

# CS110: Matrix operations in C

Lecture 10

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# Arrays

- Matrix is
  - A two dimensional array
  - Array of one dimensional arrays
- Stored in memory
  - Row-wise (Row Major order)
  - Column-wise (Column Major order)

# Storage of matrix

- A - 2X3 matrix has elements
  - $A[0][0], A[0][1], A[0][2];$
  - $A[1][0], A[1][1], A[1][2];$
  - $A[2][0], A[2][1], A[2][2];$
- Assume it is an integer matrix and  $A[0][0]$  is stored at 1000

# Row-major order

- Store the first row, then, the second, then, the third and so on...
- A - 2X3 matrix has elements
  - $A[0][0], A[0][1], A[0][2]$ ; stored at 1000-1011
  - $A[1][0], A[1][1], A[1][2]$ ; stored 1012 - 1023
  - $A[2][0], A[2][1], A[2][2]$ ; stored 1024 - 1035

# Column-major order

- Store the first column, then, the second column, then, the third and so on..
- A - 2X3 matrix has elements
  - $A[0][0], A[1][0], A[2][0]$ ; stored at 1000-1011
  - $A[0][1], A[1][1], A[2][1]$ ; stored 1012 - 1023
  - $A[0][2], A[1][2], A[2][2]$ ; stored 1024 - 1035
- The C compiler assumes Row major order.

# Lab - 3

- Write a program that reads in the entries of a 3 by 3 matrix, and prints it out in the form of a matrix. The entries could be floating point entries too.

# Solution

```
#include<stdio.h>
main() {
    int i,j, a[3][3];
    for (i = 0; i <= 2; i++) {
        for (j = 0; j <=2; j++) {
            printf("Enter a[%d][%d] \n",i,j);
            scanf("%d",&a[i][j]);
        }
    }
}
```

# What Happens

```
for (i = 0; i <= 2; i++) {  
    for (j = 0; j <= 2; j++)  
    {  
        printf("Enter a[%d][%d] \n", i, j);  
        scanf("%d", &a[i][j]);  
    }  
}
```

For every  $i$ ,  $0 \leq i \leq 2$ ,  $j$  repeats thrice taking values 0, 1 and 2. Therefore when  $i = 0$  and  $j = 0$ ,  $a[0][0]$  shall be read in, when  $i = 0$  and  $j = 1$ ,  $a[0][1]$  shall be read in and so on ....



# To Print

```
for (i = 0; i <= 2; i++) {  
    for (j = 0; j <= 2; j++)  
    {  
        printf("%d ", a[i][j]);  
    }  
    printf("\n");  
}
```

Print all elements of a given row on the same line with a space in between. After one row is printed go to the next line.

# You have to

- Repeat above for Floating point numbers
- The effort would be to
  - Format the output properly
  - Let  $p$  be a floating point number
  - `printf("%16f",..)` prints a floating point with 16 spaces - if less then blanks are introduced.
  - Use the tab escape sequence `"\t"` to align.
  - This would be true for integers too.

## Next in Lab - 3

- Write a program that reads in orders of two matrices and decides whether two such matrices can be multiplied. Print out the decision.

```
main() {  
    int row1,col1,row2,col2;  
    printf("Enter number of rows and columns in  
    Matrix 1 \n");  
    scanf("%d %d",&row1,&col1);  
    //Similarly for Matrix 2  
    if (col1 != row2) { printf("Not possible\n"); exit(0);  
} // To use exit() we need to include <stdlib.h>
```

## Next in Lab - 3

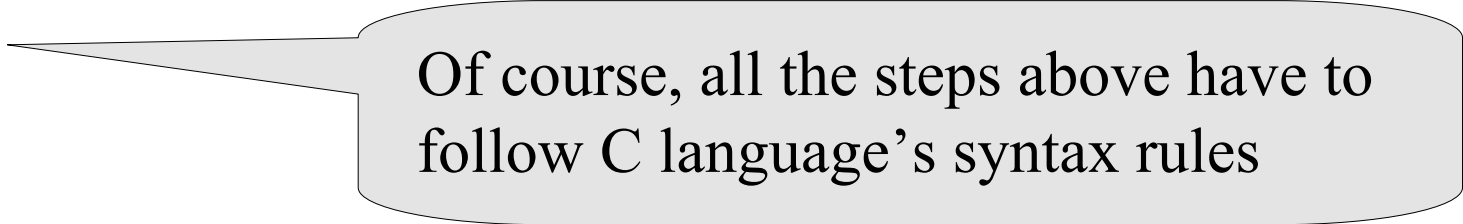
- Write a program that reads in two arbitrary matrices and multiplies them. Your output should be the two matrices and the resulting product matrix.

# Matrix Multiplication : Outline

main() {

1. declare all variables required
2. read in Matrix A
3. read in Matrix B
4. check if A and B are compatible to be multiplied
5. initialize Matrix C to have zeroes to begin with
6. multiply A and B to give C
7. print Matrix C

}



Of course, all the steps above have to follow C language's syntax rules

# Using Matrix Operations :

## Read a Matrix

```
main(){
    int a[11][11], b[11][11], c[11][11]; / *max size 10 by 10 */
    int i,j,k;
    int aRows, aCols, bRows, bCols, cRows, cCols;

    scanf("%d%d", &aRows, &aCols);
    for(int i = 1; i <= aRows; i++)
        for(int j = 1; j <= aCols; j++)
            scanf("%d", &a[i][j]);

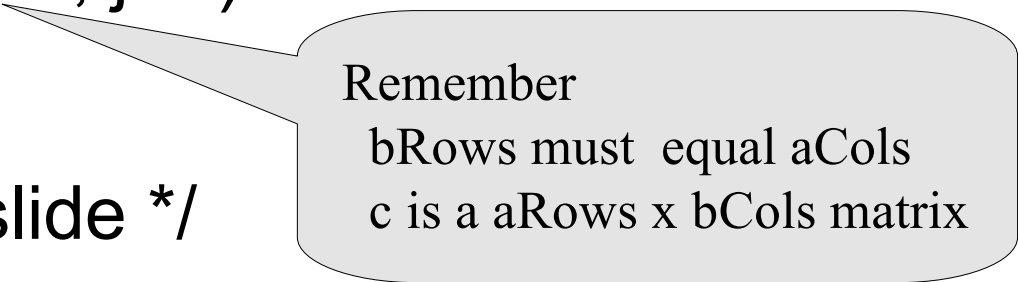
    /*continued on next slide */
```

# Read the other Matrix; Initialize the product

```
scanf("%d%d", &bRows, &bCols);  
for(i = 1; i <= bRows; i++)  
    for(j = 1; j <= bCols; j++)  
        scanf("%d", &b[i][j]);
```

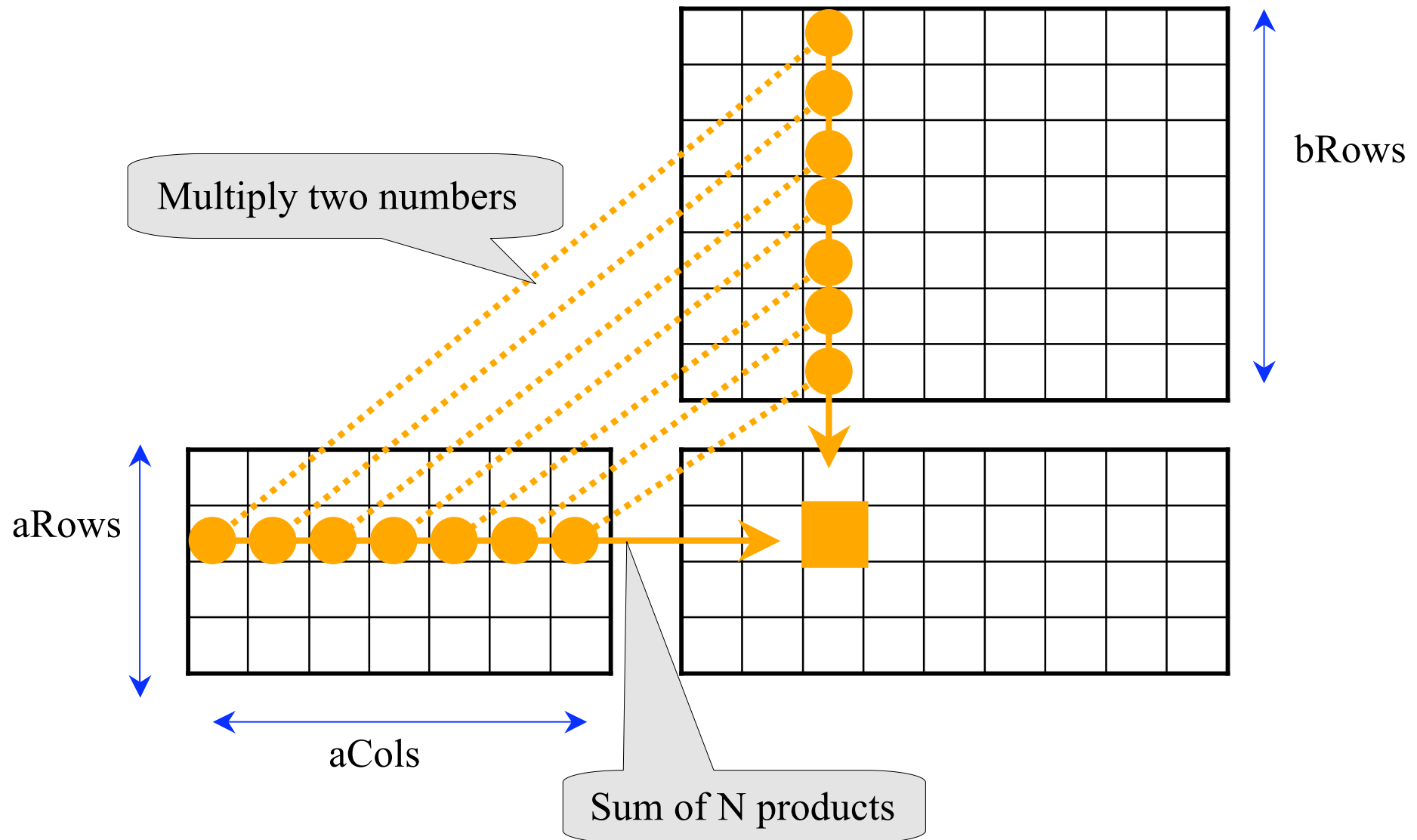
```
/* initialize entries in Matrix c to 0 */  
for(i = 1; i <= aRows; i++)  
    for(j = 1; j <= bCols; j++)  
        c[i][j] = 0;
```

```
/*continued on next slide */
```



Remember  
bRows must equal aCols  
c is a aRows x bCols matrix

# Matrix multiplication





# Multiply Matrices and Print the Result

```
/* multiply both the matrices and store in matrix c */
for(i = 1; i <= aRows; i++)
    for(j = 1; j <= bCols; j++)
        for(k = 1; k <= aCols; k++)
            c[i][j] += a[i][k]*b[k][j];

/* print matrix c */
for(i = 1; i <= aRows; i++){
    for(j = 1; j <= bCols; j++) /* print a row */
        printf("%d ", c[i][j]); /* notice missing \n */
    printf("\n"); /* print a newline at the end a row */
}
/* End of main program */
```

# Points to Ponder

- Some repetition in the program
  - Reading in matrix A seems to be similar to reading in matrix B
    - Except for changes in the number of rows and columns
- You will learn how to write functions in C later
  - Functions are written to avoid repeated actions
  - Example C program would look like

```
readMat(A, aRows, aCols);
readMat(B, bRows, bCols);
```
  - Function `readMat( )` must perform the operations desired

# Celebrity Problem

(More efficient Solution)

There are  $n$  persons in a room. A person is called a celebrity iff everyone knows him but he does not know anyone. So by definition if at all there is a celebrity there can be only one.

The problem is to find if there is a celebrity among the people in the room and if there is one identify him, by asking questions. The only kind of question allowed is “Does A know B?”.

# Celebrity Problem

(Naivest Solution)

'Naivest' Solution:

For each pair of persons A and B, ask the questions, does A know B and does B know A. Build the  $n \times n$  matrix  $\{a_{ij}\}$ .  $a_{ij}$  is 1 if the  $i$ th person knows the  $j$ th person and it is 0 otherwise. Now if there is a column  $i$ , which contains all 1s, and if row  $i$  contains all 0s except the diagonal, then the person  $i$  is a celebrity. If there is no such  $i$ , then there is no celebrity.

Number of questions asked =  $n*(n-1)$

# Celebrity Problem

(Better Solution)

Choose two persons A and B. Ask if A knows B. If answer is yes, then A cannot be the celebrity. If answer is no, then B cannot be the celebrity. So one person gets eliminated. Now if at all there is a celebrity he will be in the remaining  $n-1$  persons. Keep eliminating in this fashion. After  $n-1$  questions, only one person X, will be left. For each person  $Y \neq X$ , ask if Y knows X and X knows Y. Hence we can find if X is actually a celebrity.

Number of questions asked =  $n-1 + 2n-2 = 3n-3$

# Celebrity Problem

(Best Solution)

Simplifying assumption:  $n = 2^k$

Group the persons into pairs. So we will have  $n/2$  pairs. In each pair (A, B) eliminate one person by asking if A knows B. Now we will be left with  $n/2$  persons. Repeat the process till we are left with only one person.

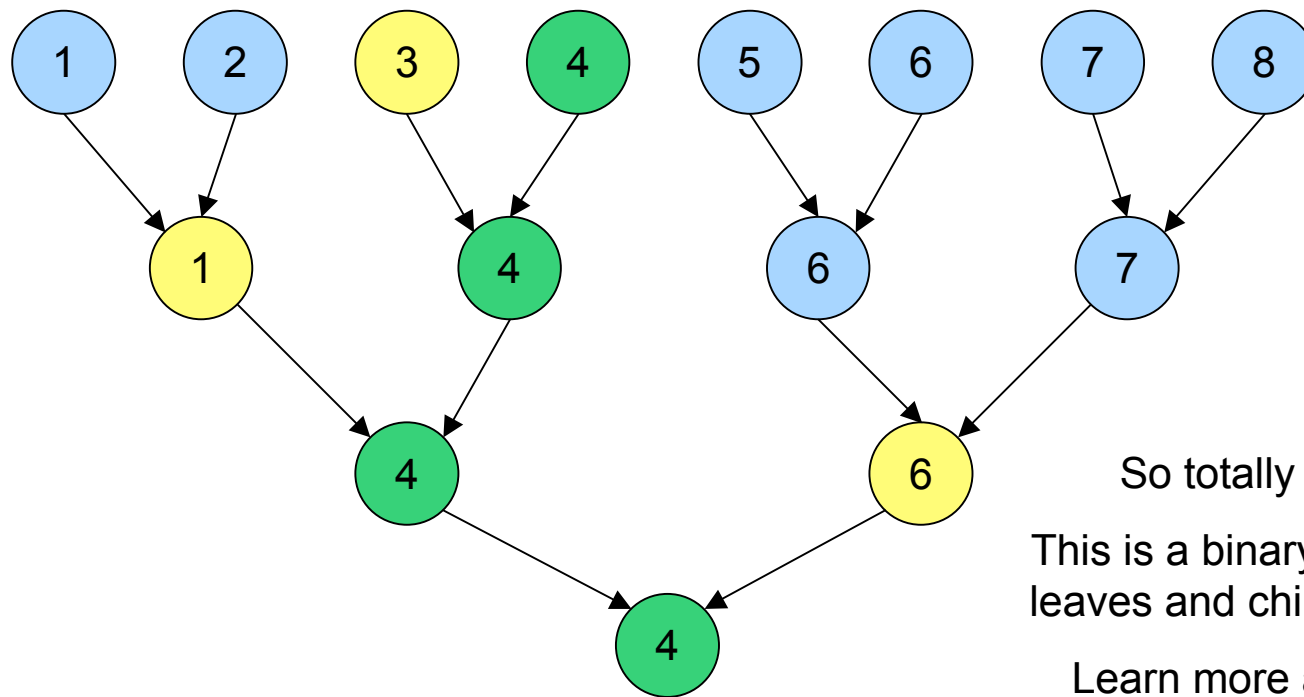
Number of times this process is repeated is obviously  $k$ , since each time the number reduces by half.

(cont.)

# Celebrity Problem

(Best Solution)

We can keep track of the tree thus obtained. Let us take a case having 8 persons. The tree looks likes something below.



So totally 7 Qs asked.

This is a binary tree with a root, leaves and children (at most 2).

Learn more about this later

# Generalizing above

- Given  $N$  people
  - $N/2$  Qs in level 1
  - $N/4$  Qs in level 2
  - 1 Question in level  $k$ , where,  $N = 2^k$
  - So total is  $N(1/2 + 1/4 + \dots 1/2^k)$
  - $N (1/2) (1 - 1/2^k)/(1/2)$
  - $N - N/2^k = N - 1$



# Celebrity Problem

(Best Solution)

Clearly, if at all there is a celebrity, it can be only 4. Also note that out of the  $2n-2 (= 6)$  questions asked to confirm whether 4 is actually a celebrity, we have already asked  $\log_2 n (= 3)$  questions (Notice the yellow circles). So once we maintain this tree, the number of additional questions that we need to ask is  $2n - 2 - \log_2 n$ .

Therefore the total number of questions asked in this case is  $n-1 + 2n - 2 - \log_2 n = 3n - 3 - \log_2 n$ .

There exists no algorithm which can do better than this, the proof of which is beyond the scope of this lecture. Extending it for  $n$  not a power of 2 is also interesting.

Thank You