

Data Structures and Algorithms

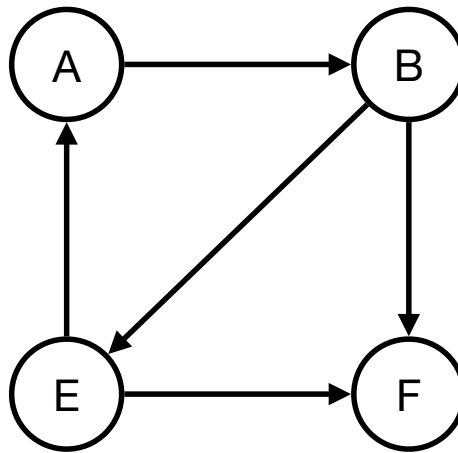
Week 10 - More on SCCs, Topological Sort

Subodh Sharma and Rahul Garg
{svs,rahulgarg}@iitd.ac.in.

Strongly Connected Components

- A graph is said to be strongly connected - If every vertex is reachable from every other vertex
- The binary relation of being strongly connected is an **equivalence relation**
 - That is it is reflexive, symmetric and transitive
- Strongly connected component of a directed graph G is also **maximal**
- **Used in Abstractions!** SCCs in a graph can be **condensed** into single vertices leading to the formation of a **DAG**

Strongly Connected Components



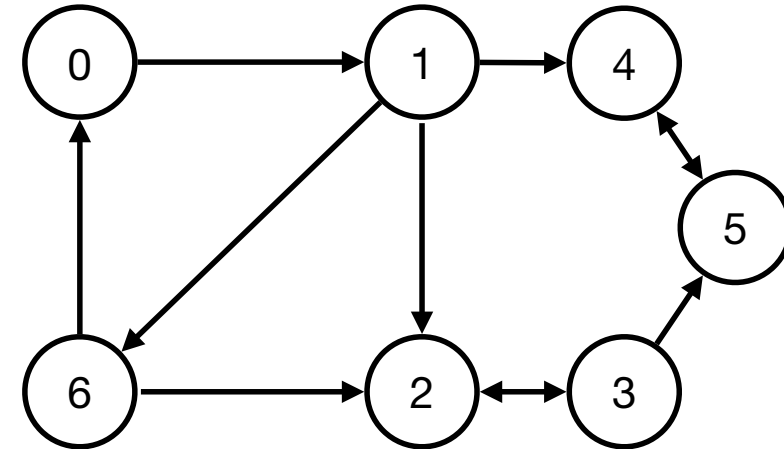
- SCC: $(\{A, B, E\}, \{\} \dots \{\})$
- Use of DFS to find SCCs — Robert Tarjan 1972 (also discovered Splay and Fibonacci Heaps)

Strongly Connected Components

Tarjan's Algorithm

- **Key Observations:**

- **Input:** Directed Graph G
- **Output:** subgraph with vertices of SCC
- Each vertex appears in exactly one SCC of the graph
- Use of DFS + idea of **low-link values**

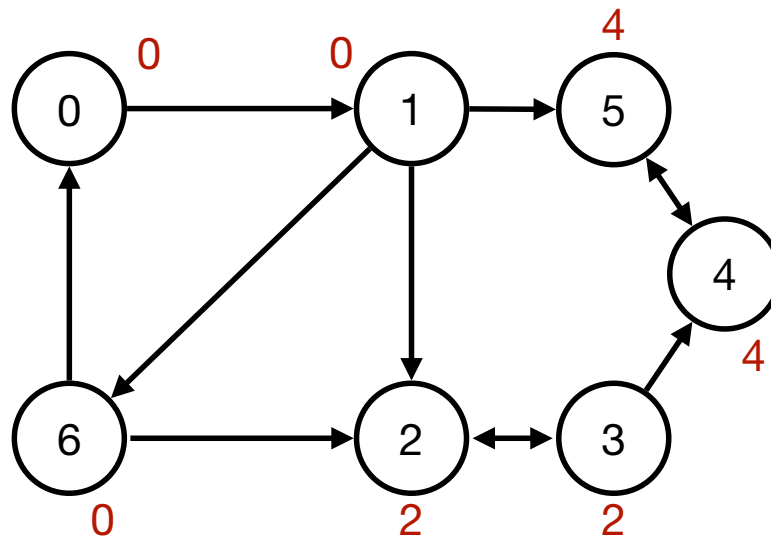


Strongly Connected Components

Tarjan's Algorithm

- **Low-link values:** An LL value of a node is the **smallest** node id reachable from that node (including itself).

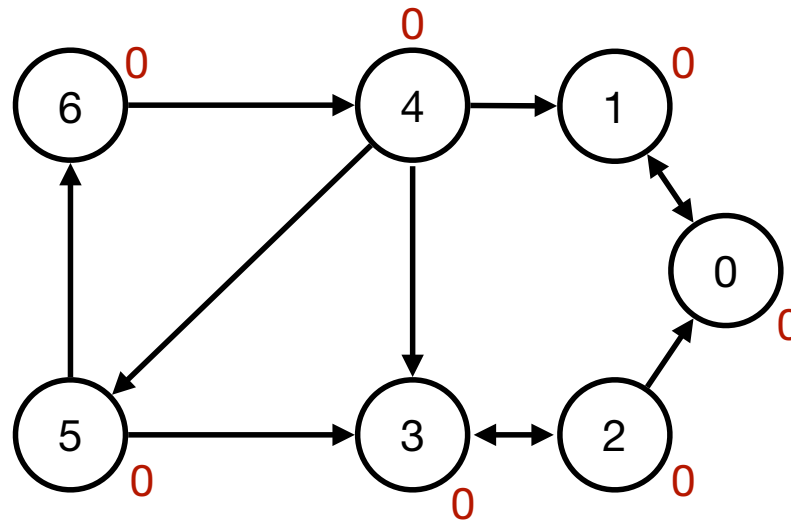
- **Time Complexity:** $O(V \cdot (V + E))$



Strongly Connected Components

Tarjan's Algorithm

- **Low-link values:** An LL value of a node is the **smallest** node id reachable from that node (including itself).
- **CAUTION:** LL values are dependent in the order of exploration

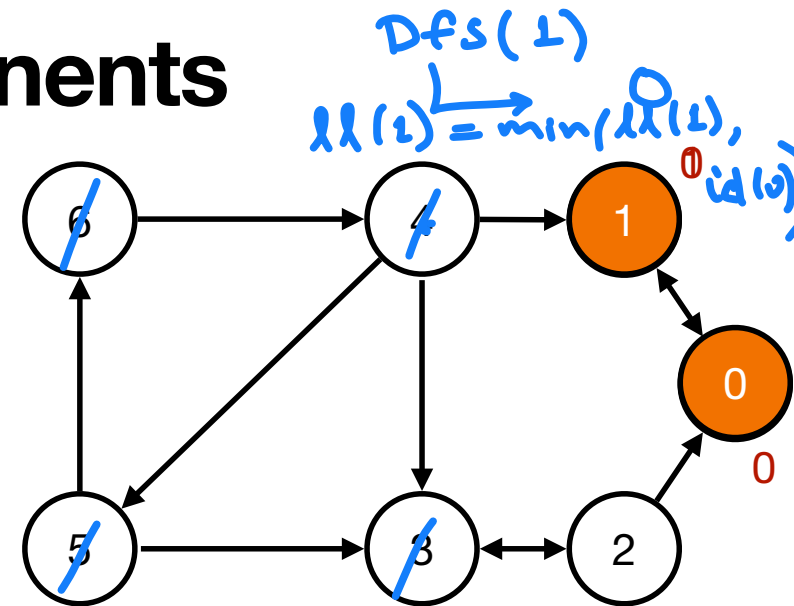


- **Incorrect SCC was computed**

Strongly Connected Components

Tarjan's Algorithm

- Algorithm:
 - Start DFS from a node
 - Upon visiting a node assign it a unique integer id and an LL value
 - Mark the node visited and then to the stack of seen nodes
 - On DFS backtrack, if the **next node** is on the stack update the LL value of the **current node** to the **minimum of the current node's and next node's LL value**
 - Allows LL values to propagate through cycles
 - If all nodes are visited and the current node starts an SCC then pop nodes of the stack until the **current** node

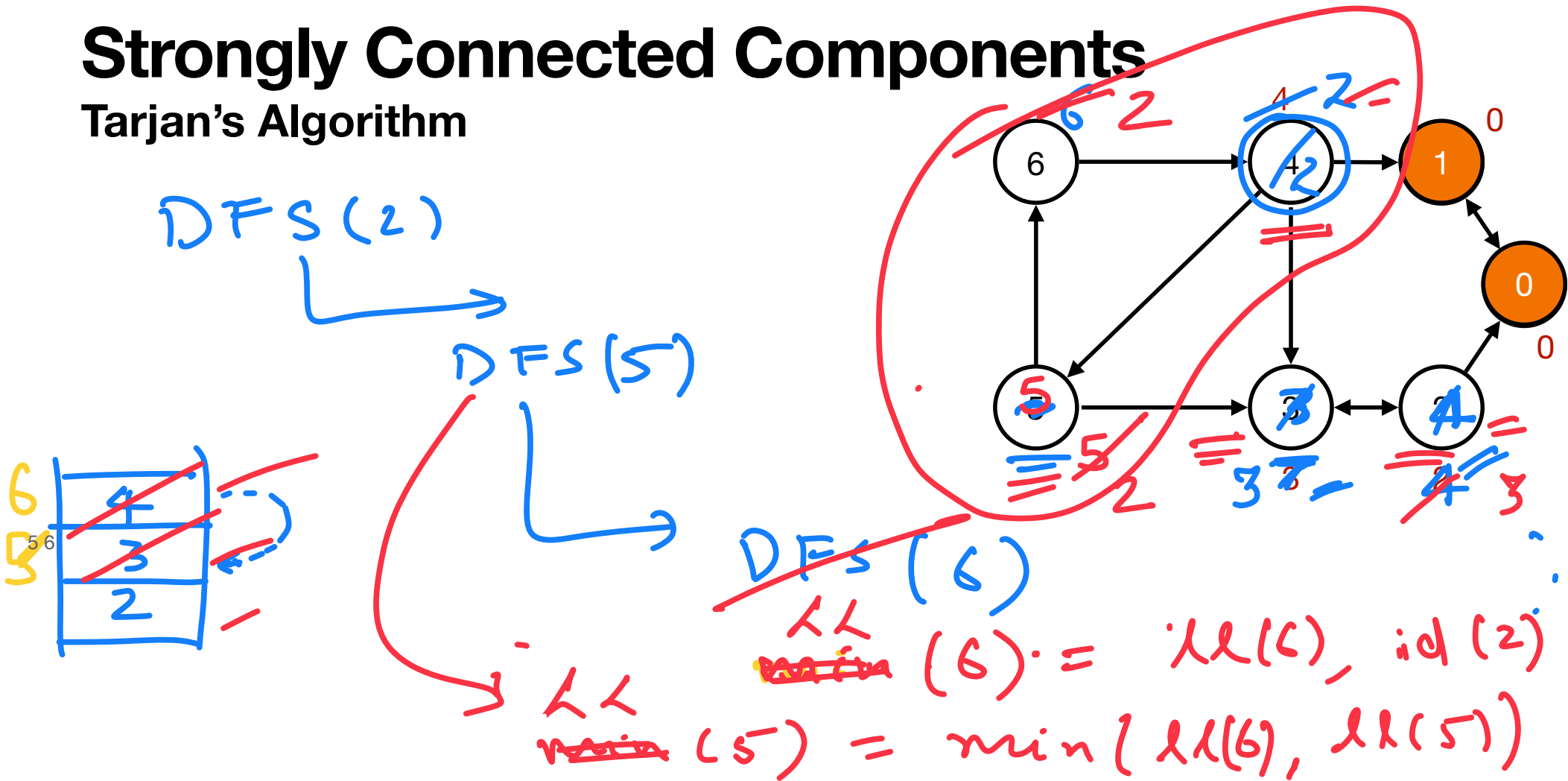


$low[1] = \min(low[1], ids[0])$

```
for (int w : adj[v]) {
    if (ids[w] == -1) { // not visited yet
        dfs(w);
        low[v] = min(low[v], low[w]);
    } else if (onstack[w]) {
        low[v] = min(low[v], ids[w]);
    }
}
```

Strongly Connected Components

Tarjan's Algorithm



Strongly Connected Components

Tarjan's Algorithm

- **Invariant of Tarjan's Alg:** A node remains on the stack **iff** there exists a path from it to a node on the stack
 - Prevents the LL values of different SCCs from interfering with each other