

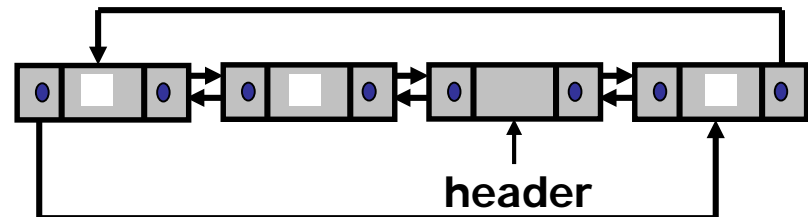


Trees

- Linear vs Nonlinear ADTs
- Tree Terminology
- Binary Tree
 - Binary Taxonomy Tree
 - Types of Binary Trees
- Binary Tree Representation
- Binary Tree Node Definition
- Binary Tree Scan Algorithms
 - Function Parameters
- Binary Search Tree

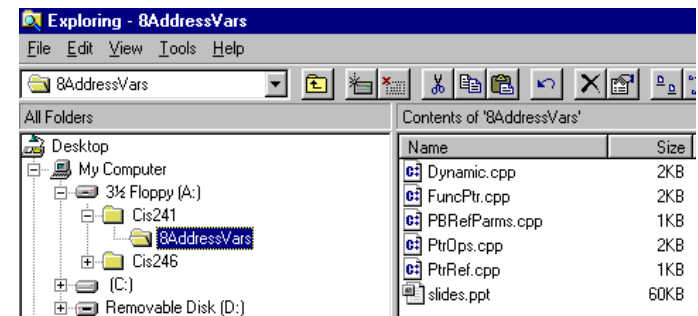
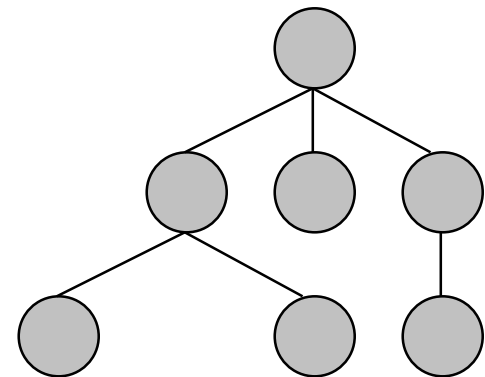
Linear ADTs

- Sequence containers store items in **positional** order
 - arrays, vectors, lists
- Linear ADTs have single successor



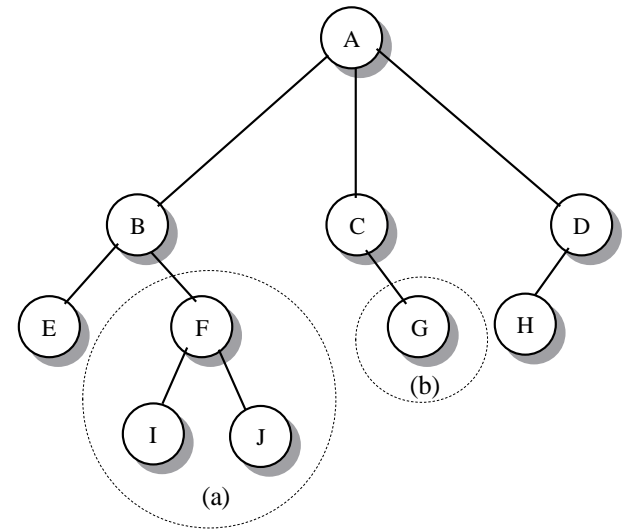
Nonlinear ADTs

- Container components have more complex relationship
 - Often allows improved efficiency for component access
- Tree is **hierarchical** structure
 - Places elements in nodes along branches that originate from root
- Tree is **nonlinear** structure
 - Each element can have multiple successors



Tree Terminology

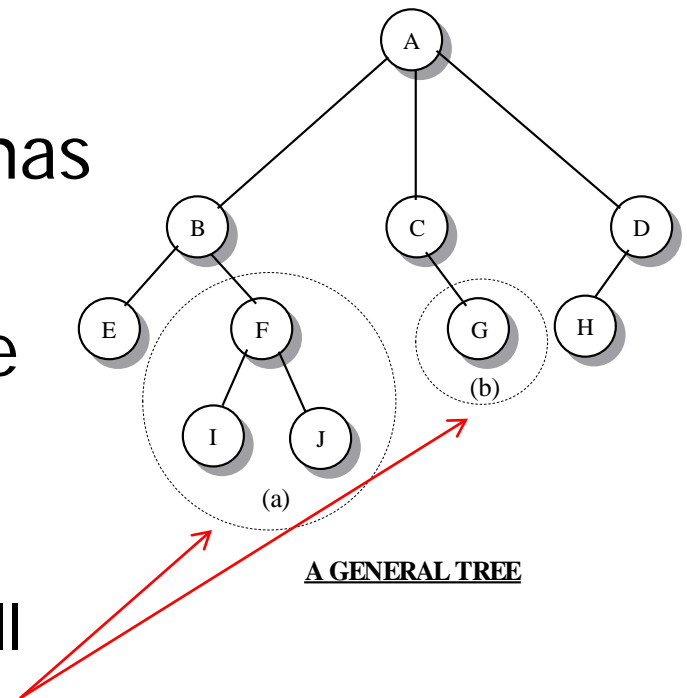
- unique component of tree is **node** that includes data and pointer references
 - tree with *no* nodes is **empty tree**
- unique node **ancestor** is called **parent**
- **root** node is the *unique* node of a tree with *no ancestor*
- node **descendants** are called **children**
- two **children** are **siblings** if they have the same **parent**
- connections from **parent** to **child** characterize the binary tree



A GENERAL TREE

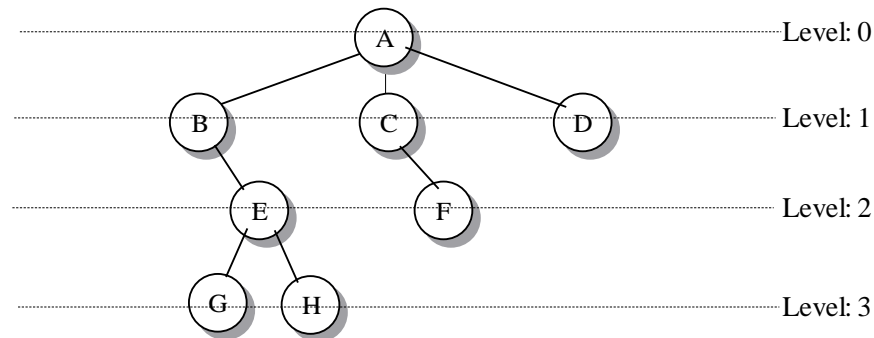
Tree Terminology

- **leaf** nodes are the nodes of a tree that have *no descendant*
- **interior** node (*nonleaf node*) has at least one child
- each node in a tree is also the **root** of a **subtree** (*recursive thinking*)
 - **subtree** describes a node and all its **descendants**



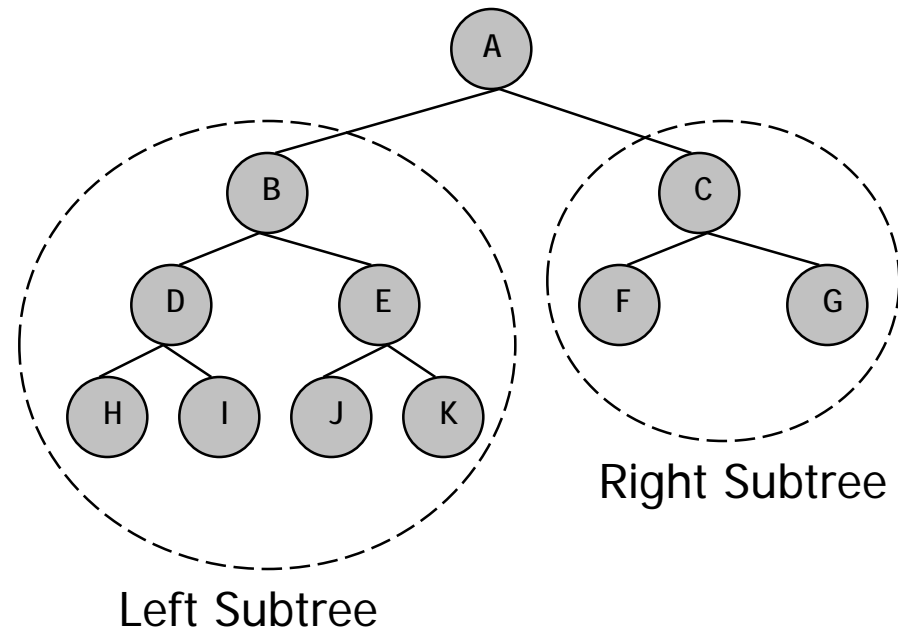
Tree Terminology

- **path** is linear subset of a tree (sequence of nodes)
 - there is a *unique* path from the root to any node
 - **level** or **depth** of a node is the length of its path from the root
- **depth** of tree is *maximum* depth of any node in the tree



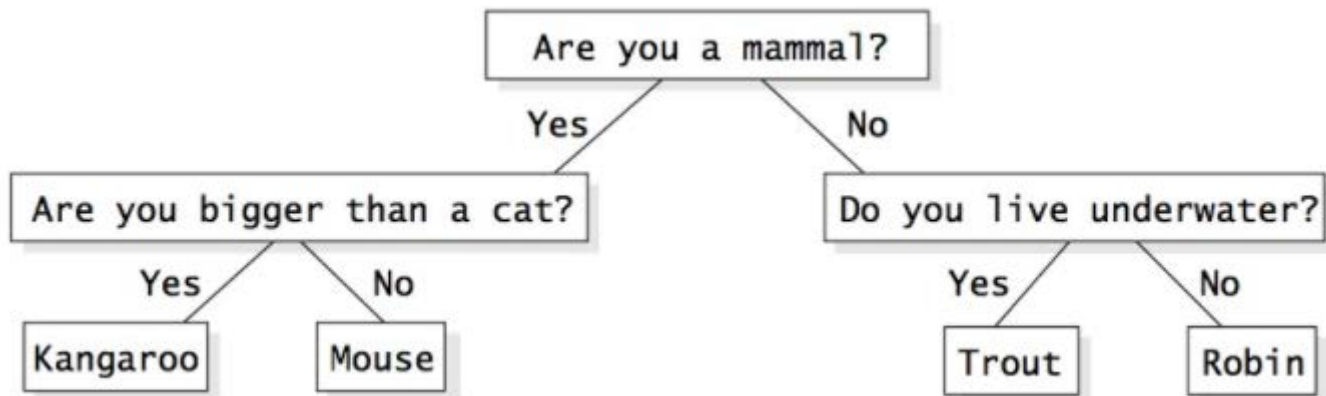
Binary Tree

- each parent has no more than two children (*descendent nodes*)
 - left child is node connected to left link
 - left subtree
 - right child is node connected to right link
 - right subtree
- each subtree is a binary tree
 - Recursive definition of binary tree
- a tree without any nodes is an empty binary tree



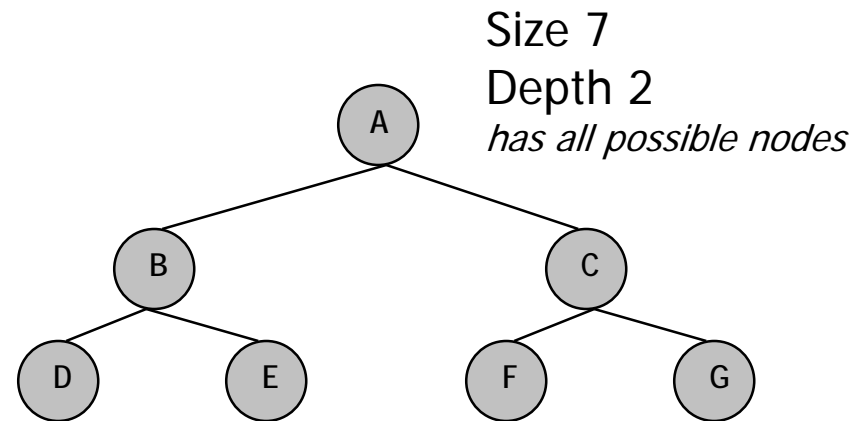
Binary Taxonomy Tree

- Stores increasing levels of detailed knowledge about a subject



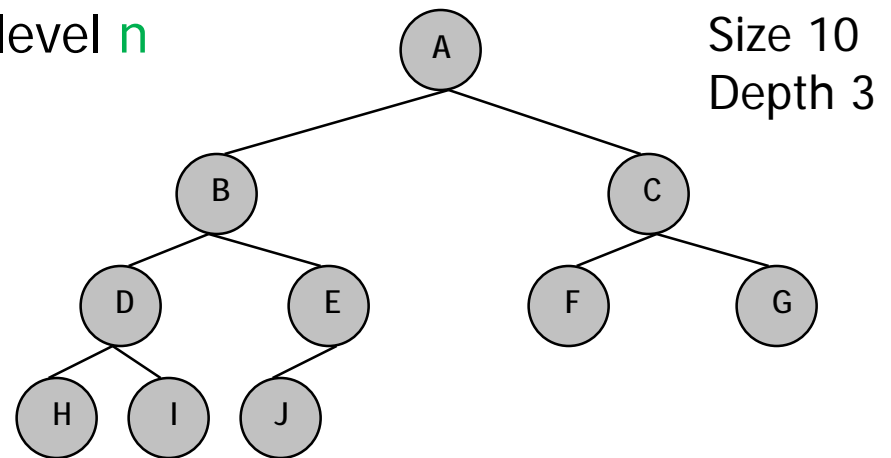
Types of Binary Trees

- full binary tree
 - every leaf node has same depth
 - every nonleaf node has two children
 - number of nodes in tree with n depth
 - $2^n + 2^{(n-1)} + \dots + 2^0$



Types of Binary Trees

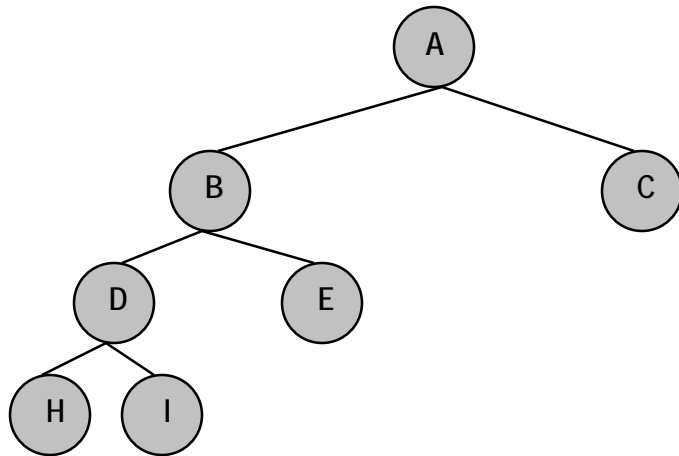
- **complete** binary tree
 - each level from 0 to $n - 1$ has all possible nodes (*full binary tree*)
 - all **leaf** nodes at depth n occupy **leftmost** position
 - number of nodes that occupy the first $n - 1$ levels
 - $2^n - 1$
 - possible number of nodes at level n
 - 1 to 2^n



Trees

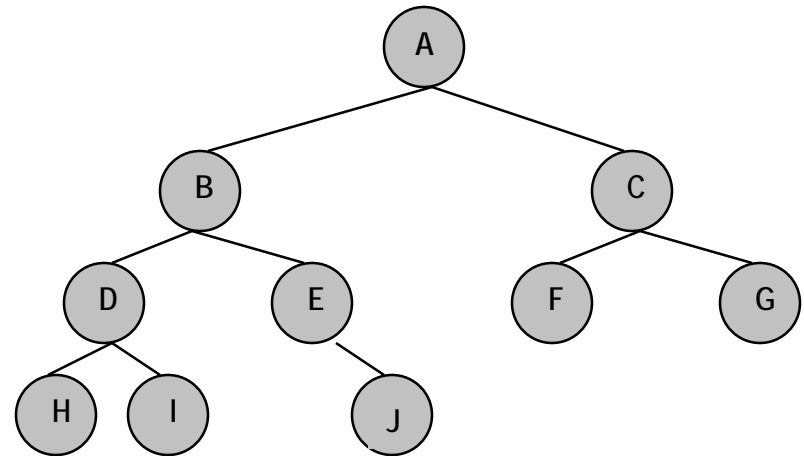
Types of Binary Trees

- non-complete binary trees



Size 7
Depth 3

Level 2 missing nodes

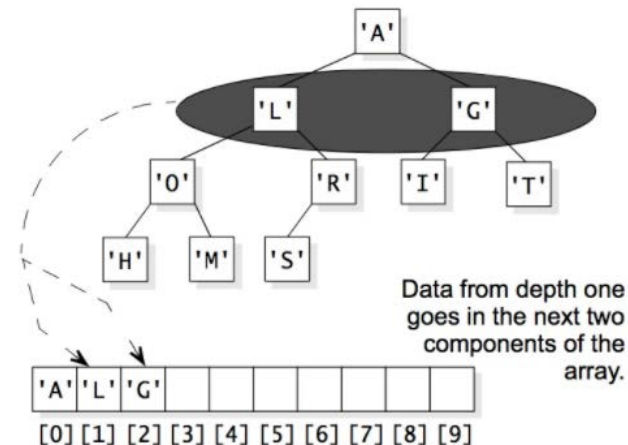
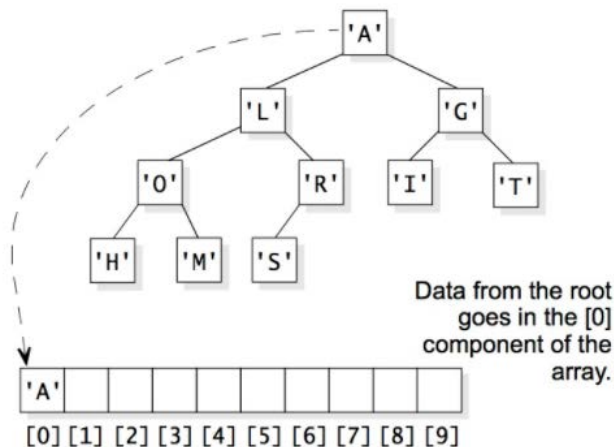


Size 10
Depth 3

Nodes at level 3 do not occupy leftmost position

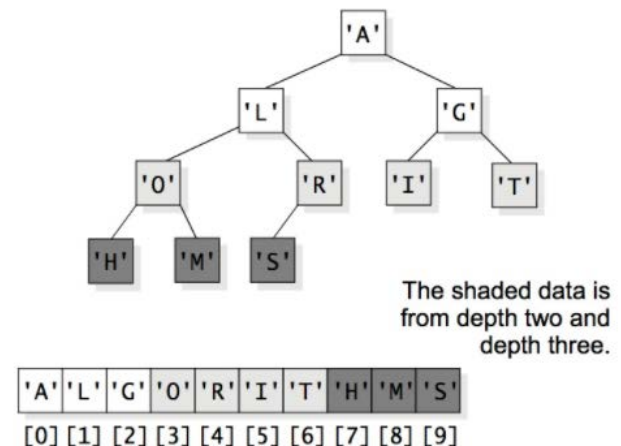
Binary Tree Representation

- **complete** binary tree can be represented using arrays



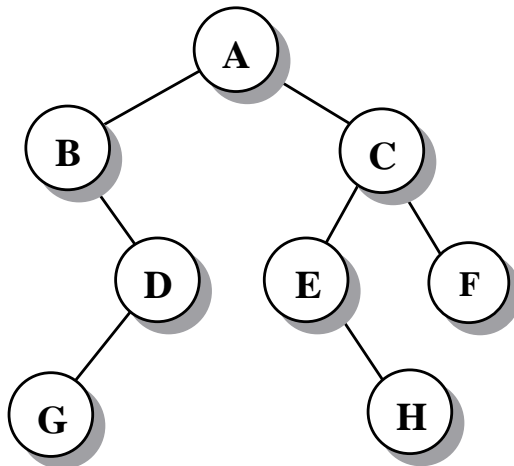
Binary Tree Representation

- **complete** binary tree can be represented using *dynamic* or *static arrays*
 - root in [0]
 - node data at [i]
 - parent at $[(i-1)/2]$
 - left child at $[2i+1]$
 - right child at $[2i+2]$

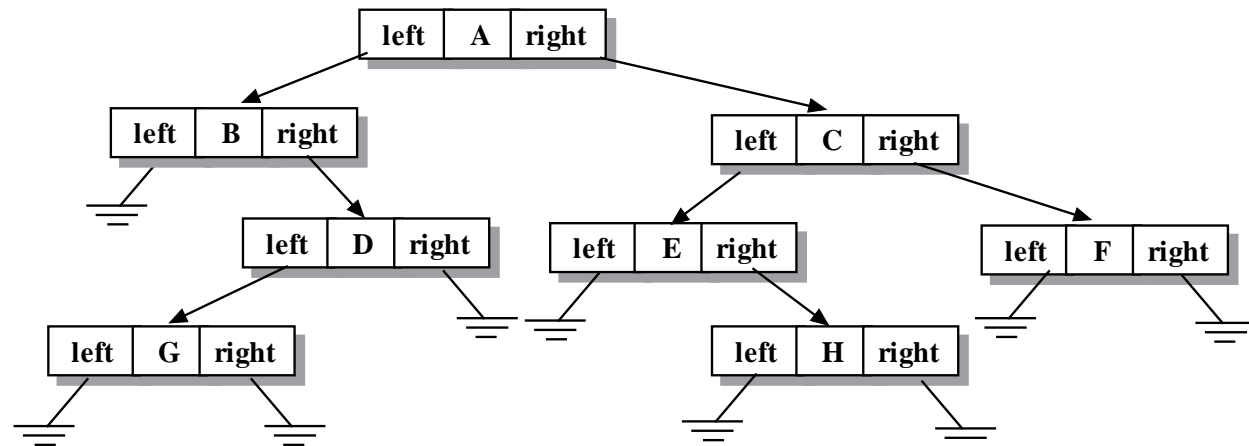


Binary Tree Representation

- *any* binary tree can be represented using connected **node** objects



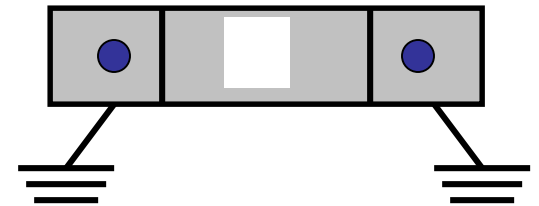
Abstract Tree Model



Tree Node Model

Binary Tree Node Definition

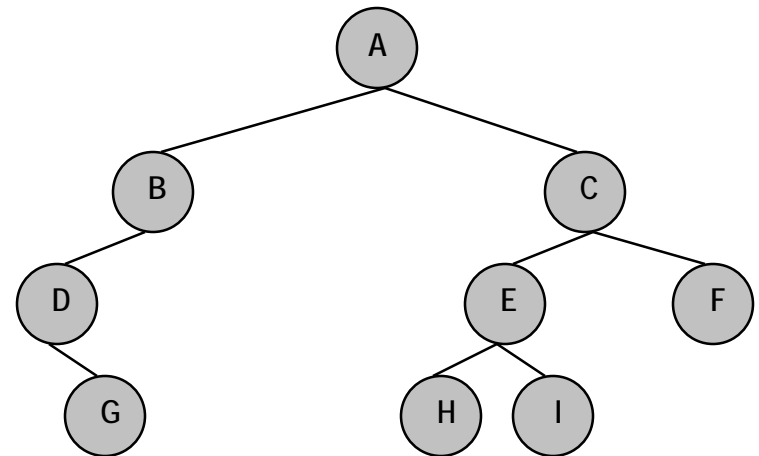
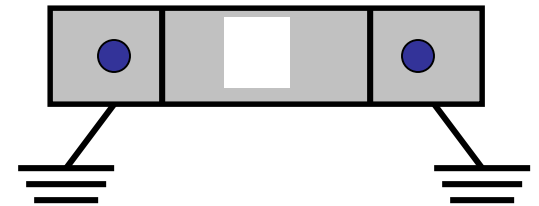
- Template `binary_tree_node` class
- Private data members include:
 - `data_field`
 - `left_field`
 - `right_field`
- Public methods include:
 - parameter constructor with defaults
 - accessor/mutator methods for private data
 - check for leaf status (*no children*)



Binary Tree Node Definition

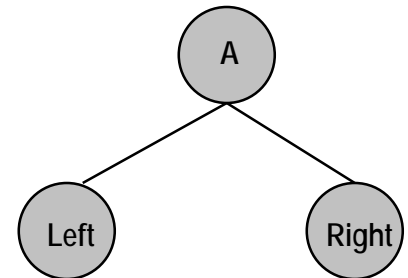
- Non-member toolkit functions:

- tree_clear
- tree_copy
- tree_size
- Tree Scanning
 - preorder
 - inorder
 - postorder
- print
- pretty_print



Binary Tree Scan Algorithms

- Since tree is **non-linear** structure, cannot scan nodes **sequentially**
- **Recursive** structure of binary tree allows traversal down subtrees
- Scanning strategies depend upon ordering of
 - Visit left child (*traverse left subtree*)
 - Perform action at node (*visit node*)
 - Visit right child (*traverse right subtree*)





Function Parameters

- Parameter to function can be function itself

```
ret_type parm_name(ptype_1, ptype_2, ...)
```

- Return type
- Parameter name
- Parentheses with parameter type

- Binary Tree Scanning Algorithms

- First parameter is void function with template parameter

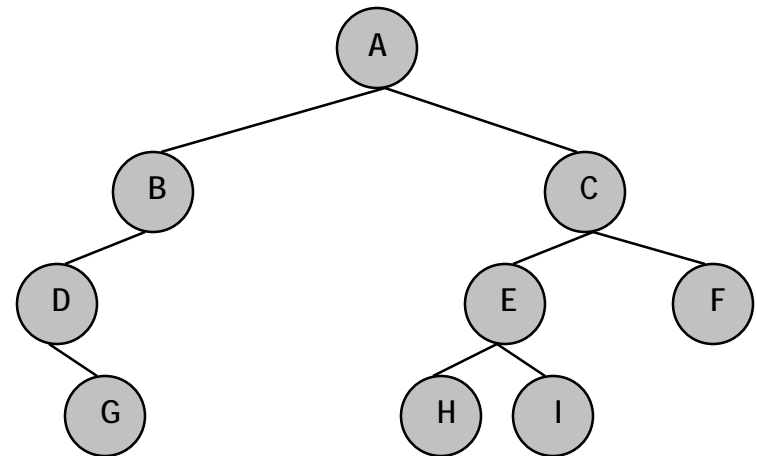
```
void func_name(Item)
```

```
void func_name(Item&)
```

Binary Tree Scan Algorithms

- **preorder** (NLR→recursive)
 - visit node
 - traverse left subtree
 - traverse right subtree

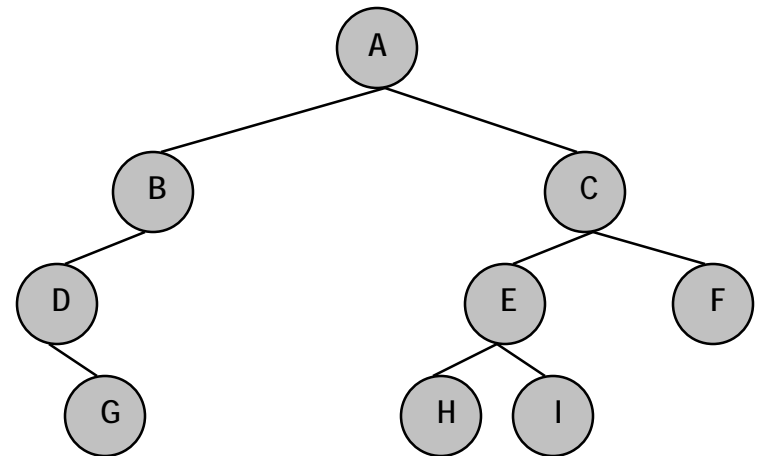
pre-order (NLR)
A B D G C E H I F



Binary Tree Scan Algorithms

- **inorder** (LNR→recursive)
 - traverse left subtree
 - visit node
 - traverse right subtree

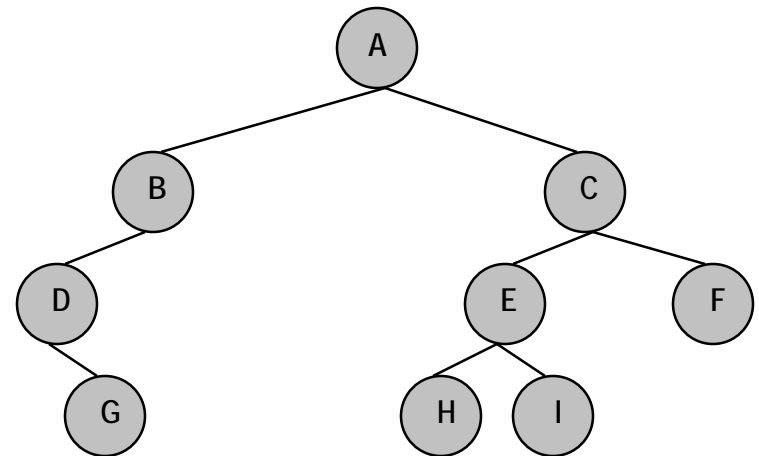
in-order (LNR)
D G B A H E I C F



Binary Tree Scan Algorithms

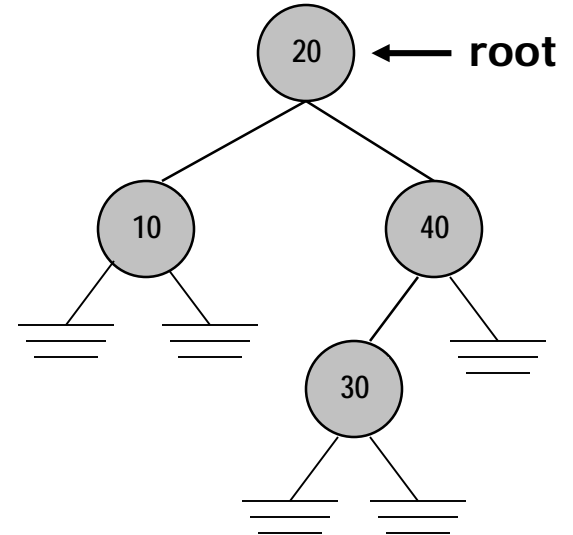
- **postorder** (LRN→recursive)
 - traverse left subtree
 - traverse right subtree
 - visit node

postorder (LRN)
G D B H I E F C A



Binary Search Tree

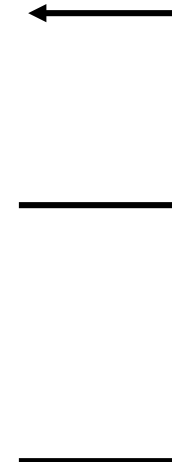
- Binary Tree whose **ordered** elements provide very efficient access and update operations
 - Elements ordered by relational operator $<$ for the template type
 - data values in **left** subtree are **less than** or **equal to** value of the node
 - data values in **right** subtree are **greater than** value of node





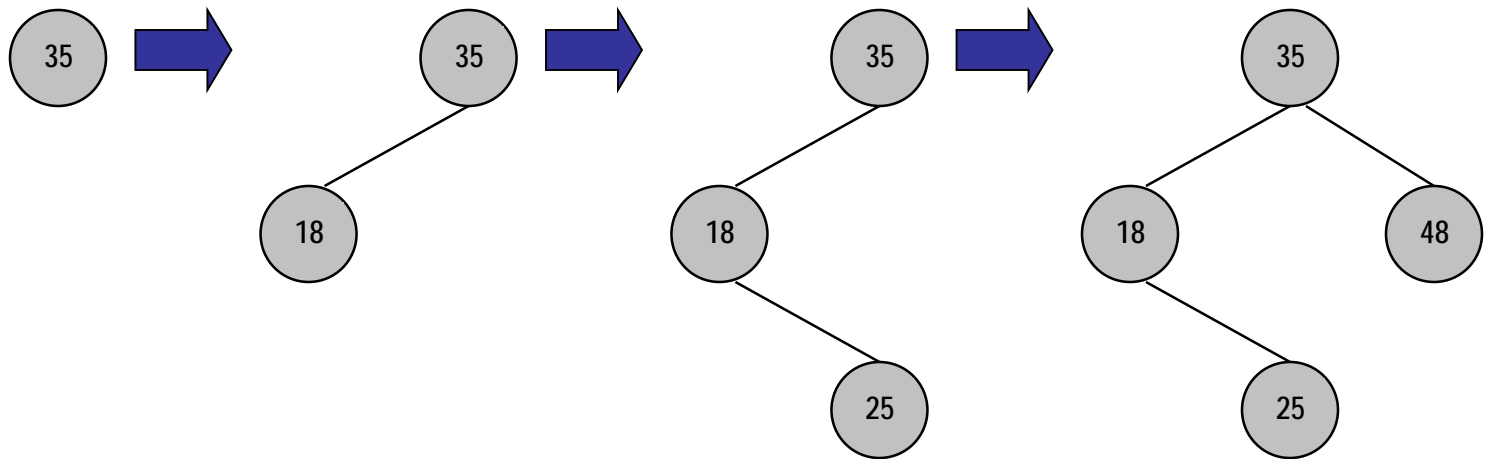
Building a BST

- elements enter BST via specific ordering strategy
- **first** element becomes **root**
- subsequent elements enter at end of path
 - if current \geq new
 - go left
 - left empty, add node
 - left not empty, go left and repeat
 - if current $<$ new
 - go right
 - right empty, add node
 - right not empty, go right and repeat



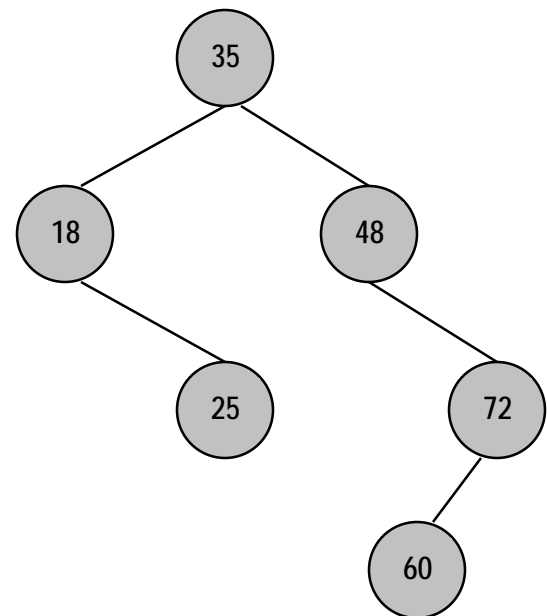
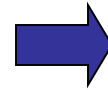
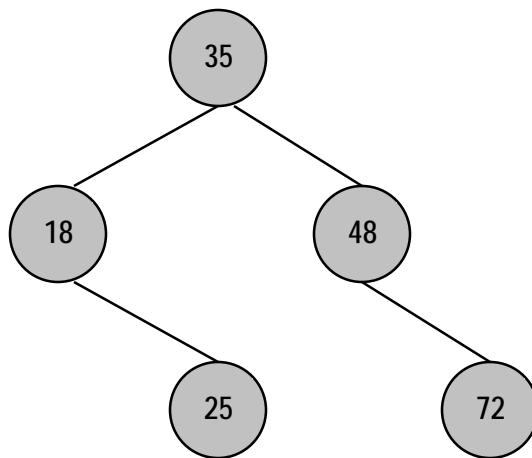
Building a BST

- Build binary search tree with elements
 - 35, 18, 25, 48, 72, 60



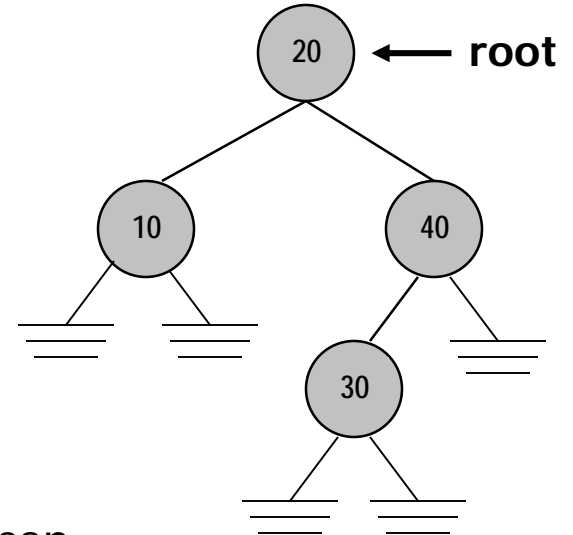
Building a BST

- Build binary search tree with elements
 - 35, 18, 25, 48, 72, 60



Binary Search Tree

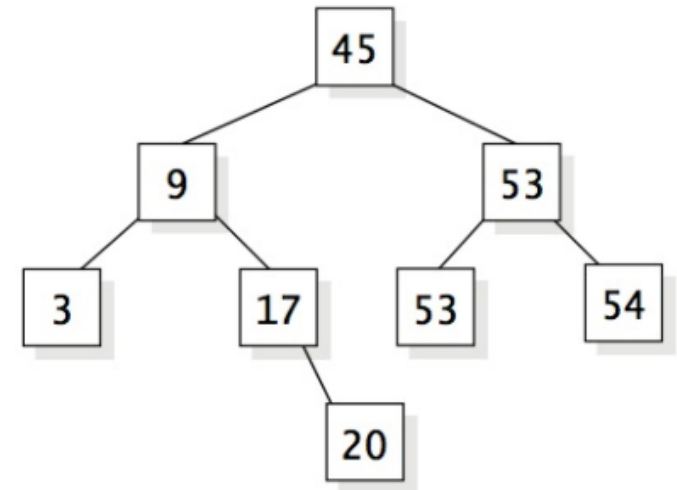
- Inorder scan (LNR) of BST ensures smaller nodes visited first
 - nodes visited in ascending order



Inorder Scan
10 20 30 40

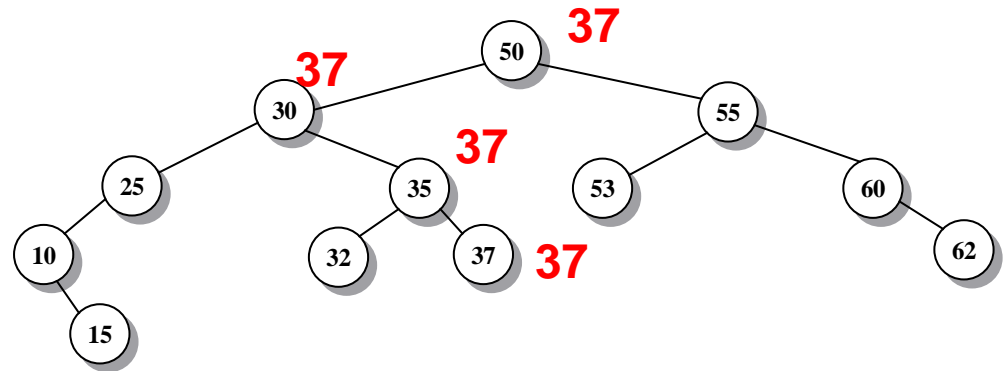
Binary Search Tree

- Improved search strategies for locating an element; path no greater than tree depth
 - scan left subtree if key is \leq
 - duplicates found in left subtree
 - otherwise scan right subtree



Locating Item in BST

- template type has operators `==` and `<` defined
- compare item with node (loop until match or all nodes visited)
 - if item `==` current
 - match found
 - else
 - if item `<` current
 - go left
 - if item `>` current
 - go right



Current Node
Root = 50

Node = 30

Node = 35

Node = 37

Action (looking for node 37)

Compare items 37 and 50

37 <= 50, move to the left subtree

Compare items 37 and 30

37 > 30, move to the right subtree

Compare items 37 and 35

37 > 35, move to the right subtree

Compare items 37 and 37. Item found.



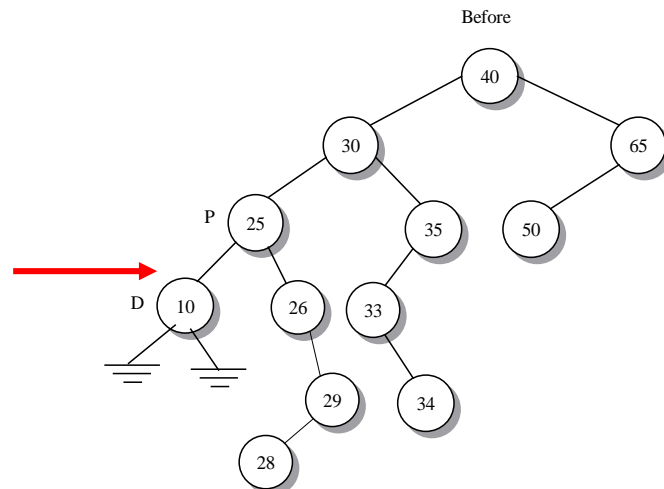
Removing Item from BST

- **must** preserve BST node relationships
 - `current == nullptr`
 - not found
 - `item == current`
 - `current_left == nullptr && current_right == nullptr`
 - `current_left == nullptr`
 - `new_root = current_right`
 - `current_left != nullptr`
 - `new_root =` largest entry in left subtree
 - remove current
 - `item < current`
 - `current = current_left`, repeat
 - `item > current`
 - `current = current_right`, repeat

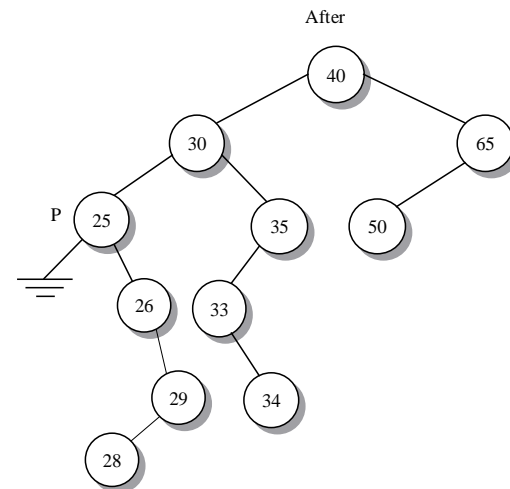


Removing Item From BST

- `item == current`
 - `current_left == nullptr && current_right == nullptr`
 - `remove current`



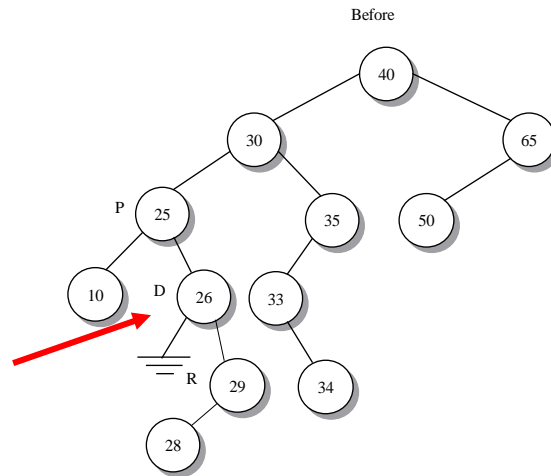
Delete leaf node 10.
pNodePtr->left is dNode



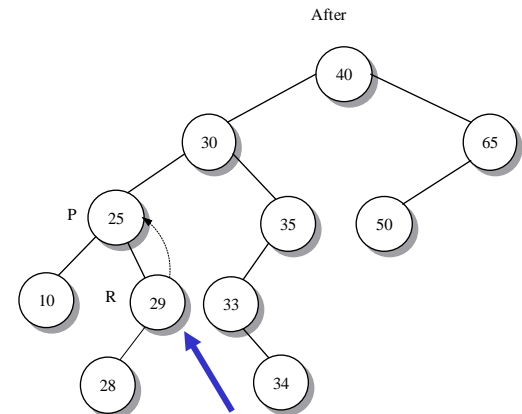
No replacement is necessary.
pNodePtr->left is NULL

Removing Item From BST

- `item == current`
 - `current_left == nullptr`
 - `new_root = current_right`
 - `remove current`



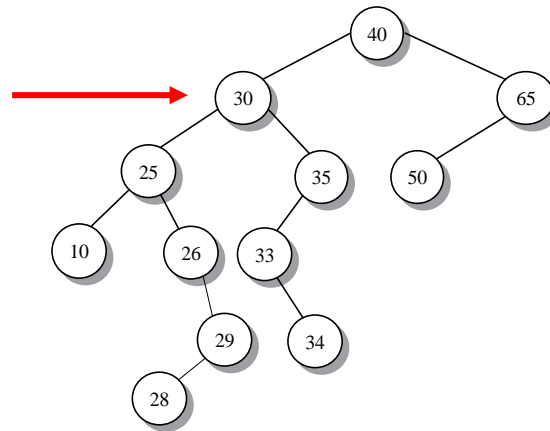
**Delete node 26 with only a right child:
Node R is the right child.**



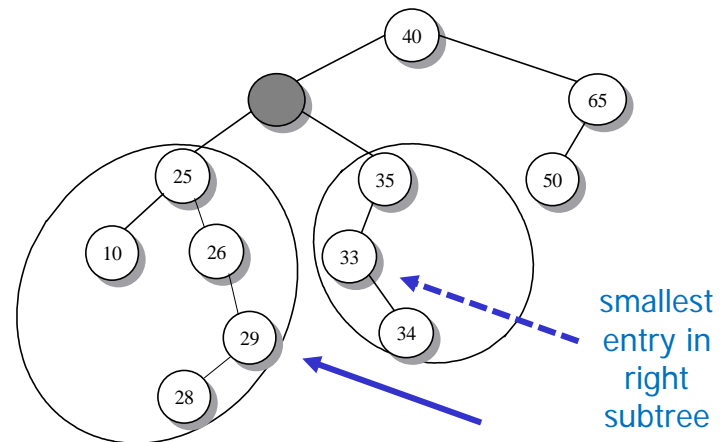
Attach node R to the parent.

Removing Item From BST

- `item == current`
 - `current_left != nullptr`
 - `new_root = largest entry in left subtree`
 - `remove current`



Delete node 30 with two children.



Orphaned subtrees.