

### 6.5.1 Intruders

**Intruder behaviour:** *Adversarial behaviour* of moving obstacle is trying to destroy avoiding our UAS. The *Intruder* UAS [1] is not trying to hurt our *UAS* actively. The *Adversarial behaviour* is neglected in this work. The non-cooperative avoidance is assumed, it can be relaxed to *cooperative avoidance* in *UTM controlled airspace*.

**Intruder information:** The *observable intruder information set* for any kind of intruder, obtained through sensor/C2 line, is following:

1. *Position* - position of intruder in *local* or *global* coordinate frame, which can be transformed into *avoidance grid coordinate frame*.
2. *Heading and Velocity* - intruder heading and linear velocity in avoidance grid coordinate frame.
3. *Horizontal/Vertical Maneuver Uncertainty Spreads* - how much can an *intruder* deviate from *original linear path* in *horizontal/vertical* plane in *Global coordinate Frame*.

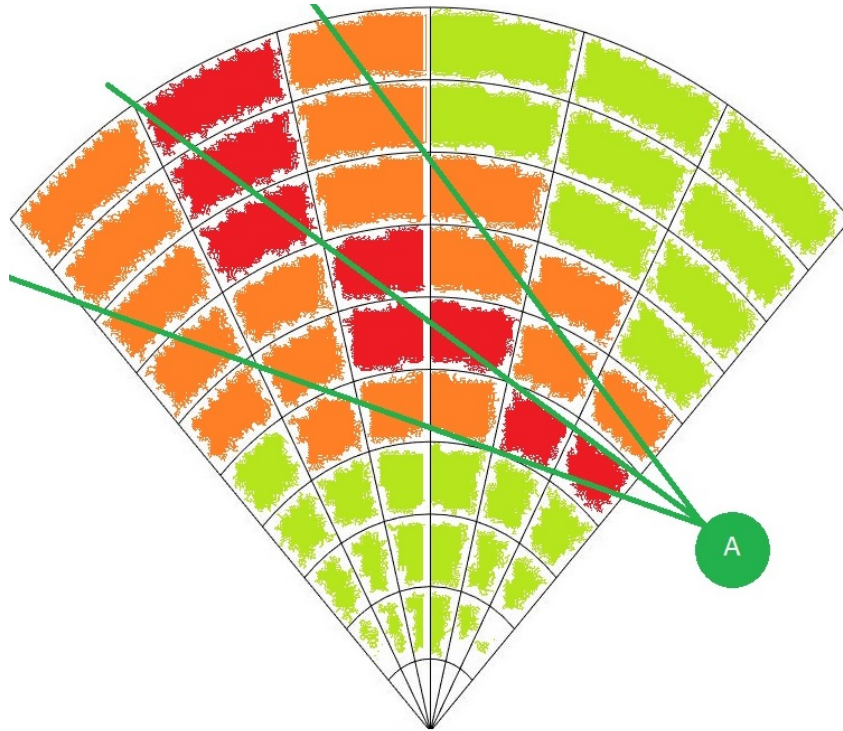


Figure 6.1: Intruder UAS intersection rate along expected trajectory.

**Example of Intruder Intersection:** Lets neglect the *time-impact* aspect on *intersection*. The *intruder* (black "I" circle) is intersecting one *avoidance grid horizontal slice* (fig. 6.1). The intruder is moving along linear path approximation based on velocity (middle green line). The *Horizontal Maneuver Uncertainty spread* is in *green line boundary area* intruder intersection rating is denoted as green-orange-red cell fill reflecting intersection

severity: red is high rate of intersection, orange is medium rate of intersection and green is low rate of intersection.

**Moving Threats:** The *UAS* can encounter following threats during the *mission execution*:

1. *Non-cooperative Intruders* - the intruders whom does not implement any approach to ensure mutual avoidance efficiency.
2. *Cooperative Intruders* - the intruders whom actively communicate or follow common agreed behaviour pattern (ex. Rules of the Air).
3. *Moving Constraints* - the constrained portion of *free* space which is shifting its boundary over time (ex. Short term bad weather).

*Note.* Our approach considers only *UAS* intruders, because *Data Fusion* considers data received through *ADS-B* messages. The *Intruders* extracted from *LiDAR* scan were not considered (ex. birds). The proposed *intruder intersection models* are reusable for other *intruder sources*.

**Approach Overview:** The *Avoidance Grid* (def. ??) is adapted to *LiDAR* sensor. The *euclidean grid intersections* are fairly simple. The *polar coordinates grid* are not. The need to keep *polar coordinates grid* is prevalent, because of fast *LiDAR* reading assessment. There are following commonly known methods to address this issue:

1. *Point-cloud Intersections* - the *threat impact area* is discredited into sufficiently thick point cloud. This point-cloud have *point impact rate* and *intersection time* assigned to each point. The *point-cloud* is projected to *Avoidance Grid*. If *impact point* hits  $cell_{i,j,k}$  the cell's impact rate is increased by amount of *point impact rate*. The final *threat impact rate* in  $cell_{i,j,k}$  is given when *all* points from point cloud are consumed. Close point problem [2] was solved by application of method [3].
2. *Polygon Intersections* - the *threat impact area* is modeled as polygon, each  $cell_{i,j,k}$  in *Avoidance Grid* is considered as *polygon*. There is a possibility to calculate cell space geometrical inclusive intersection. The *impact rate* is then given as rate between *intersection volume* and  $cell_{i,j,k}$  volume. The algorithm used for intersection selected based on:[4] the selected algorithm *Shamos-Hoey* [5].

*Note.* The *Intruder Intersection* models are based on *analytically geometry* for *cones* and *ellipsoids* taken from [6].

**Intruder Behaviour Prediction:** *Intruder Intersection Models* is about space-time intersection of *intruder body* with *avoidance Grid* and *Reach Set*:

1. The *UAS* reach set defines *time boundaries* to *enter/leave* cell in avoidance grid.

2. The *Intruder* behavioral pattern defines *rate of space intersection* with cell bounded space in avoidance grid.

The multiplication of *space intersection rate* and *time intersection rate* will give us *intruder intersection rate* for our *UAS* and intruder.

**Intruder Dynamic Model:** The definition of avoidance grid enforces the most of these methods to be numeric. Let us introduce intruder dynamic model:

$$\begin{aligned} \text{dposition}/\text{dtime} = \text{velocity} \quad | \quad & \begin{aligned} \text{position}_x(t) &= \text{position}_x(0) + \text{velocity}_x \times t \\ \text{position}_y(t) &= \text{position}_y(0) + \text{velocity}_y \times t \\ \text{position}_z(t) &= \text{position}_z(0) + \text{velocity}_z \times t \end{aligned} \end{aligned} \quad (6.1)$$

Position vector in euclidean coordinates  $[x, y, z]$  is transformed into *Avoidance Grid* coordinate frame. Velocity vector for  $[x, y, z]$  is *estimated and not changing*. The time is in interval  $[\text{entry}, \text{leave}]$ , where *entry* is intruder entry time into avoidance grid and *leave* is intruder leave time from avoidance grid.

*Note.* If *intruder* is considered, time of entry is marked as  $\text{intruder}_{\text{entry},k}$  where  $k$  is intruder identification, time of leave is marked as  $\text{intruder}_{\text{leave},k}$  where  $k$  is intruder identification.

**Cell Entry and Leave Times**  $UAS_{\text{entry}}(\text{cell}_{i,j,k})$  and  $UAS_{\text{leave}}(\text{cell}_{i,j,k})$  are depending on intersecting *Trajectories* and *bounded cell space* (eq. ??). There is *Trajectory Intersection* function from (def. ??) which evaluates *Trajectory segment* entry and leave time.

The *UAS Cell Entry* time is given as minimum of all *passing trajectory segments* entry times (eq. 6.2), if there is no *passing trajectories* the *UAS entry time* is set to 0.

$$UAS_{\text{entry}}(\text{cell}_{i,j,k}) = \min \left\{ \begin{aligned} &0, \text{entry}(\text{Trajectory}, \text{cell}_{i,j,k}) : \\ &\text{Trajectory} \in \text{PassingTrajectories} \end{aligned} \right\} \quad (6.2)$$

The *UAS Cell Leave* time is given as maximum of all *passing trajectory segments* entry times (eq. 6.3), if there is no *passing trajectories* the *UAS leave time* is set to 0.

$$UAS_{\text{leave}}(\text{cell}_{i,j,k}) = \max \left\{ \begin{aligned} &0, \text{leave}(\text{Trajectory}, \text{cell}_{i,j,k}) : \\ &\text{Trajectory} \in \text{PassingTrajectories} \end{aligned} \right\} \quad (6.3)$$

**Time Intersection Rate:** The key idea is to calculate how long the *UAS* and *Intruder* spends together in same space portion ( $\text{cell}_{i,j,k}$ ). The *Intruder* can spent some time in  $\text{cell}_{i,j,k}$  bounded by interval of *intruder* entry/leave time.

The *UAS* can spent some time, depending on *selected trajectory* from *Reach Set*. The time spent by UAS is bounded by entry (eq. 6.2) and leave (eq. 6.3).

The intersection duration of these two intervals creates *time intersection rate* numerator, the *maximal duration* of *UAS* stay gives us *denominator*. The *time intersection rate* is formally defined in (eq. 6.4).

$$\text{time} \left( \begin{array}{c} UAS, \\ Intruder, \\ cell_{i,j,k} = \circ \end{array} \right) = \frac{\left| \begin{array}{c} [intruder_{entry}(\circ), intruder_{leave}(\circ)] \\ \cap \\ [UAS_{entry}(\circ), UAS_{leave}(\circ)] \end{array} \right|}{|[UAS_{entry}(\circ), UAS_{leave}(cell_{\circ})]|} \quad (6.4)$$

**Intruder Intersection Rate:** The *Intruder Intersection Rate* (eq. 6.5) is calculated as *multiplication* of *space intersection rate* (defined later) and *time intersection rate* (eq. 6.4).

$$\text{intruder} \left( \begin{array}{c} UAS, \\ Intruder, \\ cell_{i,j,k} \end{array} \right) = \text{time} \left( \begin{array}{c} UAS, \\ Intruder, \\ cell_{i,j,k} \end{array} \right) \times \text{space} \left( \begin{array}{c} UAS, \\ Intruder, \\ cell_{i,j,k} \end{array} \right) \quad (6.5)$$

*Note.* If there is no information to derive *Intruder* entry/leave time for cells the *time intersection rate* is considered 1.

The *Intruder cell reach* time (eq. 6.6) is bounded to discrete point in intersection model [2, 3]. The intruder *entry/leave time* is calculated similar to *UAS cell entry* (eq. 6.2)/*leave* (eq. 6.3) *time*.

$$\text{pointReachTime}(Intruder, point) = \frac{\text{distance}(Intruder.initialPosition, point)}{|Intruder.velocity|} \quad (6.6)$$

**Space Intersection Rate:** The *Space Intersection Rate* reflects probability of *Intruder* intersection with portion of space bounded by  $cell_{i,j,k}$ , to be precise with intruder trajectory or vehicle body shifted along the trajectory. The principles for *space intersection rate* calculation are following:

1. *Line trajectory* - *intruder* trajectory is given by linear approximation (eq. 6.1), depending on *intruder size* the intersection with avoidance grid can be:
  - a. *Simple line* - intersection is going along the trajectory line defined by intruder model (eq.6.1).
  - b. *Volume line* - intersection is going along the trajectory line defined by intruder model (eq. 6.1) and intruder's *body radius* is considered in intersection.

2. *Elliptic cone* - initial position is considered as the top of a cone, the main cone axis is defined by intruder linear trajectory (eq. 6.1)  $time \in [0, \infty]$ . The cone width is set by horizontal and vertical spread.

**Moving Constraints:** The basic ideas is the same as in case *static constraints* (sec. ??). There is horizontal constraint and altitude constraint outlining the constrained space. The only additional concept is moving of *constraint* on horizontal plane in global coordinate system.

The constraint intersection with *avoidance grid* is done in *fixed decision Time*, for cell in *fixed cell leave time* (eq. 6.3), which means concept from static obstacles can be fully reused.

**Definition:** The *moving constraint definition* (eq. 6.7) covers minimal data scope for moving constraint, assuming linear constraint movement.

**Definition 1. Moving Constraints** The original definition (eq. ??) is enhanced with additional parameters to support constraint moving:

1. Velocity - *velocity vector on 2D horizontal plane*.
2. Detection time - *the time when constraint was created/detected, this is the time when center and boundary points position were valid*.

$$\begin{aligned} constraint = \{ & position, boundary, \dots \\ & \dots, velocity, detectionTime, \dots \\ & \dots altitude_{start}, altitude_{end}, safetyMargin \} \end{aligned} \quad (6.7)$$

**Cell Intersection:** The *intersection algorithm* follows (eq. ??), only shift of the *center and boundary points* is required.

First let us introduce  $\Delta time$  (eq. 6.8), which represents difference between the constraint detection time and expected cell leave time (eq. 6.3).

$$\Delta time = UAS_{leave}(cell_{i,j,k}) - detectionTime \quad (6.8)$$

The constraint boundary is shifted to:

$$\begin{aligned} shiftedBoundary(constraint) = \{ & newPoint = point + velocity \times \Delta time : \dots \\ & \dots \forall point \in constraint.boundary \} \end{aligned} \quad (6.9)$$

The constraint center is shifted to:

$$shiftedCenter(constraint) = constraint.center + velocity \quad (6.10)$$

*Note.* The  $\Delta time$  is calculated separately for each  $cell_{i,j,k}$ , because *UAS* is also moving and reaching cells in different times. The *cell leave time* can be calculated in advance after reach set approximation.

**Alternative Intersection Implementation:** The alternative used for intersection selected based on polygon intersection algorithms review [4], the selected algorithm is *Shamos-Hoey* [5].

The implementation was tested on *Storm scenario* (sec. ??) and it yelds same results.

# Bibliography

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