6.7.4 Computation Complexity

Introduction: The *Computation Complexity* one mission control run assessment is necessary to identify the strong and weak points of approach. Lets get through modules to assess notable calculations/algorithms complexity on high abstraction level.

Navigation Loop: I the navigation loop, the waypoint reach condition (eq. ??) is checked, this is unitary operation with worst complexity O(1). The selection process of the next Goal Waypoint can get through all waypoints in the mission if they are all unreachable the complexity is O(|waypoints|).

The *notable steps* complexity is following:

Reach Condition: O(1)

Select Next Waypoint: O(|waypoints|)

Data Fusion: The data fusion is all about threat selection.

If UAS is in *controlled airspace* it needs to iterate over received *collision Cases* to select *active ones*. The complexity of this step is linear, therefore boundary is given as O(|collisionCases|).

Thresholding *Detected Obstacles* is done by simple comparison of LiDAR ray hits in given $cell_{i,i,k}$ of $Avoidance\ Grid$.

Any loading of threats from information sources is depending on clustering. The Airspace Clustering is considered as static for our setup. Therefore the count of active airspace clusters has main impact on complexity. The count of information sources is static and not changing over mission time. Information sources usually implement Hash search function with complexity $O \ln |searchedItemSet|$.

The computation complexity boundaries for Data fusion in our setup are following:

Select Active Collision Cases: O(|collisionCases|)

Threshold Detected Obstacles: O(|cells|)

Load Map Obstacles: $O(\ln|activeClusters| \times |informationSources|)$ Load Hard Constraints: $O(\ln|activeClusters| \times |informationSources|)$ Load Soft Constraints: $O(\ln|activeClusters| \times |informationSources|)$

Note. The real-time clustering is hard non-polynomial problem [1]. Usually all information sources and sensor have polynomial complexity of processing. The controlled airspace clusters are usually set for very long period of time. Therefore Obstacle Map, Airspace Constraints, and, Weather Constraints can be considered as preprocessed

Situation Assessment: The *Situation Assessment* is evaluating triggering events. The *evaluation* is usually simple existence question without further calculations. The *complexity* of *event evaluation* for our case is O(1). There are 8 triggers. The count of *triggers* needs to be accounted in complexity boundary:

$$O(|triggers| \times eventEvaluationComplexity)$$

Note. The trigger calculation complexity needs to stay low, because the triggers are verified every Mission Control Run. The Avoidance Run trigger frequency should be very low under normal conditions.

Avoidance Run: The *Avoidance run* is most critical part of *Mission Control Run*, because *Avoidance Path* calculation. The *Navigation Path* calculation is less complex (Rule engine is not accounted), therefore *Emergency Avoidance Mode* is assumed.

The threat insertion is realized in 7^{th} to 10^{th} step. The first is Avoidance Grid filled with Static Obstacles. The Avoidance Grid is designed to separate rotary LiDAR ray space into hit count even cells. Insertion of LiDAR scan into Avoidance Grid complexity depends on total cell count. The upper boundary for insert obstacles is given like follow:

Insert Obstacles:
$$O(|cells|)$$

The intruders intersection model type impact the insertion complexity. The linear intersection (sec. ??) is going through maximum of layers count cells.

The body volume intersection model (sec. ??) can check the simple intersection condition over all Avoidance Grid in worst case, therefore complexity for this check is bounded by count of cells.

The Maneuverability Uncertainty Intersection (sec. ??) can hit all cells in Avoidance Grid. The calculation complexity boundary is exponential depending on horizontal/vertical spread in [rad]. The intersection implementation was done ad-hoc. The impact of intersection application is visible only when there is more than 4 concurrence intruders (fig. ??).

The *complexity boundary for* intruder insertion is given like follow:

Note. The intruder intersection is critical in non-controlled airspace. The main complexity gain in controlled airspace is from rule application. Our rule complexity is in worst case depending on Reach Set node count and Active Collision Cases count.

Apply Our Rules: $O(|activeCollisionCases| \times |nodes|)$

For *Hard/Soft Constraints* The algorithm used for intersection polygons was selected based on study [2], the selected algorithm *Shamos-Hoey* [3]. The *calculation complexity* boundary is given like follow:

Hard Constraints Intersection:

$$O(|cells| \times |hardConstraints| \times \max |constraintPoints|^2)$$

Soft Constraints Intersection:

$$O(|cells| \times |softConstraints| \times \max |constraintPoints|^2)$$

Each threat category application in Mission Control Run is done after each intersection in 7^{th} to 10^{th} step. All ratings (tab. ??) expect Reachibility(cell_{ij,k}) and Reachibility(Trajectory) are calculated. The calculation complexity boundary for one reachibility rating is O(1). (eq. ??, ??). The Recalculate Reachibility operation applied $4\times$ have maximal complexity boundary given as follow:

Recalculate Reachibility:
$$O(4 \times (|nodes| + |cells|))$$

Each time at the end of in 7th to 10th step the *Avoidance Path is Selected*. The *Worst Case* (expected) scenario is to *select* four paths for each *treath* application. The algorithm for *best path selection* (alg. ??) iterates over all *cells* in avoidance grid and over all *trajectories* passing through that cell. The complexity boundary for *path selection* is given as follow:

Select Path:
$$O\left(4 \times \left(|cells| + \frac{|nodes|}{|cells|}\right)\right)$$

Conclusion: Overall approach complexity is *low*. If proper *Information Sources* with efficient clustering and *intersection models for intruders* are used, the approach will stay within *non-polynomial complexity*. The average load time for *testing scenarios* is summarized in (tab. 7.1).

Note. The calculation of Reach Set is eliminated by pre-calculation for state range [4].

7.5.3 Computation Footprint

The computation footprint is summarized in computation load (tab. 7.1). The computation load (eq. ??) was calculated for each time-frame in scenarios. There is summary of minimal, maximal, average and median values.

The computational load never exceed more than 55.95% in case of emergency Head On (eq. ??), which means that every path was calculated on time.

Scenario	Computation load			
Scenario	min.	max.	avg.	med.
Building avoidance (fig. ??)	2.20%	27.40%	12.11%	13.20%
Slalom (fig. ??)	12.20%	30.50%	21.42%	21.50%
Maze (fig. ??)	24.90%	46.10%	31.51%	30.80%
Storm (fig. ??)	2.60%	26.90%	11.57%	13.90%
Emergency Converging (fig. ??)	2.75%	16.50%	5.84%	4.95%
Emergency Head On (fig. ??)	3.90%	55.95%	13.19%	6.90%
Emergency Multiple (fig. ??)	5.90%	52.35%	12.77%	8.56%
Rule-based Converging (fig. ??)	3.60%	13.50%	7.32%	5.97%
Rule-based Head on (fig. ??)	4.65%	41.60%	13.64%	9.30%
Rule-based Multiple (fig. ??)	4.37%	23.30%	11.96%	10.93%
Rule-based Overtake (fig. ??)	3.85%	13.40%	7.62%	6.70%

Table 7.1: Computation load statistics for all test cases.

Following observations can be made:

- 1. Building avoidance, Slalom, and Maze scenarios the computation load is increasing with the amount of static obstacles. The average load for Emergency avoidance mode in clustered environment is 31.51% (Maze).
- 2. Storm scenario the overall computation load is very low due the moving constraint implementation (sec. ??).
- 3. Emergency Converging/Head On/Multiple scenarios the overall computation load is quite high due the ineffective body volume intersection (sec. ??) implementation.
- 4. Rule-based Converging/Head On/Multiple scenarios the median computational load is low, because of the linear rule implementation (sec. ??)
- 5. Rule-based Overtake the average computation load is very low, because only divergence/convergence (rule. ??) waypoints are calculated and UAS stays in navigation mode.

Bibliography

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- [4] Alojz Gomola, João Borges de Sousa, Fernando Lobo Pereira, and Pavel Klang. Obstacle avoidance framework based on reach sets. In *Iberian Robotics conference*, pages 768–779. Springer, 2017.