

Introduction to E-Graphs

Rebecca Swords | Women in Compilers and Tools

Questions

What are e-graphs?

What are they good for?

How do they work?

E-Graph

A data structure representing an equivalence relation over terms

Practical Applications

Theorem proving

SMT solving

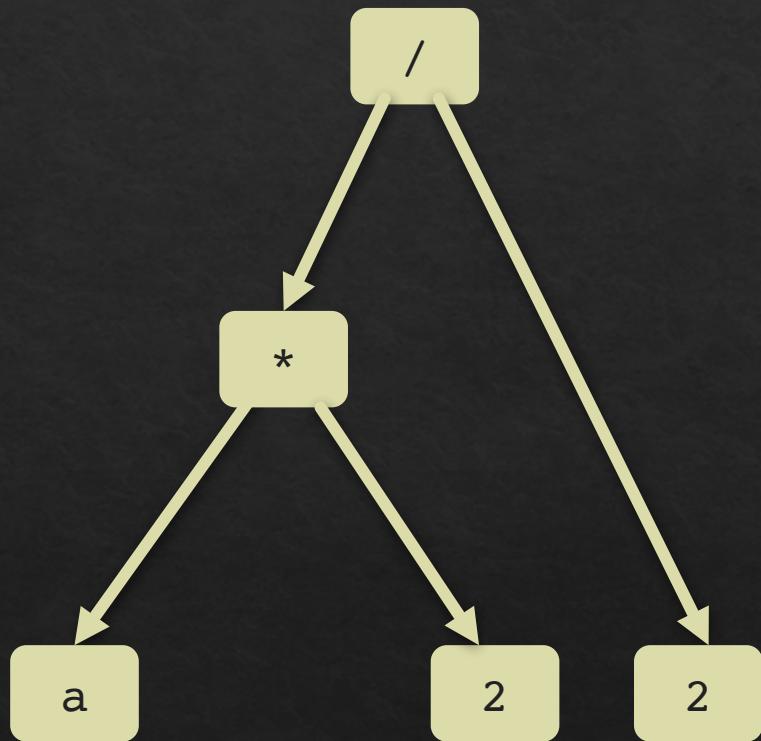
Optimization

Translation validation

Compilation

Synthesis

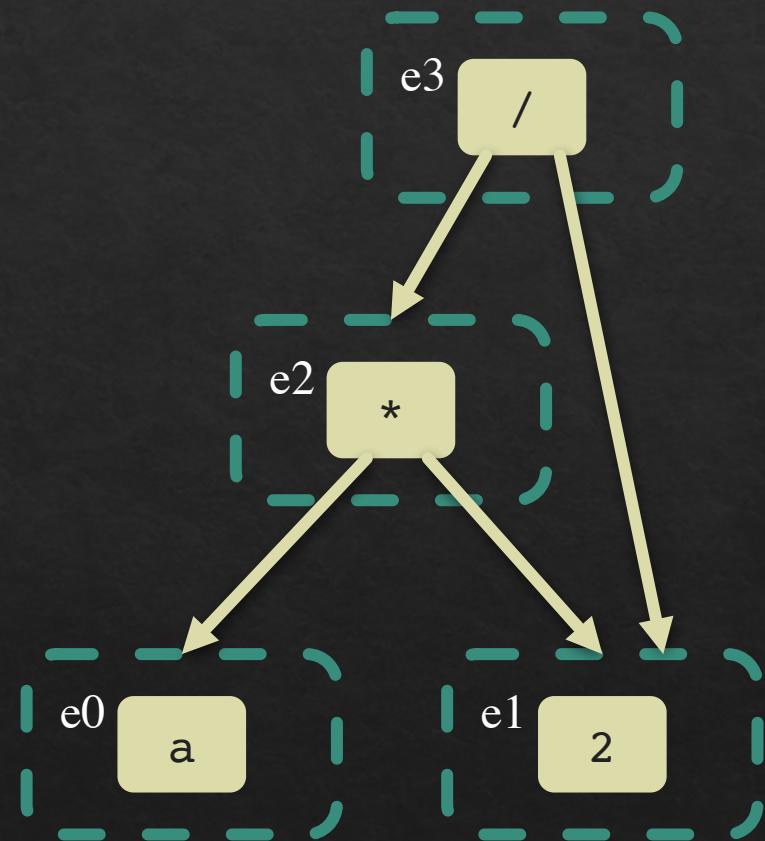
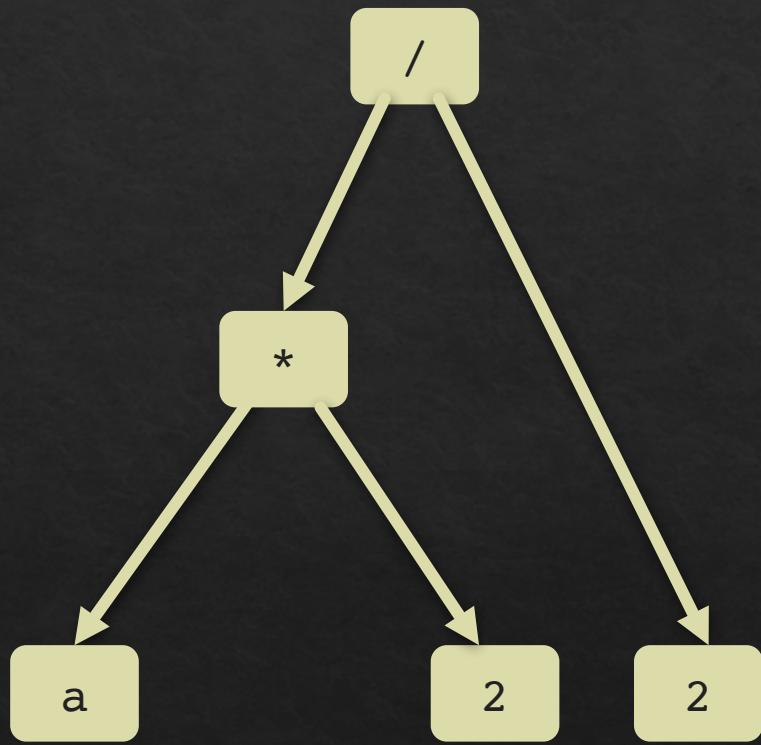
Running Example: $(a * 2) / 2$



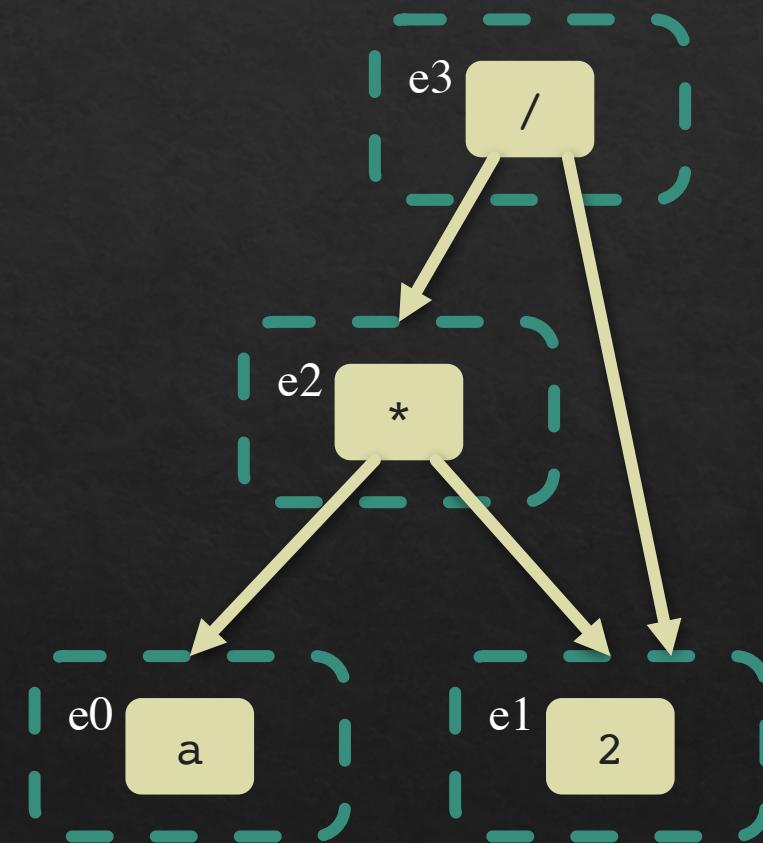
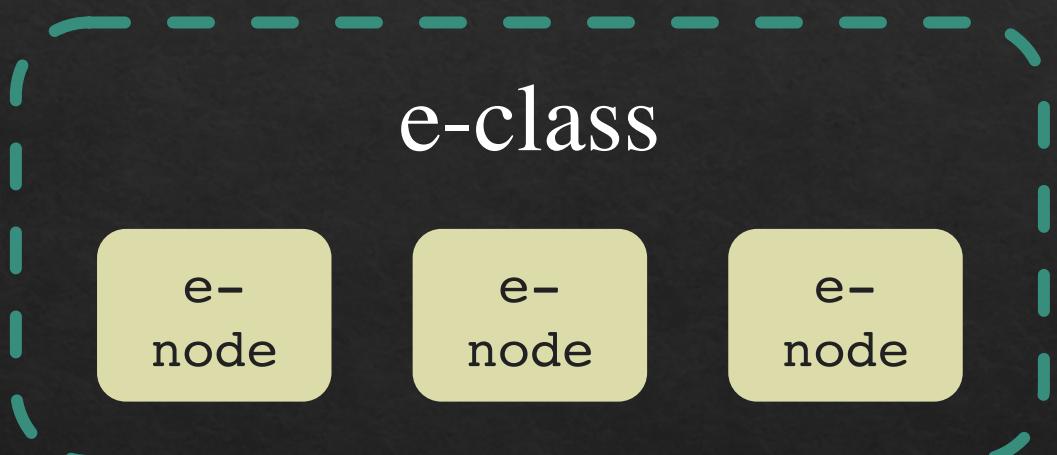
- This reduces to a
- We can use e-graphs to do it!

Start with this AST

Into an E-Graph



Into an E-Graph

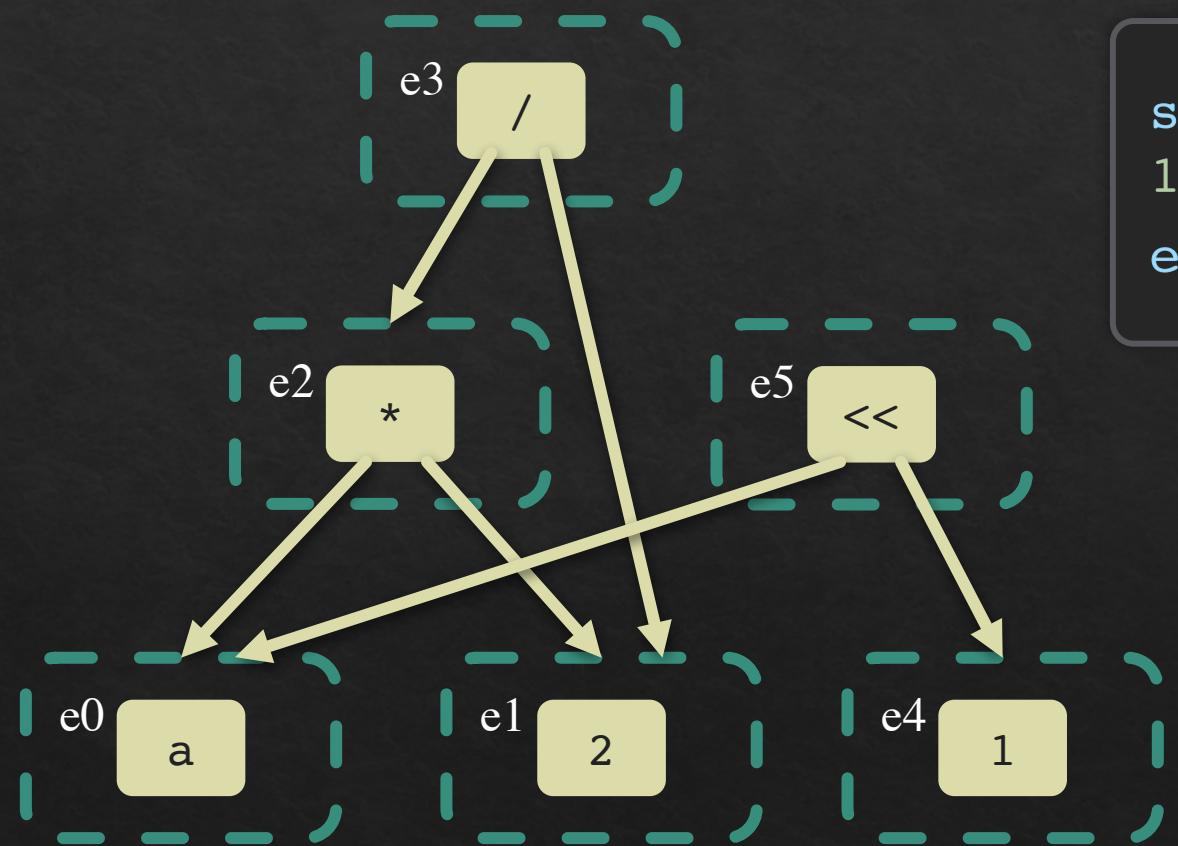


Build it with Quiche

```
1 expr = (ExprNode('a', ()) * 2) /  
2  
3 quiche_tree = ExprTree(expr)  
4  
5 egraph = EGraph(quiche_tree)
```

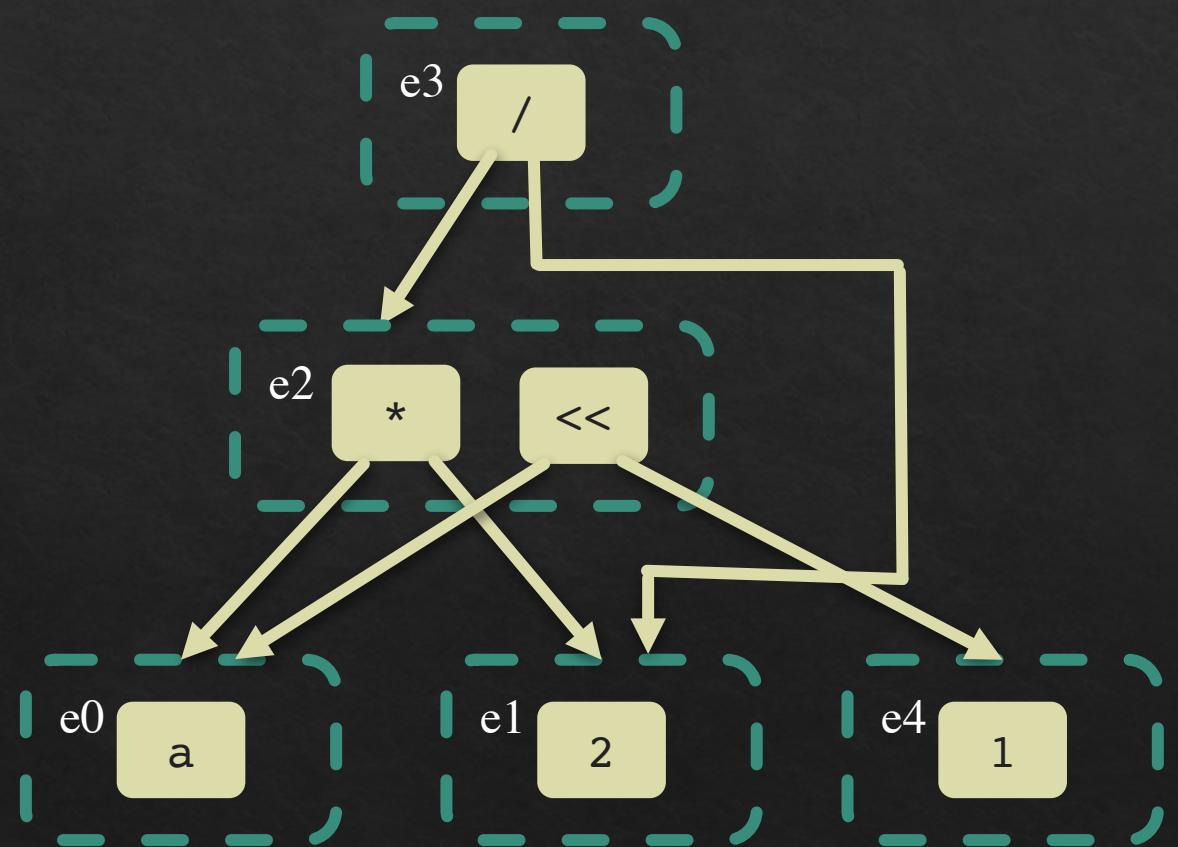
1. Parse term into the arithmetic language structure
2. Construct intermediate QuicheTree representation
3. Create e-graph from QuicheTree

Another Term: $a \ll 1$



```
shift_expr = ExprNode('a', ()) <<  
1  
egraph.add(ExprTree(shift_expr))
```

Merging Equivalent Terms



We assert:

$$a^*2 \equiv a<<1$$

It follows that:

$$(a^*2)/2 \equiv (a<<1) / 2$$

Manual Merging in Quiche

```
1 shift_eclass =  
  egraph.add(ExprTree(shift_expr))  
  times_node = ExprNode('a', ()) * 2  
  times_eclass =  
    egraph.add(ExprTree(times_node))  
  
2 egraph.merge(times_eclass, shift_eclass)  
  
3 egraph.rebuild()
```

1. Save e-class IDs for the expressions to be merged
2. Merge the two e-classes together
3. Restore e-graph invariants

E-Graphs More Formally

Structure

- ❖ E-node: an n-ary function symbol and n children (e-class IDs)
- ❖ E-class: set of e-nodes
- ❖ Union-find over e-classes: `add`, `merge`, `find` operations
- ❖ Canonical e-node: for each child, i , `find(i) = i`
- ❖ Hashcons: maps canonical e-nodes to e-classes

Invariants

- ❖ Hashcons maps all canonical e-nodes
- ❖ Equivalence closed under congruence, i.e., congruent e-nodes are in the same e-class
 - If $a = b$, then $f(a) = f(b)$

Why is this good for term rewriting?

Instead of destructive rewrites, put *all* equivalent terms in the e-graph

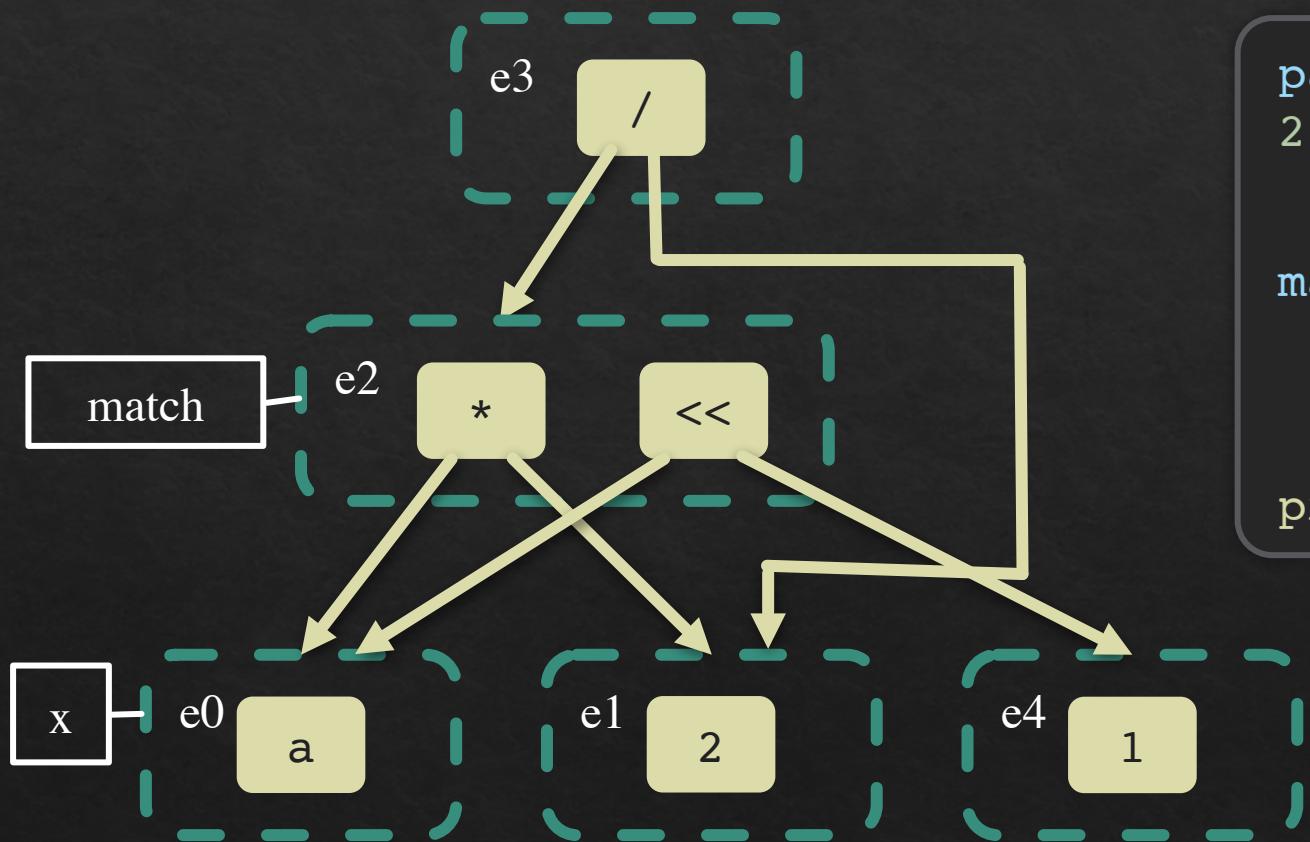
- No worries about phase ordering
- Consider all options and choose the “best” at the end

E-Matching

*Pattern matching
for e-graphs!*

- Add pattern variables to language
- `ematch` searches for a pattern and returns:
 - ◊ e-class matching the term
 - ◊ substitution from vars to e-class IDs

E-Matching Example: $x * 2$



```
pattern = ExprTree(ExprNode('x', ()) *  
2)  
  
matches = egraph.ematch(pattern,  
egraph.ecllasses())  
  
print(matches)
```

```
[ (e2, { 'x': e0 }) ]
```

Rewriting Rules: Pattern Merges

```
1 rule = ExprTree.make_rule(lambda x:  
    (x * 2, x << 1))  
  
2 Rule.apply_rules([rule], egraph)  
  
3 print("Shift e-class: ", shift_eclass)  
print("Shift e-class.find(): ",  
shift_eclass.find())
```

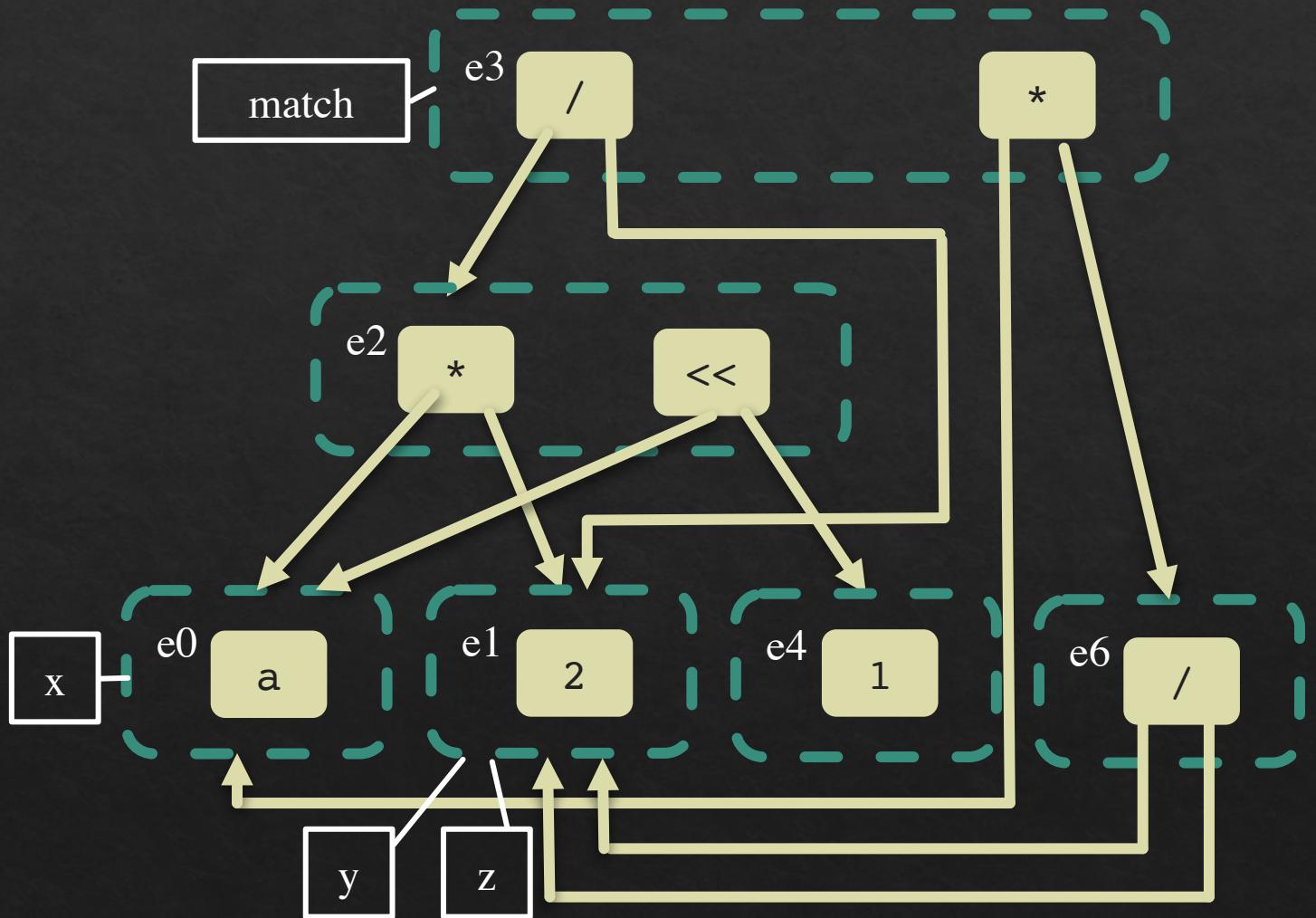
1. Create a rule:
 $x * 2 === x \ll 1$
2. Apply all rules to e-graph (and rebuild)
3. Shift e-class: e5
Shift e-class find: e2

Another rewrite:

$$(x^*y)/z$$

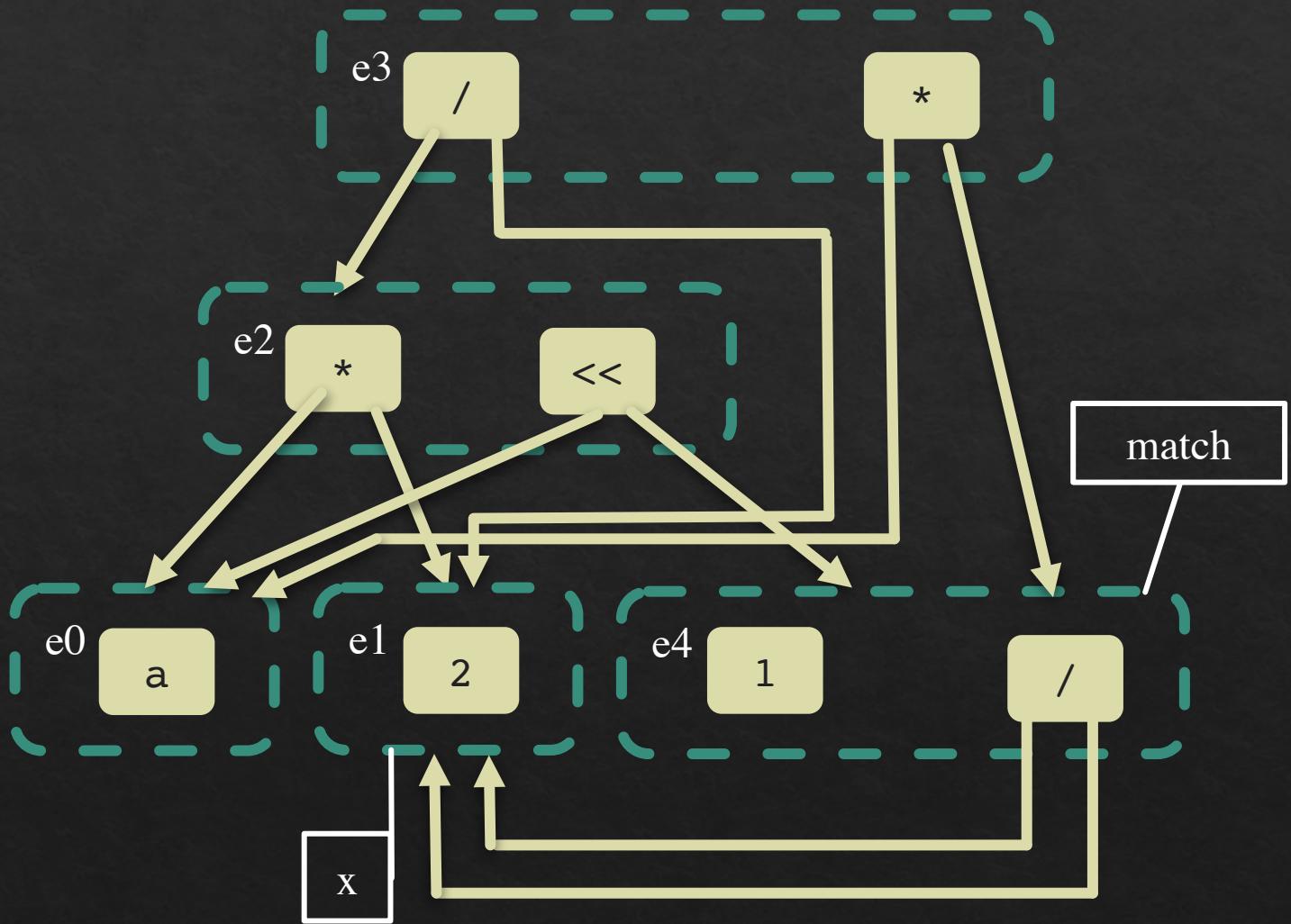
====

$$x^*(y/z)$$

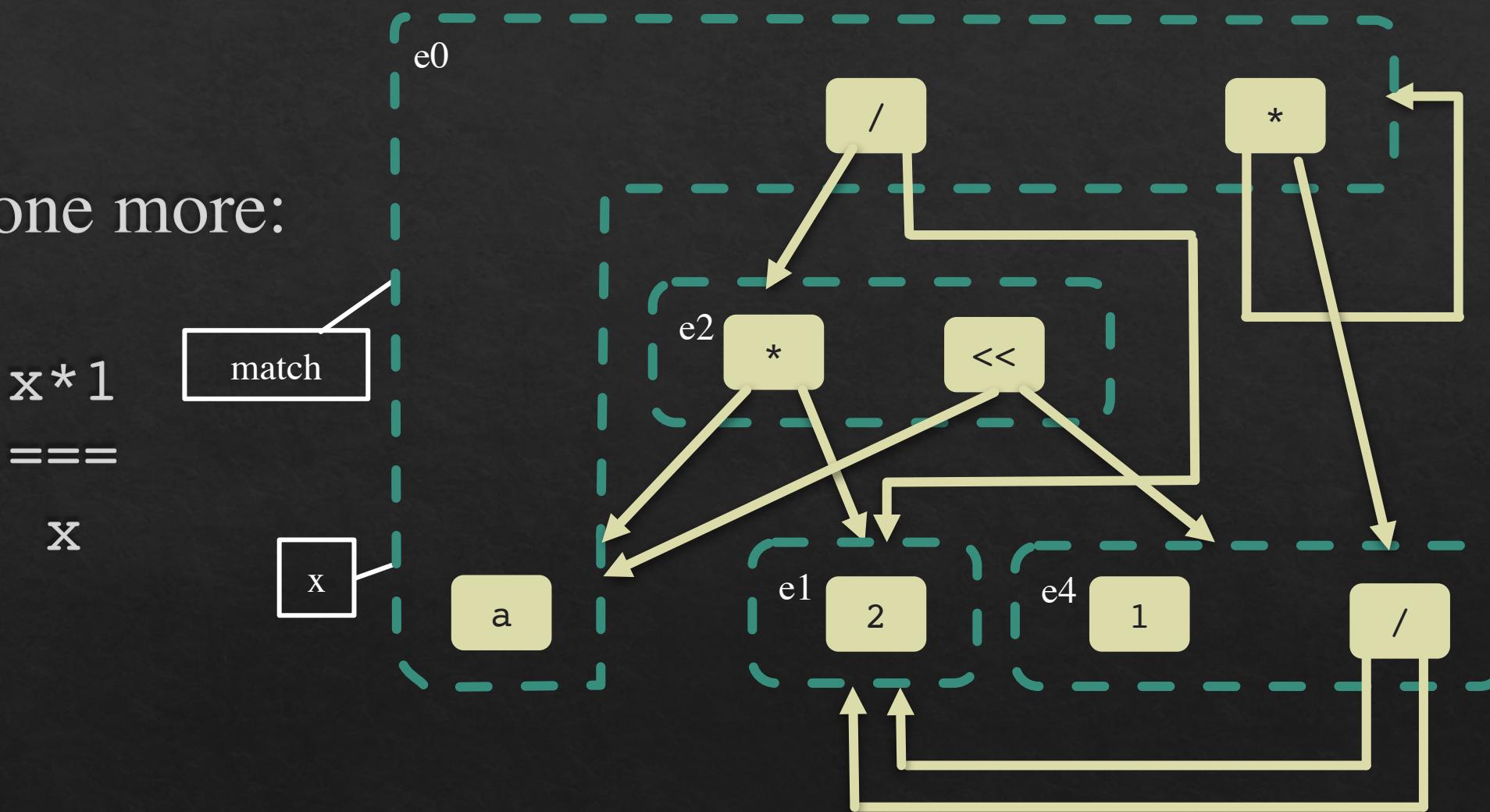


And another:

x/x
====
1



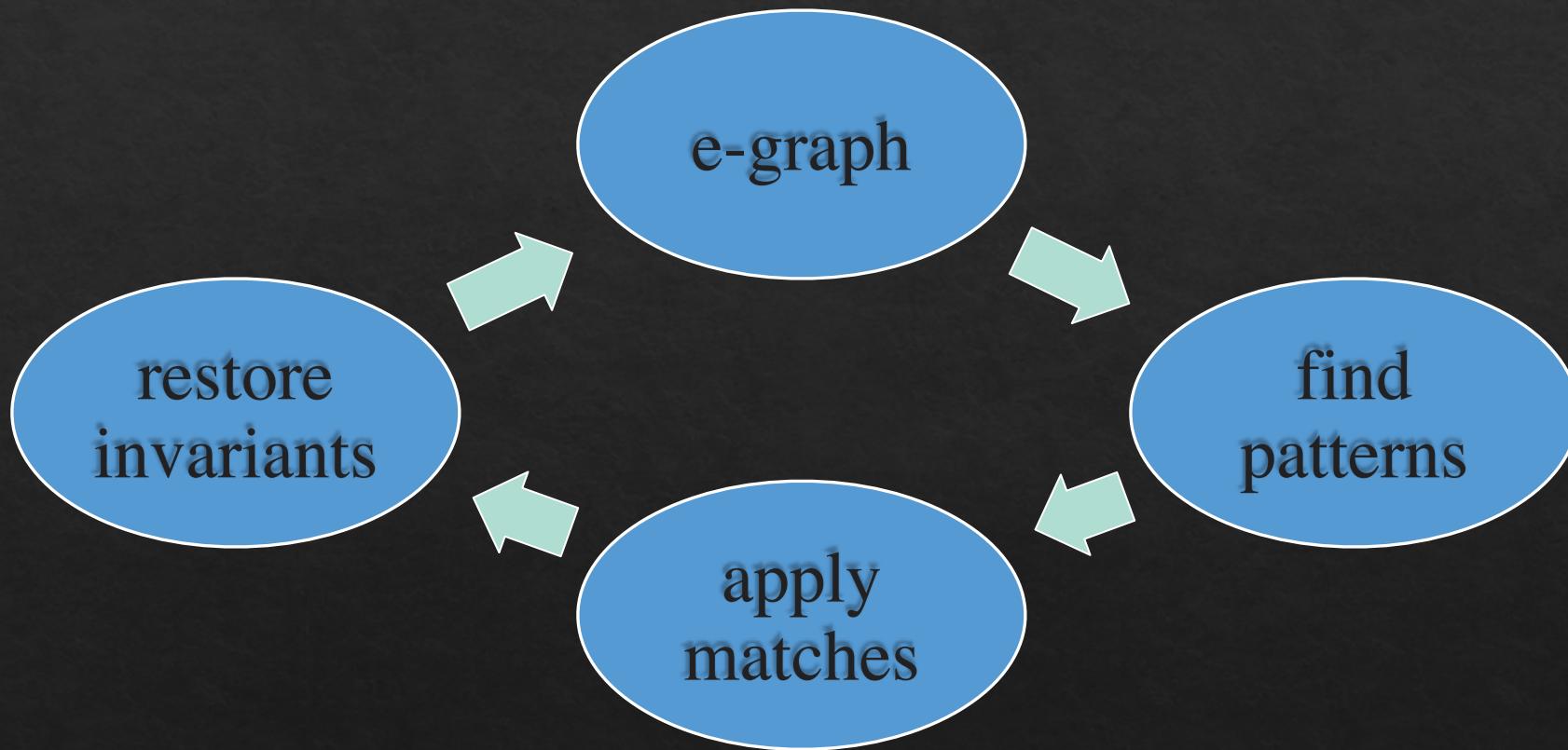
And one more:



Keep applying rewrite
rules until no new
changes are made

Equality
Saturation

Equality Saturation Loop



Apply Rules Until Saturation

1

```
rules = [
    ExprTree.make_rule(lambda x, y, z:
        ((x * y) / z, x * (y / z))),
    ExprTree.make_rule(lambda x:
        (x / x, ExprNode(1, ()))),
    ExprTree.make_rule(lambda x: (x * 1, x))
]
```

2

```
while not egraph.is_saturated():
    Rule.apply_rules(rules, egraph)
aeclasse = egraph.add(ExprTree(ExprNode('a',
    ())))
assert aeclasse.find() == egraph.root.find()
```

3

1. Same 3 rules we just applied
2. Apply rules until the e-graph is saturated
3. Verification: expect **a** to have merged with the “root” e-class

E-Class Analysis

Domain-specific e-graph extensions

- ❖ Attach datum to each e-class based on e-nodes: `make`
- ❖ Merge data when e-classes merge: `join`
- ❖ Update e-class based on datum: `modify`
- ❖ Form a join-semilattice

What Can E-Class Analyses Do?

- ❖ Program analysis
- ❖ Conditional or dynamic rewrites
- ❖ Debugging
- ❖ Pruning
- ❖ On-the-fly term extraction

Standardized interface for extending e-graphs!

Analysis Invariant

for each e-class

$$\forall c \in G. \quad d_c = \bigwedge_{n \in c} \text{make}(n) \quad \text{and} \quad \text{modify}(c) = c$$

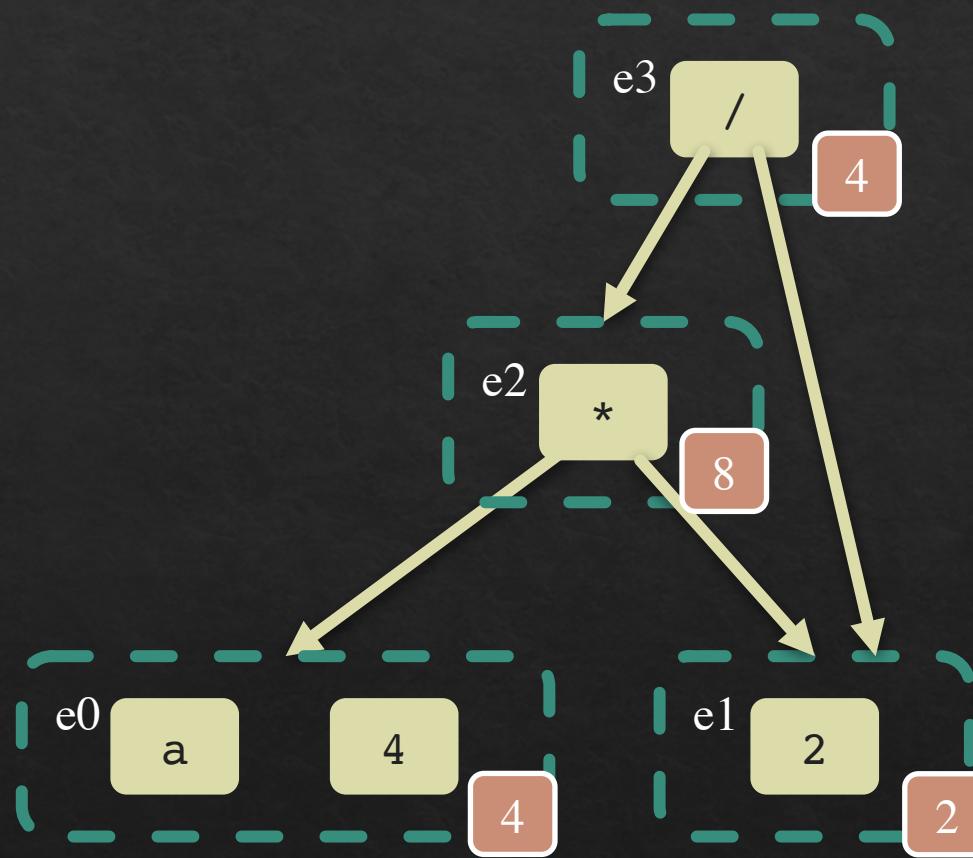
fixed point

data is the same as `make`-ing
data for each e-node and then
`join`-ing

Constant Folding

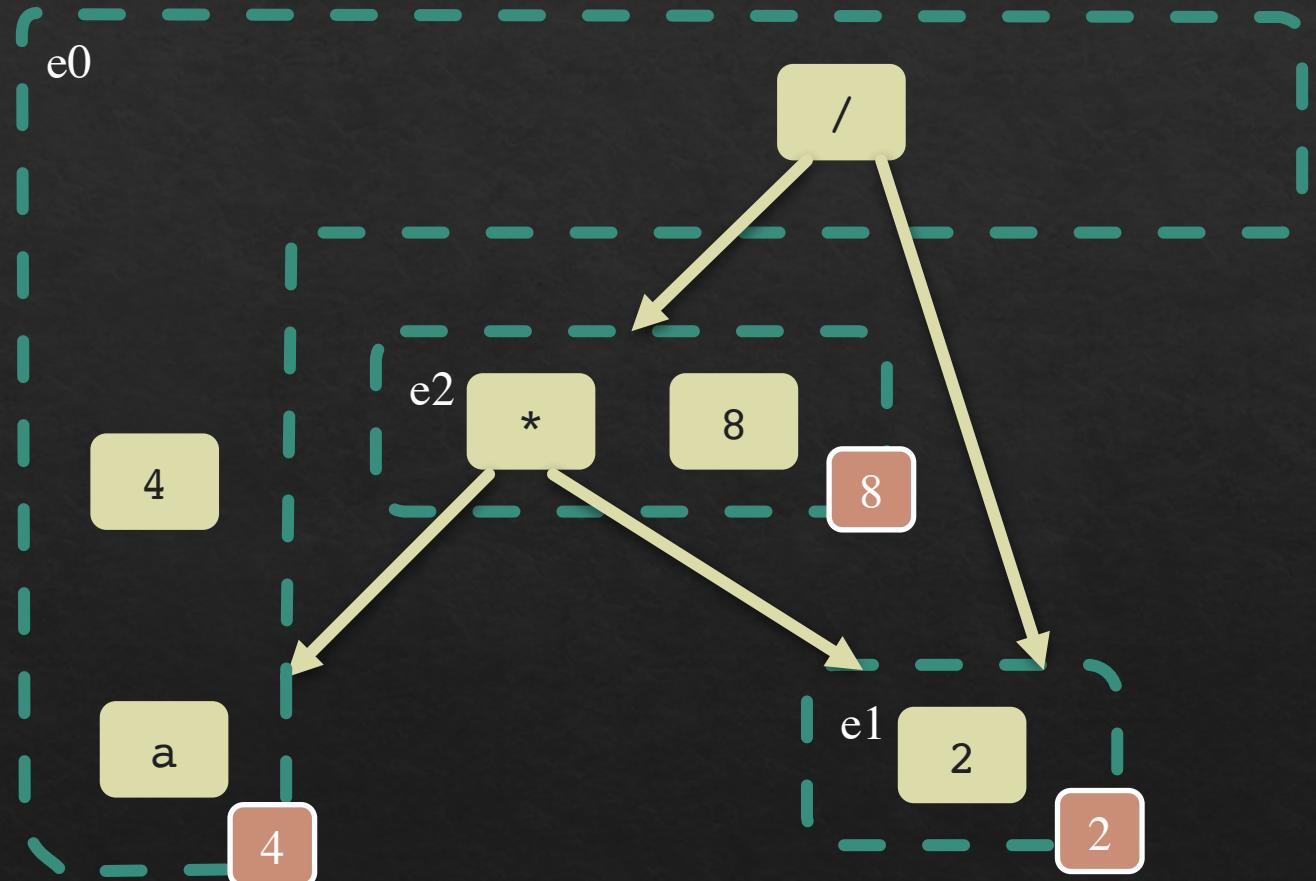
E-Class Analysis

Suppose we learn that a
===== 4



Constant Folding E-Class Analysis

Suppose we learn that a
 $\equiv 4$



Constant Folding: Usage

1

```
expr = (ExprNode('a', ()) * 2) / 2
quiche_tree = ExprTree(expr)
egraph = EGraph(quiche_tree, ExprConstantFolding())
```

2

```
four_eclasse = egraph.add(ExprTree(ExprNode(4, ())))
a_eclasse = egraph.add(ExprTree(ExprNode("a", ())))
```

3

```
egraph.merge(a_eclasse, four_eclasse)
```

4

```
egraph.rebuild()
```

5

```
assert egraph.root.data == 4
```

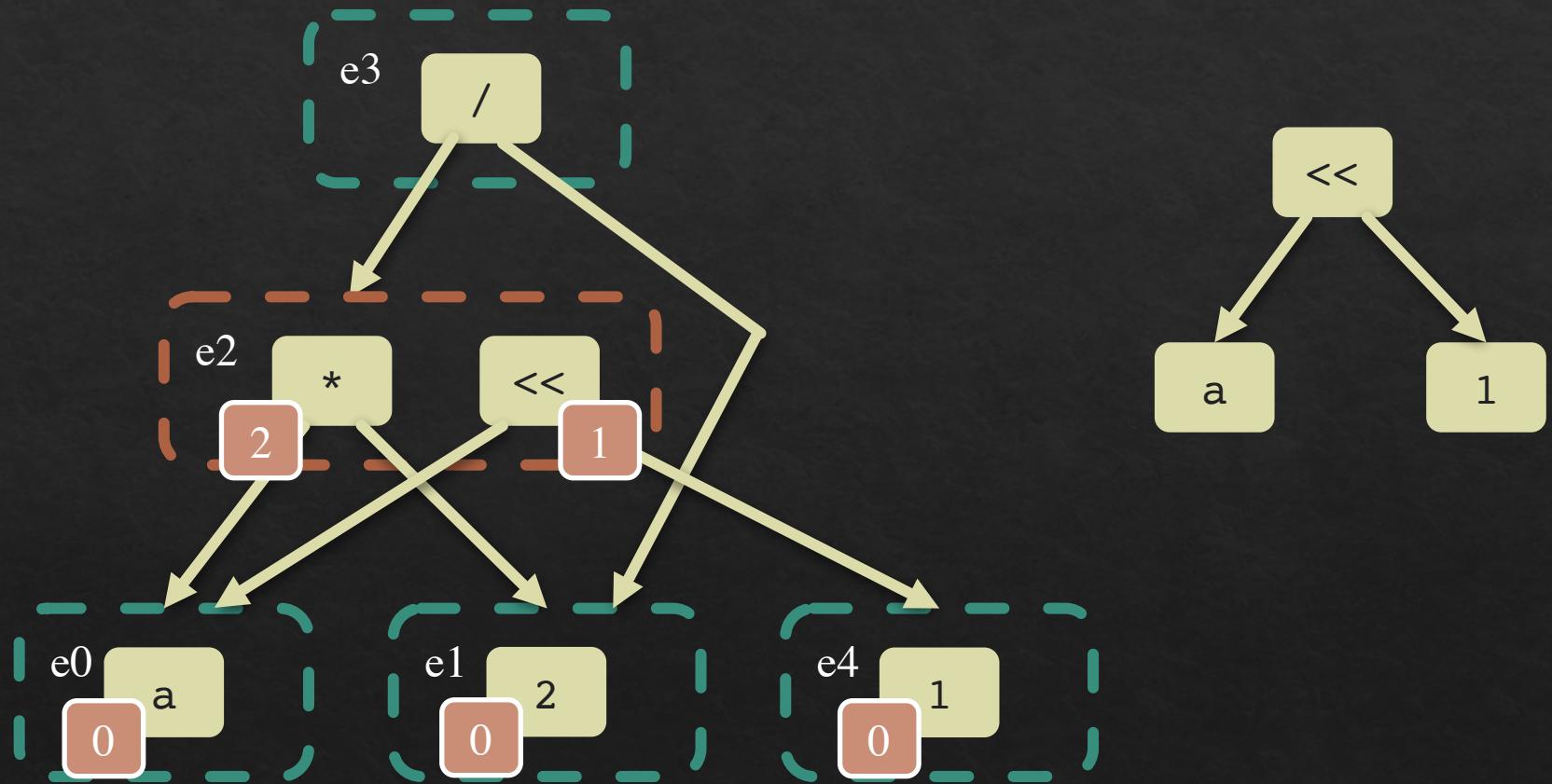
1. Create e-graph with constant folding analysis
2. Get e-class IDs
3. Merge 4 with a
4. Rebuild (update analysis)
5. Verify

Term Extraction

- ❖ Pick an e-class to extract
- ❖ Cost model assigns a cost to e-nodes
- ❖ Choose best e-node for each e-class
- ❖ Construct a term by combining the e-node values

Term Extraction Example

Operator	Cost
+	1
<<	1
*	2
/	3
default	0



Term Extraction Example

```
1 cost_model = ExprNodeCost()
2 extractor= MinimumCostExtractor()
3 extracted = extractor.extract(
4     cost_model,
5     egraph,
6     egraph.root.find(),
7     ExprTree.make_node)
8 assert str(extracted) == "a"
```

1. Initialize cost model and extractor
2. Extract the best term
3. Specify which e-class to extract
4. Function to construct ExprTree from e-node data

More on Quiche

- ❖ Add your own languages!
 - ❖ Bring your own parser, adapt your AST into a QuicheTree
- ❖ End-to-end Python rewriting!
 - ❖ Uses native Python parser (v3.7+)
 - ❖ Read/write valid Python files
- ❖ Native Python!
 - ❖ With all its pros and cons

QuicheTree

Quiche requires the user to provide a parsed tree that implements `QuicheTree` (“*bring your own parser*”).

value()

the e-node key

children()

list of the node's children

is_pattern_symbol()

for e-matching; indicates if the node is a pattern

```
class QuicheTree(ABC):  
    @abstractmethod  
    def value(self)  
  
    @abstractmethod  
    def children(self)  
  
    @abstractmethod  
    def is_pattern_symbol(self)
```

Links and References

- ❖ Quiche repo: <https://github.com/riswords/quiche>
- ❖ egg website: <https://egraphs-good.github.io/>
- ❖ egg: Fast and extensible equality saturation (POPL ‘21, Willsey, et al.): <https://dl.acm.org/doi/10.1145/3434304>
- ❖ Equality-Based Translation Validator for LLVM (CAV ‘11, Stepp, Tate, & Lerner): https://cseweb.ucsd.edu/~rtate/publications/eqsat/eqsat_stepp_cav11.pdf
- ❖ babble: Learning Better Abstractions with E-Graphs and Anti-Unification (POPL ‘23, Cao, et al.): <https://dl.acm.org/doi/10.1145/3571207>



Questions?

Additional References from Q&A

1. Link to the public E-Graphs Zulip chat: <https://egraphs.zulipchat.com/>
2. Perfect Reconstructability of Control Flow from Demand Dependence Graphs (Bahmann, et al. 2014) <https://dl.acm.org/doi/abs/10.1145/2693261>
3. E-Graphs Zulip discussion of using RVSDG representation: <https://egraphs.zulipchat.com/#narrow/stream/328976-Program-Optimization/topic/PEGs>
4. Equality Saturation for Tensor Graph Superoptimization (Yang, et al., MLSys 2014): <https://arxiv.org/abs/2101.01332>
5. Relational e-matching (Zhang, et al., POPL 2022)
6. [Logging an Egg: Datalog on E-Graphs \(EGRAPHS 2022\) - PLDI 2022 \(sigplan.org\)](#)