



Week 1 Quiz

TOTAL POINTS 10

1. According to video lesson 1.1.1, a hash table consists of three things. Which of these was NOT one of those three things? 1 point

- ☐ Encryption
- ☐ An array
- ☐ Collision handling
- ☐ A hash function

2. Given a hash function $h(\text{key})$ that produces an index into an array of size N , and given two different key values key1 and key2 , the Simple Uniform Hashing Assumption states which of the following? 1 point

- ☐ The probability that $h(\text{key1}) == h(\text{key2})$ is 0.
- ☐ If $h(\text{key1}) == h(\text{key2})$ then h needs a running time of $O(N)$ to complete.
- ☐ If $h(\text{key1}) == h(\text{key2})$ then h needs a running time of $O(\lg N)$ to complete.
- ☐ The probability that $h(\text{key1}) == h(\text{key2})$ is $1/N$.

3. According to video lesson 1.1.2, which of the following is a *good* hash function $h(\text{key})$ that translates any 32-bit unsigned integer key into an index into an 8 element array? 1 point

(Note that an expression like " $2 \& 3$ " uses the bitwise-AND operator, which gives the result of comparing every bit in the two operands using the concept of "AND" from Boolean logic; for example, in Boolean logic with binary numbers, 10 AND 11 gives 10: for the first digit, 1 AND 1 yields 1, while for the second digit, 0 AND 1 yields 0. An expression like " $4 \% 8$ " uses the remainder operator that give the remainder from integer division; for example, $4 \% 8$ yields 4, which is the remainder of $4/8$. In some cases, these two different operators give similar results. Think about why that is.)

- ☐

```
1  int h(uint key) {
2      int index = 5;
3      while (key--)
4          index = (index + 5) % 8
5      return index;
6  }
7  }
```
- ☐

```
1  int h(uint key) { return key & 7; }
```
- ☐

```
1  int h(uint key) { return rand() % 8; }
```
- ☐

```
1  int h(uint key) { return max(key,7); }
```

4. Suppose you have a good hash function $h(\text{key})$ that returns an index into an array of size N . If you store values in a linked list in the array to manage collisions, and you have already stored n values, then what is the expected run time to store a new value into the hash table? 1 point

- ☐ $O(n)$
- ☐ $O(N)$
- ☐ $O(n/N)$
- ☐ $O(1)$

5. Suppose you have a good hash function $h(\text{key})$ that returns an index into an array of size N . If you store values in a linked list in the array to manage collisions, and you have already stored n values, then what is the expected run time to find the value in the hash table corresponding to a given key? 1 point

- ☐ $O(n/N)$
- ☐ $O(N)$
- ☐ $O(n)$
- ☐ $O(1)$

6. Which one of the following four hashing operations would run faster than the others? 1 point

- ☐ Finding a value in a hash table of 4 values stored in an array of 8 elements.
- ☐ Finding a value in a hash table of 20 values stored in an array of 100 elements.
- ☐ Finding a value in a hash table of 100 values stored in an array of 1,000 elements.
- ☐ Finding a value in a hash table of 2 values stored in an array of 2 elements.

7. When storing a new value in a hash table, linear probing handles collisions by finding the next unfilled array element. Which of the following is the main drawback of linear probing? 1 point

- ☐ There may not be an available slot in the array.
- ☐ If the hash function returns an index near the end of the array, there might not be an available slot before the end of the array is reached.
- ☐ Even using a good hash function, contiguous portions of the array will become filled, causing a lot of additional probing in search of the next available unused element in the array.
- ☐ The array only stores values, so when retrieving the value corresponding to a key, there is no way to know if the value at $h(\text{key})$ is the proper value, or if it is one of the values at a subsequent array location.

8. When using double hashing to store a value in a hash table, if the hash function returns an array location that already stores a previous value, then a new array location is found as the hash function of the current array location. Why? 1 point

- ☐ Only one additional hash function is called to find an available slot in the array whereas linear probing requires an unknown number of array checks to find an available slot.
- ☐ Double hashing reduces the chance of a hash function collision on subsequent additions to the hash table.
- ☐ Double hashing reduces the clumping that can occur with linear probing.
- ☐ Since the hash function runs in constant time, double hashing runs in $O(1)$ time.

9. Which of the following data structures would be the better choice to implement a memory cache, where a block of global memory (indicated by higher order bits of the memory address) are mapped to the location of a block of faster local memory. 1 point

- ☐ An AVL tree.
- ☐ A hash table implemented with double hashing

☐ A hash table implemented with double hashing.

☐ A hash table implemented with separate chaining, using an array of linked lists.

☐ A hash table implemented with linear probing.

10. Which of the following data structures would be the better choice to implement a dictionary that not only returns the definition of a word but also returns the next word following that word (in lexical order) in the dictionary.

1 point

☐ A hash table implemented with double hashing.

☐ An AVL tree.

☐ A hash table implemented with separate chaining, using an array of linked lists.

☐ A hash table implemented with linear probing.



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