

α	type variables	
x	variables	
i	integer literals	
τ, σ	$::=$	types
	α	
	\mathbb{Z}	
	void	
	$\tau_1 \rightarrow \tau_2$	
	$\forall \alpha. \tau$	bind α in τ
	$\tau_1 \times \tau_2$	
	(τ)	S
v	$::=$	annotated values
	u^τ	
u	$::=$	values
	x	
	i	
	$\lambda(x : \tau).v$	bind x in v
	$\lambda(x : \tau).e$	bind x in e
	$v_1 v_2$	
	$v_1 \times v_2$	
	$v \cdot \mathbf{1}$	
	$v \cdot \mathbf{r}$	
	$v_1 p v_2$	
	if0 (v_1, v_2, v_3)	
	(u)	S
p	$::=$	primitives
	$+$	
	$-$	
d	$::=$	declarations
	v	
	$v \cdot \mathbf{1}$	
	$v \cdot \mathbf{r}$	
	$v_1 p v_2$	
e	$::=$	terms
	let $x = d$ in e	bind x in e
	$v'(v)$	
	if0 (v, e_1, e_2)	
	halt $[\tau]v$	
Δ	$::=$	type contexts
	Δ, α	
Γ	$::=$	value contexts

| $\Gamma, x : \tau$

$\boxed{\Delta \vdash_F \tau}$ type formation

$$\begin{array}{c}
\frac{\alpha \in \Delta}{\Delta \vdash_F \alpha} \quad \text{F_TYPE_VAR} \\
\\
\frac{}{\Delta \vdash_F \mathbb{Z}} \quad \text{F_TYPE_INT} \\
\\
\frac{\Delta \vdash_F \tau_1 \quad \Delta \vdash_F \tau_2}{\Delta \vdash_F \tau_1 \rightarrow \tau_2} \quad \text{F_TYPE_ARR} \\
\\
\frac{\Delta, \alpha \vdash_F \tau}{\Delta \vdash_F \forall \alpha. \tau} \quad \text{F_TYPE_ALL} \\
\\
\frac{\Delta \vdash_F \tau_1 \quad \Delta \vdash_F \tau_2}{\Delta \vdash_F \tau_1 \times \tau_2} \quad \text{F_TYPE_PROD}
\end{array}$$

$\boxed{\Delta; \Gamma \vdash_F v : \tau}$ annotated typing

$$\frac{\Delta; \Gamma \vdash_F u : \tau}{\Delta; \Gamma \vdash_F u^\tau : \tau} \quad \text{F_ANV_ANN}$$

$\boxed{\Delta; \Gamma \vdash_F u : \tau}$ typing

$$\begin{array}{c}
\frac{\Delta \vdash_F \tau \quad \Gamma(x) = \tau}{\Delta; \Gamma \vdash_F x : \tau} \quad \text{F_VALUE_VAR} \\
\\
\frac{}{\Delta; \Gamma \vdash_F i : \mathbb{Z}} \quad \text{F_VALUE_INT} \\
\\
\frac{\Delta \vdash_F \tau_1 \quad \Delta; \Gamma, x_1 : \tau_1 \vdash_F v : \tau_2}{\Delta; \Gamma \vdash_F \lambda(x_1 : \tau_1). v : \tau_1 \rightarrow \tau_2} \quad \text{F_VALUE_LAM} \\
\\
\frac{\Delta; \Gamma \vdash_F v_1 : \tau_1 \rightarrow \tau_2 \quad \Delta; \Gamma \vdash_F v_2 : \tau_1}{\Delta; \Gamma \vdash_F v_1 v_2 : \tau_2} \quad \text{F_VALUE_APP} \\
\\
\frac{\Delta; \Gamma \vdash_F v_1 : \tau_1 \quad \Delta; \Gamma \vdash_F v_2 : \tau_2}{\Delta; \Gamma \vdash_F v_1 \times v_2 : \tau_1 \times \tau_2} \quad \text{F_VALUE_PAIR} \\
\\
\frac{\Delta; \Gamma \vdash_F v : \tau_1 \times \tau_2}{\Delta; \Gamma \vdash_F v \cdot \mathbf{1} : \tau_1} \quad \text{F_VALUE_PRL} \\
\\
\frac{\Delta; \Gamma \vdash_F v : \tau_1 \times \tau_2}{\Delta; \Gamma \vdash_F v \cdot \mathbf{r} : \tau_2} \quad \text{F_VALUE_PRR} \\
\\
\frac{\Delta; \Gamma \vdash_F v_1 : \mathbb{Z} \quad \Delta; \Gamma \vdash_F v_2 : \mathbb{Z}}{\Delta; \Gamma \vdash_F v_1 p v_2 : \mathbb{Z}} \quad \text{F_VALUE_PRIM}
\end{array}$$

$$\frac{\Delta; \Gamma \vdash_F v_1 : \mathbb{Z} \quad \Delta; \Gamma \vdash_F v_2 : \tau \quad \Delta; \Gamma \vdash_F v_3 : \tau}{\Delta; \Gamma \vdash_F \text{if0}(v_1, v_2, v_3) : \tau} \quad \text{F_VALUE_IF0}$$

$\boxed{\Delta \vdash_K \tau}$ type formation

$$\frac{\alpha \in \Delta}{\Delta \vdash_K \alpha} \quad \text{K_TYPE_VAR}$$

$$\frac{}{\Delta \vdash_K \mathbb{Z}} \quad \text{K_TYPE_INT}$$

$$\frac{\Delta \vdash_K \tau}{\Delta \vdash_K \tau \rightarrow \text{void}} \quad \text{K_TYPE_ARR}$$

$$\frac{\Delta \vdash_K \tau_1 \quad \Delta \vdash_K \tau_2}{\Delta \vdash_K \tau_1 \times \tau_2} \quad \text{K_TYPE_PROD}$$

$\boxed{\Delta; \Gamma \vdash_K v : \tau}$ annotated typing

$$\frac{\Delta; \Gamma \vdash_K u : \tau}{\Delta; \Gamma \vdash_K u^\tau : \tau} \quad \text{K_ANV_ANN}$$

$\boxed{\Delta; \Gamma \vdash_K u : \tau}$ typing

$$\frac{\Delta \vdash_K \tau \quad \Gamma(x) = \tau}{\Delta; \Gamma \vdash_K x : \tau} \quad \text{K_VALUE_VAR}$$

$$\frac{}{\Delta; \Gamma \vdash_K i : \mathbb{Z}} \quad \text{K_VALUE_INT}$$

$$\frac{\Delta; \Gamma, x : \tau \vdash_K e}{\Delta; \Gamma \vdash_K \lambda(x : \tau).e : \tau \rightarrow \text{void}} \quad \text{K_VALUE_LAM}$$

$$\frac{\Delta; \Gamma \vdash_K v_1 : \tau_1 \quad \Delta; \Gamma \vdash_K v_2 : \tau_2}{\Delta; \Gamma \vdash_K v_1 \times v_2 : \tau_1 \times \tau_2} \quad \text{K_VALUE_PAIR}$$

$\boxed{\Delta; \Gamma \vdash_K e}$ term formation

$$\frac{\Delta; \Gamma \vdash_K v : \tau \quad \Delta; \Gamma, x : \tau \vdash_K e}{\Delta; \Gamma \vdash_K \text{let } x = v \text{ in } e} \quad \text{K_TERM_LET}$$

$$\frac{\Delta; \Gamma \vdash_K v : \tau_1 \times \tau_2 \quad \Delta; \Gamma, x : \tau_1 \vdash_K e}{\Delta; \Gamma \vdash_K \text{let } x = v \cdot 1 \text{ in } e} \quad \text{K_TERM_PRL}$$

$$\frac{\Delta; \Gamma \vdash_K v : \tau_1 \times \tau_2 \quad \Delta; \Gamma, x : \tau_2 \vdash_K e}{\Delta; \Gamma \vdash_K \text{let } x = v \cdot \mathbf{r} \text{ in } e} \quad \text{K_TERM_PRR}$$

$$\frac{\Delta; \Gamma \vdash_K v_1 : \mathbb{Z} \quad \Delta; \Gamma \vdash_K v_2 : \mathbb{Z} \quad \Delta; \Gamma, x : \mathbb{Z} \vdash_K e}{\Delta; \Gamma \vdash_K \text{let } x = v_1 \text{ p } v_2 \text{ in } e} \quad \text{K_TERM_PRIM}$$

$$\begin{array}{c}
\frac{\Delta; \Gamma \vdash_K v' : \tau \rightarrow \mathbf{void} \quad \Delta; \Gamma \vdash_K v : \tau}{\Delta; \Gamma \vdash_K v'(v)} \quad \text{K_TERM_APP} \\
\\
\frac{\Delta; \Gamma \vdash_K v : \mathbb{Z} \quad \Delta; \Gamma \vdash_K e_1 \quad \Delta; \Gamma \vdash_K e_2}{\Delta; \Gamma \vdash_K \text{if0}(v, e_1, e_2)} \quad \text{K_TERM_IF0} \\
\\
\frac{\Delta; \Gamma \vdash_K v : \tau}{\Delta; \Gamma \vdash_K \text{halt}[\tau]v} \quad \text{K_TERM_HALT}
\end{array}$$

Definition rules: 31 good 0 bad
 Definition rule clauses: 78 good 0 bad