6.4   Reach Set Approximation

Motivation: *Reach set* is strong tool for *Obstacle Avoidance* because it contains all possible *avoidance maneuvers*. The current implementations (sec. **??**) have following flaws:

1. *Realistic approximation - nonlinear systems* or *heavily constrained systems* cannot be approximated well by *linear continuous-time Reach Sets.*

2. *Finite count of possibilities* - continuous-time *Reach Set* contains infinite possibilities for *avoidance maneuvers*, the SAA system demands conflict resolution in finite time.

3. *Computationally feasible data structures* - binding related properties seems problematic, because *continuous- time reach sets* does not have unique identifier of maneuver, trajectory nor segment.

Proposed Solution Features: Our Reach set Estimation method will provide following features:

1. *System Control Interface* - implemented via *Movement Automaton*, requiring only *discrete command chain* to approximate system behaviour.

2. *Finite count of possibilities* - finite number of elements in *Reach set* will enable *scalable* calculation.

3. *Computationally feasible data structures* - approximation of Reach set as a set of trajectories, each trajectory can be split into finite number of segments. Each element will have unique identifier enabling both-side property binding.

4. *Computationally feasible data-structures* - some specific behavior, like horizontal/vertical separation, or maneuver shape can be encoded into a different types of the reach set approximation algorithms.

Discretization of Reach set: There is a need for a discrete finite *Reach Set approximation* to enable *Avoidance Strategy Evaluation* in finite time. Replacing *Continuous Control Set Inputs(t)* by *Movement Automaton* is feasible:

Definition 1 (Reach set Approximation by Movement Automaton). A trajectory *(def.* **??***) for system state* $= f(time, state, input)$ *under control of the movement automaton* $\mathcal{MA}$ *is given as execution of movement buffer (def.* **??***) with initial state of system* $state_0$*. Therefore notation* $Trajectory(state_0, buffer)$ *is used.*

The Complete Reach Set *(6.1) for system with initial state $state_0$ with existing* control strategy $control(time) \in Controls(time).$ *for time $\tau > time_0$.*

$$ReachSet(\tau, time_0, state_0) = \bigcup \{state(s) : control(s) \in Controls(s), s \in (time_0, \tau]\}$$
(6.1)

The Reach Set Approximation by Movement Automaton *(6.2) of the system under the control of the movement automation $\mathcal{MA}$ consist from the set of trajectories $Trajectory$ ($state_0$, $Buffer$), which are executed in constrained time $\tau > time_0$.*

$$ReachSet(\tau, time_0, state_0) = \left\{ Trajectory(state_0, buffer) : \begin{array}{c} duration(buffer) \\ \leq \\ (time_0 - \tau) \end{array} \right\} \quad (6.2)$$

Note. Reach Set Approximation (def. 1) is subset of *Full Reach Set* (def. **??**) in continuous space $\mathbb{R}^n$ it inherits all important properties, like *Invariance* [1].

Discretization of *Reach Set* have been achieved leaving us with *finite count* of *Trajectories*, instead of *Infinite subspace or $\mathbb{R}^N$*

Approximated Reach Set Containment: The *Approximated Reach Set* introduced in (def. 1) is constrained only by *future expansion time τ*. UAS makes space assessment in *Avoidance Grid.* There is no point to consider Trajectories outside of *Avoidance Grid*

Definition 2 (Contained Aproximated Reach Set). *For pair ($state_0$, $AvoidanceGrid_0$) at time $time_0$ and* prediction horizon $\tau = \infty$ *there is* Contained Reduced Reach Set:

$$ReachSet \begin{pmatrix} time_0, \\ state_0, \\ AvoidanceGrid_0 \end{pmatrix} = \left\{ \begin{array}{c} Trajectory(\dots) \\ \in \\ ReachSet(6.2) \end{array} : \begin{array}{c} \forall segment \in AvoidanceGrid_0, \\ segment \in Trajectory(\dots) \end{array} \right\}$$
(6.3)

Properties: Container Aproximated Reach Set *contains only trajectories where all segments belongs to* Avoidance Grid, *there are following functions:*

1. Membership function *for any* Trajectory *in* Constrained Reduced Reach set *returns* Ordered Set *of* Passing Cells.

2. Cost function *for any* Trajectory Portion *in* Constrained Reduced Reach Set *return* Cost of Execution

Passing cell: Cell *of* Avoidance Grid *which has some intersection with Trajectory.*

Note. Contained Reduced Reach Set (eq. 6.3) which is contained in *Avoidance Grid* and have an *Membership Function* enable Property transition between Reach set and *Avoidance grid.*

Example: Visibility from cells along *Trajectory* can be gathered to calculate *Trajectory's* feasibility.

Reach Set Pruning: There is a need to implement *Set Difference* between *Reach Set* and *Constraint Set.* Constraint Set cam be *Obstacle Set* from *Known World* (sec. **??**) and other different constraints.

Reach Set Trajectory Tree: (6.4) *Any Reach Set* where *Control Strategy Constraint* is implemented as *Movement Automaton*, with defined *Movements* set and for singe initial $state_0$. The *Reach Set* is given as discrete tree with root $Trajectory(state_0, \varnothing)$.

$$ReachSet(state_0, \dots) = \left\{ Trajectory(state_0, buffer) : \begin{array}{c} buffer \in Movements^i, \\ i \in \{1, \dots, k\} \end{array} \right\} \quad (6.4)$$

For each *Trajectory Segment*, there exists *intersection function* which evaluates as true if there exists at least one point in *Segment* which belongs to *Constraint Set.* Formally:

$$intersection(segment, Set) : \begin{cases} \exists point \in segment, & \\ \quad point \in Set & : true \\ Otherwise & : false \end{cases} \quad (6.5)$$

Definition 3 (Pruned Reach Set). *For* Reach set *represented as* Trajectory Tree *(eq. 6.4) and some constraint set (Set) where exist* intersection function *(eq. 6.5). The* Pruned Reach set *is given as follows:*

$$Prune(ReachSet, Set) = \left\{ Trajectory(\dots) : \begin{array}{c} \forall segment \in Trajectory, \\ \neg intersection(segment, Set) \end{array} \right\} \quad (6.6)$$

Note. Pruning(def. 3) [2] is applicable multiple times for various *Constraints Set.*

Example of *Approximated Reach set Calculation* (def. 1), *Reach Set Containment* (def. 2), and, *Pruning* is given in [3].

6.4.1   Reach Set Performance Criteria

Motivation: The need to Make *Reach Set* scalable approach. This may be a problem due the *Expansion rate. Reach set* represented as a *Trajectory Tree* (eq. 6.4) for Avoidance Grid with *layerCount* and Movement automaton with *movementCount*, the *Node count* is given as:

$$1 + \left(\sum_{i \in \{1...layerCount\}} (movementCount)^i \right) \tag{6.7}$$

This scaling is not feasible for *Avoidance Grid* with many layers (< 10) or *Movement Set* with many movements (< 9). There is need for *Reduced Reach set calculation*.

Core Performance Criteria: The scaling factor (eq. 6.7) shows that there are going to be many trajectories. The main point is that not every trajectory in *Reach Set* are giving us *maneuverability advantage*. Our expectations lies in following *Performance Requirements*:

1. *Reach set* must *Cover* maximum of the *possible unique maneuvers* in *Avoidance Grid*.

2. *Trajectories* in *Reach Set* should be smoothest possible to prevent cargo damage / UAS wear.

Trajectory footprint: Discrete space of *Avoidance Grid* is organized in cells. *Cell* is minimal space portion accessible by *property binding*. There is need to know if two trajectories contribution to *Maneuverability* in this environment.

Each trajectory passes through space in *Avoidance Grid*. If there exists a method to extract unique identifier for each *trajectory passed cells*, we can compare two trajectories *Coverage* in *Avoidance Grid*.

Definition 4 (Trajectory footprint). *For* Trajectory *from* Reach set *(def. 2) defined for* Avoidance Grid *has membership function.* Membership Function *returns* ordered set of passing cells:

$$footprint \begin{pmatrix} Trajectory, \\ AvoidanceGrid \end{pmatrix} = \left\{ \begin{array}{l} cell \in AvoidanceGrid : \\ \qquad isMember(trajectory, cell) \end{array} \right\} \tag{6.8}$$

Then we can define equality function for $Trajectory_1$ *and* $Trajectory_2$, *as comparison of their footprints in common* Avoidance Grid *as follow:*

$$isEqual \begin{pmatrix} Trajectory_1, \\ Trajectory_2, \\ AvoidanceGrid \end{pmatrix} : \left\{ \begin{array}{ll} \begin{pmatrix} footprint(Trajectory_1, \dots) \\ = \\ footprint(Trajectory_2, \dots) \end{pmatrix} & : true \\ Otherwise & : false \end{array} \right. \tag{6.9}$$

Note. Depending on *Movement Automaton's* movement set and *Avoidance Grid* parameters, there can be multiple *trajectories* which are equal.

Coverage set: Now it is possible to create set of unique *trajectory footprints* due to *footprint function* (eq. 6.8). Similarly there is a possibility to create *Reach set skeleton* containing unique trajectories, by using *equality function* (eq. 6.9). *Coverage set* is sufficient for now.

Definition 5 (Coverage Set)**.** Coverage set *(6.10) is defined for* Avoidance Grid *and* Reach Set *pair as set of unique* Trajectory footprints*:*

$$CoverageSet \begin{pmatrix} AvoidanceGrid, \\ ReachSet \end{pmatrix} = \left\{ footprint \begin{pmatrix} Trajectory, \\ AvoidanceGrid \end{pmatrix} : \begin{matrix} \forall Trajectory \\ \in ReachSet \end{matrix} \right\}$$

$$(6.10)$$

Coverage set properties: Trajectory footprint (eq. 6.8)is not *bijection*, neither *injection* for $ReachSet \rightarrow CoverageSet$. This implies following properties:

1. Equal *Reach Sets* in same *Avoidance Grid* have equal *Coverage Sets*.

2. Equal *Coverage Sets* does not imply *Reach Set* equality.

3. For two Coverage Sets there is a possibility to compare their member count to create coverage ratio.

The second *Property* gives us a preposition that there is a possibility of *Reach Set Reduction* without loosing *Coverage*.

Definition 6 (Coverage Ratio)**.** Coverage Ratio *is a ratio of* Coverage Set Member Count *between two* Reach Sets*. Reach set with* lesser count of unique Trajectories *is considered as* Reduced Reach Set*. Reach set with* greater Count of unique Trajectories *is considered as* Reference Reach Set*.*

$$referenceCoverage = |CoverageSet(ReferenceReachSet, AvoidanceGrid)|$$
$$reducedCoverage = |CoverageSet(ReducedReachSet, AvoidanceGrid)|$$
$$CoverageRatio = \frac{reducedCoverage}{referenceCoverage} \in [0,1]$$

$$(6.11)$$

Note. Reference Reach Set is usually *Full Reach Set* containing all possible trajectories in space contained by *Avoidance Grid.* In case *Full Reach Set* cannot be computed, Avoidance Grid is too large, most complex *Reach Set* is used as *Reference Reach Set.*

Trajectory smoothness: Trajectory other than straight line have some changes in *UAS* heading.

The goal is to minimize *Maneuvering* of UAS, because:

1. *Every Heading Change* needs to be reported to *UTM*.

2. *Sharp Maneuvering* can damage cargo/wear UAS.

3. *Often course changes* makes *Intruder prediction* harder for other Civil General Aviation.

For this purpose *Smoothness Metric* needs to be applied for *Reach Set* or *Trajectory*. In case of *Movement Automaton Control* two distinguish *Movement Sets* can be introduced:*Smooth* nad *Chaotic* movements set with following properties:

$$MovementSet = SmoothMovements \cup ChaoticMovements$$
$$SmoothMovements \cap ChaoticMovements = \varnothing \tag{6.12}$$
$$|SmoothMovements| > 0, \quad |ChaoticMovements| > 0$$

Then *Smoothnes clasificator* for $Trajectory(initialState, buffer)$ can be defined as *isSmooth* and *Smooth Movement Counter* function as *smoothCount* like follow:

$$isSmooth(movement) = \begin{cases} movement \in SmoothMovements & : 1 \\ movement \in ChaoticMovements & : 0 \end{cases}$$
$$smoothCount(Trajectory(\dots, buffer)) = \sum isSmooth(movement),$$
$$\forall movement \in Buffer \tag{6.13}$$

Definition 7 (Smoothness Rating for Trajectory). Smoothness *for trajectory generated by* Movement Automaton *for some* Initial State *with some* Movement Buffer, *under assumption of* Smooth and Chaotic Movement Set *split (eq. 6.12), with existing* classification *and* counter *functionals (eq. 6.13) is given as follows:*

$$Smoothness(Trajectory(\dots, buffer)) = \frac{isSmooth(Trajectory)}{movementCount(Trajectory)} \in [0, 1] \tag{6.14}$$

For Trajectory *with* $buffer = \varnothing$ Smoothness *is given as 1.*

6.4.2   Constrained Trajectory Expansion

Motivation: *Purpose* of *Navigation* is to move forward to *Goal Waypoint* in *Mission*. *Structure* of *Avoidance Grid* is designed to enable *forward* and *turning* maneuvers. The *Avoidance Grid* is organized in *Layers* characteristic by same distance from *Avoidance Grid Origin*.

Survey of motion planning algorithm was given in [4]. The ideal candidate for propagation algorithm is *Wave-front* algorithm propagating *Trajectory tree* through Layers. Due the *Avoidance Grid* onion like layers, there is possibility to implement turn maneuver through layers iterative and effectively .
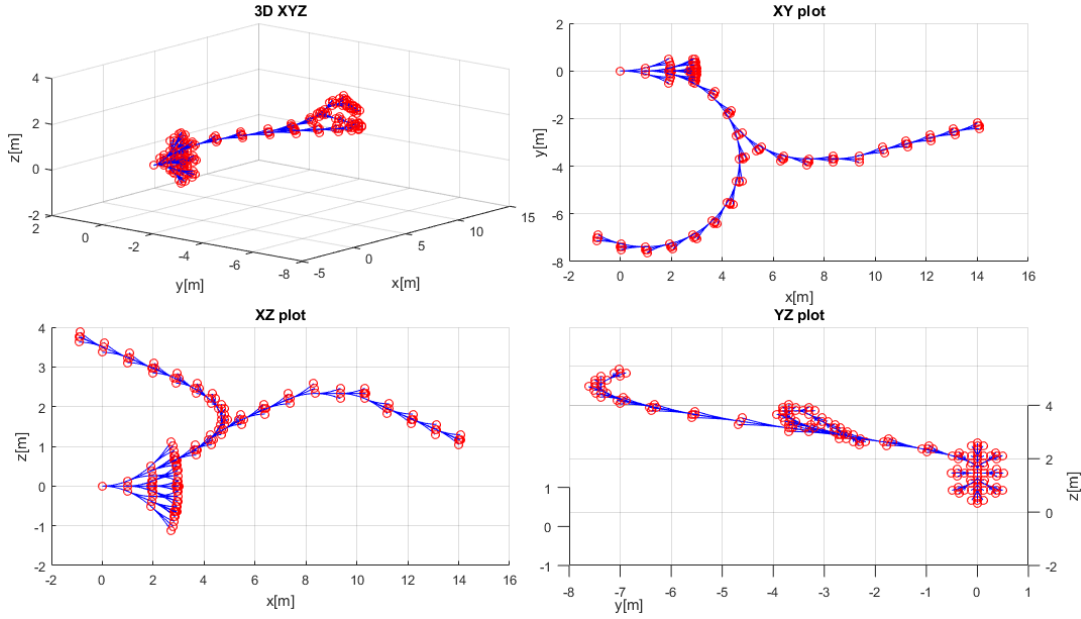
Figure 6.1: *Rapid Exploration tree as result of* Constrained trajectory expansion.

Rapid Exploration Tree (fig. 6.1) was selected, because it enables *Movement Automaton Utilization* and *Property Binding*. Similar approach was used for space exploration [5].

The example (fig. 6.1) shows a *Rapid Exploration Tree* in *Free Space* containing *Waypoint Navigation Path* and *Turn Away Path*. Both paths are starting in same *Root Node* (red circle) which was expanded with simple *Movement Automaton* (bunch of nodes originating from one node are showing way of expansion). The connection (blue line) between two nodes (red circles) represents *Trajectory portion* for *Executed Movement*.

Rapid Exploration Tree Node will contain following information:

1. *Initial state* - root entry point, used in state evolution calculation.

2. *Trajectory (state evolution)* - trajectory passing through *state space* in local coordinate frame of *Avoidance Grid*.

3. *Buffer* (applied movements) - ordered list of *executed movements* applied on *initial state* to obtain *state evolution*.

4. *Cost* - calculated for *state evolution* based on *predefined cost function*.

5. *Footprint* - ordered set of *passing cells* in *Avoidance Grid*.

6. *Parent Node Reference* - tree reference for parent node, not in case of *root node*.

7. *Other Bounded Properties* - value list of other properties, depending on *Expansion Constraints* and *Reachability* evaluation algorithm.

Wave-front propagation of Rapid Exploration Tree   is given in (alg. 6.1).

The *Avoidance Grid* have UAS with *position* $\in$ *Initial State* at the *origin*. The *Grid Layer* is a column ordered set of cells with same *Mean distance* from origin. *Grid Layers* are indexed from origin starting with 1, there is maximum of $i \geq 1$ layers.

Step: Initialization   contains base structure preparation like follows:

1. *Avoidance Grid* - Space containing *Reach set* (def. 2).

2. *Movement Automaton* - Used as *Predictor*, consuming *buffer* containing *Movements* to generate $Trajectory(initialState, buffer)$.

3. *Reach Set* - tree consisting from *Wave-frontNodes* representing the end point of $Trajectory(initialState, buffer)$ where each *Edge* represents *one Movement application*. The root is set as node containing *Initial State*.

Function *initializeReachSet(root,stack,grid,automaton)* will take the root and enforces *full wavefornt propagation* to *First Layer*.

Step: Wave-front Propagation   is forced propagation of trajectories from layer i to layer $i + 1$. The process goes as follows:

1. *Selection of Feasible candidates* - function *[candidates,leftovers] = ExpansionConstraints.select(stack)* for working layer, row and cell selects *feasible trajectory nodes* ordered by *Cost function*. The *Example of Cost Function* can be *Trajectory Smoothness* (def. 7).

2. *Expansion of Candidates* - for each *candidate* function *candidate.expandNode (automaton)* is invoked. This function will expand *Candidate Node structure* by appending *Full Trajectory Tree Evolution* until each *Leaf Trajectory* reaches *Next Layer*. Simply put *Par rent Node Node(initialState, buffer, cost, footprint)* buffer is appended by movements until the next layer is reached.

3. *Leftovers purge* - function *reachSet.purge(leftovers)* removes unexpanded *Nodes* leading to cell, effectively removing trajectories which does not lead to *next layer*.

4. *Append Reach Set* - function *reachSet.append(leafs)* puts newly created *Nodes (Trees)* into *Reach Set* structure. The *Wave-front Propagation* for one cell is finished.

Step: After Layer Propagation Purge   is covered by function *reachSet.purgeSameFootprint()* which takes trajectories with same footprint and keeps some of them based on *Selection criteria*, more in (sec. 6.4.3, 6.4.4). *Pruning methods* over *Large Decision Trees* are *fast* and *viable* [6].

Note. Reach Set is usually computed *Prior the Flight* for *some Initial State* in *Local Coordinate Frame* in *right had coordinate frame* with X^+ used as *main axis.*

Algorithm 6.1: *Wave-front propagation* of *Rapid Exploration Tree* to form *Reach Set.*

Input : Node(initialState,buffer=∅,cost=0,footprint=∅), AvoidanceGrid,
　　　　 ExpansionConstraints, MovementAutomaton(movementSet)

Output: ReachSet(AvoidanceGrid)

\# Initialization Sequence;

grid=AvoidanceGrid, automaton=MovementAutomaton, root = Node;

reachSet = initializeReachSet(root,stack,grid,automaton);

\# Main Expansion through, layers (i), rows (j), cells(k);

for *layer(1 . . . i)* in *grid* **do**

   for *row(1 . . . j* in *layer)* **do**

      for *cell(1 . . . k)* in *row* **do**

         \# apply selection criteria ;

         [candidates,leftovers] = ExpansionConstraints.select(stack);

         \# collect expansions ;

         leafs = [];

         for *candidate* in *Candidates* **do**

           | leafs= [leafs, candidate.expandNode(automaton)];

         end

         reachSet.purge(leftovers);

         reachSet.append(leafs);

      end

   end

   reachSet.purgeSameFootprint();

end

6.4.3   Chaotic Reach set

Motivation: Design of calculation method for *Reach Set Approximation* guarantying high *Maneuverability.*

Background: There is *Coverage Ratio* property of *Reach Set* (def. 6). It has been shown that creating *Reach Set* via *greedy approach* is not feasible due the *Scaling Factor. Contracted Expansion* (sec. 6.4.2) is enabling to apply selection criteria while building *Reach Set* in given *Cell.*

The *Cell* $cell_{i,j,k}$ has a center and walls from UAS viewpoint: front wall , back wall (for $layer > 1$), top wall, left wall, right wall, bottom wall. It is expected that trajectory leading close to one cell walls will continue to different cell, increasing chance to obtain more *Unique Footprints.*

Expansion Constraint Function Implementation   (alg. 6.2) is based on simple principle: *Select candidate Nodes which are closest to outer walls of Cell, with unique footprint.*

Tuning Parameters   : *Proximity to Cell outer wall* gives good chances to break into other rows or columns in *Avoidance Grid. Unique footprint* guarantees future *Unique Footprint* after appending Trajectory by *Movement application.*

1. *Considered Footprint Length* - how much last cells in footprint should be considered in unique path track, minimal value 1, default value 3, maximal value ∞. If you want to generate non redundant trajectories use ∞, it will consider full footprint.

2. *Spread Limit* - upper limit of candidates which are going to be select for further expansion, minimal value 1, default value *Count of unique Moves in Movement set*, maximal value ∞. If more than default values is selected the algorithm will generate *redundant trajectories.* If less is selected then some trajectories are omitted and *Coverage Rate* decreases sharply.

Step: Initialization   initialization of *candidate* array (return value), *leftovers* array (return Value). Node array *passing* is populated with *Nodes* which represents *end node of Trajectory* and the tip of *trajectory is constrained in* $cell_{i,j,k}$.

Step: Evaluate best trajectories with unique Footprints   following steps are executed:

1. *Best Performance Map* is created with *footprint* as key set element to ensure footprint uniqueness.

2. *Wall distance* for *test node* is calculated as a closest trajectory portion distance to *top, bottom, left, right* wall of cell $cell_{i,j,k}$

3. *Footprint* for *test node* is created with maximal length given by *Footprint Length* tuning parameter.

4. *Existence and Performance Test* is executed to ensure that best performing node is selected. If there is not key entry in the *Best Performance Map*, then new entry for *Test Node* is created. If there is key entry, the performance of *Old Node* and *Test Node* is compared and better is stored.

Step: Select candidates is executed on *Best Performance Map* records using *Wall distance* as pivot parameter, ordering by closest proximity and limited by *Search Limit* tuning parameter. The *Leftovers* are difference set between *Passing Nodes* and *Candidate Nodes*.

Algorithm 6.2: Expansion Constraint function for *Chaotic Reach Set Approximation*

    Input   : Node[] stack, Cell $cell_{i,j,k}$

    Tuning Parameters: int^+ footprintLength, int^+ spreadLimit

    Output   : Node[] candidates, Node[] leftovers

    # Initialize structures;

    Node[] candidates = [], Node[] leftovers=[];

    Node[] passing = $cell_{i,j,k}$.getFinishingTrajectories(stack);

    # Select best performing trajectories with unique footprint;

    Map<Footprint,Node> bestPerformanceMap;

    for *Node test* ∈ *passing* **do**

        wallDistance= test.minimalDistanceToWall($cell_{i,j,k}$)];

        footPrint = test.getFootprint(lastCells = footprintLength);

        if *bestPerformanceMap.contains(footPrint)* **then**

            old = bestPerformanceMap.getByKey(footprint);

            oldPerformance= old.minimalDistanceToWall($cell_{i,j,k}$);

            if *oldPerformance > wallDistance* **then**

                bestPerformanceMap.setByKey(footprint,test);

            end

        else

            bestPerformanceMap.setByKey(footprint,test);

        end

    end

    # Select best performing nodes up to *spreadLimit* count;

    candidates = bestPerformanceMap.select(count = spreadLimit).orderBy('wallDistance','Ascending');

    leftovers = passing - candidates;

    return *[candidates,leftovers]*

Example: for *Avoidance Grid* with *Distance 10 m*, *Layer count 10*, *Horizontal range* $[-45°, +45°]$, *Horizontal Cell Count 7*, *Vertical range* $[-30°, +30°]$, and *Vertical Cell Count 5*. Is given in (fig. 6.2). The UAS is at *Back-side* of *Figure* (initial state is at all *Trajectory Origins*). The *black dashed line* marks *Avoidance Grid* space boundary. Each trajectory has its own color and ends at *Front-side* of *Avoidance Grid Boundary*.
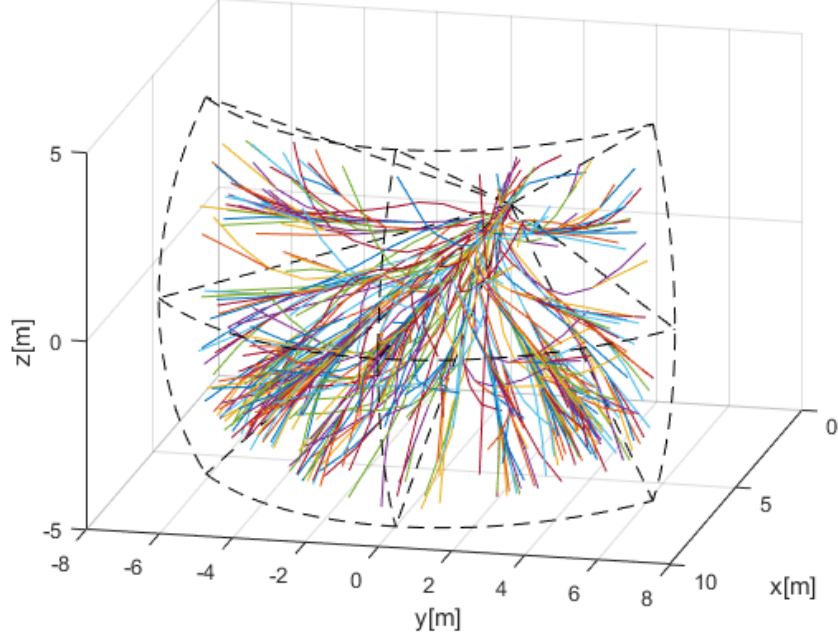

Figure 6.2: *Chaotic* reach set *approximation*.

Pros and Cons: It can be seen from example (fig. 6.2) that *Chaotic Reach Set Approximation Method* (alg. 6.2) generates a lot of *turning* and *shaky trajectories*.

High Coverage Ratio (~ 0.9) is provided, while keeping *medium node count*. The calculation complexity scales linearly with grid size. The *upper limit of trajectories* is given as follow:

$$countTrajectories(ReachSet) \leq layerCellCount \times spreadLimit$$
$$\times size(Movements) \quad (6.15)$$

The *upper limit of nodes* is given as follow:

$$countNodes(ReachSet) \leq layerCount \times layerCellCount$$
$$\times size(Movements) \times spreadLimit \quad (6.16)$$

Absence of *Smooth Trajectories* disqualifies *Chaotic Reach Set Approximation* to be used for *Navigation*. This type of reach set is feasible for *Avoidance*, because it contains variety of maneuvers.

6.4.4 Harmonic Reach set

Motivation: Imagine having an *Avoidance Grid* like (fig. **??**). There is a need of *Reach Set Approximation* which will have *Smooth Trajectories* (def. 7) going nearby *cell* centers.

Background: The *Smoothness Rating for Trajectory* (def. 7) uses two distinct sets *Smooth Movements* and *Chaotic Movements* (eq. 6.12) which are defined for our *Movement Automaton* (sec. **??**) like follow:

$$
\begin{aligned}
SmoothMovements &= \{Straight\} \\
ChaoticMovements &= Movements - SmoothMovements
\end{aligned}
\tag{6.17}
$$

Smooth Movements contains only *Straight* movement, because others are considered as extreme turning movements. *Smooth Movements* should contain only direct flight movements or slight heading correction. *Chaotic Movements* set is supplement of *Movement Automaton's Movement Set*.

The *Avoidance Grid* (fig. **??**) cell centers for fixed indexes j_{fix}, k_{fix} are linearly aligned with *initial state*. That means that cell centers of cells $cell_{1,j_{fix},k_{fix}}, \ldots, cell_{i,j_{fix},k_{fix}}$, where i is count of *layers* lies on one line. If the trajectory can achieve *cell center* on some *layer* only minor trajectory corrections are required to stay on given line. This type of trajectory gives us following advantages:

1. *Minimal steering at beginning* - the minimal steering is advantageous in *Controlled Airspace* because is diminishing the amount of communication to *UTM Service*.

2. *Additional safe space in Linear segment* - once the *center of cell* is reached, *Trajectory* sticks to line between cell centers. Each point on this line has *maximal distance* to outer walls of cell. This gives us extra space given as minimum of distance between *UAS position* and *Outer cell walls*.

Expansion Constraint Function Implementation (alg. 6.3) is based on simple principle: *Select candidate Nodes which are closest to Cell center, with unique footprint.*

Note. *Cell center* can be closely reached by *smooth movement* from previous cell or *chaotic movement* from neighbouring cell from current or previous layer. These trajectories are usually equivalent in *Smoothness*.

Tuning Parameter: *Proximity to Cell Center* gives a good chance to keep trajectory smooth or *smooth after one correction maneuver*. It has been mentioned that *Cell Center* can be reached by various trajectories. In this method full footprint length is always considered, therefore only one tuning parameter can be offered:

1. *Spread Limit* - upper limit of candidates which are going to be selected for further expansion, minimal value 1, default value *Count of unique Moves in Movement set*,

maximal value ∞. If maximal value ∞ is selected, algorithm will generate skeleton of *Reach Set* with full Coverage and with the smoothest *Trajectories*.

Step: Initialization   sets candidate *Nodes* as empty set, leftover *Nodes* as empty set. and selects all *Nodes* from *Stack* which represents *Finishing Trajectories* in working cell $cell_{i,j,k}$.

Algorithm 6.3: Expansion Constraint function for *Harmonic Reach Set Approximation*

Input	: Node[] stack, Cell cell$_{i,j,k}$
Tuning Parameters:	int$^+$ spreadLimit
Output	: Node[] candidates, Node[] leftovers

\# Initialize structures;
Node[] candidates = [], Node[] leftovers=[];
Node[] passing = cell$_{i,j,k}$.getFinishingTrajectories(stack);

\# Select unique smoothest trajectories;
Map<Buffer,Node> bestPerformanceMap;
for *Node test $\in$ passing* **do**
    centerDistance= test.getPerformance(cell$_{i,j,k}$)];
    footPrint = test.getFootprint();
    if *bestPerformanceMap.contains(footPrint)* **then**
        old = bestPerformanceMap.getByKey(footprint);
        oldPerformance= old.getPerformance(cell$_{i,j,k}$);
        if *oldPerformance > centerDistance* **then**
            bestPerformanceMap.setByKey(footprint,test);
        end
    else
        bestPerformanceMap.setByKey(footprint,test);
    end
end

\# Select best performing nodes up to *spreadLimit* count;
candidates = bestPerformanceMap.select(count =
  spreadLimit).orderBy('cellCenterDistance','Ascending');
leftovers = passing - candidates;
return *[candidates,leftovers]*

Step: Evaluate smoothest trajectories with unique Footprints   is implemented as *multi-criteria filtration*.

    First criterion is *distance to Cell Center* which is penalized by trajectory *smoothness rate* implemented in method *Node.getPerformance(Cell cell$_{i,j,k}$)* defined as follow.

$$getPerformance(Node, Cell) = \frac{distance(Node.Trajectory, Cell.Center)}{SmoothnessRate(Node.Trajectory)} \quad (6.18)$$

Distance of *Trajectory* is *enumerator*, because its considered as *base value* and is defined in interval $[0, maximalWallDistance]$. The *Smoothness Rate* is in denominator, because it is a penalization coefficient defined in interval $[0, 1]$.

Second criterion is *trajectory uniqueness* This is provided by *Best Performance Map*, where best performing *Node* belongs to one unique *trajectory footprint*. The implementation is identical to *chaotic reach set expansion* (alg. 6.2).

Step: Select candidates is executed on *Best Performance Map* records using *Penalized Cell Center Distance* as pivot parameter, ordered in ascending order and limited by *Spread Limit* tuning parameter. The *Leftovers* are difference set between *Passing Nodes* and *Candidate Nodes*.

Example: for *Avoidance Grid* with *Distance 10 m, Layer count 10, Horizontal range* $[-45°, +45°]$, *Horizontal Cell Count 7, Vertical range* $[-30°, +30°]$, and *Vertical Cell Count 5*. Is given in (fig. 6.3). The UAS is at *Back-side* of *Figure* (initial state is at all *Trajectory Origins*). The *black dashed line* marks *Avoidance Grid* space boundary. Each trajectory has its own color and ends at *Front-side* of *Avoidance Grid Boundary*. The *Spread Limit* in this case was set to 9 which is *Size of the Movement Set*.

Note. Please note *Trajectories* are organized in bundles going around *Cell Centers* smoothly. Most of the steering maneuvers are executed at the *beginning* of *Avoidance Grid*.
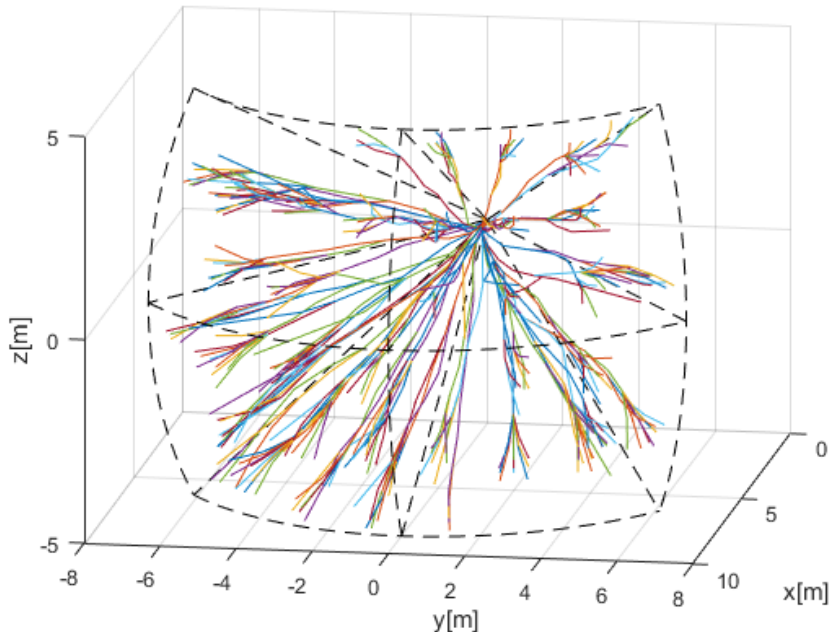

Figure 6.3: *Harmonic* reach set *approximation*.

Pros and Cons: It can bee seen from example (fig. 6.3) that *Harmonic Reach Set Approximation Method* (alg. 6.3) generates *smooth evenly spread trajectories*.

High smoothness ratio (≥ 0.9) is provided, while keeping low node count for UAS systems. The calculation complexity scales linearly with grid size. The upper limit of trajectories is given as follow:

$$countTrajectories(ReachSet) \leq layerCellCount \times spreadLimit$$
$$\times size(Movements) \quad (6.19)$$

The *upper limit of nodes* is given as follow:

$$countNodes(ReachSet) \leq layerCount \times layerCellCount \times spreadLimit \quad (6.20)$$

Absence of *High Coverage Ratio* disqualifies *Harmonic Reach Set Approximation* to be used for *Emergency Avoidance*. This type of *Reach Set* is feasible for *Open Space Navigation* or *Controlled Airspace Navigation*. Its low turning rate in contained *Trajectories* are desired for such tasks.

6.4.5   Combined Reach Set

Motivation: Harmonic reach set (sec. 6.4.4) is *efficient* for *Navigation in* Controlled Airspace. Chaotic reach set (sec. 6.4.3) is good for *Emergency avoidance*. The need to differentiation between *Navigation* and *Emergency Avoidance* mode is necessary in *Controlled Airspace.* but not in *Non-controlled Airspace*. The combination of *Harmonic* and *Chaotic* reach sets is obvious solution.

Automatic mode switch can be provided by combination of *Navigation Reach Set* and *Avoidance Reach Set* with elevated cost function. Overall having a method to merge multiple trees would be beneficial.

Background: If two *Reach Set Approximation* were calculated for same *Avoidance Grid* and *Initial State*, using same *Movement Automaton* and *UAS model* are possible to merge.

The *Reach Set Approximation* is *tree* with *Root Node* in *initail state* with movement buffer $= \varnothing$. The *movement buffer* in each node can be used as *route trace* during merging procedure. The example two reach set merge can be given as follow, where only *latest* applied movement is taken into account.

$$
\begin{bmatrix}
\text{First Reach Set} \\
\varnothing \to \left\langle \begin{array}{l} left \to \left\langle \begin{array}{l} left \\ right \end{array} \right. \\ \varnothing \end{array} \right. \\[2em]
\text{Second Reach Set} \\
\varnothing \to \left\langle \begin{array}{l} \varnothing \\ right \to \left\langle \begin{array}{l} left \\ right \end{array} \right. \end{array} \right.
\end{bmatrix}
\to
\begin{bmatrix}
\text{Combined Reach Set} \\
\varnothing \to \left\langle \begin{array}{ll} left & \to \left\langle \begin{array}{l} left \\ right \end{array} \right. \\ right & \to \left\langle \begin{array}{l} left \\ right \end{array} \right. \end{array} \right.
\end{bmatrix}
\tag{6.21}
$$

First Reach Set contains two trajectories given by buffers {*left,left*} and {*left,right*}. *Second Reach Set* contains two trajectories given by buffers {*right,left*} and {*right,right*}. The *Combined Reach Set* contains all four trajectories.

Note. The combined tree [7] does not need to have combined amount of original *Reach Sets* trajectories. There can be *Duplicity* which means that any bounded property like *Cost* must be *calculated* again.

Combined Reach Set Calculation Function (alg. 6.4) is implemented as function *NodecombinedReachSet(...)* which takes root Node with *initial State*, *Avoidance Grid* and respective parameters for each calculation method. *Harmonic spread* for *Harmonic Reach set calculation* and *Chaotic Spread*, *Footprint Length* for *Chaotic Reach set calculation*.

Separate Reach Sets are calculated using *Wave-front propagation* (alg. 6.1) using respective *Constrained Expansion* functions for *Harmonic* (alg. 6.3) and *Chaotic* (alg. 6.2) reach sets.

Combined Reach Set is created using *Node mergeTree(...)* function. Because different cost function or *Bounded Parameters Calculation* may be applied on *Original Reach Sets.*

Cost for *each node* needs to be recalculated due to original reach sets disparity. Function *combined.applyCostFunction()* will recalculate the new cost for each node.

The Goal is to have penalization for *Chaotic behaviour*, implementation of *Automatic Mode Switch* can be done like follows:

1. *Calculate Normal Cost* for Node $Cost(Node)$ for associated trajectory: $Cost(Node.Trajectory)$.

2. *Calculate Penalization for* Chaotic Behaviour, calculate *Smoothness Rating for Trajectory* (def. 7) in interval $[0,1]$, introduce penalization with base 100%.

The final $Cost(Node)$ function is applied on each *Combined Reach Set Node* and look like follows:

$$Cost(Node) = Cost(Node.Trajectory) \times \ldots$$
$$\cdots \times (1 + (1 - SmoothnessRate(Node.Trajectory))) \quad (6.22)$$

Merge Tree Function $mergeTree(\ldots)$ implements *Outer Join* operation on two trees. Example was given in (eq. 6.21). Function is applied on *root Node* iterating over *Movements in Movement Set*, because *Movement is pivot*.

Algorithm 6.4: Reach Set Merge Function and Combined Reach Set calculation

---

Tree merge function;
Node mergeTree(*Node firstNode, Node secondNode*)

   # Try to copy reference node or return null;
   Node referenceNode = (firstNode?:(secondNode?: return null));
   Node merged = new Node(referenceNode);
   merged.leafs= [];

   # Try to fetch movement nodes if exist in any sub tree;
   for *movement* $\in$ *Movements* **do**
      firstLeaf = firstNode.getLeafFor(movement);
      secondLeaf = secondNode.getLeafFor(movement);
      newLeaf = mergeTree(firstLeaf,secondLeaf);
      if *newLeaf* $\sim=$ *null* **then**
         merged.leafs.append(newLeaf);
      end
   end
   return *merged*

Combined Reach Set calculation function;
Node combinedReachSet(*Node root, AvoidanceGrid grid,int$^+$ chaoticSpread,*
  int$^+$ harmonicSpread, int$^+$ footprintLength)
   Node chaotic = chaoticReachSet(root,grid, footprintLength,chaoticSpread);
   Node harmonic = harmonicReachSet(root,grid, harmonicSpread);
   Node combined = mergeTree(chaotic,harmonic);
   combined.applyCostFunction();
   return *combined*

Example: for *Avoidance Grid* with *Distance 10 m*, *Layer count 10*, *Horizontal range* $[-45°, +45°]$, *Horizontal Cell Count 7*, *Vertical range* $[-30°, +30°]$, and *Vertical Cell Count 5*. Is given in (fig. 6.4). The UAS is at *Back-side* of *Figure* (initial state is at all *Trajectory Origins*). The *black dashed line* marks *Avoidance Grid* space boundary. Each

trajectory has its own color and ends at *Front-side* of *Avoidance Grid Boundary*. The *Chaotic Spread* was set to 8, *Footprint Length* to 3 and *Harmonic Spread* to 1.

Note. Notice there are typical trajectories from both *Harmonic* (fig. 6.3) and *Chaotic* (fig. 6.2) *Reach Set Approximations*.
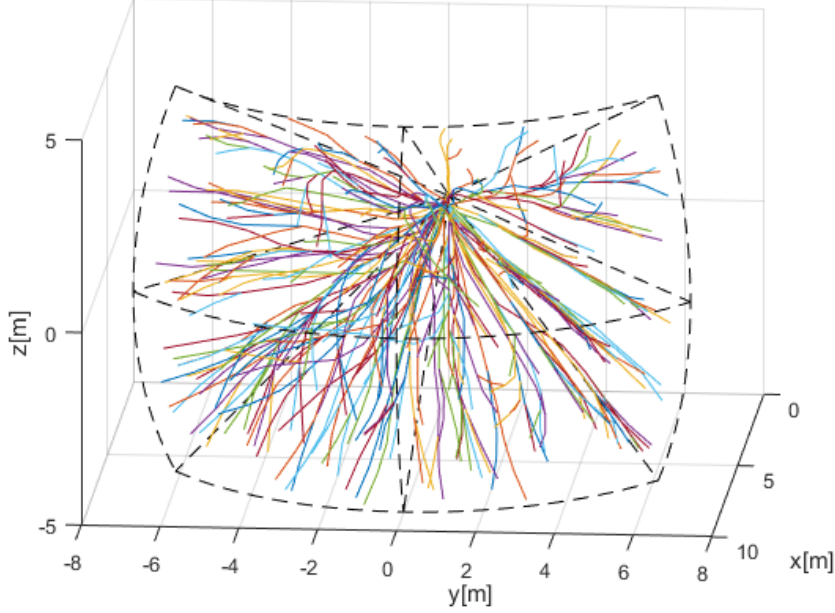

Figure 6.4: *Combined* reach set *approximation*.

Pros and Cons: It can bee seen from example (fig. 6.4) that *Combined Reach Set Approximation* (alg. 6.4) contains both types of maneuvers. *Cheaper Smooth* for navigation and *More Expensive Chaotic* for *Emergency Avoidance*. The upper limit of trajectories is given as follow:

$$countTrajectories(ReachSet) \leq countTrajectories(Chaotic)$$
$$+ countTrajectories(Harmonic) \quad (6.23)$$

The *upper limit of nodes* is given as follow:

$$countNodes(ReachSet) \leq countNodes(Chaotic) + countNodes(Harmonic) \quad (6.24)$$

Harmonic Reach Set is ideal for *Non-controlled Airspace* missions, because it contains *Automatic Mode Switch* between *Navigation* and *Emergency Avoidance*.

6.4.6 ACAS-X like Reach set

Motivation: The implementation of *ACAS-Xu* behavior in DAA system will be mandatory for *National Airspace System Integration* in United spaces [8].

Implementation of ACAS-Xu like behaviour increase usability of approach, if it can be achieved without major concept changes.

Background: The *ACAS-Xu* system on operational level has been described in [9]. The *Policy for Collision Avoidance* proposal has been given in [10].

Some behavioural patterns can be encoded into *Reach Set*. ACAS-Xu navigation part is basically *Look-up table of Maneuvers for Allowed Separations*.

The *Evasive Maneuver* selection process in ACAS-Xu is similar to our approach: *Select most energy efficient maneuver in compliance with space-time constraints*. ACAS-Xu intruder model is similar to our *Body Volume Intersection Model* (sec. **??**). The *ACAS-Xu* defines following base separations:

1. *Horizontal* - movements on *Horizontal Plane* in *Global Coordinate System*.

2. *Vertical* - movements on *Vertical Plane* in *Global Coordinate System*.

There are allowed custom separations which can be used, for further experimentation:

1. *Slash* - movement on $+45°$ *Tilted Plane to Horizontal Plane* in *Global Coordinate System*.

2. *Backslash* - movement on $-45°$ *Tilted Plane to Horizontal Plane* in *Global Coordinate System*.

For given *Movement Automaton* implementation (sec. **??**) the separations are given as follow:

$$
\begin{aligned}
Horizontal &= \{Straight, Left, Right\} \\
Vertical &= \{Straight, Up, Down\} \\
Slash &= \{Straight, UpLeft, DownRight\} \\
Backslash &= \{Straight, UpRight, DownLeft\}
\end{aligned}
\tag{6.25}
$$

For each $Node(\ldots, buffer)$ and each *separation* there is a evaluation function *isSeparation* which decides, if *Trajectory* defined by node buffer is made up only from *Separation* movements. The function $isSeparation(\ldots)$ is defined like:

$$
isSeparation(buffer, separation) = \begin{cases} \forall movement \in buffer, \\ \quad movement \in separation & : true \\ otherwise & : false \end{cases}
\tag{6.26}
$$

Following *Separation Modes* can be defined with given *separations*:

1. *Horizontal* (ACAS-X defined mode) containing *horizontal* separation.

2. *Vertical* (ACAS-X defined mode) containing *vertical* separation.

3. *Horizontal-Vertical* (ACAS-X defined mode) containing *horizontal, vertical* separations.

4. *Full* (custom defined mode) containing all *Separation Modes*.

Note. Every separation modes generates 2D trajectories set on *Respective plane*. There is no need for *Tuning parameters* for further *Expansion Constraint*.

Expansion Constraint Function Implementation (alg. 6.5) is based on simple principle: *Select only candidate Nodes which Trajectories have at least one desired Separation Mode.*

Step: Initialization sets candidate *Nodes* as empty set, leftover *Nodes* as empty set, and, select all nodes form *stack* which represents *Finishing Trajectories* in working $cell_{i,j,k}$,

Step: Candidate Selection Process is evaluated for each *test Node* from *passing Node Set*.

For each *applicable separation*, given as input parameter *separations*, The test function *isSeparation* (eq. 6.26) is applied:

1. If *test Node* trajectory belongs to at least one allowed separation it is added to candidates set.

2. Else is added to *Leftovers*.

Note. Separation sets (eq. 6.25) are not *exclusive sets* in *Movement Automaton* domain. One *Trajectory* contained by Node can belong to multiple *Separations*.

Algorithm 6.5: Expansion Constraint function for *ACAS-like Reach Set Approximation*

Input	: Node[] stack, Cell cell$_{i,j,k}$, Separation[] separations
Tuning Parameters:	*None* : ∅
Output	: Node[] candidates, Node[] leftovers

Initialize structures;
Node[] candidates = [], Node[] leftovers=[];
Node[] passing = cell$_{i,j,k}$.getFinishingTrajectories(stack);

# Select nodes containing trajectories with usable separations;
for *Node test ∈ passing* **do**
  for *separation ∈ separations* **do**

  # Get separations for Node;
  Separations[] nodeSeparations = test.getSeparations();

  # If trajectory given by buffer is on Separation plane;
  if *isIn(isSeparation(test.buffer,separation)(6.26)* **then**
  | candidates.append(test);
  end
  end

  # If there was no applicable separation, throw Node away;
  if *test ∉ candidates* **then**
  | leftovers.append(test);
  end
end

Return results;
return *[candidates,leftovers]*

Example: for *Avoidance Grid* with *Distance 10 m, Layer count 10, Horizontal range* $[-45°, +45°]$, *Horizontal Cell Count 7, Vertical range* $[-30°, +30°]$, and *Vertical Cell Count 5*. Is given in (fig. 6.5). The UAS is at *Back-side* of *Figure* (initial state is at all *Trajectory Origins*). The *black dashed line* marks *Avoidance Grid* space boundary. Each trajectory has its own color and ends at *Front-side* of *Avoidance Grid Boundary*.

Full separation mode given in (fig. 6.5a). *Horizontal-Vertical* separation mode, used in original *ACAS-Xu* testing [9], given in (fig. 6.5b). *Horizontal* separation mode given in (fig. 6.5c) is usually used by planes. *Vertical* separation mode given in (fig. 6.5d) is usually used by copters.

(a) Full.

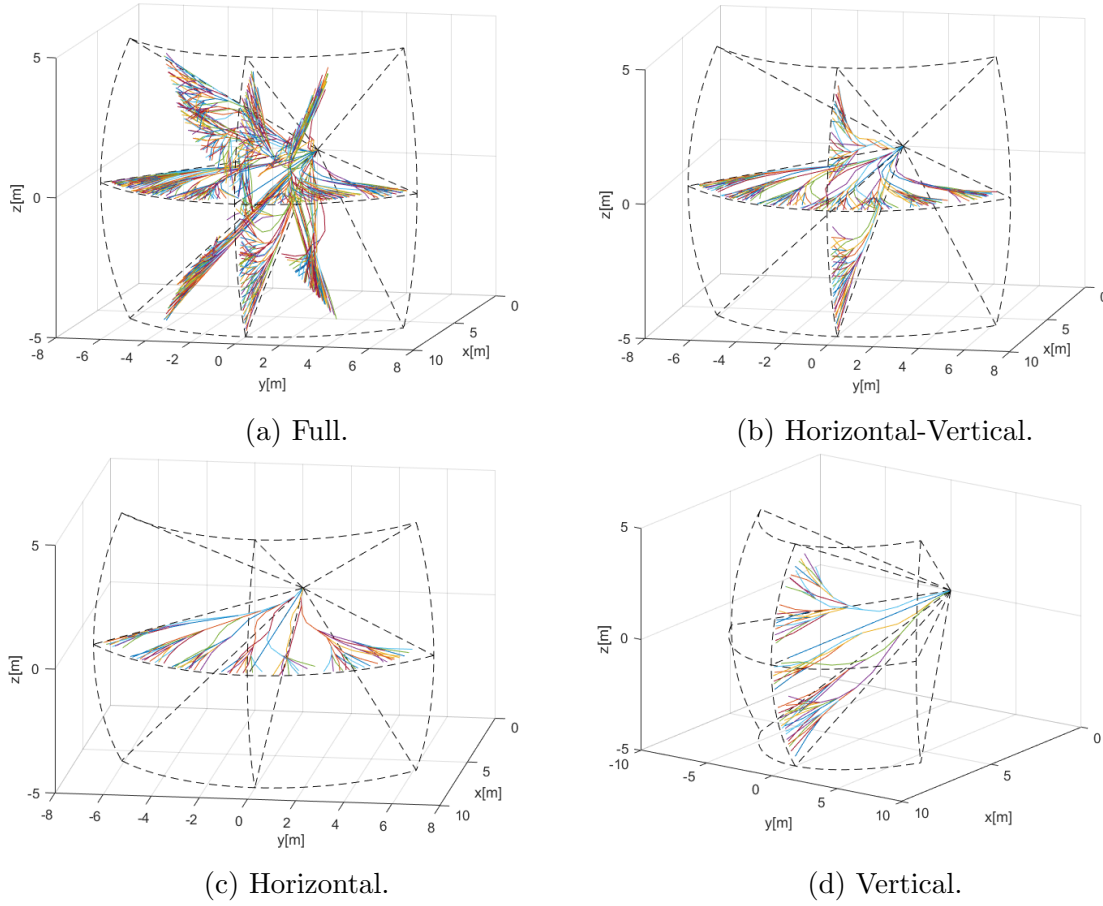(b) Horizontal-Vertical.

(c) Horizontal.

(d) Vertical.

Figure 6.5: ACAS-X imitation *reach set* approximation for various *separation modes*.

Pros and Cons: It can be seen from examples (fig. 6.5) that *ACAS-like Reach Set Approximation Method* (alg. 6.5) generates full reach set for 2D plane located in 3D space.

The *Reach Set* contains trajectories with *high coverage ratio* and *high smoothness rating* for selected 2D separation plane. Overall performance compared to full 3D reach sets (sec. 6.4.3, 6.4.4 6.4.5) is poor.

The *node* and *trajectory* count boundary was not implemented. It is common knowledge that *2D* avoidance sets does not require scaling [9]. Otherwise trajectory footprint mechanism like in *Harmonic Reach Set Approximation* (alg. 6.3) can be introduced.

This reach set implements *Planar-Separation* as native feature, it can be used for both *navigation* and *avoidance* tasks in *Controlled Airspace*. For *Non-controlled Airspace* there are far more superior *Combined Reach Set* (sec. 6.4.5).

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