: (one a $\Rightarrow A_1 aB_j$)

We can prove that $A_i aB_j \in L(G) \forall i, j \geq 0$

similarly from case 1, which means that

$f(A_i aB_j : vi, j \geq 0) \subseteq L(G)$

Case 3: (two a's $\Rightarrow A_1 aB_j aC_k$)

Case 4: (three a's $\Rightarrow A_1 aB_j aC_k aD_l$)

These two cases can also be proved similarly from case 1 and case 2.

$\Rightarrow f(A_1 aB_j aC_k aD_l : vi, j, k, l \geq 0) \subseteq L(G)$

$f(A_1 aB_j aC_k aD_l : vi, j, k, l \geq 0) \subseteq L(G)$

17. Let $\Sigma = \{a, b\}$. For each of the following languages, find a grammar that generates it.

(a) $L_1 = \{a^n b^m : n \geq 0, m < n\}$.

(b) $L_2 = \{a^{3n} b^{2n} : n \geq 2\}$.

(c) $L_3 = \{a^{n+3} b^n : n \geq 2\}$.

(d) $L_4 = \{a^n b^{n-2} : n \geq 3\}$.

(e) $L_1 L_2$.

(f) $L_1 \cup L_2$.

(g) L_1^3 .

(h) L_1^* .

(i) $L_1 - L_4$.

(a) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aS', S' \rightarrow aS'b | aS' | \lambda$

(b) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aaaaS'bbb, S' \rightarrow aaaS'b | \lambda$

(c) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aaaaS'b | \lambda$

(d) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aaaS'b, S' \rightarrow aS'b | \lambda$

(e) $G = \{f(S, A, B), f(a, b), S, P\}, P = S \rightarrow aA aaaaaB bbbb, A \rightarrow aAb | aA | \lambda, B \rightarrow aaaBbb | \lambda$

(f) $G = \{f(S, S'), f(a, b), S, P\}, P = S \rightarrow aS', S' \rightarrow aS'b | aS' | \lambda$

(g) $G = \{f(S, S', S''), f(a, b), S, P\}, P = S \rightarrow S'S'S', S' \rightarrow aS'', S'' \rightarrow aS'b | aS'' | \lambda$

(h) $G = \{f(S, A, B), f(a, b), S, P\}, P = S \rightarrow AS | \lambda, A \rightarrow aB, B \rightarrow aBb | aB | \lambda$

(i) $G = \{f(S, A, B, C, D), f(a, b), S, P\}, P = S \rightarrow aaaAb | aab | aC, A \rightarrow aBb | B, B \rightarrow aC | bD, C \rightarrow aC | \lambda, D \rightarrow bD | \lambda$

19. Find a grammar that generates the language

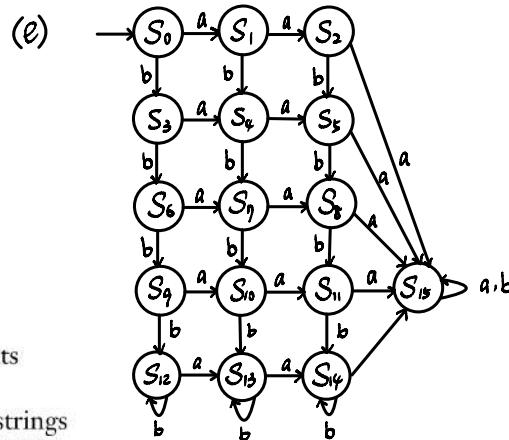
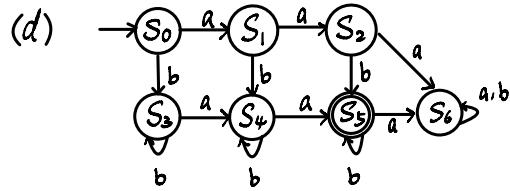
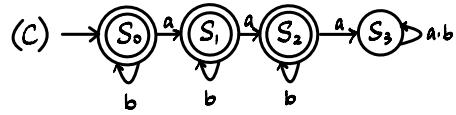
$$L \{ww^R : w \in \{a, b\}^+\}.$$

Give a complete justification for your answer.

$$G = \{S, S'\}, \{a, b\}, S, P, S' \rightarrow S \rightarrow aSa \mid bS'b, S' \rightarrow aS'a \mid bS'b \mid \lambda$$

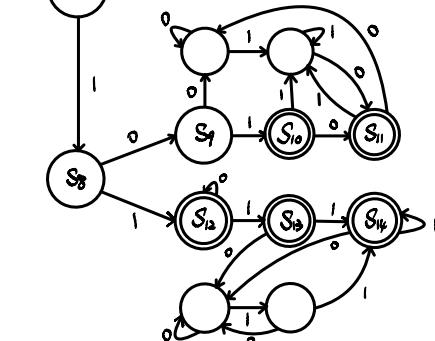
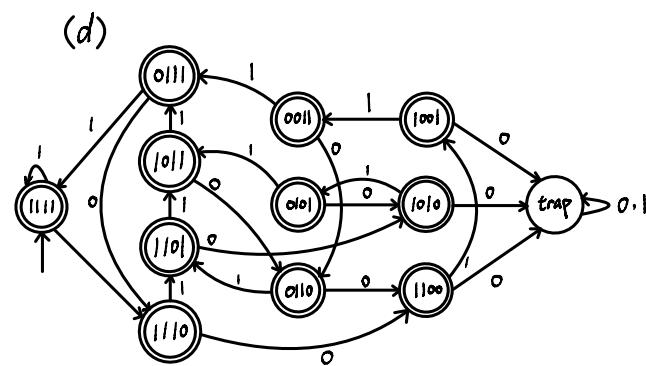
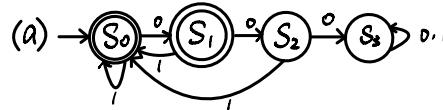
4. For $\Sigma = \{a, b\}$, construct dfa's that accept the sets consisting of

- (a) all strings with exactly one a .
- (b) all strings with at least two a 's.
- (c) all strings with no more than two a 's.
- (d) all strings with at least one b and exactly two a 's.
- (e) all the strings with exactly two a 's and more than three b 's.



11. Consider the set of strings on $\{0, 1\}$ defined by the requirements below. For each, construct an accepting dfa.

- (a) Every 00 is followed immediately by a 1. For example, the strings 101, 0010, 0010011001 are in the language, but 0001 and 00100 are not.
- (b) All strings that contain the substring 000, but not 0000.
- (c) The leftmost symbol differs from the rightmost one.
- (d) Every substring of four symbols has, at most, two 0's. For example, 001110 and 011001 are in the language, but 10010 is not because one of its substrings, 0010, contains three zeros.
- (e) All strings of length five or more in which the third symbol from the right end is different from the leftmost symbol.
- (f) All strings in which the leftmost two symbols and the rightmost two symbols are identical.

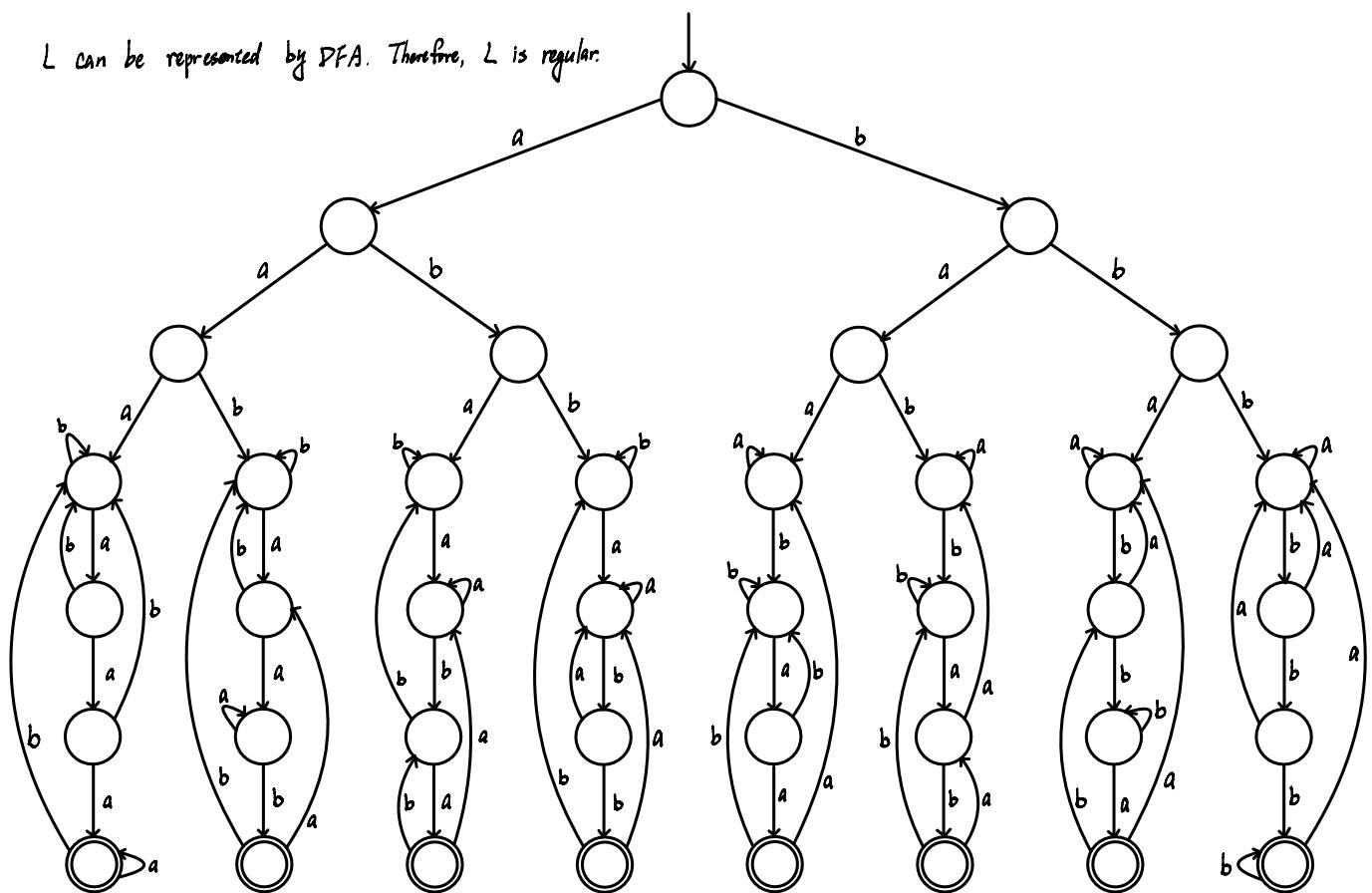


(i) $S_L = \{ \text{length } n \text{ such that } S_n = aa \text{ or } bb \in L \text{ for } n \in \{a, b\}^*\}.$
 $S_n = aS_na \text{ or } bS_nb \in L \text{ for } n \in \{a, b\}^*. \text{ Let } S_k \in L \text{ for } n \in \{a, b\}^* \Rightarrow S_{n+k} = aS_{n+k}a \text{ or } bS_{n+k}b \in L \text{ for } n \in \{a, b\}^* \Rightarrow G \subseteq L \text{ for } n \in \{a, b\}^*.$

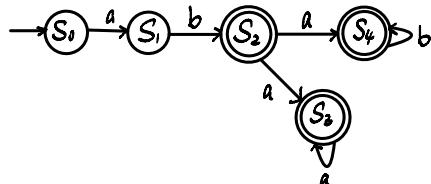
(ii) $W_L = \{ \text{length } n \text{ such that } W_n = a \text{ or } b \Rightarrow W_{n+k} = aa \text{ or } bb \in G\}.$
 $W_n = aW_na \text{ or } bW_nb \in G. \text{ Let } W_k \in G \Rightarrow W_{n+k} = aW_{n+k}a \text{ or } bW_{n+k}b \in G \Rightarrow L \text{ for } n \in \{a, b\}^* \subseteq G.$

13. Show that the language $L = \{vuv : v, w \in \{a, b\}^*, |v| = 3\}$ is regular.

L can be represented by DFA. Therefore, L is regular.



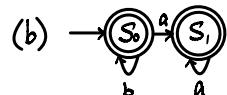
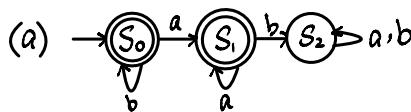
8. Design an nfa with no more than five states for the set $\{abab^n : n \geq 0\} \cup \{aba^n : n \geq 0\}$.



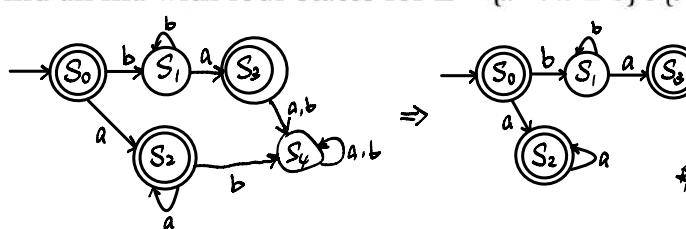
11. (a) Find an nfa with three states that accepts the language

$$L = \{a^n : n \geq 1\} \cup \{b^m a^k : m \geq 0, k \geq 0\}.$$

- (b) Do you think the language in part (a) can be accepted by an nfa with fewer than three states?

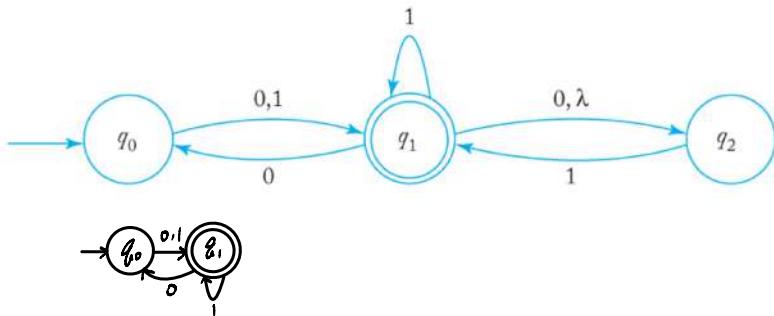


12. Find an nfa with four states for $L = \{a^n : n \geq 0\} \cup \{b^n a : n \geq 1\}$.



2. Convert the nfa in Exercise 13, Section 2.2, into an equivalent dfa.

13. Which of the strings 00, 01001, 10010, 000, 0000 are accepted by the following nfa?



4. Convert the nfa defined by

$$\delta(q_0, a) = \{q_0, q_1\}$$

$$\delta(q_1, b) = \{q_0, q_2\}$$

$$\delta(q_2, a) = \{q_1\}$$

$$\delta(q_0, \lambda) = \{q_2\}$$

with initial state q_0 and final state q_2 into an equivalent dfa.

