Introduction

Transaction Processing

- A very important component of any online information system / portal
- E-governance applications are in wide-spread use

Transaction

- A logical unit of work to be carried out
- on request by the end-user
  - transfer of specified amount of money from one account to another
  - making a reservation for a journey
  - issue a book of the library to a user ….

Transaction Processing

Input:

A number of requests for services from the end-users submitted concurrently from several input points

Action:

Carry out the requested services in a consistent manner

- no seat on a journey is reserved for more than one person
- amount debited from the party A is credited to party B

Measure:

Maintain a reasonably high throughput

- number of transactions completed per sec

What is a Transaction?

From the end-user point of view

- a logically sensible/complete piece of work

From the system point of view

- a sequence of database operations
  - reading data from tables on the disk,
    compute the updates,
    write back the data to the disk

Doing one transaction at a time, one-by-one:

- no scope of errors, but slow

Doing multiple transactions at a time

- Scope for error unless done carefully, good throughput

ACID Properties

A - Atomicity

C - Consistency

I - Isolation

D - Durability

Important properties to be satisfied by the overall system

Atomicity

- The given piece of work should be done entirely as one unit
  - Carrying out some portion of the work leads to inconsistent state of the database
    - For example, Rs.1000/-, to be transferred from A to B
    - Debited from A but not credited to B!
  - system encountered on error and stopped in the middle!
    - can not be allowed; leads to inconsistency

  If only a part of the work is done and error is encountered
  - ensure that effect of the partial work is not reflected in the database
- Responsibility of Recovery Module

Consistency

- Correctness Assumption:
  - A transaction takes the system from one consistent state to another consistent state, if executed in *isolation*
  - Responsibility of the application program developer
    - application programs or transaction programs should be carefully developed and thoroughly tested

- During the running of the transaction
  - It is possible that the DB is in an *inconsistent* state

    Ex: consider money transfer between two bank accounts

Isolation

- Let $T_1, T_2, \ldots, T_n$ be transactions submitted around the same time
- Isolation: Though the operations of T_i are interleaved with those of others, with respect to T_i,
  - any other T_j *appears* to have either completed before T_i or started after T_i finished
- The operations of T_j are completely isolated from those of T_i and hence have no effect on T_i
- Responsibility of Concurrency Control module

Durability

Upon successful completion of a transaction,

System must ensure that its effect is permanently recorded in the database

Also effect of failed transactions is not recorded

Failures

- Transaction program failures
  - internal errors – attempted division by zero etc.
  - transaction aborts
- System crashes
  - power failures etc.

Responsibility of Recovery Module

Note on Transaction Sequencing

- Suppose $T_1, T_2, \ldots, T_n$ are transactions submitted around the same time
  - The transactions may not *end* in the same order
- From the system point of view, guaranteeing atomicity and isolation is important
- The ending time of each transaction depends on
  - data items being accessed
  - computation time etc.
  - system's policy for guaranteeing atomicity/ isolation
- From the submitter's point of view
  - when the transaction finishes matters!

Concurrency Control

Interleaving the operations of transactions

- essential to achieve good throughput

However, arbitrary interleaving of operations

- leads to inconsistent database

Certain regulated interleaving so that

- two transactions involving the same data item don't conflict with each other

Concurrency control subsystem ensures this

Recovery Manager

System Crash

- due to various reasons
- some transactions might have run partially

Transaction failure

- transaction program error
- transaction was aborted by Concurrency Control Module due to violation of rules

Upon Recovery of system

- effect of partially-run transactions should not be there (atomicity)
- effect of completed transactions must be recorded (durability)

Recovery Manager and System Log

System Log:

- Details of running transactions are recorded in the log
- Important resource for recovering from crashes
- Always maintained on reliable secondary storage

Log Entries

- beginning of a transaction
- update operation details
  - old value, new value etc.
- ending of transaction
- more details later…

Transaction Operations

From the DB system point of view, only the read and write operation of a transaction are important

- Other operations happen in memory and don't affect the database on the disk

- Notation:
  - R(X) - transaction Reads item X
  - W(X) - transaction Writes item X

- Database items – denoted by X, Y, Z, …
  - Would be defined rigorously later

Transaction Operation - *Commit*

- A transaction issues a *commit* command when
  - it has successfully completed its sequence of operations
  - indicates that its effect *can be* permanently recorded in the db
- Once the system allows a transaction to commit
  - System is obliged to ensure that the effect of the transaction is permanently recorded in the database
- What exactly happens during *commit*
  - depends on the specific error recovery method used and the specific concurrency control method used
  - will become clear later on in the course

Transaction Operation - *Abort*

When a transaction issues *Abort*

- indicates that there is some internal error and
- transaction wants to terminate in the middle

System obligation/ responsibility

- ensure that the partial work done by the transaction has no effect on the database on the disk

What exactly happens during *Abort*

- depends on the specific error-recovery method and the specific concurrency control method adopted by the system

DB Model for Transaction Processing

- Database
  - Consists of several items – blocks / tuples
    - Granularity issues would be taken up later
- Transactions operate by exchanging data with DB only
  - They do not exchange messages between them
- Focus on read/write/commit/abort ops
  - Ignore in-memory ops
- Transactions are not nested

Transaction Operations

- $R_i(X)$ – read op of Txn *i,* reads db item X
  - disk block having X - copied to buffer page - if reqd
  - required value is assigned to program variable X
- $W_i(X)$ – write op of Txn *i,* writes db item X
  - disk block having X - copied to buffer page - if reqd
  - update buffer value using program variable X
  - transfer block to disk – immediately or later
- C_i – commit of Txn *i*
- A_i – abort of Txn *i*

Need for Concurrency Control

- If Txn's operations are interleaved arbitrarily
  - Several problems / anomalies arise
  - Can be classified as
    - WR anomalies
    - RW anomalies
    - WW anomalies

WR Anomalies or Dirty Reads

- Txn T_1 is in progress; updating values in a column
- Txn T_2 reads a value X updated by T_1, uses it to compute some other quantity, and finishes!
- T_1 for some reason changes X back to its original value!
  - T_2 has read *dirty* data or intermediate data

RW Anomalies or Unrepeatable Reads

- T_1 has read a value X and intends to read it again before changing it
- Between two reads of X by T_1,
  T_2 reads X, modifies it and finishes
- T_1 reads X the 2[nd] time and gets a different value!
  - *Unrepeatable read* problem
  - X could be the number of seats available for reservation

WW Anomalies or Lost Updates

- T_1 reads an item X, reduces it by 10% and wants to write it
- T_2 reads the same X before it was updated by T_1, increments it by 20% and wants to write it
- Say write of T_1 happens and then that of T_2
  - Final value of X is 1.2*X!
  - T_1's Update of X is lost – *lost update* problem

Schedules

- Sequence of interleaved operations of a Txn set
  - Ops of a Txn T appear in the same order as in T
- $R_1(X)\ R_2(Y)\ R_3(X)\ W_1(X)\ W_3(X)\ W_2(Y)$
- Serial: $R_1(X)\ W_1(X)\ R_2(Y) W_2(Y)\ R_3(X)\ W_3(X)$
  - No interleaving

Time

T1	T2	T3
R(X)		
	R(Y)	
		R(X)
W(X)		
		W(X)
	W(Y)	

Serializability

- Serial schedules
  - Do not cause any concurrency problems
  - But performance (throughput) is low
- Serializable Schedules
  - Interleaving of operations happens
  - But, in *some* sense *equivalent* to serial schedules
    - The effect of the interleaving is same as that of *some* serial schedule

Conflicting Pairs of Operations

- In a schedule, a pair of operations are said to be in *conflict* if
  - The ops belong to two different txns
  - The ops deal with the same DB item
  - One of the ops is a Write operation

    For example,

    1) R(X) of T1 conflicts with W(X) of T3
    2) R(X) of T3 conflicts with W(X) of T1
    3) W(X) of T1 conflicts with W(X) of T3

T1	T2	T3
R(X)		
	R(Y)	
		R(X)
W(X)		
		W(X)
	W(Y)	

Conflict Equivalence

A schedule S_1 is said to *conflict-equivalent* to S_2 if

The relative order of *any* conflicting pair of operations is the same in both S_1 and S_2

$R_1(X)\ R_2(Y)\ \underline{R_3(X)}\ W_1(X)\ \underline{W_3(X)}\ W_2(Y)$

is conflict equivalent to

$R_1(X)\ R_2(Y)\ \underline{R_3(X)}\ W_2(Y)\ W_1(X)\ \underline{W_3(X)}$

But is not conflict-equivalent to

$R_1(X)\ R_2(Y)\ W_1(X)\ \underline{R_3(X)}\ \underline{W_3(X)}\ W_2(Y)$

T1	**T2**	**T3**
R(X)		
	R(Y)	
		R(X)
W(X)		
		W(X)
	W(Y)	

Conflict Serializable Schedules

- A schedule is called *conflict serializable* if it is conflict equivalent to *some* serial schedule
- Or, if we can move non-conflicting ops such that
  - The relative order of ops of a Txn is intact and
  - Schedule becomes a serial
  - Schedule in picture
    - is *not* conflict-serializable
    - T1 < T3 disturbs the 2nd pair and
    - T3 < T1 disturbs the 1st and 3rd pairs

T1	T2	T3
R(X)		
	R(Y)	
		R(X)
W(X)		
		W(X)
	W(Y)	

Conflict Serializable Schedules

Suppose we make a slight change to our example:

S1

T1	T2	T3
R(X)		
	R(Y)	
		R(X)
W(X)		
		W(X)
	W(Y)	

S2

T1	T2	T3
R(X)		
	R(Y)	
W(X)		
		R(X)
		W(X)
	W(Y)	

S2 is conflict-equivalent to serial schedules

T1,T3,T2 ; T1,T2,T3 and T2,T1,T3

Non-conflicting pairs of ops -- moved to get these serial schedules

Precedence Graph of a Schedule

- Precedence graph or serialization graph
- Nodes represent transactions
- A directed arc exists from T_i to T_j if
  - An operation of T_i *precedes* an operation of T_j and *conflicts* with it
- A schedule S is conflict-serializable if and only if the precedence graph of S is *acyclic*
  - Topological sorts of the graph give equivalent serial schedules

Precedence Graph - Examples

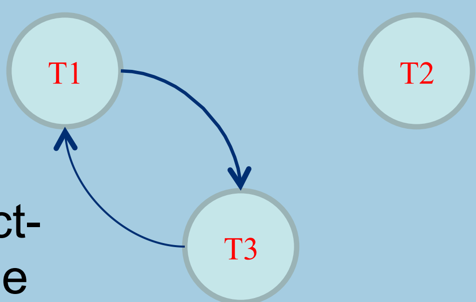
Precedence graph (or serialization graph) examples

S1

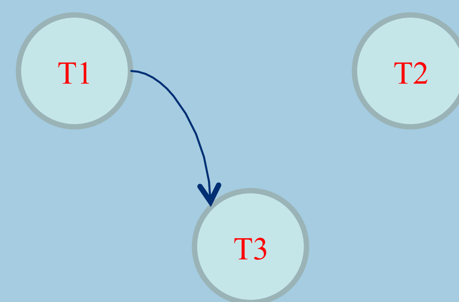
T1	T2	T3
R(X)		
	R(Y)	
		R(X)
W(X)		
		W(X)
	W(Y)	

S2

T1	T2	T3
R(X)		
	R(Y)	
W(X)		
		R(X)
		W(X)
	W(Y)	



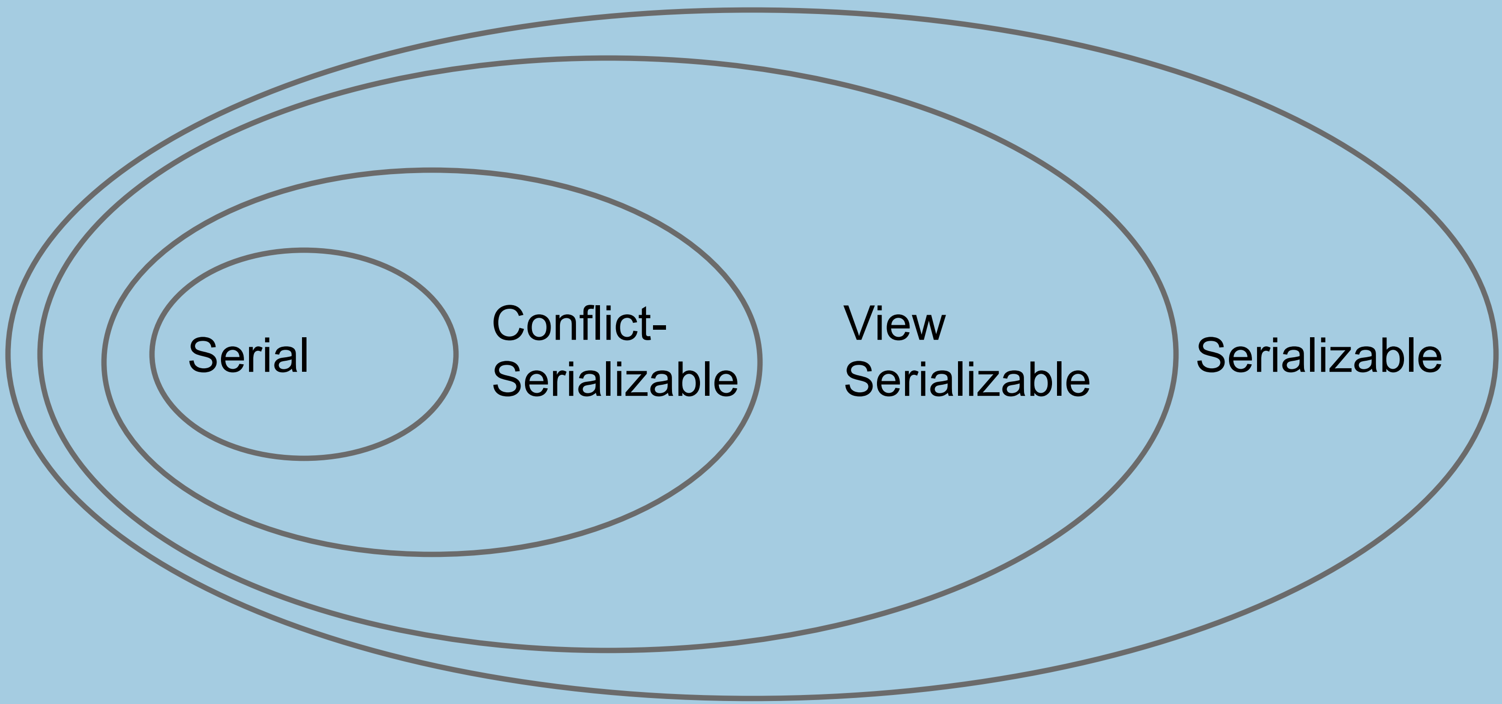
Not conflict-serializable



Conflict-serializable

Conflict Serializability and Serializability

- Conflict serializability is a sufficient condition for serializability but is not necessary!



View Serilizability

- A schedule S_1 is view-equivalent to S_2 if any
  - T_i read the initial value of a db item X in S_1,
    it does so in S_2 also.
  - T_i read a db item X written by T_j in S_1,
    it does so in S_2 also.

  and,
  - For each db item X, the txn that wrote the final value of X is same in S_1 and in S_2.
- A schedule is view-serializable
  if it is view-equivalent to some serial schedule

Concurrency Control Using Locks

- Assumptions
  - A transaction requests for a lock on a db item X before doing either *read* or *write* on X
  - A transaction unlocks X after it is done with X
- Locks - binary locks are assumed
  - Ensure mutual exclusion
  - At any time, at most one transaction holds a lock on a db item
  - Locking scheduler – ensures above
    - Keeps track of who holds lock on what item

Locking and Serializability

T1	T2	T3
L(X),R(X),U(X)		
	L(Y),R(Y),U(Y)	
		L(X),R(X),U(X)
L(X),W(X),U(X)		
		L(X),W(X),U(X)
	L(Y),W(Y),U(Y)	

- Locking alone
  - Does not guarantee serializability
  - Above schedule continues to be non-serializable

Two Phase Locking (2PL)

- 2PL Protocol
  - All *lock* requests of a transaction precede the first *unlock* request
  - Or a transaction has a *locking* phase followed by an *unlocking* phase
- If all transactions follow 2PL protocol
  - The *resulting* schedules will be conflict-serializable
- A very important and valuable result!

Why 2PL works?

- S : A schedule of *n* transactions that follow 2PL
- Let T_i be the transaction
  - that issues the first *unlock* request among all txns
- We will argue that
  - All ops of T_i can be brought to the beginning of S without passing over any conflicting ops
- We get S_1: (ops of T_i); (ops of other *n*-1 txns)
  - S_1 is conflict-equivalent to S
- Thus we can show that S is conflict-serializable

Why 2PL works?

- Suppose some op of T_i, say $W_i(X)$, is conflicting with a preceding op, say $W_j(X)$, of T_j in S

  $\ldots W_j(X), \ldots, W_i(X), \ldots$ or we have,

  $\ldots W_j(X), \ldots, U_j(X), \ldots, L_i(X), \ldots, W_i(X), \ldots$
- As T_i is the *first* txn to issue an unlock, say, $U_i(Y)$,
  - $U_i(Y)$ precedes $U_j(X)$ in S
- So, S could be

  $\ldots W_j(X), \ldots, U_i(Y), \ldots, U_j(X), \ldots, L_i(X), \ldots, W_i(X), \ldots$
- Then, T_i is not following 2PL – a contradiction

Why 2PL works?

- The argument is same for any conflicting pairs of operations involving an operation of T_i
- Thus, for all ops of T_i, there are no conflicting ops that precede it.
- So, beginning with the first op of T_i,
  - We can swap ops of T_i with the previous ops and bring them all the way to the front of schedule S
- This can be repeated among the remaining n-1 txs
- S is conflict-equivalent to a serial schedule

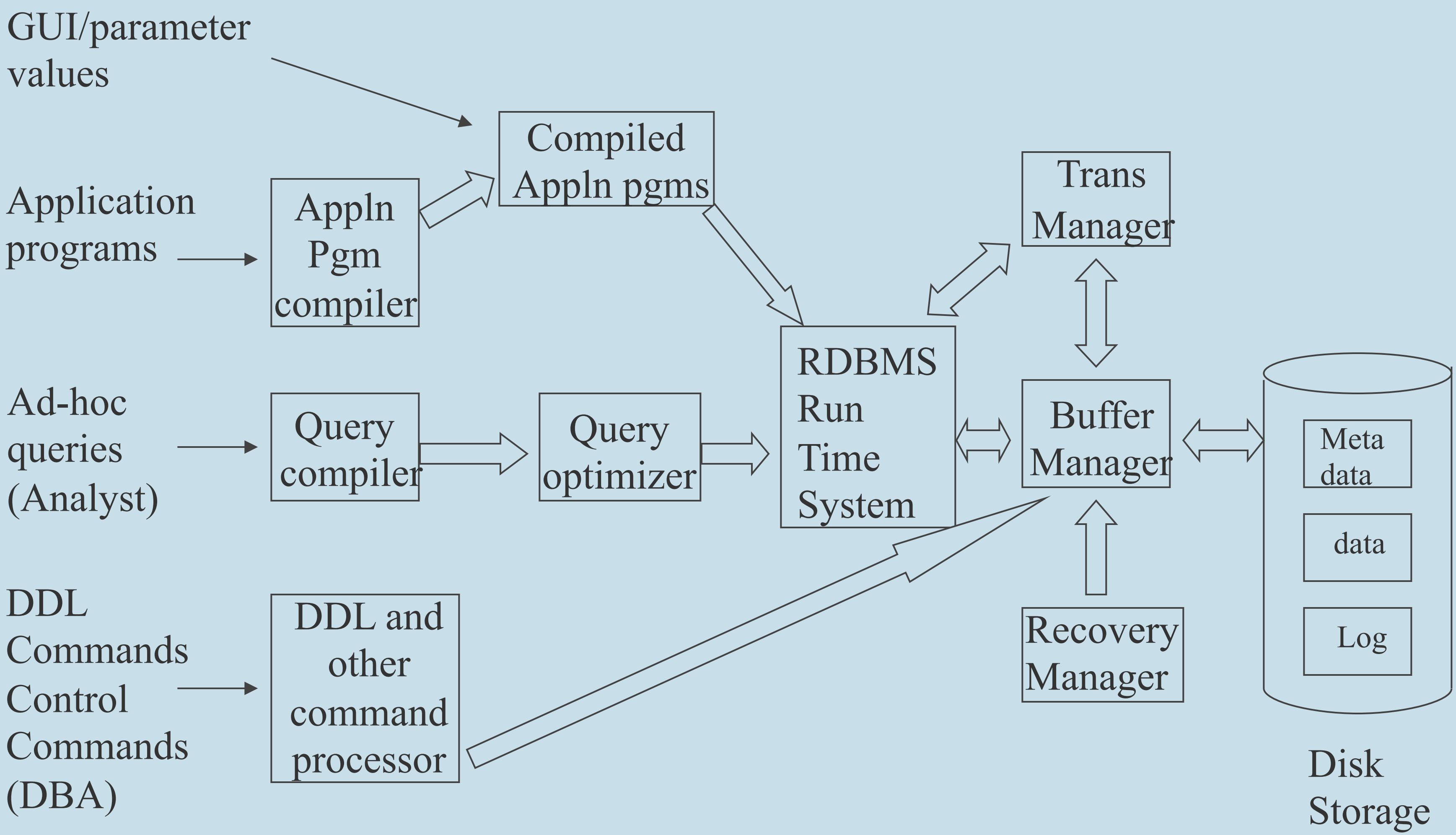
Possibility of Deadlocks

- Use of basic 2PL
  - Deadlocks may occur; Deadlock detection and resolution is adopted
- To detect
  - A graph called *wait-for* graph is maintained
  - Each running Txn – a node
  - If T_i is waiting to lock an item held by T_j
    - Directed edge from T_i to T_j
- To resolve
  - select Txn in cycles & abort/resubmit them

System Failures and Recovery

- Failures
  - Transaction errors – abrupt ending w/o completing
  - CC module may decide to abort a Txn
  - Disk crashes/power failures
- Log – the principal tool for error recovery
  - Sequential file
    - log entries reach the log on disk in same order as they are written (an important assumption)
  - Undo Logs / Redo Logs / Undo-Redo logs
  - Each with a corresponding recovery method

Architecture of an RDBMS system



Buffer Manager

- DB item – a disk block
- Disk blocks – brought into Memory Buffers
  - Modified by running transactions
  - Written to disk when transaction completes
- We will use more detailed Txn operations
  - Input(A):   get disk block with A into buffer
  - Read(A, t): t := A; do Input(A) if reqd;   t is local var
  - Write(A, t): A:= t;   do Input(A) if reqd
  - Output(A): send block with A from buffer to disk

Log Record Entries

- <Start T> -- Txn T has begun
- <End T> -- Txn T has ended
- <Commit T>
  - T has successfully completed its work
  - Changes made by T must appear in the on-disk DB
- <Abort T>
  - T has not successfully completed its work
  - Changes done by T shouldn’t be there in the on-disk DB
- Update records – specific to the logging method
- Flush Log – Force-write log entries to the disk

Undo Logging

Update Type Log Record

<T, X, V> : Txn T has changed the item X and its ***old*** value is V

Undo Logging Rules:

U_1: If a transaction T modifies db item X, then the log record <T, X, V> must be written to disk ***before*** the new value of X is written to disk.

U_2 : If a transaction T commits, then its **COMMIT** log record must be written to disk only ***after*** all db items changed by T have been written to disk, but as soon thereafter as possible.

Undo Logging

When a Txn T Commits,

DB items flow to disk as below:

for each modified DB item X

{ send update entry <T,X,V> to disk

write item X to disk }

write <Commit T> to disk

Undo Logging Example

Txn T is doing money transfer of Rs 100/- from account A to account B
Consistency Requirement: A+B is same before and after the Txn

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10	Output(B)	1600	400	1600	400	1600	
11							<commit,T>
12	Flush Log						

Recovery Using Undo Log

Examine Log and Partition Txns into:

Committed Set : Txns for which <Commit T> exists

Incomplete Set : Txns for which <Abort T> exists or <Commit T> does not exist

Examine Log in the *reverse* direction

For every update record <T,X,V>

T is Committed: ***Do nothing***; All is well due to U_2 !

T is Incomplete: Restore value of X on disk as V

-T might have changed some items on disk

-But log entries with old values are on disk due to U_1

Do <Abort T> log entry for each incomplete T & Flush Log

Undo Logging: Crash Recovery

Crash occurs sometime before the first Flush Log:
DB on disk – unchanged; T - recognized as incomplete;
Log entries of T - unsure if they are on disk; if present - used for “undo”;
no harm! <Abort T> entered; T – resubmitted;

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10	Output(B)	1600	400	1600	400	1600	
11							<commit,T>
12	Flush Log						

Undo Logging: Crash Recovery

Crash occurs sometime after the first Flush-Log but before Step 11:
DB on disk - might have changed; T - recognized as incomplete;
Log entries of T - on disk; used for undoing T;
<Abort T> entered; T resubmitted;

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10	Output(B)	1600	400	1600	400	1600	
11							<commit,T>
12	Flush Log						

Undo Logging: Crash Recovery

Crash occurs after Step 11, before Step 12: DB on disk - changed; <Commit T> - on disk: T - recognized as completed; No action reqd; <Commit T> - not on disk: T - recognized as incomplete; Log entries of T used for undoing T; <abort T> entered; T resubmitted;

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10	Output(B)	1600	400	1600	400	1600	
11							<commit,T>
12	Flush Log						

Undo Logging: Crash Recovery

Crash occurs after Step 12:
DB on disk - changed;
<Commit T> - on disk: T - recognized as completed; No action reqd;

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10	Output(B)	1600	400	1600	400	1600	
11							<commit,T>
12	Flush Log						

Redo Logging

Undo Logging:

Cancels the effect of incomplete Txns and

Ignores the committed Txns

All DB items to be sent to disk before Txn can commit

Results in lot of I/O

Redo Logging:

Ignores incomplete Txns and

Repeats the work of Committed Txns

Update log entry <T, X, V>: V is the ***new*** value

Redo Logging

Redo Logging Rule:

Before modifying any database element X on disk, it is necessary that all log records pertaining to this modification of X, including both the update record <T, X, v> and the <COMMIT T> record, ***must*** appear on disk.

Also called the ***write-ahead logging*** (WAL) rule

Items flow like this: Update Log records,
the commit entry
and then the changed DB items.

Redo Logging Example

Txn T is doing money transfer of Rs 100/- from account A to account B
Consistency Requirement: A+B is same before and after the Txn

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,400>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1600>
8							<commit,T>
9	Flush Log						
10	Output(A)	1600	400	1600	400	1500	
11	Output(B)	1600	400	1600	400	1600	

Recovery Using Redo Log

Examine Log and Partition Txns into:

Committed Set : Txns for which <Commit T> exists

Incomplete Set : Txns for which <Abort T> exists or <Commit T> does not exist

Examine Log in the *forward* direction

For every update record <T,X,V>

T is Incomplete: ***Do nothing***; DB on disk has no effects!

T is Committed: Unsure if all the effects of T are on disk

-But log entries with new values are on disk (WAL)

-Redo the change as per the log entry

Do <Abort T> log entry for each incomplete T & Flush Log

Redo Logging: Crash Recovery

Crash Occurs Before Step 8: <Commit T> not made by T;
T is recognized as incomplete; No action reqd; DB on disk - not changed (WAL); enter <Abort T>; resubmit T

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,400>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1600>
8							<commit,T>
9	Flush Log						
10	Output(A)	1600	400	1600	400	1500	
11	Output(B)	1600	400	1600	400	1600	

Redo Logging: Crash Recovery

Crash Occurs after Step 8:
<Commit T> - on disk; T is redone with the help of Log entries;
<Commit T> not on disk: T is treated as incomplete; No action reqd;
DB on disk - not changed (WAL); enter <abort T>; resubmit T

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,400>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1600>
8							<commit,T>
9	Flush Log						
10	Output(A)	1600	400	1600	400	1500	
11	Output(B)	1600	400	1600	400	1600	

Redo Logging: Crash Recovery

Crash Occurs after Step 9:
<Commit T> - surely on disk; T is redone with the help of Log entries;
Whether or not T succeeded in writing them!

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,400>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1600>
8							<commit,T>
9	Flush Log						
10	Output(A)	1600	400	1600	400	1500	
11	Output(B)	1600	400	1600	400	1600	

Undo-Redo Logging

Undo Logging:

Cancels the effect of incomplete Txns

All DB items to be sent to disk before Txn can commit

Results in a lot of I/O -- I/O can not be "bunched"

Redo Logging:

Ignores incomplete Txns and

Repeats the work of Committed Txns

Buffer might contain the blocks of several committed Txns – Buffer utilization might come down!

Undo-Redo Logging:

Better flexibility; uses more detailed logging

Undo-Redo Logging

Undo-Redo Logging Update Entry:

<T, X, U, V> :

Txn T has changed the db item X and its old value is U and new value is V

UR Logging Rule:

Before modifying any database element X on disk, because of changes made by some transaction T, it is necessary that the update record <T, X, U, V> appears on disk.

<Commit T> and disk changes – in any order!

Undo-Redo Logging Example

Txn T is doing money transfer of Rs 100/- from account A to account B
Consistency Requirement: A+B is same before and after the Txn

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500,400>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500,1600>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10							<commit,T>
11	Output(B)	1600	400	1600	400	1600	

Recovery Using Undo-Redo Logs

Examine Log and Partition Txns into:

Committed Set : Txns for which <Commit T> exists

Incomplete Set : Txns for which <Abort T> exists or <Commit T> does not exist

Recovery Method:

Redo all committed Txns – order – earliest first

Undo all incomplete Txns – order – latest first

Necessary to do both!

Undo-Redo Log: Crash Recovery

Crash before <Commit T> is on disk: Txn T – incomplete – undone.
Crash after <Commit T> is on disk: Txn T – completed – redone.

Step	Action	m	Mem-A	Mem-B	Disk-A	Disk-B	Log
1							<start T>
2	Read(A,m)	500	500		500	1500	
3	m:=m-100	400	500		500	1500	
4	Write(A,m)	400	400		500	1500	<T,A,500,400>
5	Read(B,m)	1500	400	1500	500	1500	
6	m:=m+100	1600	400	1500	500	1500	
7	Write(B,m)	1600	400	1600	500	1500	<T,B,1500,1600>
8	Flush Log						
9	Output(A)	1600	400	1600	400	1500	
10							<commit,T>
11	Output(B)	1600	400	1600	400	1600	

Dirty Data Issue and Commits

- Consider Txns T and S:
  - T has modified a db item X and it is doing some more work
  - Meanwhile, S has read X, completed its work
  - Suppose S is allowed to Commit
  - Now, T has an internal error and decides to Abort
  - S has read 'dirty' data
  - But S can't be *undone* as it was allowed to *commit*
- DB got into a trouble!

Recoverable Schedules

- A schedule S is called ***recoverable*** if no Txn T in S *commits* until all the Txns that have written an item T reads have committed
  - T needs to wait till each of the Txn from which it has read completes.
  - If all commit, then T can go ahead and commit
  - If at least one such Txn aborts, T also has to abort
  - Cascading Aborts/Rollbacks may occur

Recoverability vs Serializability

- Orthogonal concepts
- Both are important for a Transaction System!
- It is possible that a recoverable schedule is not conflict-serializable
  - Recoverability defn has no restrictions on locking
- It is also possible that a serializable schedule is not recoverable
  - Serializabilty defn has no restriction on committing

Cascadeless Schedules

- A schedule is called *cascadeless* or
  ACR(avoiding cascading rollbacks)

  if in the schedule, Txns may read only values written by committed Txns

Strict Schedules

- A schedule is called *strict* if in the schedule
  - a Txn neither reads nor writes an item X until the last Txn that writes X has terminated
- Strict Locking:
  - A Txn must not release any write locks until the Txn has either committed or aborted and the commit or abort log record has been written to disk
  - Results in strict schedules
  - Strict schedules are *cascadeless* and *serializable*