# 14. SQL queries

Principles of Data Science with R

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#### **Announcement**

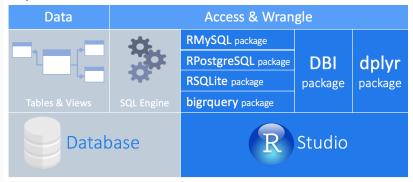
- 1. Quiz 6 this week on Canvas
- 2. HW will be released tomorrow as usual, due next Wednesday

#### Next we will see...

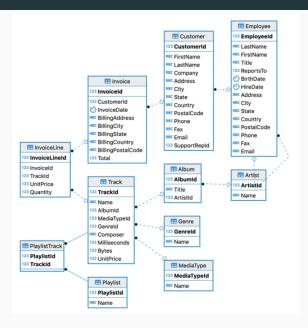
- Integrity constraints (Primary and Foreign Keys)
- More SQL queries

## **SQLite RDBMS and Rstudio**

# **Open Source Databases**



# ChinookDB Entity-Relationship Diagram (ER Diagram)



### Connecting to a DB

- (Install) and load the DBI (Database interface) package
- (Install) and load DBI complaint DataBase Connectivity driver package for the database you will be using
  - For SQLite RDBMS, we will use the SQLite() driver from the RSQLite R package
  - For Postgres RDBMS, use the Postgres() driver from the RPostgres package

chinook db is an R object that represents a connection to the database file Chinook\_Sqlite.sqlite

#### Field metadata

Unlike a data frame, there is extra information in a database table that expresses relational information between tables.

dbGetQuery(chinook\_db, "pragma table\_info(Customer)")

##		cid	name	type	${\tt notnull}$	${\tt dflt\_value}$	pk
##	1	0	CustomerId	INTEGER	1	NA	1
##	2	1	FirstName	NVARCHAR(40)	1	NA	0
##	3	2	LastName	NVARCHAR(20)	1	NA	0
##	4	3	Company	NVARCHAR(80)	0	NA	0
##	5	4	Address	NVARCHAR(70)	0	NA	0
##	6	5	City	NVARCHAR(40)	0	NA	0
##	7	6	State	NVARCHAR(40)	0	NA	0
##	8	7	Country	NVARCHAR(40)	0	NA	0
##	9	8	PostalCode	NVARCHAR(10)	0	NA	0
##	10	9	Phone	NVARCHAR(24)	0	NA	0
##	11	10	Fax	NVARCHAR(24)	0	NA	0
##	12	11	Email	NVARCHAR(60)	1	NA	0
##	13	12	SupportRepId	INTEGER	0	NA	0

#### Primary key

The **primary key** is a *unique identifier* of the rows in a table.

Two rows cannot have the same primary key

```
## CustomerId FirstName LastName City Country
## 1 1 Luís Gonçalves São José dos Campos Brazil
## 2 2 Leonie Köhler Stuttgart Germany
```

## Error: UNIQUE constraint failed: Customer.CustomerId

CustomerId is the primary key and must be unique.

## Multi-column primary key

cid

## 2

d name type
O PlaylistId INTEGER

TrackId INTEGER

Primary key's can consist of multiple columns if it takes multiple columns to identify a row in a table. But, two rows cannot have the same primary key.

```
# Single column primary key
dbGetQuery(chinook_db, "pragma table_info(Customer)")
##
      cid
                              type notnull dflt_value pk
                 name
           CustomerId
                           INTEGER.
## 1
                                                   NA 1
            FirstName NVARCHAR(40)
## 2
                                                   NA O
## 3
            LastName NVARCHAR(20)
                                                  NA O
## 4
            Company NVARCHAR(80)
                                                  NA O
## 5
            Address NVARCHAR(70)
                                                  NA O
## 6
               City NVARCHAR(40)
                                                  NA O
## 7
                State NVARCHAR(40)
                                                  NA O
              Country NVARCHAR(40)
                                                  NA O
## 8
## 9
       8 PostalCode NVARCHAR(10)
                                                  NA O
## 10
              Phone NVARCHAR(24)
                                                  NA O
                Fax NVARCHAR(24)
                                                  NA O
## 11
      10
## 12
               Email NVARCHAR(60)
                                                   NA O
     12 SupportRepId
                           INTEGER
                                                   NA O
# Multi column primary key
dbGetQuerv(chinook db, "pragma table info(PlayListTrack)")
```

NA 2

type notnull dflt\_value pk

# Primary key

Tables are not required to have a primary key, but most do. All the tables in Chinook have a primary key.



# Foreign keys

The relationship between tables is expressed by primary keys and **foreign keys**.

Remember we are working with a relational database, following a relational data model.



A foreign key field points to the primary key of another table.

#### Foreign keys

Foreign keys must either point to an existing value or be NULL.

### To enforce Foreign key constraints in SQLite RDBMS

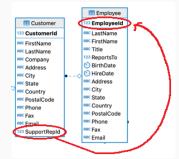
```
# Required for foreign-key support otherwise foreign keys are not enforced
dbExecute(chinook_db, "pragma foreign_keys = on")
dbGetQuery(chinook_db,
           "SELECT max(EmployeeId) FROM Employee")
    max(EmployeeId)
## 1
dbGetQuery(chinook_db,
           "SELECT max(CustomerId) FROM Customer")
    max(CustomerId)
## 1
dbExecute(chinook_db,
    "INSERT INTO Customer
    (CustomerId, FirstName, LastName, Email, SupportRepId)
    VALUES
    (59, 'Luis', 'Armstrong', 'luisArmstrong@pstat.ucsb.edu', 88)")
```

## Error: UNIQUE constraint failed: Customer.CustomerId

```
dbExecute(chinook_db,
  "INSERT INTO Customer
  (CustomerId, FirstName, LastName, Email, SupportRepId)
  VALUES
  (60, 'Luis', 'Armstrong', 'luisArmstrong@pstat.ucsb.edu', 10)")
```

## Error: FOREIGN KEY constraint failed

# Interpretation of foreign key



- Each customer in Customer table can be assigned a support representative
- The support rep is an employee at the store and therefore has a unique id, Employeeld
- This unique id, Employeeld, is the primary key of the employee table

Thus real-world relationship is encoded by the relational model using primary and foreign key relationships.

## **Integrity Constraints**

We have seen two examples of integrity constraints:

- Primary keys must be unique (and not NULL)
- Foreign keys must reference existing primary keys or be NULL

These constraints enforce the *integrity* of a database: no bad data or corrupted relationships.

Keys help maintain the integrity of the data

#### **Database Schema**

The **schema** of a database describes its *structure*:

- Names of all the tables
- Names of all fields in each table
- Primary key/foreign key relationships between tables
- Other metadata (data types of each field in each table, ...)

Basically everything other than the actual data itself.

Represented via E-R diagrams (Entity relationship)

We have been looking at parts of the schema with the pragma keyword.

```
dbGetQuery(chinook_db, "pragma table_info(customer)")
```

# More SQL queries

#### SELECT

6852860 6599424 8611245 ....

```
dbGetQuery(chinook_db,
                       "SELECT count(*) FROM track")
    count(*)
## 1
         3503
What are all the fields for every track?
dbListFields(chinook_db, "track")
## [1] "TrackId" "Name" "AlbumId" "MediaTypeId" "GenreId"
## [6] "Composer" "Milliseconds" "Bytes" "UnitPrice"
track sel <- dbGetQuerv(chinook db.
                       "SELECT * FROM track")
str(track_sel)
## 'data.frame': 3503 obs. of 9 variables:
## $ TrackId : int 1 2 3 4 5 6 7 8 9 10 ...
## $ Name : chr "For Those About To Rock (We Salute You)" "Balls to the Wall"
"Fast As a Shark" "Restless and Wild" ...
## $ AlbumTd : int 1 2 3 3 3 1 1 1 1 1 ...
## $ MediaTypeId : int 1 2 2 2 2 1 1 1 1 1 ...
## $ GenreId : int 1 1 1 1 1 1 1 1 1 ...
## $ Composer : chr "Angus Young, Malcolm Young, Brian Johnson" NA "F. Baltes,
S. Kaufman, U. Dirkscheider & W. Hoffman" "F. Baltes, R.A. Smith-Diesel, S.
Kaufman, U. Dirkscneider & W. Hoffman" ...
## $ Milliseconds: int 343719 342562 230619 252051 375418 205662 233926 210834
203102 263497
## $ Bytes : int 11170334 5510424 3990994 4331779 6290521 6713451 7636561
```

# Suppose we only want the first five records for TrackId, Name, AlbumId, Milliseconds, Bytes, UnitPrice from Track table

## 5 0.99

## SELECT, expanded

In the first line of SELECT, we can directly specify computations that we want performed

FROM table
WHERE condition
GROUP BY columns
HAVING condition
ORDER BY column [ASC | DESC]
LIMIT offset, count;

Main tools for computations:

MIN, MAX, COUNT, SUM, AVG or any math formula

## Example

To calculate the average Milliseconds, Bytes and Max UnitPrice

```
## 1 393599.2 33510207 1.99
```

To replicate this simple command on an imported data frame:

```
mean(track_sel$Milliseconds, na.rm=TRUE)

## [1] 393599.2
mean(track_sel$Bytes, na.rm=TRUE)

## [1] 33510207
max(track_sel$UnitPrice, na.rm=TRUE)

## [1] 1.99
```

We can use the GROUP BY option in SELECT to define aggregation groups

```
AlbumId AVG(Bytes)
##
         253 536359244
## 1
         229 535292434
## 2
## 3
         227 529469291
## 4
         231 514373372
         228 512231374
## 5
## 6
         254 492670102
## 7
         226 490750393
         261 453454450
## 8
         251 306109250
## 9
## 10
         249 268393262
```

(Note: the order of commands here matters; try switching the order of GROUP BY and ORDER BY, you'll get an error)

# We can use AS in the first line of SELECT to rename computed columns

```
AlbumId AvgBytes
##
          253 536359244
## 1
## 2
         229 535292434
## 3
          227 529469291
## 4
         231 514373372
## 5
         228 512231374
         254 492670102
## 6
## 7
         226 490750393
## 8
          261 453454450
## 9
          251 306109250
## 10
          249 268393262
```

## 1

We can use the WHERE option in SELECT to specify a subset of the rows to use (pre-aggregation/pre-calculation)

```
##
      AlbumId MediaTypeId AvgBytes
## 1
           50
                        1 30444082
                       1 24822832
## 2
          138
## 3
          137
                       1 19120969
## 4
          43
                       1 16221538
## 5
         97
                       1 16089011
## 6
          114
                       1 15975057
          109
                       1 15934275
## 7
## 8
          113
                       1 15521017
          127
## 9
                        1 15194926
```

AlbumId Avg(Bytes)

30444082

50

We can use the HAVING option in SELECT to specify a subset of the rows to display (post-aggregation/post-calculation)

```
AlbumId MediaTypeId AvgBytes
##
## 1
          253
                        3 536359244
## 2
          229
                        3 535292434
## 3
          227
                        3 529469291
          231
                        3 514373372
## 4
## 5
          228
                        3 512231374
## 6
          254
                       3 492670102
## 7
          226
                       3 490750393
## 8
          261
                        3 453454450
## 9
          251
                        3 306109250
          249
                        3 268393262
## 10
```

# Disconnecting from the database

After the end of a session, it is good practice to explicitly close your connection.

Does this remove the database connection chinook\_db in the R session?

#### We saw

- Integrity constraints (Primary and Foreign Keys)
- All parts of a SQL query

SELECT columns or computations
FROM table
WHERE condition
GROUP BY columns
HAVING condition
ORDER BY column [ASC | DESC]
LIMIT offset, count;

- Database tools for R
  - the R packages RSQLite, DBI
  - the database Chinook\_Sqlite.sqlite