

An example to illustrate the computation of MSNE

Theorem: A strategy profile $(\sigma_1^*, \dots, \sigma_n^*)$ is a MSNE iff $\forall i \in N$,

- (1) $u_i(\delta_i, \sigma_i^*)$ is the same $\forall \delta_i \in \delta(\sigma_i^*)$
- (2) $u_i(\delta_i, \sigma_i^*) \geq u_i(\delta_i', \sigma_i^*)$
 $\forall \delta_i \in \delta(\sigma_i^*) \quad \forall \delta_i' \in S_i \setminus \delta(\sigma_i^*)$

		2	
		A	B
1	A	2, 2	1, 2
	B	2, 1	1, 1

Define

$$b_i(\underline{\sigma}_i) = \left\{ \delta_i \in S_i : \begin{aligned} &u_i(\delta_i, \underline{\sigma}_i) \geq \\ &u_i(\delta_i', \underline{\sigma}_i) \\ &\forall \delta_i' \in S_i \end{aligned} \right\}$$

$$b_1(A) = \{A, B\}$$

$$b_1(B) = \{A, B\}$$

$$b_2(A) = \{A, B\}$$

$$b_2(B) = \{A, B\}$$

Let us compute MSNEs by considering supports one at a time

Support $\{A\} \times \{A\}$

This leads to the MSNE $((1,0), (1,0))$

Condition (1) vacuously true since

Condition (1) vacuously true since there is only one pure strategy in the support of the mixed strategies here:
 $\delta(\sigma_1^*) = \delta(\sigma_2^*) = \{A\}$

Condition (2)

Player 1:

$$u_1(A, A) = 2 \quad A \in \delta(\sigma_1^*)$$

$$u_1(B, A) = 2 \quad B \notin \delta(\sigma_1^*)$$

	A	B
A	2, 2	1, 2
B	2, 1	1, 1

Player 2:

$$u_2(A, A) = 2 \quad A \in \delta(\sigma_2^*)$$

$$u_2(A, B) = 2 \quad B \notin \delta(\sigma_2^*)$$

Note that equality is satisfied in both the cases

Thus condition (2) is satisfied.

Hence $(A, A) = ((1, 0), (1, 0))$ is a MSNE

Support $\{A\} \times \{A, B\}$

$$\sigma_1^* = (1, 0) \quad \delta(\sigma_1^*) = \{A\}$$

$$\sigma_2^* := (y, 1-y) \text{ with } 0 < y < 1 \quad \delta(\sigma_2^*) = \{A, B\}$$

Let us see if (σ_1^*, σ_2^*) leads to a MSNE

Condition (1)

Vacuously true for Player 1 since $\delta(\sigma_1^*) = \{A\}$

Player 2: $\delta(\sigma_2^*) = \{A, B\}$

$$u_2(A, A) = 2$$

$$u_2(A, B) = 2$$

	A	B
A	2, 2	1, 2
B	2, 1	1, 1

Hence condition (1) is satisfied.

Condition (2)
 vacuously true for player 2 since $\delta(\sigma_2^*) = \{A, B\}$
 Player 1: $A \in \delta(\sigma_1^*)$; $B \notin \delta(\sigma_1^*)$

$$u_1(A, (y, 1-y)) = 2y + (1-y) = y+1$$

$$u_1(B, (y, 1-y)) = 2y + (1-y) = y+1$$

Thus condition (2) is true for $\forall y: 0 < y < 1$

Hence $(A, (y, 1-y))$ is a MSNE
 $\forall 0 < y < 1$

Support $\{A, B\} \times \{A, B\}$

$$\text{Let } \sigma_1^* = (x, 1-x) \quad 0 < x < 1$$

$$\sigma_2^* = (y, 1-y) \quad 0 < y < 1$$

We explore if (σ_1^*, σ_2^*) leads to any MSNE

Condition (1)

Player 1: $\delta(\sigma_1^*) = \{A, B\}$

$$u_1(A, (y, 1-y)) = 2y + 1 - y = y+1$$

$$u_1(B, (y, 1-y)) = 2y + 1 - y = y+1$$

\therefore Satisfied

	A	B
A	2, 2	1, 2
B	2, 1	1, 1

Player 2: $\delta(\sigma_2^*) = \{A, B\}$

$$u_2((x, 1-x), A) = 2x + 1 - x = x+1$$

$$u_2((x, 1-x), B) = 2x + 1 - x = x+1$$

\therefore Satisfied

Condition (2) is vacuously true
 for both the players.

$(A, (y, 1-y))$ is a

for both the players.

Thus $((x, 1-x), (y, 1-y))$ is a MSNE $\forall 0 < x < 1 \quad \forall 0 < y < 1$

In fact, all supports here lead to MSNE.

Prisoner's Dilemma Problem

	NC	C
NC	-2, -2	-10, -1
C	-1, -10	-5, -5

We know that (C, C) is a SDSE.

Let us see if there is a MSNE:

$$((x, 1-x), (y, 1-y)) \quad \begin{array}{l} 0 < x < 1 \\ 0 < y < 1 \end{array}$$

Condition (2) vacuously satisfied

Condition (1)

Player 1:

$$u_1(\text{NC}, (y, 1-y)) = -2y - 10(1-y) = 8y - 10$$

$$u_1(\text{C}, (y, 1-y)) = -y - 5(1-y) = 4y - 5$$

If these have to be equal,

$$8y - 10 = 4y - 5 \Rightarrow 4y = 5 \Rightarrow y = \frac{5}{4}$$

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Impossible!

So, condition (1) can never be satisfied for any of the players.

Consider the support $\{c\} \times \{c\}$

Condition (1) is vacuously satisfied.

Condition (2)

Player 1:

$$u_1(c, c) = -5$$

$$u_1(NC, c) = -10$$

	NC	C
NC	-2, -2	-10, -1
C	-1, -10	-5, -5

Player 2:

$$u_2(c, c) = -5$$

$$u_2(c, NC) = -10$$

strict inequality holds in both the cases