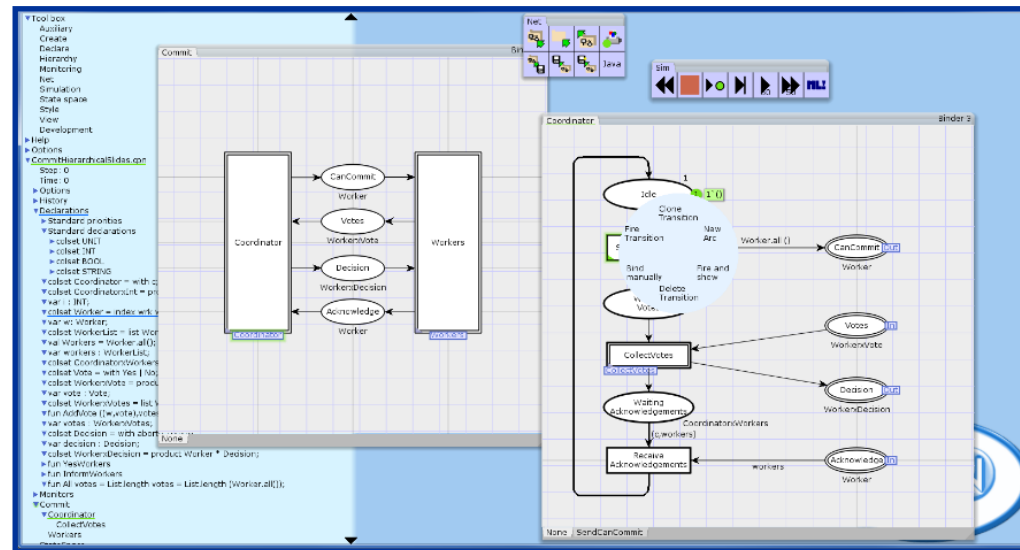


## Lecture 4

# Hierarchical Coloured Petri Nets with Modules



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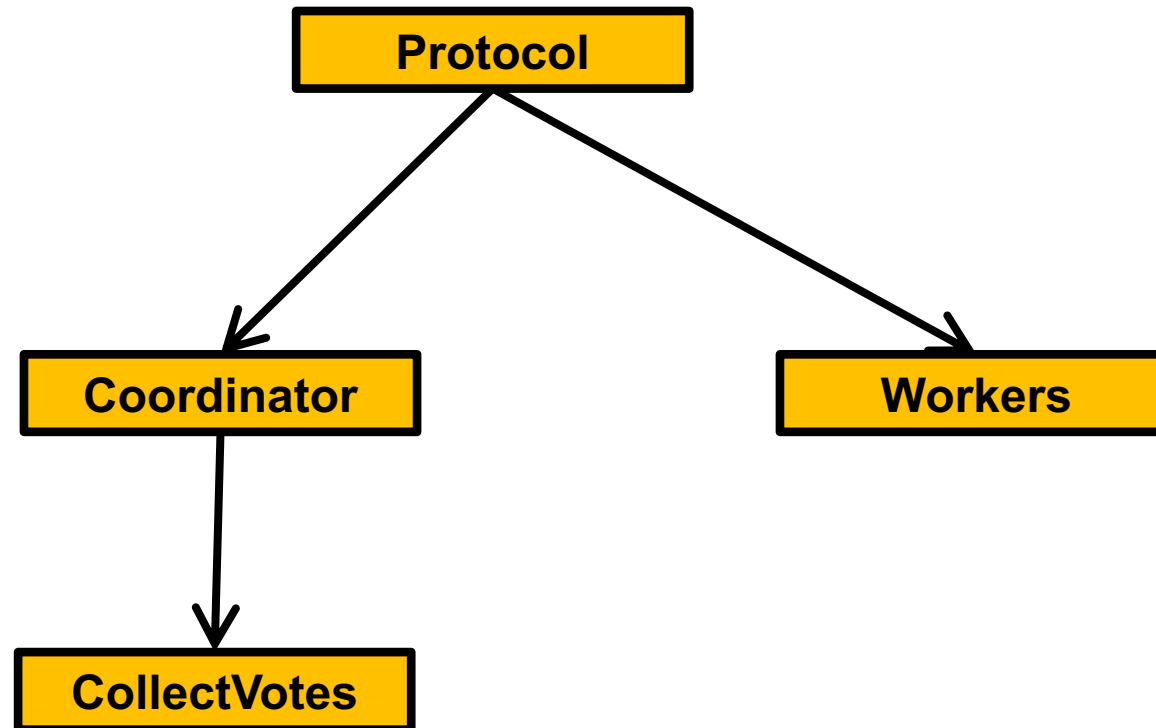
Email: [lmkr@hvl.no](mailto:lmkr@hvl.no) / WWW: [home.hib.no/ansatte/lmkr](http://home.hib.no/ansatte/lmkr)

# Introduction

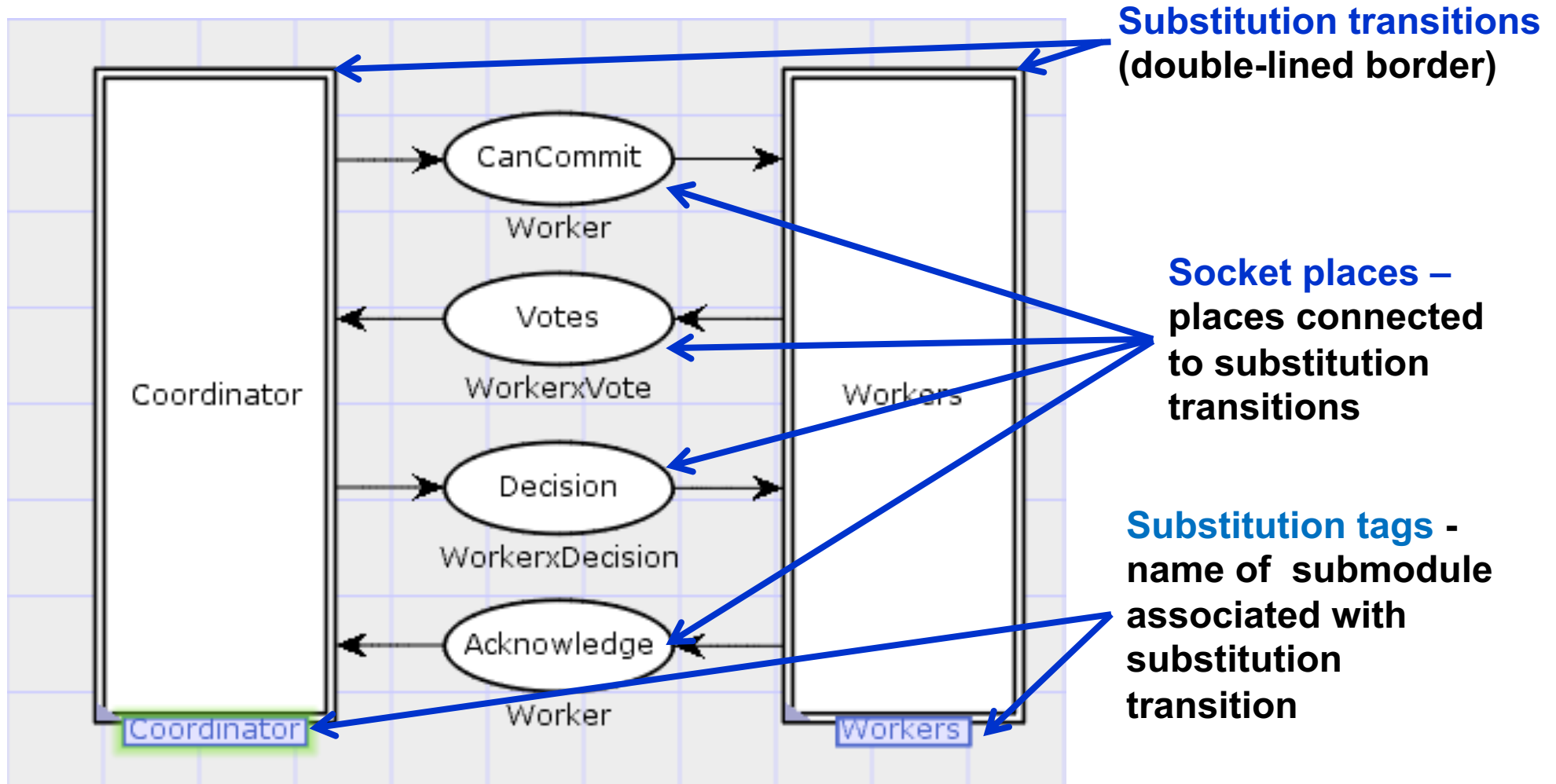
- **Important to be able to split a large CPN models into a set of modules with interfaces**
  - To support construction of large CPN models
  - To support reuse of modules and maintainability
  - To support abstraction and management of details
- **Key concepts of hierarchical CPN modules**
  - A **module** exchange tokens with its environment using input/output **port places**
  - **Substitution transitions** have associated **submodules**
  - **Port-socket relation** associates socket places of substitution transitions with the port places in the associated submodule

# Hierarchical modules

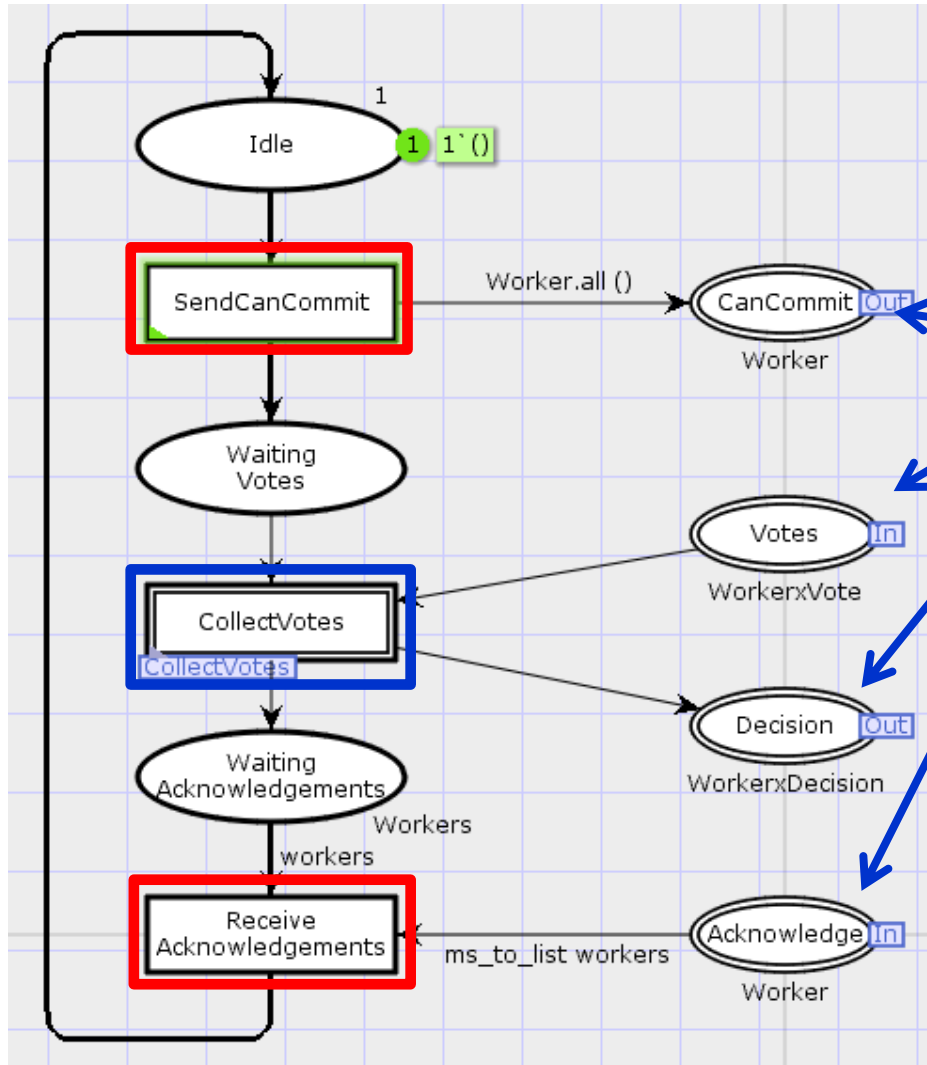
- Model is comprised of a collection of **modules** that are hierarchically organised into levels
- **Example:** the two-phase commit protocol



# Top-level: Protocol module



# Coordinator module



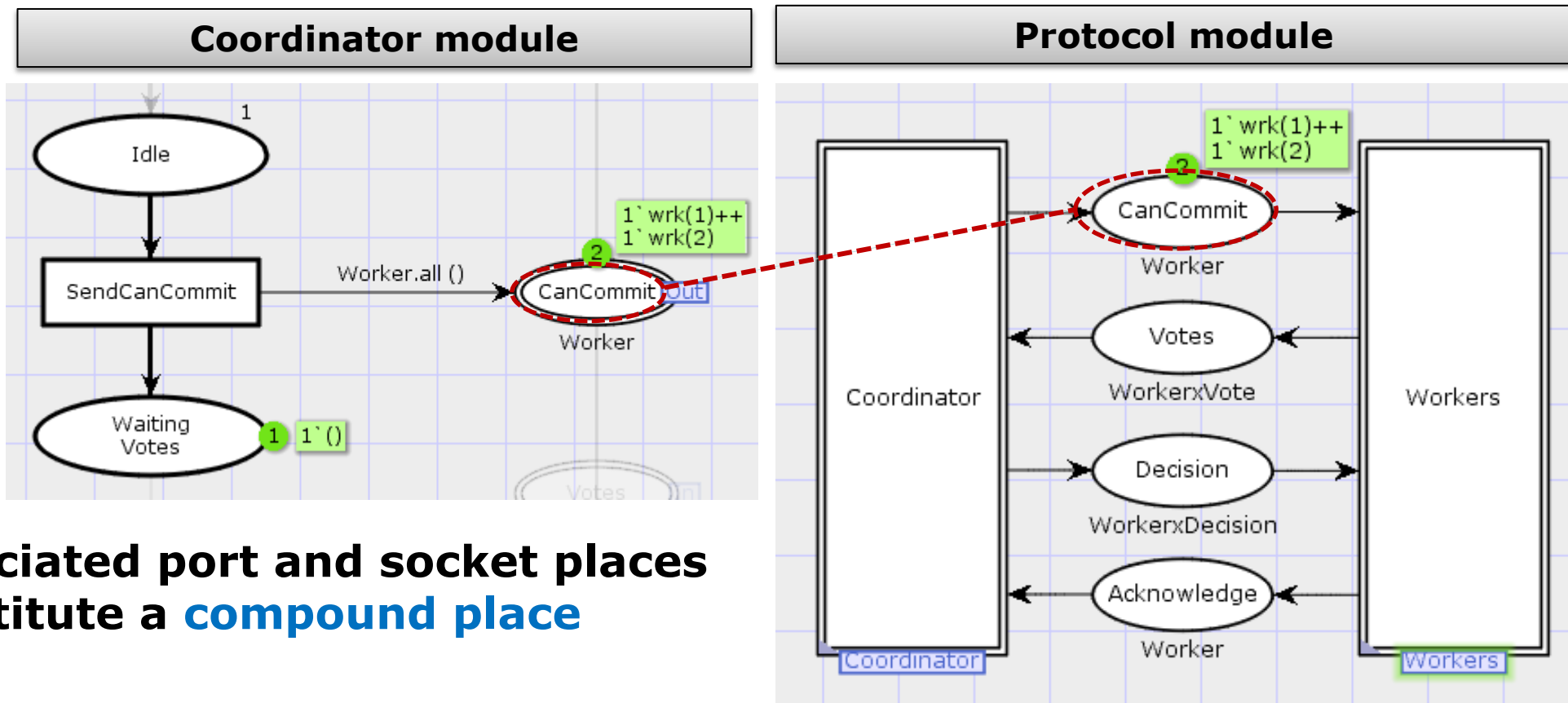
**Port place** - used for exchanging tokens with the upper-level module (IN,OUT,IN/OUT).

**SendCanCommit** and **ReceiveAcknowledgement** are ordinary transitions.

**CollectVotes** is a substitution transition

# Port-socket place relation

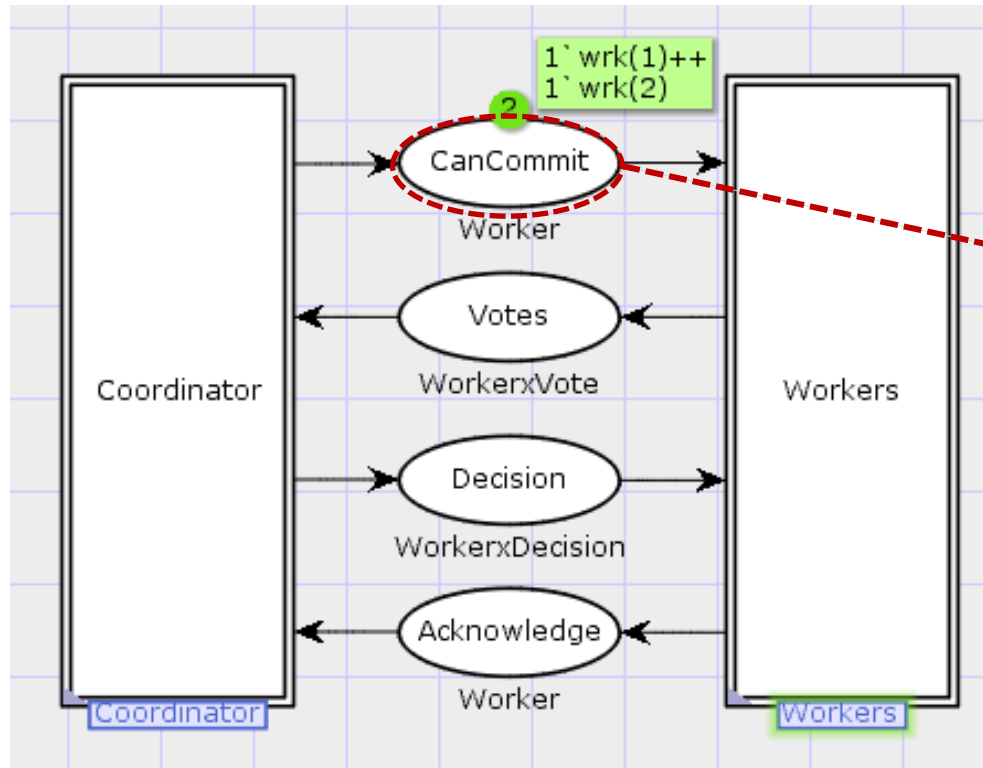
- Tokens added (removed) on a port place are added (removed) on the **associated socket place**



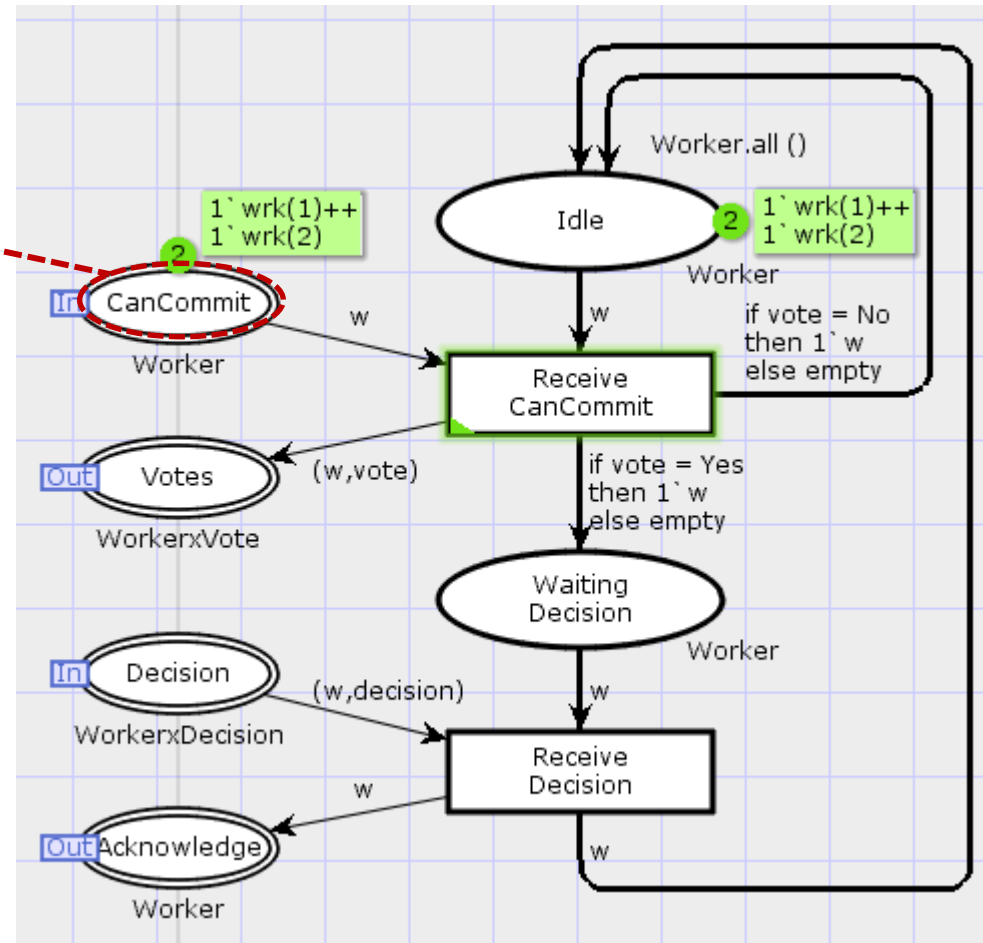
- Associated port and socket places constitute a **compound place**

# Workers Module

Protocol module



Workers module



# CPN Tools Demo

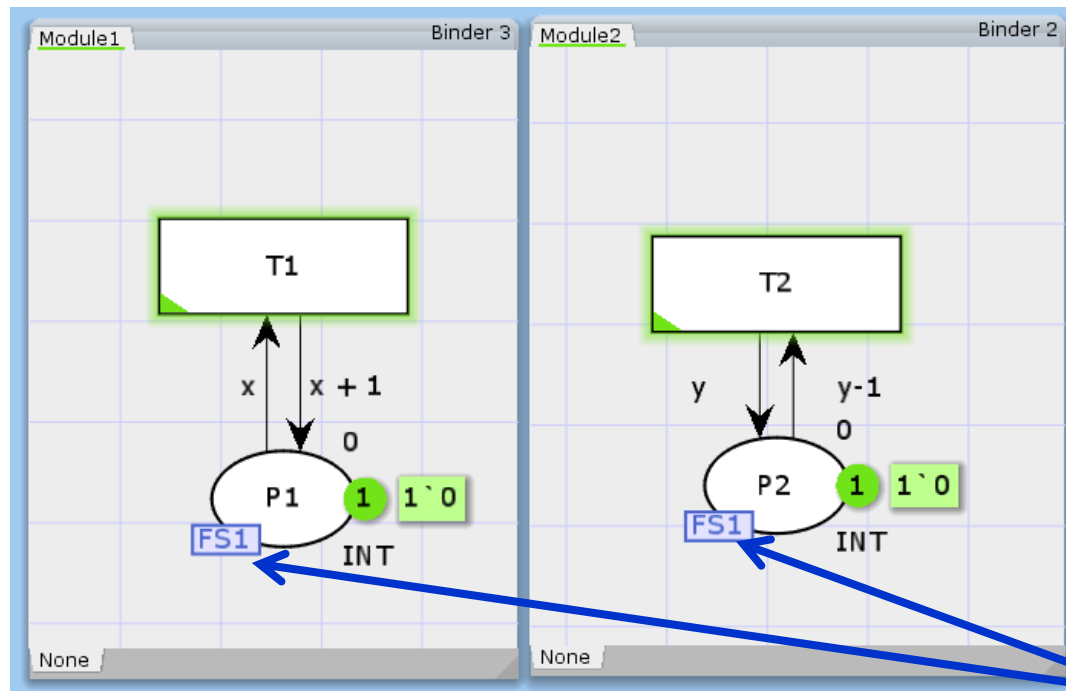
- **Hierarchical CPN models**
  - Navigating hierarchical models
  - Simulation of hierarchical models
  - Editing of modules: top-down and bottom-up development





# Place Fusion Sets

- Group of places to be treated as one compound (global) place



Any change in the marking of P1 will be reflected on P2 (and vice versa)

Similar to global variables  
- should be used with care

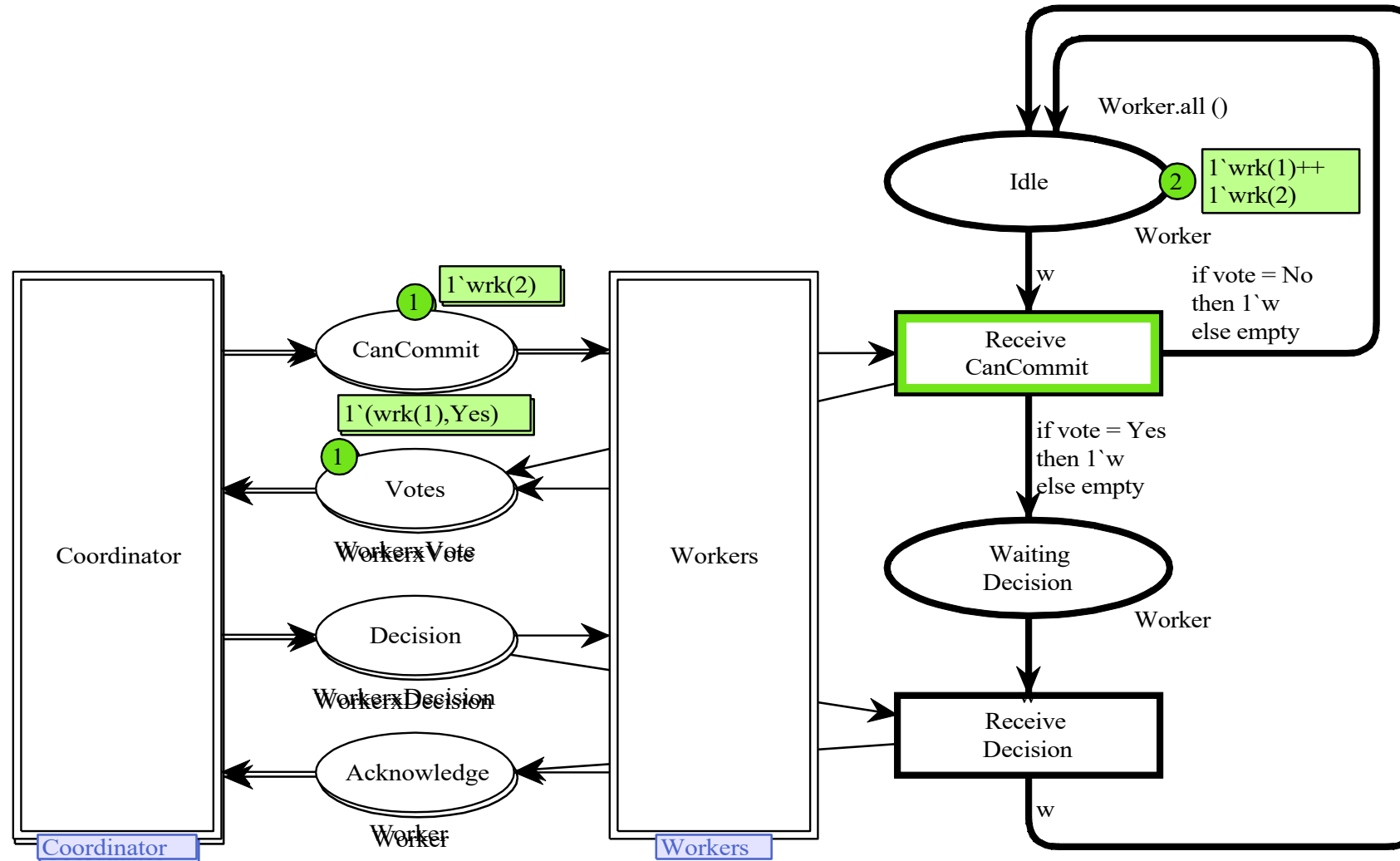
P1 and P2 are fusion places  
belonging to fusion set FS1

# Unfolding Coloured Petri Nets to Place/Transition Nets

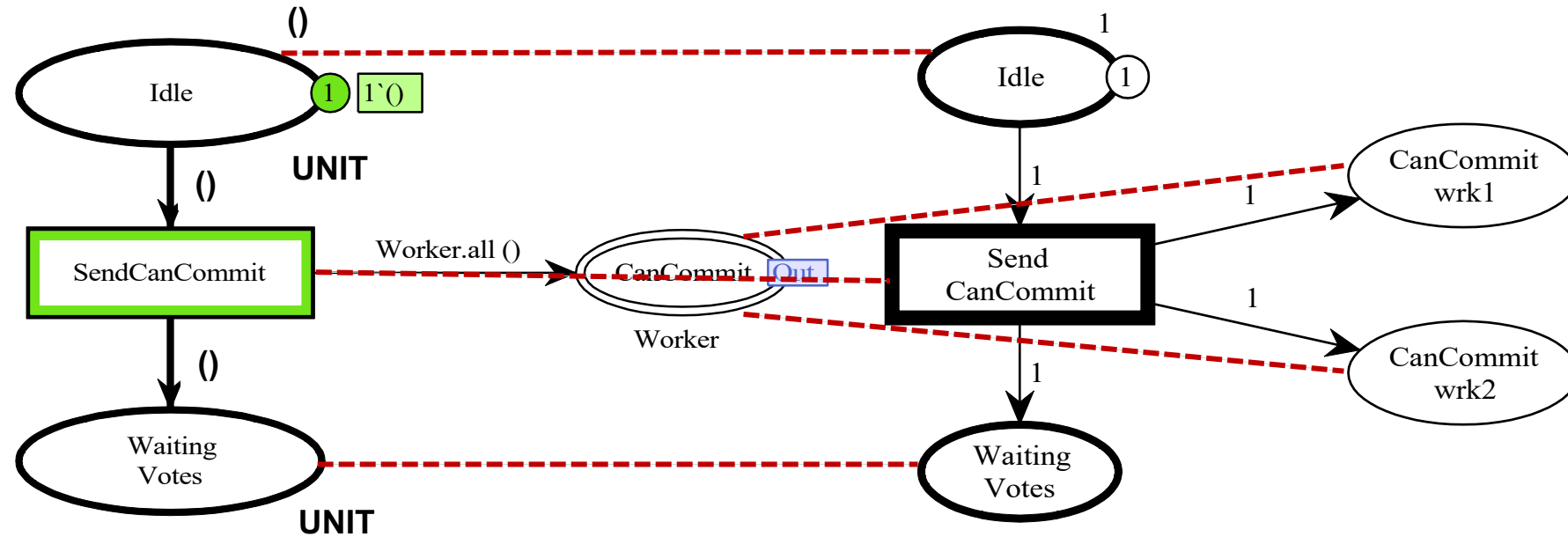
# Unfolding Coloured Petri Nets

- **A hierarchical CPN model can be unfolded to a non-hierarchical Coloured Petri Net**
  - Recursively replace each substitution transition with its associated submodule
  - Associated port and socket places are merged into a single place
- **A non-hierarchical Coloured Petri Net can be unfolded into a Place/Transition Net (PTN)**
  - Replace each CPN place with one PTN place for each colour in the colour set of the CPN place
  - Replace each CPN transition with one PTN transition for each possible binding of the CPN transition

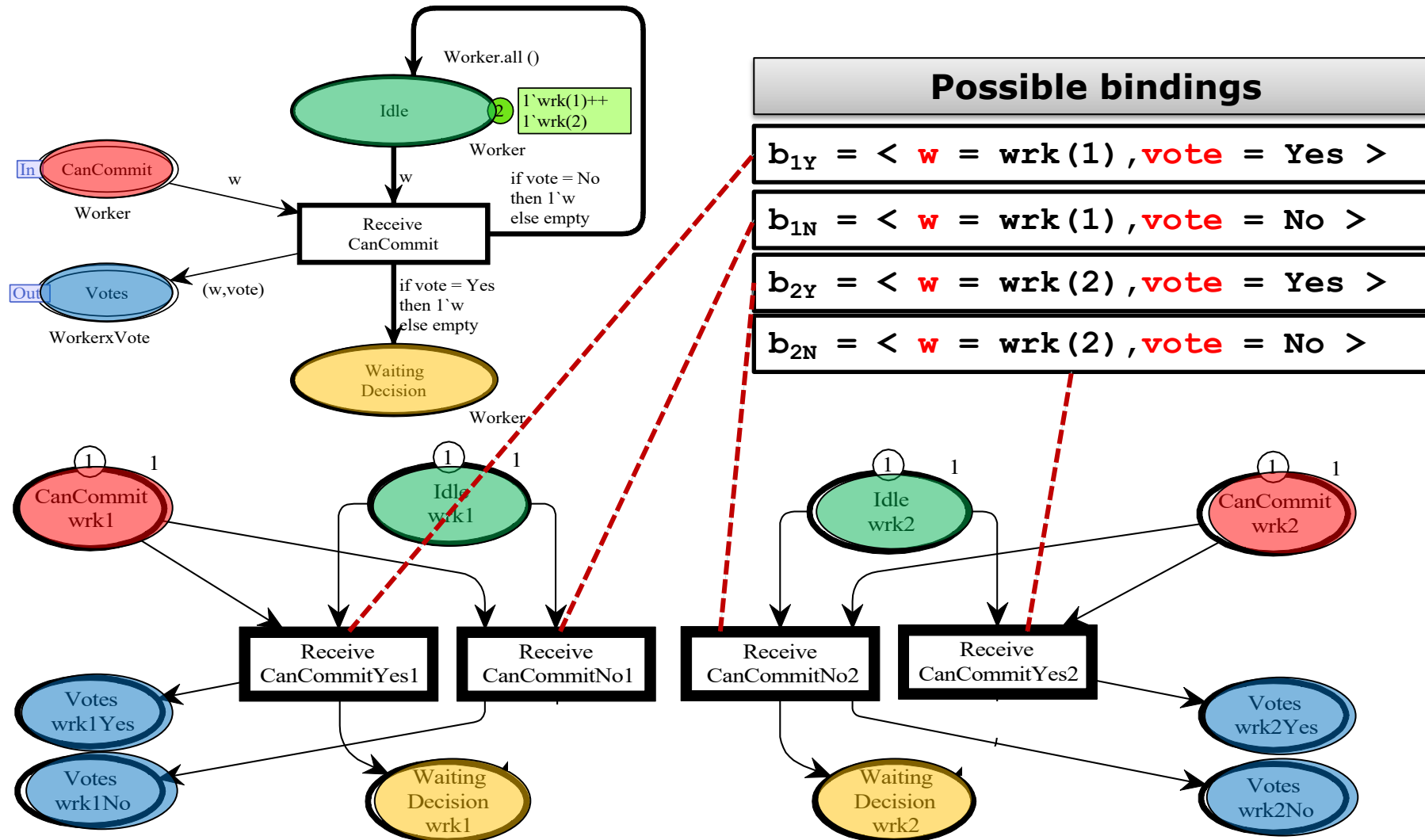
# Unfolding hierarchical CPNs



# Unfolding CPN Places



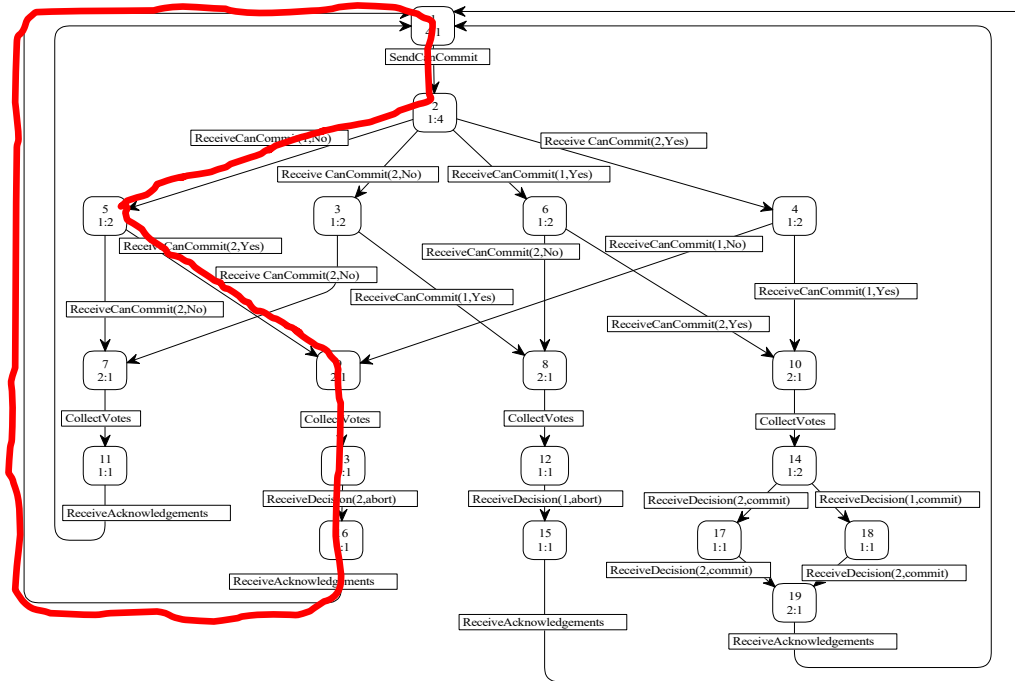
# Unfolding CPN Transitions



# CPN analysis methods

# Verification and model checking

- Formal verification of CPN models can be conducted using explicit state space exploration



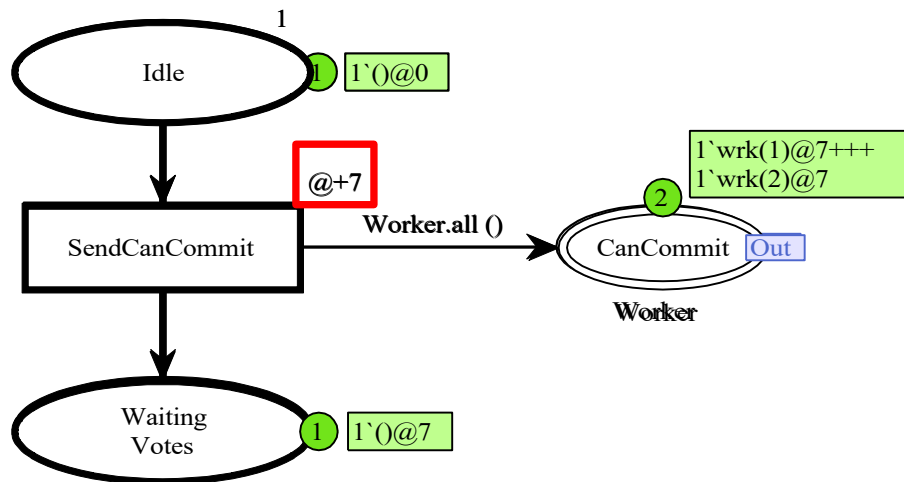
- A state space represents all possible executions of the CPN model
- Standard behavioural properties can be investigated using the state space report
- Model-specific properties can be verified using queries and temporal logic model checking

- Several advanced techniques available to alleviate the inherent state explosion problem



# Performance analysis

- CPNs include a **concept of time** that can be used to model the timed taken by activities



- A **global clock** representing the **current model time**
- Tokens carry **time stamps** describing the earliest possible model time at which they can be removed
- Time inscriptions** on transitions and arcs are used to give time stamps to the tokens produced on output places

- Random distribution functions** can be used in arc expressions (delays, packet loss, ...)
- Data collection monitors** and batch simulations can be used to compute performance figures

# Perspectives on CPNs

- **Modelling language combining Petri Nets with a programming language.**
- **The development has been driven by an application-oriented research agenda:**
- **Key characteristics**
  - Few but still powerful and expressive modelling constructs
  - **Implicit concurrency** inherited from Petri nets
    - everything is concurrent unless explicitly synchronised
  - **Verification** and **performance analysis** supported by the same modelling language

