CMPS 101

Homework Assignment 5

1. Define T(n) defined by the recurrence formula

$$T(n) = \begin{cases} 6 & 1 \le n < 3 \\ 2T(\lfloor n/3 \rfloor) + n & n \ge 3 \end{cases}$$

- a. Use the iteration method to write a summation formula for T(n).
- b. Use the summation in (a) to show that T(n) = O(n)
- c. Use the Master Theorem to show that $T(n) = \Theta(n)$
- 2. Use the Master theorem to find asymptotic solutions to the following recurrences.
 - a. T(n) = 7T(n/4) + n
 - b. $T(n) = 9T(n/3) + n^2$
 - c. $T(n) = 6T(n/5) + n^2$
 - d. $T(n) = 6T(n/5) + n\log(n)$
 - e. $T(n) = 7T(n/2) + n^2$
 - f. $S(n) = aS(n/4) + n^2$ (Note: your answer will depend on the parameter a.)
- 3. p.75: 4.3-2

The recurrence $T(n) = 7T(n/2) + n^2$ describes the running time of an algorithm A. A competing algorithm B has a running time given by $S(n) = aS(n/4) + n^2$. What is the largest integer value for a such that B is a faster algorithm than A (asymptotically speaking)? In other words, find the largest integer a such that S(n) = o(T(n)).

- 4. Let T(n) satisfy the recurrence T(n) = aT(n/b) + f(n), where f(n) is a polynomial satisfying $\deg(f) > \log_b(a)$. Prove that case (3) of the Master Theorem applies, and in particular that the regularity condition necessarily holds.
- 5. Show that the number vertices of odd degree in any graph must be even. (Hint: Use the Handshake Lemma mentioned in the Graph Theory handout.)

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