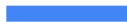


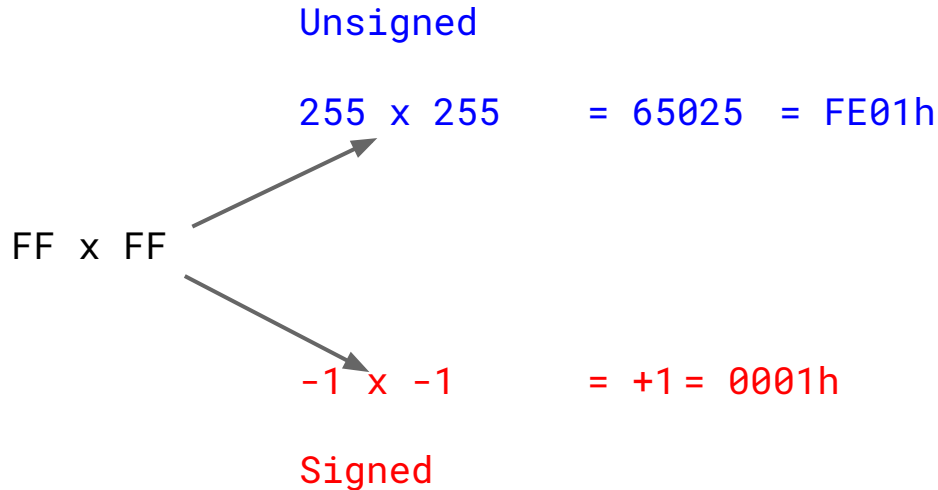
# Multiply & divide



CSC 236

# Multiply & divide

- Different from addition/subtraction
  - Separate instructions required for signed/unsigned



# Multiply & divide

- Different from addition/subtraction
  - Separate instructions required for signed/unsigned
    - User must select instruction
  - Data sizes grow (not merely overflow)

- $9 \times 9 = 81$
- $99 \times 99 = 9801$
- $999 \times 999 = 998001$
- ...
- $d \text{ digits} \times d \text{ digits} = 2d \text{ digits}$
- $d \text{ digits} \times e \text{ digits} = d+e \text{ digits}$
- 
- $\text{byte} \times \text{byte} = \text{word}$
- $\text{word} \times \text{word} = \text{double word}$

# Multiply & divide

- Different from addition/subtraction
  - Separate instructions required for signed/unsigned
    - User must select instruction
  - Data sizes grow (not merely overflow)
  - Divide is distinct operation
    - Divide is not “multiply by reciprocal”

# Instructions

- Unsigned: `mul <operand>`
- Signed: `imul <operand>`
- Only one operand
  - Multiply requires two
  - One is implicit
- Operand
  - Byte or word
  - Register or memory

# Instructions

- Unsigned: `mul <operand>`
- Signed: `imul <operand>`

al    x    byte    ↓    =    ax

# Instructions

- Unsigned: `mul <operand>`
- Signed: `imul <operand>`

al    x    byte    =    ax

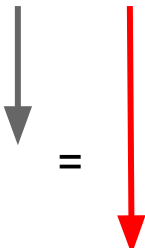
ax    x    word = double word ??

The diagram illustrates the relationship between instruction operands and register sizes. It shows two rows of text. The first row is 'al    x    byte    =    ax'. The second row is 'ax    x    word = double word ??'. A grey arrow points from the word 'byte' in the first row down to the word 'word' in the second row. A red arrow points from the 'ax' in the first row down to the words 'double word ??' in the second row.

# Instructions

- Unsigned: `mul <operand>`
- Signed: `imul <operand>`

al    x    byte    =    ax  
ax    x    word   =   dx:ax





# Two cases

- `al x byte = ax`
- `ax x word = dx:ax`

# Two cases

- $al \times \boxed{\text{byte}} = ax$
- $ax \times \boxed{\text{word}} = dx:ax$



**register or  
memory**

**cannot use  
immediate**

# Two cases

- `al x byte = ax`
- `ax x word = dx:ax`
- Overflow is not possible
- CF indicate result size
  - 0  $\Rightarrow$  significant part same size as inputs
  - 1  $\Rightarrow$  significant part larger than inputs
- `02 x 05 = 00 0A`
  - CF = 0
- `0A x 64 = 03 E8`
  - CF = 1
- **Useful for *chained* multiplies**

# Example 1

- Multiply unsigned bytes
  - 10 x 100

u10    db   10    ;0A

u100   db 100    ;64

```
mov     al,[u10]    ;ax = -- 0A
mul     [u100]      ;ax = 03 E8
                        ;CF = 1
```

## Example 2

- Multiply signed words
  - $+2 \times -1$

```
p2      dw  2      ;00 02
neg1    dw -1      ;FF FF
```

```
mov    ax,[p2]      ;dx:ax = ---- 00 02  
imul   [neg1]       ;dx:ax = FFFF FFEE  
  
;CF = ?  
;CF = 0
```

# Division

- Grammar school
  - Integers only
  - No fractions
- Divide
  - $10 / 3$
  - Answer: 3 remainder 1

# Division

- Division in high-level language
  - Quotient — `div /` : `10 / 3 = 3`
  - Remainder — `mod %` : `10 % 3 = 1`

dividend  
divisor

- Quotient: how many times can one subtract divisor from dividend
- Remainder: the amount left over (after above)

# Example

- Quotient: how many times can one subtract divisor from dividend
- Remainder: the amount left over (after these subtractions)

$$\begin{array}{r} 10 \\ \underline{3} \end{array}$$

**Can subtract 3 from 10  
3 times with  
1 left over**

**Quotient = 3  
Remainder = 1**

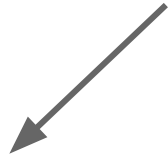


# Instructions

- Unsigned: `div <operand>`
- Signed: `idiv <operand>`
  
- Implicit operand
- Byte or word
- Register or memory

# Instructions

- Unsigned: `div <operand>`
- Signed: `idiv <operand>`



ax  
----  
byte      =    ah al  
             rem quot

# Instructions

- Unsigned: `div <operand>`
- Signed: `idiv <operand>`

ax  
----  
byte

= ah al  
rem quot

dx:ax  
-----  
word

= dx ax  
rem quot

# Overflow

- Unlike multiplication
- Can have overflow
  - If quotient > byte :  $1000/1 > \text{byte}$
  - If quotient > word:  $1000000/1 > \text{word}$

# Example 3

- unsigned\_word / byte
- 10 / 3

```
mov  ax,10      ;ax = 00 0A
mov  bl,3       ;bx = -- 03
div  bl         ;ax = 01 03
```

Remainder (ah) = 1

Quotient (al) = 3

# Notes

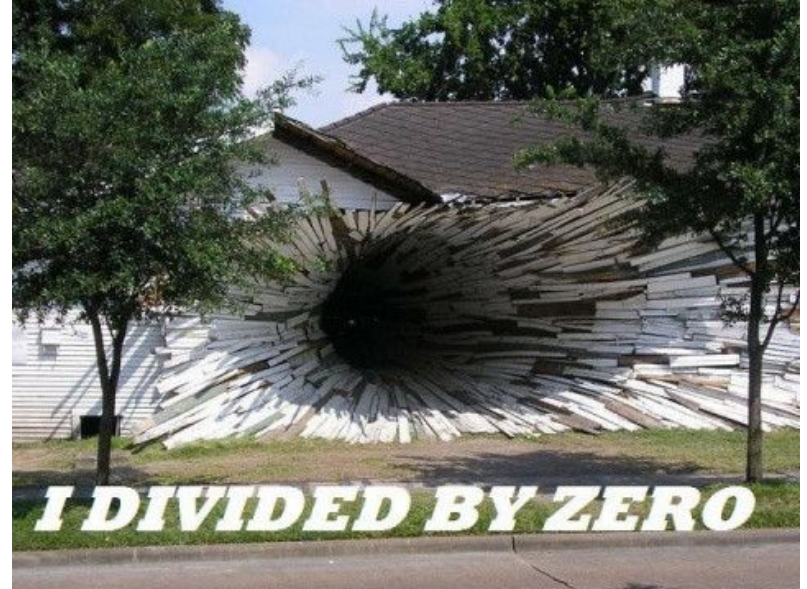
- Overflow if
  - Quotient is too big
    - alternatively : divisor too small
  - Divisor is zero

# Notes

- Overflow if
  - Quotient is too big
    - alternatively : divisor too small
  - Divisor is zero

Generates divide overflow interrupt

Program is terminated



# Notes

- Overflow if
  - Quotient is too big
    - alternatively : divisor too small
  - Divisor is zero

Generates divide overflow interrupt

Program is terminated

Not exactly: interrupt can be caught by *interrupt handler*



# Notes

- Overflow if
  - Quotient is too big
    - alternatively : divisor too small
  - Divisor is zero
- Status flags undefined after divide
- Cannot divide by immediate
  - ~~div 7~~

# Signs

Quotient & remainder are signed numbers

$$\begin{array}{r} +\underline{10} \\ + 3 \end{array}$$

$$\begin{array}{r} +\underline{10} \\ - 3 \end{array}$$

$$\begin{array}{r} -\underline{10} \\ + 3 \end{array}$$

$$\begin{array}{r} -\underline{10} \\ - 3 \end{array}$$

- Quotient
  - + for like signs
  - - for unlike signs
- Remainder
  - Sign of dividend

# Signs

Quotient & remainder are signed numbers

- Quotient
  - + for like signs
  - - for unlike signs
- Remainder
  - Sign of dividend

$$\begin{array}{r} +\underline{10} \\ + 3 \end{array}$$

$$\begin{array}{r} +\underline{10} \\ - 3 \end{array}$$

$$\begin{array}{r} -\underline{10} \\ + 3 \end{array}$$

$$\begin{array}{r} -\underline{10} \\ - 3 \end{array}$$

Q =

R =

# Signs

Quotient & remainder are signed numbers

- Quotient
  - + for like signs
  - - for unlike signs
- Remainder
  - Sign of dividend

$$\begin{array}{r} +\underline{10} \\ + 3 \end{array}$$

$$\begin{array}{r} +\underline{10} \\ - 3 \end{array}$$

$$\begin{array}{r} -\underline{10} \\ + 3 \end{array}$$

$$\begin{array}{r} -\underline{10} \\ - 3 \end{array}$$

$$Q = +3$$

$$+3$$

$$R =$$

# Signs

Quotient & remainder are signed numbers

- Quotient
  - + for like signs
  - - for unlike signs
- Remainder
  - Sign of dividend

$$\begin{array}{r} +10 \\ + 3 \end{array}$$

$$\begin{array}{r} +10 \\ - 3 \end{array}$$

$$\begin{array}{r} -10 \\ + 3 \end{array}$$

$$\begin{array}{r} -10 \\ - 3 \end{array}$$

$$Q = +3$$

$$-3$$

$$-3$$

$$+3$$

$$R =$$

# Signs

Quotient & remainder are signed numbers

- Quotient
  - + for like signs
  - - for unlike signs
- Remainder
  - Sign of dividend

$$\begin{array}{r} +10 \\ + 3 \end{array}$$

$$\begin{array}{r} +10 \\ - 3 \end{array}$$

$$\begin{array}{r} -10 \\ + 3 \end{array}$$

$$\begin{array}{r} -10 \\ - 3 \end{array}$$

$$Q = +3$$

$$-3$$

$$-3$$

$$+3$$

$$R = +1$$

$$+1$$

$$-1$$

$$-1$$

- **Quotient: how many times can one subtract divisor from dividend**
- **Remainder: the amount left over (after these subtractions)**

# Signs

Quotient & remainder are signed numbers

- Quotient
  - + for like signs
  - - for unlike signs
- Remainder
  - Sign of dividend\*

$\frac{+10}{+3}$	$\frac{+10}{-3}$	$\frac{-10}{+3}$	$\frac{-10}{-3}$
------------------	------------------	------------------	------------------

Q = +3	-3	-3	+3
--------	----	----	----

R = +1	+1	-1	-1
--------	----	----	----

**\* Exception: remainder of zero technically isn't negative, even if the dividend is.**

# Example 4

- signed\_long\_word / word
- +10 / -3

```
mov  ax,10      ;dx:ax = ---- 000A
cwd                   ;dx:ax = 0000 000A
mov  bx,-3      ;bx      =      FFFD
idiv bx         ;dx:ax = 0001 FFFD
```

Remainder (dx) = +1

Quotient (ax) = -3



# Overflow

- Addition/subtraction
  - Detect overflow after operation
- Multiply
  - Cannot overflow
- Divide
  - ??

# Overflow

- Addition/subtraction
  - Detect overflow after operation
- Multiply
  - Cannot overflow
- Divide
  - Predict overflow

# Prediction

ax      04E7      00 B3  
bl      07

ax      04E7      01 39  
bl      04

ax      04E7      02 73  
bl      02

- Which overflows?
  - If high byte (red) is not zero
- Can this be determined before divide?
  - Yes

## Legend

$04E7_{16} = 1255_{10}$   
 $1255 \text{ div } 7 = 179$   
 $1255 \text{ div } 4 = 313$   
 $1255 \text{ div } 2 = 627$

# Prediction

ax      04E7      00 B3  
bl      07

ax      04E7      01 B3  
bl      04

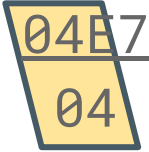
ax      04E7      02 B3  
bl      02

- Which overflows?
  - If high byte (red) is not zero
- Can this be determined before divide?
- Compare
  - High-byte of dividend and
  - Divisor

04 < 07 ⇒ no overflow

# Prediction

ax      04E7      00 B3  
bl      07

ax       04E7      01 B3  
bl      04

ax      04E7      02 B3  
bl      02

- Which overflows?
  - If high byte (red) is not zero
- Can this be determined before divide?
- Compare
  - High-byte of dividend and
  - Divisor

04 < 07 ⇒ no overflow

04 = 04 ⇒ overflow

# Prediction

ax      04E7      00 B3  
bl      07

ax      04E7      01 B3  
bl      04

ax      04E7      02 B3  
bl      02

- Which overflows?
  - If high byte (red) is not zero
- Can this be determined before divide?
- Compare
  - High-byte of dividend and
  - Divisor

04 < 07

**Rule: There will be overflow iff  
high-part-of-dividend  $\geq$  divisor**

02  $\Rightarrow$  overflow

# Example with overflow (pre) check

- Calculate quotient and remainder
  - $(a*b)/c$
  - All unsigned bytes

```
mov  al,[a]    ;a1=a
mul  [b]        ;ax=a*b
```

# Example with overflow (pre) check

- Calculate quotient and remainder

```
mov  al,[a]    ;al=a
mul  [b]        ;ax=a*b
```

- (a\*b)/c
- All unsigned bytes

$$\frac{ax}{[c]} \rightarrow \frac{ah \ al}{[c]}$$

**ah must be less than [c]**



# Example with overflow (pre) check

- Calculate quotient and remainder

- $(a*b)/c$
- All unsigned bytes

```
mov  al,[a]      ;al=a
mul  [b]          ;ax=a*b
cmp  ah,[c]       ;hi-byte ? divisor
jae  cannot      ;do not divide
div  [c]          ;ah=rem, al=quot
mov  [quot],al
mov  [rem],ah
```

# Plan ahead

```
mov  al,[a] ; ax = -- a
imul [a]    ; ax = (a)*a      = a2
imul [a]    ; ax = (a*a)*a    = a3
imul [a]    ; ax = (a*a*a)*a  = a4
```

Is this OK?

**Calculate  $a^4$**   
**where  $a$  is signed**  
**byte**  
**result is signed**  
**word**

# Plan ahead

```
mov  al,[a] ; ax = -- a
imul [a]    ; ax = (a)*a      = a2
imul [a]    ; ax = (a*a)*a    = a3
imul [a]    ; ax = (a*a*a)*a  = a4
```

```
mov  al,[a] ; ax = -- a
imul [a]    ; ax = (a)*a      = a2
imul [a]    ; ax = (a*a)*a    = a3
```

If  $a*a > \text{byte}$  this fails (overflows)

Is this OK?

# Plan ahead

```
mov  al,[a] ; ax = -- a
imul [a]    ; ax = (a)*a    = a2
imul [a]    ; ax = (a*a)*a  = a3
imul [a]    ; ax = (a*a*a)*a = a4
```

```
mov  al,[a] ; ax = -- a
imul al     ; ax = a*a   = a2
imul ax     ; ax = a2*a2 = a4
jc   error  ; cf=0 ⇒ a4 in ax
```

Simpler and safer

# Generated code

C source code

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;  
  
q1= (a * b) / c;
```

# Generated code

C source code

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;
```

```
(20*10)/3  
(200)/3
```

# Generated code

C source code

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;
```

$(20 \times 10) / 3$   
 $(200) / 3$



66

# Generated code

C source code

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;
```

```
q1= a * (b / c);
```



66



# Generated code

C source code

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;
```

```
q1= a * (b / c);
```

$20 * (10 / 3)$

$20 * (3)$



66

# Generated code

C source code

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;
```

→ 66

```
q1= a * (b / c);
```

→ 60

20\*(10/3)

20\*(3)

# Generated code 2

C source code

```
unsigned char a, b, c, q1;  
a=100; b=10; c=3;  
  
q1= (a * b) / c;
```

# Generated code 2

C source code

```
unsigned char a, b, c, q1;  
a=100; b=10; c=3;
```

```
q1= (a * b) / c;
```

$(100*10)/3 = 333$

# Generated code 2

C source code

```
unsigned char a, b, c, q1;  
a=100; b=10; c=3;
```

```
q1= (a * b) / c;
```

$(100*10)/3 = 333$



77

# Generated code 2

C source code

```
unsigned char a, b, c, q1;  
a=100; b=10; c=3;
```

```
q1= (a * b) / c;
```

$(100 * 10) / 3 = 333$   
 $= 01\ 4Dh$



77

# Generated code 2

C source code

```
unsigned char a, b, c, q1;  
a=100; b=10; c=3;
```

```
q1= (a * b) / c;
```

$(100 * 10) / 3 = 333$   
 $= 01 \text{ } \underline{4Dh}$

$$\begin{aligned} 4Dh &= 4(16) + 13 \\ &= 64 + 13 \\ &= 77 \end{aligned}$$



77

# Generated code 3

C source code

```
unsigned char a2, a3, a4, a=20;
```

```
a2=a * a;
```



# Generated code 3

C source code

```
unsigned char a2, a3, a4, a=20;
```

```
a2=a * a;
```

20 \* 20 = 400



144

# Generated code 3

C source code

```
unsigned char a2, a3, a4, a=20;
```

```
a2=a * a;
```

20 \* 20 = 400

$$400_{10} = 01\underline{90}_{16}$$

$$90_{16} = 2^7 + 2^4 = 128 + 16$$



144

# Generated code 3

C source code

```
unsigned char a2, a3, a4, a=20;
```

```
a2=a * a;
```

```
a3=a * a * a;
```

$$\begin{aligned}8000_{10} &= 20_{10} * 20_{10} * 20_{10} \\ &= 1F40_{16} \\ 40_{16} &= 64_{10}\end{aligned}$$

144

64

# Generated code 3

C source code

$$160000_{10} = 0271\underline{00}_{16}$$

```
unsigned char a2, a3, a4, a=20;
```

```
a2=a * a;
```

```
a3=a * a * a;  144
```

```
a4=a * a * a * a;  64
```

```
 0
```

# Generated code 4

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;  
q1= a * (b / c);
```

```
a=100;  
q1= (a * b) / c;
```

```
unsigned char a2, a3, a4; a=20;
```

```
a2=a * a;  
a3=a * a * a;  
a4=a * a * a * a;
```

- C source
  - 9 multiplies
  - 3 divides
- How many did compiler emit?

# Generated code 4

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;  
q1= a * (b / c);
```

```
a=100;  
q1= (a * b) / c;
```

```
unsigned char a2, a3, a4; a=20;
```

```
a2=a * a;  
a3=a * a * a;  
a4=a * a * a * a;
```

- C source
  - 9 multiplies
  - 3 divides
- How many did compiler emit?
  - > 12
  - < 12

# Generated code 4

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;  
q1= a * (b / c);
```

```
a=100;  
q1= (a * b) / c;
```

```
unsigned char a2, a3, a4; a=20;
```

```
a2=a * a;  
a3=a * a * a;  
a4=a * a * a * a;
```

- C source
  - 9 multiplies
  - 3 divides
- How many did compiler emit?
  - > 12
  - < 12
- It emitted
  - Zero (0) multiplies
  - Zero (0) divides
- ???

# Generated code 4

```
unsigned char a, b, c, q1;  
a=20; b=10; c=3;
```

```
q1= (a * b) / c;  
q1= a * (b / c);
```

```
a=100;  
q1= (a * b) / c;
```

```
unsigned char a2, a3, a4; a=20;
```

```
a2=a * a;  
a3=a * a * a;  
a4=a * a * a * a;
```

- C source
  - 9 multiplies
  - 3 divides
- How many did compiler emit?
  - > 12
  - < 12
- It emitted
  - Zero (0) multiplies
  - Zero (0) divides
- Compiler
  - Knows all values at compile time
  - Creates result itself
  - $q1=(a*b)/c \Rightarrow \text{mov } [q1], 66$



# Test 1

Cover through lecture 10

- Number systems
- Basic architecture
- 8086 architecture
- Assembler instructions
- Declaring data
- Flow control and looping
- File I/O

---

# Test 1 — Open book

- You may use
  - Any printed material from the course Moodle page.
  - Anything printed material you created (eg, notes, cheatsheets, etc)
  - Specifically, you can use
    - Old tests
    - Homeworks
    - Your own class notes
    - Class slides
    - Class text
- You may not use a calculator, phone, computing device

# Test 1 — Open book

- Most effective way to prepare
  - Take the old tests -- which are on the web site
- Test covers
  - Lectures up to and including lecture 10
  - Class notes up to and including chapter 8
- 35-40 multiple choice questions
  - ~15 number systems
  - ~5 architecture
  - ~15 8086 assembler instructions, data declarations, flow control, IO