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CONTENTS

4.1	Introduction	43
4.2	Original bat algorithm	44
	4.2.1 Description of the bat algorithm	44
	4.2.2 Pseudo-code of BA	45
	4.2.3 Parameters in the bat algorithm	45
4.3	Source code of bat algorithm in Matlab	45
4.4	Source code in C++	48
4.5	A worked example	5 0
4.6	Conclusion	52
	References	52

4.1 Introduction

The original bat algorithm (BA) was first developed by Xin-She Yang in 2010 [1], which attempts to mimic the main characteristics of echolocation of microbats. Due to its good convergence, it has been applied to solve various optimization problems [2], and it has also been extended to solve a multi-objective optimization problem [3] and chaotic variants [4]. Its applications are diverse [5, 6], and a detailed review about the earlier literature was first given by Yang and He in 2013 [2]. In addition, many variants have been developed to improve its efficiency in various ways, including a hybrid bat algorithm [7], directional bat algorithm [8, 9], discrete bat algorithm [10], and others.

The global convergence of the standard bat algorithm has recently been proved mathematically using both a Markovian framework and dynamical system theory [11].

The rest of this chapter provides all the fundamentals of the bat algorithm with the main pseudo-code, Matlab code and C++ code, followed by some simple demonstration of how it works.

4.2 Original bat algorithm

Bats, especially microbats, use echolocation for navigation, which consists of emitting a series of short, ultrasonic pulses that typically last a few milliseconds with frequencies in the range of 25 kHz to about 150 kHz. The loudness of such bursts can be up to 110 dB. When homing for prey, microbats typically increase their pulse emission rates and frequencies, which is also accompanied by variation of their loudness. The main purpose of such frequency tuning and echolocation is for navigation and hunting so as to increase the accuracy of detection and success rate of capturing the prey. Such characteristics are simulated in the bat algorithm.

4.2.1 Description of the bat algorithm

For a population of n bats, each bat (say, bat i) is associated with a location or position vector \mathbf{x}_i and a corresponding velocity vector \mathbf{v}_i . During the iterations, each bat can vary its pulse emission rate r_i , loudness A_i and its frequency f_i .

In general, a position vector is considered as a solution vector to an optimization problem in a D-dimensional search space with D independent design variables

$$\mathbf{x} = [x_1, x_2, x_3, ..., x_D],$$
 (4.1)

which will vary with the iteration or a pseudo-time counter t. So, we use the notation \boldsymbol{x}_i^t to denote the position of bat i at iteration t. Its corresponding velocity is thus denoted by \boldsymbol{v}_i^t .

Frequency tuning can be carried out by

$$f_i = f_{\min} + \beta (f_{\max} - f_{\min}), \tag{4.2}$$

where f_{\min} and f_{\max} are the minimum and maximum ranges, respectively, of the frequency f_i for each bat i. Here, β is a uniformly distributed number drawn from [0, 1]. In most applications, we should use $f_i = O(1)$. For example, we can start with $f_{\min} = 0$ and $f_{\max} = 2$.

This frequency variation is then used for updating the velocity of each bat

$$\mathbf{v}_{i}^{t+1} = \mathbf{v}_{i}^{t} + (\mathbf{x}_{i}^{t} - \mathbf{x}_{*})f_{i},$$
 (4.3)

where x_* is the best solution found at iteration t.

The position or solution vector x_i is updated by

$$x_i^{t+1} = \boldsymbol{x}_i^t + (\Delta t)\boldsymbol{v}_i^{t+1}, \tag{4.4}$$

where Δt is the iteration or time increment. As all iterative algorithms are updated in a discrete manner, we usually set $\Delta t = 1$. Thus, we can simply consider the vectors without any physical units, and then write the update equation as

$$x_i^{t+1} = x_i^t + v_i^{t+1}. (4.5)$$

The pulse emission rate r_i can monotonically increase from a lower value $r_i^{(0)}$, while the loudness can reduce from a higher value $A^{(0)} = 1$. We have

$$A_i^{t+1} = \alpha A_i^t, \quad r_i^{t+1} = r_i^{(0)} [1 - \exp(-\gamma t)], \tag{4.6}$$

where $0 < \alpha < 1$ and $\gamma > 0$ are constants. When t is large enough $(t \to \infty)$, we have $A_i^t \to 0$ and $r_i^t \to r_i^{(0)}$. To start with the simple demonstration later, we can use $\alpha = 0.97$ and $\gamma = 0.1$.

For the local modification around a member of the best solution, we can use

$$\mathbf{x}_{\text{new}} = \mathbf{x}_{\text{old}} + \sigma \epsilon_t A^{(t)},$$
 (4.7)

where ϵ_t is a random number drawn from a normal distribution N(0,1), and σ is the standard deviation acting as a scaling factor. For simplicity, we can use $\sigma = 0.1$ in our later implementation. Here, $A^{(t)}$ is the average loudness at iteration t.

4.2.2 Pseudo-code of BA

Based on the above descriptions, the main steps of the BA consists of a single loop with some probabilistic switching during the iteration. Thus, the procedure of the bat algorithm is summarized in the pseudo-code shown in Algorithm 3.

4.2.3 Parameters in the bat algorithm

Apart from the population size n, there are two parameters and four initialization values in the bat algorithm. They are $\alpha, \gamma, f_{\min}, f_{\max}, A_i^{(0)}$ and $r_i^{(0)}$. In our implementation, we have used $\alpha = 0.97, \gamma = 0.1, f_{\min} = 0, f_{\max} = 2$, and $A_i^{(0)} = r_i^{(0)} = 1$. Though n should vary for different applications, we have used n = 10 here.

Algorithm 3 Pseudo-code of the bat algorithm.

```
1: Define the objective function f(x)
 2: Initialize the bat population x_i and v_i (i = 1, 2, ..., n)
 3: Initialize frequencies f_i, pulse rates r_i and loudness A_i
 4: Set the iteration counter t=0
   while (t < t_{\text{max}}) do
       Vary r_i and A_i
6:
 7:
       Generate new solutions by adjusting frequencies
       Update velocities and locations/solutions via Eqs (4.3) and (4.5)
8:
       if rand > r_i then
9:
           Select a solution among the best solutions
10:
           Generate a local solution around the selected best solution
11:
       end if
12:
       Generate a new solution by flying randomly
13:
       if rand > A_i and f(x_i) < f(x_i) then
14:
           Accept the new solution
15:
       end if
16:
       Rank the bats and find the current best solution x_*
17:
18: end while
```

4.3 Source code of bat algorithm in Matlab

The implementation of this algorithm is relatively straightforward using Matlab. Here, we present a detailed description of a simple Matlab implementation to find the minimum of the function f(x)

minimize
$$f(\mathbf{x}) = (x_1 - 1)^2 + (x_2 - 1)^2 + \dots + (x_D - 1)^2, \quad x_i \in \mathbb{R}, \quad (4.8)$$

which has a global minimum of $x_* = (1, 1, ..., 1)$. In theory, the algorithm can search the whole domain; we always impose some limits in practice. Here, for simplicity, we use the following lower bound (Lb) and upper bound (Ub):

$$Lb = [-5, -5, ..., -5], \quad Ub = [+5, +5, ..., +5].$$
 (4.9)

The Matlab code for the bat algorithm here consists of three parts: initialization, the main loop, and the objective function. The whole codes should consist of all the lines of codes in a sequential order. However, for ease of understanding, we split them into three parts.

The first part is mainly initialization of parameter values such as α and γ , as well as the generation of the initial population of n=10 bats. The cost function is the objective function to be given later in a Matlab function.

It is worth pointing out that the simple bounds should be checked at each iteration when new solutions are generated. However, for simplicity, the implementation here does not carry out this check, and care should be taken when extending the simple computer code here.

```
1 function [best,fmin]=bat algorithm
2 % Default parameters
n = 10;
                   % Population size, typically 10 to 25
                   \%\ Loudness\ (constant\ or\ decreasing)
4 A=1;
                   % Pulse rate (constant or decreasing)
5 r = 1;
6 alpha = 0.97;
                  % Parameter alpha
7 \text{ gamma} = 0.1;
                  % Parameter gamma
8 % Frequency range
9 Freq min=0;
                   % Frequency minimum
10 Freq max = 2;
                   % Frequency maximum
11 % Max number of iterations
12 t_max=1000;
                   % Initialize iteration counter
13 t = 0:
14 % Dimension of the search variables
15 d=5;
16 % Initialization of arrays
Freq=zeros(n,1); % Frequency
v=zeros(n,d);
                      % Velocities
                      % Lower bounds
19 Lb=-5*ones(1,d);
                    % Upper bounds
20 Ub=5*ones(1,d);
21 % Initialize the population/solutions
22 for i = 1:n,
    Sol(i,:) = Lb + (Ub-Lb) .* rand(1,d);
23
24
    Fitness(i)=Fun(Sol(i,:));
25 end
26 % Find the best solution of the initial population
27 [fmin, I]=min(Fitness);
28 best=Sol(I,:);
```

Listing 4.1

BA demo initialization.

The second part is the main part, consisting a loop over the whole population at each iteration. The pulse emission rate and loudness are varied first, then new solutions are generated and evaluated. The new population is checked and the current best solution is found and updated. Though each bat can have its own pulse emission rate r_i and loudness A_i , we have used the same values of r and A for all bats. This means that the average of A_i is now simply A itself.

```
1 % Start the iterations -- Bat Algorithm
_2 while (t < t_max)
      % Varying parameters
      r=r*(1-exp(-gamma*t));
      A=alpha*A;
      % Loop over all bats/solutions
   for i=1:n,
     Freq(i)=Freq_min+(Freq_max-Freq_min)*rand;
v(i,:)=v(i,:)+(Sol(i,:)-best)*Freq(i);
9
      S(i,:) = Sol(i,:) + v(i,:);
10
      % Check a switching condition
12
      if rand>r,
      S(i,:) = best + 0.1*randn(1,d)*A;
      end
14
15
      % Evaluate new solutions
16
      F_{\text{new}} = F_{\text{un}}(S(i,:));
17
18
      % If the solution improves or not to loud
      if (Fnew <= Fitness(i)) & (rand >A)
19
     Sol(i,:)=S(i,:);
20
```

```
Fitness (i)=Fnew;
      end
23
24
    % Update the current best solution
      if Fnew<=fmin,
25
        best=S(i,:);
26
27
        fmin=Fnew;
28
    end % end of for
29
   t=t+1; % Update iteration counter
30
31 end
32 % Output the best solution
33 disp(['Best =',num2str(best),' fmin=',num2str(fmin)]);
```

Listing 4.2

Main iterations for the bat algorithm.

The third part is the objective function and ways for implementing the lower and upper bounds. This will ensure the solution vector should be within the regular bounds; however, we have not implemented the bounds here for simplicity. A new solution is evaluated by calling the objective or cost function.

```
1 % Cost or objective function
2 function z=cost(x)
3 z=sum((x-1).^2); % Solutions should be (1,1,...,1)
```

Listing 4.3

The objective function.

If we run this code with a maximum number of 1000 iterations, we can get the best minimum value as $f_{best} = 3.49 \times 10^{-29}$. Obviously, due to the random numbers used, these results are not exactly repeatable, but the order of the magnitude is easily reachable in the simulation.

4.4 Source code in C++

As the bat algorithm is simple, it can be implemented easily in any other programming language. For the same algorithm to solve the same problem as given above in Matlab, we can implement it in C++ as follows:

```
1 #include <iostream>
2 #include <time.h>
3 #include <random>
4 using namespace std;
5 //Definition of the objective function OF(.)
6 float Fun(float x[], int size_array)
7 {
      float t=0;
          for (int i=0; i < size_array; i++)
           \{t=t+(x[i]-1)*(x[i]-1);\}
9
10
      return t;
11 }
  //Generate pseudo random value from the range [0, 1]
float ra() {return (float)rand()/RAND_MAX;}
14 int main()
15 {
```

```
16 //Normal distribution generator
std::default_random_engine generator;
18 std::normal distribution < double distribution (0,1.0);
19 //Initialization of pseudo random generator
20 srand (time (NULL));
                     //Population size, typically 10 to 25
21 int n=10;
22 float A=1.0;
                     //Loudness (constant or decreasing)
                     //Pulse rate (constant or increasing)
23 float r = 1.0;
124 float alpha = 0.97; // Parameter alpha
25 float gamma=0.1; //Parameter gamma
  //Frequency range
                         //Frequency minimum
float Freq min = 0.0;
28 float Freq max = 2.0;
                         //Frequency maximum
29 //Max number of iterations
30 int t max=1000;
  //Dimension of the search variables
32 int d=5;
33 //Initialization of arrays
                    //Frequency
34 float Freq[n];
                     // Velocities
35 float v[n][d];
36 for (int i = 0; i < n; i + +)
37
38
       for (int j=0; j<d; j++)\{v[i][j]=0;}
39
40 float Lb[d]; //Lower bounds
41 float Ub[d]; //Upper bounds
                       //Solutions
42 float Sol[n][d];
43 float Fitness[n];
                       //Fitness
44 float S[n][d];
                        //New solutions
45 float rnum;
46 float Fnew;
   //Prepare arrays
47
48 for (int j=0; j < d; j++) {Lb[j]=-5.0; Ub[j]=5.0;}
49 //Initialize the population/solutions
50 for (int i = 0; i < n; i + +)
51
       for (int j=0; j< d; j++)
53
           Sol[i][j]=Lb[j]+(Ub[j]-Lb[j])*ra();
54
       Fitness [i]=Fun(Sol[i],d);
56
57 }
58 //Find the best solution of the initial population
59 int I=0; float fmin;
60 fmin=Fitness[I];
61 for (int i = 1; i < n; i + +)
62 {
       if (Fitness [i]<=fmin)
63
64
65
           fmin=Fitness[i];
66
           I=i;
67
       }
68 }
69 float best [d];
70 for (int i = 0; i < d; i + +)
71
72
           best [i]=Sol [I][i];
   //Start the iterations -- Bat Algorithm
74
75 for (int t=1; t \le max; t++)
76
  {
77
       //Varying parameters
78
       r=r*(1-exp(-gamma*t));
79
       A=alpha*A;
       //Loop\ over\ all\ bats/solutions
80
       for (int i = 0; i < n; i + +)
81
```

```
Freq[i]=Freq min+(Freq max-Freq min)*ra();
83
84
               rnum=ra();
85
                for (int j=0; j < d; j++)
86
                      \begin{array}{l} v\,[\,i\,][\,j\,] \!=\! v\,[\,i\,][\,j\,] \!+\! (\,\mathrm{Sol}\,[\,i\,][\,j\,] \!-\! \mathrm{best}\,[\,j\,]) \,*\,\mathrm{Freq}\,[\,i\,]\,;\\ S\,[\,i\,][\,j\,] \!=\! Sol\,[\,i\,][\,j\,] \!+\! v\,[\,i\,][\,j\,]\,; \end{array} 
87
                     if (rnum>r)
89
90
                     {
                           S[i][j] = best[j] + 0.1*A*distribution(generator);
91
93
                //Evaluate new solutions
94
               Fnew=Fun(S[i],d);
95
                //If the solution improves or not to loud
96
97
                if ((Fnew <= Fitness[i]) && (ra()>A))
98
                     Fitness [i]=Fnew;
99
                     for (int j=0; j< d; j++)
                           Sol[i][j]=S[i][j];
103
104
                //Update the current best solution
               if (Fnew<=fmin)
106
107
                     fmin=Fnew;
108
                     for (int j=0; j< d; j++)
                           best [ j ]=S[ i ] [ j ];
112
113
               }
115 }
116 //Output the best solution
117 cout << "#### fmin = "<< fmin << endl;
118 cout<<"#### Best = [ ";
for(int i=0; i<d; i++) {cout<<best[i]<<" ";}
120 cout << " ] " << endl;
121 getchar();
122 return 0;
123 }
```

Listing 4.4

Bat algorithm in C++.

4.5 A worked example

Let us use the bat algorithm to find the minimum of

$$f(\mathbf{x}) = (x_1 - 1)^2 + (x_2 - 1)^2 + (x_3 - 1)^2, \tag{4.10}$$

in the simple ranges of $-5 \le x_i \le 5$. This simple 3-variable problem has the global minimum solution $x_{\text{best}} = [1, 1, 1]$.

For the purpose of algorithm demonstration, we only use n = 5 bats. Suppose the initial population is randomly initialized and their objective (fitness)

values are as follows:

$$\begin{cases} x_1 = \begin{pmatrix} 5.00 & 0.00 & 5.00 \end{pmatrix}, & f_1 = f(x_1) = 33.00, \\ x_2 = \begin{pmatrix} 2.00 & 2.00 & 3.00 \end{pmatrix}, & f_2 = f(x_2) = 6.00, \\ x_3 = \begin{pmatrix} -3.00 & -2.00 & 0.00 \end{pmatrix}, & f_3 = f(x_3) = 26.00, \\ x_4 = \begin{pmatrix} -5.00 & 0.00 & 5.00 \end{pmatrix}, & f_4 = f(x_4) = 53.00, \\ x_5 = \begin{pmatrix} 3.00 & 4.00 & 5.00 \end{pmatrix}, & f_5 = f(x_5) = 29.00, \end{cases}$$
(4.11)

where we have only used two decimal places for simplicity. Though the actual random numbers generated for the initial population can be any numbers in the simple bounds, here we have used the numbers that are rounded up as the initialization; this is purely for simplicity and clarity.

Obviously, the best solution with the lowest (or best) value of the objective function is x_2 with $f_2 = 6.00$ for this population at t = 0. Thus, we have

$$\mathbf{x}_{\text{best}} = [2.00, 2.00, 3.00]. \tag{4.12}$$

In principle, each bat in the bat population can have its own pulse emission rate r_i and loudness A_i . However, for simplicity of implementation and discussion here, we can use the same r and A for all bats.

At the first iteration t=1, starting with $r^{(0)}=1$, $A^{(0)}=1$, $\alpha=0.97$ and $\gamma=0.1$, we have

$$r = r^{(0)}[1 - \exp(-\gamma t)] = 1 \times [1 - \exp(-0.1 \times 1)] = 0.0952,$$
(4.13)

$$A = \alpha A^{(0)} = 0.97 \times 1 = 0.97. \tag{4.14}$$

When carrying out the loop over all the 5 bats, if we draw a random number $\beta = 0.2$, then the frequency becomes $f_i = 0 + 0.2 \times (2 - 0) = 0.4$, so the velocity \mathbf{v}_1 can be calculated by

$$\mathbf{v}_1^1 = \mathbf{v}_1^0 + (\mathbf{x}_1 - \mathbf{x}_{\text{best}}) \times 0.4 = [4.20, -2.80, 2.80],$$
 (4.15)

where we have used $f_{\min} = 0$ and $f_{\max} = 2$. This gives the new solution

$$x_1^1(\text{new}) = x_1 + v_1 = [9.20, -2.80, 7.8],$$
 (4.16)

which in turn gives $f(x_1^1) = 127.92$. This is worse than the previous solutions. However, the next step is to generate a uniformly distributed random number. Suppose we get $\beta = 0.5$, which is greater than r = 0.0952; we carry out a local search by

$$\boldsymbol{x}_{1}^{1}(\text{new}) = \boldsymbol{x}_{\text{best}} + 0.1\epsilon A. \tag{4.17}$$

If we draw a random number vector, we get

$$\epsilon = [-2, 0.5, -1.5],$$
(4.18)

and we have

$$\mathbf{x}_{1}^{1}(\text{new}) = \mathbf{x}_{\text{best}} + 0.1\epsilon \times 0.97 = [1.8060, 2.0485, 2.8545].$$
 (4.19)

This gives $f(x_1^1) = 5.1882$, which is better than x_{best} . So we update it as the new best solution

$$\mathbf{x}_{\text{best}} = [1.8060, 2.0485, 2.8545]. \tag{4.20}$$

Similarly, we do this for the rest of the other four solutions in the population. Once the first iteration is done, we vary r and A again, and then carry out another loop over the whole population.

The bat algorithm can be efficient. For the objective function, the optimal solution is $x_* = [1, 1, 1]$. Using the above codes, we can typically get $f_{\text{best}} = 4.9 \times 10^{-16}$ after 500 iterations.

4.6 Conclusion

The bat algorithm is a simple, flexible and yet efficient algorithm for solving optimization problems. In addition, this algorithm is relatively simple to implement with a minimum use of memory, and the computation cost is also low. The bat algorithm has been extended to other forms with many variants and applications [2, 10, 11].

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Conclusion 53

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