# Econ 600: taught by Prof. Shaowei Ke

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September 2, 2015

## Disclaimer

This is a personal note of mine. I will try to follow professor Ke's lecture as close as possible. However, neither is this an official lecture note, nor will Linfeng be responsible for any errors + typos. Nevertheless, corrections and suggestions are always welcomed.

As this lecture note will be maintained on Github, PLEASE:

- Use the "Issues" feature on Github to post suggestions;
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Paragraphs starting with "Note that ..." are most likely my personal reflections. Please be aware of this.

# 1 Lecture 1: Logic, Sets and some Real Analysis<sup>1</sup>

## 1.1 Logic

**Definition 1.1. Proposition** is a sentence that is either *true* or *false*. It cannot be both true and false.

Note: "true" and "false" may not necessarily be based on any (objective/subjective) factual basis. However, to give a concrete example, contextually correct propositions are usually employed.

**Definition 1.2.** Logic Connectives:  $\wedge$  and  $\vee$ . Let P and Q be propositions

- Conjunction of P and Q is denoted as  $P \wedge Q$ ;
- Disjunction of P and Q is denoted as  $P \vee Q$ .
- Negation of P is denoted as:  $\neg P$ .

P	Q	$P \wedge Q$	$P \lor Q$	$\neg P$
1	1	1	1	0
1	0	0	1	0
0	1	0	1	0
0	0	0	0	1

Table 1: Truth Table for logic connectives

**Truth Table** is vaguely defined, with each row being a possible "state of the world". On top of this,

**Definition 1.3** (Conditionals and Biconditionals). Let P, Q, R be propositions,

- 1. Conditional of P and Q is  $P \implies Q$ ;
- 2. Bi<br/>conditional of P and Q is<br/>  $P \iff Q.$

P	Q	$P \implies Q$	$P \iff Q$
1	1	1	1
1	0	0	0
0	1	1	0
0	0	1	1

Table 2: Truth Table for Conditionals and Biconditionals

Note that, the two 1's are obtained for free. Conditional of P and Q are trivially true if P is false (thus the conditional is not entered, thereby cannot be disproved?).

←Check This.

Additionally, from an external source ( $\leftarrow$  click me!):

Conditionals are FALSE only when the first condition (if) is true and the second condition (then) is false. All other cases are TRUE.

**Definition 1.4.** Two propositions are **equivalent** if they have the same truth table, denoted using " $\equiv$ ".

**Example 1.** Claim: that  $P \implies Q$  and  $\neg Q \implies \neg P$  are equivalent.

*Proof.* Refer to table 3: that by definition, the truth table of the two conditionals are the same.  $\Box$ 

Note, (it seems that) $^a$  truth tables are the same if the two "column vectors" denoting the true/false status are the same.

<sup>a</sup>Since "truth table" was not explicitly defined.

**Definition 1.5** (Tautology). A proposition whose truth table consists only 1's is called **tautology**.

<sup>&</sup>lt;sup>1</sup>Relation, Function, Correspondence and Sequences in  $\mathbb{R}$ 

Table 3: Truth Table: equivalence of  $P \implies Q$  and  $\neg Q \implies \neg P$ 

P	Q	$P \implies Q$	$  \neg Q \implies \neg P$
1	1	1	1
1	0	0	0
0	1	1	1
0	0	1	1

**Example 2.** Claim:  $Q \implies (P \implies Q)$  is a tautology.

*Proof.* Refer to Table 4

Table 4: Truth Table: Tautology

P	Q	$P \implies Q$	$Q \implies (P \implies Q)$
1	1	1	1
1	0	0	1
0	1	1	1
0	0	1	1

### **Remark 1.6.** We introduce the following 4 types of proof:

- 1. Direct proof: to follow the direction of the statement.
  - **Proposition**: For odd integers x, y, x + y is an even integer.
- 2. Proof by contrapositive: (restate the proposition and prove the easier direction).
  - **Proposition**: If  $n^2$  is an odd integer (P), then n is an odd integer.

*Proof.* Prove instead that: "if n is an even integer, then  $n^2$  is an even integer".  $\square$ 

- 3. Proof by contradiction: (construct a structure that leads to contradiction between derived conditions and given conditions.).
  - That  $\sqrt{2}$  is rational number<sup>2</sup>.
- 4. Proving a "if and only if" statement/proposition to be true: either one of the following 4 are valid strategies:
  - (a)  $P \implies Q$  and  $Q \implies P$ ;
  - (b)  $P \implies Q$  and  $\neg P \implies \neg Q$ ;
  - (c)  $\neg Q \implies \neg P \text{ and } Q \implies P$ ;
  - (d)  $\neg Q \implies \neg P \text{ and } \neg P \implies \neg Q$ .

 $<sup>^{2}</sup>$ The set of rational numbers is denoted as Q.

### 1.2 Sets

**Remark 1.7** (Russell's paradox). The barber is a man who shaves all those and only those who do not shave themselves.

In terms of set theory, let  $R = \{x : x \notin x\}$ , then:

$$R \in R \iff R \notin R$$

which is very problematic.

**Definition 1.8** (Sets). There are two definition of sets:

- (Enumerating all elements)
   A set is a collection of objects, e.g. {1,2,...} <sup>3</sup> or {1,2} <sup>4</sup>.
- 2. (Describing properties to be satisfied by elements in the set)

  If A is a set of all objects that satisfies property P, then we can write

$$A = \{x : P(x)\}$$

where the colon means "such that", and P(x) means that x satisfies property P.

Now, we can define the following **sets** using the two definitions of sets:

- (Natural Number)  $N = \{1, 2, \ldots\};$
- (Integer)  $Z = \{x : x = n \text{ or } x = -n \text{ or } x = 0, \text{ for some } n \in N\};$
- (Rational number)  $Q = \{x : x = \frac{m}{n}, m, n \in Z\}.$

**Definition 1.9** (Set Equality). Two sets A and B are equal if they have the same elements. That is:

$$A = B$$
 if and only if  $x \in A \iff x \in B, \forall x$ 

Note, that the notion  $\forall x$  was used sloppily here. Without loss of generality, it shall better be  $\forall x \in A \bigcup B$ .

**Definition 1.10** (Set Containment). A set A is contained in a set B, denoted by  $A \subseteq B$ , if  $\forall x \in A \implies x \in B$ .

As a consequence, A = B if and only if  $A \subseteq B$  and  $B \subseteq A$ .

**Definition 1.11** (Cardinality (finite case)). If a set A has  $n \in \mathbb{N}^5$  distinct elements, then n is the cardinality of A and we call A a finite set. The **cardinality of** A is denoted by |A|.

**Definition 1.12** (Empty set  $\emptyset$ ). The empty set is the set with no element.

<sup>&</sup>lt;sup>3</sup>a countably infinite set.

<sup>&</sup>lt;sup>4</sup>a finite set.

<sup>&</sup>lt;sup>5</sup>Natural number.

**Definition 1.13** (Power set  $2^A$ ). Let A be a set. The **power set of** A is the collection of all subsets of A.

Note that, A is an arbitrary set. It could be finite, in which case  $2^A$  easy to envision; At the other extreme, it could be a uncountable set. Nevertheless, the following equality shall hold:

$$|2^A| = 2^{|A|}$$

**Example 3.** Let  $A = \{1,3\}$ , then  $2^A = \{\emptyset, \{1\}, \{3\}, \{1,3\}\}$ . In terms of notation, note that 1 is an element in A, thus  $1 \in A$ ; yet,  $\{1\}$  is a subset of A, thus  $\{1\} \subset A$ .

**Definition 1.14** (Operations on sets:  $\bigcup$ ,  $\bigcap$ ,  $\setminus$  and  $\cdot^c$ .). Let A and B be two sets:

- Union:  $A \bigcup B := \{x : x \in A \lor x \in B\};$
- Intersection:  $A \cap B := \{x : x \in A \land x \in B\};$
- A and B is disjoint if  $A \cup B = \emptyset$ ;
- Difference of A and B is defined as:  $A \setminus B := \{x \in A \land x \notin B\};$
- Complements of  $A: A^c := \{x : x \notin A\}.$

Side note: Index set I is a countable set.

$$\bigcup_{i \in I} A_i = \{x : x \in A_i \text{ for some } i \in I\}$$

**Definition 1.15** (de Morgan's law).

$$\left(\bigcup_{i\in I} A_i\right)^c = \bigcap_{i\in I} \left(A_i^c\right) \text{ and } \left(\bigcap_{i\in I} A_i\right)^c = \bigcup_{i\in I} \left(A_i^c\right)$$

**Exercise 1.16.** Prove that  $(A \bigcup B)^c = A^c \cap B^c$ .

*Proof.* Prove mutual containment using element argument.

Counters reset

### 1.3 Relation, Function and Correspondence

**Definition 1.1** (Ordered pair). For two sets A and B, an ordered pair is (a, b) such that  $a \in A$  and  $b \in B$ .

**Definition 1.2** (*n*-taple). Let there be *n* sets:  $A_1, \ldots, A_n$ , an *n*-taple is  $(a_1, \ldots, a_n)$  such that  $a_i \in A_i, \forall i = 1, 2, \ldots n$ .

**Definition 1.3** (Cartesian Product). Let  $A_1, \ldots, A_n$  be non-empty sets. Cartesian product of  $A_1, \ldots, A_n$  is  $A_1 \times \cdots \times A_n$ , defined as:

$$\Pi_{i=1}^{n} A_i = \{(a_1, \dots, a_n) : a_i \in A_i, \forall i = 1, \dots, n\}$$

**Definition 1.4** (Relation). A relation from set A to set B is a subset of  $A \times B$ , denoted by R.

$$aRb \iff (a,b) \in R$$

A relation on A is a subset of  $A \times A$ .

**Definition 1.5.** A relation  $R \subseteq A \times A$  is said to be:

- reflective if  $aRa \ \forall a \in A$ . (That is,  $(a, a) \in R, \ \forall a \in A$ .);
- complete if either aRb or bRa,  $\forall a, b \in A$ ;
- symmetric if  $\forall a, b \in A$ ,  $aRb \implies bRa$ ;
- antisymmetric if  $\forall a, b \in A$ , aRb and  $bRa \implies a = b$ .
- transitive if  $\forall a, b, c \in A$  s.t. aRb and bRc, aRc (is implied).

Table 5: Property of common relations

	<	$\leq$	$\mid$ $\in$	$\subseteq$	$\succeq$
reflective	X	1	X	1	1
complete	X	1	X	X	1
symmetric	X	X	X	X	X
antisymmetric	1	1	<b>✓</b>	1	X
transitive	1	1	X	1	1

Note that, < and  $\le$  are defined on  $\mathbb{R}$ ;  $\in$  and  $\subseteq$  are defined on sets;  $\succeq$  is preference relation that represents "weakly prefer".

Also note that, completeness implies reflectiveness.

**Definition 1.6** (Equivilence relation). An **equivalence** on set A is a relation E that is reflective, symmetric and transitive. It is denoted as  $\sim$ .

For any  $a \in A$ , the equivalence class of a with respect to  $\sim$  is defined to be the set

$$E_{\sim}(a) = \{ b \in A, b \sim a \}$$

Remark: by construction in Definition 1.4, equivalence ( $\sim$ ) is defined as "a relation on A", which is thereby defined in the Cartesian space.

**Definition 1.7** (Function: defined using Relation from A to B). A function from set A to set B is a relation f from A to B such that:

- (i)  $\forall a \in A, \exists b \in B \text{ such that } (a, b) \in f, \text{ i.e. } afb$
- (ii)  $\forall a \in A$ , if  $(a, b) \in f$  and  $(a, c) \in f$  for some  $b, c \in B$ , then b = c.

Note that, alternatively, the two conditions could be written in short as:

(iii) 
$$\forall a \in A, \exists! b \in B \text{ such that } (a, b) \in f, \text{ i.e. } afb$$

**Convention for** f: If  $(a,b) \in f$ , we write f(a) = b. And, f could be interpreted as a "mapping": " $f: A \to B$ ".

**Definition 1.8** (Domain and Rnage). If f is a function from A to B, then A is called the **domain** of f and B is the **codomain** of f. The **range** of f is the set:

$$Ran(f) = \{b \in B : \exists a \in A \text{ s.t. } f(a) = b\}.$$

**Definition 1.9** (Propoteries of functions). Let f be a function, then:

(i) f is surjective if Ran(f) = B;

onto

(ii) f is **injective** if  $a_1 \neq a_2 \in A \implies f(a_1) \neq f(a_2)$ ;

1-to-1

(iii) f is bijective if f is surjective and injective.

Side note: a indicator function is defined as following: for A being a set and  $S \subseteq A$ ,

$$\mathcal{X}_S(a) = \begin{cases} 1 & \text{if } a \in S \\ 0 & \text{otherwise} \end{cases}$$

**Definition 1.10** (Image and Preimage). For  $f:A\to B$  and  $C\subseteq A$ , the **image** of C under f is

$$f(C) = \{b \in B : \exists a \in C \text{ s.t. } f(a) = b\}$$

The **preimage** of  $D \subseteq B$  is

$$f^{-1}(D) = \{ a \in A : f(a) \in D \}$$

**Exercise.** Prove that

- 1.  $f^{-1}(f(A)) = A$ , and
- 2.  $f(f^{-1}(B)) = B$  if and only if f is surjective.

**Proposition 1.11.** Given  $f: A \to B$ , then  $f^{-1}: B \to A$  is a function if and only if f is bijective.

**Definition 1.12** (Sequence). A sequence is a function  $f: N \to A$ , denoted by  $\{a_1, a_2, \ldots\} = \{a_i\}_{i=1}^{\infty}$  i.e. the set of all sequence is the following set:

$$A^{\infty} = A \times A \times \cdots$$

**Definition 1.13** (Cardinality, for (infinite) sequences). Two sets A, B have the same cardinality if  $\exists$  a bijective function  $f: A \to B$ .

Then,  $|A| \ge |B|$  if there exists an injective function  $f: B \to A$ . (Example:  $|Z| \ge |N|$  by using identify mapping from N to Z;  $|N| \ge |Z|$  by enumerating elements in Z using N. Thus, |Z| = |N|.) Eventually, we have:

$$|\mathbb{R}^2| = |\mathbb{R}| > |Q| = |Z| = |N|$$

**Definition 1.14** (Correspondence).  $T: A \rightrightarrows B$  is a correspondence such that  $T: A \to 2^A \setminus \emptyset$ .

### 1.4 Sequences

**Definition 1.1** (Sequence in  $\mathbb{R}$ ). A sequence of real number is a function  $a: N \to \mathbb{R}$  s.t.  $a(i) = a_i$  is the *i*-th component of the sequence  $\{a_j\}_{j=1}^{\infty}$ .

**Definition 1.2** (Increasing sequence). A real sequence is increasing if  $a_{n+1} \ge a_n \ \forall n \in \mathbb{N}$ .

**Definition 1.3** (Bounded and Bounded (from) above/below). A real sequence is

- bounded above if  $\exists \bar{m} \in \mathbb{R} \text{ s.t. } a_n \leq \bar{m} \ \forall n \in \mathbb{N}$ .
- bounded below if  $\exists \underline{m} \in \mathbb{R} \text{ s.t. } a_n \geq \underline{m} \ \forall n \in \mathbb{N}$ .
- **bounded** if it is bounded above and bounded below.

**Definition 1.4** (Least upper bound).  $a \in \mathbb{R}$  is the least upper bound of a sequence  $\{a_n\}$  if

- (i) a is an upper bound;
- (ii) a is the smallest upper bound, i.e.  $\not\exists b \in \mathbb{R}$  s.t. b < a and b is a upper bound of  $\{a_n\}$ .

**Axiom 1.5** (Axiom of Real Number: completeness axiom). If S is a nonempty set of real numbers that is bounded above, then there exists a least upper bound that is also a real number.

Note, that, claiming that the upper bound is in  $\mathbb{R}$  is redundant.

**Definition 1.6** (Convergence sequences). A real sequence  $\{a_n\}$  converges to the limit  $a \in \mathbb{R}$  if  $\forall \varepsilon > 0, \exists N \text{ s.t. } \forall n \geq N$ 

$$|a_n - a| < \varepsilon$$

We write  $\lim_{n\to\infty} a_n = a$  or  $a_n \to a$ .

• If a sequence does not converge, then it diverges. (To  $+\infty$  or  $-\infty$ .)

**Theorem 1.** A monotone bounded sequence converges.

*Proof.* Discuss two cases where 1)  $\{a_n\}$  is an increasing sequence, and 2)  $\{a_n\}$  is a decreasing sequence. Then, proof is completed through using either least upper bound (for increasing sequence) or largest lower bound (for decreasing sequence).

<sup>&</sup>lt;sup>6</sup>This is an ordered set.

## 2 Lecture 2: convergence and more

### 2.1 Sequence and Convergence

**Definition 2.1.** A set  $S \subset X$  is a linearly ordered set if there is a relation " $\leq$ " on X s.t.

 $\leq$  is complete, transitive and antisymmetric.

Note that, given the linear ordering, we can define < accordingly. (For arbitrary  $a, b \in X$  and  $a \le b$ , then we say a < b if  $a \le b$  and  $a \ne b$ .)

**Definition 2.2** (Boundedness for an arbitrary set.). Let X be a linearly ordered set and  $S \subset X$ , then  $a \in X$  is the **supremum** (or *least upper bound*) of X if:

- 1. a itself is an upper bound of S, i.e.
- 2. for  $b \in X$ , b < a, then b is not an upper bound of S.

Corollary: For  $a = \sup X$ ,  $\forall \varepsilon > 0$ , there exists  $x \in S$  s.t.  $x > a - \varepsilon$ .

**Axiom 2.3** (Completeness Axiom). If S is a nonempty set of real numbers that is bounded above, then there exists a least upper bound.

**Definition 2.4** (Sequence in  $\mathbb{R}$ ). A sequence of real number is a function  $a: N \to \mathbb{R}$  s.t.  $a(i) = a_i$  is the *i*-th component of the sequence  $\{a_j\}_{j=1}^{\infty}$ .

**Remark 2.5.**  $\{a_n\}$  is bounded if a(N) is bounded.

Note, here N is the set of all natural numbers  $\{1, 2, \ldots, \}$ . Thus, we hereby define the boundedness of a sequence using the our previous definition of set-boundedness.

**Lemma 2.6.** A monotone bounded sequence converges.

**Definition 2.7** (Subsequence). A subsequence  $\{a_{n_i}\}$  of  $\{a_n\}$  is a sequence s.t.  $1 \le n_1 \le n_2 \le \ldots$  That is:

 $\exists$  conversion function  $\Phi: N \to N$  s.t.  $n_i = \Phi(i)$  and  $\Phi(i) < \Phi(j)$  whenever i < j. We can also write:  $a_{n_i} = a_{\Phi(i)}$ .

**Lemma 2.8.** Every sequence of  $\mathbb{R}$  has a monotone subsequence.

*Proof.* Proof by doodling: try to construct a decreasing sequence first, if failed (cannot identify infinitely many of elements as candidate of the sequence), construct an increasing one.

Formally: let  $S = \{i : \text{ if } j > i, \text{ then } a_j < a_i\}.$ 

- if |S| = |N| (countably infinite) <sup>1</sup>, we have found a monotone (decreasing) sequence.
- If  $|S| < \infty$ , let max S = N, then by construction,  $\exists n_1 \text{ s.t. } a_{n_1} \ge a_{N+1}$ . Since  $n_1 \notin X$ , there exists  $n_2 > n_1$  s.t.  $a_{n_2} \ge a_{n_1} \ge a_N$ .

We can construct an increasing sequence in this fashion.

<sup>&</sup>lt;sup>1</sup>Writing  $|S| = \infty$  is not rigorous enough, since uncountably infinite could also be denoted similarly.

**Theorem 2.9** (Bolzano-Weierstrass Theorem). A bounded sequence of  $\mathbb{R}$  has a convergent subsequence.

*Proof.* By Lemma 2.8, such bounded sequence of  $\mathbb{R}$  has a monotone subsequence, which inebriates the boundedness property.

Thus, by Lemma 2.6, such bounded monotone sequence converges.

**Remark 2.10** (Properties of Limits). For  $a_n \to a$  and  $b_n \to b$  (two convergent sequences):

- (i)  $c \cdot a_n \to c \cdot a$ , for  $c \in \mathbb{R}$ ;
- (ii)  $a_n + b_n \rightarrow a + b$
- (iii)  $a_n \cdot b_n \to a \cdot b$
- (iv)  $\frac{a_n}{b_n} \to \frac{a}{b}$  s.t.  $b \neq 0$  and  $b_n \neq 0 \ \forall n$ .
- (v)  $\forall n \in \mathbb{N}$ , if  $c \leq a_n$ , then  $c \leq a$ . (Note that we have defined only one linear ordering  $\leq$ .) However,  $a_n > c$  does not imply a > c. (e.g.:  $\frac{1}{n} > 0$ ,  $\forall n$ , yet  $\frac{1}{n} \to 0 = 0$ .)
- (vi)  $\forall n$ , if  $b_n \leq a_n$ , then  $b \leq a$ .

**Definition 2.11** (Cauchy sequence).  $\{a_n\}$  is a Cauchy sequence if  $\forall \varepsilon > 0, \exists N \text{ s.t. } \forall m, n \geq N, |a_m - a_n| < \varepsilon.$ 

Note that, since the definition of convergent sequence relies on knowing the limit a, when such limit is not attainable, Cauchy becomes handy.

**Theorem 2.12.** Every convergent sequence is Cauchy.

*Proof.* Given  $\{a_n\} \to a$ , thus  $\forall \frac{\varepsilon}{2} > 0 \ \exists N \text{ s.t. } |a_n - a| < \frac{\varepsilon}{2}, \ \forall n > N$ . Now, for any  $m, n \geq N$ , we have:

$$|a_m - a_n| = |a_m - a + a - a_n|$$
  

$$\leq |a_m - a| + |a_n - a| < \varepsilon$$

**Example**: Prove that  $a_{n+1} = \frac{a_n + 2a_{n-1}}{3}$  converges for  $a_1 = 0$ ,  $a_2 = 1$ .

*Proof.* Step 1 First observe that:  $a_n$  is an average of two real numbers that are in [0,1]. Thus,  $a_n \in [0,1]$ .

Step 2 Also observe that by rearranging the terms in the equality, we have:

$$\frac{a_{n+1} - a_n}{a_n - a_{n-1}} = -\frac{2}{3}$$

At this point, we check definition of Cauchy sequence by showing that: for arbitrary  $\varepsilon$ , we can find a N such that  $|a_m - a_n| < \varepsilon$ . Deriving the functional form of  $|a_m - a_n|$  suffices. (We can then use this functional form to find a proper N.)

Without loss of generality, let m > n, then:

$$|a_{m} - a_{n}| = |a_{n} - a_{n+1} + a_{n+1} - \dots - a_{m}|$$

$$\leq |a_{n} - a_{n+1}| + |a_{n+1} - a_{n+2}| + \dots + |a_{m-1} - a_{m}|$$

$$\leq \left(\frac{2}{3}\right)^{n-1} + \left(\frac{2}{3}\right)^{n} + \dots + \left(\frac{2}{3}\right)^{m-2}$$

$$= \frac{\left(\frac{2}{3}\right)^{n-1} \left(1 - \left(\frac{2}{3}\right)^{m-n+2}\right)}{1 - \frac{2}{3}}$$

$$= O\left(\left(\frac{2}{3}\right)^{n}\right)$$

By now, we can easily demonstrate that the definition of Cauchy sequence could be satisfied by choosing a proper N for any given  $\varepsilon$ .

**Lemma 2.13.** Every Cauchy sequence is bounded.

*Proof.* Let  $\{a_n\}$  be an arbitrary Cauchy sequence. Then, for arbitrary  $\varepsilon > 0$ , we know that  $\exists N_{\varepsilon} > 0$  such that  $\forall m, n > N, |a_m - a_n| < \varepsilon$ .

Now, to construct an upper bound for  $\{a_n\}$ , without loss of generality, let  $\varepsilon = 1$ . Then, we know that there exists  $N_1 > 0$  such that  $\forall n, m > N_1$ ,  $|a_n - a_m| < 1$ . Then, let  $M_1$  denote the bound (either upper or lower). Then, in absolute value, we can define it to be:

$$|M_1| = \max\{|a_1|, \dots, |a_{N_1}|, |a_{N_1+1}| + 1\}$$

Through more careful, yet unnecessary, discussions, we can derive the exact bound using the absolute value  $|M_1|$ .

Note that, the bound we found above is only one of the upper bound. It is not necessarily the sup nor inf.  $\Box$ 

**Theorem 2.14.** Every Cauchy sequence in  $\mathbb{R}^2$  converges.

*Proof.* Let  $\{a_n\}$  be an arbitrary Cauchy sequence. We want to show  $\{a_n\}$  converges to some  $a \in \mathbb{R}$ . That is to show:  $\forall \varepsilon > 0, \exists N_0 \in \mathbb{N} \text{ s.t. } \forall n \geq N_0, |a_n - a| < \varepsilon$ .

(Step 1:) For arbitrary  $\varepsilon > 0$ , given that  $\{a_n\}$  is a Cauchy sequence, for  $\frac{\varepsilon}{2} > 0$ ,  $\exists N_1 \in \mathbb{N}$  s.t.

$$|a_m - a_n| < \varepsilon, \quad \forall m, n > N_1$$

<sup>&</sup>lt;sup>2</sup>Note that, for  $\{\frac{1}{n}\}$  defined on (0,1], it does not converge in this space since  $0 \notin (0,1]$ .

(Step 2:) By Lemma 2.13, we know that every Cauchy sequence is bounded. Thus, by Bolzano-Weierstrass Theorem (Theorem 2.9), we know that  $\exists \{a_{n_i}\} \to a$  for some certain real number  $a \in \mathbb{R}$ .

 $\Leftarrow$  a is

By definition of convergence of (sub)sequence, for the arbitrary  $\varepsilon$  that we started with, Limit!  $\exists I \in \mathbb{N} \text{ s.t.}$ 

$$|a_{n_j} - a| < \frac{\varepsilon}{2}, \quad \forall j > I$$

Now, let  $N_0 = \max\{n_I, N_1\}$ , we see that  $\forall n > N_0$  and  $n_i > N_0$ , we have:

$$|a_n - a| = |a_n - a_{n_j} + a_{n_j} - a|$$
  
 $\leq |a_n - a_{n_j}| + |a_{n_j} - a| < \varepsilon$ 

Note:  $a_{n_j}$  is an arbitrary element of the subsequence  $\{a_{n_i}\}$  that we found convergent through B-W Theorem.

Also note that,  $|a_n - a_{n_j}| < \frac{\varepsilon}{2}$  follows from Step 1 that  $\{a_n\}$  is Cauchy to start with.

**Definition 2.15** (Cauchy Criterion). A sequence in  $\mathbb{R}$  is a convergent sequence if and only if it is a Cauchy Sequence.

Demonstration: Theorem 2.12 applies in  $\mathbb{R}$ , thereby convergent sequence in  $\mathbb{R}$  is Cauchy; Theorem 2.14 completes the proof.

**Remark 2.16** (Useful limits). Limits of sequences as  $n \to \infty$ :

- $\lim_{n\to\infty} \frac{n^{\alpha}}{(1+p)^n} = 0$  for p>0 and  $\alpha>0$ . (This demonstrates exponential function dominates polynomials in the limit.)
- $\lim_{n\to\infty} \sqrt[n]{n} = 1$
- $\lim_{n \to \infty} \left(1 + \frac{1}{n}\right)^n = e$ ; then  $\lim_{n \to \infty} \left(1 + \frac{t}{n}\right)^n = e^t$ .
- $\lim_{n\to\infty} \sqrt[n]{p} = 1$  if p > 0.

Refer to page 57 of [Rudin(1976)] Theorem 3.20 for detailed proofs.

**Definition 2.17** (limsup, liminf). Let  $\{a_n\}$  be a sequence in  $\mathbb{R}$ , we say:  $\limsup\{a_n\} = a$  if  $\sup S = a$ , where  $S = \{b \in \mathbb{R} : \exists \text{ subsequence } \{a_{n_i}\} \text{ s.t. } a_{n_i} \to b\}$ .

Not surprisingly, we can define

$$\lim\inf\{a_n\} = -\lim\sup\{-a_n\}$$

**Exercise:** equivalent definition of limsup Prove that  $\limsup a_n = a$  if and only if:

- (i)  $\forall \varepsilon > 0$ ,  $\exists N > 0$  s.t.  $a_n < a + \varepsilon$ ,  $\forall n > N$ ;
- (ii)  $\forall \varepsilon > 0, \forall n \in \mathbb{N}, \exists k > n \text{ s.t. } a_k > a \varepsilon.$

Note that, (i) specified a property for subsequence; and (ii) is merely about the existence of one element in the sequence, to be found for all  $(\varepsilon, n) \in \mathbb{R}_{++} \times N$ .

*Proof.* The iff statement will be established in the following three steps:

• Prove that  $\limsup a_n = a$  implies (i).

WTS:  $\forall \varepsilon > 0$ ,  $\exists N > 0$  s.t.  $a_n < a + \varepsilon$ ,  $\forall n > N$ ;

First, suppose that  $a = +\infty$ , that is  $\{a_n\}$  is not bounded from above. Then we are done. Then, suppose that  $\{a_n\}$  is bounded from above. We now prove by contradiction. Suppose that  $\exists \varepsilon > 0$  s.t. no such  $N \in \mathbb{N}$  exists. Then, we know that 1 cannot serve the role of N. So, for some  $n_1 > 1$ ,

$$a_{n_1} \ge a + \varepsilon$$

Still,  $n_1 + 1$  cannot serve the role of N, then for some  $n_2 > n_1 + 1$ ,

$$a_{n_2} \geq a + \varepsilon$$

By induction, we can construct a subsequence that is bounded from below by  $a + \varepsilon$ . Note that, the original sequence is bounded from above, by Bolzano-Weierstrass Theorem, we know that a bounded sequence converges. However, the limit of such subsequence shall be larger than a, contradicting  $\limsup a_n = a$ .

Thus what we assumed is wrong. We thereby proved the original claim in (ii).

Note that, the converse of the following claim:

 $\forall \varepsilon > 0, \quad \exists N \in \mathbb{N} \text{ s.t. } \forall n \geq N, \text{ proposition } P \text{ is true.}$ 

 $\exists \varepsilon > 0, \forall N \in \mathbb{N} \text{ s.t. } \exists n \geq N, \text{ proposition } P \text{ is } \mathbf{not} \text{ true.}$ 

• Prove that  $\limsup a_n = a$  implies (ii).

WTS: Given that  $\limsup a_n = a, \forall \varepsilon > 0, \forall n \in \mathbb{N}, \exists k > n \text{ s.t. } a_k > a - \varepsilon.$ 

Now, for arbitrary  $\varepsilon > 0$ , by definition of limsup, we know that  $\exists a' \in (a - \frac{\varepsilon}{2}, a)$  s.t.  $\exists \{a_{n_j}\}$  (a subsequence of  $\{a_n\}$ ) s.t.  $a_{n_j} \to a'$ .

For this convergent subsequence per se, given the arbitrary  $\varepsilon$  we have specified in the very beginning, we know that  $\exists J > 0$  s.t.

$$|a_{n_j} - a'| < \frac{\varepsilon}{2},$$
 for all  $j > J$ 

Now, for arbitrary  $n \in N$ , we can always find a  $k = n_i$  with i > J, such that  $a_k = a_{n_i}$  is within  $\frac{\varepsilon}{2}$  distance away from a'. Combining this fact with the construction that  $a' \in (a - \frac{\varepsilon}{2}, a)$ , it is clear the  $a_k$  we found specifically for  $\varepsilon$  and  $n \in N$  satisfies:  $a_k > a - \varepsilon$ .

• Prove that (i) and (ii) implies that  $\limsup a_n = a$ .

To prove that  $\limsup a_n = a$ , we first show that a is the limit of a subsequence of  $\{a_n\}$ ; then we show that  $\not\exists a' > a$  s.t. a' is the limit of a subsequence of  $\{a_n\}$ .

Firstly, by (i) and (ii), for arbitrary  $\varepsilon > 0$ , we can find a subsequence  $\{a_{n_j}\}$  with certain  $N \in \mathbb{N}$  such that  $a - \varepsilon < a_{n_j} < a + \varepsilon$ ,  $\forall n_j > N$ . (Step 1: by (i), we can find a  $N^{\varepsilon}$  for

arbitrary  $\varepsilon > 0$ , so that:  $a_n < a + \varepsilon \ \forall n > N^{\varepsilon}$ ; Step 2, for the  $\varepsilon$  and all  $\tilde{n} \geq N^{\varepsilon}$ , we can find a  $a_{k_{\tilde{n}}}$  s.t.  $a - \varepsilon < a_{k_{\tilde{n}}}$ . Thus, we have composed a subsequence  $\{a_{k_{\tilde{n}}}\}$ .)

Then, suppose  $\exists a' > a$  as the limsup, then  $\forall \varepsilon > 0 \ \exists N'$  s.t.  $\forall n' > N'$ ,  $|a_{n'} - a'| < \varepsilon$ . However, (i) is violated when  $\varepsilon < \frac{a'-a}{2}$ : suppose that  $a_{n_k} \to a'$ . Then,  $\exists N' > 0$  s.t.  $\forall k > N'$ ,  $|a_{n_k} - a'| < \varepsilon$ . Given that  $\varepsilon < \frac{a'-a}{2}$ , there does not exist a N that may satisfy (i). (The " $\forall n > N$ " statement is violated due to the subsequence that converges to a'.)

Alternatively, one can prove the statement using an equivalent definition of limsup:

$$\limsup_{n \to \infty} a_n = \lim_{n \to \infty} \sup_{k \ge n} a_k$$

Thus, (i) implies that  $\forall \varepsilon > 0, \exists N \text{ s.t. } \forall n > N$ :

$$\sup_{k > n} a_k < a + \varepsilon$$

Therefore,  $\limsup_{n \to \infty} a_n < a + \varepsilon \iff \limsup_{n \to \infty} a_n \le a$ ;

At the same time, (ii) implies that  $\forall \varepsilon > 0$ , for arbitrary  $n \in N$ ,  $\exists k > n$  s.t.  $a_k > a - \varepsilon$ . Then:

$$\sup_{j > n} a_j > a - \varepsilon$$

Therefore,  $\limsup_{n\to\infty} a_n > a - \varepsilon \iff \limsup_{n\to\infty} a_n \ge a$ ;

**Definition 2.18** (Infinite series). Given a sequence  $\{a_n\}$ , let  $s_n = \sum_{i=1}^n a_i$  be a sequence  $\{s_n\}$ , it is called **infinite series**. We write  $\sum_{n=1}^{\infty} a_n = a$  if  $\{s_n\}$  converges to a.

**Example 4.** For  $a_n = \frac{1}{2^n}$ , we can obtain an expression for  $\sum_{n=1}^M a_n$ ; and  $\sum_{n=1}^\infty \frac{1}{n} = \infty$ . Also note that the sum of arbitrary segment of  $\{\frac{1}{n}\}$  can be arbitrarily large if the length of such segment is long enough.

**Definition 2.19** (Rearrangement).  $\{n_i\}_{i=1}^{\infty}$  is a sequence of natural numbers in which each natural number appears exactly once. Let  $b_i = a_{n_i}$ , then  $b_i$  is a **rearrangement** of  $\{a_i\}_{i=1}^{\infty}$ .

**Definition 2.20** (Absolute convergence). If  $\sum_{n=1}^{\infty} |a_n|$  converges, we say that  $\sum_{n=1}^{\infty} a_n$  converges absolutely. (e.g. for  $a_n = (-1)^n \frac{1}{n}$ ,  $\sum_{n=1}^{\infty} < \infty$ , yet  $\sum_{n=1}^{\infty} |a_n| \to \infty$ .)

**Proposition 2.21.** If  $\sum_{n=1}^{\infty} a_n$  converges absolutely, then  $\sum_{n=1}^{\infty} b_i = \sum_{n=1}^{\infty} a_n$ , where  $\{b_n\}$  is a rearrangement of  $\{a_n\}$ .

Note that, rearranging  $\{(-1)^n\}_{n=1}^{\infty}$  can give raise to arbitrary partial sum  $\in \mathbb{Z}$ .

**Review:** subsets in  $\mathbb{R}$  Epistemic-wise, we established the construction of following sets sequentially:

- 1. N: The set of natural number; [It is countable.]
- 2.  $\mathbb{Z}$ : The set of integers; [It is also countable. In fact,  $|\mathbb{N}| = |\mathbb{Z}|$ .

- 3. Q: The set of rational number; [It is also countable, and dense.]
- 4.  $\mathbb{R}$ : The real line. [Completeness Axiom]

**Definition 2.22** (Principle of Mathematical Induction). The set of natual numbers is the smallest set that satisfies the axiom of Mathematical Induction.

**Example 5.** Prove that  $1 + 2 + \cdots + n = \frac{n(n+1)}{2}$ .

*Proof method:* To prove by induction:

- When n = 1, LHS = RHS;
- Suppose that, for some  $n_0 \in \mathbb{N}$ , LSH = RHS  $\forall n \in \mathbb{N}$  s.t.  $n \leq n_0$ , then we show that LHS = RHS for  $n = n_0 + 1$ .

2.2 Real Value Functions

**Definition 2.1.** A real valued function defined on X (an arbitrary set) is represented as following, with  $\mathbb{R}$  as the codomain:

$$f: x \to \mathbb{R}$$

**Notation 2.2.** For  $a \in \mathbb{R}$  and f, g being real value functions, "=,  $\geq$ , >,  $\gg$ , function addition and (scalar) multiplication" are defined as follows:

- If  $f(x) = a \ \forall x \in X$ , we write f = a;
- If  $f(x) \ge g(x) \ \forall x \in X$ , we write  $f \ge g$ ;
- If  $f \ge g$ , but not the other way, then f > g. (f(x) = g(x)) is permissible for some  $x \in X$ ).
- If  $f(x) > g(x) \ \forall x \in X$ , then we write  $f \gg g$ .
- (f+g)(x) := f(x) + g(x);
- $(a \cdot f)(x) \coloneqq a \cdot f(x);$
- $(f \cdot g)(x) := f(x) \cdot g(x)$ .

Note that, f > g is a "weakly hight" relationship.

**Definition 2.3** (strictly/weakly increasing/decreasing). Construction is intuitive and thereby omitted.

**Definition 2.4** (Limit of function). A function  $f: x \to \mathbb{R}$  converges to  $a \in \mathbb{R}$  as x approaches some  $x_0 \in X$  if

$$\forall \varepsilon > 0, \ \exists \delta > 0 \text{ s.t. } \forall x \in (x_0 - \delta, x_0 + \delta)$$
$$|f(x) - a| < \varepsilon$$

in which case we write  $\lim_{x\to x_0} f(x) = a$ .

**Definition 2.5** (Right limit). A function  $f: x \to \mathbb{R}$  converges to  $a \in \mathbb{R}$  from right as x approaches  $x_0$  if

$$\forall \varepsilon > 0, \ \exists \delta > 0 \text{ s.t. } \forall x \in (\underline{x_0}, x_0 + \delta),$$
$$|f(x) - a| < \varepsilon$$

We write the right limit as:  $\lim_{x \to x_0^+} f(x) = a$ .

**Proposition 2.6.** Suppose  $f: X \to \mathbb{R}$  and  $g: X \to \mathbb{R}$ , with  $\lim_{x \to x_0} f(x) = a$  and  $\lim_{x \to x_0} g(x) = b$ .

- (i)  $\lim_{x \to x_0} f(x) \pm g(x) = a + b;$
- (ii)  $\lim_{x \to x_0} f(x) \cdot g(x) = a \cdot b;$
- (iii)  $\lim_{x \to x_0} \frac{f(x)}{g(x)} = \frac{a}{b}$  if  $g \neq 0$  and  $b \neq 0$ .

**Definition 2.7** (Continuity). A function  $f: X \to \mathbb{R}$  is continuous at  $x_0 \in X$  if

$$\lim_{x \to x_0} f(x) = f(x_0)$$

Note that, we can draw definition of *limit of function* to formalize an  $\varepsilon - \delta$  argument that defines a continuous function.

# 3 Lecture 3: Linear Space and $f: X \to \mathbb{R}$ , continued

## 3.1 Linear spaces and linear algebra

**Definition 3.1** (Vector Space). **Vector space** V over a field F is a set V together with vector addition and scalar multiplication.

ullet A field F is a set with addition and multiplication operation defined among its own elements.

Example:  $\mathbb{R}$  with normal + and  $\cdot$  is a field, denoted as: " $F : \mathbb{R}, +, \cdot$ ".

Formally, a field is also established using a set of axiom. Note that field is "equipped with": 0, (-1) elements.

Axiomatically,  $\forall u, v, w \in V$  and  $a, b \in F$ , the following shall be satisfied:

**Axiom** 1 
$$u + (v + w) = (u + v) + w;$$

**Axiom** 2 u + v = v + u;

**Axiom** 3  $\exists \theta \in V \text{ s.t. } u + \theta = u;$ 

**Axiom** 4  $\exists \phi(u) \in V \text{ s.t. } u + \phi(u) = \theta;$ 

**Axiom** 5  $a \cdot (u + v) = a \cdot u + a \cdots v$ ;

**Axiom** 6  $(a+b) \cdot u = a \cdot u + b \cdot u$ ;

**Axiom** 7  $a \cdot (b \cdot u) = b \cdot (a \cdot u)$ ;

**Axiom** 8 V is closed under vector addition and scalar multiplication;

**Axiom** 9  $1 \cdot u = u$ , where 1 is the identity in F.

Note that, the last axiom was not stated in lecture.

**Proposition 3.2.** Using the axioms, we can show the following equalities hold:

- 1.  $0 \cdot u = \theta$ ;
- 2.  $\phi(u) = (-1) \cdot u;$
- 3.  $a\theta = \theta$ ;
- 4.  $\theta$  is unique.

*Proof.* Relies heavily on algebraic tricks. Omitted as of 2015-08-29 15:01:15.

**Example 6** (Example for vector spaces). 1.  $V = \mathbb{R}^n$  and  $F : R, +, \cdot$ ; 2.  $V = \{ax^2 + bx + c : a, b, c \in \mathbb{R}, x \in [0, 1]\}$ , for  $F : \mathbb{R}, +, \cdot$ .

**Definition 3.3.** A vector space cna also be called a linear space.

**Definition 3.4** (Linear subspace). For V being a linear space and  $U \subseteq V$ , if U itself is a linear space with <u>the same</u> vector additions and scalar multiplication, then we say U is a linear subspace of V.

Note that, this definition admits the case where U = V, i.e. though trivially, V is a linear subspace of itself.

### 3.1.1 \*Finite\* Linear combination, span and linear independence of vectors

From now on, we limit the discussion to the following case:

- 1. Adopt  $\mathbb{R}$  with normal addition and multiplication to be the field F;
- 2. Consider only finite operations when defining linear combination and span;
- 3. Note that: it is still permissible for V to be an arbitrary set.

**Definition 3.5** (Linear Combination). For  $U \subseteq V$ ,

(i) If  $U = \{v_1, \ldots, v_n\}$  for some  $n \in \mathbb{N}$ , i.e. U is a finite subset of V, then a linear combination of U is a new vector:

$$v = \sum_{i=1}^{n} a_i v_i, \ a_i \in \mathbb{R}, \ i = 1, \dots, n$$

(ii) If U is no longer finite, regardless of whether is is countably infinite or uncountable, a **linear combination of** U is a vector that is a linear combination of finitely many vector of U.

**Definition 3.6** (span of a set of vectors). For  $A = \{v_1, \dots, v_n\}$ ,

$$span(A) = \{ \sum_{i=1}^{n} a_i v_i : a_i \in \mathbb{R}, i = 1, \dots, n \}$$

**Proposition 3.7.** The span of any  $U \subset V$  is a linear subspace of linear space V.

Scatch of proof. Note that, by construction of span(A), arbitrary coefficient is allows. Letting all coefficients to be 0 gives raise to the  $\theta$ ; other properties may follow from standard algebra in  $\mathbb{R}$  (the field).

**Definition 3.8** (Linear independence). A (finite) set of vectors A is linearly independent if  $\exists v \in A$  can be written as linear combinations of the others. Formally,

$$A = \{v_1, \dots, v_n\}$$
 is linearly independent if  $\sum_{i=1}^n a_i v_i = \varepsilon \implies a_i = 0 \forall i$ 

*Proof.* Suppose not, that is  $\sum_{i=1}^n a_i v_i = \varepsilon$  yet  $a_j \neq 0$  for some j, then we can write:

$$-a_j v_j = \sum_{k \neq j} a_k v_k$$

where, upon simplification,  $v_j$  could be written as a linear combination of the other vectors.

**Proposition 3.9.** For  $A \subseteq V$ , span(A) is the smallest linear space that contains A.

Alternatively, one can define span(A) to be the intersection of all linear subspaces of V that contains A.

**Definition 3.10** (base and dimension of V). If  $\{v_1, \ldots, v_n\}$  are linearly independent and  $span(\{v_1, \ldots, v_n\}) = V$  (the linear space), then  $\{v_1, \ldots, v_n\}$  is called a **base of** V.

In this case, the **dimension** of V is dim(V) = n.

**Theorem 3.11.** If A and B are two bases of V and A, B are finite, then |A| = |B|.

Idea of the proof. Suppose  $A = \{u_1, u_2\}$  and  $B = \{v\}$ . Then one can write:

$$u_1 = av;$$
  $u_2 = bv$  for some  $a, b \in \mathbb{R}$ 

Therefore,  $u_1$  and  $u_2$  are not linearly independent.

Note that, it seems to me that the finiteness assumption only serves the need of simplifying the proof.

#### 3.1.2 Matrix

**Definition 3.12.** A  $m \times n$  matrix could be written as:

$$A = \begin{bmatrix} v_1 \\ \vdots \\ v_m \end{bmatrix} = [u_1 \dots u_n]$$

where  $v_i$  is a  $1 \times n$  (row) vector, and  $u_i$  is a  $m \times 1$  (column) vector.

**Definition 3.13** (Rank of a matrix). The *maximum number* of linearly independent row/column vectors denotes the rank of a matrix.

Comment: implicitly, by definition,  $rank(A) = rank(A^T)$ .

**Definition 3.13** (Linear transformation).  $T:U\to V$  is a linear transformation if

$$T(au_1 + bu_2) = aT(u_1) + bT(u_2), \quad \forall a, b \in \mathbb{R}$$

**Remark 3.14.** A  $m \times n$  matrix A is a linear transformation from  $\mathbb{R}^n \to \mathbb{R}^m$ .

#### 3.2 Real-Valued functions continued

#### 3.2.1 Continuity and its corollaries

**Definition 3.1** (Interval). An interval of  $\mathbb{R}$  is either [a,b], (a,b], [a,b) or (a,b); where  $a,b \in \mathbb{R} \bigcup \{+\infty, -\infty\}$  (the extended real line).

**Theorem 3.2** (Intermediate Value Theorem). If I is an interval of  $\mathbb{R}^{1}$ , and  $f: I \to \mathbb{R}$  is continuous, then f(I) is also an interval of  $\mathbb{R}^{2}$ .

 $<sup>^{1}</sup>I$  could be a connected set in Euclidean space ( $\mathbb{R}^{n}$ ).

<sup>&</sup>lt;sup>2</sup>Correspondingly, f(I) would be a connected set.

**Proposition 3.3.** If f is continuous and bijective (thus invertible, i.e.  $f^{-1}$  is a function), ten f is either strictly increasing or strictly decreasing.

Note that:

- Continuity forced bijections to be monotone;
- "Strictness" is used to support bijection;
- A stronger statement (yet correct) goes as follows:

Let I and J be both intervals, then  $f: I \to J$  is continuous and bijective if and only if it is strictly monotonic.

**Theorem 3.4** (Extreme Value Theorem). If  $f : [a, b] \to \mathbb{R}$  is continuous, then  $\exists x_1, x_2 \in [a, b]$  s.t.

$$f(x_1) = \sup f([a, b])$$

$$f(x_2) = \inf f([a, b])$$

Comment: using max and inf in the statement would be more precise though.

**Definition 3.5** (Uniformly continuity).  $f: x \to \mathbb{R}$  is said to be uniformly continuous if  $\forall \varepsilon > 0, \exists \delta > 0$ , s.t.

$$|f(x) - f(y)| < \varepsilon, \qquad \forall |x - y| < \delta$$

Note that:

- 1. We no longer specify a certain point  $x_0 \in X$ ;
- 2. Instead, the  $\delta$  applies to all  $x, y \in X$  as long as they are within  $\delta$  distance away.

**Exercise 3.6.** Prove that  $f(x) = \frac{1}{x}$  (x > 0) is not uniformly continuous.

*Professor's Proof.* Without loss of generality, suppose  $\varepsilon = \frac{1}{2}$ . Now we want ot show that  $\not \exists \delta > 0$  s.t. if  $|x - y| < \delta$ ,  $|f(x) - f(y)| < \frac{1}{2}$ .

By the property of f(x), we look for a threshold  $z^*(\varepsilon, \delta)$  at which:

$$\left| \frac{1}{z^*} - \frac{1}{z^* + \delta} \right| = \frac{1}{2}$$

Then, for arbitrary  $\delta > 0$ , write  $z^* = z^*(\varepsilon, \delta)$ , we have:

$$|f(z') - f(z' + \delta)| > \frac{1}{2}, \quad \forall z' < z^*$$

Thus, we see that for  $\varepsilon = \frac{1}{2}$ , there does not exist a  $\delta > 0$  that satisfies  $|f(x) - f(y)| < \frac{1}{2}$   $\forall |x - y| < \delta$ .

Comment: in professor's proof, there is a flaw: choosing z' and  $z' + \delta$  won't help disprove the original statement. This could easily be fixed as shown in the alternative proof.

Alternative proof. Without loss of generality, suppose  $\varepsilon = \frac{1}{2}$ . Now we want ot show that  $\beta > 0$  s.t. if  $|x - y| < \delta$ ,  $|f(x) - f(y)| < \frac{1}{2}$ .

By the property of f(x), we look for a threshold  $z^*(\varepsilon, \delta)$  at which:

$$\left| \frac{1}{z^*} - \frac{1}{z^* + \frac{\delta}{2}} \right| = \frac{1}{2}$$

Then, for arbitrary  $\delta > 0$ , write  $z^* = z^*(\varepsilon, \delta)$ , we have:

$$|f(z') - f(z' + \frac{\delta}{2})| > \frac{1}{2}, \quad \forall z' < z^*$$

Thus, we see that for  $\varepsilon = \frac{1}{2}$ , there does not exist a  $\delta > 0$  that satisfies  $|f(x) - f(y)| < \frac{1}{2}$   $\forall |x - y| < \delta$ .

Note that, it is the highlighted condition that has been disproved.

3.3 Differentiation

Remark 3.1. "Differentiation" is essentially a process of taking linear approximation.

**Definition 3.2** (Tangent line). The tangent line to a function y = f(x) at the point  $(x_0, f(x_0))$ , when exists, is the line through  $(x_0, f(x_0))$  with slope

$$\alpha = \lim_{x \to x_0} \frac{f(x) - f(x_0)}{x - x_0} = \lim_{\Delta x \to 0} \frac{f(x + \Delta x) - f(x)}{\Delta x}$$

When  $\alpha$  exists, the tangent line exists. It could be written as:

$$y = f(x_0) + \alpha(x - x_0)$$

**Definition 3.3** (Differentiation). The derivative of  $f: x \to \mathbb{R}$  at  $x_0 \in X$  is

$$f'(x) = \lim_{x \to x_0} \frac{f(x) - f(x_0)}{x - x_0}$$

We may also write the derivative as:

$$\left(\frac{df(x)}{dx} \middle|_{x=x_0}\right)$$

The **derivative** of f is denoted by

$$\frac{df(x)}{dx}$$

**Remark 3.4** (Properties of derivatives). For f, g as functions and  $a, b \in \mathbb{R}$ :

- (i) (af + bg)' = af' + bg'
- (ii) (fg)' = f'g + fg'

(iii) 
$$\left(\frac{f}{g}\right)' = \frac{f'g - fg'}{g^2}$$

**Proposition 3.5** (Chain Rule). If  $g: x \to \mathbb{R}$  is differentiable at  $x_0 \in X$  and  $f: Y \to \mathbb{R}$  is differentiable at  $g(x_0) \in Y$ , then f(g(x)) is differentiable at  $x = x_0$ , we write:

$$\left(\frac{df(g(x))}{dx}\Big|_{x=x_0}\right) = f'(g(x_0))g'(x_0)$$

**Proposition 3.6** (Inverse function theorem). If  $f: X \to Y$  is bijective, then derivative of  $f^{-1}: Y \to X$  is

$$\frac{df^{-1}(y)}{dy} = \frac{1}{f'(f^{-1}(y))}$$

*Proof.* Since f is bijective function, we have:  $f(f^{-1}(y)) = y$ . Differentiating w.r.t.<sup>3</sup> y gives:

$$f'(f^{-1}(y)) \cdot \frac{df^{-1}(y)}{dy} = 1 \implies \frac{df^{-1}(y)}{dy} = \frac{1}{f'(f^{-1}(y))}$$

**Definition 3.7** (local maximum). Function  $f: X \to \mathbb{R}$  has a local maximum at  $x_0 \in X$  if  $\exists \delta > 0$ , s.t.

$$f(x_0) \ge f(x), \quad \forall x \in \{x \in X : |x - x_0| < \delta\}$$

**Proposition 3.8** (Condition for interior local maximum). If  $f: X \to \mathbb{R}$  is differentiable, and has a local maximum at an interor point  $x = x_0^4$ , then  $f'(x_0) = 0$ 

*Proof.* First consider the right limit:  $\lim_{\Delta \to 0^+} \frac{f(x_0 + \Delta x) - f(x_0)}{\Delta x}$ . For  $\Delta x > 0$  and  $f(x_0 + \Delta x) - f(x_0) \le 0$  (by local maximum), we see:

$$\lim_{\Delta \to 0^+} \frac{f(x_0 + \Delta x) - f(x_0)}{\Delta x} \le 0$$

In similar spirit, we conclude that:

$$\lim_{\Delta \to 0^{-}} \frac{f(x_0 + \Delta x) - f(x_0)}{\Delta x} \ge 0$$

Thus we conclude that  $f'(x_0) = 0$  due to differentiability of f at  $x_0$ .

Note that, if  $x_0$  is at the boundary of X, whether this proposition holds (or not) depends on how we define the derivative at the boundary point.

**Theorem 3.9** (Rolle's Theorem). If  $f:[a,b]\to\mathbb{R}$  is differentiable and f(a)=f(b)=0, then  $\exists x_0\in(a,b)$  s.t.  $f'(x_0)=0$ .

<sup>&</sup>lt;sup>3</sup>with respect to

 $<sup>^4</sup>x_0$  is an interior point of X, i.e.  $\exists \delta > 0$  s.t.  $\{x \in X : |x - x_0| < \delta\} \subseteq X$ 

*Proof.* Since  $f:[a,b] \to \mathbb{R}$  is differentiable and hence continuous, if  $\sup f([a,b]) > 0$ , then we can locate a  $x_0$  as local maximum. Then, by the previous proposition,  $f'(x_0) = 0$ ;

Alternatively, if  $\inf(f[a,b]) < 0$ , we can find a  $x_1$  as local minimum. This also gives raise that  $f'(x_1) = 0$ .

Otherwise, f is flat, and 
$$f'(x) = 0 \ \forall x \in [a, b].$$

Note that, this is like reaching a plateau/basin when leaving at sea-level and reaching another point at sea-level.

**Theorem 3.10** (Mean Value Theorem). If  $f:[a,b] \to \mathbb{R}$  is differentiable, then  $\exists x_0 \in (a,b)$  s.t.

$$f'(x_0) = \frac{f(b) - f(a)}{b - a}$$

*Proof.* Subtract a line function:  $y = f(a) - \frac{f(b) - f(a)}{b - a}(x - a)$  from f(x) to get g(x), we can apply Rolle's Theorem and find a  $x_0$  that satisfies  $g'(x_0) = 0$ .

Note that, one can envision subtracting a line-function as a transformation of coordinate system.  $\Box$ 

**Theorem 3.11** (Generalized Mean Value Theorem). Let  $f:[a,b] \to \mathbb{R}$ ,  $g:[a,b] \to \mathbb{R}$  be both differentiable, then  $\exists x_0 \in [a,b]$  s.t.

$$g'(x_0)(f(b) - f(a)) = f'(x_0)(g(b) - g(a))$$

Note that, we can rationalize this theorem as: the ratio of average speed shall equal the ratio of travel speed at some point of time<sup>5</sup>.

**Theorem 3.12** (L'Hopital Rule). f and g are differentiable, with  $g'(x) \neq 0$ ,  $\forall x \in X$ . Suppose  $\lim_{x \to x_0} \frac{f'(x)}{g'(x)} = q$ , then if either of the following conditions is satisfied,  $\lim_{x \to x_0} \frac{f(x)}{g(x)} = q$ .

- (i) If  $f(x), g(x) \to 0$  as  $x \to x_0$ ;
- (ii) If  $f(x), g(x) \to \infty$  as  $x \to x_0$ .

**Proposition 3.13** (Derivative is continuous at  $x_0$ ). If  $\lim_{x\to x_0} f'(x)$  exists, then  $f'(x_0) = \lim_{x\to x_0} f'(x)$ .

*Proof.* Define  $h(x) = f(x) - f(x_0)$ , we see that  $h(x) \to 0$  as  $x \to x_0$ ; then, define  $g(x) = x - x_0$ , we also see that  $g(x) \to 0$  as  $x \to x_0$ .

Thus, by L' Hopital's rule, we have:

$$\lim_{x \to x_0} \frac{h(x)}{g(x)} = \lim_{x \to x_0} \frac{h'(x)}{g'(x)} = \frac{f'(x)}{1} = f'(x)$$

Note that, what we started with is by definition  $f'(x_0) = \lim_{x \to x_0} \frac{f(x) - f(x_0)}{x - x_0}$ . So, we are done.  $\square$ 

<sup>&</sup>lt;sup>5</sup>Though, Prof Ke did not specify which one is the "time variable".

## 4 Lecture 4: Differentiation and Linear Algebra

### 4.1 Differentiation

**Definition 4.1** (derivative). If f has a derivative f', and f' itsef is also differentiable, then we write the derivative of f', f''. The we also derive the other higher order derivatives:

$$f', f'', f^{(3)}, f^{(4)}, \dots, \frac{d^n f(x)}{dx^n} = f^{(n)}$$

**Theorem 4.2.** For  $f:[a,b]\to\mathbb{R}$ , suppose  $f^{(n-1)}$  is continuous (i.e.  $f\in C^{n-1}$ ), and  $f^{(n)}$  exists. Then for  $x_0, \bar{x}\in[a,b], \exists \tilde{x}\in(\min\{x_0,\bar{x}\},\max\{x_n,\bar{x}\})$  s.t.

$$f(\bar{x}) = f(x_0) + f'(x_0)(\bar{x} - x_0) + \frac{1}{2}f''(x_0)(\bar{x} - x_0)^2 + \frac{1}{3!}f^{(3)}(x_0)(\bar{x} - x_0)^3 + \dots + \frac{1}{(n-1)!}f^{(n-1)}(x_0)(\bar{x} - x_0)^{n-1} + \frac{1}{n!}f^{(n)}(\tilde{x})(\bar{x} - x_0)^n$$

Alternatively, one can write:

$$f(\bar{x}) = \sum_{k=0}^{n-1} \frac{f^{(k)}(x_0)}{k!} (\bar{x} - x_0)^k + \frac{1}{n!} f^{(n)}(\tilde{x}) (\bar{x} - x_0)^n$$

*Proof.* This is a direct proof, which relies on Mean Value Theorem and a tricky construction.

**Remark 4.4** (Limitation of Taylor series expansion). If  $f^{(n)}$  exists for all  $n \in \mathbb{N}$ , can we write down the Taylor series as follows?

$$f(\bar{x}) \stackrel{?}{=} \sum_{k=0}^{\infty} \frac{f^{(k)}(x_0)}{k!} (\bar{x} - x_0)^k$$

It turns out that:

- (i) RHS may not necessarily converge: for  $x_0 = 0$ ,  $f(x) = \frac{1}{1-x} \implies RHS = \sum_{k=0}^{\infty} x^k$ . If |x| > 1, RHS does not converge.
- (ii) RHS converges but LHS  $\neq$  RHS: Let  $e^{-\frac{1}{0}} = 0$ , then  $e^{-\frac{1}{x^2}}$  has a Taylor series expansion at  $x_0 = 0$ , yet RHS = 0.
- (iii) The equality may hold, for example:

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!} \implies \text{LHS} = \text{RHS}, \forall x$$

#### 4.1.1 Multi-variable Calculus

**Definition 4.5** (Differentiable multi-variate vector-valued function). For  $f: X \to \mathbb{R}^m$ ,  $X \subseteq \mathbb{R}^n$ , we say that f is differentiable at  $x_0 \in X$  if  $\exists A \ (m \text{ by } n \text{ matrix}) \text{ s.t.}$ 

$$\lim_{x \to x_0} \frac{||f(x) - [f(x_0) + A \cdot (x - x_0)]||}{||x - x_0||} = 0$$

where the norm of x is defined as:  $||x|| = \sqrt{x_1^2 + x_2^2 + \dots + x_n^2}$ .

Note that,  $f(x_0) + A \cdot (x - x_0)$  is a local approximation of f(x) at  $x_0$ . Here,  $f(x_0)$  is a  $m \times 1$  vector,  $x - x_0$  is a  $n \times 1$  vector, and A is a linear transformation from  $\mathbb{R}^n$  to  $\mathbb{R}^m$ .

#### Exercise: L'Hopital's Rule

- $\lim_{x\to 0} \frac{e^x-1}{x} = \lim_{x\to 0} \frac{e^x}{1} = 1$  (By L'Hopital's Rule).
- $\lim_{x\to 0} x \cdot ln(1+\frac{1}{x}) = \lim_{x\to 0} \frac{ln(1+\frac{1}{x})}{\frac{1}{x}} = \lim_{x\to 0} \frac{\frac{1}{1+\frac{1}{x}}\cdot\left(-\frac{1}{x^2}\right)}{-\frac{1}{x^2}} = \lim_{x\to 0} \frac{1}{1+\frac{1}{x}} = 0$  (Again, by L'Hopital's Rule)

#### 4.1.2 Partial Derivatives

**Definition 4.6** (Partial Derivatives). For  $x : x \to \mathbb{R}$  with  $x \in \mathbb{R}^n$ , the partial derivative of f at  $x = (x_1, \dots, x_n)$  with respect to  $x_i$  is defined as:

$$\frac{\partial f(x)}{\partial x_i} := f_i(x) = \lim_{\lambda \to 0} \frac{f(x + \lambda e_i) - f(x)}{\lambda}$$

where  $e_i = [0, \dots, 0, 1, 0, \dots]^T \in \mathbb{R}^n$  (1 is at the *i*-th element).

**Definition 4.7**  $(\partial x_i \partial x_j)$ . If we take the partial derivative of  $\frac{\partial f(x)}{\partial x_i}$  with respect to  $x_j$ , we have:

$$\frac{\partial \left(\frac{\partial f(x)}{\partial x_i}\right)}{\partial x_j} \coloneqq \frac{\partial^2 f(x)}{\partial x_i \cdot \partial x_j} \coloneqq f_{ij}$$

**Proposition 4.8.** If  $f_i$ ,  $f_j$  and  $f_{ij}$  exists, and  $f_{ij}$  is continuous at  $x_0 \in x$ , then  $f_{ji}$  exists at  $x_0$  and  $f_{ij} = f_{ji}$ .

#### 4.1.3 Gradient

**Definition 4.9** (Gradient  $\nabla$ ). Remember the matrix A we defined in Definition 4.5? Let  $A = \nabla f(x)$ , then  $\nabla f(x)$  is called the gradient of f at x, i.e.  $\nabla f(x)$  is a  $m \times n$  matrix s.t.

$$\lim_{x \to x_0} \frac{||f(x) - (f(x_0) + \nabla f(x_0)(x - x_0))||}{||x - x_0||}$$

It turns out that we are essentially defining  $\nabla f(x_0)$  as a  $m \times n$  matrix without transpose? Note:  $\nabla f(x_0) = [f_1(x_0), f_2(x_0), \dots, f_n(x_n)]$ . (Simon and Blume page 321 defined gradient to be a  $n \times m$  matrix.

**Definition 4.10** (Inner product). In  $\mathbb{R}^2$ , the inner product of  $(x_1, y_1)$  and  $(x_2, y_2)$  is

$$\langle (x_1, y_1), (x_2, y_2) \rangle := (x_1, y_1) \cdot \begin{pmatrix} x_2 \\ y_2 \end{pmatrix} = x_1 x_2 + y_1 y_2$$

Additionally, one can define the norm as  $||(x,y)|| = \sqrt{\langle (x,y),(x,y) \rangle}$ . Further more, one can write:

$$\langle (x_1, y_1), (x_2, y_2) \rangle = ||(x_1, y_1)|| \cdot ||(x_2, y_2)|| \cdot \cos \theta$$

where  $\theta$  is the angle from vector  $(x_1, y_1)$  to  $(x_2, y_2)$ .

**Remark 4.11** (Tangential plane example). For  $f : \mathbb{R}^2 \to \mathbb{R}$ , let  $X \times Y = \mathbb{R}^2$  be the domain and  $Z = \mathbb{R}$  be the codomain. Then, we can define a tangential plane at  $x_0 = (x', y')$  as follows:

• Fix the y component, we can define the following plane that is parallel to Z-X plane:

$$z = f(x_0) + \frac{\partial f(x_0)}{\partial x}(x - x')$$

• Fix the x component, we can define the following plane that is parallel to Z-Y plane:

$$z = f(x_0) + \frac{\partial f(x_0)}{\partial y}(y - y')$$

Then, the tangential plane is defined as:

$$f(x_0) + \left(\frac{\partial f(x_0)}{\partial x}, \frac{\partial f(x_0)}{\partial y}\right) \cdot (x - x_0)$$

**Projection** (Continued from the previous  $f: \mathbb{R}^2 \to \mathbb{R}$  function) In  $X \times Y$ , we can draw a set of contour-lines at which f(x) achieves the same Z-value.

Under the projection, the tangential plane we defined previous now becomes a tangent line. More importantly, the projection of *gradient vector* is orthogonal to such tangent line in X - Y.

In general, the gradient is also orthogonal to the tangent plane.

**Proposition 4.12.** The gradient is the direction along which f(x) increases the fastest.

**Implication of gradient** Formally, Remark 4.11 could be stated as follows:

•  $\frac{\partial f(\bar{x})}{\partial x_i}$  indicates the slope of the function f restricted to the subset of  $\mathbb{R}^n$  s.t.

$$x_j = \bar{x}_j, \quad \forall j \neq i$$

Therefore, the vector of the slopes  $\left(\frac{\partial f}{\partial x_1}, \cdots, \frac{\partial f}{\partial x_n}\right)$  denotes the tangent plane's "slope".

• The gradient  $\left(\frac{\partial f}{\partial x_1}, \dots, \frac{\partial f}{\partial x_n}\right)$  points out the direction in  $\mathbb{R}^n$  along which f increases the fastest.

### Example interpreting properties of gradient: z = 2x + y

- 1. Plot it in x y z;
- 2. Calculate the gradient at (1,2) as:  $\nabla f(1,2) = (2,1)$  through definition.
- 3. Demonstrate the solution to the following optimization problem

$$\lim_{\Delta x, \Delta y} 2(\bar{x} + \Delta x) + (\bar{y} + \Delta y)$$
$$s.t. ||(\Delta x, \Delta y)|| = 1$$

#### 4.2 Vector & Matrix Differentiation

**Definition 4.1.** For  $f: \mathbb{R}^n \to \mathbb{R}^m$ , write  $f(x) = \begin{bmatrix} f^{(1)}(x) \\ \vdots \\ f^{(m)}(x) \end{bmatrix}$ . Define the derivative of f as

$$f'(x) = Df(x) = D\begin{bmatrix} f^{(1)}(x) \\ \vdots \\ f^{(m)}(x) \end{bmatrix} = \begin{bmatrix} f_1^{(1)}(x) & \cdots & f_n^{(1)}(x) \\ \vdots & \ddots & \vdots \\ f_1^{(m)}(x) & \cdots & f_n^{(m)}(x) \end{bmatrix}, \text{ where } f'(x) \text{ is defined to be a}$$

 $m \times n$  matrix.

Convention here is different from what we used for multi-variable single-valued functions, where  $f^{(n)} := \frac{d^n f}{dx^n}$  for  $f : \mathbb{R}^n \to \mathbb{R}$ . Here,  $f^{(m)}(x)$  is the m-th in the codomain.

Example: for 
$$f(x,y) = \begin{bmatrix} 2x+y \\ 3x^2 \end{bmatrix}$$
, then  $Df(x,y) = \begin{bmatrix} 2 & 1 \\ 6x & 0 \end{bmatrix}$ .

**Definition 4.2** (Higher order (partial) derivatives for single-valued function). The Higher order (partial) derivatives for  $f : \mathbb{R}^n \to \mathbb{R}$  is a  $1 \times n$  vector defined as:

$$\frac{df}{dx} = Df = [f_1(x), \cdots, f_n(x)].$$

For the second order derivative:

$$\frac{d}{dx}\left(\frac{df}{dx}\right) = D^2 f = \begin{bmatrix} f_{11}(x) & \cdots & f_{1n}(x) \\ \vdots & \ddots & \vdots \\ f_{n1}(x) & \cdots & f_{nn}(x) \end{bmatrix} = \frac{\partial^2 f}{\partial x \partial x^T}$$

Here, x is a vector, and  $\partial x \partial x^T \simeq x^2$  in  $\mathbb{R}^1$ , where  $\simeq$  is defined in terms of equivalent expression.

Note that, the goal here is to let  $f' \cdot x$  yield a scalar, instead of a matrix.

Exercise 4.3. (Example: for taking derivative of composite functions.)

1. f(x,y) is a real-valued function, and  $x = u + \log v$  and  $y = u - \log v$ . Show that:

$$\frac{\partial^2 f}{\partial u^2} = \frac{\partial^2 f}{\partial x^2} + 2\frac{\partial^2 f}{\partial x \partial y} + \frac{\partial^2 f}{\partial y^2}$$

*Proof.* Write first taht f(x(u, v), y(u, v)) where  $x(u, v) = u + \log v$ ;  $y(u, v) = u - \log v$ .

$$\frac{\partial f(x(u,v),y(u,v))}{\partial u} = f_x \frac{\partial x}{\partial u} + f_y \frac{\partial y}{\partial u} = f_x + f_y$$

(note that,  $\frac{\partial x}{\partial u} = 1$  and  $\frac{\partial y}{\partial u} = 1$ , by evaluating the functional form.)

Now, take this simplified functional form for  $\frac{\partial f}{\partial u} = f_x + f_y$ , look into:

$$\frac{\partial^2 f}{\partial u^2} = f_{xx} \cdot \frac{\partial x}{\partial u} + f_{xy} \cdot \frac{\partial y}{\partial u} + f_{yx} \cdot \frac{\partial x}{\partial u} + f_{yy} \cdot \frac{\partial y}{\partial u}$$

Simplify this by evaluating  $\frac{\partial x}{\partial u} = 1 = \frac{\partial y}{\partial u}$  and get the expression we wanted to show.  $\square$ 

2.  $f(x,y) = g(\frac{x}{y})$ , g is differentiable, show that:

$$x \cdot \frac{\partial f}{\partial x} + y \frac{\partial f}{\partial y}$$

*Proof.* Take the original equality and differentiated w.r.t x and y, we have:

$$\frac{\partial f}{\partial x} = g' \frac{\partial \left(\frac{x}{y}\right)}{\partial x} = \frac{1}{y} \cdot g'$$

$$\frac{\partial f}{\partial y} = g' \frac{\partial \left(\frac{x}{y}\right)}{\partial y} = \frac{x}{-y^2} g'$$

**Definition 4.4** (Inner product). For matrices and vectors that are made of real numbers, inner product is defined as:

$$a^T \cdot x = \sum_{i=1}^n a_i x_i,$$

where  $x = (x_1, \ldots, x_n)^T$  (as a  $n \times 1$  (column) vector); and  $a = (a_1, \ldots, a_n)^T$  (as a  $n \times 1$  (column) vector).

# References

[Rudin(1976)] Walter Rudin. Principles of mathematical analysis. New York: McGraw-Hill, [1976], 1976.