

## "Chrome V8 Source Code" 27. The mysterious and simple dispatch\_table\_

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### 1 abstract

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This article is the third of the Builtin topic, explaining the execution of Bytecode, data structure and Dispatch. `dispatch_table_` is the connection The link between Bytecodes, which records the address of each Bytecode handler, Ignition searches and executes the corresponding Bytecode. The content of this article is organized as follows: Bytecode execution and data structure (Chapter 2); Bytecode scheduling (Chapter 3).

### 2 Bytecode execution

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In V8, the interpreter responsible for executing Bytecode is Ignition. When Ignition executes Bytecode, it needs to make a lot of complicated preparations. These "preparations" Work" follow-up article explains, we focus on the execution of Bytecode.

Bytecode is generated at the granularity of JavaScript functions and stored in Bytecode array, that is, Bytecode array is an array that stores Bytecode.

The source code is as follows:

```
1. class BytecodeArray : public FixedArrayBase {
2. public:
3.     static constexpr int SizeFor(int length) {
4.         return OBJECT_POINTER_ALIGN(kHeaderSize + length);
5.     }
6.     inline byte get(int index) const;
7.     inline void set(int index, byte value);
8.     inline Address GetFirstBytecodeAddress();
9. //Omit.....
10. };
```

We only explain two points relevant to this article:

(1) Line 3 of the code `SizeFor(int length)` calculates the length of the Bytecode array. The value of the parameter `length` is after compiling the JavaScript function. The number of Bytecodes obtained, the space required by `length+Bytecode` array is equal to the length of Bytecode array. When creating Bytecode array When, use `SizeFor(int length)` to calculate the length of the requested memory;

(2) Line 8 of code `GetFirstBytecodeAddress()` obtains the first address of Bytecode. Copy the Bytecode generated by Parser to This function is used when Bytecode array.

In `Factory::NewBytecodeArray()`, use the return value of `SizeFor(int length)` to apply for memory, and use `CopyBytes()` to

Bytecode is copied to the first address. The following is a piece of Bytecode source code:

```
1. a7          StackCheck
2. 12 00       LdaConstant [0]
3. 15 01 00    StaGlobal [1], [0]
4. 13 01 02    LdaGlobal [1], [2]
5. 26 f9       Star r2
6. 29 f9 02    LdaNamedPropertyNoFeedback r2, [2]
7. 26 fa       Star r1
8. 0c 05       LdaSmi [5]
9. Constant pool (size = 6)
10. 0000005EAE403019: [FixedArray] in OldSpace
11.   - map: 0x03d0be000169 <Map>
12.   - length: 6
13.           0: 0x005eae402f59 <String[#22]: ignoreCase here we go>
14.           1: 0x038ee90c3cc1 <String[#1]: a>
15.           2: 0x01b92bc2bde1 <String[#9]: substring>
16.           3: 0x038ee90c3fa9 <String[#1]: b>
17.           4: 0x01b92bc33839 <String[#7]: console>
18.           5: 0x01b92bc32e79 <String[#3]: log>
```

Line 2 of code `12 00 LdaConstant [0]`: 12 is the number of `LdaConstant`, which is also the enumeration value of `LdaConstant`, that is `Bytecode[0x12]=kLdaConstant`, the source code is as follows:

```
enum class Bytecode : uint8_t {

    kWide, kExtraWide, kDebugBreakWide, kDebugBreakExtraWide, kDebugBreak0,
    kDebugBreak1, kDebugBreak2, kDebugBreak3, kDebugBreak4, kDebugBreak5,
    kDebugBreak6, kLdaZero, kLdaSmi, kLdaUndefined, kLdaNull, kLdaTheHole, kLdaTrue,
    kLdaFalse, kLdaConstant, kLdaGlobal, kLdaGlobalInsideTypeof, kStaGlobal,
    kPushContext, kPopContext, kLdaContextSlot, kLdaImmutableContextSlot,
    kLdaCurrentContextSlot, kLdaImmutableCurrentContextSlot, kStaContextSlot,
    kStaCurrentContextSlot, kLdaLookupSlot, kLdaLookupContextSlot,
    kLdaLookupGlobalSlot, kLdaLookupSlotInsideTypeof,
    kLdaLookupContextSlotInsideTypeof, kLdaLookupGlobalSlotInsideTypeof,
    kStaLookupSlot, kLdar, kStar, kMov, kLdaNamedProperty,
    kLdaNamedPropertyNoFeedback, kLdaKeyedProperty, kLdaModuleVariable,
    kStaModuleVariable, kStaNamedProperty, kStaNamedPropertyNoFeedback,
    kStaNamedOwnProperty, kStaKeyedProperty, kStaInArrayLiteral,
    kStaDataPropertyInLiteral, kCollectTypeProfile, kAdd, kSub, kMul, kDiv, kMod,
    kExp, kBitwiseOr, kBitwiseXor, kBitwiseAnd, kShiftLeft, kShiftRight,
```

```
kShiftRightLogical, kAddSmi, kSubSmi, kMulSmi, kDivSmi, kModSmi, kExpSmi,  
kBitwiseOrSmi, kBitwiseXorSmi, kBitwiseAndSmi, kShiftLeftSmi, kShiftRightSmi,  
kShiftRightLogicalSmi, kInc, kDec, kNegate, kBitwiseNot, kToBooleanLogicalNot,  
kLogicalNot, kTypeOf, kDeletePropertyStrict, //Omit.....  
}
```

V8 stipulates: fb represents register R0, fa represents register R1, and so on. At 29 f9 02 LdaNamedPropertyNoFeedback r2, In [2], f9 represents register R2, and 02 represents the constant pool [2]. When executing LdaNamedPropertyNoFeedback, Ignition obtains Get the base address of dispatch\_table, and then get the LdaNamedPropertyNoFeedback through base address+0x29

handler, the source code is as follows:

```
// Calls the GetProperty builtin for <object> and the key in the accumulator.  
IGNITION_HANDLER(LdaNamedPropertyNoFeedback, InterpreterAssembler) {  
    TNode<Object> object = LoadRegisterAtOperandIndex(0);  
    TNode<Name> name = CAST(LoadConstantPoolEntryAtOperandIndex(1));  
    TNode<Context> context = GetContext();  
    TNode<Object> result =  
        CallBuiltin(Builtins::kGetProperty, context, object, name);  
    SetAccumulator(result);  
    Dispatch();  
}
```

## 3 Dispatch

Dispatch\_table is an array of pointers. The enumeration value of Bytecode represents its position in the array, and the corresponding Bytecode is stored in this position.

The address of the handler. The initialization of Dispatch\_table is as follows:

```
1. void Interpreter::Initialize() {  
2.     Builtins* builtins = isolate_>builtins();  
3.     // Set the interpreter entry trampoline entry point now that builtins are  
4.     // initialized.  
5.     Handle<Code> code = BUILTIN_CODE(isolate_, InterpreterEntryTrampoline);  
6.     DCHECK(builtins->is_initialized());  
7.     DCHECK(code->is_off_heap_trampoline() ||  
8.             isolate_>heap()->IsImmovable(*code));  
9.     interpreter_entry_trampoline_instruction_start_ = code->InstructionStart();  
10.    // Initialize the dispatch table.  
11.    Code illegal = builtins->builtin(Builtins::kIllegalHandler);  
12.    int builtin_id = Builtins::kFirstBytecodeHandler;  
13.    ForEachBytecode([=, &builtin_id](Bytecode bytecode,  
14.                                     OperandScale operand_scale) {  
15.        Code handler = illegal;  
16.        if (Bytecodes::BytecodeHasHandler(bytecode, operand_scale)) {  
17. #ifdef DEBUG  
18.         std::string builtin_name(Builtins::name(builtin_id));  
19.         std::string expected_name =  
20.             Bytecodes::ToString(bytecode, operand_scale, "") + "Handler";
```

```
twenty one:         DCHECK_EQ(expected_name, builtin_name);
22. #endif

twenty three:         handler = builtins->builtin(builtin_id++);
twenty four:         }
25.         SetBytecodeHandler(bytecode, operand_scale, handler);
26.     });
27.     DCHECK(builtin_id == Builtins::builtin_count);
28.     DCHECK(IsDispatchTableInitialized());
29. }
```

The above 13-26 lines of code are anonymous functions, of which 25 lines of code initialize Dispatch\_table. The source code is as follows:

```
1. void Interpreter::SetBytecodeHandler(Bytecode bytecode,
2.                                     OperandScale operand_scale, Code handler)
{
3.     DCHECK(handler.kind() == Code::BYTECODE_HANDLER);
4.     size_t index = GetDispatchTableIndex(bytecode, operand_scale);
5.     dispatch_table_[index] = handler.InstructionStart();
6. }
7. //.....Separator line.....
8. size_t Interpreter::GetDispatchTableIndex(Bytecode bytecode,
9.                                           OperandScale operand_scale) {
10.     static const size_t kEntriesPerOperandScale = 1u << kBitsPerByte;
11.     size_t index = static_cast<size_t>(bytecode);
12.     return index + BytecodeOperands::OperandScaleAsIndex(operand_scale) *
13.               kEntriesPerOperandScale;
14. }
```

The above 5th line of code `dispatch_table_` is the member variable that stores the dispatch table that we have been thinking about for a long time; the 4th line of code `GetDispatchTableIndex()` calculates the position of the Bytecode handler in `dispatch_table`. This position is consistent with the enum class `Bytecode` is the same. Figure 1 shows the call stack of `SetBytecodeHandler`.

# Figure 1

```

87         return builtins->builtin(builtin_index);
88     }
89
90     void Interpreter::SetBytecodeHandler(BytecodeHandler* handler) {
91         // Set the handler for the current operation.
92         DCHECK(handler.kind() == Code::BYTECODE_HANDLER);
93         size_t index = GetDispatchTableIndex(handler);
94         dispatch_table_[index] = handler;
95     }
96
97     // static
98     size_t Interpreter::GetDispatchTableIndex(BytecodeHandler* handler) {
99         return builtin_index_map_[handler->kind()];
100     }

```



The `dispatch_table_` in the 11th line of code above is a member variable of `Interpreter`. `Interpreter` is a member variable of `Isolate`, the source code is as follows

Down:

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```
6.     Zone* compiler_zone_ = nullptr;
7.     CompilerDispatcher* compiler_dispatcher_ = nullptr;
8.     friend class heap::HeapTester;
9.     friend class TestSerializer;
10.    DISALLOW_COPY_AND_ASSIGN(Isolate);
11. };
```

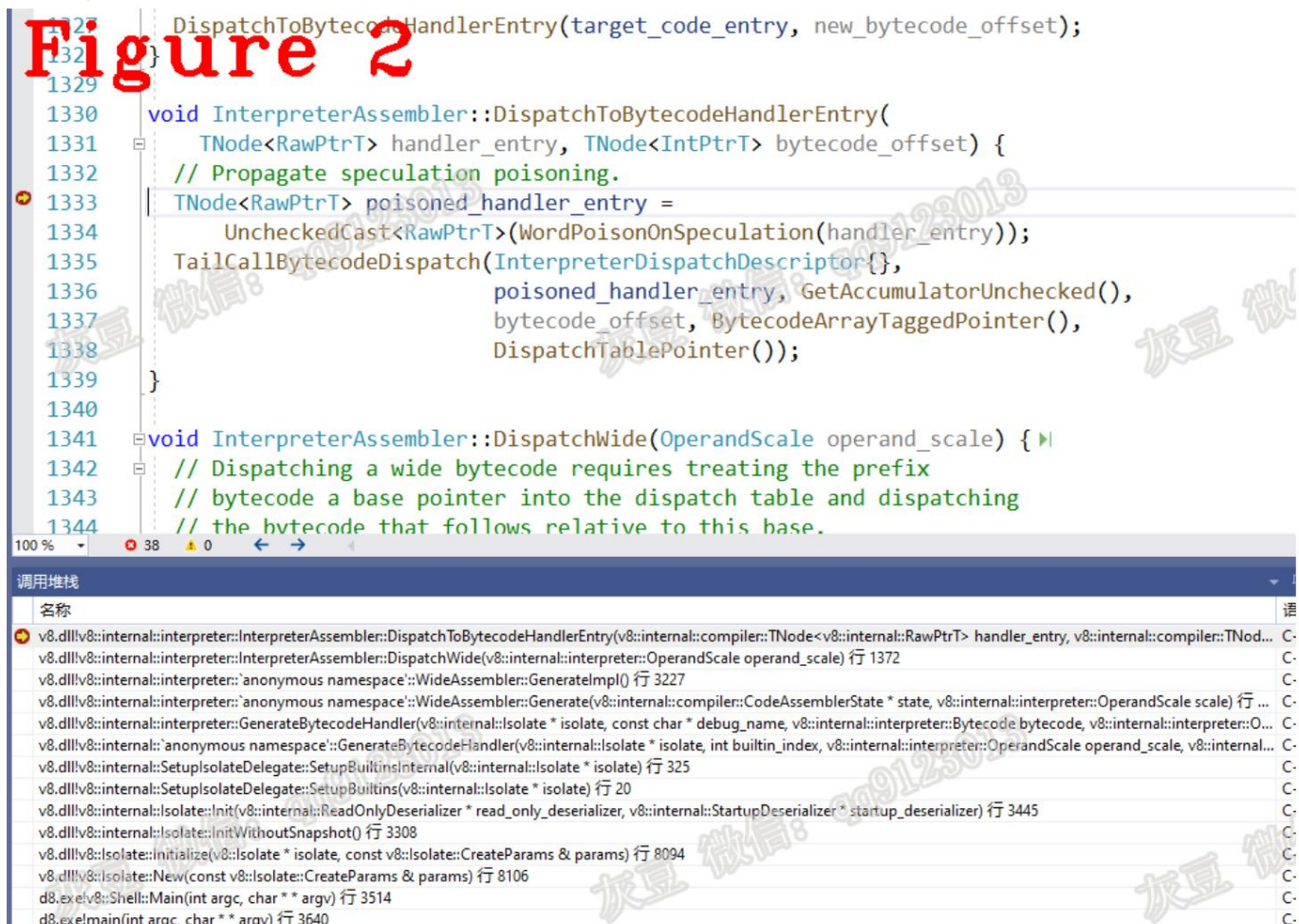
It can be seen from the above code: `Isolate->interpreter_>dispatch_table_gets dispatch_table_`. Below is the Source code of `Dispatch()` called in Bytecode handler:

```
1. void InterpreterAssembler::Dispatch() {
2.     Comment("===== Dispatch");
3.     DCHECK_IMPLIES(Bytecodes::MakesCallAlongCriticalPath(bytecode_),
made_call_);
4.     TNode<IntPtrT> target_offset = Advance();
5.     TNode<WordT> target_bytecode = LoadBytecode(target_offset);
6.     if (Bytecodes::IsStarLookahead(bytecode_, operand_scale_) {
7.         target_bytecode = StarDispatchLookahead(target_bytecode);
8.     }
9.     DispatchToBytecode(target_bytecode, BytecodeOffset());
10. }
11. //.....Divider line.....
12. void InterpreterAssembler::DispatchToBytecode(
13.     TNode<WordT> target_bytecode, TNode<IntPtrT> new_bytecode_offset) {
14. if (FLAG_trace_ignition_dispatches) {
15.     TraceBytecodeDispatch(target_bytecode);
16. }
17.     TNode<RawPtrT> target_code_entry = Load<RawPtrT>(<
18.         DispatchTablePointer(), TimesSystemPointerSize(target_bytecode));
19.     DispatchToBytecodeHandlerEntry(target_code_entry, new_bytecode_offset);
20. }
21. //.....Separation line.....
22. void InterpreterAssembler::DispatchToBytecodeHandlerEntry(
twenty three.     TNode<RawPtrT> handler_entry, TNode<IntPtrT> bytecode_offset) {
twenty four.     // Propagate speculation poisoning.
25.     TNode<RawPtrT> poisoned_handler_entry =
26.         UncheckedCast<RawPtrT>(WordPoisonOnSpeculation(handler_entry));
27.     TailCallBytecodeDispatch(InterpreterDispatchDescriptor{,
28.         poisoned_handler_entry,
GetAccumulatorUnchecked(),
29.         bytecode_offset, BytecodeArrayTaggedPointer(),
30.         DispatchTablePointer());
31. }
32. //.....Separator line.....
33. void CodeAssembler::TailCallBytecodeDispatch(
34. const CallInterfaceDescriptor& descriptor, TNode<RawPtrT> target,
35.     TArgs... args) {
36.     DCHECK_EQ(descriptor.GetParameterCount(), sizeof...(args));
37.     auto call_descriptor = Linkage::GetBytecodeDispatchCallDescriptor(
38.         zone(), descriptor, descriptor.GetStackParameterCount());
39.     Node* nodes[] = {target, args...};
```



```
40.     CHECK_EQ(descriptor.GetParameterCount() + 1, arraysize(nodes)); raw_assembler()-  
41.     >TailCallN(call_descriptor, arraysize(nodes), nodes);  
42. }
```

The above three methods jointly implement Bytecode dispatch. The 5th line of code calculates `target_bytecode`; the 17th line of code calculates `target_bytecode_entry`; the 27th line of code starts to jump; the 34th line of code creates a call descriptor; the 41st line of code generates a Node `node` and adds the node to the end of the current basic block, the jump is completed. For a detailed explanation of `TailCallN()`, see the eleventh article. Figure 2 shows the call stack of `Dispatch()`.



#### Technical summary

(1) The number of Bytecode is the subscript of the Bytecode handler in the array `dispatch_table_`; (2)

The initialization of `dispatch_table_` is completed when Isolate starts;

(3) The advantage of using fixed physical registers to save `dispatch_table_` is to avoid unnecessary Pushing and popping simplifies the design of Bytecode and improves the efficiency

of Dispatch; Tip: When I debug V8, `dispatch_table_` is always saved in the physical register R15. For debugging methods, see article 18.

Okay, that's it for today, see you next time.

Personal abilities are limited, there are shortcomings and

mistakes, criticisms and corrections are welcome WeChat: qq9123013 Note: v8 Communication Zhihu: <https://www.zhihu.com/people/v8blink>

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