

HW2 CMPS 101

1. Prove $f(n) + g(n) = \Theta(\max(f(n), g(n)))$.

assuming knowing $f(n)$ and $g(n)$ are asymptotically non-negative functions.

assume $\max(f(n), g(n)) = \Theta(f(n) + g(n))$

\Rightarrow this is only true if and only if there exists $c_1, c_2, n_0 > 0$

so \Rightarrow

$$0 \leq c_1(f(n) + g(n)) \leq \max(f(n), g(n)) \leq c_2(f(n) + g(n)) \text{ for all } n \geq n_0$$

we also get $f(n) + g(n) \geq f(n) \geq 0$ & $f(n) + g(n) \geq g(n) \geq 0$

combining the two inequalities it becomes $f(n) + g(n) \geq \max(f(n), g(n))$

thus $\max(f(n), g(n)) \leq c(f(n) + g(n))$ w/ $n \geq n_0$ & $c=1$

it is also true that $\max(f(n), g(n)) \geq f(n)$ $\max(f(n), g(n)) \geq g(n)$ $n \geq n_0$

$$\begin{aligned} 2 \max(f(n), g(n)) &\geq (g(n) + f(n)) \\ &= \max(f(n), g(n)) \geq \frac{1}{2}(g(n) + f(n)) \text{ for all } n \geq n_0 \end{aligned}$$

$$\therefore \max(f(n), g(n)) = \Omega(g(n) + f(n)) \text{ w/ constant } c = \frac{1}{2}$$

2. Explain why "The running time of algorithm A is at least $O(n^2)$ " is meaningless.

To say that statement is the running time, $T(n)$, could be ~~asymptotically~~ = or asymptotically $>$ than $O(n^2)$
ex) ~~A~~ A rocket ship goes at least 2 meters off the ground.

3 determine true or false.

$\triangleright 2^{n+1} = O(2^n)$

$\hookrightarrow 2^{n+1} = 2(2^n) \Rightarrow 2(2^n) \leq C \cdot 2^n$

\Rightarrow True if $C \geq 2$

$\triangleright 2^{2n} = O(2^n)$

\hookrightarrow is $2^{2n} \leq C \cdot 2^n$

\Rightarrow false due to ratio

$\hookrightarrow \frac{2^{2n}}{C \cdot 2^n} = \frac{(2^n)^2}{C \cdot 2^n} = \frac{2^n}{C}$

$C \cdot 2^n$ can never be an upper bound

4	A	B	O	o	Ω	ω	Θ
a	$\lg^k n$	n^c	yes	yes	no	no	no
b	n^k	C^n	yes	yes	no	no	no
c	\sqrt{n}	$n^{\sin n}$	no	no	no	no	no
d	2^n	$2^{n/2}$	no	no	yes	yes	no
e	$n^{\lg c}$	$c \log n$	yes	no	yes	no	yes
f	$\lg(n!)$	$\lg(n^n)$	yes	no	yes	no	yes

a) using L'Hospital's rule $\lim_{n \rightarrow \infty} \frac{(\lg n)^k}{n^c} = \lim_{n \rightarrow \infty} \frac{k(\lg n)^{k-1}}{c n^{c-1}}$

\swarrow must do k times together here

$= \lim_{n \rightarrow \infty} \frac{k(k-1)(\lg n)^{k-2} \frac{1}{n}}{c^2 n^{c-2}} \leftarrow \dots$

$\hookrightarrow = \lim_{n \rightarrow \infty} \frac{k(k-1)(k-2) \dots 1}{c^k n^k} = 0$

b) L'Hospital $\Rightarrow \lim_{n \rightarrow \infty} \frac{n^k}{c^n} = 0$ applying the rule repeatedly to see $n^k = o(c^n)$

c) I can see B, $n^{\sin n}$, contains $\sin n$ which goes between -1 & 1 . Maxing out at its maximum value $n^{\sin n} > c\sqrt{n}$ and so $n^{\sin n} \neq O(\sqrt{n})$. Also when at its minimum $n^{\sin n} < c\sqrt{n}$ thus $n^{\sin n} \neq \Omega(\sqrt{n})$.

d. $\lim_{n \rightarrow \infty} \frac{2^n}{2^{\frac{n}{2}}} = \infty$ meaning $2^n = \omega(2^{\frac{n}{2}})$

e. $n^{b^c} = c^{b^a}$ equal

f. $\lg(n^n) = n(\lg n)$ & using stirling's formula $\lg(n!) = \Theta(n \lg n)$

5. d. $f(n) = O(g(n))$ implies $2^{f(n)} = O(2^{g(n)})$

false. if $f(n) = 2n$; $g(n) = n \Rightarrow f(n) = O(g(n))$ but $2^{2n} = 4^n \neq O(2^n)$

e. $f(n) = O((f(n))^2)$

false to make this true, there must be positive constants c and n_0 such that for all $n \geq n_0$
 $0 \leq f(n) \leq c(f(n))^2$

unless $f(n) = \frac{1}{n}$ then it doesn't hold. since
 there no positive constant c and no such n_0
 all $n \geq n_0$

$$0 \leq \frac{1}{n} \leq c \frac{1}{n^2}$$

h. $f(n) + o(f(n)) = \Theta(f(n))$

in definition $0 \leq o(f(n)) \leq f(n)$, so $f(n) \leq f(n) + o(f(n)) \leq 2f(n)$

$\therefore f(n) + o(f(n)) = \Theta(f(n))$

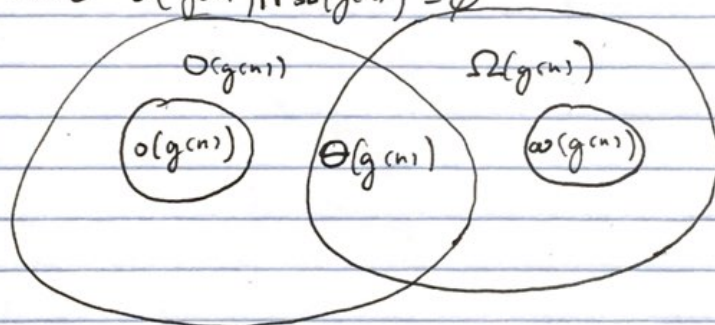
6. $f(n) = \Theta(n)$, Prove $\sum_{i=1}^n f(i) = \Theta(n^2)$

$$f(n) = \Theta(n) \text{ iff } \lim_{n \rightarrow \infty} \frac{f(n)}{n} = c$$

$$f(n) = o(n) \text{ if and only if } \forall c > 0, \exists n_0 > 0, \forall n \geq n_0: 0 \leq \frac{f(n)}{n} < c$$

7. Let $g(n)$ be asymptotically non-negative function.

$$\text{Prove } o(g(n)) \cap \Omega(g(n)) = \emptyset$$



no element in $o(g(n))$ cannot be in $\Omega(g(n))$

Let $f(n) = o(g(n))$ then for any $c > 0$, there exists a constant $n_0 > 0$ such that $0 \leq f(n) < cg(n)$ for all $n \geq n_0$. Meaning there are no positive constants c and n_0 such that $0 \leq cg(n) \leq f(n)$ for all $n \geq n_0$

$$\text{so } o(g(n)) \cap \Omega(g(n)) = \emptyset$$