

Multiple Source Shortest Path with unit weights

1 Introduction

Given: Let G be a directed graph (V, \vec{E}) , embedded on a surface with genus g . All edge weights are unit.

Find: Consider boundary f of G . $\forall v \in f$, find a shortest path to $\forall u \in V$.

Let T be the BFS (Breadth first search) tree of G , and C be the BFS co-tree in G . Then there is exactly $2g$ leftover edges $L = e_1, e_2, \dots, e_{2g}$.

There exists a unique cycle λ_i in $C \cup e_i$, and $(\lambda_1, \lambda_2, \dots, \lambda_{2g}) = \Lambda$ defines homology basis.

We define homological signature of an edge as follows:

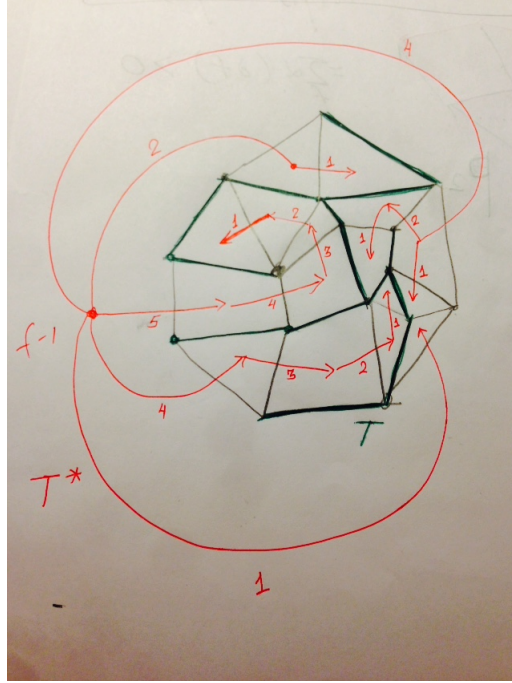
$$[e]_i = \begin{cases} 1 & , \text{if } e \in \lambda_i \\ -1 & , \text{if } \text{rev}(e) \in \lambda_i \\ 0 & , \text{otherwise} \end{cases}$$

Furthermore, we define α recursively as follows:

$$\alpha(\vec{e}^*) = \begin{cases} 1 & , \text{if } \vec{e}^* \text{ is a leaf edge of } C \\ \sum_{\text{tail}(\vec{e}^*) = \text{head}(\vec{e}')} \alpha(\vec{e}') & , \text{otherwise} \end{cases}$$

We can extend above definition with $\alpha(\vec{e}) = \alpha(\vec{e}^*)$ and $\alpha(e)^* = -\alpha(\text{rev}(\vec{e}^*))$.

Let $\tilde{w}(\vec{e}) = \langle 1, [\vec{e}] \rangle$ be new weight vector for each edge in G .



Def: An edge \vec{e} is "holier" than \vec{e}' , if $\tilde{w}(\vec{e}) < \tilde{w}(\vec{e}')$ in lexicographic comparison. Therefore, we can define "holiness" of any $S \subset G$ as follows:

$$\text{Ho}(S) = \sum_{\vec{e} \in S} \tilde{w}(\vec{e})$$

Holiest tree is a spanning tree with minimal "holiness". We build Holiest tree rooted at r , using slight tweak in the Bellman-Ford algorithm for finding shortest path tree rooted at r .

Observation-1: Since first components are unit for all edges,

BuildHoliestTree(G, \tilde{w}, r):

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Set  $d[r] \leftarrow \langle 0, [0], 0 \rangle$ 
 $\text{pred}(r) \leftarrow \text{NULL}$ 
for all  $v : v \neq r$ 
   $d[v] \leftarrow \langle \infty, [\infty], \infty \rangle$ 
   $\text{pred}(v) \leftarrow \text{NULL}$ 
put  $r$  into queue
while queue is not Empty:
  Let  $s \leftarrow$  dequeue item
  for all  $u : s \rightarrow u$ 
    if  $u$  is not marked
      mark  $u$  and put in the queue
    if  $\text{isTense}(s \rightarrow u)$ 
       $\text{relax}(s \rightarrow u)$ 

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isTense($s \rightarrow u$):

return $d[s] + \tilde{w}(s \rightarrow u) < d[u]$

relax($s \rightarrow u$):

$d[u] \leftarrow d[s] + \tilde{w}(s \rightarrow u)$
 $\text{pred}[u] \leftarrow s$

Observation-2: Each edge will be added once to the queue.

Corollary: Each edge will be relaxed at most once.

Lemma-1: If there is no tense edge in G , then for each $v : r \rightarrow \dots \rightarrow \text{pred}(\text{pred}(v)) \rightarrow \text{pred}(v) \rightarrow v$ is the holiest path from r to v .

Proof: Let's prove it by induction on $d[v]$ distance from the root r .

Base: $d[v] = 0$, then $v = r$, so the claim is definitely true.

Induction Step: Suppose the claim is true for all vertex $v \in V$ such that $d[v] < d$ for some d . Consider vertex v such that $d[v] = d$. By induction hypothesis, all vertices with $d[u] = d - 1$ have "holiest" path correctly updated. By definition, $d[v] = \min_{s \rightarrow v} d[s] + \tilde{w}(s \rightarrow v)$, here $d[s] = d - 1$. By Induction hypothesis, $d[s]$ is "holiest" path, so if there is no tense edge in G then $d[v] = \min_{s \rightarrow v} d[s] + \tilde{w}(s \rightarrow v)$ holds. \square

Corollary: The algorithm will produce "holiest" tree rooted at r in linear time.

We now have produced our initial "Holiest" tree.

2 Minimum Cost Flow Problem and methods to solve them

There are several ways to define minimum cost flow and with different types of flows (non-negative and skew-symmetric), capacity or upper bound, edge demand or lower bound, edge cost, and with flow balance. The standard way to define:

$$\begin{array}{l}
 \text{min } \langle \text{cent}, \text{flow} \rangle \\
 \text{s.t } \sum_{u \rightarrow v} \text{flow}(u \rightarrow v) - \sum_{v \rightarrow w} \text{flow}(v \rightarrow w) = 0 \\
 \text{flow}(u \rightarrow v) = u(u \rightarrow v) \\
 \text{flow} \geq 0
 \end{array}$$

The easiest solution we can find is using augmenting cycle:

We can run any maximum flow algorithm to find feasible solution to the problem (neglecting the cost). Let the augmenting cycle be a cycle with negative cost. Then sending a flow through this cycle would reduce the total cost, while maintaining the feasible property.

We can also define reduced cost as follows:

Let $\phi(v)$ be a any potential function on a vertex. Then $\bar{\text{cost}}(u \rightarrow v) = \phi(u) - \phi(v) + \bar{\text{cost}}(u \rightarrow v)$ satisfies the condition that $\text{cost}(C) = \bar{\text{cost}}(C)$ for any cycle C .

Lemma: A feasible flow f is optimal \leftrightarrow there is no augmenting cycle in the residual graph.

Proof: \rightarrow Suppose the flow is optimal. If there is an augmenting cycle C , then we can send flow through C and reduce the cost of current flow, contradicting the f is optimal.

\leftarrow Suppose there is no augmenting cycle. Let $\phi(v) = [\text{Shortest path from } s \text{ to } v \text{ with respect to the cost function}]$. Then $\bar{\text{cost}}(u \rightarrow v) = \phi(u) - \phi(v) + \text{cost}(u \rightarrow v) \geq 0$. For the new cost function, sending more flow in a new residual graph would increase the cost of any flow, therefore f is optimal. \square

We can find the augmenting cycle in a graph systematically as follows:
Let T be any fixed spanning tree of G . Define new potential function for each vertex

$$\text{slack}_T(u \rightarrow v) = \begin{cases} 0 & , \text{if } u \rightarrow v \in T \\ \sum_{e \in \text{cycle in } T \cup \{u \rightarrow v\}} \text{cost}(u \rightarrow v) & , \text{Otherwise} \end{cases}$$

Above definition preserves the property that $\text{slack}_T(C) = \text{cycle}(C)$ for any cycle C in G . Therefore, we can essentially find negative reduced cost edge in residual graph and push flow through it and do pivoting. Spanning tree T will be updated as follows:

Update T:

- Find bottleneck capacity in a cycle
- Push the flow amount equal to bottleneck capacity through the cycle
- Pivot out the bottleneck capacity edge, and pivot in the dart with negative reduced cost
- Recompute vertex potentials

3 Finding pivot quickly

- What data structure do we maintain in the G^* ? Finding shortest path in network can also be understood as a Linear Programming problem as follows:
- How do we find next pivot quickly using above structure?

There are two ways to represent the flow in the graph:

- $f(u \rightarrow v) = -f(v \rightarrow u)$, for all edges
- $f(u \rightarrow v) \geq 0$ and $f(u \rightarrow v) = 0$ if $f(u \rightarrow v) > 0$

Transshipment problem

$$\begin{array}{l} \min < f, \$ > \\ \text{s.t } \partial f = b \\ f \geq 0 \end{array}$$

Under generic assumption on $f, \$$:

- Basis spanning tree T , there is a unique flow $flow_T$ that satisfies the condition $\partial flow_T = b, flow_T(e)$ is nonzero only for edges in T .
- There exist a unique spanning tree T_{OPT} with $flow_{T_{\text{OPT}}}$ is optimal.

Subsequently, we can define slack as follows:
For fixed spanning tree T , there is unique cycle $C = T \cup \{e\}$ for edge e not in T . $slack(e) = \sum_{l \in C} s(l)$, and 0 otherwise. **Observe that slack is not negative, since otherwise there is no optimal solution to our LP**

The main LP is:

$$\boxed{\begin{array}{l} \min < f, slack_T > \\ \text{s.t } \partial f = \partial flow_T \\ f \geq 0 \end{array}}$$

$$\boxed{\begin{array}{l} \min < s, flow_T > \\ \text{s.t } \partial s = \partial slack_T \\ s \geq 0 \end{array}}$$

And in the case of non-planar embedded graph with genus g :

$$\boxed{\begin{array}{l} \min < f, slack_T > \\ \text{s.t } \partial f = \partial flow_T \\ f \geq 0 \end{array}}$$

$$\boxed{\begin{array}{l} \min < s, flow_T > \\ \text{s.t } \partial s = \partial slack_T \\ [s] = [slack_T] \\ s \geq 0 \end{array}}$$

Consider the primal LP, We can rewrite the constraints as follows:

$$\begin{aligned} \partial f = \partial flow_T &\iff \sum_u f(u \rightarrow v - f(v \rightarrow u)) \\ &\iff \gamma(u) \sum_u f(u \rightarrow v) - f(v \rightarrow u) = \sum_{u \rightarrow v} f(u \rightarrow v)(\gamma(u) - \gamma(v)) = \sum_{u \rightarrow v} f(u \rightarrow v)\gamma(u \rightarrow v) \end{aligned}$$

Observe here that $\gamma(u \rightarrow v) = -\gamma(v \rightarrow u)$ We can define s in terms of γ by setting the negative values to be 0, and we will get the exact dual program defined above.

4 Analysis

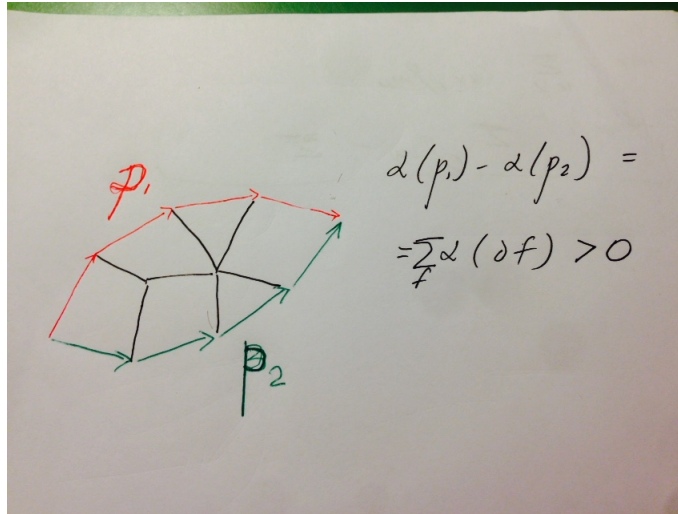
- Building initial tree
- Pivoting
- Number of times each edge is pivoted
- Overall running time

Couple questions regarding the slack and flow:

- If the fixed tree T is arbitrary tree (not necessarily the Holiest Tree, then the flow in the answer does not have to be optimal) but solution to the linear program $\min < f, slack_T >$ is equal to the answer from fixed tree T , not necessarily the optimal solution. That is because for each vertex, the demand satisfies the constraint and if we consider the tree T , then the value of $\min < f, slack_T >$ would be 0, implying it is the optimal solution. (Sum cannot be negative since otherwise there is no optimal solution)
- The reason we picked the slack as the way we defined is due to the fact that slack is not negative, ensuring that the nothing bad happens.
- What makes the non-planar case special with $2g$ extra constraints?
- How does the slack in dual representation help us to find the pivots quickly?

5 Additional

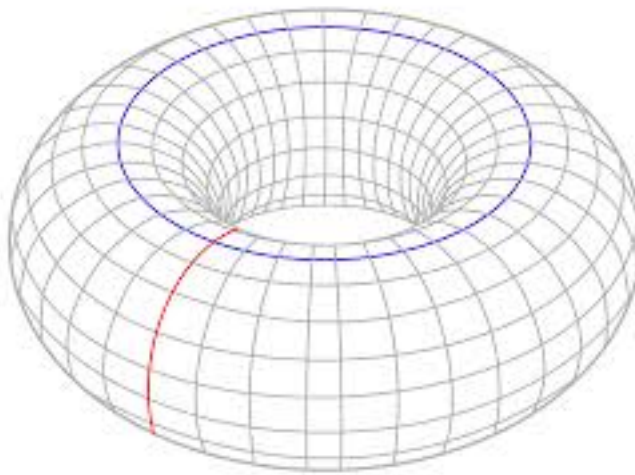
NOTE: Necessity of the α definition on the edges for "Holiness"
Consider the following picture:



By the definition of *alpha*:

$$\alpha(p_1) - \alpha(p_2) = \sum_f \alpha(df) > 0$$

This will ensure that any two paths p_1, p_2 , whose $w(p_1) = w(p_2)$ and $[p_1]_\Lambda = [p_2]_\Lambda$, has $\alpha(p_1) \neq \alpha(p_2)$

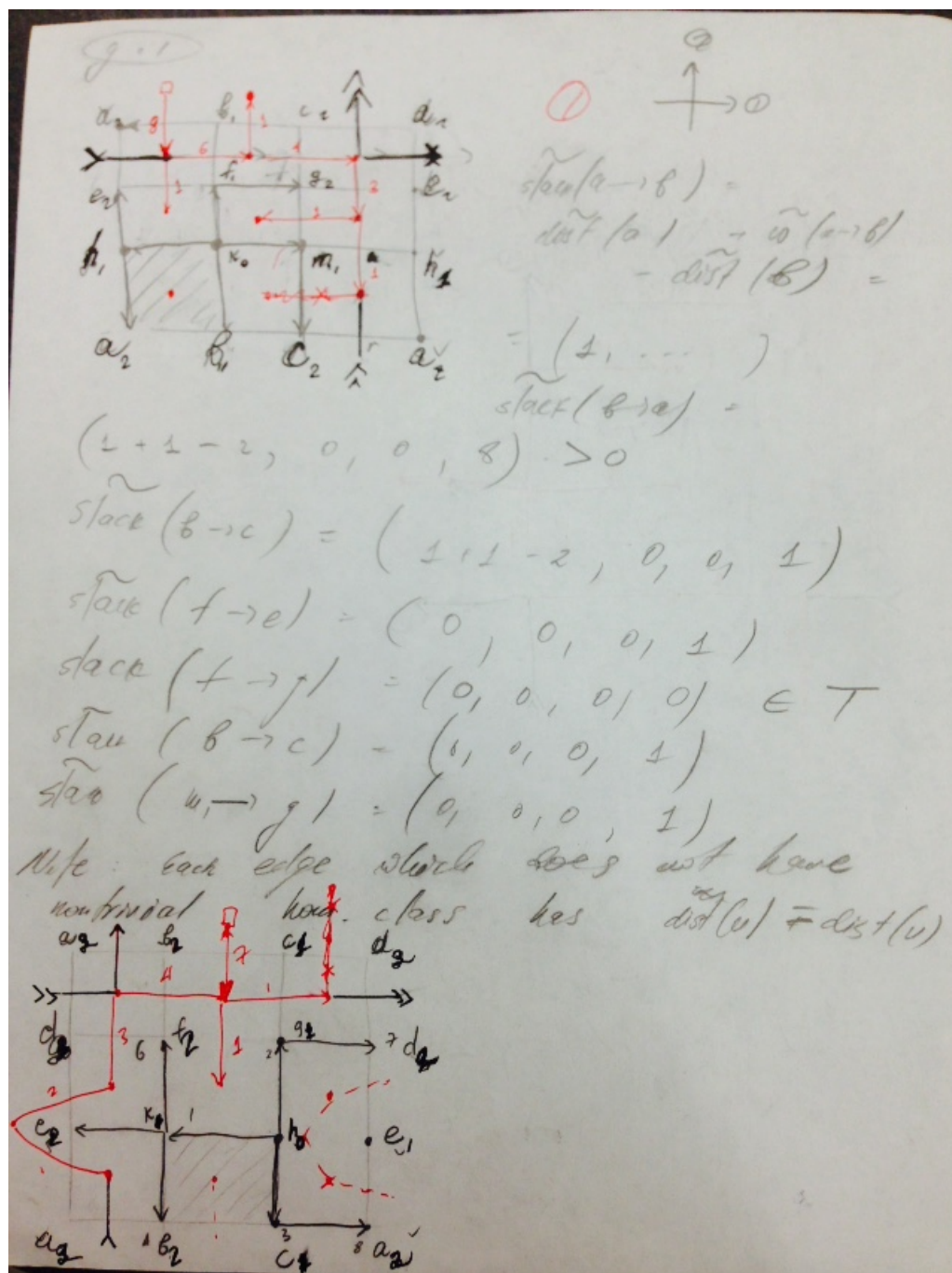


[http : //en.wikipedia.org/wiki/Homology\(mathematics\)](http://en.wikipedia.org/wiki/Homology(mathematics))

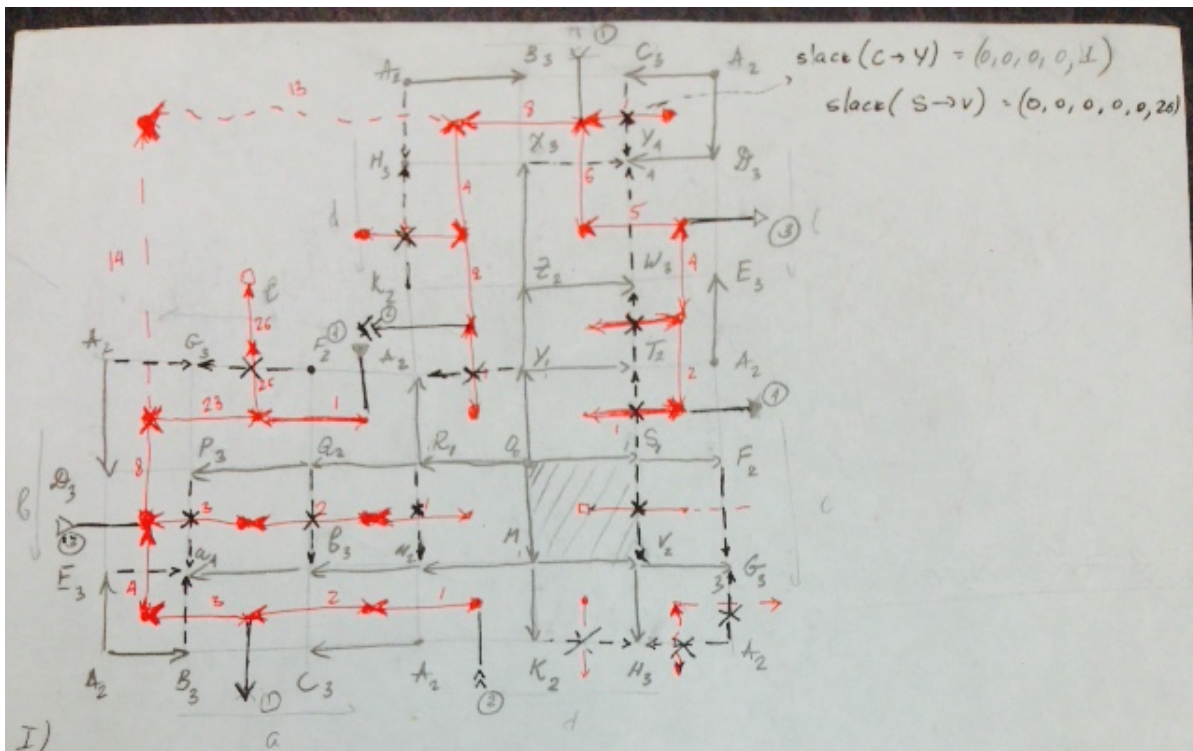
6 Working on examples:

Difference of Holiest Tree and leftmost tree:

On genus $g = 1$ surface:



On genus $g = 2$ surface:



$$s_{k\alpha}(b_5 > c) > 0$$

$$\text{stack}(A \rightarrow H) = (0, -1, -1, -1, -1, -13)$$

$$\widetilde{\text{stack}}(B \rightarrow a) = (0, -1, -1, 0, 0, -3)$$

for 1, 2, 3, 4 cycles.

11) state $(A \rightarrow 11) = (0, 1$

$$\widehat{\text{face}}(w-ry) = (0, 0, 0, 5, 1, 5)$$

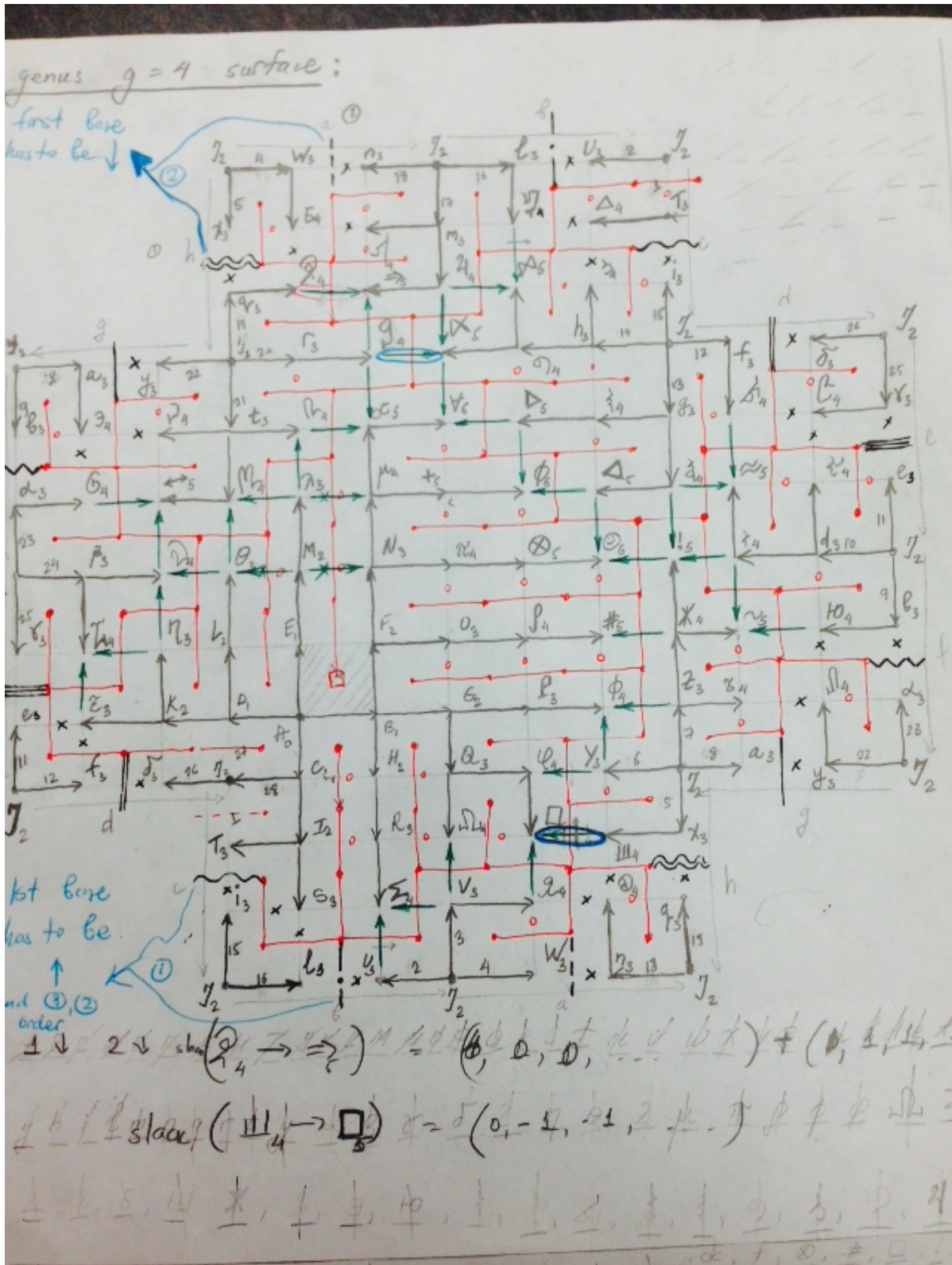
$$\text{span} \{A, B\} = (0, 4, -1, -1, 1, 14)$$

$$\widehat{A} (E - A) = \begin{pmatrix} 0 & 1 & -1 & 0 & 0 & 4 \end{pmatrix}$$

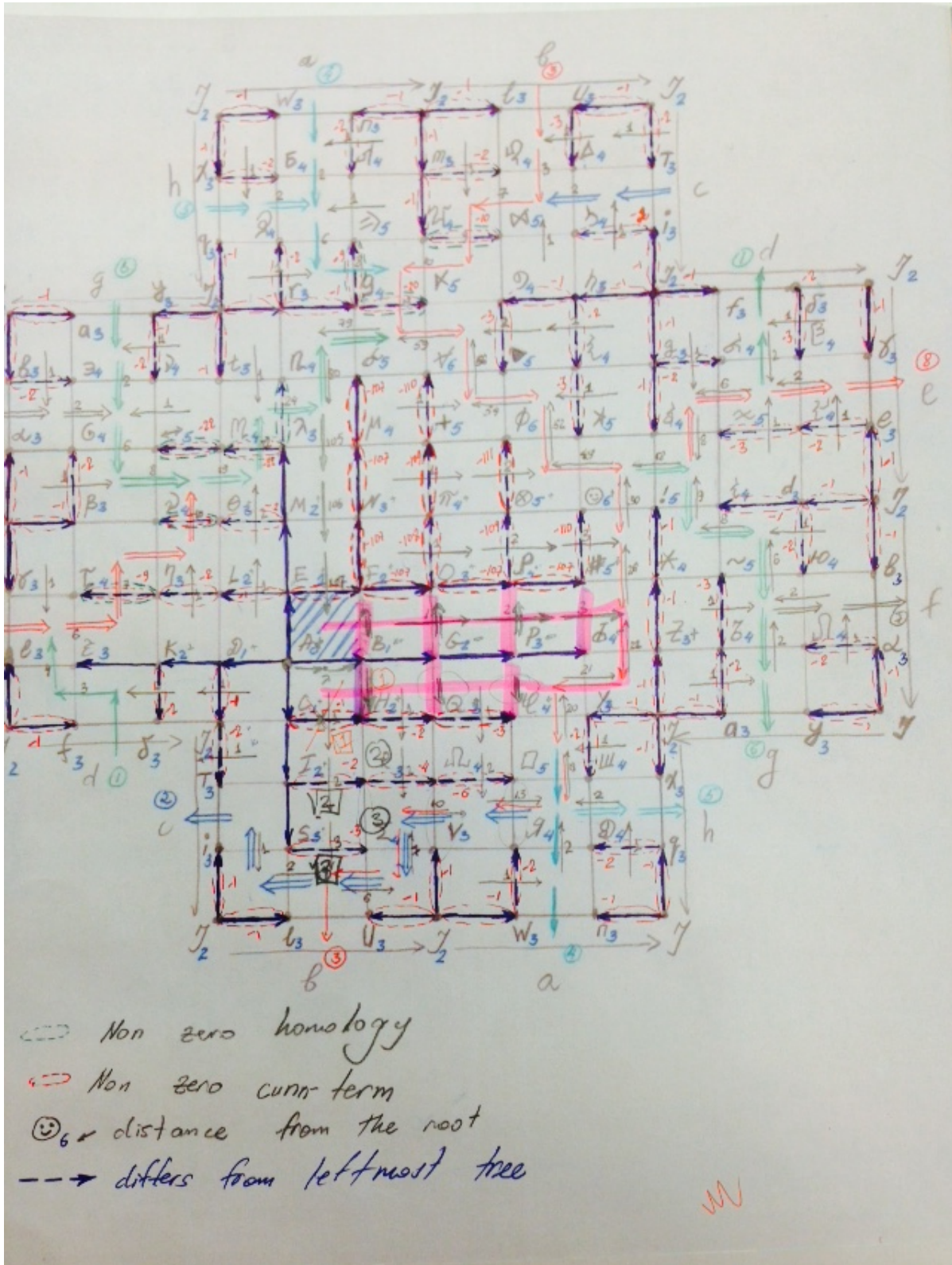
$$\text{slak}(B \rightarrow a) = (0, 1, -1, 0, 0, 3)$$

So we might be able to argue \exists ordering and direction s.t. $\text{stock}(u \rightarrow v) \geq 0$ for all $u \rightarrow v \notin T$.

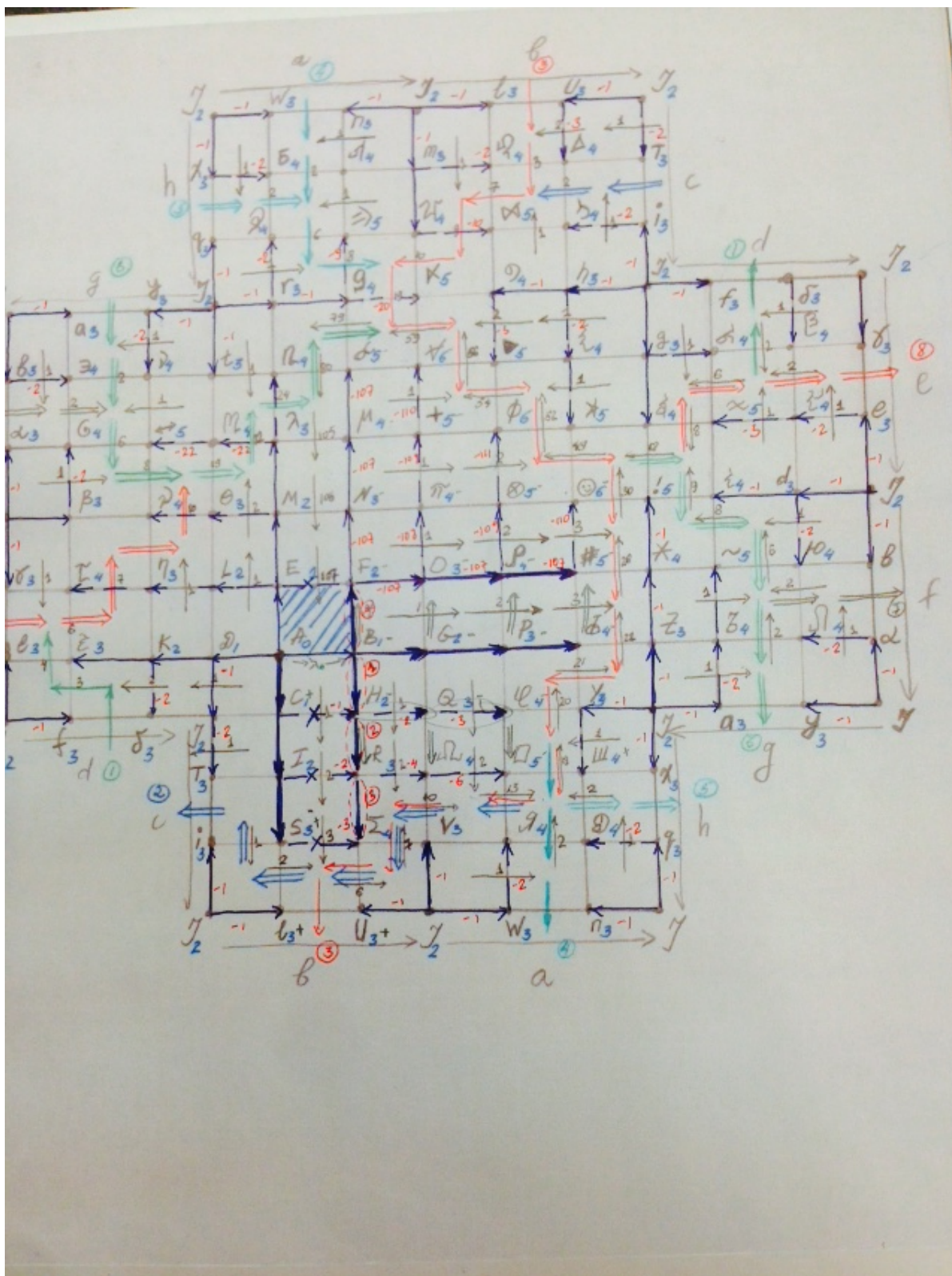
On genus $g = 4$ surface with initial Holy Tree build:



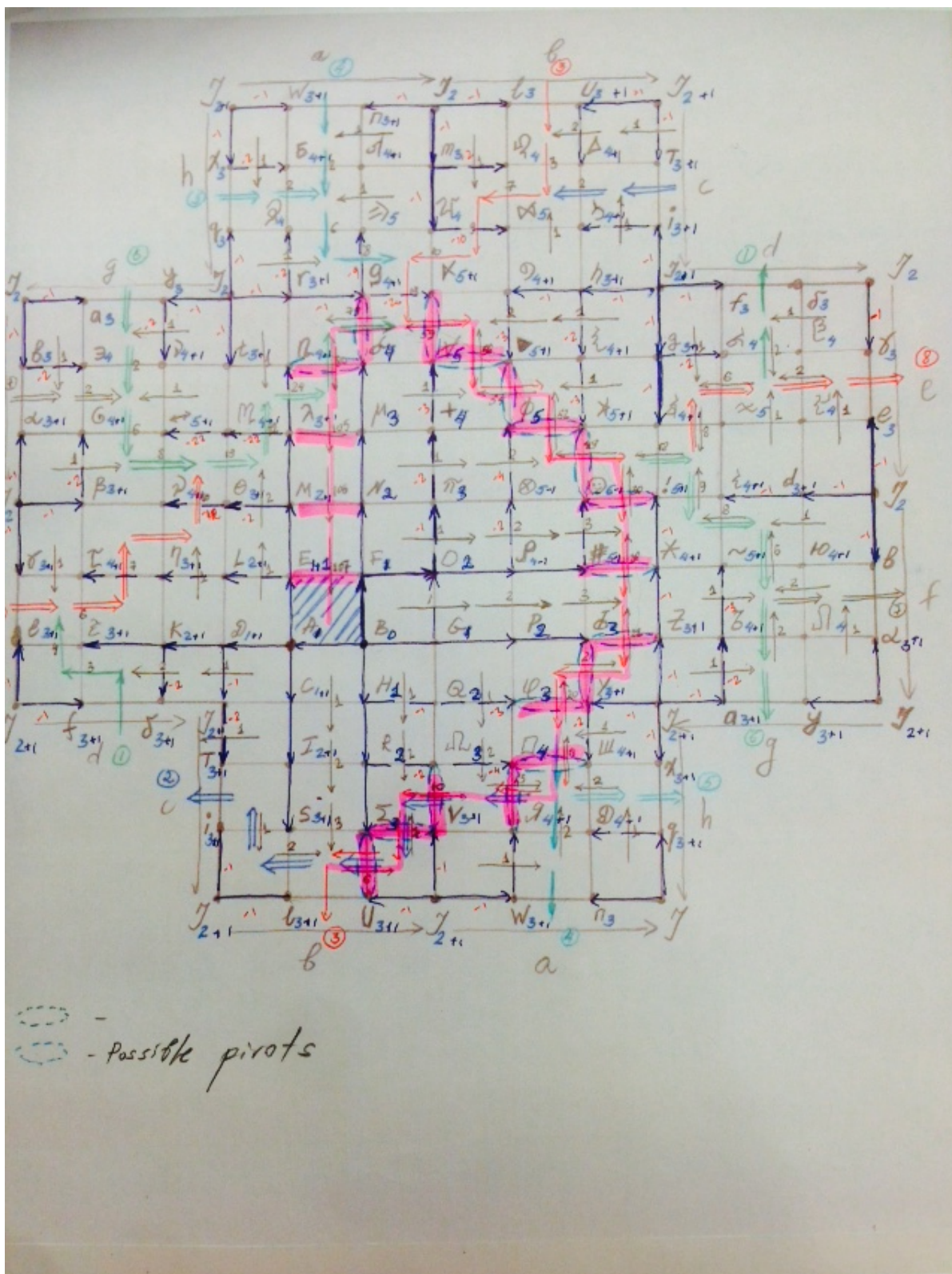
On genus $g = 4$ surface with initial pivot:



On genus $g = 4$ surface with initial pivot:



On genus $g = 5$ surface with initial pivot:



7 References:

- Cabello, Sergio, Erin W. Chambers, and Jeff Erickson. "Multiple-source shortest paths in embedded graphs." *SIAM Journal on Computing* 42.4 (2013): 1542-1571.
- Eisenstat, David, and Philip N. Klein. "Linear-time algorithms for max flow and multiple-source shortest paths in unit-weight planar graphs." *Proceedings of the forty-fifth annual ACM symposium on Theory of computing*. ACM, 2013.
- Erickson, Jeff. "Maximum flows and parametric shortest paths in planar graphs." *Proceedings of the twenty-first annual ACM-SIAM symposium on Discrete Algorithms*. Society for Industrial and Applied Mathematics, 2010.