

1. Determine whether the following arguments in propositional logic are valid:

(a) $p \vee (q \rightarrow r) \models p$

(b) $p \rightarrow (s \wedge q), r \rightarrow (q \vee \neg s), p \models \neg r$

2. Consider the following translation key:

p : Having a belief is having a disposition to act in some ways.

q : Non-living objects can be disposed to act in some ways.

r : Non-living objects can have beliefs.

s : Non-living objects represent the world being some way or another.

Now, translate the following arguments to propositional logic and determine whether they are valid:

(a) Neither can non-living objects represent the world as being some way or another, nor is having a belief having a disposition to act in some ways only if non-living objects can have beliefs. In conclusion, non-living objects cannot represent the world as being some way or another unless non-living objects can have beliefs.

Goal: to translate an argument that is hard to parse, even if it seems bad.

(b) If having a belief is having a disposition to act in some ways, and if non-living objects can be disposed to act in some ways, then non-living objects can have beliefs. But non-living objects can have beliefs only if they represent the world as being some way or another, and non-living objects do not represent the world as being some way or

Goal: to translate an argument that is easier to parse and seems good.

another. Therefore, non-living objects cannot be disposed to act in some ways.

3. Fill out the missing rules in the following proof:

1	$p \wedge q$	Assumption
2	r	Assumption
3	p	_____ 1
4	$p \wedge r$	_____ 2, 3

4. Fill out the missing lines in the following proof:

1	$p \rightarrow q$	Assumption
2	$p \rightarrow r$	Assumption
3	p	Assumption
4	q	$\rightarrow E$ _____
5	r	$\rightarrow E$ _____
6	$q \wedge r$	$\wedge I$ _____

5. **Extra:** Offer a proof of the following arguments:

(a) $(p \wedge q) \wedge (r \wedge s) \vdash r \wedge s$

(b) $(p \wedge r) \wedge (s \wedge q) \vdash r \wedge s$

(c) $p \rightarrow r, s \wedge p \vdash r \wedge s$

A *derivation* or *proof* is a finite sequence of formulas each of which is either an assumption or the outcome of an application of a rule of inference to prior formulas. The last formula in the sequence is the conclusion of the proof.

Here's a reminder of the rules needed for these derivations:

l	$\varphi \wedge \psi$	
n	φ	$E\wedge 1$
l	$\varphi \wedge \psi$	
m	ψ	$E\wedge 1$
l	ϕ	
m	ψ	
n	$\phi \wedge \psi$	$I\wedge 1, m$
l	$\phi \rightarrow \psi$	
m	ϕ	
n	ψ	$E\rightarrow 1, m$