## Dynamic Programming

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# Order Isomorphism

## Order Isomorphism

#### Definition 1

A surjective map F from poset  $(V, \lesssim)$  to poset  $(\hat{V}, \leq)$  is called an

- Order isomorphism if  $v \lesssim w \iff Fv \leq Fw$
- Order anti-isomorphism if  $v \lesssim w \iff Fw \leq Fv$

Comment: F under this definition is bijective.

### Exercise 3.1.1.

Given  $h \in \mathbb{R}^X$ , let  $Fh = \exp(\theta h)$ . Show that F is an order isomorphism from  $\mathbb{R}^X$  to  $(0,\infty)^X$  whenever  $\theta > 0$ .

### Proof.

Fix  $\theta > 0$ . We know that  $\exp(\theta x)$  is a bijective function. Let  $h_1, h_2 \in \mathbb{R}^X$  such that  $h_1 < h_2$ . This implies

$$\theta h_1 \le \theta h_2$$

As  $\exp(\cdot)$  is order preserving, hence, we have

$$Fh_1 = \exp(\theta h_1) \le \exp(\theta h_2) = Fh_2$$

Let  $k_1, k_2 \in (0, \infty)^X$ ,  $k_1 < k_2$  by surjectivity, we have

$$k_1 = F(q_1) = \exp(\theta q_1), k_2 = F(q_2) = \exp(\theta q_2), q_1, q_2 \in \mathbb{R}^X$$

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### Exercise 3.1.1. Continue

### Proof.

We have

$$q_1 = \frac{\ln k_1}{\theta} \le \frac{\ln k_2}{\theta} = q_2$$

as ln is order preserving. Therefore, by definition, F is an order isomorphism from  $\mathbb{R}^X$  to  $(0,\infty)^X$ 

Exercise 3.1.2.

Let  $V = M^X$  and  $\hat{V} = \hat{M}^X$ ,  $M, \hat{M} \subset \mathbb{R}$ . Let  $\varphi$  be a map from M onto  $\hat{M}$  and let  $Fv = \varphi \circ v$ . Prove if  $\varphi$  is an order isomorphism from M to  $\hat{M}$ , then F is an order isomorphism from V to  $\hat{V}$ .

### Proof.

 $\varphi$  is order isomorphism then  $\varphi$  is bijective, order preserving with order preserving inverse. Hence apply this dim X times, we get F is bijective, order preserving with order preserving inverse.

### Exercise 3.1.3

Let  $V, \hat{V}$  be posets. Show that every order isomorphism F is a bijection. Show that every order anti-isomorphism is also a bijection.

### Proof.

Let  $v_1, v_2 \in \hat{V}$  such that  $v_1 = v_2$ . By surjectivity, we have

$$v_1 = F(w_1), \quad v_2 = F(w_2), \quad w_1, w_2 \in V$$

Hence, we have

$$F(w_1) \le F(w_2) \implies w_1 \lesssim w_2$$

$$F(w_2) \le F(w_1) \implies w_2 \lesssim w_1$$

Hence,  $w_1 = w_2$ . This proves that F is injective.

### Exercise 3.1.4.

Let F be a bijection from  $(V, \preceq)$  to  $(\hat{V}, \leq)$ . Show that

- $\bullet$  F is an order isomorphism if and only if F and  $F^{-1}$  are order preserving

### Proof.

Skip

### Lemma 3.1.1.

Let F be an order isomorphism from  $(V, \preceq)$  to  $(\hat{V}, \leq)$ . If the supremum of  $\{v_{\alpha}\}_{{\alpha} \in \Lambda} \subset V$  exists in V, then

$$\bigvee_{\alpha} Fv_{\alpha}$$
 exists in  $\hat{V}$  and  $\bigvee_{\alpha} Fv_{\alpha} = F\bigvee_{\alpha} v_{\alpha}$ 

### Proof.

Let  $v := \bigvee_{\alpha} v_{\alpha} \in V$ . Let  $\hat{w}$  be any upper bound of  $\{Fv_{\alpha}\}$ , i.e.,  $Fv_{\alpha} \leq \hat{w}$  for all  $\alpha \in \Lambda$ . By surjectivity, we let  $\hat{w} = F(w)$ , and by order isomorphism, we have

$$v_{\alpha} \preceq w$$
 for all  $\alpha \in \Lambda$ 

Hence, w is an upper bound of  $\{v_{\alpha}\}$ , this implies  $v \lesssim w$ . Hence,

$$F(v) \le F(w) = \hat{w}$$

This implies  $F(v) = F \bigvee_{\alpha} v_{\alpha}$  is the least upper bound of  $\{Fv_{\alpha}\}$ .

### Exercise 3.1.6

Let  $V, \hat{V}$  be posets and let  $(v_n)$  be a sequence in V. And let F be a map from V to  $\hat{V}$ . Prove the following

- If F is an order isomorphism, then  $v_n \uparrow v$  if and only if  $Fv_n \uparrow Fv$  in V
- $\bullet$  If F is an order anti-isomorphism, then  $v_n \uparrow v$  if and only if  $Fv_n \downarrow Fv$  in  $\hat{V}$ .

### Proof.

$$v_n \uparrow v \implies v_1 \le v_2 \le \dots \le v \text{ and } v = \bigvee_n v_n.$$

Hence, by order isomorphism, we have

$$Fv_1 \le Fv_2 \le \cdots \le Fv$$

and  $Fv = \bigvee_n Fv_n$  Moreover, F is order isomorphism implies  $F^{-1}$  is order isomorphism. Hence, the other direction follows.

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### Exercise 3.1.7

#### Prove

- If  $V, \hat{V}$  are order isomorphic, then V is totally ordered if and only if  $\hat{V}$  is totally ordered
- F is an order anti-isomorphism from V to  $\hat{V}$  if and only if F is an order isomorphism from V to its dual  $\hat{V}^{\partial}$

## Proof.

Skip.

## Conjugate Dynamics

- We start with the definition of conjugacy between dynamical systems  $((V, T_{\sigma}))$  is a dynamical system with state space V and evolution  $T_{\sigma}$ ).
- ullet Then, we go to the most basic structure, V as a poset, and upgrade conjugacy to order conjugacy.
- This prepares for the later upgrade from dynamical system to ADP

### **Definitions**

### Definition 2

We call a **discrete time dynamical system** is a pair (V, S), where V is any set, and S is a self-map on V.

### Definition 3

Two dynamical systems (V,S) and  $(\hat{V},\hat{S})$  are said to be **conjugate under** F if

F is a bijection from V into  $\hat{V}$  and  $F \circ S = \hat{S} \circ F$  on V

or we can write it as

$$S = F^{-1} \circ \hat{S} \circ F$$

## Proposition 3.1.2.

If (V, S) and  $(\hat{V}, \hat{S})$  are conjugate, then

- ② v is a fixed point of S if and only if Fv is a fixed point of  $\hat{S}$
- $\bullet$   $\hat{v}$  is fixed point of  $\hat{S}$  if and only if  $F^{-1}\hat{v}$  is a fixed point of S
- v is the unique fixed point of S in V if and only if Fv is the unique fixed point of  $\hat{S}$  in  $\hat{V}$ .

#### Proof.

Let v be the unique fixed point of S, i.e., Sv = v. Hence,

$$F(v) = F(Sv) = \hat{S}(Fv)$$

Hence, F(v) is a fixed point of  $\hat{S}$ . Let  $\hat{w} = F(w)$  be a fixed point  $\hat{S}$ , then by part 2, w is the fixed point of S. Hence, w = v, and this implies F(w) = F(v).

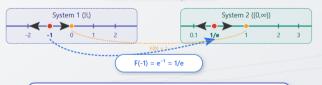
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#### **Conjugate Dynamical Systems**

Example 3.1.2



#### **Phase Lines and Dynamics**



#### Conjugacy Verification: $F \circ S = \hat{S} \circ F$

$$F(S(x)) = e^{(2x + 1)} = e^{1} \cdot e^{(2x)} = e \cdot (e^{x})^{2} = \hat{S}(F(x))$$

## Order conjugacy

### Definition 4

Consider two dynamical systems (V, S) and  $(\hat{V}, \hat{S})$ , where  $V, \hat{V}$  are posets. We call these systems **order conjugate under** F if they are conjugate under F, and, F is an order isomorphism.

### Exercise 3.1.9.

Prove that order conjugacy is an equivalence relation on the set of dynamical systems over partially ordered set.

### Proof.

We denote  $(V, S) \sim (\hat{V}, \hat{S})$  is they are order conjugate. We need to show this relation is reflexive, symmetric and transitive.

• (Reflexivity) Let F = Id which is a bijection. We have

$$F \circ S = S = S \circ F$$

Moreover, we have  $F = F^{-1}$  is order preserving. Hence,

$$(V,S) \sim (V,S)$$



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### Exercise 3.1.9 Continue

### Proof.

- (Symmetry) Let  $(V, S) \sim (\hat{V}, \hat{S})$  under F.
  - F is bijection implies  $F^{-1}$  is bijection
  - $F \circ S = \hat{S} \circ F \implies F^{-1} \circ \hat{S} = S \circ F^{-1}$

Hence, (V, S) and  $(\hat{V}, \hat{S})$  are conjugate under  $F^{-1}$ .

• F is order preserving with order preserving inverse implies  $F^{-1}$  is order preserving with order preserving inverse

Hence,  $F^{-1}$  is order isomorphism. Hence  $(\hat{V}, \hat{S}) \sim (V, S)$ 



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### Exercise 3.1.9 Continue

### Proof.

- (Transitive) Let  $(V_1, S_1) \sim (V_2, S_2)$  under F and  $(V_2, S_2) \sim (V_3, S_3)$  under G.
  - F, G are bijective implies  $H := (G \circ F)$  is bijective
  - $\bullet$   $F \circ S_1 = S_2 \circ F, G \circ S_2 = S_3 \circ G \implies (G \circ F) \circ S_1 = G \circ S_2 \circ F = S_3 \circ (G \circ F)$

Hence,  $(V_1, S_1)$  and  $(V_3, S_3)$  are conjugate under H.

- F, G are order preserving with order perserving inverses
- $G \circ F$  are order preserving with order preserving inverses

Hence,  $(V_1, S_1) \sim (V_3, S_3)$  under  $(G \circ F)$ .



### Lemma 3.1.3.

If (V, S) and  $(\hat{V}, \hat{S})$  are order conjugate under F, then S is order stable on V if and only if  $\hat{S}$  is order stable on  $\hat{V}$ .

#### Proof.

- $(\Longrightarrow)$  Suppose S is order stable on V. This implies
- (S1) S has a unique fixed point  $v^* \in V$
- (S2)  $v \in V, v \preceq v^* \implies v \preceq Sv \text{ and } v \in V, v^* \preceq v \implies Sv \preceq v$
- (S1) implies  $\hat{S}$  has a unique fixed point  $\hat{v}^* := F(v) \in \hat{V}$  by Proposition 3.1.2. Moreover we have
  - For  $\hat{v} := F(v) \in \hat{V}, \hat{v} \le \hat{v}^* \underset{o.i.}{\Longrightarrow} v \lesssim v^* \underset{S2}{\Longrightarrow} v \lesssim Sv \underset{o.i.}{\Longrightarrow} Fv \le \underbrace{FSv = \hat{S}\hat{v}}_{conjugate}$

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## Isomorphic ADP

In this section, we will see how one ADP can be transformed into another ADP (or ADPs are equivalent up to a transformation).

Such transformation will result

- Simpler form of Bellman equation
- Tailored to solve for some problems (Exponential BE)
- etc.

### Definitions

#### Definition 5

Let  $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  be ADPs with policy sets  $\mathbb{T} := \{T_{\sigma} : \sigma \in \Sigma\}$  and  $\hat{\mathbb{T}} := \{\hat{T}_{\sigma} : \sigma \in \Sigma\}$ . We call these ADPs **isomorphic** under F if

- F is an order isomorphism from V to  $\hat{V}$
- 2 these two ADPs have the same policy set  $\Sigma$
- **3**  $(V, T_{\sigma})$  and  $(V, \hat{T}_{\sigma})$  are order conjugate under F for all  $\sigma \in \Sigma$ .

## Example 3.1.3. Fei et al. (2021) Exponential Bellman Equation

Exponential risk-sensitive Q-factor Bellman equation (ADP:  $((0, \infty)^G, \mathbb{M})$ )

$$M_{\sigma}h = \exp(\theta r + \beta \ln P_{\sigma}h), \quad P_{\sigma}h(x, a) := \sum_{x'} h(x', \sigma(x'))P(x, a, x')$$

Risk-sensitive Q-factor policy operator (ADP:  $(\mathbb{R}^G, \mathbb{T})$ )

$$T_{\sigma}f = r + \frac{\beta}{\theta} \ln \left[ P_{\sigma} \exp(\theta f) \right], \quad P_{\sigma} \exp(\theta f)(x, a) := \sum_{x'} \exp(\theta f(x', \sigma(x')) P(x, a, x'))$$

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## Example 3.1.3. Continue

Let  $\theta > 0$ , and

$$Fh = \exp(\theta h)$$

is an order isomorphism from  $\mathbb{R}^G$  to  $(0, \infty)^G$ . For conjugacy, we have

$$(F \circ T_{\sigma})(h) = \exp\left(\theta(r + \frac{\beta}{\theta}) \ln\left[P_{\sigma} \exp(\theta h)\right]\right)$$
$$= \exp(\theta r + \beta \ln P_{\sigma}(Fh))$$
$$= (M_{\sigma} \circ F)(h)$$

Hence,  $((0,\infty)^G, \mathbb{M})$  and  $(\mathbb{R}^G, \mathbb{T})$  are isomorphic

## Example 3.1.4. RDP

Let  $(\Gamma, V, B)$  and  $(\Gamma, \hat{V}, \hat{B})$  be two RDPs with identical state sapce X, action space A, and feasible correspondence  $\Gamma$ . Let  $V = M^X, \hat{V} = \hat{M}^X$ , where  $M, \hat{M} \subset \mathbb{R}$ . If there exists an order isomorphism  $\varphi$  from M to  $\hat{M}$  such that

$$B(x, a, v) = \varphi^{-1}[\hat{B}(x, a, \varphi \circ v)]$$
 for all  $v \in V$  and  $(x, a) \in G$ 

then  $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  are isomorphic. From exercise 3.1.2, F is an order isomorphism from V to  $\hat{V}$ , and

$$T_{\sigma} = F^{-1} \circ \hat{T}_{\sigma} \circ F$$

### Lemma 3.1.4.

#### Lemma 6

Isomorphism between ADPs is an equivalence relation on the set of ADPs.

### Proof.

Let  $\mathbb{A}$  be the set of ADPs. We denote  $(V_1, \mathbb{T}_1) \sim (V_2, \mathbb{T}_2)$  if there are isomorphic. We need to prove that  $\sim$  is reflexive, symmetric and transitive.

- (Reflexivity) Let  $(V, \mathbb{T}) \in \mathbb{A}$ , as the ADP has the same policy set as itself and by Exercise 3.1.9, we get reflexivity.
- (Symmetry) Let  $(V_1, \mathbb{T}_1) \sim (V_2, \mathbb{T}_2)$ , then they have the same policy set. We use Exercise 3.1.9 get symmetry
- (Transitivity) Let  $(V_1, \mathbb{T}_1) \sim (V_2, \mathbb{T}_2)$  and  $(V_2, \mathbb{T}_2) \sim (V_3, \mathbb{T}_3)$ , hence these three ADPs have the same policy set. We use Exercise 3.1.9. get transitivity.

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Isomorphisms and Optimality

### Notations

#### We take

- $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  be two ADPs
- $\mathbb{T} := \{ T_{\sigma} : \sigma \in \Sigma \}; \ \hat{\mathbb{T}} := \{ \hat{T}_{\sigma} : \sigma \in \Sigma \}$
- $v_{\sigma}(resp.\hat{v}_{\sigma})$  be the unique fixed point of  $T_{\sigma}(resp.\hat{T}_{\sigma})$
- $T(resp.\hat{T})$  be the Bellman operator of  $(V, \mathbb{T})$  (resp.  $(\hat{V}, \hat{\mathbb{T}})$ )
- $v^*$  (resp.  $\hat{v}^*$ ) be the value function of  $(V, \mathbb{T})$  (resp.  $(\hat{V}, \hat{\mathbb{T}})$ )

Theorem 3.1.5.

### Theorem 7

If  $(V, \mathbb{T})$  and  $(\hat{V}, \hat{\mathbb{T}})$  are isomorphic under F, then

- $\bullet \ \sigma \ is \ v\text{-}greedy \ for \ (V,\mathbb{T}) \iff \sigma \ is \ Fv\text{-}greedy \ for \ (\hat{V},\hat{\mathbb{T}})$
- $\bullet$   $\sigma$  is optimal for  $(V, \mathbb{T}) \iff \sigma$  is optimal for  $(\hat{V}, \hat{\mathbb{T}})$
- 3 Regularity, well-posedness, and order stability is preserved under isomorphism.