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CSE13s

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All Sorts Of C Code Design Document

General Idea

The overarching plan of this project is to implement several very common sorting algorithms into C code. The algorithms to be implemented are Bubble Sort, Shell Sort, Quick Sort, and Heap Sort. These algorithms will then be run through a test harness with random values created using the Mersenne Twister random number generator, in order to gauge their speed and efficacy. In general, the way the code will function will be to use arrays to categorize the data, and then use the specified sorting method to organize the data by size, in ascending order.

Bubble Sort

Bubble sort functions by iterating through all the elements in the array, and comparing them as pairs to see which element is larger. The larger element is then moved into the rightmost position, ensuring that with each pass the largest element "bubbles" to the top of the array.

Pseudocode:

Define function, accepting one parameter, "arr"

Create a for loop with var i, iterating as many times as the array has indices

Set a checker to 0, to test after each pass if any elements were swapped

Create a for loop with var j, starting at the number of elements, ending at i,
decrementing one with each pass

```
If arr[j] < arr[j-1]

Swap arr[j] and arr[j-1]

Set checker to 1

If checker == 0;

Break

Return the array
```

Shell Sort:

The idea of shell sort is to first pair the elements in an array with elements that are far away from each other. Then, compare the two elements, and swap them if the left is larger than the right. Then, lowering the gap between the elements, shell sort creates a new subarray of multiple elements, and sorts it. This continues until the gap is zero. The idea is to create partially sorted arrays that make a different sorting method, in this case insertion sort, much more efficient.

```
Pseudocode:
Gap calculation (generator in c)
Define gap calc(n)
       If N == 1:
               Return 0
       Else if If N \le 2
               N = 1
       Else:
               N = (N * 5) // 11
       Return N;
Def Shell Sort(array);
       Start = (length(array) *5) // 11
       Gap = start
       For Loop, until gap <=0, incrementing gap calc(gap) each pass
               I = gap
               For loop, iterating until i == len(array) - 1, incrementing 1 each pass
                       J = I
                       For loop, iterating until j = gap - 1, decrementing by gap each pass
                              If array[i] < array[i - step],
                                      Swap the two values array[j-step] = array[j]
```

Quicksort:

The idea of quick sort is pick a pivot point at some point in the array, then sort the data into sections based on whether they are greater than or less than the pivot. Then, the function performs this again on the smaller arrays, splitting them and sorting them. This happens until the elements are split down into arrays that are sufficiently small, which are then sorted by shell sort.

Pseudocode:

```
SMALL = 8 (arbitrary number of where we stop sorting)
Define Quicksort (pointer, length of array)
       Length = length of array
       array = pointer to where the array is
       If length < SMALL;
               shellsort(array, length)
               Return 0;
       Pivot = array[(length/2)]
While(true)
       Lswap addr = length/2
       Rswap addr= length/2
       Swapped = false
       For L in range(array[0], Pivot)
               If array[L] >= pivot
                      L \text{ swap} = L
                      break
       For R in range(length, pivot, -1)
               If array[R] < pivot
                      R \text{ swap} = R
                      Break
       If 1 swap != r swap
               Swap array[1] and array[r]
               Set swapped = true
       If swapped = false;
               break
Quicksort((pointer to array[0]), pivot)
Quicksort((pointer to pivot, pivot)
Returns 0
```

Heapsort:

The idea of heapsort is to store data in a structure known as a "heap", which follows basic rules that ensure that the data is an understandable pattern and is sorted. For a min heap, in which the smallest element is at the top, the rules are as follows; Every element must be smaller than its two children. The smallest of any two children should be in the leftmost position. To add elements to the array, add them on the bottom right. If the added element is smaller than its parent, swap the two. Continue until the min heap rules are obtained. To remove an element (in a sorted fashion), remove the top element, replacing it with the smaller of its two children. Continue this down the heap until the heap rules are obtained. Once every element has been removed, a sorted array will have been produced.

Pseudocode;

```
//address of left child
1 child(addr)
       Return (addr *2) +1
//address of right child
r child(addr)
       Return (addr * 2) +2
//address of parent(addr)
parent(addr)
       Return (n-1) // 2
//function to move element up the heap
//moving something up the heap
Up heap(addr, array)
       While addr > 0 and array[addr] < array[parent(addr)]
              Temp = array[addr]
              Array[addr] = array[parent(addr)]
              array[parent(addr)] = temp
              Addr = parent(addr)
//moving an element down the heap
Down heap (array, heap size):
       Addr = 0
```

```
While 1 child(addr) > heap size;
                      If 1 child(addr) == heap size //if there's no left right child, the left is
smaller
                              Smaller = r child(addr)
                      Else:
                              If l child(addr) \leq r child(addr)
                                      Smaller = 1 \ child(addr)
                              Else:
                                      Smaller = r \ child(addr)
                      If array[addr] < array[smaller]
                              Break
                      Temp = array[addr]
                      array[addr] = array[smaller]
                      array[smaller] = temp
                      Addr = smaller
       //building the heap
       Build heap(array, length of array)
               Length = length of array
               Allocate space for the size of the array
               Set that allocated space to heap
               For addr in range(0, length -1)
                      Heap[addr] = array[addr] //copy array to the heap
                      up heap(addr, heap)
               Return a pointer pointing to heap
       //sort the aforementioned heap
       sort heap(array, length)
               Build heap with array and length
               Create a new sorted array with size length
               For arr in range(0, length)
                      Sorted[arr] = heap[0]
```

```
Heap[0] = heap[length - arr - 1]
down_heap(heap, length - arr)
```

Return pointer to sorted

Test Harness Code

The test harness is designed to test and catalog the performance of the sorting algorithms, as well as to display the efficacy of all the algorithms when pitted against one another.

Specify the possible command line inputs: "asbqhr:n:p:H"

Define main function, with arguments (int argc, char **argv)

```
//Set default values
A arg = 0
S arg = 0
B arg = 0
Q arg = 0
h arg = 0
R_arg = 13371453
N arg = 100
P arg = 100
H arg = 0
Set = empty set() //Make a set using empty set()
//parse arguments
For loop, with arg = 0, until arg = (argc-1), incrementing 1
       insert set(argv[arg], set) //put argument from argv into the set
Another loop, looping until counter = argc
       If (member set(a, set) = true)
               A arg = 1
               delete set(a, set)
       If (member set(s, set) = true
               s arg = 1
               delete set(s, set)
       If (member set(b, set) = true
               b arg = 1
```

```
If (member set(q, set) = true
               q arg = 1
               delete_set(q, set)
       If (member set(h, set) = true
               h arg = 1
               delete set(h, set)
       If (member_set(r, set) = true)
               r_arg = strtoul(optarg,NULL,10)
               delete set(r, set)
       If (member set(n, set) = true)
               n_arg = atoi(optarg)
               delete set(n, set)
       If (member set(p, set) = true
               p arg = atoi(optarg)
               delete set(p, set)
       If (member set(H, set) = true)
               H arg = 1
               delete set(H, set)
//execution of code
mtrand_seed(s_arg)
allocate space(n arg)
Set pointer ptr to allocated space
For loop, counter = 0, until counter = n arg-1
        *(ptr + counter) = mtrand rand64()
If a arg = 1:
       print(bubble(prt, n arg))
       print(shell(prt, n_arg)))
       print(quick_sort(prt, n_arg)))
       print(heap_sort(prt, n_arg)))
If b arg = 1:
```

delete set(b, set)