Dynamic Programming

Chapter 5: Markov Decision Processes

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Topics

- Introduction to Markov decision processes
- Lifetime value
- Optimality
- Value function iteration
- Howard policy iteration
- Optimistic policy iteration
- Applications
- Refactoring dynamic programs

Markov Decision Processes (MDPs)

A class of well-behaved dynamic programs

- broad enough to encompass many economic applications
- includes optimal stopping problems as a special case
- admits a clean, powerful theory

A foundation stone for

- reinforcement learning
- artificial intelligence
- operations research, etc.

States and Actions

We take as given

- 1. a finite set X called the **state space** and
- 2. a finite set A called the action space

We study a controller who, at each integer $t\geqslant 0$

- 1. observes the current state $X_t \in X$
- 2. responds with an action $A_t \in A$

Her aim is to maximize

$$\mathbb{E}\sum_{t>0}\beta^t r(X_t, A_t) \quad \text{ given } X_0 = x_0$$

Actions restricted by a **feasible correspondence** Γ

- $\Gamma(x)$ is a nonempty subset of A for each $x \in X$
- interpretation: $\Gamma(x) = \text{actions available in state } x$

Given Γ , we set

$$\mathsf{G} = \{(x, a) \in \mathsf{X} \times \mathsf{A} : a \in \Gamma(x)\}\$$

called the set of feasible state-action pairs

Reward r(x, a) is received at feasible state-action pair (x, a)

A stochastic kernel from G to X is a map $P \colon G \times X \to \mathbb{R}_+$ satisfying

$$\sum_{x' \in \mathsf{X}} P(x,a,x') = 1 \quad \text{ for all } (x,a) \text{ in } \mathsf{G}$$

Interpretation

• next period state x' is drawn from distribution $P(x,a,\cdot)$

Now let's put it all together:

Given X and A, an MDP is a tuple $\mathfrak{M}=(\Gamma,\beta,r,P)$ where

- 1. Γ is a nonempty correspondence from $X \to A$
- 2. β is a constant in (0,1)
- 3. r is a function from G to \mathbb{R}
- 4. P is a stochastic kernel from G to X

In what follows,

- β is called the **discount factor**
- r is called the reward function

MDP dynamics:

The Bellman equation is

$$v(x) = \max_{a \in \Gamma(x)} \left\{ r(x, a) + \beta \sum_{x' \in \mathsf{X}} v(x') P(x, a, x') \right\}$$

Key idea: reduce infinite horizon problem to a two period problem

— "Bellman's principle of optimality"

In the two period problem, the controller trades off

- 1. current rewards and
- 2. expected discounted value from future states

Example. Rust (1987 ECMA) examined an engine replacement problem for a bus workshop

At each t, the supervisor decides whether or not to replace a bus engine

Replacement is costly but delaying risks unexpected failure

We consider a version with binary action A_t

- $A_t = 1 \implies$ state resets to some fixed renewal state \bar{x}
 - e.g., mileage resets to zero when an engine is replaced
- $A_t = 0 \implies$ state updates according to $Q \in \mathcal{M}(\mathbb{R}^{\mathsf{X}})$
 - · e.g., mileage increases stochastically when not replaced

Current reward is r(x, a) at state x and action a

• $x \in X = \text{some finite set (e.g., possible milage values)}$

The discount factor is $\beta \in (0,1)$

The Bellman equation is

$$v(x) = \max \left\{ r(x,0) + \beta \sum_{x' \in \mathsf{X}} v(x') Q(x,x'), \ r(x,1) + \beta v(\bar{x}), \right\}$$

 $= \max \{ \text{value of action 0}, \text{ value of action 1} \}$

Ex. Show that the renewal problem can be framed as an MDP

Proof: We set

- X = as above
- $A = \{0, 1\}$
- $\Gamma(x) = A$ for all $x \in X$
- r = as above

We define

$$P(x, a, x') := a\mathbb{1}\{x' = \bar{x}\} + (1 - a)Q(x, x')$$

You can check that P is a stochastic kernel from G to X

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You can check that P is a stochastic kernel from G to X

By construction, $\mathfrak{M}=(\Gamma,\beta,r,P)$ is an MDP Moreover.

$$\begin{split} v(x) &= \max_{a \in \Gamma(x)} \left\{ r(x,a) + \beta \sum_{x' \in \mathsf{X}} v(x') P(x,a,x') \right\} \\ &= \max_{a \in \{0,1\}} \left\{ r(x,a) + \beta \left[av(\bar{x}) + (1-a) \sum_{x' \in \mathsf{X}} v(x') Q(x,x') \right] \right\} \\ &= \max \left\{ r(x,1) + \beta v(\bar{x}), \ r(x,0) + \beta \sum_{x' \in \mathsf{X}} v(x') Q(x,x') \right\} \end{split}$$

We have recovered the renewal Bellman equation on slide 11

Example. Cake eating, where

$$W_{t+1} = R(W_t - C_t)$$
 $(t = 0, 1, ...)$

- ullet investing d dollars today returns Rd next period
- $C_t, W_t \geqslant 0$ are current consumption and wealth
- (W_t) restricted to finite grid $\mathsf{W} \subset \mathbb{R}_+$

The agent seeks to maximize $\mathbb{E} \sum_{t\geqslant 0} \beta^t u(C_t)$ given $W_0=w$

Bellman equation:

$$v(w) = \max_{0 \leqslant w' \leqslant Rw} \left\{ u \left(w - \frac{w'}{R} \right) + \beta v(w') \right\}$$

This model can be framed as an MDP with W as the state space

- action is $A_t = W_{t+1} =$ savings
- action space is also W
- feasible correspondence is

$$\Gamma(w) = \{ a \in \mathsf{W} : a \leqslant Rw \}$$

• since $A_t = R(W_t - C_t)$, current reward is

$$r(w, a) = u(w - a/R)$$
 $(a \in \Gamma(w))$

• stochastic kernel is $P(w, a, w') = \mathbb{1}\{w' = a\}$

Example. Discrete time optimal stopping problems can be framed as MDPs

See the text for details

This is important from a theoretical perspective

illustrates the generality of MDPs

But expressing optimal stopping problems as an MDP requires an extra state variable

• status $S_t = 1$ {already stopped}

This is why we treated them separately (Ch. 4)

Policies

Time travel is forbidden

• A_t cannot depend on X_{t+j} for any $j\geqslant 1$

While A_t can depend on the history

$$H_t := (A_0, X_0, \dots, A_{t-1}, X_{t-1}, X_t),$$

for MDPs, conditioning on X_t is sufficient

In other words,

• A_t depends only on X_t

Below we write $A_t = \sigma(X_t)$ and call σ a policy function

A feasible policy is a

$$\sigma \in \mathsf{A}^\mathsf{X}$$
 such that $\sigma(x) \in \Gamma(x)$ for all $x \in \mathsf{X}$

• Let $\Sigma :=$ the set of all feasible policies

Choosing
$$\sigma \in \Sigma \implies$$

respond to state X_t with action $A_t := \sigma(X_t)$ at $\underline{\mathsf{all}}\ t \geqslant 0$

If $A_t := \sigma(X_t)$ at all $t \geqslant 0$, then

$$X_{t+1} \sim P(X_t, \sigma(X_t), \cdot)$$
 for all $t \geqslant 0$

Thus, X_t is P_{σ} -Markov for

$$P_{\sigma}(x, x') := P(x, \sigma(x), x') \qquad (x, x' \in \mathsf{X})$$

- Fixing a policy "closes the loop" in the state dynamics
- Solving an MDP means choosing a Markov chain!

Rewards

Under the policy σ , rewards at x given by $r(x, \sigma(x))$

Let

- $r_{\sigma}(x) := r(x, \sigma(x))$
- $\mathbb{E}_x := \mathbb{E}[\cdot \mid X_0 = x]$

Now

$$\mathbb{E}_x r(X_t, A_t) = \mathbb{E}_x r_{\sigma}(X_t) = \sum_{x'} r_{\sigma}(x') P_{\sigma}^t(x, x') = (P_{\sigma}^t r_{\sigma})(x)$$

The **lifetime value of** σ starting from x is

$$v_{\sigma}(x) := \mathbb{E}_{x} \sum_{t \geqslant 0} \beta^{t} r_{\sigma}(X_{t})$$
$$= \sum_{t \geqslant 0} \mathbb{E}_{x} \left[\beta^{t} r_{\sigma}(X_{t}) \right]$$
$$= \sum_{t \geqslant 0} \beta^{t} (P_{\sigma}^{t} r_{\sigma})(x)$$

Since $\beta < 1$, the spectral radius of βP is < 1, so

$$v_{\sigma} = \sum_{t \ge 0} (\beta P_{\sigma})^t r_{\sigma} = (I - \beta P_{\sigma})^{-1} r_{\sigma}$$

Policy Operators

How should we compute v_{σ} given σ ?

We saw above that

$$v_{\sigma} = (I - \beta P_{\sigma})^{-1} r_{\sigma}$$

- Computationally helpful when X is small
- Problematic for large problems

Example. If
$$|\mathsf{X}| = 10^6$$
, then $I - \beta \, P_\sigma$ is $10^6 \times 10^6$

Matrices of this size are difficult invert—or even store in memory

Another way to compute v_{σ} : use the **policy operator**

$$(T_{\sigma} v)(x) = r(x, \sigma(x)) + \beta \sum_{x' \in \mathsf{X}} v(x') P(x, \sigma(x), x')$$

- Defined at all $v \in \mathbb{R}^X$
- ullet Analogous to T_σ for the optimal stopping problem

In vector notation, we can write

$$T_{\sigma} v = r_{\sigma} + \beta P_{\sigma} v$$

• T_{σ} is order-preserving on \mathbb{R}^{X} — why?

Ex. Show that T_{σ} is a contraction of modulus β on \mathbb{R}^{X}

For any v, w in \mathbb{R}^{\times} we have

$$|T_{\sigma}v - T_{\sigma}w| = \beta |P_{\sigma}v - P_{\sigma}w|$$

$$= \beta |P_{\sigma}(v - w)|$$

$$\leq \beta P_{\sigma} |v - w|$$

$$\leq \beta P_{\sigma} ||v - w||_{\infty} \mathbb{1}$$

$$= \beta ||v - w||_{\infty} \mathbb{1}$$

Now use $|a| \leq |b|$ implies $||a||_{\infty} \leq ||b||_{\infty}$

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Ex. Show that v_{σ} is the unique fixed point of T_{σ} in \mathbb{R}^{X}

Proof: Since $\beta < 1$, we have

$$v = T_{\sigma} v \iff v = r_{\sigma} + \beta P_{\sigma} v$$
 $\iff v = (I - \beta P_{\sigma})^{-1} r_{\sigma}$
 $\iff v = v_{\sigma}$

Hence

v is a fixed point of $T_{\sigma} \iff v = v_{\sigma}$

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Greedy Policies

Fix $v \in \mathbb{R}^{X}$

A policy σ is called v-greedy if

$$\sigma(x) \in \operatorname*{argmax}_{a \in \Gamma(x)} \left\{ r(x, a) + \beta \sum_{x'} v(x') P(x, a, x') \right\}$$

for all $x \in X$

Ex. Prove: at least one v-greedy policy exists in Σ

<u>Proof</u>: Immediate because $\Gamma(x)$ is finite and nonempty at all x

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The Bellman Operator

Recall: the Bellman equation is

$$v(x) = \max_{a \in \Gamma(x)} \left\{ r(x, a) + \beta \sum_{x'} v(x') P(x, a, x') \right\}$$

The **Bellman operator** for $\mathfrak M$ is the self-map on $\mathbb R^X$ defined by

$$(Tv)(x) = \max_{a \in \Gamma(x)} \left\{ r(x, a) + \beta \sum_{x'} v(x') P(x, a, x') \right\}$$

By construction, $Tv = v \iff v$ satisfies the Bellman equation

Ex. Prove: σ is v-greedy if and only if

$$T_{\sigma} v = Tv$$

Proof: σ is v-greedy if and only if

$$\sigma(x) \in \operatorname*{argmax}_{a \in \Gamma(x)} \left\{ r(x, a) + \beta \sum_{x'} v(x') P(x, a, x') \right\} \qquad \forall \, x \in \mathsf{X}$$

This is equivalent to

$$r(x, \sigma(x)) + \beta \sum_{x'} v(x') P(x, \sigma(x), x') = (Tv)(x) \quad \forall x \in X$$

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Ex. Prove that, for all $v \in \mathbb{R}^{X}$,

$$Tv = \bigvee_{\sigma \in \Sigma} T_{\sigma} v$$

Proof:

Fix $v \in \mathbb{R}^X$

For any v-greedy $\sigma \in \Sigma$, we have $T_{\sigma} v = Tv$

Also, for all $\sigma \in \Sigma$ and $x \in X$,

$$(T_{\sigma}v)(x) = r(x,\sigma(x)) + \beta \sum_{x'} v(x')P(x,\sigma(x),x') \leqslant (Tv)(x)$$

$$\therefore Tv = \bigvee_{\sigma} T_{\sigma} v$$

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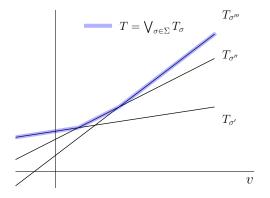


Figure: Visualization in one dimension

Ex. Prove: T is a contraction of modulus β on \mathbb{R}^X

Proof: Recall that, for finite S and any $f,g \in \mathbb{R}^S$,

$$|\max_{s \in S} f(s) - \max_{s \in S} g(s)| \leqslant \max_{s \in S} |f(s) - g(s)|$$

Hence, for any v, w in \mathbb{R}^X , we have

$$|(Tv)(x) - (Tw)(x)| = \left| \max_{\sigma \in \Sigma} (T_{\sigma} v)(x) - \max_{\sigma \in \Sigma} (T_{\sigma} w)(x) \right|$$

$$\leq \max_{\sigma \in \Sigma} |(T_{\sigma} v)(x) - (T_{\sigma} w)(x)|$$

$$\therefore |Tv - Tw| \leqslant \max_{\sigma \in \Sigma} |T_{\sigma} v - T_{\sigma} w| \leqslant \beta ||v - w||_{\infty}$$

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$$\leq \max_{\sigma \in \Sigma} |(T_{\sigma} v)(x) - (T_{\sigma} w)(x)|$$

$$\therefore |Tv - Tw| \leqslant \max_{\sigma \in \Sigma} |T_{\sigma} v - T_{\sigma} w| \leqslant \beta ||v - w||_{\infty}$$

Optimality

The value function is defined by $v^* := \vee_{\sigma \in \Sigma} v_{\sigma}$ More explicitly,

$$v^*(x) := \max_{\sigma \in \Sigma} v_{\sigma}(x) \qquad (x \in \mathsf{X})$$

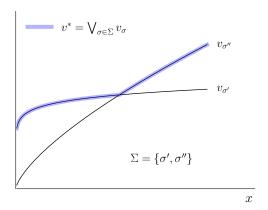
Thus, $v^*(x) = \underline{\text{maximal lifetime value}}$ from state x

A policy $\sigma \in \Sigma$ is called **optimal** if

$$v_{\sigma} = v^*$$

Thus, σ is optimal \iff lifetime value is maximal at each state

Example. Consider



In this case there is \underline{no} optimal policy

Indeed, v^* differs from both $v_{\sigma'}$ and $v_{\sigma''}$

Hence no σ satisfies $v^*=v_\sigma$

Below we show that such an outcome is **not** possible for MDPs

In other words, an optimal policy always exists

This leads to our next slide. . .

Theorem. If $\mathcal M$ is an MDP with Bellman operator T and value function v^* , then

- 1. v^* is the unique fixed point of T in \mathbb{R}^X
- 2. A feasible policy is optimal if and only it is v^* -greedy
- 3. At least one optimal policy exists

Remark: Point (2) is called Bellman's principle of optimality

We prove the theorem below through some exercises

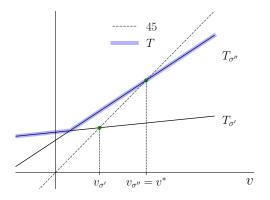


Figure: Illustration of optimality for MDPs

Ex. Show that Bellman's principle of optimality implies existence of an optimal policy

Proof

Let Bellman's principle of optimality hold

For each $v \in \mathbb{R}^{X}$, a v-greedy policy exists

Hence a v^* -greedy policy exists

Hence an optimal policy exists

Ex. Show that Bellman's principle of optimality implies existence of an optimal policy

Proof:

Let Bellman's principle of optimality hold

For each $v \in \mathbb{R}^{X}$, a v-greedy policy exists

Hence a v^* -greedy policy exists

Hence an optimal policy exists

We know that T is a contraction mapping on \mathbb{R}^{X}

Hence T is globally stable on \mathbb{R}^{X} with unique fixed point $\bar{v} \in \mathbb{R}^{\mathsf{X}}$

To prove (1) of the above theorem, we have to show $\bar{v}=v^*$

Ex. Show $\bar{v} \leqslant v^*$

$\underline{\mathsf{Proof}} \text{: Let } \sigma \in \Sigma \text{ be } \bar{v}\text{-}\mathsf{greedy}$

Then
$$T_{\sigma}\bar{v}=T\bar{v}=\bar{v}$$

Hence \bar{v} is a fixed point of T_{σ}

But the only fixed point of T_{σ} in \mathbb{R}^{X} is v_{σ}

We have shown existence of a σ such that $\bar{v}=v_{\sigma}$

But then $\bar{v} \leqslant v^*$ (why?)

Ex. Show $v^* \leqslant \bar{v}$

Proof: Fix $\sigma \in \Sigma$ and note that

$$T_{\sigma} v \leqslant T v$$
 for all $v \in \mathbb{R}^{\mathsf{X}}$

Moreover, T is order-preserving and globally stable

As we saw in Ch. 2, this implies that the fixed points are likewise ordered

In particular, $v_{\sigma} \leqslant \bar{v}$

Hence \bar{v} is an upper bound of $\{v_{\sigma}\}_{\sigma\in\Sigma}$

Therefore $v^* \leqslant \bar{v}$

Ex. Show
$$v^* \leqslant \bar{v}$$

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Hence \bar{v} is an upper bound of $\{v_{\sigma}\}_{\sigma \in \Sigma}$

Therefore $v^* \leqslant \bar{v}$

We have now shown that $v^* = \bar{v}$

Thus, v^* is a fixed point of T in \mathbb{R}^X

But T is globally stable on \mathbb{R}^{X}

Hence v^* is the only fixed point of T in \mathbb{R}^{X}

In other words,

 $v^*=$ the unique solution to the Bellman equation in \mathbb{R}^{X}

To prove Bellman's principle of optimality, we use $Tv^*=v^*$

We have

$$\sigma$$
 is v^* -greedy \iff $T_{\sigma} v^* = Tv^* \iff$ $T_{\sigma} v^* = v^*$

The last statement is equivalent to $v_{\sigma} = v^*$ (why?)

Hence

$$\sigma$$
 is v^* -greedy $\iff v^* = v_\sigma \iff \sigma$ is optimal

In other words, Bellman's principle of optimality holds

Algorithms

Previously we used value function iteration (VFI) to solve optimal stopping problems

Here we

- 1. present a generalization suitable for arbitrary MDPs
- 2. introduce two other important methods

The two other methods are called

- 1. Howard policy iteration (HPI) and
- 2. Optimistic policy iteration (OPI)

Algorithm 1: VFI for MDPs

```
\begin{split} & \text{input } v_0 \in \mathbb{R}^{\mathsf{X}}, \text{ an initial guess of } v^* \\ & \text{input } \tau, \text{ a tolerance level for error} \\ & \varepsilon \leftarrow \tau + 1 \\ & k \leftarrow 0 \\ & \text{while } \varepsilon > \tau \text{ do} \\ & & | v_{k+1}(x) \leftarrow (Tv_k)(x) \\ & & \text{end} \\ & \varepsilon \leftarrow \|v_k - v_{k+1}\|_{\infty} \\ & & k \leftarrow k + 1 \end{split}
```

end

Compute a v_k -greedy policy σ

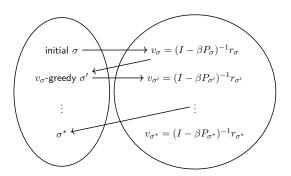
return σ

VFI is

- easy to understand
- easy to implement
- globally convergent
- relatively robust

However, we can often find faster methods with a bit of effort

Howard Policy Iteration



Iterates between computing the value of a given policy and computing the greedy policy associated with that value

Algorithm 2: Howard policy iteration for MDPs

```
\begin{array}{c} v_0 \leftarrow v_\sigma \text{ and } k \leftarrow 0 \\ \textbf{repeat} \\ & \sigma_k \leftarrow \text{a } v_k\text{-greedy policy} \\ & v_{k+1} \leftarrow (I - \beta P_{\sigma_k})^{-1} r_{\sigma_k} \\ & \textbf{if } v_{k+1} = v_k \textbf{ then break} \\ & k \leftarrow k+1 \end{array}
```

return σ_k

input $\sigma \in \Sigma$

Proposition. HPI returns an exact optimal policy in a finite number of steps

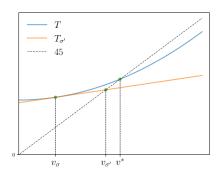
Also, rate of convergence is faster than VFI

In fact HPI is analogous to gradient-based Newton iteration on ${\it T}$

Details are in the text

In general, for a given fixed point problem,

- 1. Newton iteration yields a quadratic rate of convergence
- 2. Successive approximation yields a linear rate of convergence



- σ' is v_{σ} -greedy if $T_{\sigma'}v_{\sigma}=Tv_{\sigma}$
- $v_{\sigma'}$ is the fixed point of $T_{\sigma'}$

Optimistic Policy Iteration

OPI is a "convex combination" of VFI and HPI

Similar to HPI except that

- HPI takes current σ and obtains v_{σ}
- ullet OPI takes current σ and iterates m times with T_{σ}

Recall that, for any $v \in \mathbb{R}^{X}$, we have $T_{\sigma}^{m}v \to v_{\sigma}$ as $m \to \infty$

Hence OPI replaces v_{σ} with an approximation

Algorithm 3: Optimistic policy iteration for MDPs

```
input v_0 \in \mathbb{R}^\mathsf{X}, an initial guess of v^* input \tau, a tolerance level for error input m \in \mathbb{N}, a step size k \leftarrow 0 \varepsilon \leftarrow \tau + 1 while \varepsilon > \tau do  \mid \sigma_k \leftarrow \text{a } v_k\text{-greedy policy}   v_{k+1} \leftarrow T^m_{\sigma_k} v_k   \varepsilon \leftarrow \|v_k - v_{k+1}\|_\infty   k \leftarrow k + 1
```

end

return σ_k

Proposition. For all values of m we have $v_k \to v^*$

Ex. Show that $\mathsf{OPI} = \mathsf{VFI}$ when m = 1

Proof. Suppose we are at step k with $v_k \in \mathbb{R}^{\mathsf{X}}$

Now we take σ_k to be $v_k\text{-greedy}$ and set $v_{k+1}=T_{\sigma_k}v_k$

Since σ_k is v_k -greedy, we also have $T_{\sigma_k}v_k=Tv_k$

Hence v_{k+1} is also the next step for VFI

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Now we take σ_k to be $v_k\text{-greedy}$ and set $v_{k+1}=T_{\sigma_k}v_k$

Since σ_k is v_k -greedy, we also have $T_{\sigma_k}v_k=Tv_k$

Hence v_{k+1} is also the next step for VFI

Intuitively, if $m=\infty$, OPI is identical to HPI

This is because $\lim_{m\to\infty}T^m_{\sigma_k}v_k=v_{\sigma_k}$

Without paralleization, an intermediate value of m is usually better than both

More rules of thumb:

- parallelization tends to favor HPI
- ullet VFI works well when eta is small and optimization is cheap

Application: Optimal Inventories

Previously we analyzed S-s inventory dynamics

our aim was to understand Markov chains

But are such dynamics realistic?

We now investigate whether S-s behavior arises naturally in optimizing model

firm chooses its inventory path to maximize firm value

We assume for now that the firm only sells one product

Given a demand process $(D_t)_{t\geqslant 0}$, inventory $(X_t)_{t\geqslant 0}$ obeys

$$X_{t+1} = f(X_t, A_t, D_{t+1})$$

where

- $f(x, a, d) = (x d) \lor 0 + a$
- ullet A_t is units of stock ordered this period
- ullet The firm can store at most K items at one time

The state space is $X := \{0, \dots, K\}$

We assume $(D_t) \stackrel{\text{IID}}{\sim} \varphi \in \mathfrak{D}(\mathbb{Z}_+)$

Profits are given by

$$\pi_t := X_t \wedge D_{t+1} - cA_t - \kappa \mathbb{1}\{A_t > 0\}$$

- output price $\equiv 1$
- orders in excess of inventory are lost
- c is unit product cost (and unit sales prices = 1)
- κ is a fixed cost of ordering inventory

With $\beta := 1/(1+r)$ and r > 0, the value of the firm is

$$V_0 = \mathbb{E} \sum_{t > 0} \beta^t \pi_t$$

Expected current profit is

$$r(x,a) := \sum_{d \geqslant 0} (x \wedge d)\varphi(d) - ca - \kappa \mathbb{1}\{a > 0\}$$

The set of feasible order quantities at x is

$$\Gamma(x) := \{0, \dots, K - x\}$$

The Bellman equation is

$$v(x) = \max_{a \in \Gamma(x)} \left\{ r(x, a) + \beta \sum_{d \geqslant 0} v(f(x, a, d)) \varphi(d) \right\}$$

The model is an MDP $\mathfrak{M}=(\Gamma,\beta,r,P)$ with

- state space X and action space A := X
- Γ , r and β as above
- stochastic kernel

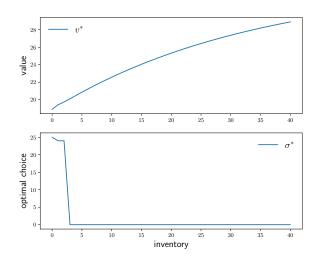
$$P(x, a, x') := \mathbb{P}\{f(x, a, D) = x'\}$$
$$= \sum_{d \ge 0} \mathbb{1}\{f(x, a, d) = x'\}\varphi(d)$$

Since ${\mathfrak M}$ is an MDP, all optimality results on slide 35 apply

using Distributions

```
f(x, a, d) = max(x - d, 0) + a # Inventory update
function create inventory model(; β=0.98, # discount factor
                                   K=40, # maximum inventory
                                   c=0.2, k=2, # cost paramters
                                    p=0.6) # demand parameter
    \phi(d) = (1 - p)^d * p # demand distribution
    x vals = collect(0:K) # set of inventory levels
    return (; \beta, K, c, \kappa, p, \phi, x vals)
end
"The function B(x, a, v) = r(x, a) + \beta \sum x' v(x') P(x, a, x')."
function B(x, a, v, model; d max=100)
    (; \beta, K, c, \kappa, p, \phi, x \text{ vals}) = \text{model}
    revenue = sum(min(x, d) * \phi(d)  for d in 0:d max)
    current profit = revenue - c * a - \kappa * (a > 0)
    next value = sum(v[f(x, a, d) + 1] * \phi(d) for d in 0:d max)
    return current profit + β * next value
end
```

```
"The Bellman operator."
function T(v, model)
     (; \beta, K, c, \kappa, p, \phi, x \text{ vals}) = \text{model}
     new v = similar(v)
     for (x idx, x) in enumerate(x vals)
          \Gamma x = 0: (K - x)
          new_v[x_idx], _ = findmax(B(x, a, v, model) for a in \Gamma x)
     end
     return new v
end
"Get a v-greedy policy."
function get greedy(v, model)
     (; \beta, K, c, \kappa, p, \phi, x vals) = model
    \sigma star = zero(x vals)
     for (x idx, x) in enumerate(x vals)
          \Gamma x = 0: (K - x)
          , a idx = findmax(B(x, a, v, model) for a in \Gammax)
          \sigma \operatorname{star}[x \operatorname{id}x] = \Gamma x[a \operatorname{id}x]
     end
     return σ_star
end
```



Ex. Try to replicate these plots

- Use the code given above (or at least the same parameters)
- Use value function iteration

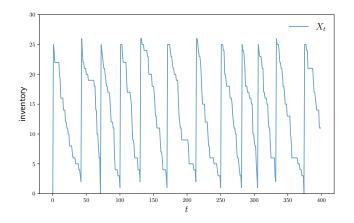


Figure: Optimal inventory dynamics

Optimal Savings with Labor Income

Wealth evolves according to

$$W_{t+1} = R(W_t + Y_t - C_t)$$
 $(t = 0, 1, ...)$

- (W_t) takes values in finite set $W \subset \mathbb{R}_+$
- ullet (Y_t) is Q-Markov chain on finite set ${\sf Y}$
- $C_t, W_t \geqslant 0$

The household maximizes

$$\mathbb{E}\sum_{t\geq 0}\beta^t u(C_t)$$

The model is an MDP with state space $X := W \times Y$

- action = gross savings = next period wealth
- The feasible correspondence is

$$\Gamma(w,y) = \{ s \in \mathsf{W} : s \leqslant R(w+y) \}$$

current reward is

$$r(w, y, s) = u(w + y - s/R)$$
 (utility of consumption)

stochastic kernel is

$$P((w,y), s, (w',y')) = 1\{w' = s\}Q(y,y')$$

Bellman operator:

$$(Tv)(w,y) = \max_{w' \in \Gamma(w,y)} \left\{ u(w+y-w'/R) + \beta \sum_{y' \in \mathsf{Y}} v(w',y')Q(y,y') \right\}$$

The policy operator for given $\sigma \in \Sigma$ is

$$(T_{\sigma} v)(w, y) =$$

$$u(w + y - \sigma(w, y)/R) + \beta \sum_{y' \in Y} v(\sigma(w, y), y')Q(y, y')$$

How to solve $v_{\sigma} = (I - \beta P_{\sigma})^{-1} r_{\sigma}$?

Set

$$P_{\sigma}((w,y),(w',y')) := \mathbb{1}\{\sigma(w,y) = w'\}Q(y,y')$$

and

$$r_{\sigma}(w,y) := u(w+y-\sigma(w,y)/R)$$

Either

- reshape for matrix algebra (compress (w, y) to single index)
- use linear operator routines

How to solve $v_{\sigma} = (I - \beta P_{\sigma})^{-1} r_{\sigma}$?

Set

$$P_{\sigma}((w,y),(w',y')) := \mathbb{1}\{\sigma(w,y) = w'\}Q(y,y')$$

and

$$r_{\sigma}(w,y) := u(w+y-\sigma(w,y)/R)$$

Either

- reshape for matrix algebra (compress (w,y) to single index)
- use linear operator routines

```
using OuantEcon, LinearAlgebra, IterTools
function create savings model(; R=1.01, \beta=0.98, \nu=2.5,
                                                                                                                                                                              w \min_{0.01} w \max_{0.01} w \sup_{0.01} w \sup_{0.01
                                                                                                                                                                              \rho=0.9, \nu=0.1, y size=5)
                     w grid = LinRange(w min, w max, w size)
                      mc = tauchen(y_size, \rho, v)
                      y grid, Q = exp.(mc.state values), mc.p
                      return (; β, R, γ, w grid, y grid, Q)
end
 "B(w, v, w', v) = u(R*w + v - w') + \beta \sum v' v(w', v') O(v, v')."
function B(i, j, k, v, model)
                      (; \beta, R, \gamma, w grid, y grid, Q) = model
                     w, y, w' = w \operatorname{grid}[i], y \operatorname{grid}[j], w \operatorname{grid}[k]
                     u(c) = c^{(1-\gamma)} / (1-\gamma)
                     c = w + v - (w' / R)
                     @views value = c > 0 ? u(c) + \beta * dot(v[k, :], 0[j, :]) : -Inf
                      return value
end
```

```
"The Bellman operator."
function T(v. model)
    w idx, y idx = (eachindex(g) for g in (model.w grid, model.y grid))
    v new = similar(v)
    for (i, j) in product(w idx, y idx)
         v \text{ new}[i, j] = \text{maximum}(B(i, j, k, v, model) \text{ for } k \text{ in } w \text{ idx})
    end
    return v new
end
"The policy operator."
function T \sigma(v, \sigma, model)
    w idx, y idx = (eachindex(g) for g in (model.w grid, model.y grid))
    v new = similar(v)
    for (i, j) in product(w idx, y idx)
         v \text{ new[i, i]} = B(i, i, \sigma[i, i], v. \text{ model})
    end
    return v new
end
```

```
include("finite_opt_saving_0.jl")
"Compute a v-greedy policy."
function get_greedy(v, model)
    w_idx, y_idx = (eachindex(g) for g in (model.w_grid, model.y_grid))
    σ = Matrix{Int32}(undef, length(w_idx), length(y_idx))
    for (i, j) in product(w_idx, y_idx)
        _, σ[i, j] = findmax(B(i, j, k, v, model) for k in w_idx)
    end
    return σ
end
```

"Get the value v σ of policy σ ." function get value(g. model) # Unpack and set up (; β , R, γ , w grid, γ grid, Q) = model w idx, y idx = (eachindex(q) for q in (w grid, y grid)) wn, yn = length(w idx), length(y idx)n = wn * yn $u(c) = c^{(1-\gamma)} / (1-\gamma)$ # Build P σ and r σ as multi-index arrays $P \sigma = zeros(wn, yn, wn, yn)$ $r \sigma = zeros(wn, yn)$ for (i, i) in product(w idx, v idx) w, y, w' = w grid[i], y grid[j], $w grid[\sigma[i, j]]$ $r \sigma[i, j] = u(w + y - w'/R)$ for j' in y idx $P \sigma[i, j, \sigma[i, j], j'] = Q[j, j']$ end end # Reshape for matrix algebra $P \sigma = reshape(P \sigma, n, n)$ $r \sigma = reshape(r \sigma, n)$ # Apply matrix operations --- solve for the value of σ $v \sigma = (I - \beta * P \sigma) \setminus r \sigma$ # Return as multi-index array return reshape(v σ, wn, yn) end

```
"Value function iteration routine."
function value iteration(model, tol=1e-5)
    vz = zeros(length(model.w grid), length(model.v grid))
    v star = successive approx(v -> T(v, model), vz, tolerance=tol)
    return get greedy(v star, model)
end
"Howard policy iteration routine."
function policy iteration(model)
    wn, yn = length(model.w grid), length(model.y grid)
    \sigma = ones(Int32, wn, vn)
    i, error = 0, 1.0
    while error > 0
        v \sigma = get value(\sigma, model)
        \sigma new = get greedy(v \sigma, model)
        error = maximum(abs.(\sigma new - \sigma))
         \sigma = \sigma \text{ new}
        i = i + 1
         println("Concluded loop $i with error $error.")
    end
    return o
end
```

```
"Optimistic policy iteration routine."
function optimistic_policy_iteration(model; tolerance=le-5, m=100)
   v = zeros(length(model.w_grid), length(model.y_grid))
   error = tolerance + 1
   while error > tolerance
        last_v = v
        σ = get_greedy(v, model)
        for i in 1:m
            v = T_σ(v, σ, model)
        end
        error = maximum(abs.(v - last_v))
end
   return get_greedy(v, model)
end
```

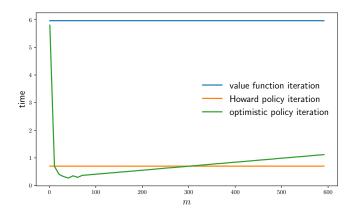


Figure: Timings for alternative algorithms

Investment with Adjustment Costs

A monopolist faces an inverse demand function of the form

$$P_t = a_0 - a_1 Y_t + Z_t$$

- a_0, a_1 are positive parameters
- Y_t is output
- P_t is price and
- the demand shock Z_t follows

$$Z_{t+1} = \rho Z_t + \sigma \eta_{t+1}, \qquad \{\eta_t\} \stackrel{\text{IID}}{\sim} N(0,1)$$

Current profits are given by

$$\pi_t := P_t Y_t - c Y_t - \gamma (Y_{t+1} - Y_t)^2$$

- $\gamma (Y_{t+1} Y_t)^2$ represents adjustment costs
- ullet rapid changes to capacity are expensive when γ large

Objective: maximize value

$$\mathbb{E}\sum_{t=0}^{\infty}\beta^t\pi_t$$

• $\beta = 1/(1+r)$, where r > 0 is a fixed interest rate

Building intuition: If $\gamma=0$ then no intertemporal trade-off

Therefore maximize current profit each period:

$$\max_{Y_t} \{ P_t Y_t - c Y_t \} = \max_{Y_t} \{ (a_0 - a_1 Y_t + Z_t) Y_t - c Y_t \}$$

Ex. Show that the maximizer is

$$\bar{Y}_t := \frac{a_0 - c + Z_t}{2a_1}$$

On the other hand, if γ is very large then $(Y_t)_{t\geqslant 0}$ should be almost constant

Thus, we expect the following:

- $\gamma \approx 0 \implies Y_t$ tracks \bar{Y}_t closely
- ullet γ large $\implies Y_t$ will be smoother than $ar{Y}_t$

In short,

more adjustment costs \implies smoother time path for $(Y_t)_{t\geqslant 0}$

Implementation as an MDP

Let $Y \subset \mathbb{R}_+$ be a grid containing output values

We discretize (Z_t) using Tauchen's method:

 (Z_t) is Q-Markov on finite set $\mathsf{Z} \subset \mathbb{R}$

The state space $X := Y \times Z$

The action space is Y

The feasible correspondence is $\Gamma(x) = Y$ for all x

choice of output is not restricted by the state

The set Σ is all $\sigma \colon \mathsf{Y} \times \mathsf{Z} \to \mathsf{Y}$

The current reward function is current profits:

$$r(y, z, \hat{y}) = (a_0 - a_1 y + z - c)y - \gamma(\hat{y} - y)^2$$

• $\hat{y} = action = choice of next-period output$

The stochastic kernel is

$$P((y,z), \hat{y}, (y',z')) = \mathbb{1}\{y' = \hat{y}\}Q(z,z')$$

Now the problem defines an MDP

all of the optimality theory for MDPs applies

Optimal Investment

The Bellman and policy operators for this problem are

$$(Tv)(y,z) = \max_{y' \in \mathbb{R}} \left\{ r(y,z,y') + \beta \sum_{z' \in \mathsf{Z}} v(y',z') Q(z,z') \right\}$$
$$(T_{\sigma} v)(y,z) = r(y,z,\sigma(y,z)) + \beta \sum_{z' \in \mathsf{Z}} v(\sigma(y,z),z') Q(z,z')$$

On \mathbb{R}^X , both are

- order-preserving
- ullet contractions of modulus eta

A v-greedy policy is a $\sigma \in \Sigma$ that obeys

$$\sigma(y,z) = \operatorname*{argmax}_{y' \in \mathsf{Y}} \left\{ r(y,z,y') + \beta \sum_{z' \in \mathsf{Z}} v(y',z') Q(z,z') \right\}$$

By our results for MDPs

- ullet v^* -greedy policies = optimal policies
- OPI / HPI / VFI converge

Implications for output can be studied by

- 1. generating a $Q ext{-Markov}$ chain $(Z_t)_{t=1}^T$
- 2. simulating optimal output via $Y_{t+1} = \sigma^*(Y_t, Z_t)$

```
using QuantEcon, LinearAlgebra, IterTools
include("s approx.jl")
function create investment model(;
        r=0.04.
                                               # Interest rate
        a 0=10.0, a 1=1.0,
                                               # Demand parameters
        y=25.0, c=1.0,
                                               # Adjustment and unit cost
        y_min=0.0, y_max=20.0, y_size=100,  # Grid for output
        \rho = 0.9, \nu = 1.0,
                                               # AR(1) parameters
        z size=25)
                                               # Grid size for shock
    \beta = 1/(1+r)
    y grid = LinRange(y min, y max, y size)
    mc = tauchen(z_size, \rho, v)
    z grid, Q = mc.state values, mc.p
    return (; β, a 0, a 1, γ, c, y grid, z grid, Q)
end
```

```
0.00
```

The aggregator B is given by

$$B(y, z, y') = r(y, z, y') + \beta \Sigma_z' v(y', z') Q(z, z')."$$

where

$$r(y, z, y') := (a_0 - a_1 * y + z - c) y - \gamma * (y' - y)^2$$

0.0.0

```
function B(i, j, k, v, model)  (; \ \beta, \ a\_0, \ a\_1, \ \gamma, \ c, \ y\_grid, \ z\_grid, \ Q) \ = \ model \\ y, \ z, \ y' \ = \ y\_grid[i], \ z\_grid[j], \ y\_grid[k] \\ r \ = \ (a\_0 \ - \ a\_1 \ * \ y + \ z \ - \ c) \ * \ y \ - \ \gamma \ * \ (y' \ - \ y)^2 \\ return \ @views \ r \ + \ \beta \ * \ dot(v[k, :], \ Q[j, :]) \\ end
```