

Algebraic Statistics

P1: Intro to Algebraic Geometry

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Solution: $f \in I \iff r = 0$ for some monomial ordering.

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Look at: Groebner basis

Definition: Fix a monomial order. Suppose I is an ideal such that $\langle LT(I) \rangle = \langle LT(g_1), LT(g_2), \dots, LT(g_t) \rangle$. Then, $G = \{g_1, g_2, \dots, g_t\}$ is a Groebner basis of I .

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Hilbert Basis Theorem: Every ideal $I \subset \mathbb{K}[\vec{x}]$ is finitely generated.

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Then, every nonzero ideal $I \subset \mathbb{K}[\vec{x}]$ has a finite Groebner basis.

Importance: Given a Groebner basis G of ideal I and ordering:

1. The division algorithm gives a unique remainder r .

2. $r = \text{Rem}(f, G) = 0 \iff f \in I$.

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How to find a Groebner basis? Can use Buchberger's algorithm.
First, need to introduce S-polynomials.

S-polynomials:

$$S(f_1, f_2) = x^\gamma \left[\frac{f_1}{LT(f_1)} - \frac{f_2}{LT(f_2)} \right], \text{ where } x^\gamma = LCM(LM(f_1), LM(f_2))$$

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Example: $f_1 = y - x^2$, $f_2 = z - x^3$, $y \underset{lex}{>} z \underset{lex}{>} x$

$$\begin{aligned} S(f_1, f_2) &= yz \left[\frac{y - x^2}{y} - \frac{z - x^3}{z} \right] \\ &= -zx^2 + yx^3 \end{aligned}$$

Buchberger's Criterion:

$G = \{g_1, g_2, \dots, g_t\}$ is a Groebner basis for I
 $\iff \text{Rem}(S(g_i, g_j), G) = 0$, for all $i \neq j$.

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Note: $\text{Rem}(S', G) \neq 0$ under $x \underset{\text{lex}}{>} y \underset{\text{lex}}{>} z$.

Thus, we have a way to determine if $\{g_1, \dots, g_t\}$ is a Groebner basis.

But **how do we find a Groebner basis?**

Buchberger's Algorithm

Input: $F=(f_1, f_2, \dots, f_s)$

Output: $G=(g_1, g_2, \dots, g_t)$

$G=F$

Repeat

$G'=G$

 For each pair $\{i, j\}$, $i \neq j$ in G' , do

$g=\text{Rem}(S(i, j), G)$

 If $S \neq 0$

 then $G=G \cup \{g\}$

Until $G=G'$

Sample steps:

| G | Rem(S(..), G) |
|--------------------------|----------------------|
| $\{g_1, g_2\}$ | g_3 |
| $\{g_1, g_2, g_3\}$ | g_4 |
| $\{g_1, g_2, g_3, g_4\}$ | 0 |

Problem: Groebner basis is not unique and contains redundancies.
e.g. $\{y - x^2, z - x^3\}$ for $\langle y - x^2, z - x^3 \rangle$ w.r.t. $y \underset{lex}{>} z \underset{lex}{>} x$

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Importance: Reduced Groebner basis is unique!

Thank you! :) :)