Circumscription

Revision: February 13, 2018; Rendered: May 25, 2021

Definition of predicate circumscription. Formalized with the PIE system.

1 Definition of Predicate Circumscription

xcirc(S, F)

Defined as

$$F \wedge \neg xraise(S, F)$$
.

A version of parallel predicate circumscription [Lif94]. The F parameter is the circumscribed formula. The S parameter specifies the roles of the predicates in the circumscription. It is a list of specifiers of the following form, where p is a predicate and the second component indicates positive, negative, or both polarities:

p is to be minimized p-n p is to be maximized p-p p is varying p-pn p is fixed p-pn p is not mentioned in p-pn

In some cases it might be necessary to specify explicitly also the arity of the respective predicates, e.g. p/1-n.

The S argument is considered complementary to the scope argument to circ in [Wer12] (like in forgetting instead of projection).

xraise(S, F)

Defined as

$$\exists Q \, (F_1 \wedge T_1 \wedge \neg T_2),$$

where

 $F_1 := F[S \mapsto Q],$ $S_2 := S$ with arities from F, $T_1 := \text{transfer clauses } S_2 \to Q,$ $T_2 := \text{transfer clauses } Q \to S_2.$

The second-order subformula on which xcirc is based, Similar to raise [Wer12], however the scope argument S is considered complementary (like in forgetting instead of projection).

2 Some Examples

These are examples from [Lif94].

Input: xcirc([p-n], pa). Result of elimination:

$$pa \wedge \forall x (px \rightarrow x = a).$$

Input: $xcirc([p-n], \neg pa)$. Result of elimination:

 $\forall x \neg px.$

Input: $xcirc([p-n], (pa \land pb))$.

Result of elimination:

$$\mathsf{pa} \wedge \mathsf{pb} \wedge \forall x \, (\mathsf{p}x \to x = \mathsf{a} \vee x = \mathsf{b}).$$

 $\mathrm{Input:}\ \mathit{xcirc}([p\text{-}\mathsf{n}],(\mathsf{pa}\vee\mathsf{pb})).$

Result of elimination:

$$(\mathsf{pa} \lor \mathsf{pb}) \land \forall x (\mathsf{p}x \to (\mathsf{pa} \to x = \mathsf{a}) \land (\mathsf{pb} \to x = \mathsf{b})).$$

This formula is valid: $last result \leftrightarrow \forall x (px \leftrightarrow x = a) \lor \forall x (px \leftrightarrow x = b).$

Input: $xcirc([p-n], (\neg pa \lor pb))$.

Result of elimination:

$$\forall x \neg px.$$

Input: $xcirc([p-n], (pa \lor (pb \land pc)))$.

Result of elimination:

$$(\mathsf{pa} \lor (\mathsf{pb} \land \mathsf{pc})) \qquad \land \\ \forall x \, (\mathsf{p}x \to (\mathsf{pa} \to x = \mathsf{a}) \land (\mathsf{pb} \land \mathsf{pc} \to x = \mathsf{b} \lor x = \mathsf{c})).$$

This formula is valid: $last_result \leftrightarrow \forall x \, (px \leftrightarrow x = a) \lor (\forall x \, (px \leftrightarrow x = b \lor x = c) \land a \neq b \land a \neq c)$.

Input: $xcirc([p-n], \forall x px)$.

Result of elimination:

 $\forall x \, \mathsf{p} x.$

Input: $xcirc([p-n], \forall x (qx \rightarrow px)).$

Result of elimination:

$$(\forall x (\mathsf{q} x \to \mathsf{p} x) \to \forall x (\mathsf{p} x \to \mathsf{q} x)) \land \forall x (\mathsf{q} x \to \mathsf{p} x).$$

Input: $xcirc([p-n], \exists x px)$.

Result of elimination:

$$\forall x (px \to \forall y (py \to x = y)) \land \exists x px.$$

This formula is valid: $last result \leftrightarrow \exists x \, \forall y \, (py \leftrightarrow x = y).$

Input: $xcirc([p-n], \forall x pxx)$.

Result of elimination:

$$(\forall xy \, (\mathsf{p} xy \to x = y) \lor \exists x \, \neg \mathsf{p} xx) \land \forall x \, \mathsf{p} xx.$$

Input: $xcirc([p-n, q-pn], \forall x (qx \rightarrow px)).$

Result of elimination:

$$\forall x (qx \to px) \land \forall x \neg px.$$

This formula is valid: $last result \leftrightarrow \forall x \neg px \land \forall x \neg qx$.

$block_axioms$

Defined as

$$\begin{array}{ll} \forall x \, (\mathsf{block}(x) \wedge \neg \mathsf{ab}(x) \to \mathsf{ontable}(x)) & \wedge \\ \neg \mathsf{ontable}(\mathsf{b}_1) & \wedge \\ \mathsf{block}(\mathsf{b}_1) & \wedge \\ \mathsf{block}(\mathsf{b}_2) & \wedge \\ \mathsf{b}_1 \neq \mathsf{b}_2. \end{array}$$

Input: $xcirc([ab-n, ontable-pn], block_axioms)$.

Result of elimination:

$$\begin{array}{ll} \mathsf{b}_1 \neq \mathsf{b}_2 & \wedge \\ \neg \mathsf{ontable}(\mathsf{b}_1) & \wedge \\ \mathsf{block}(\mathsf{b}_1) & \wedge \\ \mathsf{block}(\mathsf{b}_2) & \wedge \\ (\mathsf{ab}(\mathsf{b}_1) \to \forall x \, (\mathsf{ab}(x) \to x = \mathsf{b}_1 \wedge \mathsf{block}(x))) \wedge \\ \forall x \, (\mathsf{block}(x) \to \mathsf{ab}(x) \vee \mathsf{ontable}(x)). \end{array}$$

This formula is valid: $last_result \leftrightarrow block_axioms \land \forall x (ab(x) \leftrightarrow x = b_1).$

2.1 Auxiliary Macros

last_result

Defined as

X,

where

last ppl result(X).

References

- [Lif94] Vladimir Lifschitz. Circumscription. In Dov M. Gabbay, C. J. Hogger, and J. A. Robinson, editors, Handbook of Logic in Artificial Intelligence and Logic Programming, volume 3, pages 298–352. Oxford University Press, 1994.
- [Wer12] Christoph Wernhard. Projection and scope-determined circumscription. *Journal of Symbolic Computation*, 47:1089–1108, 2012.

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