XM3 Assembler User's Guide Larry Hughes, PhD 1 May 2020

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1 The Assembler

The X-Makina assembler is a two-pass assembler for the XM3 machine: the first pass builds the symbol table (from records in an XM3 assembly-language source module), while the second generates an executable load module (a file) for the XM3 machine. A list file is produced containing the records from the assembly-language file at the end of the first pass (if errors are detected during the first pass) or the end of the second pass (whether or not errors are detected during the second pass).

This document describes how to use the XM3 assembler and the internals of the assembler.

2 The assembly language file

An assembly language module is any text file (for example, created through Notepad or Wordpad) consisting of *records*. Each record has a specific format, described below.

2.1 Record format

An assembly-language file consists of one or more *records* containing instructions and data for the assembler to translate into machine-readable records. The record is *free format*, meaning

that it has no fixed fields. All records have the same format, defined as follows (**bold** indicates terminal symbols):¹

```
Record = (Label) + ([Instruction | Directive]) + (Operand) + (; Comment)

Label = Alphabetic + 0 {Alphanumeric} 30

Instruction = * An instruction mnemonic, see section 2.3 *

Directive = * An assembler directive, described in section 2.4 *

Operand = * The operand(s) associated with the Instruction or Directive, see section 2.5 *

Comment = * Text associated with the record – ignored by the assembler, see section 2.6 *

Alphabetic = [A..Z | a..z | _ ]

Alphanumeric = [A..Z | a..z | 0..9 | _ ]
```

Instructions and *Directives* are treated as case insensitive (that is, the instruction or directive can be upper case, lower case, or some combination thereof). However, a *Label*, if it exists, is case sensitive, meaning that, for example, the label **Alpha** is not the same as the label **ALPHA**.

2.2 Labels

A *Label* is a text-string of up to 32 characters. Each label must begin with an alphabetic character and be followed by zero or more alphabetic or numeric characters (i.e., alphanumeric). *Labels* are case sensitive.

Valid labels are stored in the symbol table with either the value of the location counter (i.e., where the instruction will be placed in memory, if there is an instruction on the record or if the remainder of the record is blank or contains a comment) or the value associated with the equate directive (see 2.4, Directives).

2.3 Instructions

The assembler supports XM3's Instruction Set Architecture (ISA), found in *Introduction to the XM3 Instruction Set Architecture*. A summary of the ISA can be found in section 7.

Instructions are case insensitive. The assembler produces the same executable code for the following instructions (LD - load register):

```
LD R2,R3; R3 = R2
Ld R2,R3; R3 = R2
1D R2,R3; R3 = R2
1d R2,R3; R3 = R2
```

Regardless of case, *Instructions* are **reserved words** and cannot be used as *Labels* or *Operands*.

2.4 Directives

A directive (or *pseudo-instruction*) is a command in a record that is processed by the assembler (i.e., it directs the assembler to do something). It has no corresponding machine instruction,

¹ The example is shown in a Data Dictionary format that can be used to define data structures: '(' ')' – optional; '[' '|' ']' – choice; 'LB{' '}UB' – sequence from LB (default 1) to UB (default ∞); and '+' – AND or concatenate.

although it can produce machine code. The directives currently supported by XM3 are (*Operand* is described in section 2.5, Operands):

ALIGN

Increments the location counter to the next even-byte address if the location counter is odd. If the address is odd, the executable file will contain a zero-value in the .XME file. **ALIGN** does not take an operand.

BSS Operand

The **BSS** (Block Started by Symbol) directive reserves a block of memory of *Operand* bytes is reserved. If there is a *Label* associated with the **BSS**, it is stored in the symbol table with the value of the location counter. The location counter is increased by the specified number of bytes. The label can be omitted.

BYTE Operand

One byte is stored in the memory location associated with the location counter. An *Operand* larger than 8-bits in length is an error. If there is a *Label*, it is stored in the symbol table along with the location counter. The location counter is increased by one.

END (Operand)

Denotes the end of the program. Any records that follow the END record are ignored. If an *Operand* is supplied, it must refer to a *Label* in the symbol table or an actual address; this is the starting address used by the loader.

EQU Operand

The **EQU** (equate) record's *Label* is equated with (i.e., assigned to) the *Operand*. The *Label* and the value of the *Operand* are stored in the symbol table. A *Label* is required. The location counter is not incremented.

ORG Operand

The **ORG** (origin) directive changes current location counter value to the address specified in *Operand*.

WORD Operand

Two bytes are stored in consecutive memory locations, starting at the location specified by the current value of the location counter. If there is a *Label*, it is stored in the symbol table along with the location counter. The location counter is increased by two bytes. Note that 16-bit quantities should fall on even-byte boundaries. The **ALIGN** directive can ensure proper alignment.

Directives are case insensitive. The assembler makes no distinction between the following directives:

BYTE	#FF
Byte	#FF
bvte	#FF

Regardless of case, Directives are reserved words and cannot be used as Labels or Operands.

2.5 Operands

An *Operand* contains up to three *Values* (separated by commas without spaces) required by the *Instruction* or *Directive*. It is defined as:

$$Operand = Value + 0{"," + Operand}3$$

A Value is either a numeric value or a label (section 2.2):

Numeric and *Label* values are distinguished using a prefix, with '\$', ''', and '#' denoting a numeric *Value* (*Alphanumeric* is defined above and terminal symbols or values are in **bold**):

The *Escaped Alphanumeric* value is limited to the following C-escape sequences ("Unknown" refers to an *Alphanumeric* character not in the list of supported characters):

Character	Converted valued	Meaning			
'\b'	#08	BS - Backspace			
'\t'	#09	TAB			
'\n'	#0a	Linefeed, Newline			
'\r'	#0d	Carriage return			
'\0'	#00	NUL			
'\\'	#5c	Backslash			
'\"	#27	Single quote			
'\'''	#22	Double quote			
'\Unknown'	#3f ('?')	Unknown character			

2.6 Comments

A comment is any text after a semicolon (";") to the end-of-record (delimited by a NUL character). Comments are ignored by the assembler.

2.7 Notes

- The location counter is incremented by the number of bytes associated with the **ALIGN**, **BSS**, **BYTE**, **ORG**, or **WORD** directive.
- Directives and instructions are reserved words and cannot be used as labels.

- Duplicate labels are not permitted.
- Characters begin and end with the single quote character "'.".
- Decimal values are prefixed with "\$".
- Hexadecimal numbers are prefixed with "#".
- Hexadecimal numbers cannot be signed.
- Unsigned values can be associated with the "+" sign (that is, -32768 to +65535)
- Any WORD value that exceeds its range is truncated to the least significant two bytes (four nibbles).
- Any BYTE value that exceeds its range is truncated to the least significant byte (two nibbles).

In addition, the assembler does not support:

- External references (i.e., references to *Labels* in other files)
- Include files (i.e., external files to be "copied into" the program for assembly)
- In-line arithmetic expressions (i.e., operands containing arithmetic operators)

3 Built-in symbols

There are eight built-in symbols, representing XM3's eight CPU registers (R0, R1, R2, R3, R4, R5, R6, and R7). Other symbols can be equated as registers, for example:

```
; LR is equated with R5
LR
     equ
          R5
     equ R6
                   ; r6 is equated with R6
r6
                    ; SP is equated with r6 (equated with R6)
SP
    equ r6
SP
    EQU
          R6
рс
    equ
          R7
                  ; Subroutine return
    mov LR,pc
    mov R5,R7
                   ; Equivalent to the previous record
```

4 Examples

The following example shows how a problem (assigning the ASCII characters 'A' through 'Z' to a block of memory starting at location #800) could be solved using the XM3 assembler.

4.1 First pass errors - .LIS file

The first attempt at solving the problem is as follows:

```
; Sample XM3 program
; Initialize a block of memory to 'A' through 'Z'
; ECED 3403
; 1 May 2020
SIZE equ
          $26
CAP A equ 'A'
CAP Z equ 'Z'
; Start of data area
  org #800
Base bss
          SIZE
                      ; Reserve SIZE bytes
; Start of code area
  org #1000
Start
 movlz CAP A, RO; RO = 'A'
 movlz
          CAP^{T}Z,R1 ; R1 = 'Z'
; R2 = Base (Base address to store characters)
  movl Base, R2 ; R2 = LSByte(Base) or \#00 movh Base, R2 ; R2 = MSByte(Base) or \#08
Loop
  st.b R0,R2+ cmp.b R0,R1
                       ; [R2] = R0; R2 = R2 + 1
                       ; R0 = R1 ('Z')
          NE,$3,$0
                       ; If not equal, do next 3 instruction, otherwise skip
  cez
           $1,R0
  add.b
                       ; No: R0 = R0 + 1 (next ASCII char)
  swap R0,R0
                       ; NOP to test count
                       ; Repeat loop
 bra Loop
; End of program
Done
          '*',R1
  movlz
  bra Done
  end Start
              ; End of program - first executable address is "Start"
```

The program is dragged-and-dropped onto the assembler, which stops at the end of the first pass, indicating that first-pass errors occurred. The corresponding .LIS file contains the following:

```
XM3 Assembler - Version 3.0 (10 April 2020)
Input file name: ArrayInit.asm
Time of assembly: Fri 1 May 2020 14:02:22
 1
  2
                  ; Sample XM3 program
  3
                  ; Initialize a block of memory to 'A' through 'Z'
  4
                  ; ECED 3403
  5
                  ; 1 May 2020
  6
  7
                  SIZE equ
                              $26
  8
                  CAP A equ
                              'A'
  9
                  CAP Z equ
                              'Z'
10
11
                  ; Start of data area
12
13
                        org
                              #800
14
                  Base bss
                              SIZE
                                          ; Reserve SIZE bytes
15
16
                  ; Start of code area
17
18
                              #1000
                        orq
19
 20
                                          ; R0 = 'A'
                  Start movlz CAP A, R0
                        movlz CAP Z,R1
 21
                                          ; R1 = 'Z'
 22
                  ; R2 = Base (Base address to store characters)
 23
                        movl Base, R2
                                          ; R2 = LSByte(Base) or #00
 24
                        movh Base, R2
                                          ; R2 = MSByte(Base) or #08
 25
 26
                  Loop
 27
                        st.b R0, R2+
                                          ; [R2] = R0; R2 = R2 + 1
 28
                        cmp.b R0,R1
                                          ; R0 = R1 ('Z')
                        cez NE,$3,$0
 29
                                          ; If not equal, do next 3 instruction, otherwise skip
              ** Expecting INST or DIR
 30
                        add.b $1,R0
                                           ; No: R0 = R0 + 1 (next ASCII char)
 31
                                           ; NOP to test count
                        swap R0, R0
 32
                        bra
                              Loop
                                          ; Repeat loop
 33
                  ; End of program
 34
                  Done movlz '*',R1
 35
                        bra
                              Done
 36
                  ;
 37
                                          ; End of program - first executable address is "Start"
                        end
                              Start
```

First pass errors - assembly terminated

** Symbol table **

** Symbol table **			
Name	Type	Value	Decimal
Done	LBL	100C	4108
cez	LBL	100C	4108
Loop	LBL	1008	4104
Start	LBL	1000	4096
Base	LBL	0800	2048
CAP_Z	LBL	005A	90
CAP_A	LBL	0041	65
SIZE	LBL	001A	26
R7	REG	0007	7
R6	REG	0006	6
R5	REG	0005	5
R4	REG	0004	4
R3	REG	0003	3
R2	REG	0002	2
R1	REG	0001	1
R0	REG	0000	0

The assembler indicates that an instruction or directive was expected on record 29 of the .ASM file ("**** Expecting INST or DIR"). The instruction "cex" is taken as a label by the assembler since it is not a valid instruction, meaning that the operand "Done" is causing the error.

4.2 Second pass - .LIS file

The invalid instruction can be corrected (changing "cez" to "cex"). By running the corrected assembler file through the assembler a second time, the assembler indicates that the assembly was successful.

Two new files are in the directory, one the list file (.LIS) and the other, the executable (.XME). The .LIS file contains the name of the input (.ASM) file (without its full path), a listing of the assembled program (from left-to-right: the record number, the machine address, the instruction or data to be stored in that location, and the original .ASM record), the symbol table (from left-to-right: the name or label, the type [LBL – label or REG – register], and its value), and the name of the .XME file (with its full path):

```
XM3 Assembler - Version 3.0 (10 April 2020)
Input file name: ArrayInit.asm
Time of assembly: Fri 1 May 2020 14:19:55
 1
  2
                  ; Sample XM3 program
  3
                  ; Initialize a block of memory to 'A' through 'Z'
  4
                 ; ECED 3403
  5
                  ; 1 May 2020
  6
  7
                 SIZE equ
                             $26
  8
                 CAP A equ
                              'A'
  9
                  CAP Z equ
                              'Z'
10
11
                  ; Start of data area
12
13
                        org
                              #800
14
                                         ; Reserve SIZE bytes
      0800 0000
                 Base bss
                             SIZE
15
16
                  ; Start of code area
17
18
                             #1000
                        orq
19
 20
     1000 6A08 Start movlz CAP A, RO
                                        ; R0 = 'A'
                       movlz CAP Z,R1
 21
     1002 6AD1
                                        ; R1 = 'Z'
 22
                 ; R2 = Base (Base address to store characters)
 23
     1004 6002
                       movl Base, R2
                                               ; R2 = LSByte(Base) or #00
 24
     1006 7842
                       movh Base, R2
                                               ; R2 = MSByte(Base) or #08
 25
 26
                 Loop
 27
     1008 5CC2
                       st.b R0,R2+
                                               ; [R2] = R0; R2 = R2 + 1
     100A 4541
 28
                       cmp.b R0,R1
                                          ; R0 = R1 ('Z')
 29
     100C 2458
                       cex NE, $3, $0
                                         ; If not equal, do next 3 instruction, otherwise skip
 30
     100E 40C8
                       add.b $1,R0
                                         ; No: R0 = R0 + 1 (next ASCII char)
 31
     1010 4C80
                       swap R0, R0
                                         ; NOP to test count
 32
     1012 23FA
                       bra
                             Loop
                                          ; Repeat loop
 33
                 ; End of program
     1014 6951 Done movlz '*',R1
 34
 35
     1016 23FE
                       bra
                             Done
 36
                  ;
 37
                                          ; End of program - first executable address is "Start"
                        end
                             Start
```

Successful completion of assembly

** Symbol table **			
Name	Type	Value	Decimal
Done	LBL	1014	4116
Loop	LBL	1008	4104
Start	LBL	1000	4096
Base	LBL	0800	2048
CAP_Z	LBL	005A	90
CAP_A	LBL	0041	65
SIZE	LBL	001A	26
R7	REG	0007	7
R6	REG	0006	6
R5	REG	0005	5
R4	REG	0004	4
R3	REG	0003	3
R2	REG	0002	2
R1	REG	0001	1
R0	REG	0000	0

.XME file: C:\Users\larry\OneDrive\Courses\ECED 3403 - 2020\XM3\XM3 - Test files\Completed XM3 tests\ArrayInit.xme

4.3 Second pass - .XME file

The .XME file is the executable file produced by the assembler from the .ASM file at the end of the successful second pass. The file consists of <u>S-records</u>:

```
S0100000ArrayInit.asmED
S104080000F3
S11B1000086AD16A02604278C25C41455824C840804CFA235169FE231F
S9031000EC
```

The assembler supports three types of S-record:

- **S0**: The header record containing the name of the .ASM file from which the executable was obtained. The file name is the name of the file found in the directory; the full path is omitted. In this example, the S0-record fields are as follows:
 - so: SO record indicator
 - 10: Length of record (address bytes [2], filename bytes [13], and checksum [1]: 16 bytes)
 - 0000: Address field
 - ArrayInit.asm: The name of the file
 - ED: The checksum
- **S1**: Data and instructions are stored in S1-records. In this example, there are two S1 records.

The first record's fields are:

- s1: S1 record indicator
- 04: Length of the record (address bytes, data/instruction bytes, and checksum: 4 bytes)
- 0800: Address field indicates location of first byte, as specified by the origin record (address #0800).
- 00: First data/instruction byte. Since this came from a BSS, the assembler produces only the first byte.
- F3: The checksum of the length byte, address bytes, and data/instruction bytes.

The second record's fields are:

- s1: S1 record indicator
- 1B: Length of the record (address bytes, data/instruction bytes, and checksum: 27 bytes)
- 1000: Address field indicates location of first byte, as specified by the origin record (address #1000)
- 086AD16A02604278C25C41455824C840804CFA235169FE23: The data/instruction bytes produced by the assembler from the original .ASM file. In this example, the first byte, #08, is stored in location #1000, the second, #6A, in location #1001, the third, #D1, in location #1002, and so on. The byte ordering used is *little-endian* (least-significant byte is stored in the first byte address and the most-significant byte is stored in the second byte address). In the case of #086A, #08 is stored first, then #6A. This pattern can be seen by comparing the output of the .LIS file with that of the .XME file.
- 1F: The checksum of the length byte, address bytes, and data/instruction bytes.

S9: The starting address of the program, used when the program is loaded into XM3's memory. If an address is specified as part of the **END** record in the .ASM file, it is used here. If the END record is omitted or there no starting address specified, the assembler defaults to zero. The data/instruction field is omitted. In this example, the program included a starting address, #1000:

s9: S9 record indicator

03: Length of the record (3 bytes)

1000: Starting address (#1000)

EC: The checksum of the length byte and address bytes.

5 Internals of the assembler

The XM3 assembler is two-pass: The first pass checks each non-commented record for validity, stores the label in the symbol table (if present), and increments the location counter, while the second pass rereads the file, generating the corresponding machine code for each non-comment record.

The assembler has a *location counter* that indicates where the next instruction or data value is to be loaded (i.e., resides) in memory. The location counter is incremented by 2 for all instructions and the number of bytes associated with the directive (if **BSS**, **BYTE**, or **WORD**). The location counter is incremented by 0 or 1 (**ALIGN**, depending on the current value of the location counter), it is incremented by the size of a reserved block of memory (**BSS**), and is assigned an entirely new value (if **ORG**).

5.1 First pass

The first pass repeats the following until end-of-file or an END directive is found:

- 1. Read the next record.
- 2. Comment records are ignored and the location counter is not to be incremented. *Labels* are stored in the symbol table along with the value of the location counter (i.e., the *Label's* address). A duplicate label is an error.
- 3. The next field must be an *Instruction*, *Directive*, *Comment*, or nothing. Anything else is an error.
- 4. If there is an *Instruction* or *Directive* and it requires an *Operand* field, the operand must exist and be correct. An invalid or unexpected *Operand* will cause an error diagnostic to be issued. It is possible that the *Operand* is undefined until the second pass (if it is a *Label* that is a forward reference). This does not affect the location counter; it is simply incremented by the number of bytes required by the instruction or directive.
- 5. *Comments*, if they exist, are ignored.

If an error is found in a record, the subsequent records are processed until the end-of-file or **END**. The file and any errors are written to the .LIS file.

5.2 Second pass

The second pass is performed if one or more errors are detected on the first pass. If not errors are found, the .LIS file is rewound for output of the listing file and the .XME file is opened.

The second pass repeats the following until end-of-file or an **END** directive is found:

- 1. Read the next record.
- 2. If the record is a comment, it should be ignored and step 1 repeated.
- 3. Ignore the Label if there is one (it was handled during the first pass).
- 4. A record containing an *Instruction* has the instruction and operand(s) extracted. The *Instruction*'s corresponding opcode is found and the *Operand*(s) are determined from supplied *Value* or *Label*. The opcode and any values are combined according to the rules associated with the Instruction's format to create the machine instruction. The machine instruction is then emitted, along with the current value of the location counter. The location counter is incremented by 2.
- 5. If the record contains a *Directive*, directive-specific steps are performed; for example, ORG changes the location counter and **BYTE** or **WORD** writes the value to the .XME file.
 - If the Directive is **END**, the assembler stops reading the file. If errors were detected, the file is removed.

5.3 Notes

• Error messages are generated for missing *Operands* (the number depends upon the *Instruction* or *Directive*) or a supplied *Operand* (if *Operands* are not required by the *Instruction* or *Directive*).

6 Known limitation

The assembler has the following known limitation as of 1 May 2020:

1. The MOVH instruction takes the most significant 8-bits of the 16-bit value. Using an 8-bit value will result in the assembler using a value of #00:

```
MOVH #FFEE,R0 ; R0 = #FF
MOVH #FF,R0 ; R0 = #00
```

If errors are found, please contact Dr. Hughes, supplying the .ASM program and any other supporting documentation to allow for the correction of the error.

7 XM3 Instruction Set

Bit value definitions for XM3 Instruction Set (last page)

0	1	Instruction opcode bit values								
PRPO		Pre- or post-increment or pre- or post-decrement (Load and Store).								
DEC		Decrement the register (before or after the instruction is executed).								
INC		Increment the register (before or after the instruction is executed).								
W/B		Word (16-bits) or byte (8-bits) addressing or register size.								
R/C		Register (0) or Constant (1).								
S		Source register bit (one of 3).								
D		Destination register bit (one of 3).								
В		Bit (one of 8) in MOVL, MOVLZ, MOVLS, and MOVH instructions.								
OFF		A bit used in an offset (in LDR, STR, and branching instructions).								
S/C		Source register or constant value (see below)								
SA		SVC (Service Call) vector address (#0 through #F).								
С		Conditional execution code (#0 to #E)								
Т		THEN (True) count (#0 to #7)								
F		ELSE (False) count (#0 to #7)								
V, SLP, I	N, Z, C	Condition code values (oVerflow, Sleep, Negative, Zero, and Carry).								

Register and Constant values for R/C and SRC bits

R	R/C									
0	1	Encoding								
Register	Constant	(bits 3-5)								
R0	0	000								
R1	1	001								
R2	2	010								
R3	4	011								
R4	8	100								
R5/LR	16	101								
R6/SP	32	110								
R7/PC	-1	111								

Conditional execution codes and descriptions for CEX instruction

Code	Description	True state	Code	Description	True state
0000	Equal / equals zero	Z	1000	Unsigned higher	C and !Z
0001	Not equal	!Z	1001	Unsigned lower or same	!C or Z
0010	Carry set / unsigned higher or same	С	1010	Signed greater than or equal	N == V
0011	Carry clear / unsigned lower	!C	1011	Signed less than	N != V
0100	Minus / negative	N	1100	Signed greater than	!Z and (N == V)
0101	Plus / positive or zero	!N	1101	Signed less than or equal	Z or (N != V)
0110	Overflow	V	1110	Always (default)	any
0111	No overflow	!V	1111	Invalid	

XM3 Instruction Set

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	Mnemonic	Instruction
0	0	0	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	BL	Branch with Link
0	0	1	0	0	0	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	OFF	BRA	Unconditional branch (branch always)
0	0	1	0	0	1	С	С	С	С	Т	Т	Т	F	F	F	CEX	Conditional execution
0	0	1	0	1	0	0	0	0	0	0	0	0	PR	PR	PR	SETPRI	Set current priority
0	0	1	0	1	0	0	0	0	0	0	1	SA	SA	SA	SA	SVC	Control passes to address specified in vector [SA]
0	0	1	0	1	0	0	0	0	0	1	V	SLP	N	Z	С	SETCC	Set PSW bits (1 = set)
0	0	1	0	1	0	0	0	0	1	0	V	SLP	N	Z	С	CLRCC	Clear PSW bits (1 = clear)
0	1	0	0	0	0	0	0	R/C	W/B	S/C	S/C	S/C	D	D	D	ADD	Add: DST = DST + SRC/CON
0	1	0	0	0	0	0	1	R/C	W/B	S/C	S/C	S/C	D	D	D	ADDC	Add with Carry: DST = DST + (SRC/CON + Carry)
0	1	0	0	0	0	1	0	R/C	W/B	S/C	S/C	S/C	D	D	D	SUB	Subtract: DST = DST + (¬SRC/CON + 1)
0	1	0	0	0	0	1	1	R/C	W/B	S/C	S/C	S/C	D	D	D	SUBC	Subtract with Carry: DST = DST + (¬SRC/CON + Carry)
0	1	0	0	0	1	0	0	R/C	W/B	S/C	S/C	S/C	D	D	D	DADD	Decimal-add: DST = DST + (SRC/CON + Carry)
0	1	0	0	0	1	0	1	R/C	W/B	S/C	S/C	S/C	D	D	D	CMP	Compare: DST + (¬SRC/CON + 1)
0	1	0	0	0	1	1	0	R/C	W/B	S/C	S/C	S/C	D	D	D	XOR	Exclusive OR: DST = DST ⊕ SRC/CON
0	1	0	0	0	1	1	1	R/C	W/B	S/C	S/C	S/C	D	D	D	AND	AND: DST = DST & SRC/CON
0	1	0	0	1	0	0	0	R/C	W/B	S/C	S/C	S/C	D	D	D	BIT	Bit test: DST & SCR/CON
0	1	0	0	1	0	0	1	R/C	W/B	S/C	S/C	S/C	D	D	D	BIC	Bit clear: DST = DST & ~SRC/CON
0	1	0	0	1	0	1	0	R/C	W/B	S/C	S/C	S/C	D	D	D	BIS	Bit set: DST = DST SRC/CON
0	1	0	0	1	1	0	0	0	W/B	S	S	S	D	D	D	MOV	DST = SRC
0	1	0	0	1	1	0	0	1	0	S	S	S	D	D	D	SWAP	Swap SRC and DST
0	1	0	0	1	1	0	1	0	W/B	0	0	0	D	D	D	SRA	Shift DDD right (1 bit) arithmetic
0	1	0	0	1	1	0	1	1	W/B	0	0	0	D	D	D	RRC	Rotate DDD right (1 bit) through carry
0	1	0	0	1	1	1	0	0	0	0	0	0	D	D	D	SWPB	Swap bytes in DDD
0	1	0	0	1	1	1	0	0	0	0	0	1	D	D	D	SXT	Sign-extend byte to word in DDD
0	1	0	1	1	0	PRPO	DEC	INC	W/B	S	S	S	D	D	D	LD	DST = mem[SRC plus addressing]
0	1	0	1	1	1	PRPO	DEC	INC	W/B	S	S	S	D	D	D	ST	mem[DST plus addressing] = SRC
0	1	1	0	0	В	В	В	В	В	В	В	В	D	D	D	MOVL	DST.Low byte = BBBBBBBBB; DST.High byte unchanged
0	1	1	0	1	В	В	В	В	В	В	В	В	D	D	D	MOVLZ	DST.Low byte = BBBBBBBB; DST.High byte = 00000000
0	1	1	1	0	В	В	В	В	В	В	В	В	D	D	D	MOVLS	DST.Low byte = BBBBBBBBB; DST.High byte = 11111111
0	1	1	1	1	В	В	В	В	В	В	В	В	D	D	D	MOVH	DST.Low byte unchanged; DST.High byte = BBBBBBBB
1	0	OFF	OFF	OFF	OFF	OFF	OFF	OFF	W/B	S	S	S	D	D	D	LDR	DST = mem[SRC + sign-extended 7-bit offset]
1	1	OFF	OFF	OFF	OFF	OFF	OFF	OFF	W/B	S	S	S	D	D	D	STR	mem[DST + sign-extended 7-bit offset] = SRC