

ChE352
Numerical Techniques for Chemical Engineers
Professor Stevenson

Lecture 14

Linear algebra is easy in numpy

```
b = np.array([1, 1, 1]) # Define vector b
A = np.array([[4, -1, 1], [-1, 4.25, 2.75],
[1, 2.75, 3.5]]) # Define matrix A
x = np.linalg.solve(A, b) # Find  $Ax = b$ 
C = A + B # Find A plus B
D = A @ B # Find A times B (matrix multiply)
E = A * B # Find A times B COMPONENT-WISE
Bsq = B**2 # Square each element of B ( $\neq B@B$ )
L = np.linalg.cholesky(A) # Cholesky factor
Ainv = np.linalg.inv(A) # Inverse (unwise)
Apnv = np.linalg.pinv(A) # pseudo-inverse
```

Linear algebra is easy(?) in numpy

by default, a 1-D np.array is "flat"

a = np.array([1, 2, 3]) # not row OR col

a2 = np.array([[1, 2, 3]]) # row vector

b = a.reshape(-1, 1) # column vector

c = a.reshape(1, -1) # row vector

d = b.flatten() # flat version again

shape affects matrix operations:

one_number = c @ b # [[14]]

nine_numbers = b @ c # 3x3 matrix

Writing out a large matrix in numpy

```
A = np.array([
    [9, 1.5, 0, 2.5, 0.5],
    [1.5, 10, 1.5, 0.5, 2 ],
    [0, 1.5, 11, 0, 2 ],
    [2.5, 0.5, 0, 8, 1 ],
    [0.5, 2, 2, 1, 5 ]],
)
```

```
# Aligning the columns makes it
# easier to spot errors
```

How do you confirm your data?

```
import hashlib  
# Get the unique hash of the data  
sha = hashlib.sha256()  
sha.update(A.dumps()) # dump string  
print(sha.hexdigest()[8])  
# then you can compare your hash to  
# someone else's to see immediately  
# if you have the same data
```

What does the [8] do? Why does it help?

How do you confirm an answer?

```
# example system of linear equations
```

```
b = np.array([1.4, 1.7, 2.0])
```

```
A = np.array([[ 4, -1,  1 ],  
              [-1, 4.25, 2.75],  
              [ 1, 2.75, 3.5]])
```

```
# first we'll solve, then test the solution
```

```
x = np.linalg.solve(A, b) #  $Ax = b$ : find  $x$ 
```

```
# check whether  $x$  is close (default=1e-8)
```

```
Ax = A @ x # matrix multiply  $A$  times  $x$ 
```

```
print(np.allclose(Ax, b)) # test  $Ax = b$ 
```

Matrices are also linear operators

- Any linear operation over vectors of length N can be represented as an $N \times N$ matrix
 - *Linear operator* means: maps vector x to vector y such that all entries of y are weighted sums of entries of x
 - Example: *rotate a vector*
- This means we can treat an **operator** / transformation (a function from one vector to another) as **data** ($N \times N$ grid of numbers)

Rotation matrices

- Any rotation or scaling of a vector can be represented with a special matrix
- A rotation matrix is one that changes the *direction* of a vector but not its *magnitude*

Simple 2D
rotation matrix
examples

$$\begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}, \quad \begin{bmatrix} -1 & 0 \\ 0 & -1 \end{bmatrix}, \quad \begin{bmatrix} 0 & 1 \\ -1 & 0 \end{bmatrix}$$

for 90°, 180°, and 270° counter-clockwise rotations.

- A 3D rotation can be represented with Euler angles (\mathbf{R}^3) or quaternions (\mathbf{R}^4), but a rotation matrix is easiest: just multiply

Matrix vocabulary

You should know what all of these things are:

- Scalar (dot) product, Matrix product
- Square, Diagonal, Identity matrices
- Upper and Lower Triangular matrices
- Transpose, Symmetric

$$A^{-1}, A^T, |A|, M_{ij}$$

You should know what all of these things are:

- Matrix inverse, existence of inverse (When?)
- Singular/noninvertible v. nonsingular/invertible
- Determinant
- Eigenvector / eigenvalue

Nonsingular Matrices are Nice

The following statements are equivalent for A in $\mathbb{R}^{n \times n}$:

1. $\text{rank}(A) = n$
2. A is nonsingular
3. A^{-1} exists
4. The rows and columns of A are lin. indep.
5. $\det(A) \neq 0$
6. $\text{range}(A) = \mathbb{R}^n$
7. $\text{nullspace}(A) = \{0\}$
8. $Ax = b$ has a unique solution x^* for each b in \mathbb{R}^n
9. The only solution to $Ax = 0$ is $x^* = 0$
10. Zero is not an eigenvalue of A

iff, \Leftrightarrow

What if it isn't square?

How could a matrix be **close** to singular?

Factoring Matrices

- Solving a linear system $Ax = b$ requires $O(n^3)$ operations with Gaussian Elim. for A in $R^{n \times n}$
- If n is large (>1000) or if we need $x = A^{-1}b$ for many different choices of b , this is expensive
- If we can find triangular matrices L & U such that $A = LU$, then we can find x differently:

$$Ax = b \rightarrow LUx = b, \text{ let } y = Ux \rightarrow Ly = b$$

- If U and L are triangular, solving $Ly = b$ for y and $Ux = y$ for x takes only $O(n^2)$ operations
- If $n=1000$, how much faster is this than $O(n^3)$?

LU solution for $Ax = b$

1. Let $A = LU \rightarrow L Ux = b$
2. Let y be a new n -vector: $y = Ux \rightarrow Ly = b$
3. Since L is lower triangular, the first equation in $Ly = b$ says that $y_1 = b_1 / L_{11} \rightarrow 1$ flop for y_1
4. y_1 is now known and the second equation involves only y_1 and $y_2 \rightarrow$ Calculate y_2
5. Repeat until y_n **This takes $2n^2$ flops total**
6. y is now known; repeat the same process for $Ux = y$, starting now with x_n and going up to x_1 since U is upper triangular.

Activity: $Ly = b$

5 min to do, 5 min discuss

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 4 & 0 \\ 5 & 2 & 3 \end{bmatrix}, \quad b = \begin{bmatrix} 0 \\ 4 \\ 6 \end{bmatrix}$$

Solve for y if $Ly = b$. It should only take $O(n^2) \approx 9$ flops.

Is this faster than Gauss. Elimination?

Answer: $Ly = b$

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 4 & 0 \\ 5 & 2 & 3 \end{bmatrix}, \quad b = \begin{bmatrix} 0 \\ 4 \\ 6 \end{bmatrix}$$

$$y_1 = 0$$

$$y_2 = 1$$

$$y_3 = 4/3$$

NOTE: getting LU form takes $O(n^3)$ flops

Special Matrices

- Sparse matrices – Have many more zeros than non-zeros, can save computation (What would this mean in a physical system?)
- Symmetric matrix: $A = A^T$ (What shape is A ?)
- Positive definite matrices: $x^T P x > 0$ for any non-zero vector x (square, symmetric)
- Negative definite matrices: $x^T N x < 0$ for any non-zero vector x (square, symmetric)
- Positive semidefinite: $x^T S x \geq 0$ for any x
- Negative semidefinite: $x^T D x \leq 0$ for any x

Properties of PD Matrices

If a matrix P in $\mathbb{R}^{n \times n}$ is PD, it implies the following:

1. P is non-singular
2. All diagonal entries of P are positive
3. $-P$ is negative definite
4. Solving $Px = b$ has stable growth of error
5. All leading principle minors (1×1 , 2×2 , \dots , $n \times n$) must be positive (Sylvester's criterion)
6. $P = U^T U = LL^T$ for some upper triangular U with positive diagonal entries. We call this the Cholesky Decomposition: $L = U^T$, $U = L^T$

Vector Norms = *Kinds of Length*

$\forall x \in \mathbb{R}^n, \quad \exists \|\cdot\|_p : \mathbb{R}^n \rightarrow \mathbb{R} \quad \text{such that :}$

$$\|x\|_p = \begin{cases} \sqrt[p]{|x_1|^p + |x_2|^p + \dots + |x_{n-1}|^p + |x_n|^p} & 1 \leq p < \infty \\ \max(|x_1|, |x_2|, \dots, |x_{n-1}|, |x_n|) & p = \infty \end{cases}$$

e.g. $\|x\|_2 = \sqrt{x^T x}$ (always > 0 if $x \neq 0$)

Every p corresponds to a kind of vector length

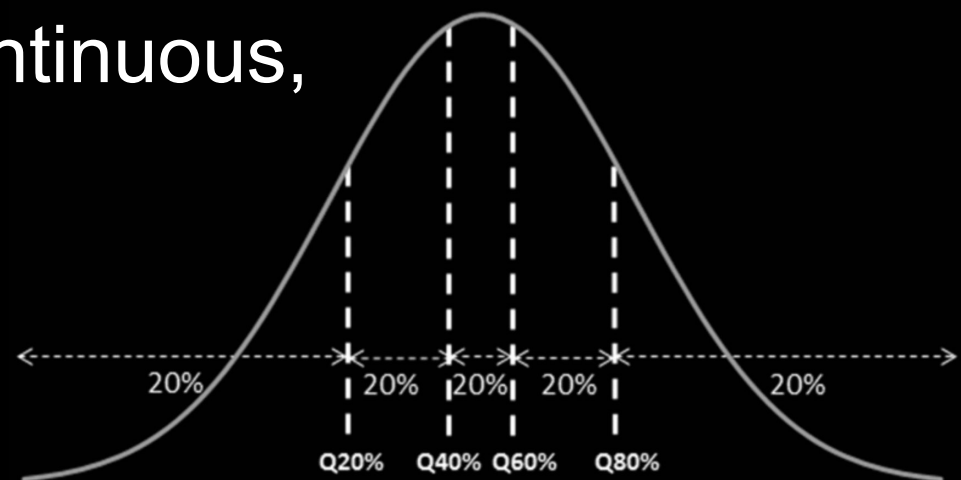
Inner product: (how is it like the 2-norm?)

$$\langle x, y \rangle \equiv x^T y = y^T x \quad \forall x, y \in \mathbb{R}^n$$

Example: Probability Vectors

- For any event, the set of probabilities of all outcomes form a "probability vector" with a 1-norm (Manhattan norm) of exactly 1
- Some matrices can act on probability vectors and give new probability vectors
- For non-finite outcomes, the 1-norm sum \sum is continuous, aka an integral \int

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} p \\ 1-p \end{pmatrix} = \begin{pmatrix} 1-p \\ p \end{pmatrix}$$



Vector Spaces

The following are true for any norm and any elements x and y of a normed vector space:

1. $\|x\| \geq 0$

2. $\|x\| = 0 \iff x = 0$

3. $\|\alpha x\| = |\alpha| \|x\| \quad \forall \alpha \in \mathbb{R}$

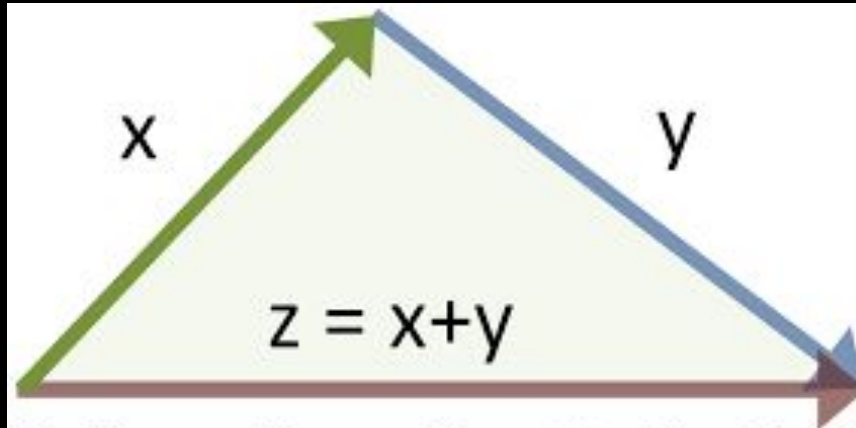
4. $\|x + y\| \leq \|x\| + \|y\|$ (triangle inequality)

5. $|\langle x, y \rangle| \leq \|x\| \|y\|$ (Cauchy-Schwarz inequality)

Norms are a
measure of
distance between
two points/vectors

What vector spaces do we use in this class?

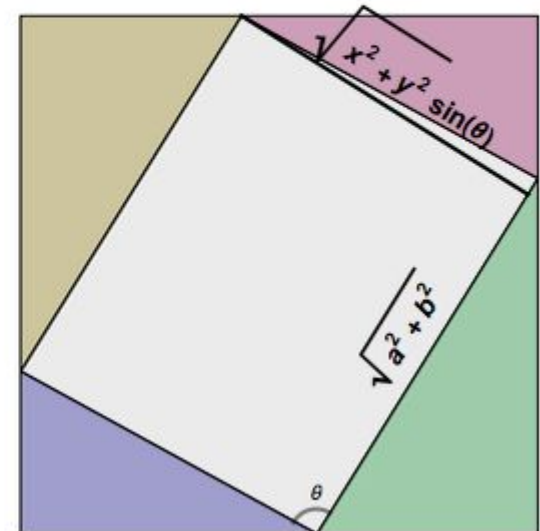
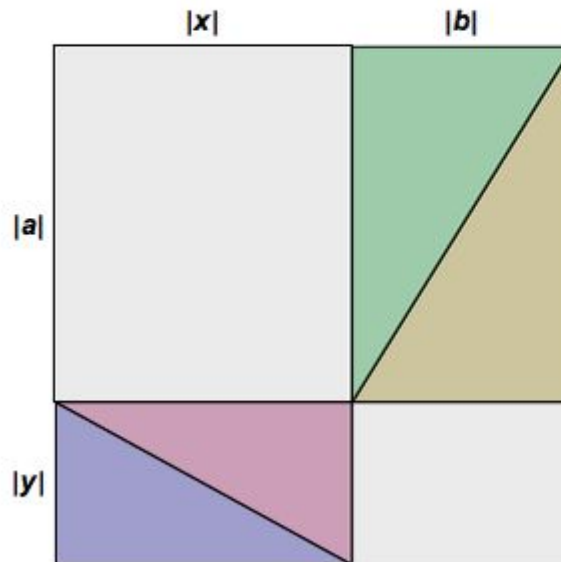
Vector norm inequalities



Triangle inequality

$$|x + y| \leq |x| + |y|$$

Cauchy-Schwarz inequality
 $|x \cdot y| \leq |x| \cdot |y|$



Activity: Vector Norms

5 min to do, 5 min discuss

Find the 1-norm, 2-norm (Euclidean norm), and the infinity-norm of the following vectors:

$$x = \begin{bmatrix} 1 \\ 4 \\ 9 \\ 1 \\ 15 \end{bmatrix}, \quad y = \begin{bmatrix} 2 \\ 0 \\ -7 \\ 0 \\ 0 \end{bmatrix}$$

Then show that the triangle & Cauchy-Schwarz inequalities hold for the 2-norm.

Cauchy-Schwarz

$$|x \cdot y| \leq |x| \cdot |y|$$

Triangle

$$|x + y| \leq |x| + |y|$$

Answer: Vector Norms

$$\|x\|_1 = |1| + |4| + |9| + |1| + |15| = 30,$$

$$\|y\|_1 = |2| + |0| + |-7| + |0| + |0| = 9$$

$$\|x\|_2 = \sqrt{|1|^2 + |4|^2 + |9|^2 + |1|^2 + |15|^2} = 18,$$

$$\|y\|_2 = \sqrt{|2|^2 + |0|^2 + |-7|^2 + |0|^2 + |0|^2} = \sqrt{53}$$

$$\|x\|_\infty = \max_{i=1\dots n} |x_i| = 15, \quad \|y\|_\infty = \max_{i=1\dots n} |y_i| = 7$$

$$\|x + y\|_2 = \left\| \begin{bmatrix} 3 & 4 & 2 & 1 & 15 \end{bmatrix}^T \right\|_2$$

$$= \sqrt{255} \approx 16 \leq 18 + \sqrt{53} = \|x\|_2 + \|y\|_2 \approx 25$$

$$|\langle x, y \rangle| = |2 - 63| = 61 \leq 18\sqrt{53} = \|x\|_2 \|y\|_2 \approx 131$$

Which norm is
"the" norm?

How is the
 ∞ -norm related
to infinity?

Eigenvalues and Eigenvectors

Heard of these before?

What is the characteristic polynomial?

What can we do with eigen stuff?

$$Ax = \lambda x$$

$$p(\lambda) = \det(A - \lambda I)$$

“Eigen” means
“characteristic”,
or “own” as used
in the phrase
“my own room”

Eigenvalues and Eigenvectors

Eigenvalues & eigenvectors make the most sense when you think of matrices as *operators*, aka linear transformations, not just grids of numbers (though both ideas are true)

Eigenvectors

Eigenvalues

$$A\vec{v} = \lambda\vec{v}$$

How do we find λ_i numerically?

$$A = \begin{bmatrix} 5 & -1 & 3 \\ 2 & 8 & 0 \\ 3 & -1 & 11 \end{bmatrix}$$

We can find the eigenvalues by solving for λ in $\det(A - \lambda I) = 0$ (the "characteristic polynomial"). This is roughly $O(N^3)$, like matrix multiplication.

We can find the eigenvectors by substituting each eigenvalue λ_i into the definition $A\mathbf{v}_i = \lambda_i\mathbf{v}_i$

Finding eigenvalues

"Characteristic
polynomial"

$$p(\lambda) = \det(A - \lambda I) = \det \begin{bmatrix} 5 - \lambda & -1 & 3 \\ 2 & 8 - \lambda & 0 \\ 3 & -1 & 11 - \lambda \end{bmatrix}$$

– λI means
subtract λ from
the diagonal
(Why?)

$$= (5 - \lambda)(8 - \lambda)(11 - \lambda) - 0 - 6 - 0 + 2(11 - \lambda) - 3^2(8 - \lambda)$$

$$= (40 - 13\lambda + \lambda^2)(11 - \lambda) + 7\lambda - 56$$

$$= -\lambda^3 + 24\lambda^2 - 176\lambda + 384 = 0 = (\lambda - 4)(\lambda - 8)(\lambda - 12) \rightarrow$$

$\lambda_1 = 4$
$\lambda_2 = 8$
$\lambda_3 = 12$

Polynomial is cubic in λ ,
because A is a rank-3 matrix

Cubic polynomial, so 3 roots

Eigenvector v^1

Finding eigenvectors

$$\lambda_1 = 4 \rightarrow Av^1 = \lambda_1 v^1 = 4v^1 \rightarrow Av^1 - 4v^1 = 0 \rightarrow$$

$$(A - 4I)v^1 = 0 \quad (Bx = 0 \text{ means } \text{rank}(B) < n) \rightarrow$$

$$\begin{bmatrix} 1 & -1 & 3 \\ 2 & 4 & 0 \\ 3 & -1 & 7 \end{bmatrix} \begin{bmatrix} v_1^1 \\ v_2^1 \\ v_3^1 \end{bmatrix} \rightarrow \begin{array}{l} v_1^1 - v_2^1 + 3v_3^1 = 0 \\ 2v_1^1 + 4v_2^1 = 0 \\ 3v_1^1 - v_2^1 + 7v_3^1 = 0 \end{array} \rightarrow \begin{array}{l} 3v_3^1 = v_2^1 - v_1^1 \\ v_1^1 = -2v_2^1 \\ 3v_1^1 + 7v_3^1 = v_2^1 \end{array}$$

$$\rightarrow v^1 = \begin{bmatrix} 2 \\ -1 \\ -1 \end{bmatrix} \quad OR \quad v^1 = \begin{bmatrix} \sqrt{6}/3 \\ -\sqrt{6}/6 \\ -\sqrt{6}/6 \end{bmatrix}$$

A system of linear equations $Ax = b$, where b is all zeros

Vector quantum mechanics

- Probability vectors have 1-norm = 1.0
- What if we used **2-norm** = 1.0 instead?
- Generalized probability vectors = quantum state vectors
- Matrices which keep the 2-norm are called *Hermitian* matrices (in QM, *operators*)

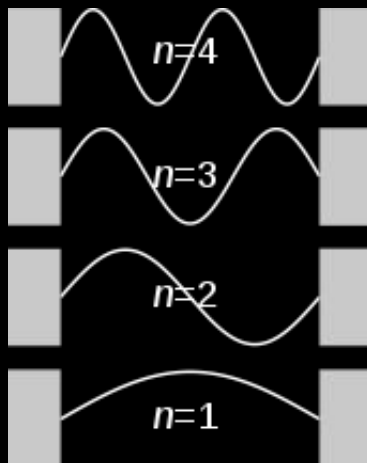
$$A \text{ Hermitian} \iff A = \overline{A^T}$$

A matrix which maintains the 2-norm is its own conjugate transpose

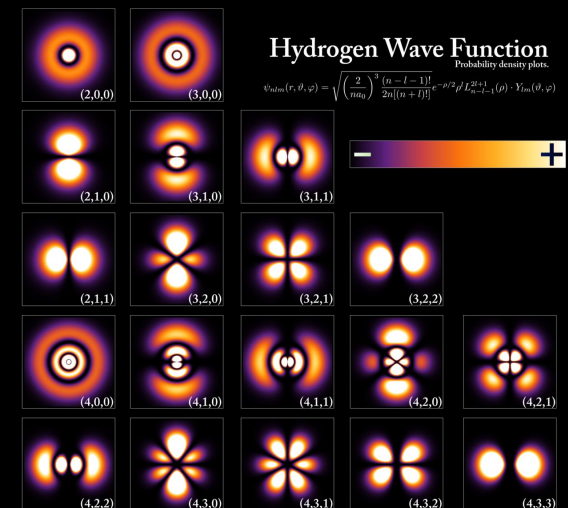
- Energy is just an eigenvalue: $H\psi = E\psi$
- ~100% of quantum chemistry is with vectors

How do we build a quantum state?

- For discrete properties, like spin $\uparrow\downarrow$, it's easy: vector contains all discrete possibilities
 - For one particle, spin state vector = $[a, b]$ where a & b are amplitudes of $|\uparrow\rangle$ & $|\downarrow\rangle$
- For continuous properties, like position, we can pick a set of functions that approximate it



The state vector consists of an amplitude for each basis function



Eigenvalues in numpy

```
from numpy.linalg import (det, diag, eig,  
                           matrix_rank, norm)  
  
d = det(A)    # Determinant of A  
  
r = matrix_rank(A)    # Rank of A  
  
# if A is Hermitian, can use faster eigh(A)  
  
eigenvalues, eigenvectors = eig(A)  
  
# norm works for p = 1, 2, np.inf  
  
x = norm(A, p)    # Matrix norm
```

Next time: Optimization

“World domination is such an ugly phrase.
I prefer to call it *world optimization*.”

– Eliezer Yudkowsky, author

(PNM is weak here, so all F&B this time)

Iterative search, steepest descent,
nonlinear solvers, Newton, quasi-Newton:
F&B 7.6 and 10.1-4