

# Continuation-Passing Style

## Syntax

$x$	identifiers	
$e, K$	$::=$	terms
	$\mid$	$t$
	$\mid$	$s$
$t$	$::=$	trivial terms, <i>i.e.</i> values
	$\mid$	$x$
	$\mid$	$\lambda x. e$ bind $x$ in $e$
$s$	$::=$	serious terms, <i>i.e.</i> computations
	$\mid$	$e_0 e_1$

## Transformation

### Call-by-value

$$\begin{aligned} \llbracket x \rrbracket K &= K \ x \\ \llbracket \lambda x. e \rrbracket K &= K \ (\lambda x \ k. \llbracket e \rrbracket k) \\ \llbracket t_0 \ t_1 \rrbracket K &= \llbracket t_0 \rrbracket \llbracket t_1 \rrbracket K \\ \llbracket t_0 \ s_1 \rrbracket K &= \llbracket s_1 \rrbracket (\lambda x_1. \llbracket t_0 \rrbracket x_1 \ K) \\ \llbracket s_0 \ t_1 \rrbracket K &= \llbracket s_0 \rrbracket (\lambda x_0. x_0 \llbracket t_1 \rrbracket K) \\ \llbracket s_0 \ s_1 \rrbracket K &= \llbracket s_0 \rrbracket (\lambda x_0. \llbracket s_1 \rrbracket (\lambda x_1. x_0 \ x_1 \ K)) \end{aligned}$$

### Call-by-name

$$\begin{aligned} \llbracket x \rrbracket K &= K \ x \\ \llbracket \lambda x. e \rrbracket K &= K \ (\lambda x \ k. \llbracket e \rrbracket k) \\ \llbracket t_0 \ e_1 \rrbracket K &= \llbracket t_0 \rrbracket \llbracket e_1 \rrbracket K \\ \llbracket s_0 \ e_1 \rrbracket K &= \llbracket s_0 \rrbracket (\lambda x_0. x_0 \llbracket e_1 \rrbracket K) \end{aligned}$$