

The Capacity of the Kanerva Associative Memory is Exponential

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Abstract:

The capacity of an associative memory is defined as the maximum number of words that can be stored and retrieved reliably by an address within a given sphere of attraction. It is shown by sphere packing arguments that as the address length increases, the capacity of any associative memory is limited to an exponential growth rate of $1 - h_2(0)$, where $h_2(0)$ is the binary entropy function in bits, and 0 is the radius of the sphere of attraction. This exponential growth in capacity can actually be achieved by the Kanerva associative memory, if its parameters are optimally set. Formulas for these optimal values are provided. The exponential growth in capacity for the Kanerva associative memory contrasts sharply with the sub-linear growth in capacity for the Hopfield associative memory.

ASSOCIATIVE MEMORY AND ITS CAPACITY

Our model of an associative memory is the following. Let (X, Y) be an (address, datum) pair, where X is a vector of n bits and Y is a vector of m bits, and let $(X(1), Y(1)), \dots, (X(M), Y(M))$ be M (address, datum) pairs stored in an associative memory. If an address X that is close to some stored address $X(j)$ is presented at the input, then it should produce at the output a word Y that is close to the corresponding contents $Y(j)$. To be specific, let us say that an associative memory can correct fraction ϵ errors if an address X within Hamming distance ϵn of $X(j)$ retrieves Y equal to $Y(j)$. The Hamming sphere around each $X(j)$ will be called the sphere of attraction, and ϵn will be called the radius of attraction. If the associative memory can correct fraction ϵ errors, then the capacity of this associative memory is the maximum number of words that it can store while correcting fraction ϵ errors. Unfortunately, this notion of capacity is ill-defined, because it depends on exactly which (address, datum) pairs have been stored. Clearly, no associative memory can correct fraction ϵ errors for every sequence of stored (address, datum)

pairs. Consider.

for example, a sequence in which several different words are written to the same address. No memory can reliably retrieve the contents of the overwritten words. At the other extreme, any associative memory ' can store an unlimited number

of words and retrieve them all reliably, if their contents are identical. A useful definition of capacity must lie somewhere between these two extremes, that for most sequences of addresses $X(1), \dots, X(M)$ and most sequences of data

$y(1), \dots, y(M)$, the memory can correct fraction θ errors. We define In this paper, we are interested in

the largest M such that This work was supported by the National Science Foundation under NSF grant IST-8509860 and by an

IBM Doctoral Fellowship. © American Institute of Physics 1988 185 I most sequences' in a probabilistic sense, as some set

of sequences y_i with total probability greater than say, .99. When all sequences are equiprobable, this reduces to the deterministic

version: sequences.