

CH3. LINKED LIST AND RECURSION

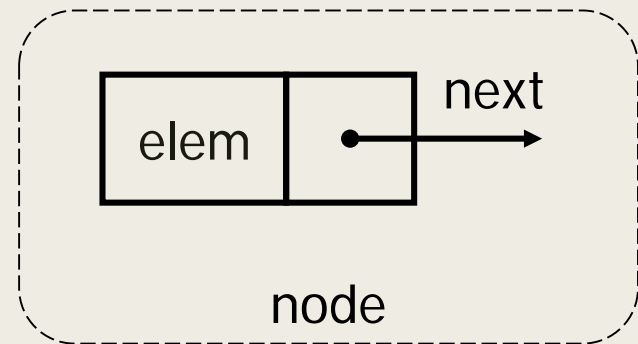
CSED233 Data Structure

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POSTECH

Singly Linked List (§ 3.2)

- A singly linked list is a concrete data structure consisting of a sequence of nodes
- Each node stores
 - *element*
 - *link to the next node*



```

class StringNode {                                // a node in a list of strings
private:
    string elem;                                   // element value
    StringNode* next;                             // next item in the list

    friend class StringLinkedList;                // provide StringLinkedList access
};

```

Code Fragment 3.13: A node in a singly linked list of strings.

```

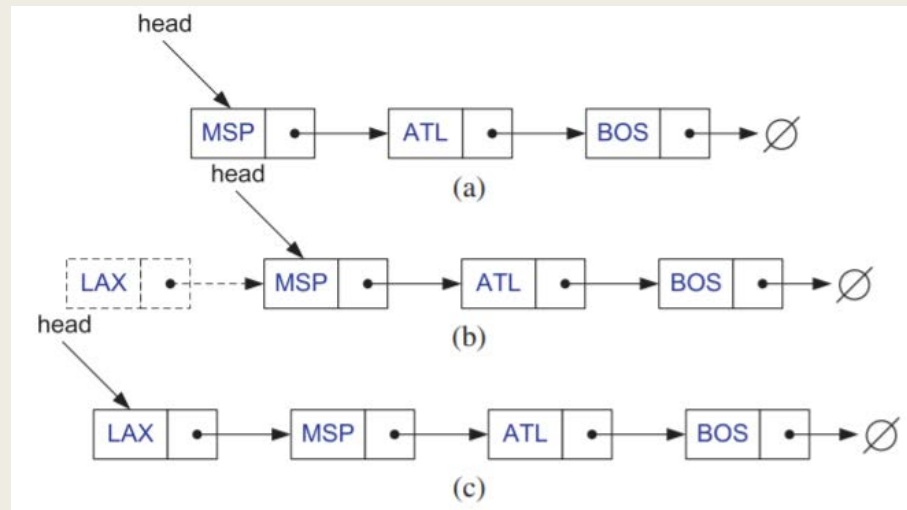
class StringLinkedList {                          // a linked list of strings
public:
    StringLinkedList();                           // empty list constructor
    ~StringLinkedList();                          // destructor
    bool empty() const;                           // is list empty?
    const string& front() const;                   // get front element
    void addFront(const string& e);                // add to front of list
    void removeFront();                            // remove front item list
private:
    StringNode* head;                             // pointer to the head of list
};

```

Code Fragment 3.14: A class definition for a singly linked list of strings.

Inserting at the Head

1. Create a new node
2. Set its *elem* value
3. Set its *next* link to the current head
4. Update head to point to new node

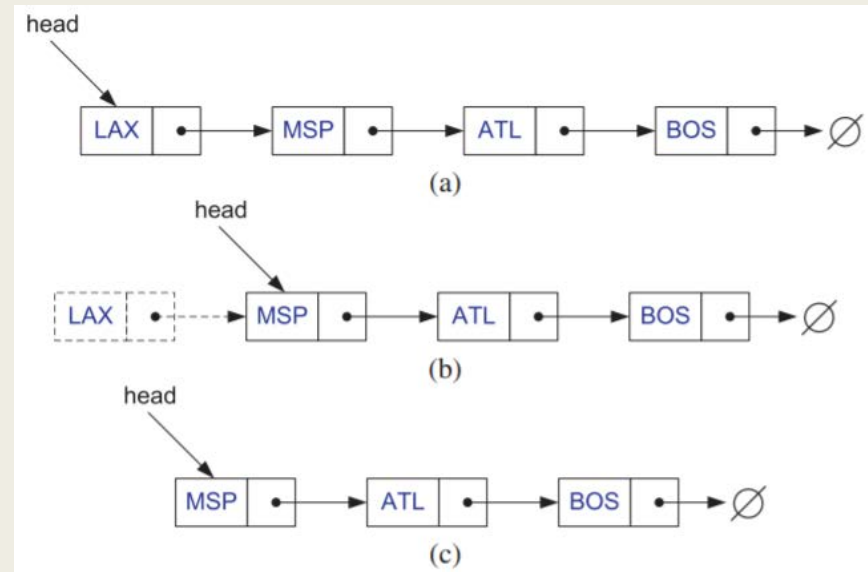


```
void StringLinkedList::addFront(const string& e) { // add to front of list
    StringNode* v = new StringNode;                // create new node
    v->elem = e;                                     // store data
    v->next = head;                                  // head now follows v
    head = v;                                        // v is now the head
}
```

Code Fragment 3.16: Insertion to the front of a singly linked list.

Removing at the Head

1. Update head to point to next node in the list
2. Remove the former first node



```
void StringLinkedList::removeFront() {  
    StringNode* old = head;           // remove front item  
    head = old->next;                 // save current head  
    delete old;                       // skip over old head  
}
```

Code Fragment 3.17: Removal from the front of a singly linked list.

Generic Singly Linked List

```
template <typename E>
class SNode {
private:
    E elem;
    SNode<E>* next;
    friend class SLinkedList<E>;
};
```

```
template <typename E>
class SLinkedList {
public:
    SLinkedList();
    ~SLinkedList();
    bool empty() const;
    const E& front() const;
    void addFront(const E& e);
    void removeFront();
private:
    SNode<E>* head;
};
```

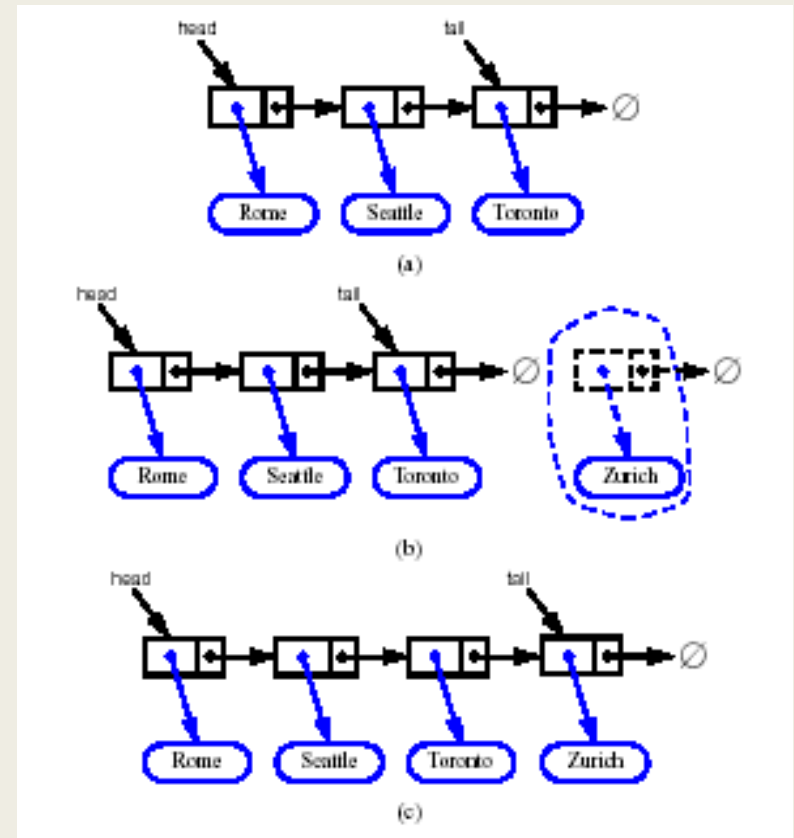
```
template <typename E>
void SLinkedList<E>::addFront(const E& e) {
    SNode<E>* v = new SNode<E>;
    v->elem = e;
    v->next = head;
    head = v;
}
```

```
template <typename E>
void SLinkedList<E>::removeFront() {
    SNode<E>* old = head;
    head = old->next;
    delete old;
}
```

```
SLinkedList<string> a;
a.addFront("MSP");
// ...
SLinkedList<int> b;
b.addFront(13);
```

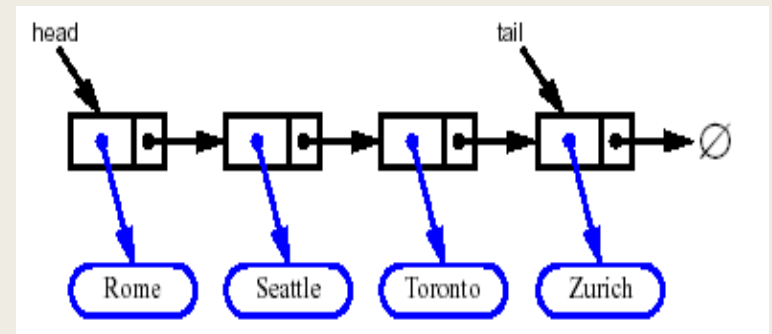
Inserting at the Tail (need to have “tail”)

1. Create a new node
2. Set its elem value
3. Set its next link to NULL
4. Have old last node point to new node
5. Update tail to point to new node



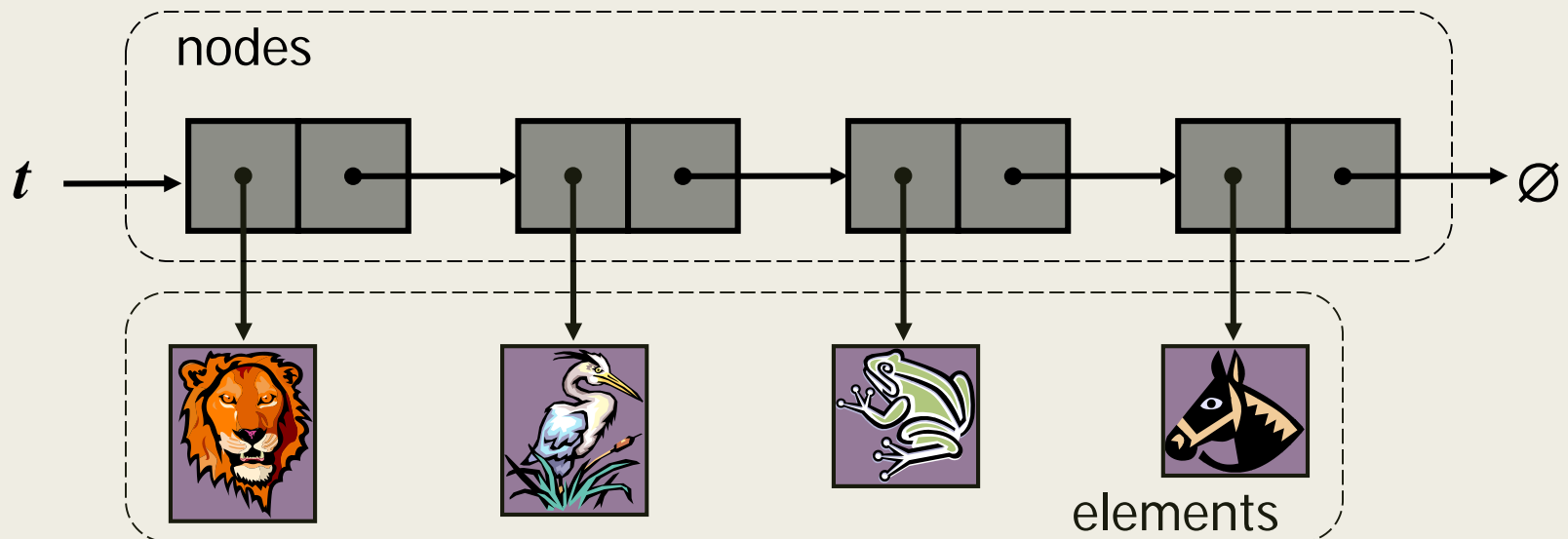
Removing at the Tail (need to have “tail”)

- Removing at the tail of a singly linked list is not efficient! Why?
- There is no constant-time way to update the tail to point to the previous node



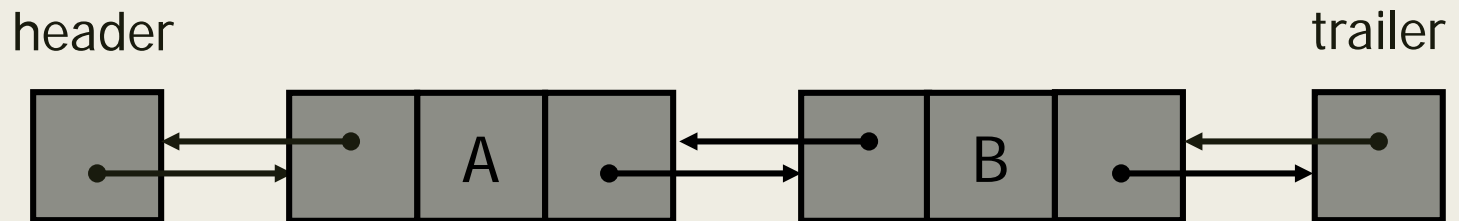
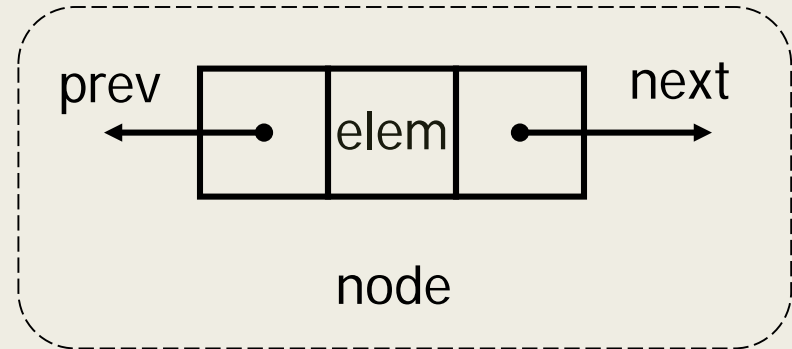
Stack as a Linked List (§ 5.1.3)

- We can implement a stack with a singly linked list
- The top element is stored at the first node of the list
- The space used is $O(n)$ and each operation of the Stack ADT takes $O(1)$ time



Doubly Linked List (§ 3.3)

- Each node stores
 - *link to the prev node*
 - *element*
 - *link to the next node*



```

typedef string Elem;           // list element type
class DNode {                 // doubly linked list node
private:
    Elem elem;                // node element value
    DNode* prev;              // previous node in list
    DNode* next;              // next node in list
    friend class DLinkedList; // allow DLinkedList access
};

```

Code Fragment 3.22: C++ implementation of a doubly linked list node.

```

class DLinkedList {           // doubly linked list
public:
    DLinkedList();            // constructor
    ~DLinkedList();           // destructor
    bool empty() const;       // is list empty?
    const Elem& front() const; // get front element
    const Elem& back() const;  // get back element
    void addFront(const Elem& e); // add to front of list
    void addBack(const Elem& e);  // add to back of list
    void removeFront();          // remove from front
    void removeBack();          // remove from back
private:                        // local type definitions
    DNode* header;              // list sentinels
    DNode* trailer;
protected:                     // local utilities
    void add(DNode* v, const Elem& e); // insert new node before v
    void remove(DNode* v);           // remove node v
};

```

Code Fragment 3.23: Implementation of a doubly linked list class.

Inserting

```

// insert new node before v
void DLinkedList::add(DNode* v, const Elem& e) {
    DNode* u = new DNode; u->elem = e; // create a new node for e
    u->next = v;                        // link u in between v
    u->prev = v->prev;                  // ...and v->prev
    v->prev->next = v->prev = u;
}

void DLinkedList::addFront(const Elem& e) // add to front of list
{ add(header->next, e); }

void DLinkedList::addBack(const Elem& e) // add to back of list
{ add(trailer, e); }
```

Code Fragment 3.26: Inserting a new node into a doubly linked list. The protected utility function `add` inserts a node z before an arbitrary node v . The public member functions `addFront` and `addBack` both invoke this utility function.

Removing

```
void DLinkedList::remove(DNode* v) {           // remove node v
    DNode* u = v->prev;                         // predecessor
    DNode* w = v->next;                         // successor
    u->next = w;                                // unlink v from list
    w->prev = u;
    delete v;
}

void DLinkedList::removeFront()                // remove from front
{ remove(header->next); }

void DLinkedList::removeBack()                 // remove from back
{ remove(trailer->prev); }
```

Code Fragment 3.27: Removing a node from a doubly linked list. The local utility function `remove` removes the node `v`. The public member functions `removeFront` and `removeBack` invoke this utility function.

Circularly Linked List and List Reversal (§ 3.4)

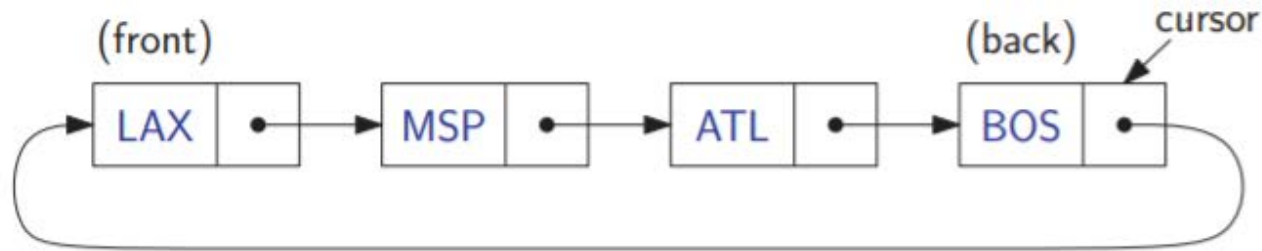


Figure 3.15: A circularly linked list. The node referenced by the cursor is called the back, and the node immediately following is called the front.

```
void listReverse(DLinkedList& L) {           // reverse a list
    DLinkedList T;                          // temporary list
    while (!L.empty()) {                    // reverse L into T
        string s = L.front(); L.removeFront();
        T.addFront(s);
    }
    while (!T.empty()) {                    // copy T back to L
        string s = T.front(); T.removeFront();
        L.addBack(s);
    }
}
```

Code Fragment 3.35: A function that reverses the contents of a doubly linked list *L*.

Recursion (§ 3.5)

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot (n-1) \cdot (n-2) \cdots 3 \cdot 2 \cdot 1 & \text{if } n \geq 1. \end{cases}$$

$$\text{factorial}(n) = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot \text{factorial}(n-1) & \text{if } n \geq 1. \end{cases}$$

```
int recursiveFactorial(int n) {           // recursive factorial function
    if (n == 0) return 1;                 // basis case
    else return n * recursiveFactorial(n-1); // recursive case
}
```

Code Fragment 3.36: A recursive implementation of the factorial function.

Linear Recursion (§ 3.5.1)

Algorithm LinearSum(A, n):

Input: A integer array A and an integer $n \geq 1$, such that A has at least n elements

Output: The sum of the first n integers in A

if $n = 1$ **then**

return $A[0]$

else

return LinearSum($A, n - 1$) + $A[n - 1]$

Code Fragment 3.38: Summing the elements in an array using linear recursion.

Algorithm ReverseArray(A, i, j):

Input: An array A and nonnegative integer indices i and j

Output: The reversal of the elements in A starting at index i and ending at j

if $i < j$ **then**

 Swap $A[i]$ and $A[j]$

 ReverseArray($A, i + 1, j - 1$)

return

Code Fragment 3.39: Reversing the elements of an array using linear recursion.

Binary Recursion (§ 3.5.2)

- Fibonacci numbers
 - 0, 1, 1, 2, 3, 5, 8, 13, ..

$$\begin{aligned}F_0 &= 0 \\F_1 &= 1 \\F_i &= F_{i-1} + F_{i-2} \quad \text{for } i > 1\end{aligned}$$

Algorithm BinaryFib(k):

Input: Nonnegative integer k

Output: The k th Fibonacci number F_k

if $k \leq 1$ **then**

return k

else

return BinaryFib($k - 1$) + BinaryFib($k - 2$)

Code Fragment 3.42: Computing the k th Fibonacci number using binary recursion.

Fibonacci execution

- Let n_k be the number of recursive calls by BinaryFib(k)
 - $n_0 = 1$
 - $n_1 = 1$
 - $n_2 = n_1 + n_0 + 1 = 1 + 1 + 1 = 3$
 - $n_3 = n_2 + n_1 + 1 = 3 + 1 + 1 = 5$
 - $n_4 = n_3 + n_2 + 1 = 5 + 3 + 1 = 9$
 - $n_5 = n_4 + n_3 + 1 = 9 + 5 + 1 = 15$
 - $n_6 = n_5 + n_4 + 1 = 15 + 9 + 1 = 25$
 - $n_7 = n_6 + n_5 + 1 = 25 + 15 + 1 = 41$
 - $n_8 = n_7 + n_6 + 1 = 41 + 25 + 1 = 67.$
- Note that n_k at least doubles every other time
- That is, $n_k > 2^{k/2}$. It is exponential!

A Better Fibonacci Algorithm

- Use linear recursion instead

Algorithm LinearFibonacci(k):

Input: A nonnegative integer k

Output: Pair of Fibonacci numbers (F_k, F_{k-1})

if $k = 1$ *then*

return $(k, 0)$

Else

$(i, j) = \text{LinearFibonacci}(k - 1)$

return $(i + j, i)$

- LinearFibonacci makes $k-1$ recursive calls