Transform Order by Reversals : ILP vs SAT-solver

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For Integer Linear Programming, we have the following conditions:
1.
For each k from 1 to $n-1$:

$$NOP(k) + \sum_{(p,q):p<q} R(p,q,k) = 1$$

For an equivalent dnf, we see that there is one and only one 1 among all these values.

$$(NOP(k) \wedge \neg R_1 \wedge \neg R_2 \wedge ......) \vee (\neg NOP(k) \wedge R_1 \wedge \neg R_2 \wedge ......) \vee ......$$

We want a cnf that is not too complicate. Logic: One of them is true and none of the pairs can be both true.

$$(NOP(k) \vee R_1 \vee R_2 \vee ...) \wedge (\neg NOP(k) \vee \neg R_1) \wedge (\neg NOP(k) \vee \neg R_2) \wedge (\neg R_1 \vee \neg R_2) \wedge ...$$

2
For i from 1 to n:

$$\sum_{q=1}^{n} X(1,i,i,q) = 1$$

For an equivalent cnf, we see that there is one and only one 1 among all these values.

$$(X_1 \vee X_2 \vee X_3 \vee ...) \wedge (\neg X_1 \vee \neg X_2) \wedge (\neg X_1 \vee \neg X_3) \wedge (\neg X_2 \vee \neg X_3) \wedge ...$$

3
For i from 1 to n:

$$\sum_{p=1}^{n} X(n-1,i,p,Q_2(i)) = 1$$

For an equivalent cnf, we see that there is one and only one 1 among all these values.

$$(X_1 \vee X_2 \vee X_3 \vee ...) \wedge (\neg X_1 \vee \neg X_2) \wedge (\neg X_1 \vee \neg X_3) \wedge (\neg X_2 \vee \neg X_3) \wedge ...$$

4
For each node q at a level k from 2 to $n-1$:

$$\sum_{p=1}^{n}[\sum_{i=1}^{n} X(k-1,i,p,q)] = 1$$

We already know there is at most one 1. Hence for cnf:

$$(X_1 \vee X_2 \vee X_3 \vee ...)$$

5

For each node p at a level k from 2 to $n-1$:

$$\sum_{q=1}^{n}[\sum_{i=1}^{n} X(k,i,p,q)] = 1$$

We already know there is at most one 1. Hence for cnf:

$$(X_1 \vee X_2 \vee X_3 \vee ...)$$

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