OS HW2

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- 文章位址: https://hackmd.io/@lohsuan/rkdDdjdNp
- 題目敘述: Assignment 2: Scheduling Policy Demonstration Program
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Source Part

Github: Source Code

```
sche test outcome
sudo ./sched test.sh ./sched demo ./sched demo 312551105
Running testcase 1: ./sched_demo -n 1 -t 0.5 -s NORMAL -p -1 ......
Result: Success!
Running testcase 2: ./sched_demo -n 2 -t 0.5 -s FIF0,FIF0 -p 10,20 ......
Result: Success!
Running testcase 3: ./sched_demo -n 3 -t 1.0 -s NORMAL,FIF0,FIF0 -p -1,10,30 ......
Result: Success!
```

1. Describe how you implemented the program in detail.

Parse program arguments

Report Part

使用 getopt parse 傳入的 argv, getopt 會根據 "n:t:s:p:" optstring 制定的規則做 parse。

```
while ((opt = getopt(argc, argv, "n:t:s:p:")) != -1) {
    switch (opt) {
   case 'n':
       if (atoi(optarg) != num_threads) {
            printf("Please put -n argument as the first passing value\n");
            printf("Usage: sudo ./sche -n 1 -t 0.5 -s NORMAL -p -1\n");
            exit(EXIT_FAILURE);
       break;
    case 't':
        time_wait = atof(optarg);
       break;
    case 's':
        char *s_string = optarg;
        char *s_token = strtok(s_string, ",");
        for (int i = 0; i < num_threads; i++) {</pre>
            policies[i] = s_token;
            s_token = strtok(NULL, ",");
       break;
    case 'p':
        char *p_string = optarg;
        char *p_token = strtok(p_string, ",");
        for (int i = 0; i < num_threads; i++) {</pre>
            priorities[i] = atoi(p_token);
            p_token = strtok(NULL, ",");
       break;
    default:
        fprintf(stderr, "input usage error\n");
        exit(EXIT_FAILURE);
```

下面的程式宣告 cread thread 會用到的 attribute,每個變數的作用以註解展示。

Create < num threads > worker thread info

```
struct sched_param param[num_threads]; // thread parameters
 pthread_t thread[num_threads];
                                 // thread identifiers
 thread_info_t thread_info[num_threads]; // thread information
Set CPU affinity
```

• 將所有 thread 設定跑在相同 CPU (ex: cpu_id=0),細節在註解中。

- (若無加上 #define _GNU_SOURCE 在設定affinity 時會失敗)
- #define _GNU_SOURCE

```
int cpu_id = 0;
 cpu_set_t cpuset;
 CPU_ZERO(&cpuset);
                      // clear cpuset
 CPU_SET(cpu_id, &cpuset); // set CPU 0 on cpuset
 sched_setaffinity(0, sizeof(cpuset), &cpuset);
Set the attributes to each thread
```

• 根據上面 parse 的 parameter 設定 thread_info 中各個 field 的值。 • 並根據 policy 是 FIFO 或 NORMAL 決定 pthread_create 需不需要設定並傳入前面宣告

- 之attr。 • 設定 attr 之步驟細節在註解中說明。
- pthread_create(&thread[i], &attr[i], thread_func, &thread_info[i]):
- 。 4th parameter: &thread_info[i] 是傳入 thread_func 的參數

。 3rd parameter: thread_func 是讓此 thread 執行的程式

for (int i = 0; i < num_threads; i++) {</pre> thread_info[i].thread_id = i; thread_info[i].time_wait = time_wait; if (strcmp(policies[i], "FIFO") == 0) { thread_info[i].sched_policy = SCHE_FIFO; // initialize thread attributes pthread_attr_init(&attr[i]); // set the scheduling inheritance to explicit pthread_attr_setinheritsched(&attr[i], PTHREAD_EXPLICIT_SCHED); // set the scheduling policy - FIFO if (pthread_attr_setschedpolicy(&attr[i], SCHE_FIF0) != 0) { printf("Error: pthread_attr_setschedpolicy\n"); exit(EXIT_FAILURE); // set the priority param[i].sched_priority = priorities[i]; // set priority // set the scheduling paramters if (pthread_attr_setschedparam(&attr[i], ¶m[i]) != 0) { printf("Error: pthread_attr_setschedparam\n"); exit(EXIT_FAILURE); } // create the thread if (pthread_create(&thread[i], &attr[i], thread_func, &thread_info[i]) != 0) { printf("Error: FIFO pthread_create\n"); exit(EXIT_FAILURE); } else if (strcmp(policies[i], "NORMAL") == 0) { thread_info[i].sched_policy = SCHE_NORMAL; pthread_create(&thread[i], NULL, thread_func, &thread_info[i]); } else {

• 將 pthread_barrier_init 至於 thread_func 最開始,等到所有thread ready 才繼續執行 task

Start all thread_info at once

thread and main thread)

printf("Unexpected input \n");

pthread_barrier_t barrier; // initialize barrier threads + main thread = num_threads + 1 pthread_barrier_init(&barrier, NULL, num_threads + 1); pthread_barrier_wait(&barrier);

• 利用 | pthread_barrier_init | 設定需要等待多少 thread , 數量為num_threads + 1 (num of

pthread_barrier_destroy(&barrier);

```
• pthread_barrier_*() 把先後到達的 threads 擋在同一欄桿前,直到所有 thread 到齊才撤
     下欄桿同時放行。
   • Ref: https://www.cnblogs.com/liyulong1982/p/5480678.html
Worker thread: Do the task
```

• clock_gettime 可以取得 wall-clock time 或程式的 CPU time, 其所傳回的時間是 timespec

• 透過每次更新 end_time 並計算與 start_time 的差值取得 thread 現在已使用多長時間

• 當 thread 執行的時間超過了 start_time - end_time 時就會 break ,也就是等待了 busy time 的時間

while (1) {

MICRO_6;

- for (int i = 0; i < 3; i++) { printf("Thread %d is running\n", (int)thread_info->thread_id);
- /* Busy for <time_wait> seconds */ struct timespec start_time, end_time; clock_gettime(CLOCK_THREAD_CPUTIME_ID, &start_time); int start_time_msec = start_time.tv_sec * MILLI_3 + start_time.tv_nsec /

clock_gettime(CLOCK_THREAD_CPUTIME_ID, &end_time);

int end_time_msec = end_time.tv_sec * MILLI_3 + end_time.tv_nsec / MICRO_6; if ((end_time_msec - start_time_msec) >= busy_time_msec) { break; sched_yield(); // yield the CPU Wait for all thread_info to finish 使用 pthread_join 等待所有 thread ter,inate for (int i = 0; i < num_threads; i++) {</pre> pthread_join(thread[i], NULL);

this ensure the real-time threads can fully utilize the CPU resource without being preempted by fair-time threads

Compile, Run, and Test

test with TA code

Ensure real-time threads can fully utilize the CPU

• Set sysctl -w kernel.sched_rt_runtime_us=1000000 in terminal

\$ sudo ./sched_test.sh ./sched_demo ./sched_demo_312551105

compile \$ gcc sched_demo_312551105.c -lpthread -o sched_demo_312551105 # run (run with "sudo" to have previlege to create_thread) \$ sudo ./sched_demo_312551105 -n 3 -t 1.0 -s NORMAL,FIF0,FIF0 -p -1,10,30 \$ sudo ./sched_demo_312551105 -n 4 -t 0.5 -s NORMAL, FIFO, NORMAL, FIFO -p -1, 10, -1, 30

2. Describe the results of ./sched_demo -n 3 -t 1.0 -s NORMAL, FIFO, FIFO -p -1, 10, 30 and what causes that

Thread 1 is running Thread 1 is running Thread 0 is running Thread 0 is running

Thread 3 is running Thread 3 is running

struct

- FIFO (thread 1, 2) 因為是 real time policy 比 NORMAL (thread 0) 有更高的優先權,且 thread 2比 thread 1的 priority 要大 • 因此先執行 thread 2 接著 thread 1 最後是 NORMAL 的 thread 0 Scheduling Policy git:(main) make run sudo ./sched demo 312551105 -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30 Thread 2 is running Thread 2 is running Thread 2 is running Thread 1 is running
- Thread 0 is running 3. Describe the results of ./sched_demo -n 4 -t 0.5 -s NORMAL, FIFO, NORMAL, FIFO -p -1, 10, -1, 30, and what causes that
- FIFO (thread 1, 3) 因為是 real time policy 比 NORMAL (thread 0, 2) 有更高的優先權,且 thread 3 比 thread 1 的 priority 要大 • 因此先執行 thread 3 接著換 thread 1, 最後 thread 0, 2 因沒有 priority 優先所以交替進行。

sudo ./sched demo 312551105 -n 4 -t 0.5 -s NORMAL, FIFO, NORMAL, FIFO -p -1, 10, -1, 30

- Thread 3 is running Thread 1 is running Thread 1 is running Thread 1 is running Thread 2 is running Thread 0 is running Thread 2 is running Thread 0 is running Thread 2 is running Thread 0 is running
- 4. Describe how did you implement n-second-busy-waiting? • clock_gettime 可以取得 wall-clock time 或程式的 CPU time, 其所傳回的時間是 timespec
- 當 thread 執行的時間超過了 start_time end_time 時就會 break ,也就是等待了 busy time 的時間

• 透過每次更新 end_time 並計算與 start_time 的差值取得 thread 現在已使用多長時間

```
for (int i = 0; i < 3; i++) {
    printf("Thread %d is running\n", (int)thread_info->thread_id);
    /* Busy for <time_wait> seconds */
    struct timespec start_time, end_time;
    clock_gettime(CLOCK_THREAD_CPUTIME_ID, &start_time);
    int start_time_msec = start_time.tv_sec * MILLI_3 + start_time.tv_nsec /
MICRO_6;
    while (1) {
        clock_gettime(CLOCK_THREAD_CPUTIME_ID, &end_time);
        int end_time_msec = end_time.tv_sec * MILLI_3 + end_time.tv_nsec /
MICRO_6;
        if ((end_time_msec - start_time_msec) >= busy_time_msec) {
            break;
   sched_yield(); // yield the CPU
```

```
CLOCK_THREAD_CPUTIME_ID: 程式單一 thread 所耗費的時間。
 struct timespec {
  time_t tv_sec; /* seconds */
  long tv_nsec; /* nanoseconds */
```