Distributed Systems Andrew Tanenbaum + Maarten Steen

## Building Microservices Ch4##

Communication styles, pros/cons

Osynchronous blocking (a) asynchronous nonblocking

nequest-nesponse Collab Devent-driven collab In-Process, calls win a single process

inter-process, calls blu different processes across a network

a) Performance

- in-process > inter-process, lu compiler can make optimizations may cause you to network your API if you're company in vs. inter-

- Need to be aware often size of your params ble thuy will be passed over a network blu microservices & senalited / deservatived

- Need to be aware of unen a call is happening oner a network

W Changing Interfaces

- The Jisenice exposing the interface & the Jisenice using the interface are deployed suparately so breaking rollouts to the interface need to either.

c) Error Handling

- in-process le errors are welly deterministic & relatively straightform - inter-process you can have a host of down stream or network issues

\* crash failure (server crash)

\* omission failure (sent something, no response or events stop coming) \* timing failure (too late or too early)

\* nesponse failure (rusponse received, but incomplete or wrong info)

\* arbitrary failure (axa Byzantine failure, something has gone wrong but no one can agree on if/uly the failure occurred

-> may be transient errors

or ofcuruise so huat clients can respond correctly

internal restrict

request - response

RESTOVEY: - Queue based

HTTP - PPC

ex: Topic -based brokers

common data . W. db orfilesystem

- REST OVER HTTP

Asynchronous nonblocking.

inter-uservice comms contid. · Synchronous blocking, make call + wait To Async & nonblocking, send call edon't wait · Request-nesponse, make call, wait, expectnesponse · Event-driven, send event & don't cane who listens Lo Common data, shane a data source understand needs around · reliable comms · acceptable latency · vol. of comms. · security nucls Ofirst choose b/w nequest-nesponce & event driven & you can mix & match Communication styles Sync blocking + easily understood - temporal coupling blw instances of wervices - potential cascading issues ble of outages - resource contention in chained blocking caus async nonblocking Snequest-response: make nequest then any instance can receive nesponse Common data: upstream juserice Changes common data to be used by Vitream Jusenice Went-driven: Upstream broadcasts event, downstream can listen t temporal decoupling - 1 complexity + handy if triggered work is slow - 1 choices \*async lawait functions use a sync blocking call impl: common data data war lake, dump into in data wanehouse, uservice would need to know smicture of wavehouse > for both data is assumed to run in one direction upstream to downstream files, in policy to check for new Mes - Matercy + simple - potential coupling + interoperability (old & new systems, etc.) - potially more point of failure lev. Mu filesystem It good if you're sending a lot itsey) ot data at once

Pattern: neguest / response

can be blocking or nonblocking (ex-winotification responses)

implementation: sync us async

Sync., open network connection ende until response is received.

async: send neguest toquene, workis done by consumer austrum publishes to another quene teratille work is complete. The senices both need to know when to head esend responses to. The queue can hold many nequests until the consumer is neady for them. Can't nely on talking of the same usenice instance that cheater the nequest so menant state inforward need to be housed blockness.

\* you will need to handle timeouts for sync or async

\* can run calls in parallel vs. in sequence to Ilatency, asynchawait can hap

- good for when aresponse is needed before more work can be done

you need to know of a call didn't work, can netry

## pattern: Event-driven comms

· Mservice emits event about something leat happened. Others may or may not listen.

· event: statement about something that has occurred, typically win the emitting usen.

· pushes the nesponsibility of knowing must to do inter face of an ement to the listening menices

· <u>luent</u> is a tuing that happened, a <u>message</u> is a tuing me send over an asyne comms mehanism

The message is the medium, the event is the payload

implementation

· Producers use an API to publish an event to the broker

Broker handus subscriptions & possibly want misg. In Consumer has seen

- "Keep your middleware dumb + keep the marts in the endpoints
- · Can also use an HTTP Feed, but no good for I latency lex. Atom thom
- \* be wany of the sunk cost fallacy, try Atomityon mud it but suiten to a message broker if you need to
- · What should be in the event? just an ID, which may result in a barrage of neguests for more information

or Fully Detailed events, too everysuing you'd be happy to share via an API

Can & coupling + act as a historical necord. Just water how large we mig is

Enter prise Integration
Patterns by Gregor
Pompe & Bobby Wood

## matshould be in the event contid

could break others

· You could also send two events, one we sensitive into e one wo (ox. Pl)

· need to manage visibility + facult tolerance (dying to before sending 2nd mag)

· once you put data in an event it's a part of your contract. Removing it

lume to use it?

But 1 complexity, you can use a mix of styles though.

a message hospital (or dead letter queue) + max netries may help avoid Catastrophic failovers were your workers die one by one trying to handle a bad huseage lunt crashes them

and the second of the second o

· ensure & monitoring & Lusing correlation IDs