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Assignment 5

## Public Key Cryptography

### General Idea:

The general idea of this assignment is to implement a basic cryptography library using the RSA algorithm, that can encode and decode messages based on numerical values.

The RSA algorithm will choose two random, arbitrary prime numbers  $p$  and  $q$ , and multiply them together to get  $n$ . Then, an encryption key  $E$  will be chosen, such that  $E$  is coprime (does not evenly divide)  $n$ . Then, our decryption key  $D$  is calculated, based on the equation  $(d * e) \text{ is identical to } 1 \pmod{\text{totient}(n)}$ . Totient is the number of integers between 1 and  $N$  that are coprime to  $n$ .

### Pseudo Code:

#### Greatest Common Denominator

The point of the following function is to compute the greatest common denominator or divisor of the two numbers inputted. This is used in several computations throughout the RSA algorithm.

GCD(storage\_var, a, b) -computes the GCD of a and b, storing it in the storage variable

While (b != 0)

$T = b$

$B = a \% b \text{ //(} a \bmod b \text{)}$

$A = t$

Set storage\_var = a

#### Modular Exponentiation

Modular Exponentiation is a form of calculating exponents that is more efficient than simply repeated multiplication. This works by repeatedly squaring a number, starting at  $n^1$ , and multiplying each result by  $(\bmod n)$  to prevent absurdly large numbers.

Modular\_exponents(a,d,n)

$V = 1$

```

P = a
Counter = d
While counter > 0
    If counter is odd;
         $V = (v * p) \% n$ 
         $P = (p * p) \% n$ 
         $D = d / 2$ 
Return v

```

### Miller-Rabin PseudoPrime Test

This function is designed to test, with a great deal of accuracy, whether a number is prime or not. It uses the Miller-Rabin pseudoprime algorithm, which works by testing whether or not a number is composite, and repeating this process enough times to say that it is in all likelihood prime.

```

is_prime(n,k);
// this first sequence is calculating r and s, such that  $2^s * r = n - 1$ 
F = 4
While (((n-1) % f) % 2 = 0) // while r is not odd
    S = f
     $r = ((n-1) \% f)$ 
     $F = f * 2$  //increase S
For (i = 1, until i = k)
    Rand = random(from 2 to n - 2)
    Y = power_mod(a, r, n)
    If  $y \neq 1$  and  $y \neq n - 1$ 
        J = 1
        While  $j \leq s-1$  and  $y \neq n-1$ 
             $Y = \text{power\_mod}(y, 2, n)$ 
            If  $y = n$ 
                Return false
        J = j + 1

```

```

        If  $y \neq n-1$ 
            Return false
    Return true

```

## Modular Inverses

Computes a value  $i$  such that  $ai = 1 \pmod n$ .

`mod_inverse(i, a, n)`

```

    R = n
    rp = a
    T = 0
    tp = 1
    While rp != 0
        Q = r / rp
        Temp = rp
        Rp = r - (q * rp)
        R = temp
        Temp2 = tp
        tp = t - (q * tp)
        T = temp2
    If r > 1
        Deref i = 0 //no inverse, i = 0
    If t < 0
        T += n
    Set i = t

```

## Make Prime

Generates a random prime number, with a specified bit size, putting it into P address. This takes advantage of the fact that prime numbers can be expressed as  $6(n) + 1$  or  $6(n) - 1$ .

`Make_prime(pointer P, bits, iters)`

```

    R_int = modular_exponents(2, bits) // raise 2 to the power of bits

```

R\_int -= 1 //compute maximum number held in that many bits, thus giving us our bottom range

Rand = random number from r\_int to (r\_int \* 100)

Prime = (rand \* 6) - 1

If is\_prime(prime, iters) == True

Pointer P = Prime

Else if is\_prime((prime+2), iters) == True

Pointer P = (prime + 2)

RSA Make Pub

Makes an RSA public key

rsa\_make\_pub(p, q, n, e, bits, iters)

Iters = random() //random number

P\_Bits = random() in range (bits/4, 3\*bits / 4)

Q\_bits = (3 \* bits / 4) - P\_bits

make\_prime(p, P\_bits)

make\_prime(q, Q\_bits)

Lam\_n = ((P-1) \* (Q-1) / gcd(P-1, Q-1)

For(int e, while gcd(e, lam\_n) != 1)

E = mpz\_randomb(bits)

Return E

Rsa Make Priv

Makes a new private key based off of public key

rsa\_make\_priv(d,e,p,q)

$\text{Lam\_n} = ((P-1) * (Q-1) / \text{gcd}(P-1, Q-1))$

Declare  $d = 0$

$\text{modular\_inverse}(d, e, \text{lam\_n})$

Return  $d$

Writes a public key to a file

$\text{rsa\_write\_pub}(n, e, s, \text{username}, \text{filename})$

$\text{FILE} = \text{fopen}(\text{filename}, \text{w})$

$\text{gmp\_fprintf}(\text{FILE}, \%X, n)$

$\text{fprintf}(\text{FILE}, \%c, \backslash n)$

$\text{gmp\_fprintf}(\text{FILE}, \%X, e)$

$\text{fprintf}(\text{FILE}, \%c, \backslash n)$

$\text{gmp\_fprintf}(\text{FILE}, \%X, s)$

$\text{fprintf}(\text{FILE}, \%c, \backslash n)$

$\text{fprintf}(\text{FILE}, \text{username})$

$\text{fprintf}(\text{FILE}, \%c, \backslash n)$

$\text{fclose}(\text{FILE})$

Reads public key from a file

$\text{rsa\_read\_pub}(n, e, s, \text{username}, \text{filename})$

$\text{FILE} = \text{fopen}(\text{filename}, \text{r})$

$n = \text{gmp\_fscanf}(\text{FILE}, \%X, n)$

$e = \text{gmp\_fscanf}(\text{FILE}, \%X, e)$

$s = \text{gmp\_fscanf}(\text{FILE}, \%X, s)$

$\text{username} = \text{fscanf}(\text{FILE}, \%s, \text{username})$

```
fprintf(FILE, \n)
```

```
fclose(FILE)
```

Writes a private key to a file

```
rsa_write_priv(n, d, filename)
```

```
FILE = fopen(filename, w)
```

```
gmp_fprintf(FILE, %X, n)
```

```
fprintf(FILE,%c, \n)
```

```
gmp_fprintf(FILE, %X, d)
```

```
fprintf(FILE,%c, \n)
```

Reads private key from a file

```
rsa_read_priv(n, d, filename)
```

```
n = gmp_fscanf(FILE, %X, n)
```

```
e = gmp_fscanf(FILE, %X, d)
```

```
username = fscanf(FILE, %s, username)
```

```
return
```

Encrypts a block of data using the c, m, e, n values passed in.

```
rsa_encrypt(c, m, e, n)
```

```
C = modular_exponent(m to the e) mod n
```

```
Return c
```

Encrypts a whole file using the key information provided.

```
Rsa_encrypt_file(FILE *infile, FILE *outfile, n, e)
```

```
//k is the block size
```

```
K = log(n) / log(2)
```

$K = (k - 1) / 8$

Arr = Calloc(uint8\_t array of size k)

Arr[0] = 0xFF

for(char in INFILE, while char != End Of File)

for(j = 1 ; j <= k -1 and char != End Of File; )

Arr[j] = char

Increment char

mpz\_import(store in message, size j, 1, 1, 0, take from block) //number

here represent the data order and nails, etc

rsa\_encrypt(store in cypher, message, e, n)

gmp\_fprintf(OUTFILE, c)

Clear MPZ variables

return;

RSA decrypt- decrypts message m using private exponent d.

Void rsa\_decrypt(message , cypher, d, n)

pow\_mod(store in message: cypher to the d, mod n)

Return;

RSA decrypts a whole file, using rsa\_decrypt

Void rsa\_decrypt\_file(INFILE, OUTFILE, n, d)

J = 0; //used to store size of bytes used later

K = bits\_in(n) //real function is gmp\_sizeinbase()

$K = (k - 1) / 8$

Block = calloc(k, sizeof(char)) // make a block with k elements

Text = fscanf(INFILE)

Do

    decrypt(message m, c,d,n)

    gmp\_export(store in message, cypher text, size j, 1,1,0, take from block)

    for(until counter = j, counter++)

        fprintf(OUTFILE, message);

while(text != EOF)

RSA Sign - makes rsa signature

Void rsa\_sign(s,m, d, n)

    pow\_mod(store in s, m to the d mod n)

RSA Verify- checks if signature matches username

Bool rsa\_verify(m, s, e, n)

    Make mpz\_t t;

    pow\_mod(store in t, s to the e mod n)

    if (t = m)

        Return true;

    If (t != m)

        Return false;