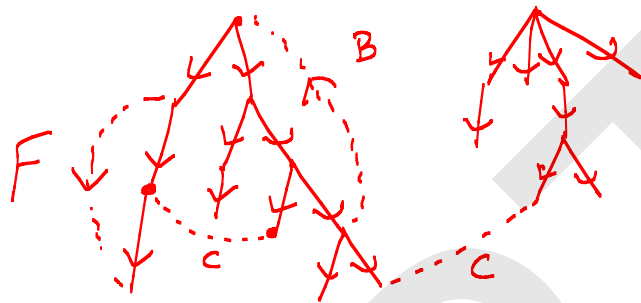
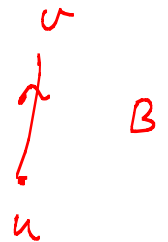
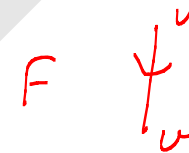


DFS in Directed Graph.



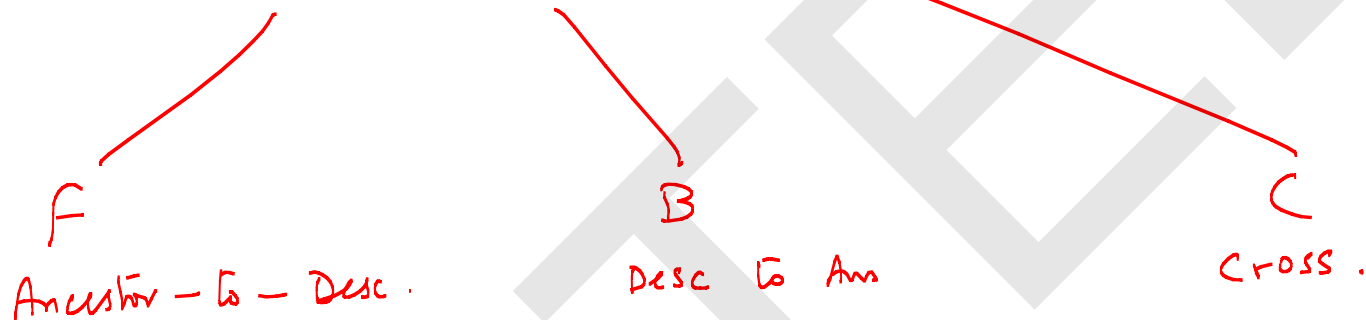
$$e = (u, v)$$



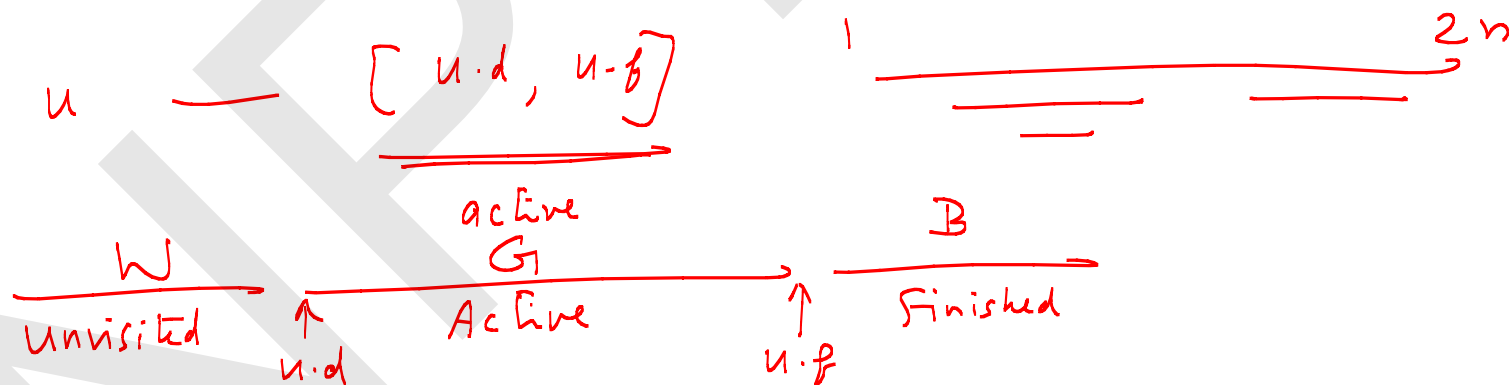
Visited node \longrightarrow unvisited node — Tree edge.

Visited node \longrightarrow visited node — Non Tree edges.

Non Tree edges



Time — 1 to $2n$ consecutive integers.



u

v

$u.d$ u $u.f$

$v.d$ v $v.f$



$$u.d < v.d < v.f < u.f$$

v is desc of u .

DFS (G)

For each $u \in V$
 $u.col = white$
 $u.p = NULL$.

Time = 0

For each $u \in V$
 if ($u.col == white$)
 DFS (G, u).

DFS (G, u)

Time = Time + 1;

$u.d = Time$;

$u.col = Gray$;

For each $v \in Adj[u]$

 if ($v.col == white$)

$v.p = u$;

$\parallel (u, v)$ is a Tree
 edge .

 DFS (G, v).

 else $\parallel (u, v)$ is a
 Non-Tree edge .

elaborations
of
Control point
related to
Non-Tree edges.

else $\parallel (u, v)$ is a non tree edge.

if $(v.col == Gray)$

$\parallel (u, v)$ is a back edge. \leftarrow

if $(v.f < u.d)$

$\parallel (u, v)$ is a cross edge. \leftarrow

else

\parallel forward edge.

Finding strongly connected components.
of a Directed Graph. }