Muth 5600

Recall the QR algorithm

For k = 1,2,...

A = QR

A = RQ

equivalent to a simultaneous orthogonal iterations.

If the eigenvulues are strictly ordered

 $|\lambda_1| > |\lambda_2| > \cdots > |\lambda_n|$

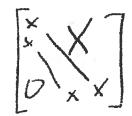
(*)

this process converges to a triangular matrix that is similar to A and has the visenessure of A along the diagonal.

However, we need to make the algorithms move efficient and we need to get it to work without the assumption (*).

We will start with transforming A to upper Hessenburg form.

A is u.H (=) i-j>1 = aij=0



$$H_o = U_o^T A U_o$$

UT = WO

How do we do this?

1st column

$$H = \begin{bmatrix} 1 & 0 \\ 0 & I - 2vv^T \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \\ 0 & H \end{bmatrix} \begin{bmatrix} a_{11} & y^T \\ z & A_{22} \end{bmatrix}_{\eta-1} = \begin{bmatrix} a_{11} & y^T \\ Hz & HA_{22} \end{bmatrix}$$

- pich H such that Hz 13 a multiple of e,
- multiplying with HT from the right to get a similarity transform) does not undo this effect.

- Then repeat the process on HAzz HT as usual.
- Note that creating zeros below the diagonal gets undone by multiplication from the right.

- upper Hessenberg is the best we can
 do.
- of course if we could get triangular form we'd solve in a finite # of steps

Does the QR iteration preserve the upper Hessenberg structure?

yes

$$H = QR$$

$$Q \text{ is upper Hessenbrg}$$

$$\begin{bmatrix} x \times x & x \\ x \times x \\ x \times x \\ 0 & 0 & x \end{bmatrix} \begin{bmatrix} x \times x \times x \\ x \times x \times x \\ 0 & 0 & x \end{bmatrix} R$$

$$Q$$

$$Q$$

$$H$$

use induction on the columns

- so now we have the algorithm

$$H_{k-1} = U_k R_k \qquad QR_1 \quad U_k \quad u.H-$$

$$H_k = R_k \quad U_k$$

- We would like this to converge to upper quesi-triangular form

- quasi-triangular = block triangular with 1x1 and 2x2 triangular blocks

$$Q^{T}AQ = \begin{bmatrix} R_{11} & R_{1m} \\ R_{22} & \\ \\ R_{mm} \end{bmatrix}$$

- the For this to be useful, u.H. QR factorization of an matrix should take cally O(n2) eperations.
- Use Givens Rotations.
- explain for 2x?

$$G = \begin{bmatrix} \omega_{0} + \omega_{0} \\ -\sin \theta \end{bmatrix}$$
 $C = \omega_{0} + \delta$
 $S = \sin \theta$

- rotates by d

$$\begin{bmatrix} c & -5 \\ 5 & c \end{bmatrix}$$

$$- G^{-1} = G^{T} \text{ since } \begin{bmatrix} c & 5 \\ -5 & c \end{bmatrix} \begin{bmatrix} c^{2}is^{2} & 0 \\ 0 & c^{2}is^{2} \end{bmatrix} = \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$$

- pich of such that
$$\begin{bmatrix} c & s \\ -s & c \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} x \\ 0 \end{bmatrix}$$

$$-5x + cy = 0$$

$$\frac{s}{c} = \frac{\gamma}{x}$$

So we embedd Givens Rotations into the diagonal as usual

$$\begin{bmatrix} \times \times \\ \times \times \\ \end{bmatrix} \begin{bmatrix} \times \times \times \times \\ 0 \times \times \times \\ 0 & 0 \times \times \end{bmatrix}$$

$$\begin{bmatrix} 1 & & & & \\ & \times & \times & \\ & & \times & \times \end{bmatrix} \begin{bmatrix} \times & \times & \times & \\ 0 & \times & \times & \times \\ 0 & 0 & 0 & \times \end{bmatrix}$$

QR factorization

Q: product of Givens Rotations.

general Gr. R.

requires only O(4?) operations.

- critical aspect of QR algorithm.

Decoupling.

$$H = \begin{bmatrix} H_{11} & H_{12} \\ O & H_{22} \end{bmatrix}$$

H,, H22 upper Hessenseng

- An uH matrix is unreduced if all its subdinguisal estries are non-are

$$H - MI = UR$$
 QR fact.
 $H = RU + MI$

- This generals a sequence of similar

- How do we pick pe

- decoupling occurs when M is an inveduced.

H-MI = UR is singular

U is orthogonal => R is singular

- In fact, rnn = 0

- n hy?

- For an unreduced upper Hessenseng matrix H, the first n-1 colections of H-MI are linearly independent

- So the n-th cohemn of H-MI must be a linear combination of the first n-1, since H-MI is singular

H-M= UR

- the first n-1 columns of H-MI

are linear combinations of the first

N-1 columns of U, 50 is the n-th

column

- ran is zeo