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Low Power Hardware Accelerated Internet of Things Cryptography

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Abstract

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Introduction

Over the last few years there has been a shift in type of devices connected to the internet from just servers, PC's and later smartphones, to small embedded processors that can control many devices. The idea of connecting such devices to internet has been dubbed 'Internet of Things' or 'IoT' and has the aim to make our lives simpler. These IoT devices can range from industrial applications, such as automated factories, to public infrastructure, like a smart electrical grid, or more homely appliances in a smart connected home. Thus due to the wide range of products that a connected IoT device can be applied to improve the efficiency and/or usefulness, it has been predicted that billions of devices will be in use by [insert year/reference]. This also means that the complexity of the devices varies greatly with simple light switches and even kettles being given the IoT treatment[insert reference], but at the other end of the scale a network of connected self-driving cars is being considered [insert reference].

The one thing that all of these devices have in common though is that they need to be secure as they communicate sensitive and private data through an open channel on the internet between the user and the device. Hence, to keep the potential adversaries from accessing the data and possibly controlling numerous connected devices, maliciously or not, an encryption algorithm can be used with only the specified users having access to the decryption key. There are many encryption algorithms that perform this function and most can be implemented in both software and dedicated hardware such as a Application Specific Integrated Circuit (ASIC) or a Field Programmable Gate Array (FPGA). As a majority of IoT devices are implemented on small embedded processors which have limited resources, the hardware option might possibly be a better solution for IoT devices. However, due to the fact that most IoT devices are always on, power consumption is a very important factor when considering options for adding hardware accelerated encryption and for battery powered devices it is often more critical than the actual encryption.

The goals of this project are to explore various encryption algorithms and compare their performance based on data throughput, accuracy, security and power consumption when implemented in software and hardware. To evaluate these parameters the same algorithms can be coded in C or C++ for the software versions and a Hardware Development Language (HDL) such as System Verilog can be first simulated in ModelSim insert reference, before programming a FPGA for the hardware version.

These comparisons can then be used to match the algorithms to the appropriate IoT device as they all have different requirements for relative security level and power consumption, as for example a light switch does not necessarily need to be protected from the same level of attack as a set of digital locks or private data storage. In order for the hardware to work with IoT devices it will also need a communication protocol like I^2C or SPI to work with embedded processors, and possibly Ethernet or WiFi to act as the gateway to internet for the device. Some of these protocols are available on FPGA development boards but can be implemented in System Verilog code.

Background Research and Literature

2.1 Internet of Things

As mentioned in chapter 1 there are many IoT devices that require varying levels of security and have to be protected against different of attacks, like side channel attacks. The purpose of this section is to discover what these devices are and what properties are required in the encryption algorithm in terms of throughput, power or energy consumption, software or hardware restrictions and security.

[Insert Iot Products/services and Power/Resource constraints]

2.2 Cryptography

After evaluating the conditions required to provide the appropriate level of security in section 2.1 this section explores the cryptographic principles and algorithms available that satisfy those conditions. The primary objective of cryptography is to convert, or encrypt, a readable message known as plaintext into an unreadable form, ciphertext, so that adversaries cannot read the contents, but over the years the scope of cryptography has widened [insert reference]. Cryptography is therefore the study of encryption and these techniques, its counterpart: the study of breaking the encryption to find the original message, is known as cryptanalysis. Eventually, for most encryption techniques a weakness is found, and subsequently exploited, so more complex techniques are conceived and with the invention of computers the complexity of the algorithms has increased greatly. However, the computer power is also available for cryptanalysis so the cycle of continuous improvement of the algorithms hasn't stopped.

In cryptography there are two main concepts that the algorithms are based on: symmetric and asymmetric keys[insert reference]. In asymmetric key cryptography a unique key is used to encrypt data and different, but a related key, is used to decrypt it. The relationship between the two keys is often defined by maths problem that is very difficult to solve which is the basis of the encryption[insert reference]. On the

other hand, symmetric key cryptography uses the same key for both encryption and decryption, hence symmetric, with the security usually provided by a combination of simple logic operations [insert reference].

2.2.1 Asymmetric Key

Asymmetric key cryptography can be referred to as public key due to the fact that one of the related keys can be publicly available without compromising the security of the encrypted data. This is because the keys are usually generated based on mathematical problems that have no solution or the solution is impossible for a computer to solve efficiently, such that solving it takes longer than an exhaustive key search[insert reference]. There are many problems that fit this criteria but the most popular in use today are the integer factorization and elliptic curve problems used by the RSA[insert reference] and the family of Elliptic-curve cryptography (ECC) techniques[insert reference] respectively. Public key cryptography can be used in two different modes as if data is encrypted with the intended recipients public key only they can decrypt it with their private key, thus encryption. However, if a private key is used for encryption then using the public key to decrypt it ensures the senders identity, authentication[insert reference].

2.2.2 Symmetric Key

Similar to the symmetric/private comparison symmetric key cryptography is also known as private key as in order to keep the encrypted data secure the key used must be kept secret. There are two main types of private key algorithms that operate on the plaintext differently: block ciphers which uses a fixed number of bits, block; or stream ciphers which encrypts data bit by bit[insert reference].

Modern block ciphers are based on Claude Shannons work on product ciphers in [insert reference], in which he suggested that iterating a cipher for multiple rounds, with subkeys, improves the security. Hence, the cipher to be iterated didn't need to be complex operations and simple logic operations such as XOR, substitution or permutation of the plaintext could be used [insert reference]. The base cipher that is iterated is known as the round function and it takes as an input a block of plaintext and a subkey, which is generated from the main key by a separate key expansion function, and outputs a block of ciphertext. The output is usually the result of the round function XORed with the subkey. The round functions are mostly designed using either a Feistel network (F network) or a Substitution Permutation network (SP network) insert reference.

The Fesitel network was named after physicist Horst Feistel who was a integral part of the team at IBM that developed the early block cipher Lucifer. It works by splitting the input plaintext into equal words and applying the round function to it right, or LSB, word before swapping the words for the next iteration. This functionality, for both encryption and decryption, is described in [Insert Figure]. A more complete description is available in [Insert Appendix].

On the other hand, Substitution Permutation networks operate on the whole plaintext block using S-boxes for substitution and P-boxes for permutation[insert reference]. There can be more steps that treat the block slightly differently but those are the basic steps and when combined they are enough to provide Shannon's confusion and diffusion[insert reference]. The basic structure of an SP network is in [Insert Figure] but a more detailed description is available in [Insert Appendix].

Stream ciphers were initially designed to approximate the One Time Pad (OTP) cipher that was proved to be completely unbreakable by Claude Shannon[insert reference]. The OTP works by combining each digit of the plaintext with a completely random keystream. Stream ciphers work by generating a pseudo-random keystream to combine with the plaintext[insert reference]. Because the keystream is pseudo-random and not completely random a stream cipher is breakable[insert reference]. The keystream is created by a pseudo-random number generated with a crypto-graphic key used as a seed.

There are modes of operation for block ciphers, in [insert reference], that provide better security by using feedback of the ciphertext to the next block. Some of these modes of operation also allow block ciphers to behave similar to stream ciphers as they encrypt an initialization vector with the key and the resulting ciphertext can be combined with the plaintext.

2.2.3 Decisions

Due to the fact that asymmetric key algorithms are hard to solve they require complex hardware or software to implement which is undesirable for this project. Also, with the exception of ECC the key sizes needed for the security can be very large so with the limited IO pins available on FPGAs they could prove difficult to program. On the other hand, many private key algorithms are designed to be efficient in hardware especially Feistel networks as an inverted round function isn't required. While a stream cipher can be useful to encrypt serial data that will most likely be the source, although the modes of operation available for block ciphers provide more flexible functionality including stream cipher modes. Therefore, the algorithm chosen for this project will most likely be a block cipher with a Feistel network.

2.3 Conventional Algorithms

There are many block ciphers that are considered very secure and therefore popular, they include: DES, AES, Blowfish [insert reference]. DES operates on a block of 64 bits for 16 rounds using a key length of 64 bits but it has an effective key length of 56 bits as 8 bits were used for parity [insert reference]. AES, an upgrade to DES, is far more secure as uses a 128 bit block and has the flexibility of using three different key lengths: 128, 192 and 256. The number of rounds that AES iterates depends on the key length with 10 rounds used for a 128 bit key, 12 for 192, and 14 for the largest key [insert reference]. Blowfish, like DES, operates on a 64 bit block and iterates for 16 rounds, but it can use a variable key length in the range 32 to 448 bits [insert reference].

2.3.1 Standardization

DES, which stands for Data Encryption Standard, is one the earliest block ciphers used in the computer age and it was developed by IBM in the 1970s based on their earlier cipher Lucifer [insert reference]. As with Lucifer it was designed around a Feistel network but the round function used also has a SP network structure to it [insert reference] It has the name Data Encryption Standard as it was accepted as the standard encryption algorithm by the US National Bureau of Standards (NBS), now the National Institute of Standards and Technology (NIST), in 1977 after it was altered by the National Security Agency (NSA), which caused some controversy [insert reference].

DES was used for about two decades but in the 1990s several successful attacks proved its weakness [insert reference] so in 1997 NIST started a selection process to find its replacement. Many algorithms were submitted as candidates for the standard but the finalists were: MARS, RC6, Rijndael, Serpent, and Twofish [insert reference]. It took three years to decide on the algorithm to be set as the standard which was announced as Rijndael in 2000 and the standard was et in 2001, with the 128, 192, 256 bit keys being used in the standard [insert reference]. Unlike DES, AES uses a SP network as it is efficient, in time, in both hardware and software [insert reference]. The round function treats the block as a 4x4 byte matrix and performs multiple steps on the data: sub bytes; shift rows; mix columns and add round key [insert reference].

[Insert Attacks and FPGA implementations of AES]

Since its standardization in 2001 AES has been used

2.3.2 Other Algorithms

The US standardized algorithms quickly became very popular and can be considered the unofficial global standard insert reference. Although, there are many other algorithms that are considered secure and are commonly used. These algorithms might be used because there is still some distrust of the NSAs involvement in the algorithms and they are more open source. This is the case of the Blowfish algorithm as it is unpatented and can therefore be used in any product without legal consequence. Blowfish was designed by Bruce Schneier in the early 1990s as he, and many others, noticed the insecurity of DES particularly with the 56 bit key length making a brute force attack more plausible insert reference.

[Insert Attacks and FPGA implementations of Blowfish]

2.3.3 Decisions

2.4 Lightweight Algorithms

After deciding that the conventional algorithms might not be suitable for the low power devices targeted by this project, some more lightweight algorithms were found including: PRESENT, PRINCE and the SIMON and SPECK algorithms insert reference. However, as IoT is an emerging technology and is the main reason for lightweight cryptography there isn't a standard set by NIST or any other organisation insert reference.

2.4.1 SIMON & SPECK

The SIMON and SPECK family of algorithms are the lightweight techniques proposed by the NSA that were designed to perform well in both software and hardware while still being secure; and to be flexible in terms of block and key size insert reference, listed in Table 2.1. The different algorithms are similar but SIMON was optimised for hardware implementations and SPECK for software. As with AES the number of rounds iterated depends on the key size but as the block size varies as well it also has an effect as shown in Table 2.1. The structure of both algorithms is a Feistel network and thus it works on words of n bits, where 2n is the block size, and with a key of mn bits.

| Block Size | Key Size | n | m | SIMON Rounds | SPECK Rounds |
|------------|----------|----|---|--------------|--------------|
| 32 | 64 | 16 | 4 | 32 | 22 |
| 48 | 72 | 24 | 3 | 36 | 22 |
| 48 | 96 | 24 | 4 | 36 | 23 |
| 64 | 96 | 32 | 3 | 42 | 26 |
| 64 | 128 | 32 | 4 | 44 | 27 |
| 96 | 96 | 48 | 2 | 52 | 28 |
| 96 | 144 | 48 | 3 | 54 | 29 |
| 128 | 128 | 64 | 2 | 68 | 32 |
| 128 | 192 | 64 | 3 | 69 | 33 |
| 128 | 256 | 64 | 4 | 72 | 34 |

Table 2.1: A table of the modes of operation for the SIMON & SPECK Algorithms. Adapted from [insert reference].

As SIMON was designed primarily for hardware it only makes use of XOR (\oplus) , AND (&) and circular rotate operations $(R^j[x])$ on the n bit wide words. The encryption and decryption functions take the Feistel network form described in [Insert Figure] with the round function Equation 2.1.

$$F(x) = (R^{1}[x] \& R^{8}[x]) \oplus R^{2}[x]$$
(2.1)

SPECK on the other hand, being optimised for software implementations uses XOR (\oplus) , modulo 2^n addition (+) and circular rotate operations $(R^j[x])$ The encryption and decryption functions take a slightly different form to the basic Feistel network, [Insert Figure], and are shown in Equation 2.2 and 2.3 where $\alpha = 7$ and $\beta = 2$ if n = 16, but $\alpha = 8$ and $\beta = 3$ otherwise.

$$L_{i+1} = (R^{-\alpha}[L_i] + R_i) \oplus K_i$$
 (2.2)

$$R_{i+1} = R^{\beta}[R_i] \oplus (R^{-\alpha}[L_i] + R_i) \oplus K_i = R^{\beta}[R_i] \oplus L_{i+1}$$
 (2.3)

[Insert Attacks and FPGA implementations of SIMON & SPECK]

${\bf 2.4.2}\quad {\bf Other\ Algorithms}$

2.4.3 Decisions



Progress

Based on the research explored in chapter 2 I chose the SIMON algorithm from the NSA, mainly because of its flexibility in security levels, with different key lengths, that could be applied to the different devices explored in section 2.1. This means that multiple algorithms don't need to be developed and I could concentrate on making my code as efficient as possible. As the aim of this project is to compare how an algorithm performs in hardware and software and due to the fact that there isn't a standard library or program for SIMON, a software version was required for benchmarking. Starting in software also increased by familiarity of the algorithm and provide a good starting point for System Verilog development.

For all versions of the algorithm (table 2.1) efficiency, in terms of power consumption and resource use, is very important but time efficiency is not an initial priority. All versions of this algorithm were developed not only with the description but also the pseudocode provided in [insert reference]. That document also has a set of test vectors for all modes that define the ciphertext that the algorithm should produce from the given key and plaintext [insert reference]. Due to the Feistel network structure of this algorithm encryption can be done in parallel to the key expansion but for decryption requires the keys to be pre expanded. For this reason, and because some modes of operation only require encryption, two variants of the algorithm were developed: one for just encryption and one full functionality.

3.1 Hosted C

As there are ten versions of this algorithm that all differ slightly with the block key sizes they could all be developed individually. However, as ten versions of the similar functions would be required I felt it was better to have one version that is flexible. However, checking which mode the program is before most operations would be very inefficient and as the program would rarely changing mode during runtime the decisions could be made at compile time using preprocessor macros as in [insert listing][insert reference]. The macros were used to define the variable type used for the words based on the values for n with the uintn-t type. Also the value of m was used with macro if statements in the key expansion functions where extra code is required if m = 4.

3.2 System Verilog

After the Hosted C version was working with the test vectors the System Verilog development began. Similar to the function used in the software developed in section 3.1 modules were created and tested with individual testbenches before being combined in the top level control modules. Most of the modules, including the rotate operations, were done in combination logic blocks for initial simplicity to ensure correct simulations. While this might work in simulations, it could be unreliable and inefficient when implemented in an FPGA. Also parameters were used so the different modes could use the same code but the only testing done so far is with the SIMON32/64 version.

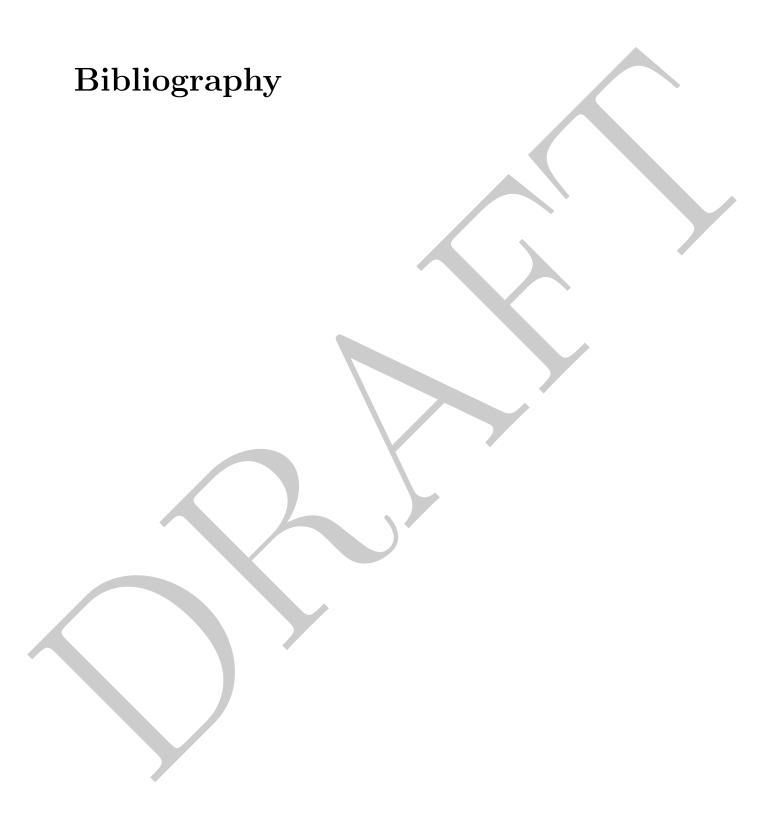
[Insert Listings, test vectors, simulations and synthesis]



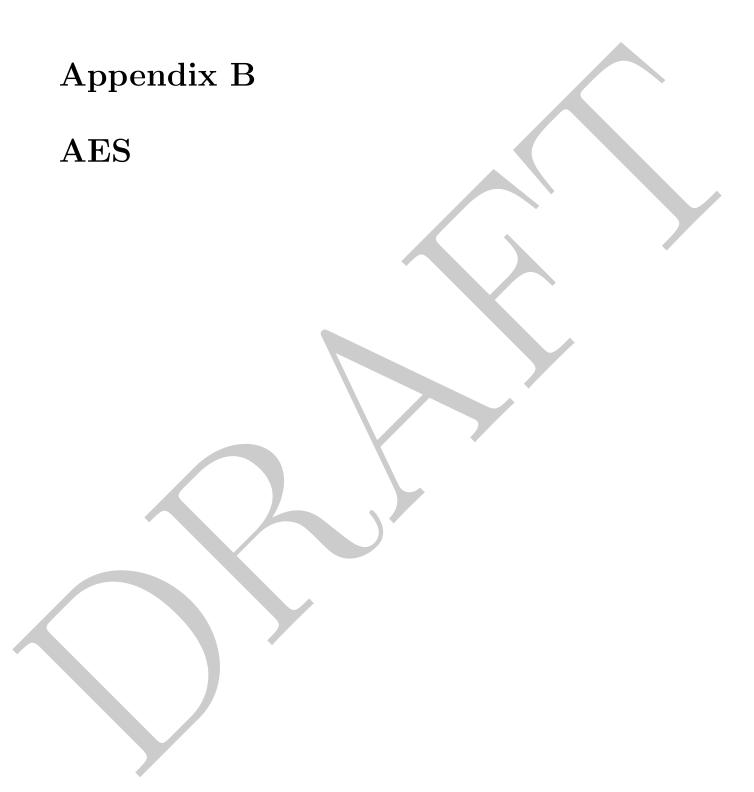
Future Plan

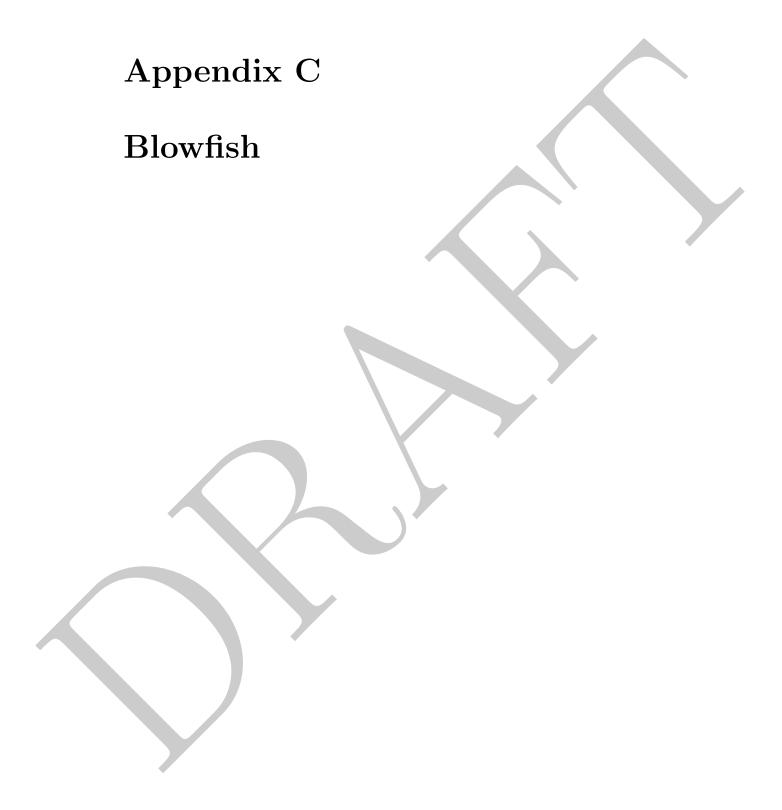
To continue this project the System Verilog code needs to be improved with more sequential operations to improve reliability and efficiency. Serializing the algorithm at the byte level could also be explored which could have some interesting results similar to the bit serial version presented in [insert reference]. Work will also have to be done to implement the System Verilog top module onto an FPGA and interface it with communication module, that could be provided by the FPGA development board used or from an OpenCores.org project[insert reference]. When the FPGA is functioning correctly tests will need to be developed to determine the throughput in hardware and software, preferably operating at similar clock frequencies to ensure the results will be as comparable as possible. The hosted C program, section 3.1, should also be adapted to be programmed onto an 8 or 16 bit microcontroller and then the same parameters will be tested and compared.

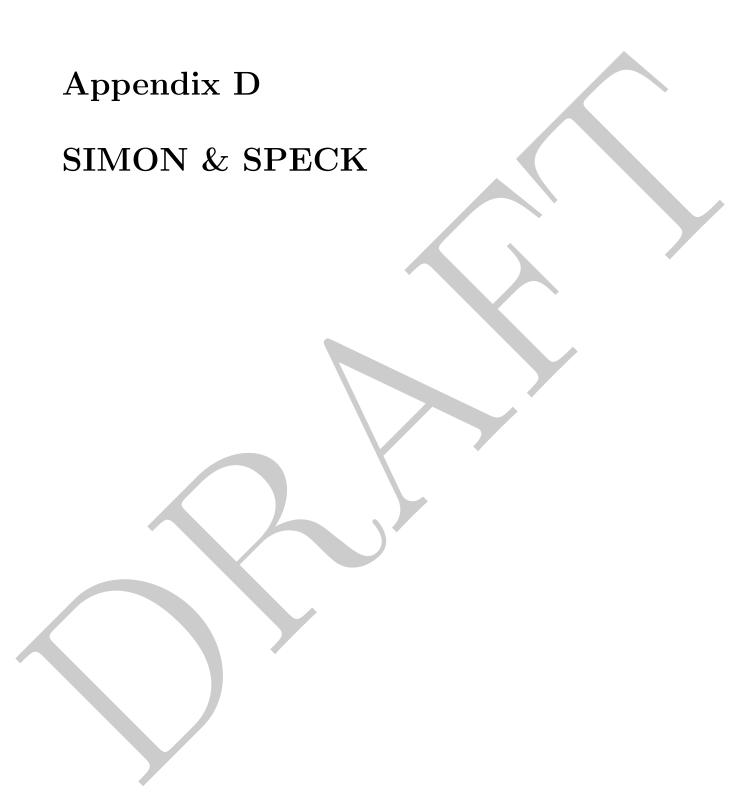
[Insert Contingencies, Processor/FPGA connected]

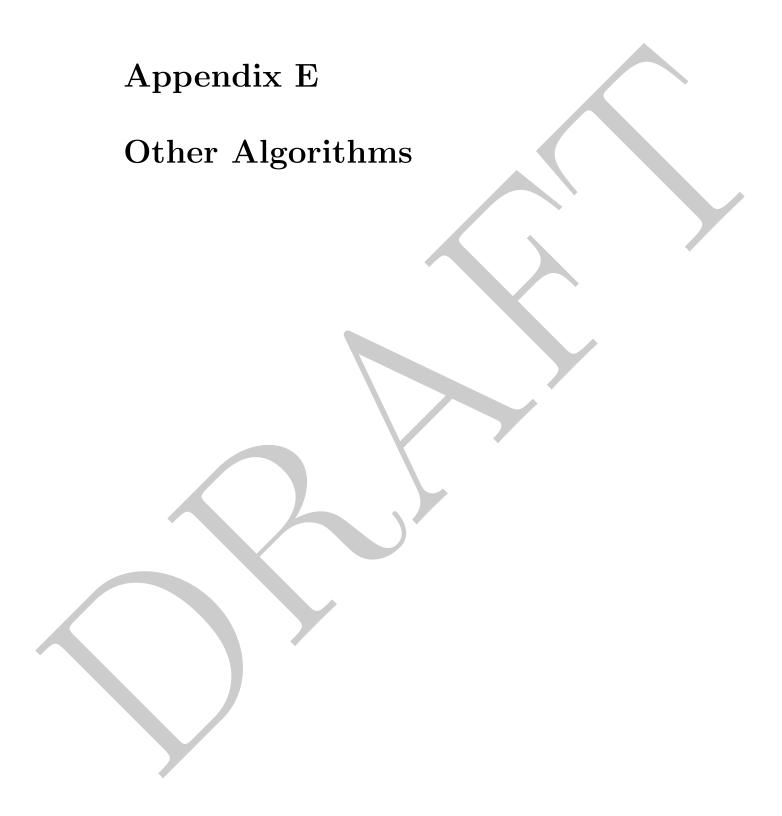












Appendix F

Table of FPGA and Software Data