Kristopher Blair Normalization Assignment

#1

Consider the following relations for an order-processing application database for Acme Products, Wile E. Coyote, owner. Determine whether the relations are in BCNF and, if not, decompose them.

```
Order(<u>Order#</u>, Order_date, Customer#, Total_amount)
Order_Item(<u>Order#</u>, <u>Item#</u>, Quantity_ordered, Price_each, Total_price,
Discount%)
```

Assume that each item has a different discount. The Price_each refers to one item, Total_price is the cost per item multiplied by the number of items purchased, Order_date is the date on which the order was placed, and the Total amount is the amount of the order.

ANSWER:

The relations are not in BCNF. There is a partial dependency in the Order Item relation.

- Order(Order#, Order date, Customer#, Total amount)
- Order Item(Order#, Item#, Quantity ordered, Total price)
- Item(<u>Item#</u>, Price each, Discount)

This shows the correct BCNF relation for this data.

#2

```
Determine whether the following relation is in BCNF and, if not, decompose it.

Rx(<u>Doctor#</u>, <u>Patient#</u>, <u>Date</u>, Diagnosis, Treatment_code, Charge)
```

A tuple describes a visit of a patient to a doctor along with a treatment code and a charge. Assume that each treatment code has a fixed charge (regardless of patient).

ANSWER:

The relation shown is not in BCNF. It should be as follows:

- Rx(<u>Doctor#</u>, <u>Patient#</u>, <u>Date</u>, Diagnosis, Treatment_code)
- Doc Treatments(<u>Doctor#</u>, <u>Treatment code</u>, Charge)

#3

Consider the relation R (A, B, C, D, E, F, G, H, I, J) and the set of functional dependencies {AB \rightarrow C, A \rightarrow DE, B \rightarrow F, F \rightarrow GH, D \rightarrow IJ}. What is the key for R? Decompose R into 2NF and then 3NF. Repeat the decomposition into 2NF and then 3NF starting with the original version of R and using the following set of functional dependencies instead: {AB \rightarrow C, BD \rightarrow EF, AD \rightarrow GH, A \rightarrow I, H \rightarrow J}.

ANSWER:

The key for R is AB. Decomposing gives you the following:

- R(A, B, C)
- R $A(\underline{A}, D, E)$
- $R_B(\underline{B}, F, G, H)$
- R D(\underline{D} , I, J)

Using the second set of functional dependencies, it decomposes to the following:

- $R(\underline{A}, \underline{B}, C)$
- R BD(\underline{B} , \underline{D} , E, F)
- R AD(\underline{A} , \underline{D} , G, H)
- R $A(\underline{A}, I)$
- R $H(\underline{H}, J)$

#4

Consider the relation DISK_DRIVE (Serial_number, Manufacturer, Model, Batch, Capacity, Retailer). Each tuple in the relation DISK_DRIVE contains information about a disk drive with a unique serial number, made by a manufacturer, with a particular model number, released in a certain batch, which has a certain storage capacity, and is sold by a certain retailer. For example, the tuple

```
Disk drive('1978619', 'Acme Drives', 'A2235X', '765324', 500, 'CompuMax')
```

specifies that Acme Drives made a 500GB disk drive with serial number 1978619 and model number A2235X; it was released in batch 765324 and is sold by CompuMax.

Translate each of the following into a FD:

- 1. The manufacturer and serial number uniquely identifies the drive.
- 2. A model number is registered by a manufacturer and therefore can not be used by another manufacturer.
- 3. All disk drives in a particular batch are the same model.
- 4. All disk drives of a certain model of a particular manufacturer have exactly the same capacity.

After stating the above FDs, decompose DISK DRIVE into 3NF.

ANSWERS:

FDs

- 1. Serial number, Manufacturer → Model, Batch, Capacity, Retailer
- 2. Manufacturer \rightarrow Model
- 3. Batch \rightarrow Model
- 4. Model → Capacity

Decomposition into 3NF:

- DISK DRIVE(Serial number, Manufacturer, Batch, Retailer)
- Manufacturer Models(Manufacturer, Model)
- Batches(Batch, Model)
- Models(Model, Capacity)