▼ Eligenvalue Problems

Aim: Given a square matrix A, find eigenvalues/ eigenvectors of A. (nxn matrix)

1 Review/definitions

• Given a square matrix A, we define the eigenvectors v & associated eigenvalues A the solutions of Av = Av

The set of all eigenvalues is the spectrum of A

- one eigenvalue may have more than I eigenvector associated to it. The subspace spanned by these agenvectors is called the eigenspace associated to A.
- To find $A & \underline{V}$, we must solve $(A IA)\underline{V} = 0$

which only has non-mirial solutions if det (A-IZ)=c

This yields an n-th order polynomial (if t is nxn) equation called the characteristic equation

$$P_A(A) = 0$$

Note that if A is real it may still have complex eigenvalues.

 $P_{A}(x) = K(\lambda - \lambda_{1}) - \cdots (\lambda - \lambda_{n})$ where $A: \in \mathbb{C}$ and we allow for $\lambda_{i} = \lambda_{j}$ who define the algebraic multiplicity of a particular eigenvalue λ_{k} as the number of factors $(\lambda - \lambda_{k})$ in $P_{A}(\lambda)$.
We define the geometric multiplicity of λ_{k} as the dimension of the eigenspace associated with λ_{k} .

The algebraic multiplicity is always greater or equal to the geometric multiplicity. If the algebraic multiplicity is strictly greater than the geometric multiplicity than the eigenvalue Ax is called defective.

A matrix which has at least one defective eigenvalue is called a <u>defective matrix</u>

Example: The matrix $\begin{pmatrix} 1 & 1 \\ -1 & 3 \end{pmatrix}$ has an

étaphralue $\lambda = 2$ with algebraic multiplicity 2 bout only one eigenvector v = (1)

• Rank 2 nullity: The <u>nullity</u> of A is the demension of the eigenspace of the null eigenvalue. The rank of A is (n-nullity)

2 Analytics vs numerics

When seeking Evalues & Evectors analytically, one usually first solves the characteristic equation to find the Evalues and then solve, for each ax,

AVK = AKVK

to find the Evector (or subspace) associated with the

Solving high-order polynomical equations numerically in practise, is actually quite difficult. Threaver, once a solution is found, its algebraic multiplicity remains unknown.

=> Finding Evalues & Evectors numerically requires thinking about the porobern entirely differently.

3 Similarly transformations/change of basis

A matrix actually represents an internet geometric operation on vectors; as such, the properties of that geometric operation should be the same regardless of the coordinate system it is expressed "th.

- => A change of basis (of coordinates) is usually referred to as a simularity transformation;
- > The proporties of the transformed matrix (i.e. the matrix expressed in the new coordinate system) are invariant, and in particular
 - ⊕ the determinant 2 the trace are invariant
 - 1 the characterstic polynomial, the eigenvalues 2 eigenvectors are invariant The rank 8 nullity are invariant The rank 2 nullity are invariant

nathemanically:

- A set of n vectors form a basis for R" provided they are linearly independent (note that they don't need to be orthogonal).
- Suppose we have such a set $\{e_1, \dots, e_n\}$ then a vector \underline{x} has coordinates (x_1, \dots, x_n) in the barns formed by $\{e_i\}$ implies that $\underline{X} = \sum_{i=1}^{n} x_i e^{ix}$

Suppose we now have another (new) basis
$$g_1, \dots, g_n g$$
, and the vector g_1 has

the new coordinates $\{x_i', \ldots, x_n'\}$ in the new booss

However,
$$n_{k} = \sum_{i=1}^{n} (n_{k} \cdot e_{i}) e_{i} = \sum_{i=1}^{n} (n_{k}^{T}e_{i}) e_{i}$$

(effectively the j-th coordinate of n_{k} in the fe_{i} ? basis

$$X = \sum_{i=1}^{n} x_{i} e_{i} = \sum_{k=1}^{n} x_{k}' \sum_{i=1}^{n} (n_{k}^{T}e_{i}) e_{i}$$

$$x_{i} = \sum_{k} x_{k}' (n_{k}^{T}e_{i})$$

Let N be the matrix with coefficient $N_{ij} = n_{i}^{T}e_{i}$

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$$x_{i} = \sum_{k} x_{k}' (n_{k}^{T}e_{i})$$
The advance of N_{ij} are the coordinates of N_{ij} and N_{ij} are N_{ij} are the coordinates of N_{ij} are N_{ij}

A and A', share all of the properties listed earlier

= $A' = N^T A N$

system provided

 $A' = N^T AN$

=> If A & A' shore the same Evectors & Evalues, then by finding an appropriate coordinate change we may be able to "reveal" the eigenvalues of A'.

A well-known bout not very useful example

. Given a real, symmetric matrix A, if we knew the eigenvectors (which we don't), we could unemed the matrix

$$V = \begin{pmatrix} 1 & 1 \\ V_1 & V_2 \end{pmatrix}$$
 = matrix with column vectors the Evectors of A

then $V^TAV = A' = \begin{pmatrix} \lambda_1 \\ 0 \end{pmatrix}$ change of coordinate of the new basis formed $\begin{pmatrix} \lambda_1 \\ 0 \end{pmatrix}$

where quite donously, the eigenvalues of A are the same as the eigenvalues of A.

ASIDE . Note the following theorems

A is non-defective

$$\begin{pmatrix} A_1 & O \\ O & A_n \end{pmatrix} = V^{-1}AV$$

where V is formed by the normalized E-vectors of A (noe λ : and v: can be complex)

- 2) A is non-defective & AA* A*A (=> VV* = I
- A is non-defective & AAT-ATA (=) VVT = I (₹)

(4) The Schur factorization.

. If A is an nxn matrix with Evalues $\{\lambda_i\}$ $\{\lambda_i \in \mathcal{C}\}$ then there exists a unitary matrix P such that

where
$$T = \begin{pmatrix} 2_1 \times \cdots \times \\ & &$$

- To prove this statement, we will construct an algorithm to create this factorization. But note that once T is known, the $\{\lambda_i\}$ are on the diagonal and can directly be read. Moreover this reveals the algebraic multiplicity of the Evalues.
- Another factorization is also very verful if we only want to work with real matrices: the real Schur factorization $A = Q T Q^T$ where Q is a real

orthogonal matrix and it is a real upper mangular matrix except for ax2 blocks around the diagranal containing the real and imaginary ports of the complex Evalues

$$\tilde{T} = \begin{pmatrix} A_1 \times X \times X \times X \\ Re\lambda_2 & Im\lambda_2 \times X \end{pmatrix}$$
 (eg here with A_i real $Im\lambda_2 & Re\lambda_2 \times X \end{pmatrix}$ except for A_2 and its complex conjugated A_3)

- . Note that both factorizations still work if the eigenvalues are <u>defective</u>
- . To construct the Schur factorization we will first need to learn about the Hessenberg form.

The Hessenberg form

. It is always possible to find an oithogonal coordinate mansformation

where
$$A = PHP^{T} \quad (\text{equivalently} \quad H = P^{T}AP)$$

$$Where \quad H = \begin{pmatrix} x & x & --- & x \\ x & x & x \\ 0 & x & x \\ 0 & x & x \end{pmatrix} \quad (i.e. \ H \text{ has} \quad Hij = 0 \text{ if} \quad (i.e. \ H \text{ has} \quad H \text{ has} \quad H \text{ has} \quad H \text{ has} \quad (i.e. \ H \text{ has} \quad H \text{ ha$$

To prove this statement, we will construct an algorithm to put A into a Hebbenberg form.

(5) Searching for ligenvalues: a 2-step process

- ⇒ We will from now on construct an algorithm which reveals the (complex) eigenvalues of any matrix A by reduction onto a Scher form (of the real kind).
- > The standard algorithm for doing this counsts in 2 steps

Step 1: Direct reduction of A to its Hessenberg form

Step 2: Iterative convergence of the initial Hessenberg form towards the real Schur form.

Why do this 2-step process?
One could thank of vsing thouseholder matrices to reduce A to Schur form directly
of first column
PPA - of second column
\$ n-1 times
then multiply to right by
P(n-1) P(n) P(n) P(n+1) T
Problem. The multiplication to the right repopulate the zeros with non-zero dements!
Example PDA acts on the whole matrix A:
$\begin{pmatrix} s_1 \times \times \\ 0 \times \end{pmatrix}$
influence region of Pivon A
$\sqrt{0} \times \times \sqrt{1}$
so PAP also acts on whole married
(but PT is not the correct thouseholder mater)
for P'A so it just repopulates the zeros
Idoa: Instead consmot the first thouseholder
matrix to zero elements only below the
Subdiagonal
PA (XXXXXX) < untouched row
mfluence block of
10 Y X Pinon A

This constroic straightforwardly leads to the reduction of A through to a Hessenberg matix which will then need further work to get to Schur form.

6 Algorithm(s) for Hessenberg reduction

There are Two possibilities for doing this, in fact:

- Procedure analogous to Gaussian elimination (see Numerical Recipes). Faster, but not necessarly always stable
 - · Householder reduction (as above). Almays stable.

The algorithm for reduction into Hessenberg form using the Hotseholder matrices is quite similar to the one used in QR decomposition. The main difference are

- . A must be successively multiplied by $\widetilde{P}^{(k)}$ on the left and by $\widetilde{P}^{(k)}$ on the right.

 The $\widetilde{P}^{(k)}$ are slightly different from the
 - P(K) of the QR decomposition.
- Since the actual $\stackrel{\sim}{P}$ are not needed later (only H is), we return the full Hessenberg form into A.

Algorithm (for an nxn mathx
$$A$$
)

do $k = 1$, $n-2$

consider $x' = (0,0,...,0, a_{kH,k}, ..., a_{N,k})$
 $Sk = \pm 1x^{(k)}1$
 $Ck = Sk(Sk-a_{kH,k})$
 $A = \widetilde{\rho}^{(k)}A$
 $A = A\widetilde{\rho}^{(k)}T$

(**)

endo.

The step (*) depends on whether A is on ext entirely replaced by H (in which case the info on the \widetilde{W} vectors is lost) or M as in the QR

The step (*) depends on whether A is on extention replaced by H (in which case the info on the M vectors is lost) or if, as in the OR factoritation the W vectors are strict below the subdiagonal of H whereas the subdiagonal H self is returned as separate vector.

$ \begin{pmatrix} x & x & x & - & - & x \\ S_1 & x & & & & \\ 0 & S_2 & & & & & \\ & & & & & & \\ \end{pmatrix} $	t A Is
	X\
	·
] X /
	i Te January Santa Japanese P
ũ vector	

and the subdigonal of H

(S, -- Sn-1) is returned

Separately.

$$\begin{array}{c} \text{(it):} & \text{$a_{k+1}, k = 8_{k}$} \\ \text{$a_{k+2:n, k} = 0$} \\ \text{$doj = k+1, n$} \\ \text{$doi = k+1, n$} \\ \text{$a_{ij} = a_{ij} - 2_{ij}$} \\ \text{$a_{ik} = a_{mk} = a_{mk}$} \\ \text{$a_{ik} = a_{mk}$} \end{array}$$

Note that neither six non cir needs to be returned.

option 2 $\begin{cases} a_{k+1}, k = a_{k+1}, k - s_k \\ do j = k+1, n \\ do t = k+1, n \end{cases}$ $\begin{vmatrix} a_{k+1}, k = s_k \\ do j = k+1, n \\ a_{k+1}, n = s_k \end{vmatrix}$ $\begin{vmatrix} a_{k+1}, k - s_k \\ do j = k+1, n \\ a_{k+1}, n = s_k \end{vmatrix}$ $\begin{vmatrix} a_{k+1}, k - s_k \\ do j = k+1, n \\ a_{k+1}, n = s_k \end{vmatrix}$ $\begin{vmatrix} a_{k+1}, k - s_k \\ do j = k+1, n \\ a_{k+1}, n = s_k \end{vmatrix}$ $\begin{vmatrix} a_{k+1}, k - s_k \\ do j = k+1, n \\ a_{k+1}, n = s_k \end{vmatrix}$ (Note that sk must be stored & returned as separate rector, simularly for ck) Step (**) do i=1, n for both $|a_j| = a_{ij} - \sum_{m=k_{ij}}^{n} a_{mk} a_{jk} a_{im}$ All of the above applied to general matrices Jf A is a real symmetric matrix their the Hessenberg form is inclusional and the Schur form is the deaponal matrix of the eigenvalues. In that case too P is the matrix of the orthogonal eigenvectors