

## Lecture 2b: Transistors, Gates, and Boolean Logic

**CSCI11: Computer Architecture  
and Organization**

Chapter 3.1-3.3.3 in [Patt and Patel:  
Introduction to Computing Systems...](#)

All of [AspenCore Boolean Algebra](#)

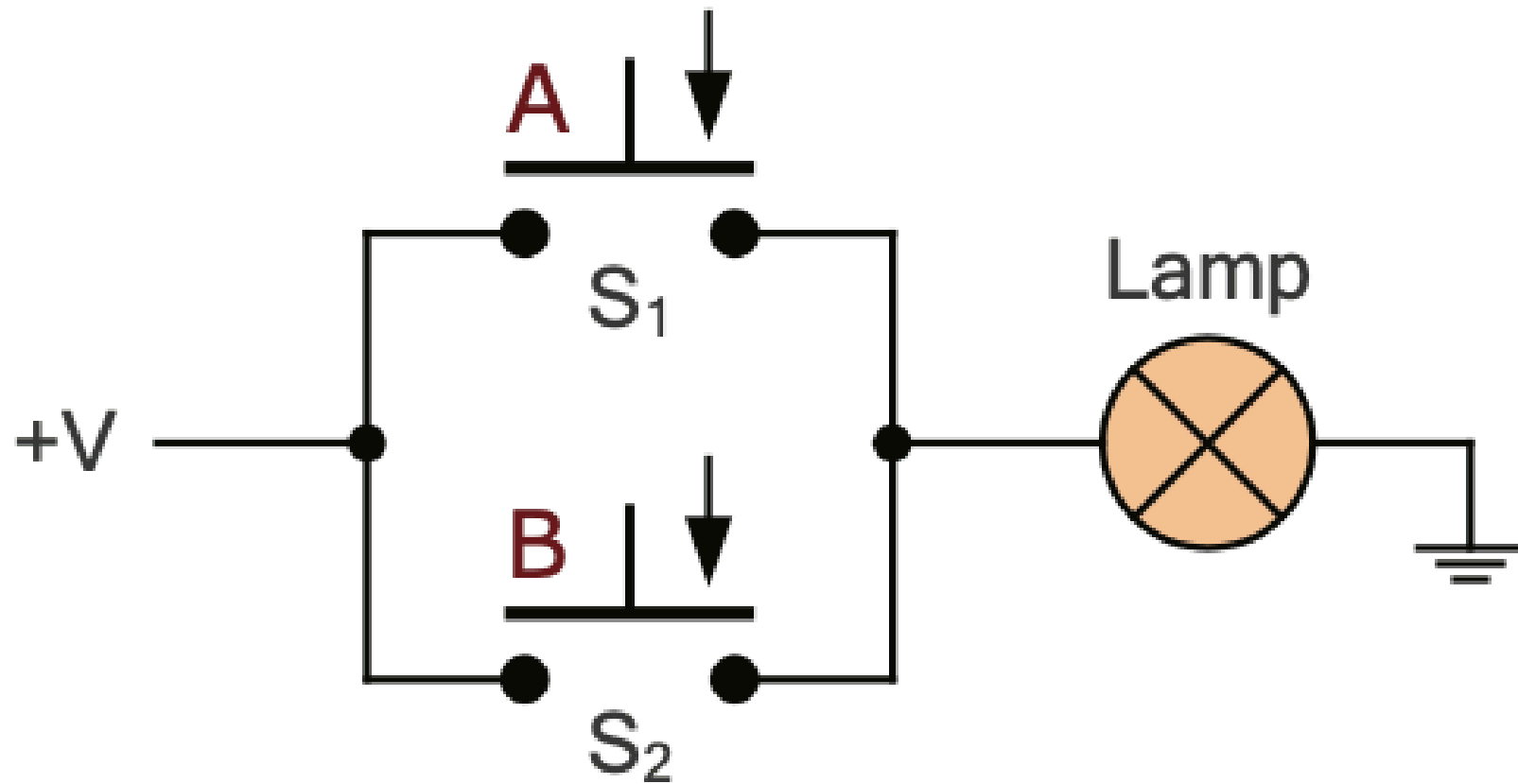
# Everything is a zero or a one

- 0 or 1
- pristine or hole, in a punch card
- gnd or VCC
- 0 or 5V
- ground or power
- open or closed

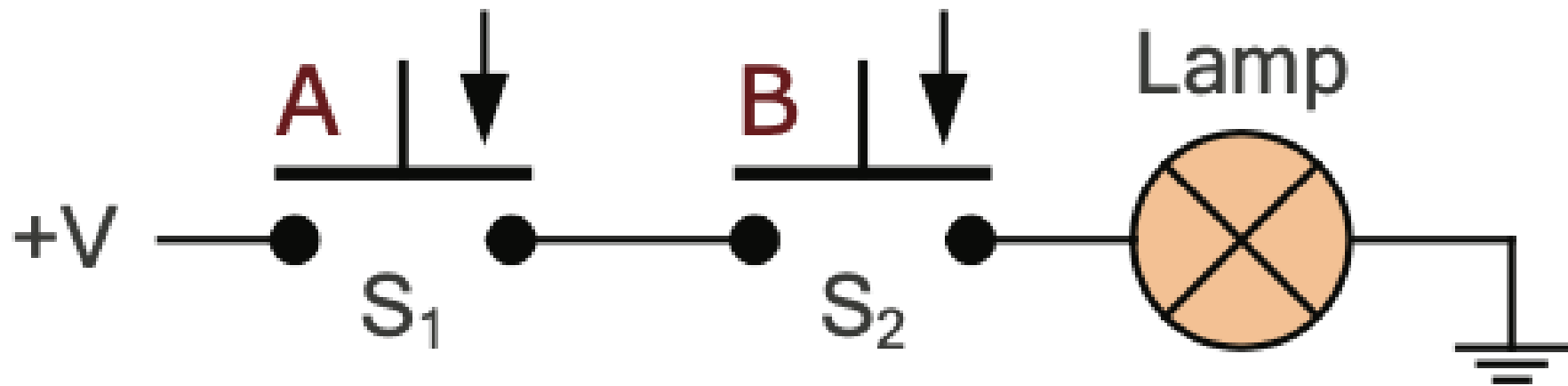
# How do we do this electronically?

- Transistors
- Lots of them...
- 100's of Billions in latest processors

## FIGURE 3. TWO SWITCHES IN PARALLEL



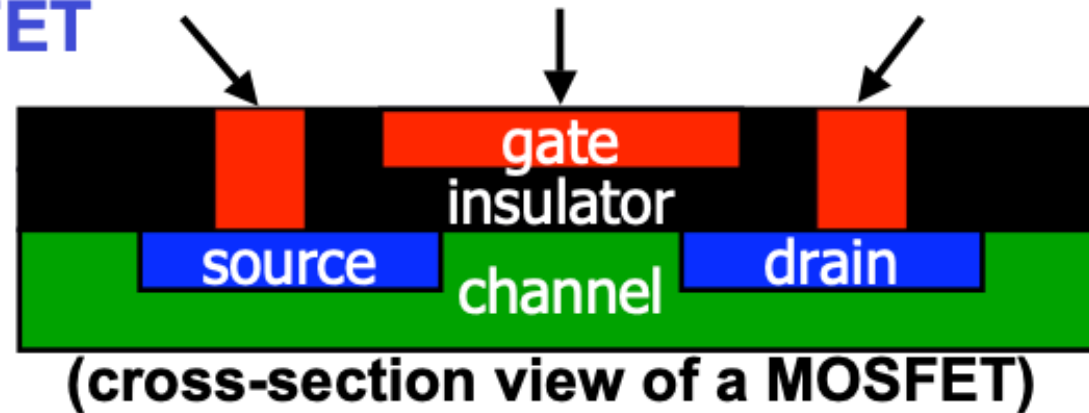
## FIGURE 2. TWO SWITCHES IN SERIES



## What if we had switches which were:

- really fast!
- really small!
- cheap!
- low power consumption!
- easy to manufacture and test (*relatively*)

## MOS + FET



**MOS:** three materials needed to make a transistor

- **Metal** (Al, W, Cu): conductor
- **Oxide** ( $\text{SiO}_2$ ): insulator
- **Semiconductor** (doped Si): conducts under certain conditions

**FET:** field effect (the mechanism) transistor

- Voltage on gate: current flows source to drain (transistor on)
- No voltage on gate: no current (transistor off)

**Recall, two types of MOSFET: n and p**

## Why electrons flow, or why current flows:

- applied gate voltage creates an **electric field** that attracts charge carriers (electrons or holes), or what we know as *current*
- forming a conductive path (channel) between *source* and *drain*
- The type of charge carriers depends on whether it's an **n-type** or **p-type** MOSFET
- **The gate voltage effectively acts like an electronic switch, controlling current flow through the channel**



# MOSFETs are ideal for digital electronics as they can be:

- very efficient switches
- consume almost no power in their off state
- provide good conductivity when on
- can be made extremely small\* (*measured in 10's nm*)
- can be produced and tested in high volumes

**\*25,400,000 nanometers (nm) in 1"**

**\*645,160,000,000,000 square nanometers (nm<sup>2</sup>) in 1 square inch**

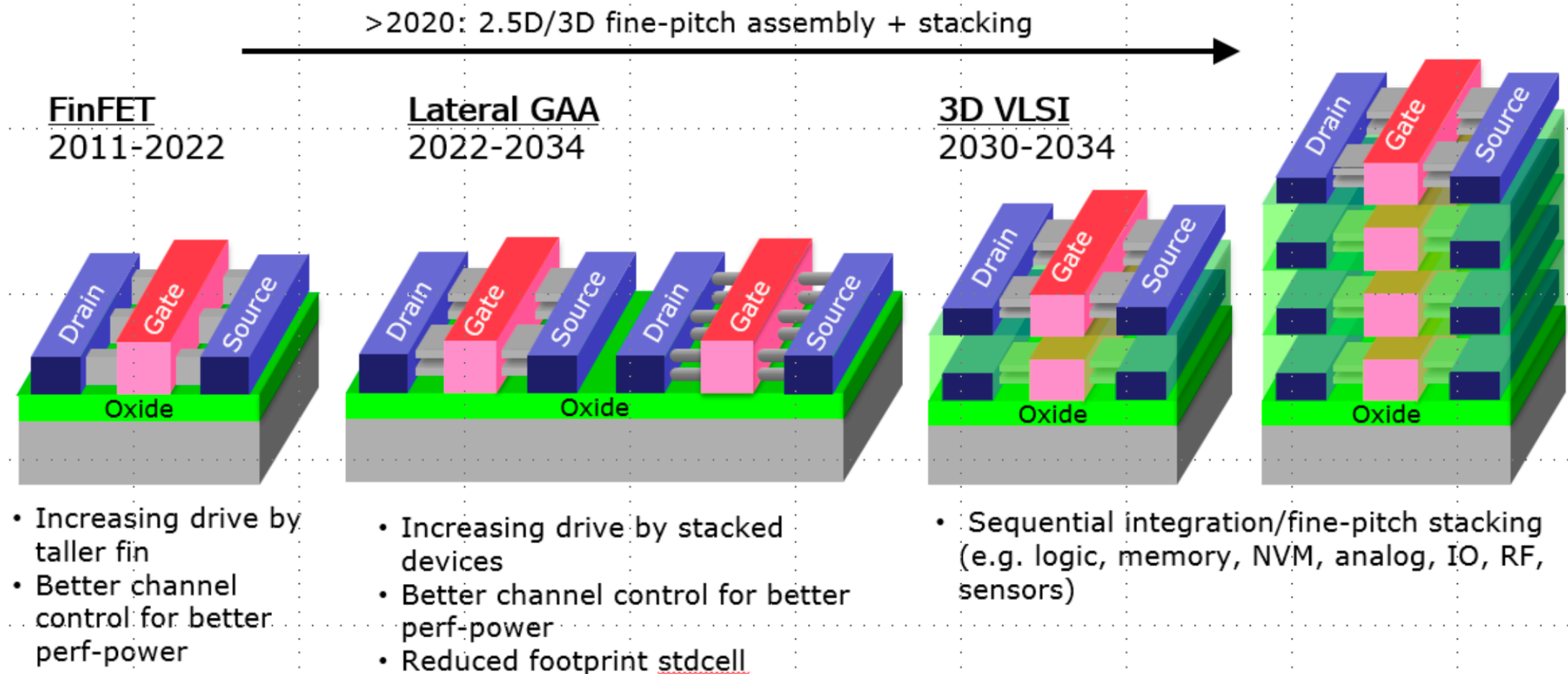
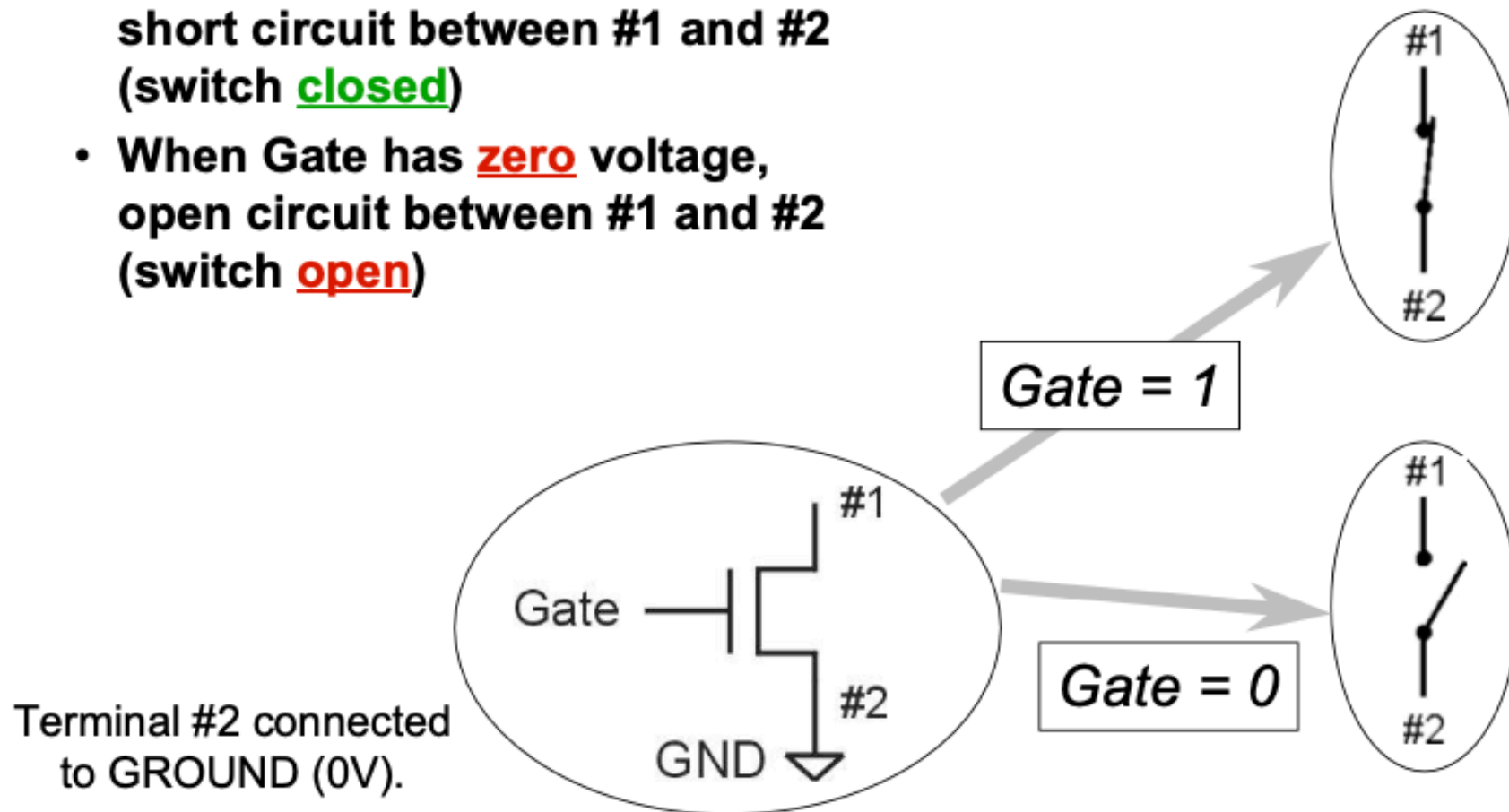


Figure MM-5

Evolution of device architectures in the IRDS More Moore roadmap

## N-type MOS Transistor

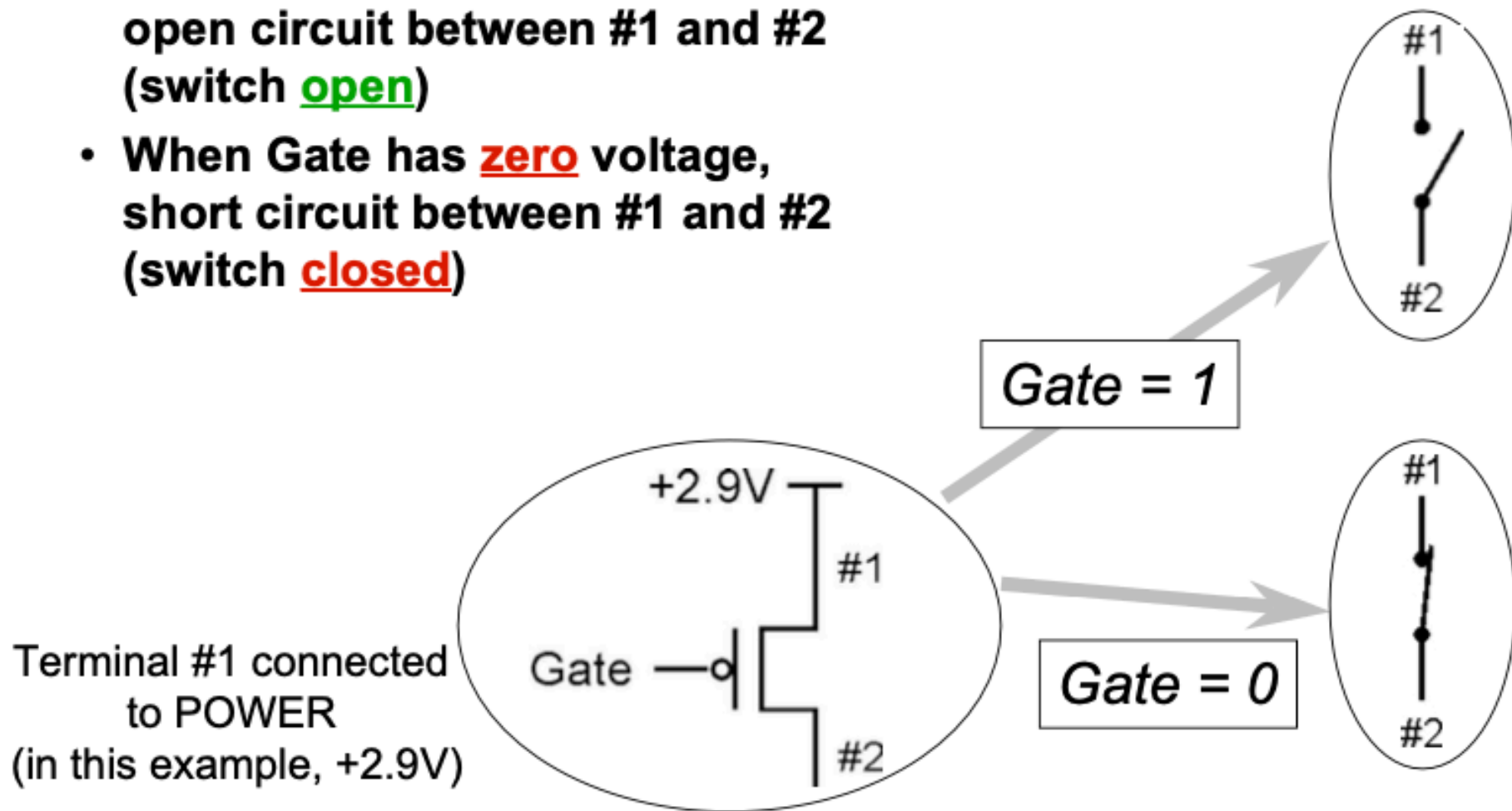
- When Gate has **positive** voltage, short circuit between #1 and #2 (switch **closed**)
- When Gate has **zero** voltage, open circuit between #1 and #2 (switch **open**)



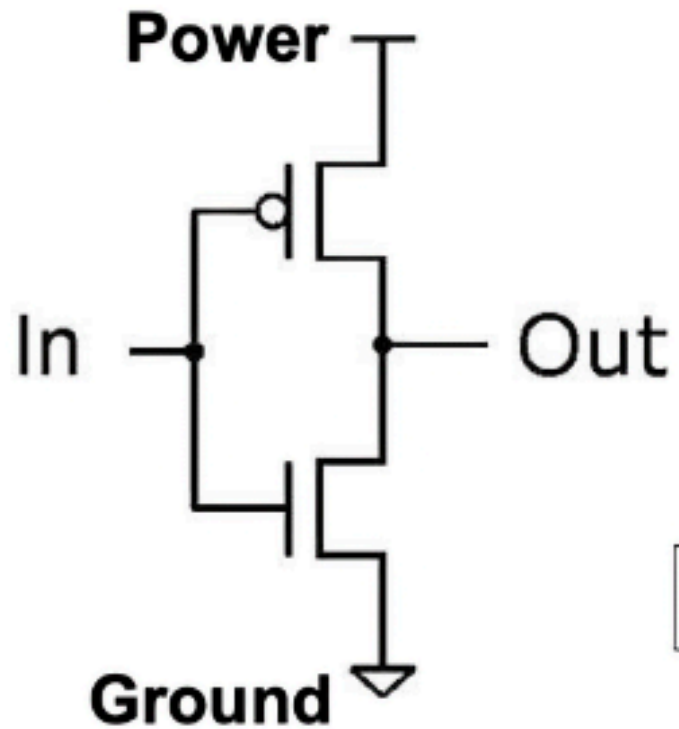
# P-type MOS Transistor

**P-type** is *complementary* to n-type

- When Gate has **positive** voltage, open circuit between #1 and #2 (switch **open**)
- When Gate has **zero** voltage, short circuit between #1 and #2 (switch **closed**)

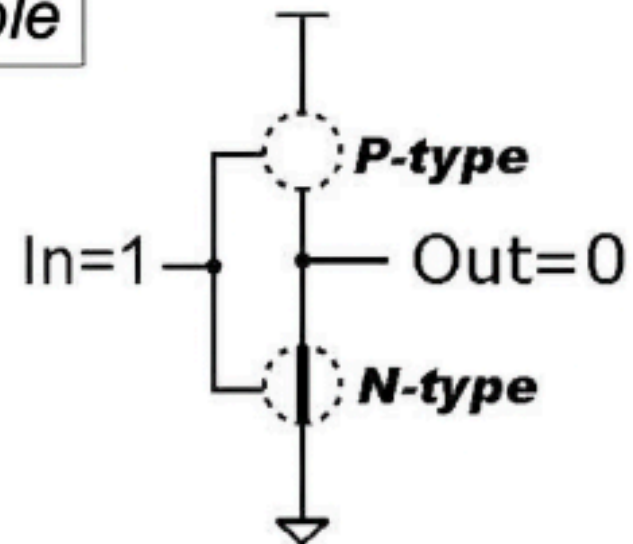
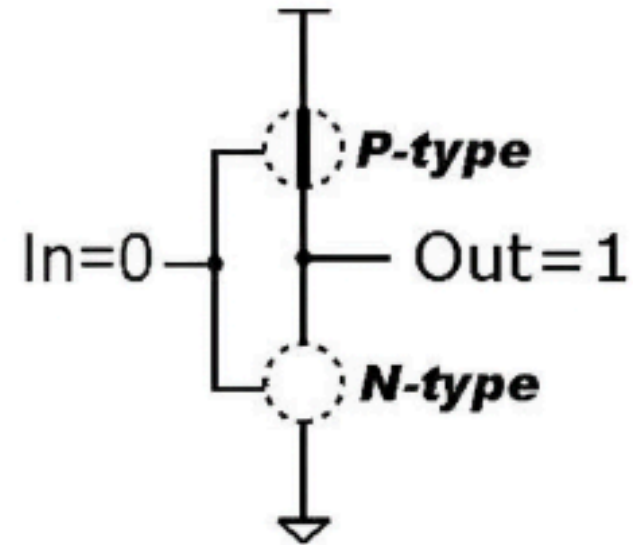


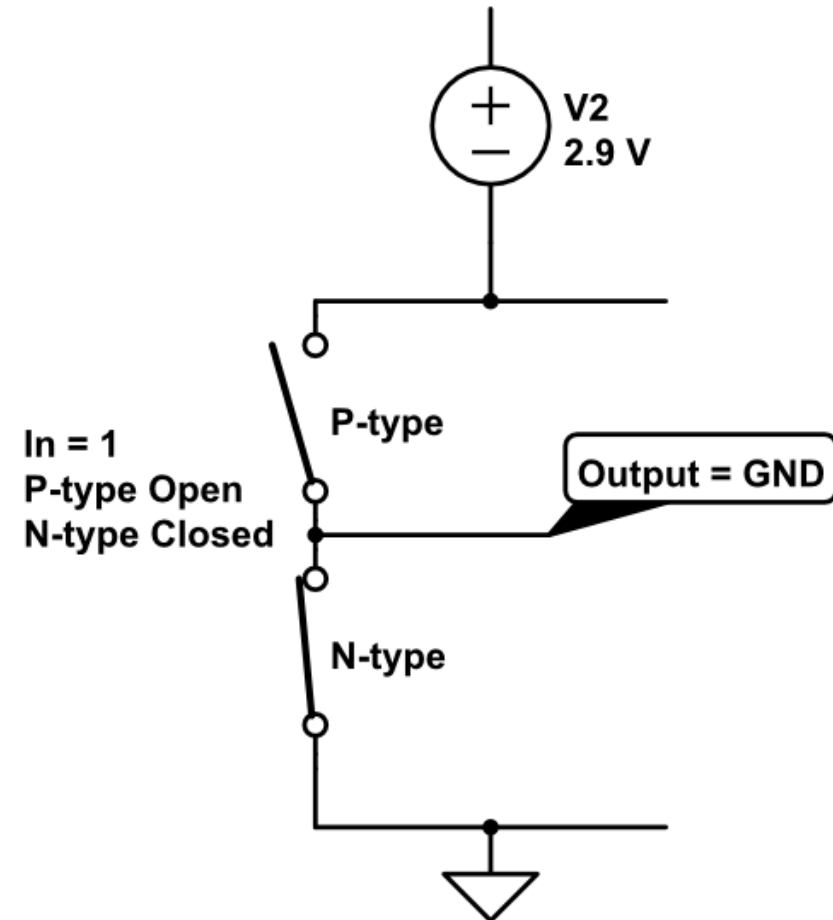
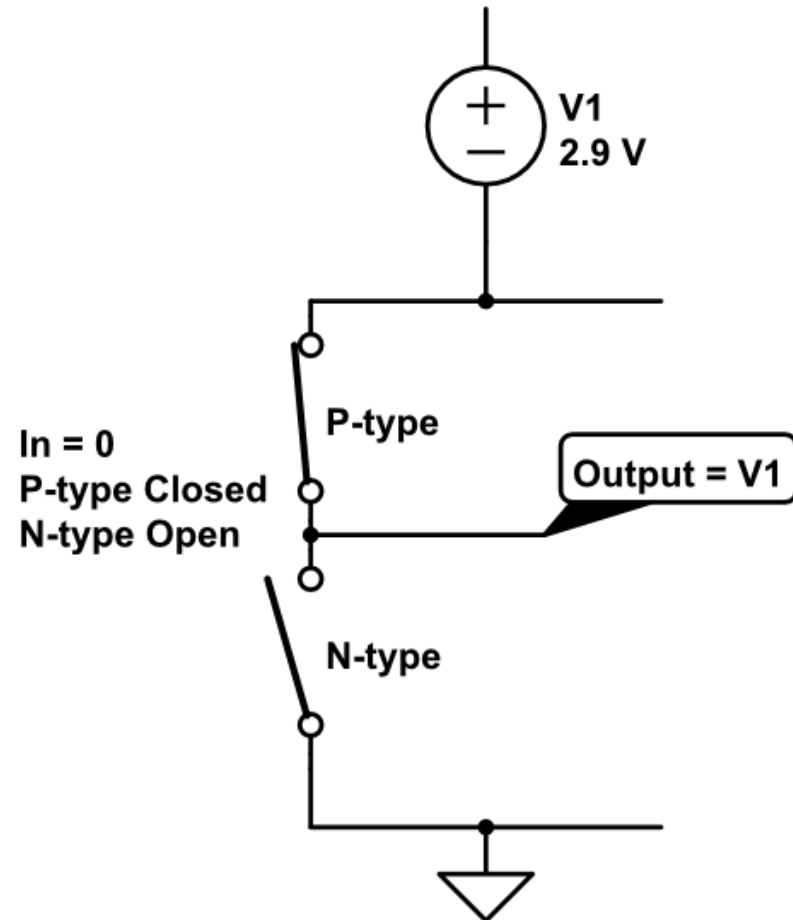
## Inverter (NOT Gate)



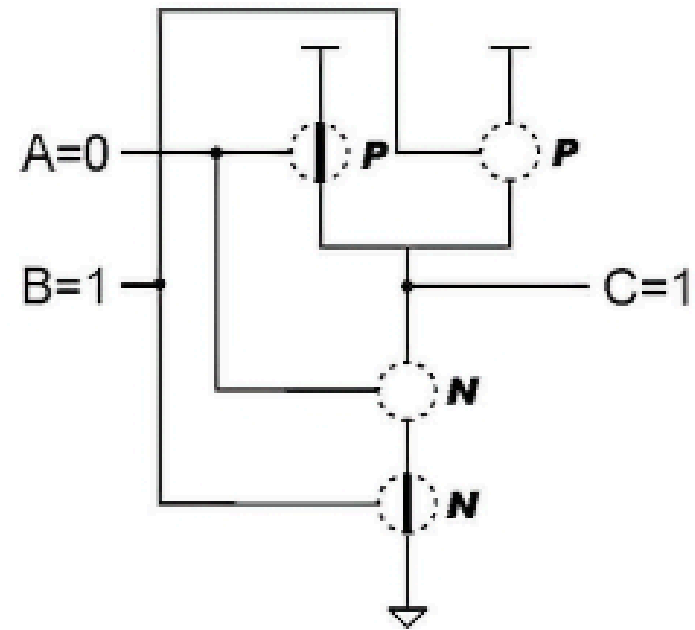
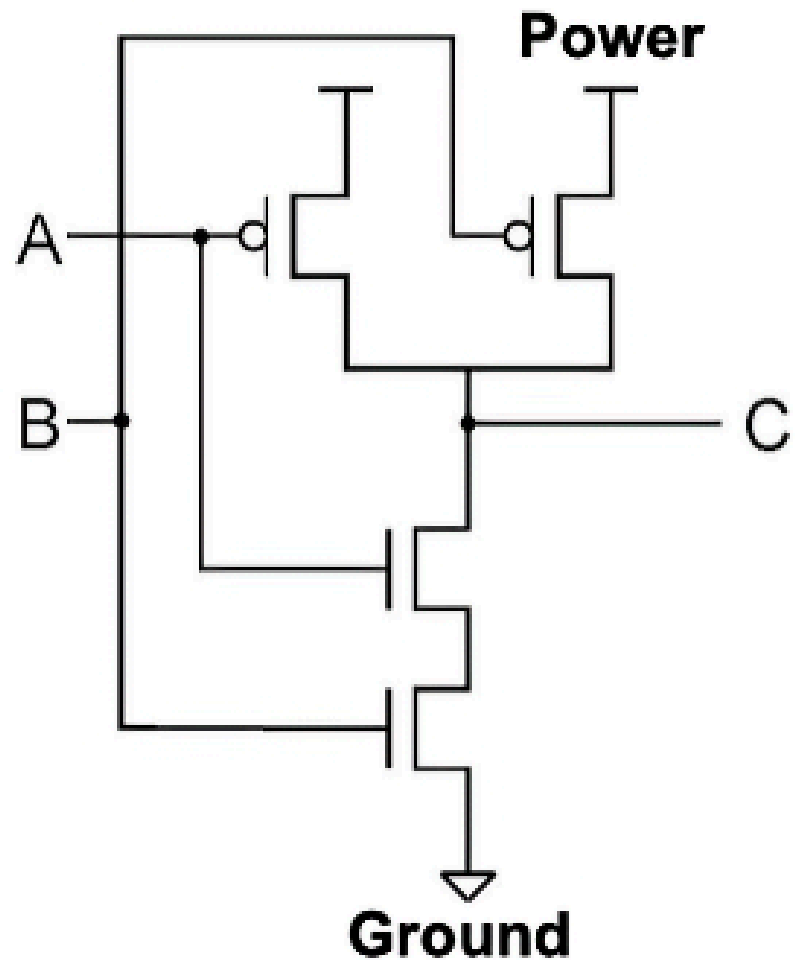
Truth table

In	Out
0	1
1	0



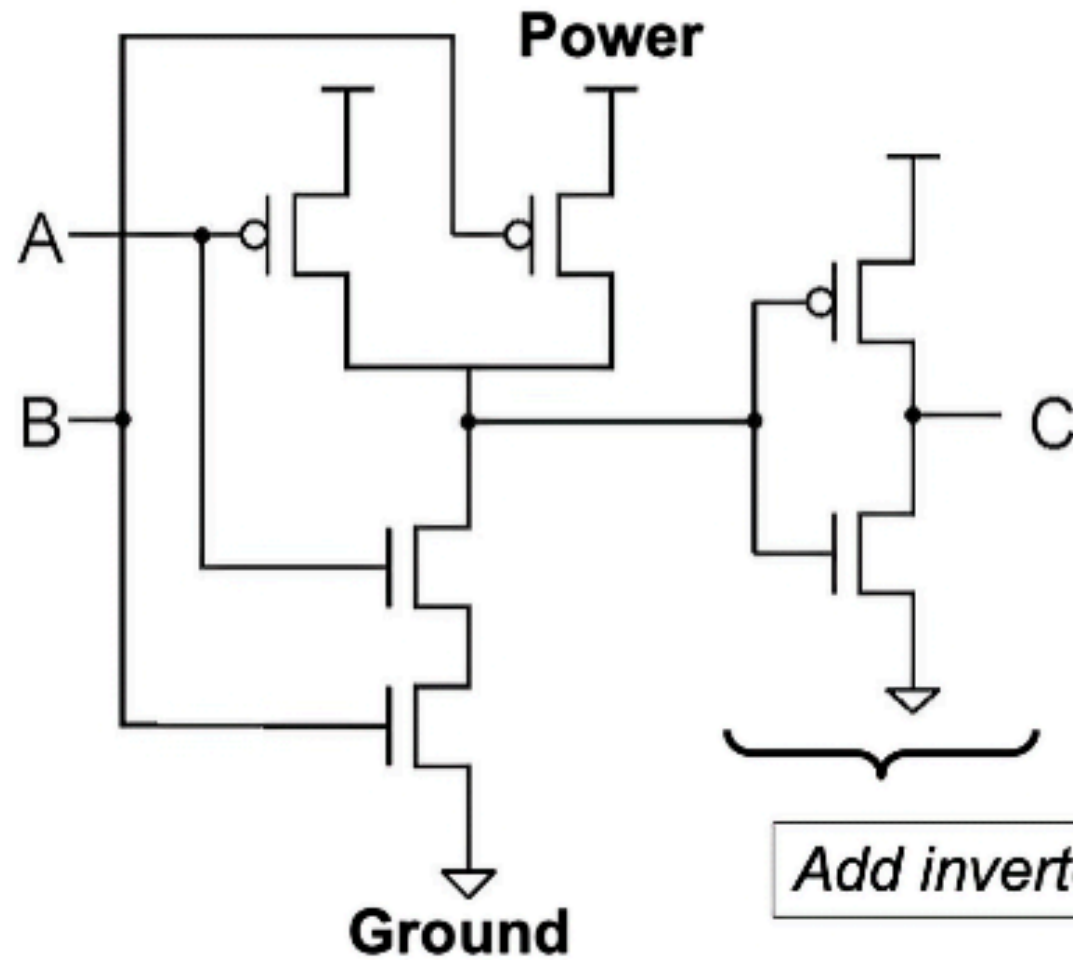


## NAND Gate (NOT-AND)



A	B	C
0	0	1
0	1	1
1	0	1
1	1	0

## AND Gate

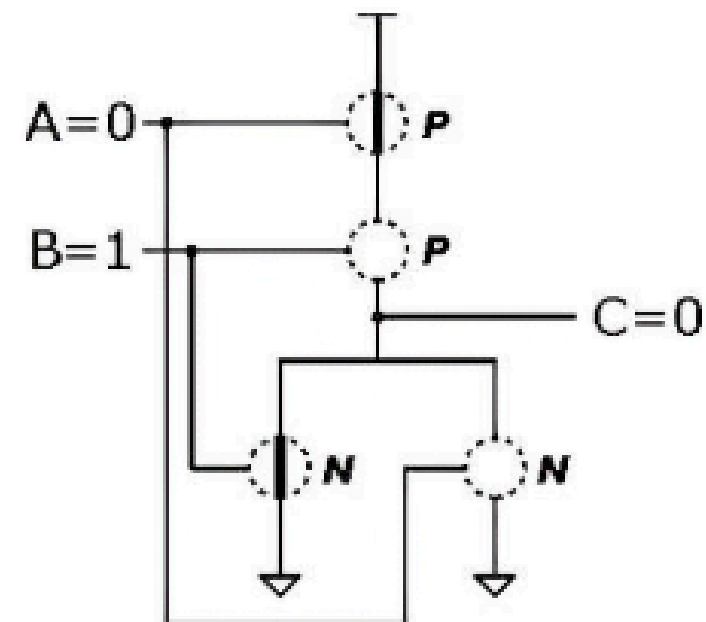
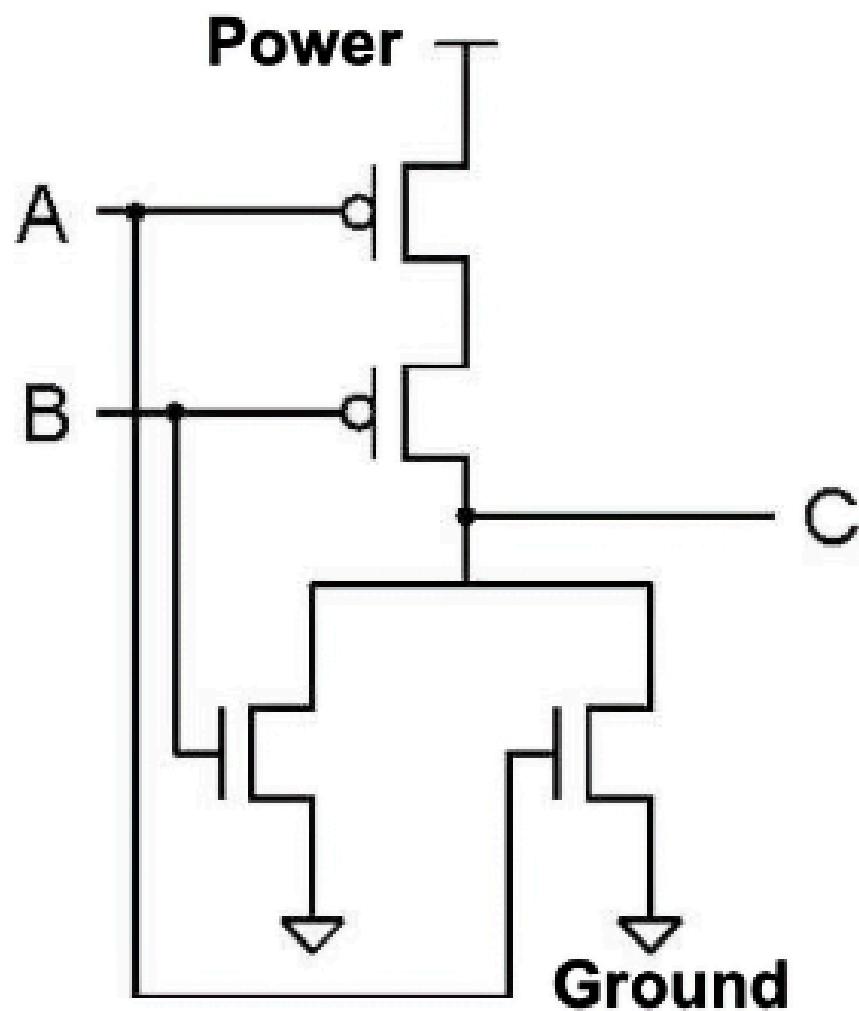


A	B	C
0	0	0
0	1	0
1	0	0
1	1	1

*Add inverter to NAND.*

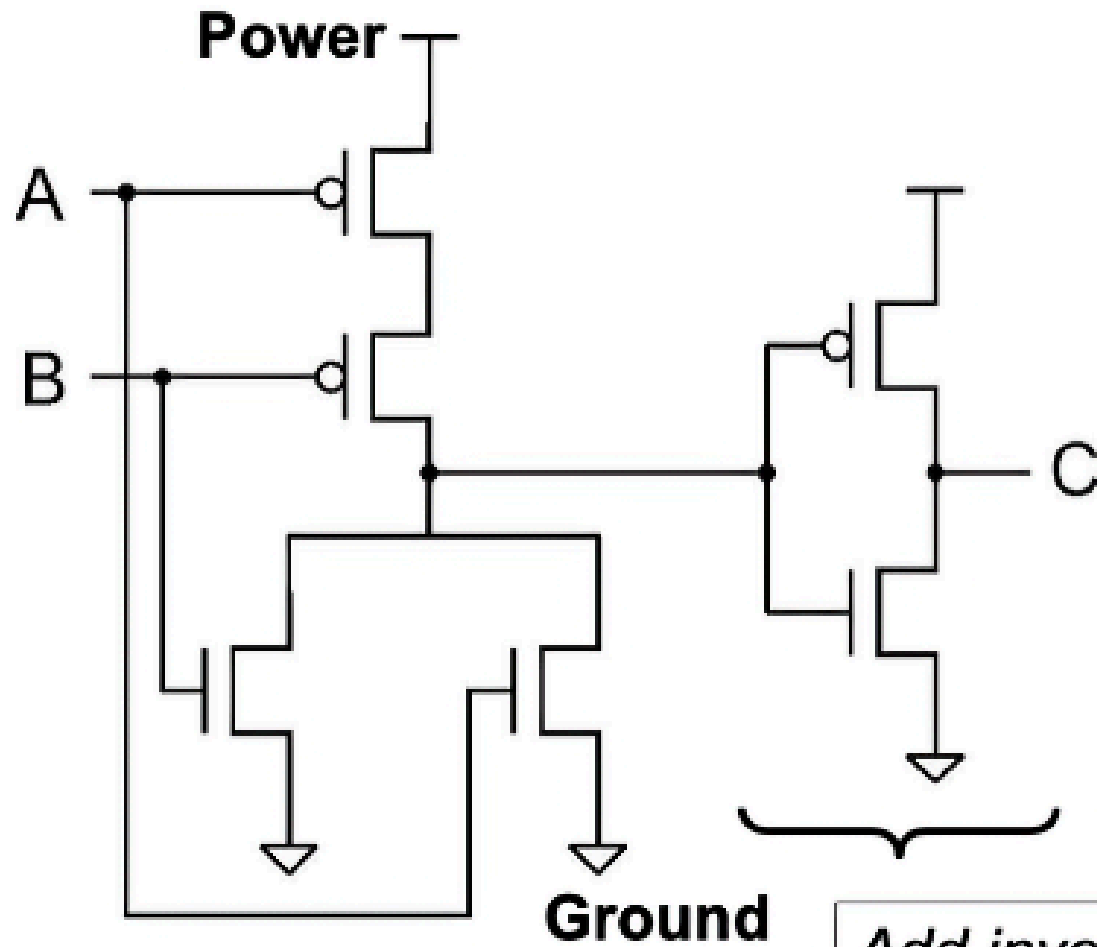


# NOR Gate (NOT-OR)



A	B	C
0	0	1
0	1	0
1	0	0
1	1	0

## OR Gate



A	B	C
0	0	0
0	1	1
1	0	1
1	1	1

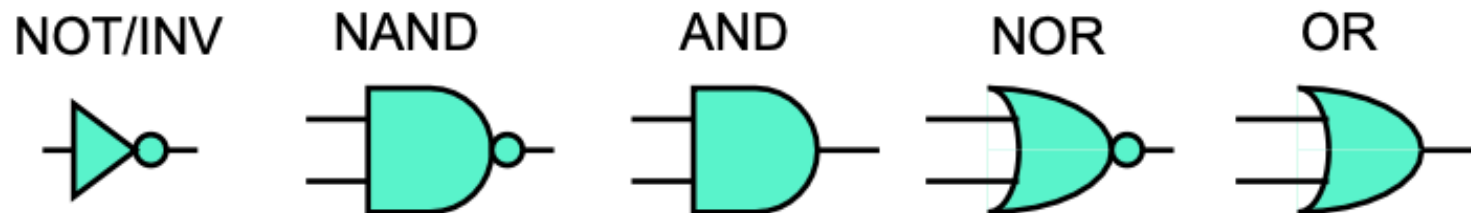
## Basic Gates

### From Now On... Gates

- Covered transistors mostly so that you know they exist
- Note: “Logic Gate” not related to “Gate” of transistors

### Will study implementation in terms of gates

- Circuits that implement Boolean functions

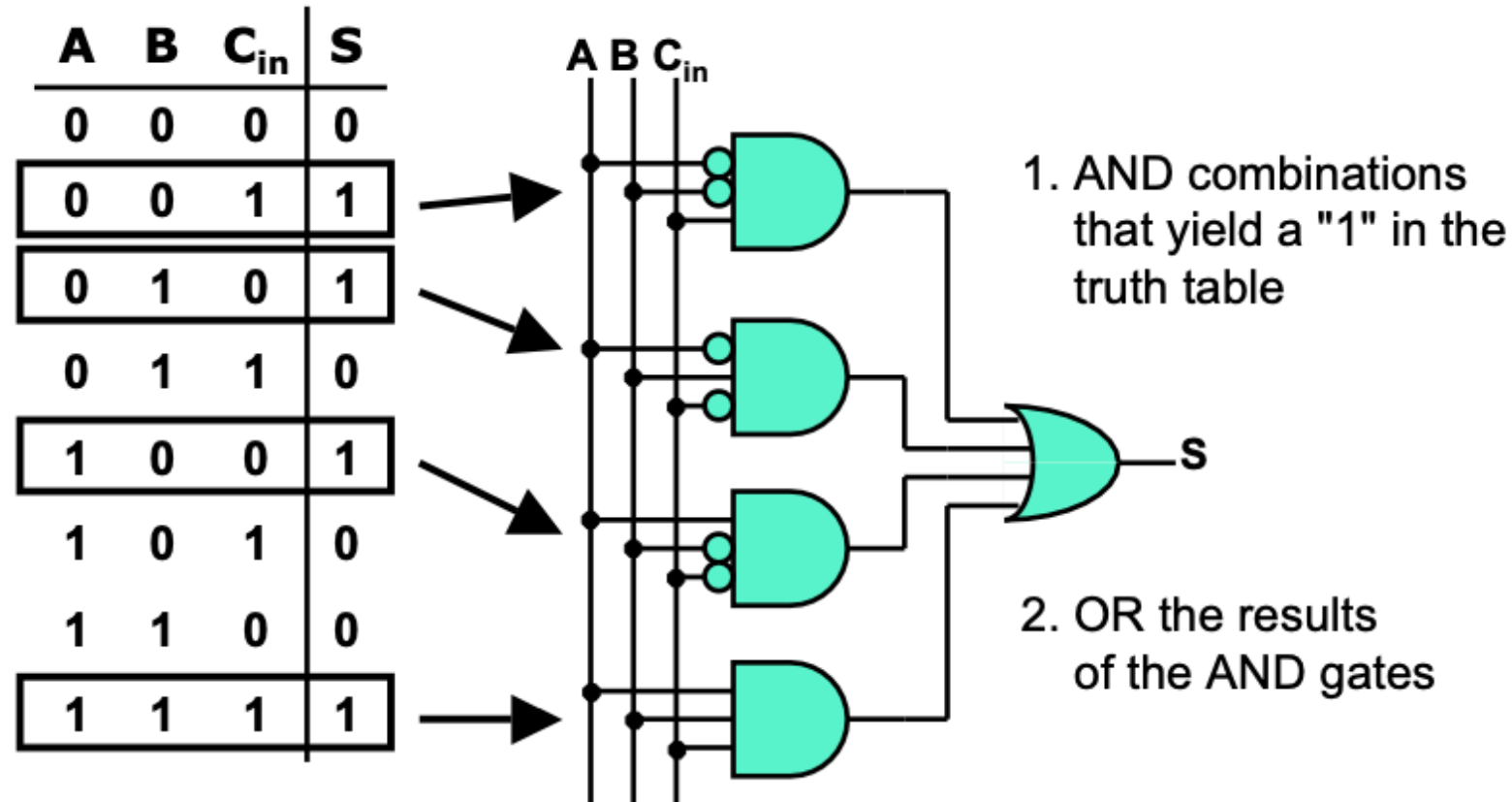


### More complicated gates from transistors possible

- XOR, Multiple-input AND-OR-Invert (AOI) gates

## Logical Completeness

AND, OR, NOT can implement ANY truth table



**Now What?**

**Before you go make a microprocessor, let's add the ability to [simplify](#)**

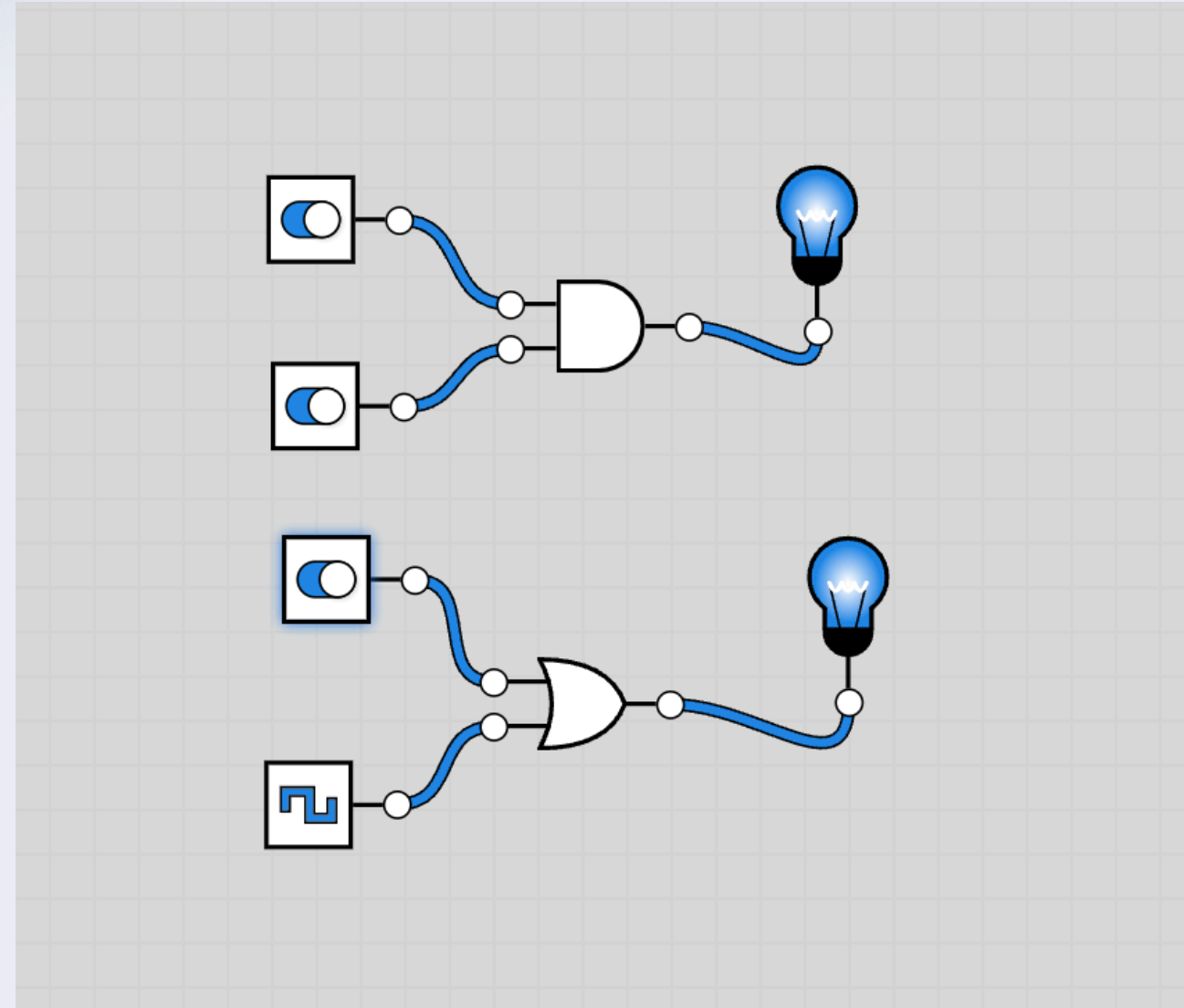
# Boolean Algebra

A series of laws which help you to understand how to *simplify* and *solve* boolean expressions.

# Digital Logic Simulator

Digital Logic Simulator\*

*\*(otherwise, known as don't take my word for it.)*



# Process

1. I'll describe the boolean algebraic law
2. You implement in the simulator
3. Once you believe, you have a successful implementation, screenshot and save to a homework folder



# Idempotent

A fundamental law in binary logic where an operation repeated on itself yields the same result. For example:

- $A \text{ AND } A = A$
- $A \text{ OR } A = A$

This means performing the same logical operation multiple times doesn't change the outcome.

# Annulment

Also known as the dominance law, this principle states that:

- $A \text{ AND } 0 = 0$
- $A \text{ OR } 1 = 1$

When combined with these special values, the result is predetermined regardless of the other operand.

# Identity

The identity law states that certain operations with specific values leave the operand unchanged:

- $A \text{ AND } 1 = A$
- $A \text{ OR } 0 = A$

These values (1 for AND, 0 for OR) act as neutral elements in the operation.

# Complement Law of Switches

This law deals with complementary operations:

- $A \text{ AND } (\text{NOT } A) = 0$
- $A \text{ OR } (\text{NOT } A) = 1$

When a value is combined with its complement using AND or OR, the result is always definite.

## Double Complement Law

States that negating a value twice returns the original value:

- $\text{NOT}(\text{NOT } A) = A$

This is similar to how two negatives make a positive in regular mathematics.

## Commutative Law

The order of operands doesn't affect the result:

- $A \text{ AND } B = B \text{ AND } A$
- $A \text{ OR } B = B \text{ OR } A$

You can swap the position of values and get the same outcome.

## Distributive Law

Similar to algebra, one operation distributes over another:

- $A \text{ AND } (B \text{ OR } C) = (A \text{ AND } B) \text{ OR } (A \text{ AND } C)$
- $A \text{ OR } (B \text{ AND } C) = (A \text{ OR } B) \text{ AND } (A \text{ OR } C)$

This allows for restructuring complex expressions.

# Absorptive Law

Describes how certain combinations of operations can be simplified:

- $A \text{ AND } (A \text{ OR } B) = A$
- $A \text{ OR } (A \text{ AND } B) = A$

The additional term gets "absorbed" into the expression.



## Associative Law

The grouping of operations doesn't affect the result:

- $(A \text{ AND } B) \text{ AND } C = A \text{ AND } (B \text{ AND } C)$
- $(A \text{ OR } B) \text{ OR } C = A \text{ OR } (B \text{ OR } C)$

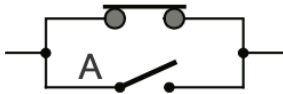
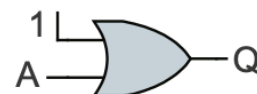
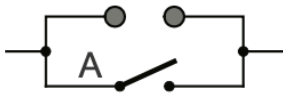
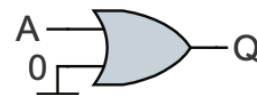
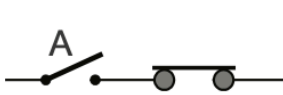
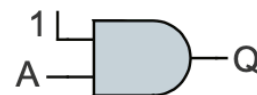
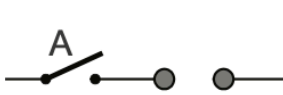
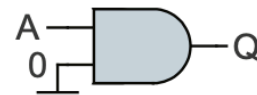


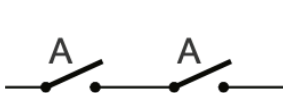
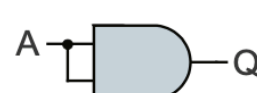
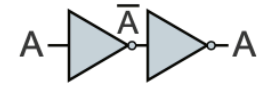

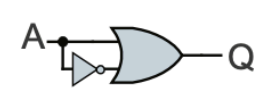
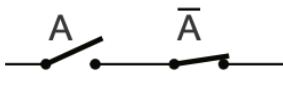
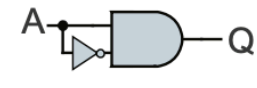
You can regroup terms without changing the outcome.

# De Morgan's Theorem

A fundamental principle for simplifying logical expressions:

- $\text{NOT}(A \text{ AND } B) = (\text{NOT } A) \text{ OR } (\text{NOT } B)$
- $\text{NOT}(A \text{ OR } B) = (\text{NOT } A) \text{ AND } (\text{NOT } B)$

This law shows how negation distributes across AND/OR operations.

Boolean Expression	Description	Equivalent Switching Circuit	Equivalent Logic gate	Boolean Algebra Law or Rule
$A + 1 = 1$	A in parallel with closed = CLOSED			Annulment
$A + 0 = A$	A in parallel with open = A			Identity
$A \cdot 1 = A$	A in series with closed = A			Identity
$A \cdot 0 = 0$	A in series with open = OPEN			Annulment
$A + A = A$	A in parallel with A = A			Idempotent
$A \cdot A = A$	A in series with A = A			Idempotent
$\text{NOT}(\text{not-}A) = A$	NOT NOT A = A (double negative)			Double Negation
$A + \bar{A} = 1$	A in parallel with NOT-A = CLOSED			Complement
$A \cdot \bar{A} = 0$	A in series with NOT-A = OPEN			Complement

## Credits:

Patt and Patel: "*Introduction to Computing Systems*"

AspenCore: "*Boolean Algebra*"

UPenn CSE240 Fall 2006 Fassell: "*Chapter 3 Lecture*"

IEEE 2021: "*More Moore*"