

Recitation

CFG and CKY

11/12/2020

Param P

CFG

- Set of nonterminal symbols N
- Set of terminal symbols Σ
- A start symbol $S \in N$
- Set R of productions of the form $A \rightarrow B$ where $A \in N$ and $B \in (\Sigma \cup N)^*$ ie. B is a string of terminals and non-terminals
- Given G is a CFG, language of G is defined as set of all terminal strings that can be derived from the start symbol
- **Context free:** applicability of a rule depends only on a nonterminal symbol, not on its context

Regular Grammar

- Set of nonterminal symbols N
- Set of terminal symbols Σ
- A start symbol $S \in N$
- Set R of productions of the form $A \rightarrow aB$ or $A \rightarrow a$ where $A, B \in N$ and $a \in \Sigma$
- Can be implemented as finite state automata
- Set of all regular languages is strictly smaller than set of all context-free languages
- Problem: regular languages cannot capture long-distance dependencies

CKY Parsing

- Syntactic parsing is the task of recognizing a sentence and assigning a syntactic structure to it (i.e. adding POS tags).
- The **CKY algorithm** is a bottom-up approach to figuring out if a sentence/sequence of words is in the language/grammar provided, because it starts with the terminals and sees which subtrees it can build, eventually making those subtrees into larger subtrees until you reach S (or don't, in which case it's not in the language).
- To use the CKY algorithm, the CFG must be in **CNF**.

CNF

Chomsky Normal Form

- A CFG $G=(N, \Sigma, R, S)$ is in Chomsky Normal Form (CNF) if the rules take one of the following forms:
 - $A \rightarrow B C$, where $A \in N, B \in N, C \in N$.
 - $A \rightarrow b$, where $A \in N, b \in \Sigma$.

$S \rightarrow NP VP$	$V \rightarrow \text{saw}$
$VP \rightarrow V NP$	$P \rightarrow \text{with}$
$VP \rightarrow VP PP$	$D \rightarrow \text{the}$
$PP \rightarrow P NP$	$N \rightarrow \text{cat}$
$NP \rightarrow D N$	$N \rightarrow \text{tail}$
$NP \rightarrow NP PP$	$N \rightarrow \text{student}$

Any CFG can be converted to an equivalent grammar in CNF that expresses the same language.

Cannot have null rules and unreachable rules as well

CNF Example

CFG \rightarrow CNF

• $S \rightarrow bC$

• $C \rightarrow B$

• $B \rightarrow bAE$

• $A \rightarrow bE$

• $E \rightarrow e$

CFG \rightarrow CNF

- $S \rightarrow FC$

- $C \rightarrow FAE$

- $A \rightarrow FE$

- $E \rightarrow e$

- $F \rightarrow b$

Replace 'b' by non-terminal 'F' by adding a new rule

Merge B into C since C has only 1 non-terminal
On the right side.

CFG \rightarrow CNF

- $S \rightarrow FC$
- $C \rightarrow FG$
- $G \rightarrow AE$
- $A \rightarrow FE$
- $E \rightarrow e$
- $F \rightarrow b$

C had 3 non-terminals on the right. Take the last \exists 2
And make a new rule 'G' for them.

Now each NT has only 2 Non-Terminals on the right
or only 1 Terminal.

This is final CNF form

CKY Example

CFG

$S \rightarrow NP \quad VP$

$NP \rightarrow D \quad N$

$VP \rightarrow V \quad Adj$

$VP \rightarrow V \quad NP$

$D \rightarrow the$

$N \rightarrow \text{snow} \quad snow$

$V \rightarrow snow$

$N \rightarrow clouds$

$V \rightarrow clouds$

$N \rightarrow sky$

$Adj \rightarrow sky$

Sentence: the snow clouds the sky

	the	snow	clouds	the	sky	
	0	1	2	3	4	5
0						
1						
2						
3						
4						
5						

		the	snow	clouds	the	sky	
		0	1	2	3	4	5
0	X						
1	X						
2	X						
3	X						
4	X						
5	X	X	X	X	X	X	

CFG

$S \rightarrow NP \quad VP$
 $NP \rightarrow D \quad N$
 $VP \rightarrow V \quad Adj$
 $VP \rightarrow V \quad NP$
 $D \rightarrow the$
 $N \rightarrow \text{snow} \quad snow$
 $V \rightarrow snow$
 $N \rightarrow clouds$
 $V \rightarrow clouds$
 $N \rightarrow sky$
 $Adj \rightarrow sky$

		0	1	2	3	4	5
0	X	D					
1	X						
2	X						
3	X						
4	X						
5	X	X	X	X	X	X	X

CFG

$S \rightarrow NP \quad VP$
 $NP \rightarrow D \quad N$
 $VP \rightarrow V \quad Adj$
 $VP \rightarrow V \quad NP$
 $D \rightarrow the$
 $N \rightarrow \text{snow} \quad snow$
 $V \rightarrow snow$
 $N \rightarrow clouds$
 $V \rightarrow clouds$
 $N \rightarrow sky$
 $Adj \rightarrow sky$

		the	snow	clouds	the	sky	
	0	1	2	3	4	5	
0	X	D					
1	X	X	N, V				
2	X	X	X	N, V			
3	X	X	X	X	D		
4	X	X	X	X	X	N, Adj	
5	X	X	X	X	X	X	

CFG

S	→	NP	VP
NP	→	D	N
VP	→	V	Adj
VP	→	V	NP
D	→	the	
N	→	snow	snow
V	→	snow	
N	→	clouds	
V	→	clouds	
N	→	sky	
Adj	→	sky	

			the	snow	clouds	the	sky	
		0	1	2	3	4	5	
0	X	D	→	NP				
1	X	X		N, V	X			
2	X	X		X	N, V	X		
3	X	X		X	X	D	→	NP
4	X	X		X	X	X		N, Adj
5	X	X		X	X	X		X

CFG

S	→	NP	VP
NP	→	D	N
VP	→	V	Adj
VP	→	V	NP
D	→	the	
N	→	snow	snow
V	→	snow	
N	→	clouds	
V	→	clouds	
N	→	sky	
Adj	→	sky	

		the	snow	clouds	the	sky
	0	1	2	3	4	5
0	X	D → NP	X			
1	X	X	N, V	X	X	
2	X	X	X	N, V	X	VP
3	X	X	X	X	D → NP	
4	X	X	X	X	X	N, Adj
5	X	X	X	X	X	X

CFG

S	→	NP	VP
NP	→	D	N
VP	→	V	Adj
VP	→	V	NP
D	→	the	
N	→	snow	snow
V	→	snow	
N	→	clouds	
V	→	clouds	
N	→	sky	
Adj	→	sky	

		0	1	2	3	4	5
			the	snow	clouds	the	sky
0	X		D → NP	X	X		
1	X	X	N, V	X	X	X	
2	X	X	X	N, V	X	VP	
3	X	X	X	X	D → NP		
4	X	X	X	X	X	N, Adj	
5	X	X	X	X	X	X	

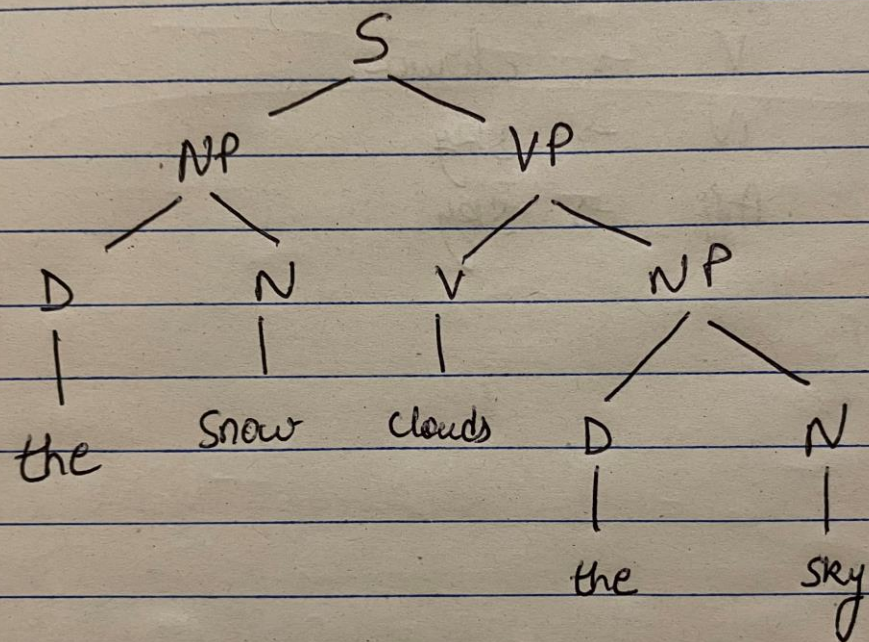
CFG

S	→	NP	VP
NP	→	D	N
VP	→	V	Adj
VP	→	V	NP
D	→	the	
N	→	snow	snow
V	→	snow	
N	→	clouds	
V	→	clouds	
N	→	sky	
Adj	→	sky	

		0	1	2	3	4	5
			the	snow	clouds	the	sky
0	X	D	→ NP	X	X	S	
1	X	X	N, V	X	X	X	
2	X	X	X	N, V	X	VP	
3	X	X	X	X	D	→ NP	
4	X	X	X	X	X	N, Adj	
5	X	X	X	X	X	X	

CFG

S	→	NP	VP
NP	→	D	N
VP	→	V	Adj
VP	→	V	NP
D	→	the	
N	→	snow snow	
V	→	snow	
N	→	clouds	
V	→	clouds	
N	→	sky	
Adj	→	sky	



CFG

$S \rightarrow NP \quad VP$
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 $N \rightarrow sky$
 $Adj \rightarrow sky$