

# Railroad Ink

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## Last Move

### Part A

#### Sets

$I$  – set of all inside squares, as (row,col) pairs

$O$  – set of all outside (start) squares, as (row, col) pairs

$S$  – set of all squares,  $S = I \cup O$

$P$  – set of all pieces

$T$  – set of all possible tiles that can be placed, with a piece, rotation and reflection

$E$  – set of the edge types  $E = \{\text{railway, highway}\}$

$P_s$  – set of all special pieces,  $P_s \subseteq P$

$T_j$  – set of all junction tile,  $T_j \subseteq T$

#### Data

$numstarts$  – the number of start pieces

$MP$  – the longest possible path that can be constructed

#### Functions

$adj(s)$  – returns the squares adjacent to square  $s$

$loose-ends(t, s)$  – returns the count of loose ends on internal edges if tile  $t \in T$  is placed at square  $s \in S$

$centre(s)$  – returns whether square  $s$  is a centre square

$clash(e)$  – returns the edge type that clashes with the given edge type

$tile-at(s)$  – returns the tile currently on the board at square  $s \in S$  if there is one

$free(s)$  – returns whether a piece can be placed at square  $s \in S$

$variations(p)$  – returns all the possible tiles associated with piece  $p$

$texttt{hand} - count(p)$  – how many of piece  $p$  are in the current hand

## Variables

### board variables

$x_{t,s}$  – 1 if tile  $t \in T$  is placed at square  $s \in S$

$y_{s,s',e}$  – whether there is a connection between  $s \in S$  and  $s' \in \text{adj}(s)$  s.t  $s < s'$  of type  $e \in E$

### connecting flow variables

$f_{s,s',e}$  – connecting flow from  $s \in S$  to  $s' \in \text{adj}(s)$  of type  $e \in E$  for joining the start squares

$f'_{s,e}$  – connecting flow from edge type  $e \in E$  to the other edge type (for junctions) at square  $s \in I$

$g_s$  – connecting flow from start square  $s \in O$  to the super sink

$h_s$  – 1 if there is any connecting flow from square  $s \in O$  to the super sink

$j$  – 1 if the bonus point for connecting all start squares is earned

### longest path variables

$a_{s,e}$  – whether square  $s \in I$  is the start square of the longest railway of type  $e \in E$

$b_{s,s',e}$  – longest path flow from  $s \in I$  to  $s' \in \text{adj}(s)$  s.t  $s' \in I$  of type  $e \in E$

$c_{s,s',e}$  – whether longest path flow of type  $e \in E$  is permitted from  $s \in I$  to  $s' \in \text{adj}(s)$  s.t  $s' \in I$

$d_{s,e}$  – whether square  $s \in I$  is part of the longest path of type  $e \in E$

## Objective

Maximise:

$$\left( 48 - 4 \cdot \sum_{s \in O} h_s + j \right) + \left( \sum_{\substack{s \in I \\ \text{centre}(s)}} x_{t,s} \right) + \left( \sum_{\substack{s \in I \\ e \in E}} d_{s,e} \right) - \left( \sum_{\substack{t \in T \\ s \in I}} x_{t,s} \cdot \text{loose-ends}(t, s) - \sum_{\substack{(s,s',e) \in y \\ s,s' \in I}} 2 \cdot y_{s,s',e} \right)$$

This objective function is the number of points earned, there are four components, in order:

- 1) Points earned from connections between start squares. Any square without a super sink connection is connected and scores points.  $j$  is the bonus point for connecting all start squares.
- 2) Points earned from placing pieces on the centre squares.
- 3) Points earned from the longest railway and highway.  $d_{s,e}$  variables are only set if a square is on the longest path.
- 4) Points lost due to errors. Points lost to errors can be counted as the total number of railways/highways on internal edges minus the pairs that are joined together.

## Constraints

### Piece placement constraints

$$\sum_{t \in T} x_{t,s} \leq 1 \quad \forall s \in S \quad (1)$$

$$x_{\text{tile-at}(s),s} = 1 \quad \forall s \in S \mid \sim \text{free}(s) \quad (2)$$

$$\sum_{\substack{s \in S \mid \text{free}(s) \\ t \in \text{variations}(p)}} x_{t,s} = \text{hand-count}(p) \quad \forall p \in P \quad (3)$$

$$\sum_{\substack{s \in S \mid \text{free}(s), \\ t \in \text{variations}(p)}} x_{t,s} \leq 1 \quad \forall p \in P_s \quad (4)$$

$$\sum_{\substack{s \in S \mid \text{free}(s), \\ p \in P_s, t \in \text{variations}(p)}} x_{t,s} \leq 1 \quad (5)$$

$$\sum_{\substack{s \in S, p \in P_s \\ t \in \text{variations}(p)}} \leq 3 \quad (6)$$

$$\sum_{\substack{t \in T \mid \\ t.R=e}} x_{t,s} + \sum_{\substack{t \in T \mid \\ t.L=\text{clash}(e)}} x_{t,s'} \leq 1 \quad \forall s = (r, c) \in S, e \in E \mid s' = (r, c+1) \in S \quad (7)$$

$$\sum_{\substack{t \in T \mid \\ t.B=e}} x_{t,s} + \sum_{\substack{t \in T \mid \\ t.T=\text{clash}(e)}} x_{t,s'} \leq 1 \quad \forall s = (r, c) \in S, e \in E \mid s' = (r+1, c) \in S \quad (8)$$

$$2 \cdot y_{s,s',e} \leq \sum_{\substack{t \in T \mid \\ t.R=e}} x_{t,s} + \sum_{\substack{t \in T \mid \\ t.L=e}} x_{t,s'} \quad \forall s = (r, c) \in S, e \in E \mid s' = (r, c+1) \in S \quad (9)$$

$$2 \cdot y_{s,s',e} \leq \sum_{\substack{t \in T \mid \\ t.B=e}} x_{t,s} + \sum_{\substack{t \in T \mid \\ t.T=e}} x_{t,s'} \quad \forall s = (r, c) \in S, e \in E \mid s' = (r+1, c) \in S \quad (10)$$

### Joining constraints

$$\sum_{s' \in \text{adj}(s)} f_{s,s',e} + f'_{s,e} = \sum_{s' \in \text{adj}(s)} f_{s',s,e} + f'_{s,\text{clash}(e)} \quad \forall s \in I, e \in E \quad (11)$$

$$1 + \sum_{\substack{s' \in \text{adj}(s) \\ e \in E}} f_{s',s,e} = \sum_{\substack{s' \in \text{adj}(s) \\ e \in E}} f_{s,s',e} + g_s \quad \forall s \in O \quad (12)$$

$$g_s \leq \text{numStarts} \cdot h_s \quad \forall s \in O \quad (13)$$

$$f_{s,s',e} \leq \text{numStarts} \cdot y_{s,s',e} \quad \forall s \in S, s' \in \text{adj}(s), e \in E \mid s' > s \quad (14)$$

$$f_{s,s',e} \leq \text{numStarts} \cdot y_{s',s,e} \quad \forall s \in S, s' \in \text{adj}(s), e \in E \mid s' < s \quad (15)$$

$$f'_{s,e} \leq \text{numStarts} \cdot \sum_{t \in T_j} x_{t,s} \quad \forall s \in I, e \in E \quad (16)$$

$$(numStarts - 1) \cdot j \leq \sum_{s \in O} (1 - h_s) \quad (17)$$

### Longest path constraints

$$\sum_{s \in I} a_{s,e} = 1 \quad \forall e \in E \quad (18)$$

$$c_{s,s',e} \leq y_{s,s',e} \quad \forall s \in S, e \in E, s' \in \text{adj}(s) \mid s' \in I \wedge s < s' \quad (19)$$

$$c_{s,s',e} \leq y_{s',s,e} \quad \forall s \in S, e \in E, s' \in \text{adj}(s) \mid s' \in I \wedge s > s' \quad (20)$$

$$\sum_{\substack{s' \in \text{adj}(s) \\ s' \in I}} b_{s',s,e} + MP \cdot a_{s,e} = \sum_{\substack{s' \in \text{adj}(s) \\ s' \in I}} b_{s,s',e} + d_{s,e} \quad \forall s \in I, e \in E \quad (21)$$

$$b_{s,s',e} \leq MP \cdot c_{s,s',e} \quad \forall s \in S, e \in E, s' \in \text{adj}(s) \mid s' \in I \quad (22)$$

$$\sum_{\substack{s' \in \text{adj}(s) \\ s' \in I}} c_{s,s',e} \leq 1 \quad \forall s \in I, e \in E \quad (23)$$

These constraints are designed as follows:

- (1) Only one tile can be placed at each square.
- (2) All squares that are already occupied must finish with that piece there.
- (3) Use pieces as required in your hand.
- (4) Each special piece can only be played once.
- (5) Only one special piece can be played this turn.
- (6) Only three specials can be played total.
- (7, 8) Cannot place two pieces next to each other that directly clash with each other (i.e have a railway run into a highway).
- (9, 10) Can only set the  $y_{s,s',e}$  variable (connecting two squares with edge type  $e$ ) if there are matching edges between the squares.
- (11) The joining inflow at internal squares (from adjacent squares or converted through a junction) must equal the outflow (to adjacent squares or passed through the junction).
- (12) The joining inflow at start squares (inflow of 1 plus from adjacent square) must equal the outflow (flow to adjacent square and to the super sink).
- (13) The  $h_s$  variable must be set if there is any flow to the super sink, i.e  $g_s > 0$ .
- (14, 15) There can only be joining flow when there is a connection between the two squares of that edge type.
- (16) If there is a junction at a given square, joining flow can readily flow between the two edge types. This is only not permitted for overpasses.

- (17) An extra point can be awarded if all start pieces are joined together.
- (18) The longest path of each type can only have one start square.
- (19, 20) The longest path of a certain type can only travel on an edge that has a connection of that type.
- (21) The inflow of the longest path (from adjacent squares, plus a large inflow at the start square) must be greater than the outflow (to adjacent squares, plus points allocated to this square)
- (22, 23) These constraints prevent branching of the longest path. There may only be outflow in one direction from any given square.