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5
        Java Fast Input
6
7
    class InputReader {
        private final BufferedReader reader;
8
        private StringTokenizer tokenizer;
8
        public InputReader(InputStream stream) {
9
            reader = new BufferedReader(new InputStreamReader(stream));
            tokenizer = null;
10
       }
10
        public String nextLine() {
11
            try {
                return reader.readLine();
12
            } catch (IOException e) {
                throw new RuntimeException(e);
13
13
       }
14
        public String next() {
            while (tokenizer == null | !tokenizer.hasMoreTokens()) {
15
                tokenizer = new StringTokenizer(nextLine());
15
            return tokenizer.nextToken():
16
17
        public int nextInt() {
17
            return Integer.parseInt(next());
```

```
Simple Max Matching
bool dfs(int u) {
    if (mx[u] == T) return false:
    mx[u] = T;
    for(int v : ke[u]) {
        if (!my[v] || dfs(my[v])) {
            my[v] = u;
            return true;
        }
    }
    return false;
}
int main() {
    For(i,1,n) {
        T++:
        res += dfs(i);
    }
    // choose my & i
    Konig
void konig(){
   queue <int> qu;
   f1(i,m) if (!Assigned[i]) qu.push(i);
   f1(i,n) if (!Assigned[N-i]) qu.push(N-i);
   while (qu.size()){
      int u=qu.front(); qu.pop();
      for (int i=0; int v=a[u][i]; i++)
      if (!(Choosed[v]++)) qu.push(Assigned[v]);
   }
   f1(i,m) if (Assigned[i] && !Choosed[i] && !Choosed[Assigned[i]])
   Choosed[i]=true:
}
```

4 Hopcroft Karp Max Matching algorithm

```
// Worse Case: E * sqrt(v)
const int MAXN = 50005, MAXM = 50005;
vector < int > gph[MAXN];
int dis[MAXN], 1[MAXN], r[MAXM], vis[MAXN];
void clear() {
    for (int i = 0; i < MAXN; i++)
        gph[i].clear();
void add_edge(int 1, int r) {
    gph[1].push_back(r);
bool bfs(int n) {
    queue < int > que;
    bool ok = 0;
    memset(dis, 0, sizeof(dis));
    for (int i = 0; i < n; i++) {
        if (1[i] == -1 && !dis[i]) {
            que.push(i);
            dis[i] = 1;
        }
    while (!que.empty()) {
        int x = que.front();
        que.pop();
        for (auto &i : gph[x]) {
            if (r[i] == -1)
                 ok = 1;
            else if (!dis[r[i]]) {
                 dis[r[i]] = dis[x] + 1;
                que.push(r[i]);
    return ok;
bool dfs(int x) {
    for (auto &i : gph[x]) {
        if (r[i] == -1 \mid | (!vis[r[i]] && dis[r[i]] == dis[x] + 1 &&
                            dfs(r[i]))) {
            vis[r[i]] = 1;
            1 \lceil x \rceil = i:
            r[i] = x;
```

```
return 1;
        }
    }
    return 0;
int match(int n) {
    memset(1, -1, sizeof(1));
    memset(r, -1, sizeof(r));
    int ret = 0;
    while (bfs(n)) {
        memset(vis, 0, sizeof(vis));
        for (int i = 0; i < n; i++)
            if (1[i] == -1 && dfs(i))
                ret++;
    }
    return ret:
bool chk[MAXN + MAXM];
void rdfs(int x, int n) {
    if (chk[x])
        return:
    chk[x] = 1;
    for (auto &i : gph[x]) {
        chk[i + n] = 1;
        rdfs(r[i], n);
    }
}
vector<int> getcover(int n, int m) {
    // solve min. vertex cover
    match(n):
    memset(chk, 0, sizeof(chk));
    for (int i = 0; i < n; i++)
        if(1[i] == -1)
            rdfs(i. n):
    vector <int> v;
    for (int i = 0; i < n; i++)
        if (!chk[i])
            v.push_back(i);
    for (int i = n; i < n + m; i++)
        if (chk[i])
            v.push_back(i);
    return v;
```

5 Max matching min cost

```
// numbered from 0. i \rightarrow mx[i]
const int V = 1000. INF = 1e9:
int g[V][V], mx[V], my[V], fx[V], fy[V], d[V], ar[V], tr[V], p;
int slack(int u, int v) {
    return g[u][v] - fx[u] - fy[v];
int augment(int s) {
    queue < int > q;
    q.push(s);
    fill_n(tr, p, -1);
   for(int i = 0; i < p; ++i) d[i] = slack(s, i), ar[i] = s;
    while(true) {
        while(!q.empty()) {
            int u = q.front();
            q.pop();
            for(int v = 0; v < p; ++v) if(tr[v] == -1) {
                    int w = slack(u, v);
                    if(w == 0)
                        tr[v] = u;
                        if(my[v] == -1) return v;
                        q.push(my[v]);
                    if(d[v] > w) d[v] = w, ar[v] = u;
        int delta = INF;
        for(int v = 0; v < p; ++v) if(tr[v] == -1) delta =
                    min(delta. d[v]):
        fx[s] += delta:
        for(int v = 0; v < p; ++v)
            if(tr[v] == -1) d[v] -= delta;
            else fx[my[v]] += delta, fy[v] -= delta;
        for (int v = 0; v < p; ++v) if (tr[v] == -1 && d[v] == 0) {
                tr[v] = ar[v];
                if(my[v] == -1) return v;
                q.push(my[v]);
   }
void maxMatchMinCost() {
   fill_n(mx, p, -1);
    fill_n(my, p, -1);
```

```
for(int i = 0; i < p; ++i) fx[i] = *min_element(g[i], g[i]+p);
for(int s = 0; s < p; ++s) {
    int f = augment(s);
    while(f != -1) {
        int x = tr[f], nx = mx[x];
        mx[x] = f;
        my[f] = x;
        f = nx;
    }
}</pre>
```

6 Ganeral Matching

```
class MatchingGraph {
public:
    vector <vector <int> > adj;
    vector <bool> blossom;
    vector <int> parent;
    vector <int> base;
    vector <int> match;
    int n;
    MatchingGraph() {
        n = 0;
    void addEdge(int x, int y) {
        adj[x].push_back(y);
        adj[y].push_back(x);
    void clearGraph() {
        int i:
        for (i=0; i < SZ(adj); ++i)
            adj[i].clear();
        fill(blossom.begin(),blossom.end(),false);
        fill(parent.begin(),parent.end(),-1);
        for (i=0; i < n; ++i)
            base[i] = i;
        for (i=0; i<n; ++i)
            match[i] = -1:
    void setN(int newn) {
        n = newn:
        adj.resize(n);
```

```
blossom.resize(n);
    base.resize(n):
    match.resize(n):
    parent.resize(n);
    clearGraph();
int lca(int x, int y) {
    vector <bool> fy;
    fy.resize(n);
    fill(fy.begin(),fy.end(),false);
    while (true) {
        x = base[x];
        fv[x] = true;
        if (match[x] == -1)
            break:
        x = parent[match[x]];
    while (true) {
        y = base[y];
        if (fy[y])
            return y;
        y = parent[match[y]];
    return -1;
void path(int now, int child, int curbase) {
    while (base[now] != curbase) {
        blossom[base[now]] = blossom[base[match[now]]] = true;
        parent[now] = child;
        child = match[now]:
        now = parent[match[now]];
   }
int augmentPath(int x) {
    int i, j;
    for (i=0; i<n; ++i)
        base[i] = i;
    for (i=0; i<n; ++i)
        parent[i] = -1;
    queue <int> bfs;
    vector <bool> sudah;
    sudah.resize(n);
    fill(sudah.begin(), sudah.end(), false);
    sudah[x] = true:
```

```
bfs.push(x);
    while (!bfs.empty()) {
        int now = bfs.front():
        bfs.pop();
        for (i=0; i<SZ(adj[now]); ++i) {
            int next = adj[now][i];
            if (base[next] == base[now] || match[next] == now);
            else if (next == x || (match[next]!=-1 &&
                                    parent[match[next]]!=-1)) {
                int curbase = lca(now,next);
                fill(blossom.begin(),blossom.end(),false);
                path(now, next, curbase);
                path(next, now, curbase);
                for (j = 0; j < n; ++j)
                    if (blossom[j]) {
                        base[j] = curbase;
                        if (!sudah[j]) {
                             sudah[j] = true;
                             bfs.push(j);
            } else if (parent[next] == -1) {
                parent[next] = now;
                if (match[next] == -1)
                    return next:
                sudah[match[next]] = true;
                bfs.push(match[next]);
            }
        }
    return -1;
}
int edmondsMatch() {
    int i:
    int res = 0:
    for (i=0; i<n; ++i) {
        if (match[i]==-1) {
            int x = augmentPath(i);
            while (x > = 0) {
                int p = parent[x];
                int pp = match[p];
                match[x] = p;
                match[p] = x;
                x = pp;
```

```
for (i=0; i < n; ++i)
            if (match[i]!=-1)
                ++res:
        return res >> 1;
   }
};
    Dinic MaxFlow
class DinicFlow {
private:
    vector<int> dist, head, work;
    vector < int > point, flow, capa, next;
    int n, m;
    bool bfs(int s, int t) {
        For(i, 1, n) dist[i] = -1;
        queue <int> q;
        dist[s] = 0;
        q.push(s);
        while (!q.empty()) {
            int u = q.front();
            q.pop();
            for (int i = head[u]; i >= 0; i = next[i])
                if (flow[i] < capa[i] && dist[point[i]] < 0) {</pre>
                    dist[point[i]] = dist[u] + 1;
                    q.push(point[i]);
        return dist[t] >= 0:
   }
    int dfs(int s, int t, int f) {
        if (s == t) return f;
        for (int &i = work[s]; i >= 0; i = next[i])
            if (flow[i] < capa[i] && dist[point[i]] == dist[s] + 1) {</pre>
                int d = dfs(point[i], t, min(f, capa[i] - flow[i]));
                if (d > 0) {
                    flow[i] += d:
                    flow[i ^ 1] -= d:
```

```
return d;
                }
        return 0;
    }
public:
    DinicFlow(int n = 0) {
        this \rightarrow n = n;
        this -> m = 0;
        dist.assign(n + 7, 0);
        head.assign(n + 7, -1);
        work.assign(n + 7, 0);
    }
    void addEdge(int u, int v, int c1, int c2 = 0) {
        point.push_back(v);
        capa.push_back(c1);
        flow.push_back(0);
        next.push_back(head[u]);
        head[u] = m++:
        point.push_back(u);
        capa.push_back(c2);
        flow.push_back(0);
        next.push_back(head[v]);
        head[v] = m++;
    }
    int maxFlow(int s, int t) {
        int totFlow = 0:
        while (bfs(s, t)) {
            For(i, 1, n) work[i] = head[i];
            while (true) {
                int d = dfs(s, t, cmax);
                if (d == 0) break;
                totFlow += d;
            }
        return totFlow;
    }
```

Mincost MaxFlow SPFA

```
Min Cost Max Flow - SPFA
Index from 0
edges cap changed during find flow
Lots of double comparison --> likely to fail for double
MinCostFlow mcf(n);
mcf.addEdge(u, v, cap, cost);
cout << mcf.minCostFlow() << endl;</pre>
template < class Flow = int, class Cost = int >
struct MinCostFlow {
    const Flow INF_FLOW = 1000111000;
    const Cost INF_COST = 1000111000111000LL;
    int n, t, S, T;
   Flow totalFlow:
    Cost totalCost;
    vector < int > last, visited;
    vector < Cost > dis:
    struct Edge {
        int to;
       Flow cap;
        Cost cost;
        int next;
        Edge(int to, Flow cap, Cost cost, int next) :
                to(to), cap(cap), cost(cost), next(next) {}
   };
    vector < Edge > edges;
   MinCostFlow(int n): n(n), t(0), totalFlow(0), totalCost(0), last(n,
        -1), visited(n, 0), dis(n, 0) {
        edges.clear();
   }
    int addEdge(int from, int to, Flow cap, Cost cost) {
        edges.push_back(Edge(to, cap, cost, last[from]));
        last[from] = t++:
        edges.push_back(Edge(from, 0, -cost, last[to]));
        last[to] = t++;
        return t - 2;
   }
    pair < Flow , Cost > minCostFlow(int _S , int _T) {
```

```
S = _S; T = _T;
        SPFA():
        while (1) {
            while (1) {
                REP(i,n) visited[i] = 0;
                if (!findFlow(S, INF_FLOW)) break;
            }
            if (!modifyLabel()) break;
        return make_pair(totalFlow, totalCost);
    }
private:
    void SPFA() {
        REP(i,n) dis[i] = INF_COST;
        priority_queue < pair < Cost, int > > Q;
        Q.push(make_pair(dis[S]=0, S));
        while (!Q.empty()) {
            int x = Q.top().second;
            Cost d = -Q.top().first;
            Q.pop();
            // For double: dis[x] > d + EPS
            if (dis[x] != d) continue:
            for(int it = last[x]; it >= 0; it = edges[it].next)
                if (edges[it].cap > 0 && dis[edges[it].to] > d + edges[it
                    Q.push(make_pair(-(dis[edges[it].to] = d + edges[it].
                        cost), edges[it].to));
        Cost disT = dis[T]: REP(i,n) dis[i] = disT - dis[i]:
    }
    Flow findFlow(int x, Flow flow) {
        if (x == T) {
            totalCost += dis[S] * flow;
            totalFlow += flow;
            return flow;
        visited[x] = 1;
        Flow now = flow;
        for(int it = last[x]; it >= 0; it = edges[it].next)
            // For double: fabs(dis[edges[it].to] + edges[it].cost - dis[x
                1) < EPS
            if (edges[it].cap && !visited[edges[it].to] && dis[edges[it].
```

```
to] + edges[it].cost == dis[x]) {
                Flow tmp = findFlow(edges[it].to. min(now. edges[it].cap))
                edges[it].cap -= tmp;
                edges[it ^ 1].cap += tmp;
                now -= tmp;
                if (!now) break;
            }
        return flow - now;
   }
    bool modifyLabel() {
        Cost d = INF_COST;
        REP(i,n) if (visited[i])
            for(int it = last[i]; it >= 0; it = edges[it].next)
                if (edges[it].cap && !visited[edges[it].to])
                    d = min(d, dis[edges[it].to] + edges[it].cost - dis[i
                        1):
        // For double: if (d > INF_COST / 10)
                                                   INF COST = 1e20
        if (d == INF COST) return false:
        REP(i,n) if (visited[i])
            dis[i] += d:
        return true:
};
```

9 Upper Lower

```
For each edge in original flow:

Add edge with cap = upper bound - lower bound.

Add source s, sink t.
Let M[v] = (sum of lowerbounds of ingoing edges to v) - (sum of lower bounds of outgoing edges from v).
For all v, if M[v] > 0, add (s, v, M), else add (v, t, -M).
If all outgoing edges from S are full --> feasible flow exists, it is flow + lower bounds.

Feasible flow in network with upper + lower constraint, with source & sink :

Add edge (t, s) with capacity [0, INF].
Check feasible in network without source & sink.
```

```
Max flow with both upper + lower constraints, source s, sink t: add edge (
    t. s. +INF).
- Binary search lower bound, check whether feasible flow exists WITHOUT
    source / sink
     Two sat
10
int n, m, g[maxn];
bool cx[maxn];
vector <int> listV, ke[maxn], K[maxn];
int cal(int x) {
    if (x\%2 == 0) return x - 1;
    else return x + 1;
}
void add(int u, int v) {
    ke[u].pb(v);
    K[v].pb(u);
}
void dfs(int u) {
    cx[u] = true;
    for(int v : ke[u])
        if (!cx[v]) dfs(v):
    listV.pb(u);
}
void dfs(int u, int x) {
    g[u] = x:
    for(int v : K[u])
        if (g[v] == 0) dfs(v,x);
}
int main() {
    cin >> m >> n;
    n += n;
    For(i,1,m) {
        int u. v:
        cin >> u >> v;
        u *= 2;
        v *= 2:
        if (u < 0) u = cal(abs(u));
```

```
if (v < 0) v = cal(abs(v));
        add(cal(u).v):
        add(cal(v).u):
   listV.pb(0);
    For(i,1,n)
    if (!cx[i]) dfs(i);
    int ng = 0;
    Ford(i,n,1) {
       int u = listV[i];
        if (g[u] == 0) dfs(u, ++ng);
    for(int i = 2; i \le n; i += 2)
        if (g[i] == g[i-1]) NO;
    YES:
    vector <int> result;
    for(int i = 2; i <= n; i += 2)
        if (g[i] > g[i-1]) result.pb(i>>1);
    Alternative Tree
int n, m, 1, q, t, res, test,
    a[maxn], tin[maxn], tout[maxn], mark[maxn], terror[maxn], f[maxn][20];
vector < int > adj[maxn], _adj[maxn];
stack<int> stk:
void visit(const int &u) {
    tin[u] = ++t:
    for(int i = 1; i \le 1; ++i) f[u][i] = f[f[u][i-1]][i-1];
    for(auto v : adj[u])
       if (v != f[u][0]) {
           f[v][0] = u;
            visit(v);
        }
    tout[u] = ++t;
bool anc(const int &u. const int &v) {
    return tin[u] <= tin[v] && tout[u] >= tout[v];
int lca(int u, int v) {
```

if (anc(u,v)) return u;

```
if (anc(v,u)) return v:
    for(int i = 1; i >= 0; --i)
        if (!anc(f[u][i],v)) u = f[u][i];
    return f[u][0];
}
void query() {
    cin >> m;
    for(int i = 1; i <= m; ++i) {
        cin >> a[i];
        _adj[a[i]].clear();
        mark[a[i]] = test;
        terror[a[i]] = test;
    }
    sort(a+1,a+m+1,cmp);
    for(int i = 1; i < m; ++i) {
        int tmp = lca(a[i],a[i+1]);
        if (mark[tmp] < test) {</pre>
            mark[tmp] = test;
            a[++m] = tmp;
            _adj[tmp].clear();
    // sort theo tin
    sort(a+1,a+m+1,cmp);
    while (!stk.empty()) stk.pop();
    stk.push(a[1]);
    for(int i = 2; i <= m; ++i) {
        while (tout[stk.top()] < tout[a[i]]) stk.pop();
        _adj[stk.top()].push_back(a[i]);
        stk.push(a[i]);
    }
    res = 0;
    check(a[1]);
    cout << res << "\n";
}
int main() {
   1 = log2(n);
    cin >> q;
    f[1][0] = 1;
    visit(1);
    for(test = 1; test <= q; ++test) query();
```

12 Max Clique

```
class MaxClique {
public:
    static const int MV = 210;
    int V;
    int el[MV][MV/30+1];
    int dp[MV];
    int ans:
    int s[MV][MV/30+1];
    vector < int > sol;
    void init(int v) {
        V = v; ans = 0;
        FZ(el); FZ(dp);
   }
    /* Zero Base */
    void addEdge(int u, int v) {
        if(u > v) swap(u, v);
        if(u == v) return;
        el[u][v/32] = (1 << (v\%32));
    bool dfs(int v, int k) {
        int c = 0, d = 0;
        for(int i=0; i<(V+31)/32; i++) {
            s[k][i] = el[v][i];
            if(k != 1) s[k][i] &= s[k-1][i];
            c += __builtin_popcount(s[k][i]);
       }
        if(c == 0) {
            if(k > ans) {
                ans = k;
                sol.clear();
                sol.push_back(v);
                return 1;
            return 0;
```

- - for all u, |in_deg[u] - out_deg[u] | <= 1

- - At most 1 vertex with in_deg[u] - out_deg[u] = 1

```
for (int i=0; i<(V+31)/32; i++) {
            for(int a = s[k][i]: a : d++) {
                if(k + (c-d) <= ans) return 0:
                int lb = a\&(-a), lg = 0;
                a ^= lb;
                while(lb!=1) {
                    lb = (unsigned int)(lb) >> 1;
                    lg ++;
                }
                int u = i*32 + lg;
                if(k + dp[u] <= ans) return 0;
                if(dfs(u, k+1)) {
                    sol.push_back(v);
                    return 1;
                }
            }
        return 0;
    }
    int solve() {
        for(int i=V-1; i>=0; i--) {
            dfs(i, 1);
            dp[i] = ans;
        return ans;
    }
};
     Euler Path
 - When choosing starting vertex (for calling find_path), make sure deg[
     start] > 0.
 - If find Euler path, starting vertex must have odd degree.
 - Check no solution: SZ(path) == nEdge + 1.
 If directed:
 - Edge --> int
 - add_edge(int a, int b) { adj[a].push_back(b); }
 - Check for no solution:
```

```
- - At most 1 vertex with out_deg[u] - in_deg[u] = 1 (start vertex)
 - BFS from start vertex. all vertices u with out deg[u] > 0 must be
     visited
struct Edge {
    int to;
    list < Edge > :: iterator rev;
    Edge(int to) :to(to) {}
};
const int MN = 100111;
list < Edge > adj[MN];
vector <int> path; // our result
void find_path(int v) {
    while(adj[v].size() > 0) {
        int vn = adj[v].front().to;
        adj[vn].erase(adj[v].front().rev);
        adj[v].pop_front();
        find_path(vn);
    path.push_back(v);
void add_edge(int a, int b) {
    adj[a].push_front(Edge(b));
    auto ita = adj[a].begin();
    adj[b].push_front(Edge(a));
    auto itb = adj[b].begin();
    ita->rev = itb:
    itb -> rev = ita;
14 Aho Corasick
const int NODE = (int) 1e6 + 1;
const int NC = 26;
int nextNode[NODE][NC]:
int chr[NODE];
int parent[NODE];
int prefix[NODE];
```

int numNodes;

```
set < int > match [NODE];
int getPrefix(int);
int go(int u, int c) {
    if (nextNode[u][c] != -1) return nextNode[u][c]:
    if (u == 0) return 0:
    return nextNode[u][c] = go(getPrefix(u), c);
}
int getPrefix(int u) {
    if (prefix[u] != -1) return prefix[u];
    if (u == 0 || parent[u] == 0) return prefix[u] = 0;
    return prefix[u] = go(getPrefix(parent[u]), chr[u]);
}
void add(const string &s, int id) {
    int u = 0:
    for (int i = 0; i < (int) s.size(); ++i) {
        int c = s[i] - A':
        if (nextNode[u][c] == -1) {
            nextNode[u][c] = numNodes;
            fill(nextNode[numNodes], nextNode[numNodes] + NC, -1);
            chr[numNodes] = c;
            parent[numNodes] = u;
            prefix[numNodes] = -1;
            match[numNodes].clear();
            match[numNodes].insert(-1);
            ++numNodes:
        u = nextNode[u][c];
    match[u].insert(id);
}
set<int>& getMatch(int u) {
    if (match[u].count(-1) == 0) return match[u];
    const set<int> &foo = getMatch(getPrefix(u));
    match[u].insert(foo.begin(), foo.end());
    match[u].erase(-1):
    return match[u];
void init() {
```

```
fill(nextNode[0], nextNode[0] + NC, -1);
    numNodes = 1:
    Suffix Array
struct SuffixArray {
    const int L;
    string s;
    vector<vector<int> > P;
    vector < pair < pair < int , int > , int > > M;
    SuffixArray(const string &s) : L(s.length()), s(s), P(1,
                vector < int > (L, 0)), M(L) {
       for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
        for (int skip = 1, lv = 1; skip < L; skip *= 2, lv++) {
            P.push_back(vector < int > (L, 0));
            for (int i = 0; i < L; i++)
                M[i] = make_pair(make_pair(P[lv-1][i], i + skip < L
                                            ? P[lv-1][i + skip] : -1000), i
                                                ):
            sort(M.begin(), M.end());
            for (int i = 0; i < L; i++)
                P[lv][M[i].se] = (i > 0 && M[i].fi == M[i-1].fi) ?
                                 P[lv][M[i-1].se] : i:
        }
    vector < int > GetSuffixArray() {
        return P.back();
// returns the length of the longest common prefix of s[i...L-1]
    and s[j...L-1]
    int LongestCommonPrefix(int i, int j) {
        int len = 0:
       if (i == j) return L - i;
        for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
            if (P[k][i] == P[k][i]) {
                i += 1 << k;
                j += 1 << k;
                len += 1 << k:
        return len:
```

16 Suffix Array O(n)

};

```
#include <bits/stdc++.h>
#define FOR(i,a,b) for (int i=(a),b=(b);i \le b;i=i+1)
#define REP(i,n) for (int i=0,_n=(n);i < n;i=i+1)
#define MASK(i) (1LL << (i))
#define BIT(x,i) (((x)>>(i))&1)
#define tget(i) BIT(t[(i) >> 3], (i) & 7)
#define tset(i, b) { if (b) t\lceil (i) \rangle > 3\rceil \mid = MASK((i) \& 7): else t\lceil (i) \rangle > 3\rceil
     &= \text{^{\sim}MASK((i) & 7):} 
#define chr(i) (cs == sizeof(int) ? ((int *)s)[i] : ((unc *)s)[i])
#define isLMS(i) ((i) > 0 && tget(i) && !tget((i) - 1))
typedef unsigned char unc;
class SuffixArray {
    public:
    int *sa, *lcp, *rank, n;
    unc *s:
    void getbuckets(unc s[], vector<int> &bkt, int n, int k, int cs, bool
        end) {
        FOR(i, 0, k) bkt[i] = 0;
        REP(i, n) bkt[chr(i)]++;
        int sum = 0:
        FOR(i, 0, k) {
             sum += bkt[i];
             bkt[i] = end ? sum : sum - bkt[i];
        }
    }
    void inducesal(vector < unc > &t, int sa[], unc s[], vector < int > &bkt,
        int n, int k, int cs, bool end) {
        getbuckets(s, bkt, n, k, cs, end);
        REP(i, n) {
             int j = sa[i] - 1;
            if (j >= 0 && !tget(j)) sa[bkt[chr(j)]++] = j;
        }
    }
    void inducesas(vector < unc > &t. int sa[]. unc s[]. vector < int > &bkt.
        int n, int k, int cs, bool end) {
        getbuckets(s, bkt, n, k, cs, end);
        FORD(i, n - 1, 0) {
             int j = sa[i] - 1;
```

```
if (i >= 0 \&\& tget(i)) sa[--bkt[chr(i)]]=i;
}
void build(unc s[], int sa[], int n, int k, int cs) {
    vector < unc > t = vector < unc > (n / 8 + 1, 0):
    tset(n - 2, 0);
    tset(n - 1, 1);
    FORD(i, n - 3, 0) tset(i, chr(i) < chr(i+1) || (chr(i) == chr(i+1))
         && tget(i+1)));
    vector < int > bkt = vector < int > (k + 1, 0);
    getbuckets(s, bkt, n, k, cs, true);
    REP(i, n) sa[i] = -1;
    REP(i, n) if (isLMS(i)) sa[--bkt[chr(i)]] = i;
    inducesal(t, sa, s, bkt, n, k, cs, false);
    inducesas(t. sa. s. bkt. n. k. cs. true):
    bkt.clear():
    int n1 = 0;
    REP(i, n) if (isLMS(sa[i])) sa[n1++] = sa[i];
    FOR(i, n1, n - 1) sa[i] = -1;
    int name = 0:
    int prev = -1;
    REP(i, n1) {
        int pos = sa[i];
        bool diff = false:
        REP(d, n) {
            if (prev < 0 || chr(prev + d) != chr(pos + d) || tget(prev
                 + d) != tget(pos + d)) {
                diff = true;
                break:
            else if (d > 0 && (isLMS(prev + d) || isLMS(pos + d)))
                break;
        if (diff) {
            name++;
            prev = pos;
        sa[n1 + pos / 2] = name - 1;
    }
    FORD(i, n - 1, n1) if (sa[i] >= 0) sa[j--] = sa[i];
```

```
if (name < n1) build((unc *)s1, sa1, n1, name-1, sizeof(int));
    else REP(i, n1) sa1[s1[i]] = i:
    bkt.assign(k + 1, 0):
    getbuckets(s, bkt, n, k, cs, true);
    REP(i, n) if (isLMS(i)) s1[j++] = i;
    REP(i, n1) sa1[i] = s1[sa1[i]];
    FOR(i, n1, n - 1) sa[i] = -1;
    FORD(i, n1 - 1, 0) {
        j = sa[i];
        sa[i] = -1;
        sa[--bkt[chr(j)]] = j;
    inducesal(t, sa, s, bkt, n, k, cs, false);
    inducesas(t, sa, s, bkt, n, k, cs, true);
    bkt.clear():
    t.clear():
void calc_lcp(void) {
    FOR(i,1,n) rank[sa[i]] = i;
    int h = 0:
    REP(i, n) if (rank[i] < n) {</pre>
        int j = sa[rank[i] + 1];
        while (s[i + h] == s[j + h]) h++;
        lcp[rank[i]] = h;
        if (h > 0) h - -:
    }
}
SuffixArray() {
    n = 0:
    sa = lcp = rank = NULL;
    s = NULL;
SuffixArray(string ss) {
    n = ss.size();
    sa = new int[n + 7];
    lcp = new int [n + 7];
    rank = new int [n + 7];
    s = (unc *)ss.c_str();
    build(s, sa, n + 1, 256, sizeof(char));
    calc_lcp();
}
```

};

```
//Sorted suffices are SA[1] to SA[N]. The values of SA[1], SA[2], ..., SA[
    Nl are 0. 1. ... N - 1
//The\ longest\ common\ prefix\ of\ SA[i]\ and\ SA[i+1]\ is\ LCP[i]
int main(void) {
    string s = "mississippi";
    SuffixArray suffixArray(s);
    FOR(i, 1, 11) printf("%d %s %d\n", suffixArray.sa[i], s.substr(
        suffixArray.sa[i]).c_str(), suffixArray.lcp[i]);
     Manacher
void manacher() {
    memset(p,0,sizeof p);
    int center = 0, right = 0, mi;
    for (int i = 1; i < n; i++) {
        mi = 2 * center - i;
        if (right > i) p[i] = min(right - i, p[mi]);
        while (a[i+(1+p[i])] == a[i-(1+p[i])]) p[i]++;
       //printf("%d:%d\n",i,p[i]);
       if (i + p[i] > right) {
            right = i+p[i];
            center = i:
       }
   }
    DP knuth
http://codeforces.com/blog/entry/8219
Original Recurrence:
   dp[i][j] = min(dp[i][k] + dp[k][j]) + C[i][j] for k = i+1...j-1
Necessary & Sufficient Conditions:
  A[i][j-1] \le A[i][j] \le A[i+1][j]
   with A[i][j] = smallest k that gives optimal answer
Also applicable if the following conditions are met:
  1. C[a][c] + C[b][d] \le C[a][d] + C[b][c] (quadrangle inequality)
  2. C[b][c] <= C[a][d]
                                             (monotonicity)
  for all a <= b <= c <= d
To use:
   Calculate dp[i][i] and A[i][i]
```

```
FOR(len = 1..n-1)
     FOR(i = 1..n-len) {
       j = i + len
       FOR(k = A[i][j-1]..A[i+1][j])
         update(dp[i][j])
     }
// OPTCUT
#include "../template.h"
const int MN = 2011;
int a[MN], dp[MN][MN], C[MN][MN], A[MN][MN];
int n;
void solve() {
    cin >> n; FOR(i,1,n) { cin >> a[i]; a[i] += a[i-1]; }
    FOR(i,1,n) FOR(j,i,n) C[i][j] = a[j] - a[i-1];
    FOR(i,1,n) dp[i][i] = 0, A[i][i] = i;
    FOR (len, 1, n-1)
        FOR(i,1,n-len) {
            int j = i + len;
            dp[i][j] = 2000111000;
            FOR(k, A[i][j-1], A[i+1][j]) {
                int cur = dp[i][k-1] + dp[k][j] + C[i][j];
                if (cur < dp[i][j]) {
                    dp[i][j] = cur;
                    A[i][j] = k;
                }
            }
    cout << dp[1][n] << endl;</pre>
}
     DP divide conquer
```

```
http://codeforces.com/blog/entry/8219
Divide and conquer optimization:
Original Recurrence
  dp[i][j] = min(dp[i-1][k] + C[k][j]) for k < j
Sufficient condition:
```

```
A[i][j]  <= A[i][j+1]
   where A[i][j] = smallest k that gives optimal answer
   // compute i-th row of dp from L to R. optL \leftarrow A[i][L] \leftarrow A[i][R] \leftarrow A[i][R]
   compute(i, L, R, optL, optR)
       1. special case L == R
       2. let M = (L + R) / 2. Calculate dp[i][M] and opt[i][M] using 0(
           optR - optL + 1)
       3. compute(i, L, M-1, optL, opt[i][M])
       4. compute(i, M+1, R, opt[i][M], optR)
const int MN = 4011;
const int inf = 1000111000;
int n, k, cost[MN][MN], dp[811][MN];
inline int getCost(int i, int j) {
    return cost[j][j] - cost[j][i-1] - cost[i-1][j] + cost[i-1][i-1];
void compute(int i, int L, int R, int optL, int optR) {
    if (L > R) return;
    int mid = (L + R) >> 1, savek = optL;
    dp[i][mid] = inf;
    FOR(k,optL,min(mid-1, optR)) {
        int cur = dp[i-1][k] + getCost(k+1, mid);
        if (cur < dp[i][mid]) {
            dp[i][mid] = cur;
            savek = k;
        }
    compute(i, L, mid-1, optL, savek);
    compute(i, mid+1, R, savek, optR);
void solve() {
    cin >> n >> k;
    FOR(i,1,n) FOR(j,1,n) {
        cin >> cost[i][j];
        cost[i][j] = cost[i-1][j] + cost[i][j-1] - cost[i-1][j-1] + cost[i]
            ][i];
   }
```

```
dp[0][0] = 0;
    FOR(i,1,n) dp[0][i] = inf;
    FOR(i,1,k) {
        compute(i, 1, n, 0, n);
    }
    cout << dp[k][n] / 2 << endl;</pre>
}
      Convex Hull
struct Point {
    long long x, y;
    bool operator < (const Point &v) const {
        return x == v.x ? y < v.y : x < v.x;
    }
    long long cross(const Point &p, const Point &q) const {
        return (p.x - x) * (q.y - y) - (p.y - y) * (q.x - x);
    }
};
vector<Point> convexHull(vector<Point> p) {
    sort(p.begin(), p.end());
    int k = 0, n = p.size();
    vector < Point > poly (2 * n);
    for(int i = 0; i < n; ++i) {
        while (k \ge 2 \&\& poly[k-2].cross(poly[k-1], p[i]) < 0) --k;
        poly[k++] = p[i];
    }
    for(int i = n-2, t = k+1; i >= 0; --i) {
        while (k \ge t \&\& poly[k-2].cross(poly[k-1], p[i]) < 0) --k;
        poly[k++] = p[i];
    poly.resize(min(n, max(0, k - 1)));
    return poly;
}
     Geometry's tricks
// hash func overload
struct point{
bool operator == (const point &p) const{ return x == p.x && y == p.y; }
```

```
};
struct hasher {
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
unordered_map <point, int, hasher > hsh;
const double eps = 1e-9;
bool equal(const double &x, const double &y) {
    return fabs(x - y) <
           eps;
struct Point {
    double x, y;
    Point(double x = 0, double y = 0): x(x), y(y) {}
    Point operator + (const Point &p) const {
        return \{x + p.x, y +
                р.у
               };
    Point operator - (const Point &p) const {
        return {x - p.x, y -
                р.у
               };
    Point operator * (double t) const {
        return {x * t, y * t};
    double operator * (const Point &p) const {
        return x * p.x + y *
               р.у;
    double operator % (const Point &p) const {
        return x * p.y - y *
               p.x;
    bool operator == (const Point &p) const {
        return equal(x, p.x)
               && equal(y, p.y);
    double operator ~ () const {
        return sqrt(*this **this);
};
struct Comparator {
    Point a, b;
```

```
Comparator (Point a, Point b): a(a), b(b) {}
    bool operator () (const Point &p, const Point &q) {
        return (p-a) * (b-a) < (q-a) * (b-a);
    }
};
bool between (double x, double 1, double r) {
    if (1 > r) swap(1, r);
    return x + eps > 1 && x - eps < r;
bool inside(Point q, const vector < Point > &p) {
    int n = p.size();
    for (int i = 0; i < n; i++) {
        int j = i + 1 < n ? i + 1 : 0;
        if (fabs((q - p[i]) \% (p[j] - p[i])) > eps) continue;
        if ((q - p[i]) * (p[j] - p[i]) < -eps) continue;
        if ((q - p[j]) * (p[i] - p[j]) < -eps) continue;
        return true:
    }
    int fl = 0;
    for (int i = 0; i < n; i++) {
        int j = i + 1 < n ? i + 1 : 0;
        Point a = p[i], b = p[j];
        if (equal(a.x, b.x)) continue;
        if (a.x > b.x) swap(a, b);
        if (q.x < a.x - eps) continue;
        if (q.x > b.x - eps) continue;
        if ((q - a) \% (b - a) > 0) fl ^= 1;
    }
    return fl;
}
void intersect (Point p, Point q, Point a, Point b, vector < Point >
               &ints) {
    double na = (a - p) \% (q - p), nb = (b - p) \% (q - p);
    if (na * nb > eps) return;
    if (equal(na, nb)) return;
    ints.push_back(a + (b - a) * (na / (na - nb)));
void intersectCircleLine() {
    double r, a, b, c;
    double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
    if (c*c > r*r*(a*a+b*b)+EPS) puts ("no points");
    else if (abs (c*c - r*r*(a*a+b*b)) \leq EPS) {
        puts ("1 point");
        cout << x0 << ', ' << y0 << '\n';
```

```
} else {
        double d = r*r - c*c/(a*a+b*b):
        double mult = sqrt (d / (a*a+b*b));
        double ax, av, bx, by;
        ax = x0 + b * mult;
        bx = x0 - b * mult;
        ay = y0 - a * mult;
        by = y0 + a * mult;
        puts ("2 points");
        cout << ax << ', ' << ay << '\n' << bx << ', ' << by << '\n';
   }
22
    FFT
const double PI = acos(-1.0);
typedef complex < double > Complex;
#define MASK(i) (1LL<<(i))
#define BIT(x,i) (((x) >> (i)) & 1)
#define LOG 17
Complex fftRoot[MASK(LOG)], invRoot[MASK(LOG)];
#define REP(i, n) for (int i = 0, _n = (n); i < _n; i = i + 1)
void initFFT(void) {
    REP(i, MASK(LOG)) {
        double alpha = 2 * PI / MASK(LOG) * i;
        fftRoot[i] = Complex(cos(alpha), sin(alpha));
        invRoot[i] = Complex(cos(-alpha), sin(-alpha));
   }
unsigned roundUp(unsigned v) {
    --v:
   REP(i, 5) v = v >> MASK(i);
    return v + 1:
int reverse(int num, int lg) {
    int res = 0;
    REP(i, lg) if (BIT(num, i)) res \mid= MASK(lg - i - 1);
vector < Complex > fft(vector < Complex > a, bool invert) {
    int n = a.size(), lg = 0;
    while (MASK(lg) < n) lg++;
    vector < Complex > roots(n);
```

```
REP(i, n) roots[i] = invert ? invRoot[MASK(LOG) / n * i] :
                          fftRoot[MASK(LOG) / n * i]:
    REP(i, n) {
        int rev = reverse(i, lg);
        if (i < rev) swap(a[i], a[rev]);</pre>
    }
    for (int len = 2; len <= n; len <<= 1)
        for (int i = 0; i < n; i += len)
            for (int j = 0; j < (len >> 1); j++) {
                 Complex u = a[i + j], v = a[i + j + (len >> 1)] *
                                            roots[n / len * j];
                a[i + j] = u + v;
                a[i + j + (len >> 1)] = u - v;
            }
    if (invert) REP(i, n) a[i] /= n;
    return a:
vector < long long > multiply(const vector < int > &a, const vector < int >
                            &b) {
    int n = roundUp(size(a) + size(b) - 1);
    vector < Complex > pa (n), pb (n);
    for(int i = 0; i < size(a); ++i) pa[i] = a[i];
    for(int i = 0; i < size(b); ++i) pb[i] = b[i];
    pa = fft(pa, false);
    pb = fft(pb, false);
    for(int i = 0; i < n; ++i) pa[i] *= pb[i];
    pa = fft(pa, true);
    vector < long long > res (n);
    for(int i = 0; i < n; ++i) res[i] = round(real(pa[i]));</pre>
}
     NTT
23
const int MODULO = 998244353;
const int ROOT = 3; // Primitive root
void fft(vector<int> &a, bool invert) {
    int n = a.size();
    assert((n & (n - 1)) == 0):
    int lg = __builtin_ctz(n);
    for (int i = 0; i < n; ++i) {
        int i = 0:
```

for (int k = 0; k < lg; ++k) if ((i&1 << k)!=0) j = 1 <<

```
(lg-k-1);
        if (i < j) swap(a[i], a[j]);
   }
    for (int len = 2; len <= n; len *= 2) {
        int wlen = power(ROOT, (MODULO - 1) / len);
        if (invert) wlen = inverse(wlen);
        for (int i = 0: i < n: i += len) {
            int w = 1;
           for (int j = 0; j < len / 2; ++j) {
                int u = a[i + i];
               int v = 1LL * a[i + j + len / 2] * w % MODULO;
               a[i + j] = (u + v) \% MODULO;
               a[i + j + len / 2] = (u - v + MODULO) % MODULO;
               w = 1LL * w * wlen % MODULO;
    if (invert) {
        int mul = inverse(n);
        for (auto &x : a) x = 1LL * x * mul % MODULO;
   }
    998244353 = 119 * 2^23 + 1. Primitive root: 3.
    985661441 = 235 * 2^2 + 1. Primitive root: 3.
    1012924417 = 483 * 2^21 + 1. Primitive root: 5
24 Primitive Root
int generator(int p) {
    vector < int > fact;
    int phi = p-1, n = phi;
    for (int i=2; i*i <= n; ++i) if (n % i==0) {
        fact.push_back(i);
        while (n \% i == 0) n /= i:
    if (n > 1) fact.push_back(n);
    for (int res=2; res<=p; ++res) {
        bool ok = true:
        for (size t i=0: i<fact.size() && ok: ++i)
            ok &= powmod (res, phi / fact[i], p) != 1;
        if (ok) return res;
   }
    return -1;
```

25 Range Prime Counting

}

```
// Primes up to 10^12 can be counted in ~1 second.
const int MAXN = 1000005; // MAXN is the maximum value of sqrt(N) +
bool prime [MAXN];
int prec[MAXN];
vector<int> P:
void init() {
    prime[2] = true;
    for (int i = 3; i < MAXN; i += 2) prime[i] = true;
    for (int i = 3; i*i < MAXN; i += 2) {
        if (prime[i]) {
            for (int j = i*i; j < MAXN; j += i+i) prime[j] = false;</pre>
    }
    for(int i=1; i<MAXN; i++) {</pre>
        if (prime[i]) P.push_back(i);
        prec[i] = prec[i-1] + prime[i];
lint rec(lint N, int K) {
    if (N <= 1 | | K < 0) return 0;
    if (N \le P[K]) return N-1;
    if (N < MAXN && 111 * P[K]*P[K] > N) return N-1 - prec[N] +
                prec[P[K]];
    const int LIM = 250:
    static int memo[LIM*LIM][LIM];
    bool ok = N < LIM*LIM:
    if (ok && memo[N][K]) return memo[N][K];
    lint ret = N/P[K] - rec(N/P[K], K-1) + rec(N, K-1);
    if (ok) memo[N][K] = ret;
    return ret;
lint count_primes(lint N) { //less than or equal to
    if (N < MAXN) return prec[N];</pre>
    int K = prec[(int)sqrt(N) + 1];
    return N-1 - rec(N, K) + prec[P[K]];
}
```

26 Knight's shortest path

```
int KSP(int x,int y) {
   if (x < y) swap(x, y);
   if (x == 1 && y == 0) return 3;
   if (x == 2 && y == 2) return 4;
   int d = x - y;
   if (y > d) return 2*((y-d+2)/3)+d;
   return d-2*((d-y)/4);
}
```

27 Extended Euclid

```
Gia su ket qua la (x0. y0), ho nghiem la (x_0 + k * b / d, y_0 - k * a/d
    )
Phuong trinh ax + by = d co nghiem khi va chi khi d chia het cho gcd(a, b
    )
a x + b y = gcd(a, b)
int extgcd(int a, int b, int &x, int &y) {
    int g = a; x = 1; y = 0;
    if (b != 0) g = extgcd(b, a % b, y, x), y -= (a / b) * x;
    return g;
}
```

28 Factorial Mod

```
int factmod (int n, int p) { // n!, excluding p^k of course
  int res = 1;
  while (n > 1) {
    res = (res * ((n/p) % 2 ? p-1 : 1)) % p;
    for (int i=2; i<=n%p; ++i)
        res = (res * i) % p;
    n /= p;
  }
  return res % p;
}</pre>
```

29 Sqrt Mod

```
| // Jacobi Symbol (m/n), m,n >= 0 and n is odd
```

```
#define NEGPOW(e) ((e) % 2 ? -1 : 1)
int jacobi(int a. int m) {
    if (a == 0) return m == 1 ? 1 : 0;
    if (a \% 2) return NEGPOW((a-1)*(m-1)/4)*jacobi(m\%a, a);
    else return NEGPOW((m*m-1)/8)*jacobi(a/2, m);
}
int invMod(int a, int m) {
    int x, y;
    if (extgcd(a, m, x, y) == 1) return (x + m) \% m;
                                          return 0; // unsolvable
}
// No solution when: n(p-1)/2 = -1 \mod p
int sqrtMod(int n, int p) { //find x: x2 = n \pmod{p} p \ is \ prime
    int S, Q, W, i, m = invMod(n, p);
    for (Q = p - 1, S = 0; Q \% 2 == 0; Q /= 2, ++S);
    do { W = rand() % p; } while (W == 0 || jacobi(W, p) != -1);
    for (int R = powMod(n, (Q+1)/2, p), V = powMod(W, Q, p); ;) {
        int z = R * R * m \% p;
        for (i = 0; i < S && z % p != 1; z *= z, ++i);
        if (i == 0) return R;
        R = (R * powMod(V, 1 << (S-i-1), p)) % p;
    }
}
int powMod (int a, int b, int p) {
    int res = 1:
    while (b)
        if (b & 1)
            res = int (res * 111 * a % p), --b;
        else
            a = int (a * 111 * a % p), b >>= 1;
    return res;
}
      Interval line
// template Interval line Min
#define mid ((lo + hi)>>1)
class Line {
public:
    11 a, b;
```

Line (ll x = cmax, ll y = cmax) {

```
a = x, b = y;
   }
   ll get(int x) {
        return 1LL * val[x] * a + b;
   }
};
const Line oo = Line(cmax, cmax);
class ILTree {
    int m;
public:
    Line t[maxn*4];
    ILTree(int last = 200000) {
        m = last:
        init(1,1,m);
   }
    void init(int i, int lo, int hi) {
        t[i] = oo;
        if (lo == hi) return;
        init(i * 2, lo, mid);
        init(i * 2 + 1, mid + 1, hi);
   }
    void update(int i, int lo, int hi, int l, int r, Line d) {
        if (1 > hi || r < lo) return:
        if (lo >= 1 && hi <= r) {
            // t[i] hoan toan nam duoi d
            if (t[i].get(lo) <= d.get(lo) && t[i].get(hi) <= d.get(hi))
                return:
            //t[i] hoan to an nam trend thi cap nhap t[i] = d
            if (t[i].get(lo) >= d.get(lo) && t[i].get(hi) >= d.get(hi)) {
                t[i] = d;
                return;
            }
            //nua dau cua d tot hon
            if (t[i].get(lo) >= d.get(lo) && t[i].get(mid) >= d.get(mid))
                update(i * 2 + 1, mid + 1, hi, l, r, t[i]);
                t[i] = d:
```

```
return;
            }
            // nua dau cua t[i] tot hon
            if (t[i].get(lo) <= d.get(lo) && t[i].get(mid) <= d.get(mid))
                update(i * 2 + 1, mid + 1, hi, l, r, d);
                return;
            }
            // nua sau cua d tot hon
            if (t[i].get(mid + 1) >= d.get(mid + 1) && t[i].get(hi) >= d.
                get(hi)) {
                update(i * 2, lo, mid, l, r, t[i]);
                t[i] = d;
                return;
            }
            // nua sau cua t[i] tot hon
            if (t[i].get(mid + 1) <= d.get(mid + 1) && t[i].get(hi) <= d.
                get(hi)) {
                update(i * 2, lo, mid, l, r, d);
                return:
            }
        }
        update(i * 2, lo, mid, l, r, d);
        update(i * 2 + 1, mid + 1, hi, l, r, d);
    }
    ll get(int i, int lo, int hi, int pos) {
        if (lo > pos || hi < pos) return llmax;
        ll res = t[i].get(pos);
        if (lo == hi) return res;
        res = min(res, get(i * 2, lo, mid, pos));
        res = min(res, get(i * 2 + 1, mid + 1, hi, pos));
        return res;
    }
};
#undef mid
#undef oo
     BIT 2D
class query{
public:
```

```
int x, y, id, type;
         void next_x(){
                  if (type == 1) x += x \& -x;
                 else x - = x & -x;
        }
         void next_y(){
                 if (type == 1) y += y \& - y;
                 else y - = y & - y;
        }
};
const int maxqsize=100000;
vector <query> q;
deque <int> start[maxqsize+1];
vector <query> temp[n+1];
class node{
public:
        int value, y;
        int next, prev;
        node(int y){
                 this ->y=y;
                 value=0;
                 next = prev = -1;
        }
};
vector <node> ft[n+1];
int done[n+1];
int pos[n+1];
int main() {
        // your code goes here
        for(query a: q){
                 while ((a.x) & & (a.x <= n)) {
                          temp[a.x].pb(a);
                          a.next_x();
                 }
        }
        FOR(i, 1, n){
                 for(query a: temp[i]){
                          while((a.y)&&(a.y<=n)){
                                   if (done [a.y]!=i) {
                                            done[a.y]=i;
                                            pos[a.y]=ft[i].size();
                                            ft[i].pb(node(a.y));
                                   }
                                   a.next_y();
```

```
}
}
for(node &a: ft[i]){
    int yy=a.y+(a.y&(-a.y));
    if(yy<=n){
        if(done[yy]==i) a.next=pos[yy];
    }
    yy=a.y-(a.y&(-a.y));
    if(done[yy]==i) a.prev=pos[yy];
}

for(query a: temp[i]){
    if(a.type==1)//up
        start[a.id].pb(pos[a.y]);
    else start[a.id].push_front(pos[a.y]);
}
return 0;
}</pre>
```

32 Heavy-Light Decomposition

```
if (v != parent[u]) {
                        if (mxVtx == -1 || nChild[v] > nChild[mxVtx]) {
                                mxVtx = v:
               }
       }
    //New tim ra dinh con dac biet (w khong phai la dinh la) thi di chuyen
         den dinh do
       if (mxVtx > -1)
               hld(mxVtx);
    // Sau khi di het mot chuoi thi tang nChain len va bat dau mot chuoi
       for (int i = 0; i < adj[u].size(); i++) {
               int v = adj[u][i];
                if (v != parent[u] && v != mxVtx) {
                        nChain++;
                        hld(v);
void update(int u, int a) {
    // uchain chuoi hien tai cua u
    // achain chuoi hien tai cua a
    int uchain = chainInd[u], achain = chainInd[a];
    while (1) {
        // Neu u va a cung nam tren mot chuoi thi update doan tu u den a
           vaket thuc
         if (uchain == achain) {
               updateIntervalTree(..., posInBase[a], posInBase[u], ...);
              break:
        // Neu u va a khong nam tren cung mot chuoi thi update doan tu u
            den dinh dau cua chuoi hien tai
          updateIntervalTree(..., posInBase[chainHead[uchain]], posInBase[
             u], ...);
        // Nhay len dinh cha cua dinh dau hien tai
         u = parent[chainHead[uchain]];
          uchain = chainInd[u]:
```

```
}
      Splay Tree
/**
 * Problem's query: insert a[i] after a[j]
struct node {
    node *par, *left, *right;
    int value, cnt;
}:
int n, Q, tree[N];
node *root, *nilT;
void calc(node *x) {
     if (x != nilT)
         x \rightarrow cnt = x \rightarrow left \rightarrow cnt + x \rightarrow right \rightarrow cnt + 1;
}
node *FindPosition(node *cur, int pos) {
    node *x = cur;
    while (x != nilT) {
         int ord = x \rightarrow left \rightarrow cnt + 1;
         if (ord == pos)
              return x;
         if (ord > pos)
              x = x - > left;
         else
              x = x - right, pos -= ord;
    }
    return nilT;
void SetL(node *parent, node *child) {
    if (child != nilT)
         child->par = parent;
    if (parent != nilT)
         parent -> left = child;
}
void SetR(node *parent, node *child) {
    if (child != nilT)
         child->par = parent;
    if (parent != nilT)
         parent -> right = child;
```

```
void Uptree(node *x) {
    if (x == root)
        return;
    node *y = x - par, *z = y - par;
    if (y->left == x) {
        SetL(y, x->right);
        SetR(x, y);
   } else {
        SetR(v, x->left);
        SetL(x, y);
   }
    if (z == nilT)
        root = x, x->par = nilT;
    else if (z \rightarrow left == y)
        SetL(z, x);
    else
        SetR(z, x);
    calc(y);
    calc(x):
void Splay(node *x) {
    while (1) {
        node *y = x - > par;
        if (y == nilT)
             break;
        node *z = y - par;
        if (z != nilT) {
             if ((y == z \rightarrow left) == (y \rightarrow left == x))
                 Uptree(y);
             else
                 Uptree(x);
        Uptree(x);
    }
    root = x;
void Split(node *r, int i, node *&r1, node *&r2) {
    if (i == 0) {
        r1 = nilT, r2 = r;
        return;
    node *x = FindPosition(r, i);
    Splay(x);
```

```
r2 = x - > right;
    r1 = x:
    r1->right = nilT;
    calc(r1), calc(r2);
node *Join(node *r1, node *r2) {
    if (r1 == nilT)
        return r2;
    while (r1->right != nilT)
        r1 = r1 - > right;
    Splay(r1);
    SetR(r1, r2);
    calc(r1);
    return r1;
void Insert(int i. int val) {
    node *x = new node:
    x \rightarrow value = val;
    x - par = nilT;
    node *r1, *r2;
    Split(root, i - 1, r1, r2);
    SetL(x, r1), SetR(x, r2);
    calc(x);
    root = x;
}
void Delete(int i) {
    node *x = FindPosition(root, i);
    Splay(x);
    node *r1 = x - > left, *r2 = x - > right;
    r1 - par = nilT; r2 - par = nilT;
    delete x;
    root = Join(r1, r2);
}
vector<int> arr;
void GetArray(node *x) {
    if (x == nilT)
        return:
    GetArray(x->left);
    arr.push_back(x->value);
    GetArray(x->right);
    delete x;
}
```

```
int main() {
    nilT = new node:
    nilT \rightarrow value = nilT \rightarrow cnt = 0:
    nilT->left = nilT->right = nilT->par = nilT;
    root = new node;
    root -> value = root -> cnt = 1:
    root -> left = root -> right = root -> par = nilT;
    cin >> n >> Q;
    for (int i = 2; i \le n; i++)
        Insert(i, i);
    for (int i = 1; i <= Q; i++) {
        cin >> u >> v;
        int val = FindPosition(root, u) -> value;
        Delete(u);
        Insert(v, val);
    GetArray(root);
    int ans = 0;
    for (int i = 0; i < n; i++) {
        int F = 1:
        for (int x = arr[i]; x; x -= x & (-x))
            F = max(F, tree[x] + 1);
        ans = max(ans, F);
        for (int x = arr[i]; x <= n; x += x & (-x))
            tree[x] = max(tree[x], F);
   }
    cout << n - ans << endl;
/**
   Geometry - Tungluu18
   Duong tron di qua 3 diem cho truoc Let A = (0, 0) centers are Cy(Bx^2+
Bu^{2}) – Bu(Cx^{2} + Cu^{2})/D
   and (Bx(Cx^2 + Cy^2) - Cx(Bx^2 + By^2))/D where D = 2(BxCy - ByCx).
Diem trong tam giac
bool isInside(const Vector &P) {
        Vector a = C - A, b = B - A, c = P - A;
        T \text{ under } = a.x*b.y - b.x*a.y;
        T u = (c.x*b.y-c.y*b.x);
```

```
T v = (a.x*c.y-a.y*c.x);
                                                                                      return (x1, y1) + (a*e, b*e) + h*(-b, a),
        return u >= 0 && v >= 0 && u+v <= under || u<=0
                                                                                              (x1, v1) + (a*e, b*e) - h*(-b, a):
        && v \le 0 && u+v >= under:
                                                                                 Template
} //remove equalities if not want the boundary
                                                                              double eps = 1e-8;
Pick's theorem (So diem trong da giac co dinh nguyen)
        T = A - B/2 + 1
                                                                              struct point2d {
        where
                                                                                      double x, y;
                A is the area of a lattice polygon,
                                                                                      point2d operator - (point2d v) { return (point2d){x-v.x, y-v.y}; }
                I is number of lattice points inside it,
                                                                                      point2d operator+(point2d v) { return (point2d){x+v.x, y+v.y}; }
                B is number of lattice points on the boundary.
                                                                              };
                Number of lattice points minus one on a line segment
                from (0, 0) and (x, y) is gcd(x, y).
                                                                              struct point3d {
   Tich chap 3d: a \times b = (aybz - azby, azbx - axbx, axby - aybx)
                                                                                      point3d operator-(point3d v) { return (point3d)\{x-v.x, y-v.y, z-v.\}
   Distance from line AB to P (for any dimension): |(A-P)x(B-P)|/(A-P)
                                                                                          z}: }
B
                                                                              };
   Khoảng cách từ điểm đến đoan thẳng if (dot(B-A, P-A) < 0) return
dist(A,P); if (dot(A-B, P-B) < 0) return dist(B,P); return fabs(cross(P,A,B)ouble trip(point3d a, point3d b, point3d c) {
                                                                                      return a.x * ( b.y * c.z - b.z * c.y )
/ dist(A,B)); dot - tích vô hướng, cross - tích chập
                                                                                              - a.y * ( b.x * c.z - b.z * c.x )
   Hình chiếu: Hình chiếu của C trên đường thẳng AB [dot(AB, AC) / dot(AB,
                                                                                              + a.z * ( b.x * c.y - b.y * c.x );
AB) AB
   Đường tròn cắt đường tròn
                                                                              double len(point2d a) { return sqrt(a.x*a.x+a.y*a.y); }
// (x1, y1, r1) and (x2, y2, r2)
                                                                              double len(point3d a) { return sqrt(a.x*a.x*a.y*a.y*a.z*a.z); }
        a = x2 - x1;
        b = y2 - y1;
                                                                              point3d cross(point3d a, point3d b) {
        c = [(r1^2 - x1^2 - y1^2) - (r2^2 - x2^2 - y2^2)] / 2;
                                                                                      return (point3d){
        d = sqrt(a^2 + b^2);
                                                                                              a.y*b.z - a.z*b.y,
                                                                                              a.z*b.x - a.x*b.z
        if not |r1 - r2| <= d <= |r1 + r2|, return "no solution"
                                                                                              a.x*b.y - a.y*b.x
        if d == 0, circles are concentric, a special case
                                                                                      }:
   Đường tròn cắt đường thẳng
// (x1, y1, r1) and ax + by = c
                                                                              double dot(point3d a, point3d b) { return a.x*b.x + a.y*b.y + a.z*b.z; }
                                                                              double dot(point2d a, point2d b) { return a.x*b.x + a.y*b.y; }
Normalize line:
        d = sqrt(a^2+b^2)
                                                                              double cross(point2d a, point2d b) {
        a /= d:
                                                                                      return a.x*b.v - a.v*b.x:
        b /= d;
        c /= d;
                                                                              point2d rotate(point2d c, double angle) {
        e = c - a*x1 - b*v1:
                                                                                      return (point2d) {
        h = sqrt(r1^2 - e^2); // check if r1 < e for circle-line test
                                                                                              c.x*cos(angle)-c.y*sin(angle),
```

```
c.x*sin(angle)+c.y*cos(angle)
        }:
}
// distance AB to C
double linePointDist(point2d A, point2d B, point2d C, bool isSegment) {
        double dist = cross(B-A,C-A) / len(B-A);
        if(isSegment){
                if( dot(C-B,B-A) > eps) return len(B-C);
                if (dot(C-A,A-B) > eps) return len(A-C);
        return fabs(dist);
}
struct Line { double A,B,C; }; //Ax + By = C
Line makeline(point2d a1, point2d a2)
        Line ret = (Line) \{ a2.y-a1.y, a1.x-a2.x \};
        ret.C = ret.A * a1.x + ret.B * a1.y;
        return ret:
}
double dist(Line 1, point2d p)
        return fabs(p.x * 1.A + p.y * 1.B - 1.C)/sqrt(1.A*1.A+1.B*1.B);;
Line rot90 (Line 1, point2d p)
{
        Line ret = (Line) { -1.B, 1.A };
        ret.C = ret.A * p.x + ret.B * p.y;
        return ret;
point2d intersect(Line 11, Line 12)
        double det = 11.A*12.B - 12.A*11.B;
        if(fabs(det) < eps) det=0; //zero means parallel
        return (point2d) { (12.B*11.C - 11.B*12.C)/det,
        (11.A*12.C - 12.A*11.C)/det};
}
//for segment segment intersection, check additionally
//min(x1,x2) <= x <= max(x1,x2)
//min(y1,y2) <= x <= max(y1,y2)
bool segmentsIntersect( point2d A, point2d B, point2d C, point2d E)
```

```
{
    point2d in = intersect( makeline(A,B), makeline(C,E) );
    return linePointDist(A,B,in,true) < eps
    && linePointDist(C,E,in,true) < eps;
}

// get a line passing between two points
Line getmidline(point2d a, point2d b)
{
    point2d mid(a+b); mid.x/=2; mid.y/=2;
    return rot90( makeline(a,b), mid );
}

//reflect a point into it's "mirror" with repect to a line
point2d reflectPoint(Line 1, point2d p)
{
    Line r = rot90(1, p);
    point2d Y=intersect(1,r);
    return Y-(p-Y);
}</pre>
```