



Figure 1.4: In pink an *horizontal non-separating* closed curve on the torus, in red a *vertical* one.

Being non-separating is a topological property, we will work with a characterization more suited for our setting.

We call an *horizontal loop* a simple closed curve on the torus that is composed only of horizontal links, like the curve in pink in figure 1.4. Their vertical counterpart is called *vertical loop*, see the curve in red in figure 1.4.

The characterization will exploit the fact that if a closed curve is separating, then the number of crossings with any horizontal and vertical loops  $l$  is even, call this number  $Cross_l(\gamma)$ .

**Proposition 1.3.3.** A closed curve  $\gamma$  on the torus is non-separating if and only if we can find an horizontal/vertical loop  $l$  such that  $Cross_l(\gamma)$  is odd.

*Proof.* If  $\gamma$  is separating, meaning that it's the boundary between two path disconnected regions of the torus, then  $Cross_l(\gamma)$  is always even for any horizontal/vertical loop  $l$ . This can be proven using the same argument used in the proof of the Jordan's curve theorem for polygons.  $\square$

**Proposition 1.3.4.** Let  $m \in \Sigma_G$  with  $G = (V_n, E_n^p)$  be a state consisting of only separating curves. If given a square transformation  $X_s$  the new state  $m' = X_s(m)$  contains one non-separating curve  $c$ , then it contains two.

*Proof.* We need to consider only transformation that add more links, namely uniform +1 and single +2. If the new state  $m' = X_s(m)$  contains a non-separating curve  $\gamma$ , then one or two of its edges was introduced by the transformation  $X_s$ . This means that in the original state  $m$ , all the vertices in  $\gamma$  were already in the same connected component (a circuit is 2-connected), then by the Euler theorem for multigraphs we can find in  $m$  a circuit  $C$  which visits every vertex in  $\gamma$ . This implies that there exists exactly two edge-disjoint paths joining any pair of vertices. We will employ these paths to build two non-separating curves. We need to check two cases:

- $X_s$  is a *single link* +2. Call  $e_1$  and  $e_2$  the two added links with endpoints  $x$  and  $y$ . Call  $\gamma$  the non-separating curve, which we can write as the union of a path  $P_{xy}$  from  $x$  to  $y$  contained in the circuit  $C$  and one of the new edges  $e_1$ :

$$\gamma = P_{xy} \cup \{e_1\}$$

Call  $\tilde{P}_{xy}$  the other  $xy$ -path in  $C$ , then the closed curve  $\tilde{\gamma} = \tilde{P}_{xy} \cup \{e_2\}$  is also non-separating, since given an horizontal/vertical loop  $l$ :

$$Cross_l(\tilde{\gamma}) = Cross_l(C) - Cross_l(\gamma)$$

where  $Cross_l(C)$  is even, since it's a closed curve contained in the original state  $m$ , and  $Cross_l(\gamma)$  is odd since  $\gamma$  is non-separating, this implies that  $Cross(\tilde{\gamma})$  is also odd.

- $X_s$  is a uniform  $+1$ . Call  $e_1, e_2, e_3$  and  $e_4$  the new added edges,  $x$  and  $y$  the two vertices on  $\gamma$  that also are affected by the transformation. Again we can write:

$$\gamma = P_{xy} \cup \{e_1, e_2\}$$

$$\tilde{\gamma} = \tilde{P}_{xy} \cup \{e_3, e_4\}$$

the same parity of crossings argument applies here for  $\tilde{\gamma}$ .

□