

软件分析

南京大学

计算机科学与技术系

程序设计语言与
静态分析研究组

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Static Program Analysis

Data Flow Analysis — Applications

Nanjing University

Yue Li

2022

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2. Preliminaries of Data Flow Analysis
3. Reaching Definitions Analysis
4. Live Variables Analysis
5. Available Expressions Analysis

Data Flow Analysis

Data Flow Analysis

How Data Flows on CFG?

Data Flow Analysis

How Data Flows on CFG?

How Data
Flows through the

CFG (a program)?

Data Flow Analysis

How Data Flows on CFG?

How application-specific Data
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Data Flow Analysis

How Data Flows on CFG?

How application-specific Data
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Nodes (BBs/statements) and
Edges (control flows) of
CFG (a program)?

Recall 1st
lecture

Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction

Over-approximation → Flows through the
Nodes (BBs/statements) and
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Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction



Over-approximation → Flows through the

Nodes (BBs/statements) and
Edges (control flows) of
CFG (a program)?

for most static analyses
(may analysis)

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lecture

Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction



Over-approximation → Flows through the

Nodes (BBs/statements) and
Edges (control flows) of
CFG (a program)?

may analysis:

outputs information that may be true (over-approximation)

must analysis:

outputs information that must be true (under-approximation)

Over- and under-approximations are both for safety of analysis

Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction

Safe-approximation → Flows through the
may analysis: safe=over Nodes (BBs/statements) and
must analysis: safe=under Edges (control flows) of
 CFG (a program)?

Data Flow Analysis

How Data Flows on CFG?

How application-specific Data ← Abstraction

Safe-approximation → Flows through the

Transfer function → Nodes (BBs/statements) and

Control-flow handling → Edges (control flows) of
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How Data Flows on CFG?

+,-,0, \perp , \top

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$+ \ op \ - = - ; + \ op \ + = +$

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union the signs at merges

different data-flow analysis applications have
different data abstraction and
different flow safe-approximation strategies, i.e.,
different transfer functions and control-flow handlings

Recall 1st
lecture

Data Flow Analysis

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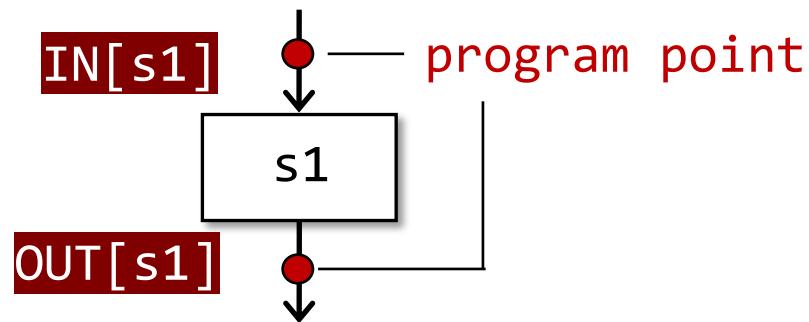
union the signs at merges

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different transfer functions and control-flow handlings

Preliminaries of Data Flow Analysis

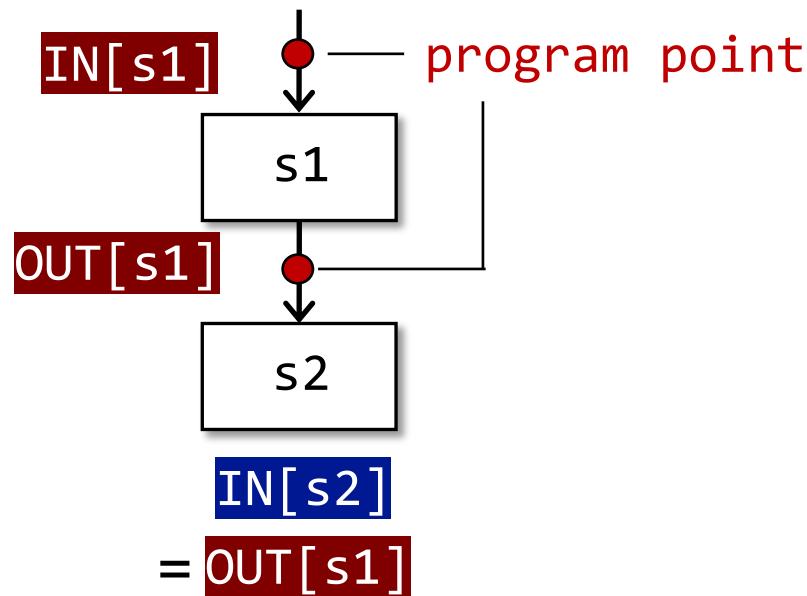
Input and Output States

- Each execution of an IR statement transforms an **input state** to a new **output state**
- The input (output) state is associated with the **program point** before (after) the statement



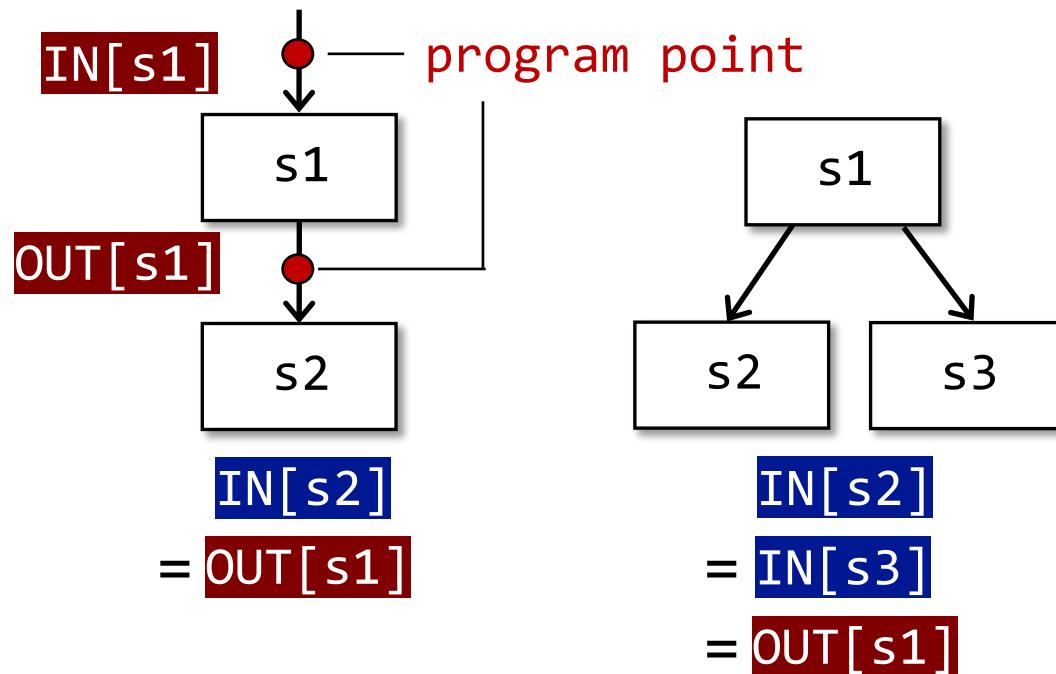
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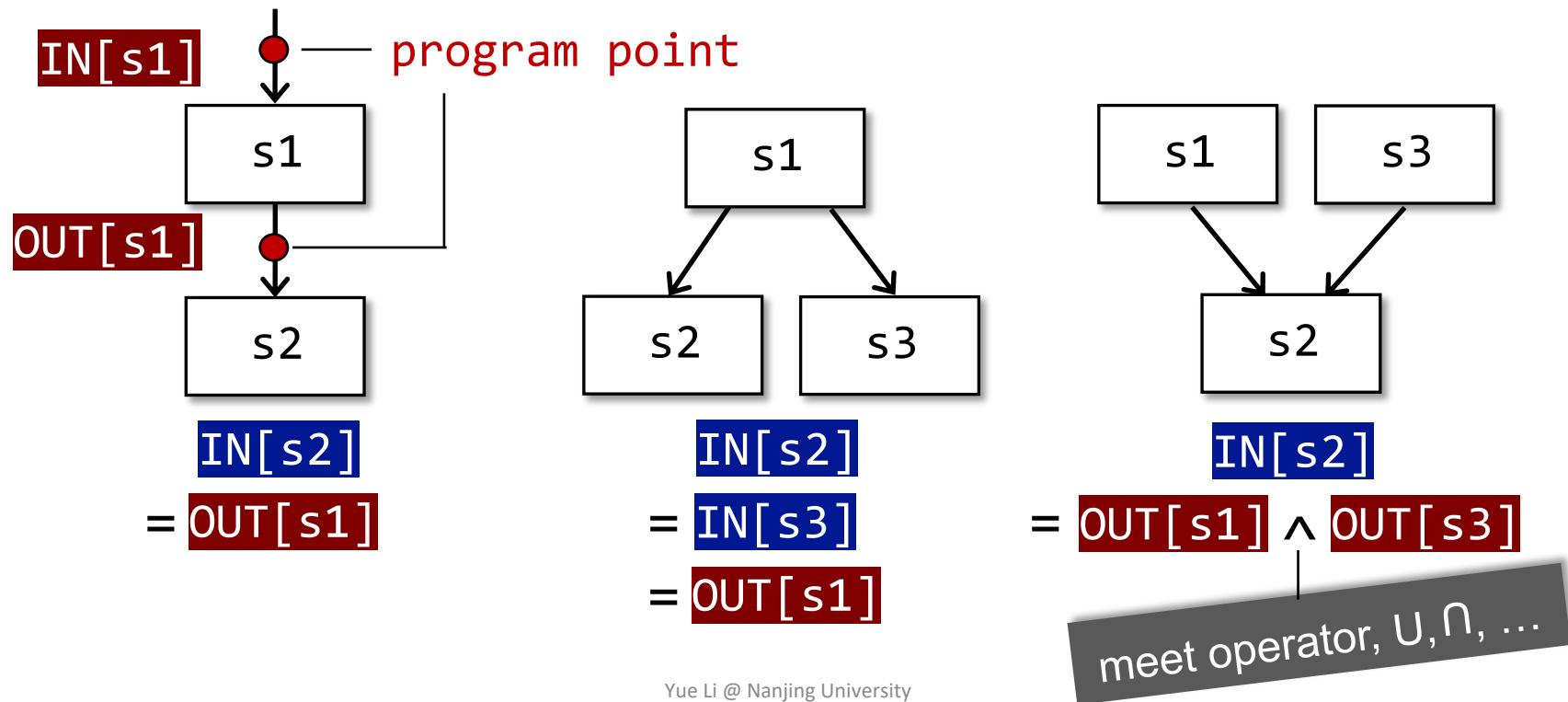
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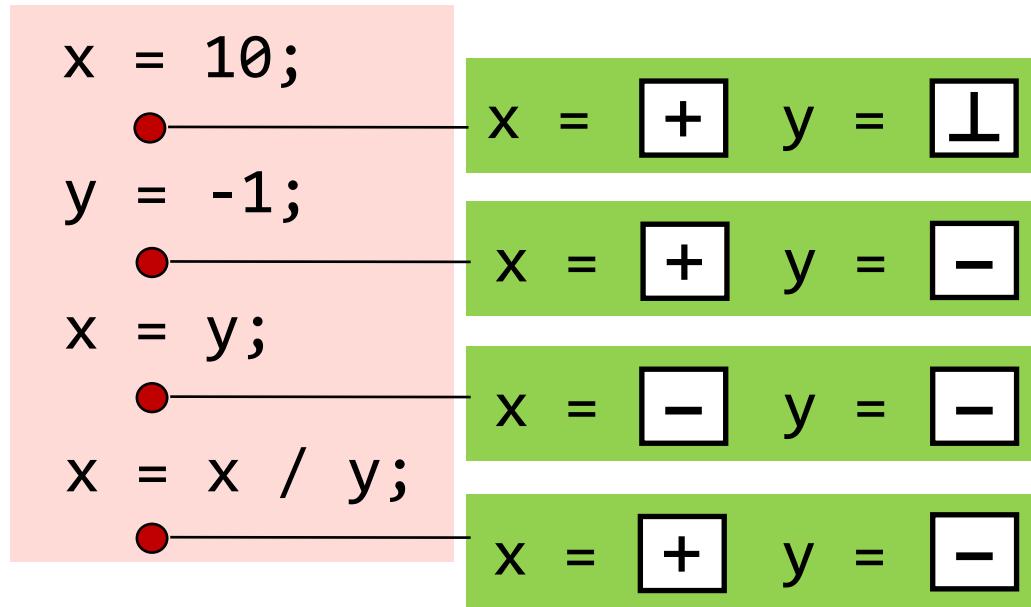
In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.

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```
x = 10;  
●  
y = -1;  
●  
x = y;  
●  
x = x / 6;  
●
```

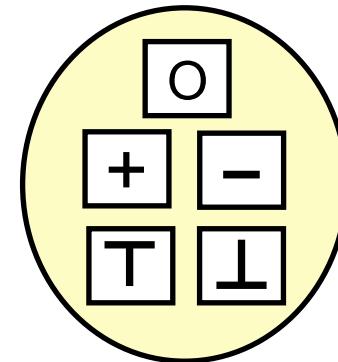
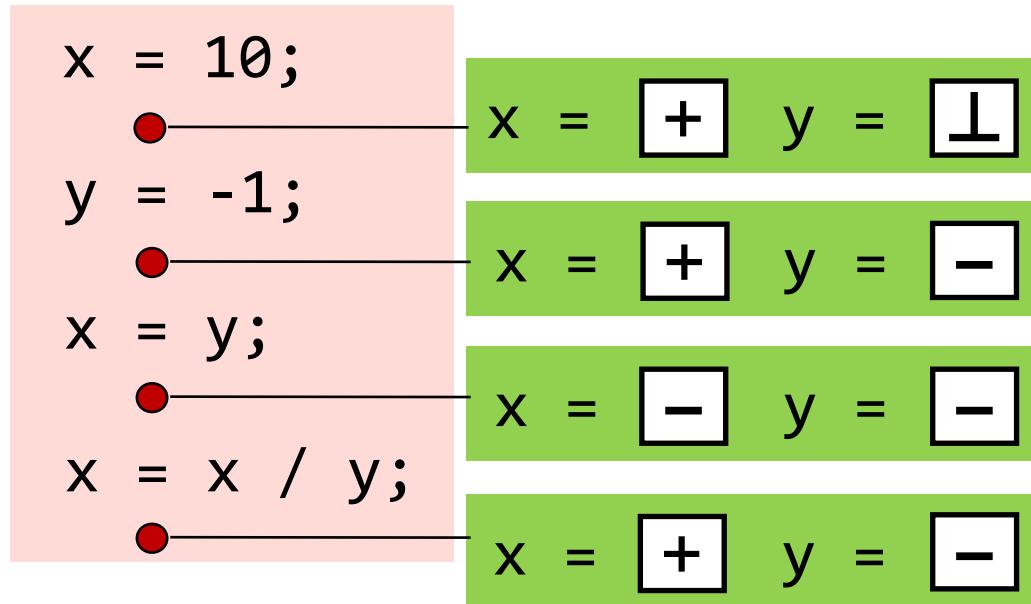
Recall 1st
lecture

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Recall 1st lecture

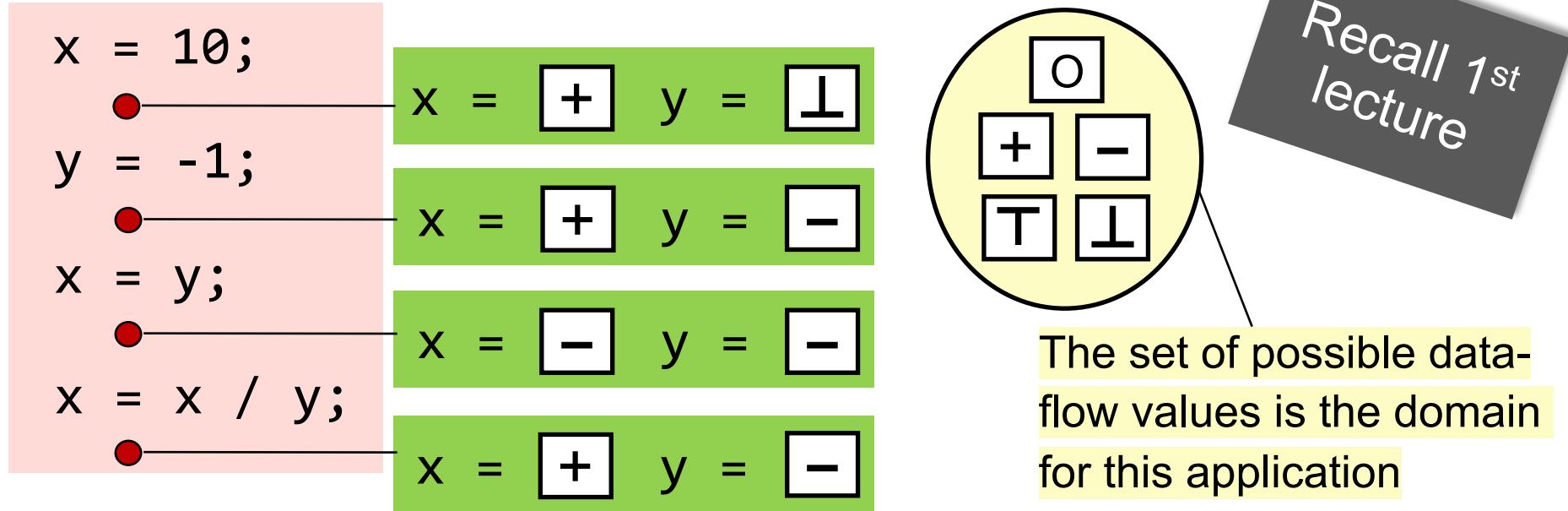
In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.



Recall 1st lecture

The set of possible data-flow values is the domain for this application

In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.



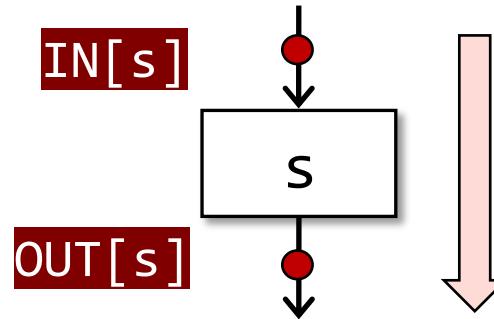
Data-flow analysis is to **find a solution** to a set of *safe-approximation-directed constraints* on the $\text{IN}[s]$'s and $\text{OUT}[s]$'s, for **all statements** s .

- *constraints* based on semantics of statements (*transfer functions*)
- *constraints* based on the *flows of control*

Notations for Transfer Function's Constraints

- Forward Analysis

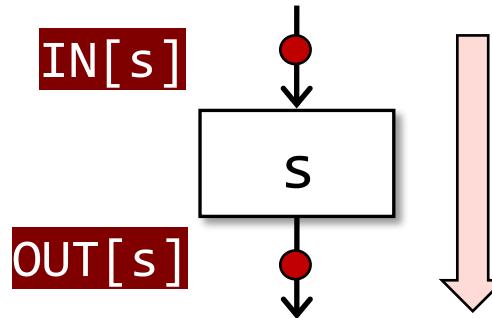
$$\text{OUT}[s] = f_s(\text{IN}[s])$$



Notations for Transfer Function's Constraints

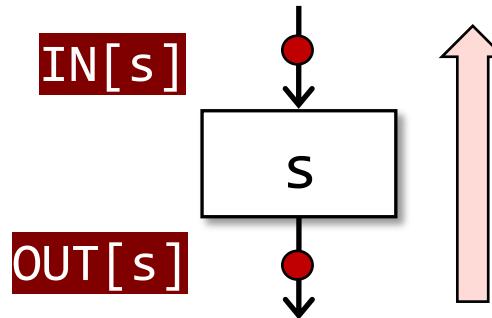
- Forward Analysis

$$\text{OUT}[s] = f_s(\text{IN}[s])$$



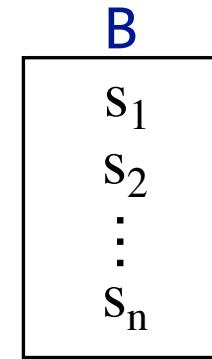
- Backward Analysis

$$\text{IN}[s] = f_s(\text{OUT}[s])$$



Notations for Control Flow's Constraints

- Control flow within a BB
- Control flow among BBs

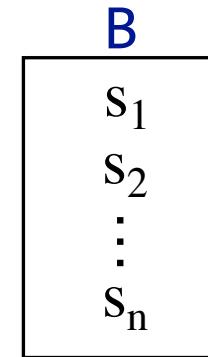


Notations for Control Flow's Constraints

- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs



Notations for Control Flow's Constraints

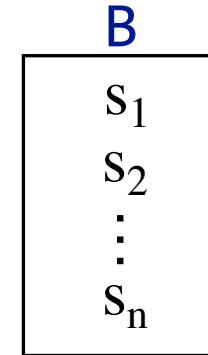
- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs

$$\text{IN}[B] = \text{IN}[s_1]$$

$$\text{OUT}[B] = \text{OUT}[s_n]$$



Notations for Control Flow's Constraints

- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs

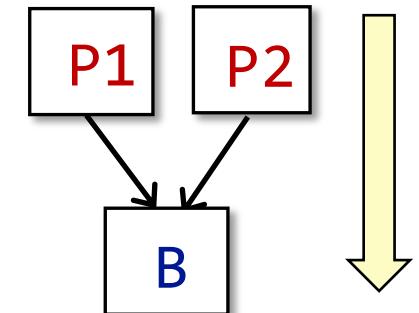
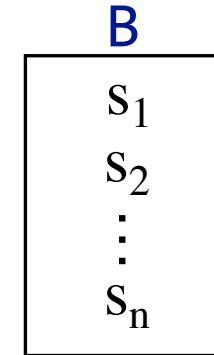
$$\text{IN}[B] = \text{IN}[s_1]$$

$$\text{OUT}[B] = \text{OUT}[s_n]$$

$$\text{OUT}[B] = f_B(\text{IN}[B]), \quad f_B = f_{s_n} \circ \dots \circ f_{s_2} \circ f_{s_1}$$

$$\text{IN}[B] = \bigwedge P \text{ a predecessor of } B \text{ OUT}[P]$$

The meet operator \bigwedge is used to summarize the contributions from different paths at the confluence of those paths



Notations for Control Flow's Constraints

- Control flow within a BB

$$\text{IN}[s_{i+1}] = \text{OUT}[s_i], \text{ for all } i = 1, 2, \dots, n-1$$

- Control flow among BBs

$$\text{IN}[B] = \text{IN}[s_1]$$

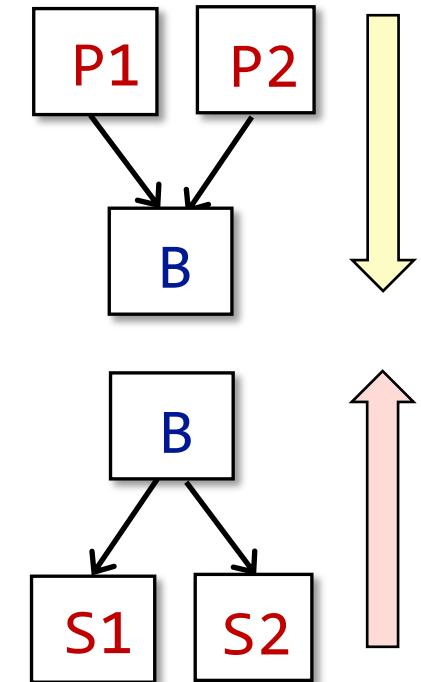
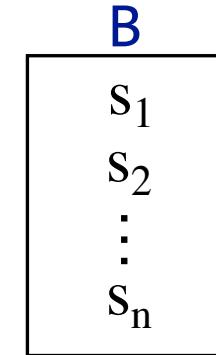
$$\text{OUT}[B] = \text{OUT}[s_n]$$

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$$\text{OUT}[B] = \bigwedge_S S \text{ a successor of } B \text{ } \text{IN}[S]$$



Data Flow Analysis Applications

(I) Reaching Definitions Analysis

(II) Live Variables Analysis

(III) Available Expressions Analysis

Issues Not Covered

- Method Calls
 - Intra-procedural CFG
 - Will be introduced in lecture: Inter-procedural Analysis
- Aliases
 - Variables have no aliases
 - Will be introduced in lecture: Pointer Analysis

Reaching Definitions

A **definition d** at program point p *reaches* a point q if there is a path from p to q such that **d** is not “killed” along that path

Reaching Definitions

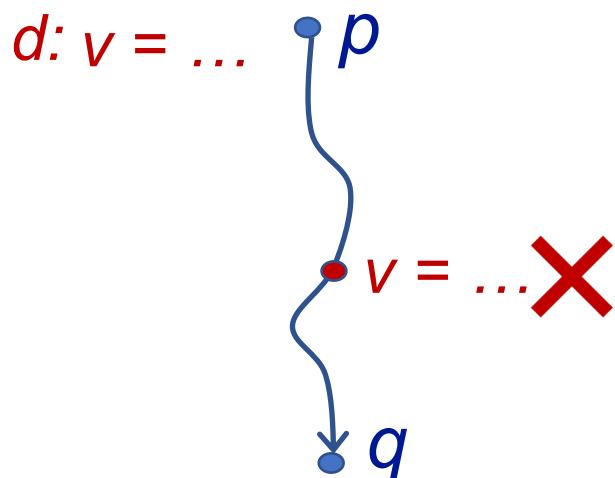
A **definition d** at program point p *reaches* a point q if there is a path from p to q such that **d** is not “killed” along that path

- A **definition of a variable v** is a statement that assigns a value to v

Reaching Definitions

A **definition d** at program point p *reaches* a point q if there is a path from p to q such that **d** is not “killed” along that path

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- Translated as: definition of variable v at program point p reaches point q if there is a path from p to q such that no new definition of v appears on that path



Reaching Definitions

A **definition d** at program point p *reaches* a point q if there is a path from p to q such that **d** is not “killed” along that path

- A **definition of a variable v** is a statement that assigns a value to v
- Translated as: definition of variable v at program point p reaches point q if there is a path from p to q such that no new definition of v appears on that path
- Reaching definitions can be used to **detect possible undefined variables**. e.g., introduce a dummy definition for each variable v at the entry of CFG, and if the dummy definition of v reaches a point p where v is used, then v may be used before definition (*as undefined reaches v*)

Understanding Reaching Definitions

- Data Flow Values/Facts
 - The definitions of all the variables in a program

Abstraction

Understanding Reaching Definitions

- Data Flow Values/Facts
 - The definitions of all the variables in a program
 - Can be represented by bit vectors
- e.g., D1, D2, D3, D4, ..., D100 (100 definitions)

A diagram showing a sequence of bits: 00000...0. A bracket underneath the first 100 bits is labeled "100 bits".

Bit i from the left represents definition Di

Abstraction

Understanding Reaching Definitions

Safe-approximation

- Transfer Function
- Control Flow

Understanding Reaching Definitions

D: $v = x \ op \ y$

Safe-approximation

This statement “**generates**” a definition D of variable v and “**kills**” all the other definitions in the program that define variable v, while leaving the remaining incoming definitions unaffected.

- Transfer Function
- Control Flow

Understanding Reaching Definitions

D: $v = x \ op \ y$

Safe-approximation

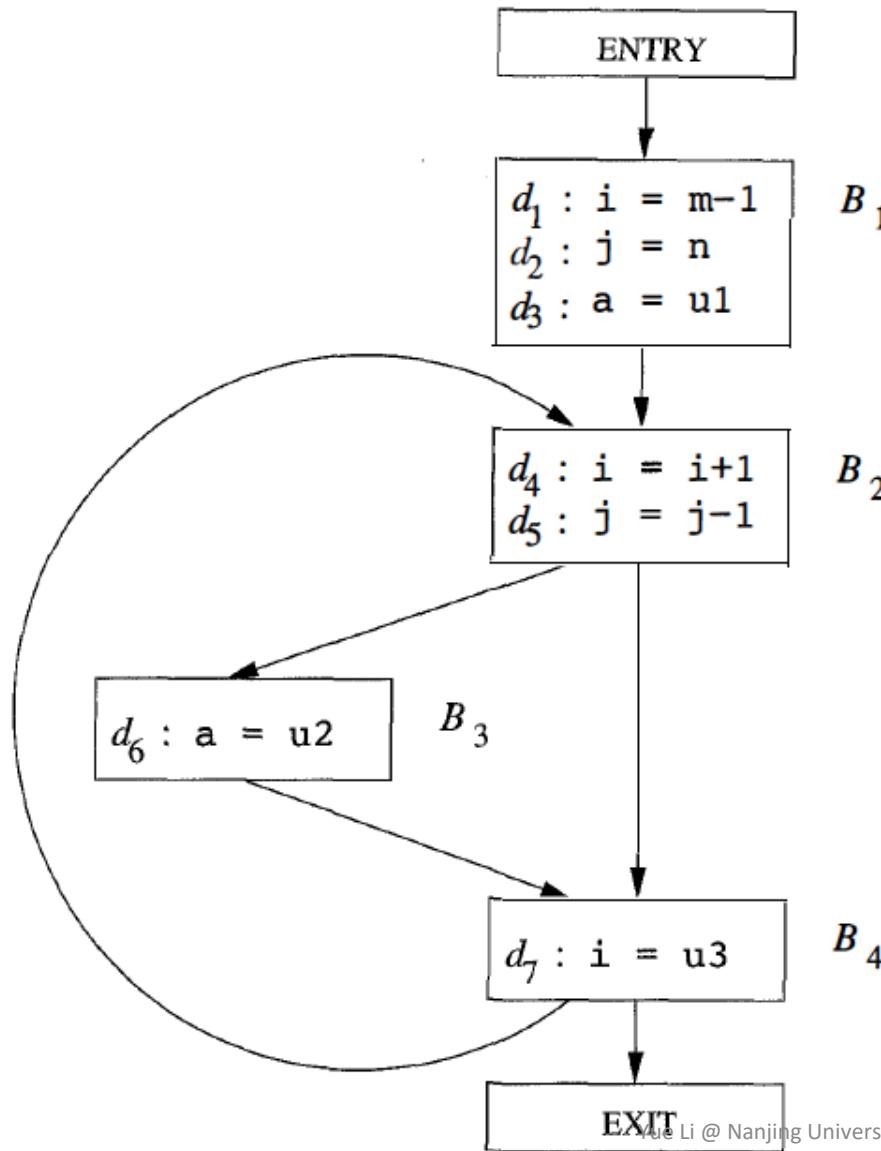
This statement “**generates**” a definition D of variable v and “**kills**” all the other definitions in the program that define variable v, while leaving the remaining incoming definitions unaffected.

- Transfer Function

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$

- Control Flow

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



$$\text{gen}_{B_1} = \{ d_1, d_2, d_3 \}$$

$$\text{kill}_{B_1} = \{ d_4, d_5, d_6, d_7 \}$$

$$\text{gen}_{B_2} = \{ d_4, d_5 \}$$

$$\text{kill}_{B_2} = \{ d_1, d_2, d_7 \}$$

$$\text{gen}_{B_3} = \{ d_6 \}$$

$$\text{kill}_{B_3} = \{ d_3 \}$$

$$\text{gen}_{B_4} = \{ d_7 \}$$

$$\text{kill}_{B_4} = \{ d_1, d_4 \}$$

Understanding Reaching Definitions

D: $v = x \ op \ y$

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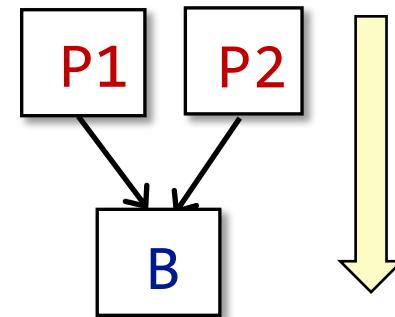
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$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

- Control Flow

$$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P]$$



Understanding Reaching Definitions

D: $v = x \ op \ y$

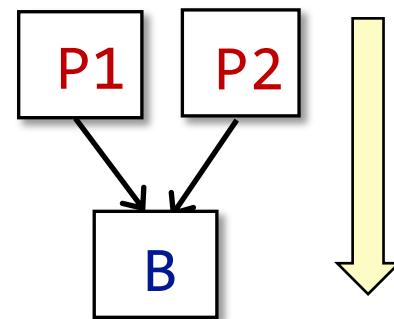
Safe-approximation

This statement “**generates**” a definition D of variable v and “**kills**” all the other definitions in the program that define variable v, while leaving the remaining incoming definitions unaffected.

A definition reaches a program point as long as there exists at least one path along which the definition reaches.

- Control Flow

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Understanding Reaching Definitions

D: $v = x \ op \ y$

Safe-approximation

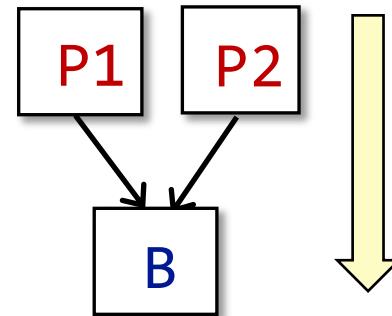
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- Control Flow

$$\text{IN}[B] = \bigcup_{P \text{ a predecessor of } B} \text{OUT}[P]$$



Algorithm of Reaching Definitions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[B] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P];$   
            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }  
    }
```

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    }
```

Why entry is excluded?

Algorithm of Reaching Definitions Analysis

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Algorithm of Reaching Definitions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

$OUT[entry] = \emptyset;$

for (each basic block $B \setminus entry$)

$OUT[B] = \emptyset;$

while (changes to any OUT occur) ?

for (each basic block $B \setminus entry$) {

$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P];$

$OUT[B] = gen_B \cup (IN[B] - kill_B);$

}

*Why this iterative algorithm
can finally stop?*

Algorithm of Reaching Definitions Analysis

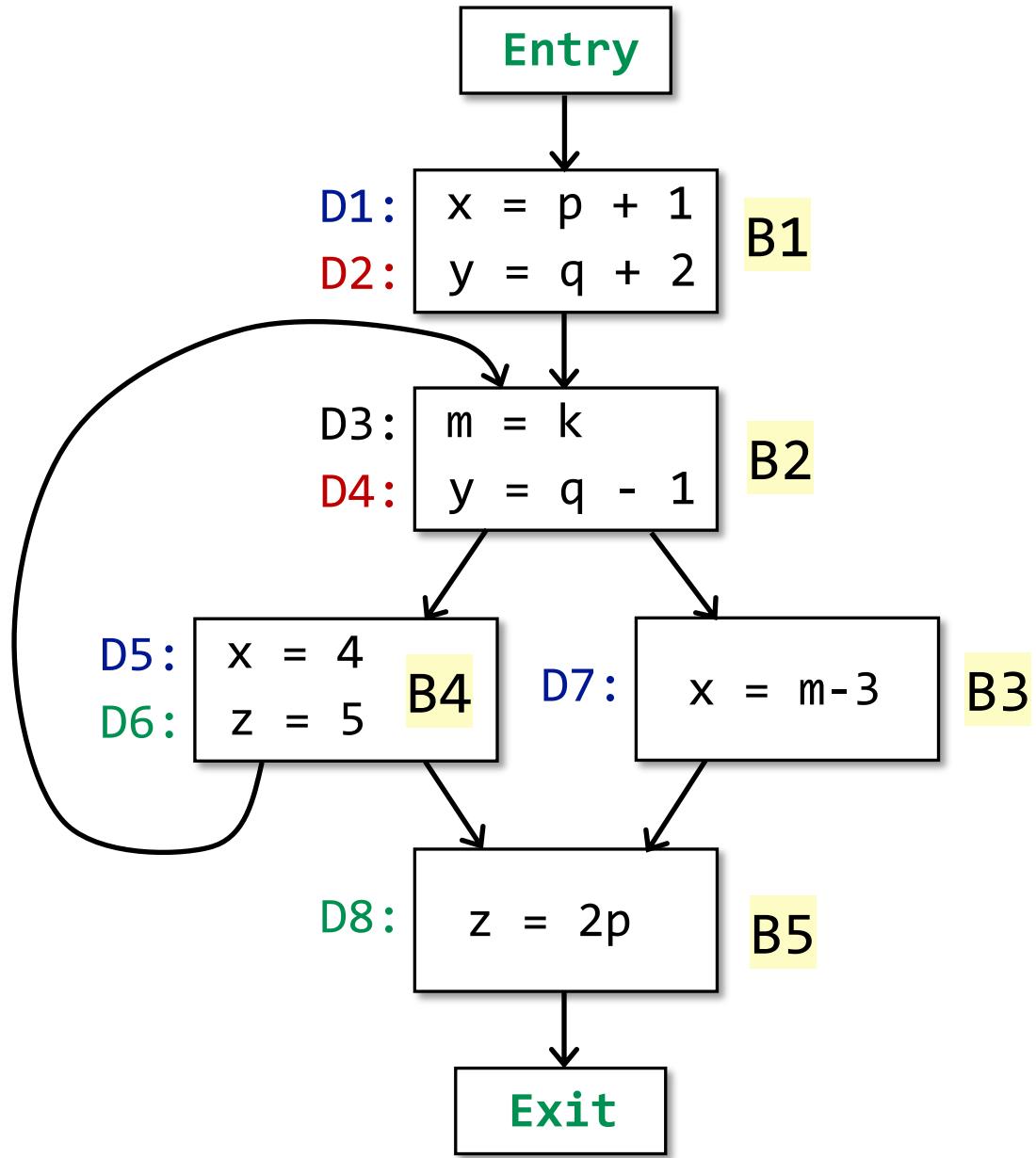
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        }
```



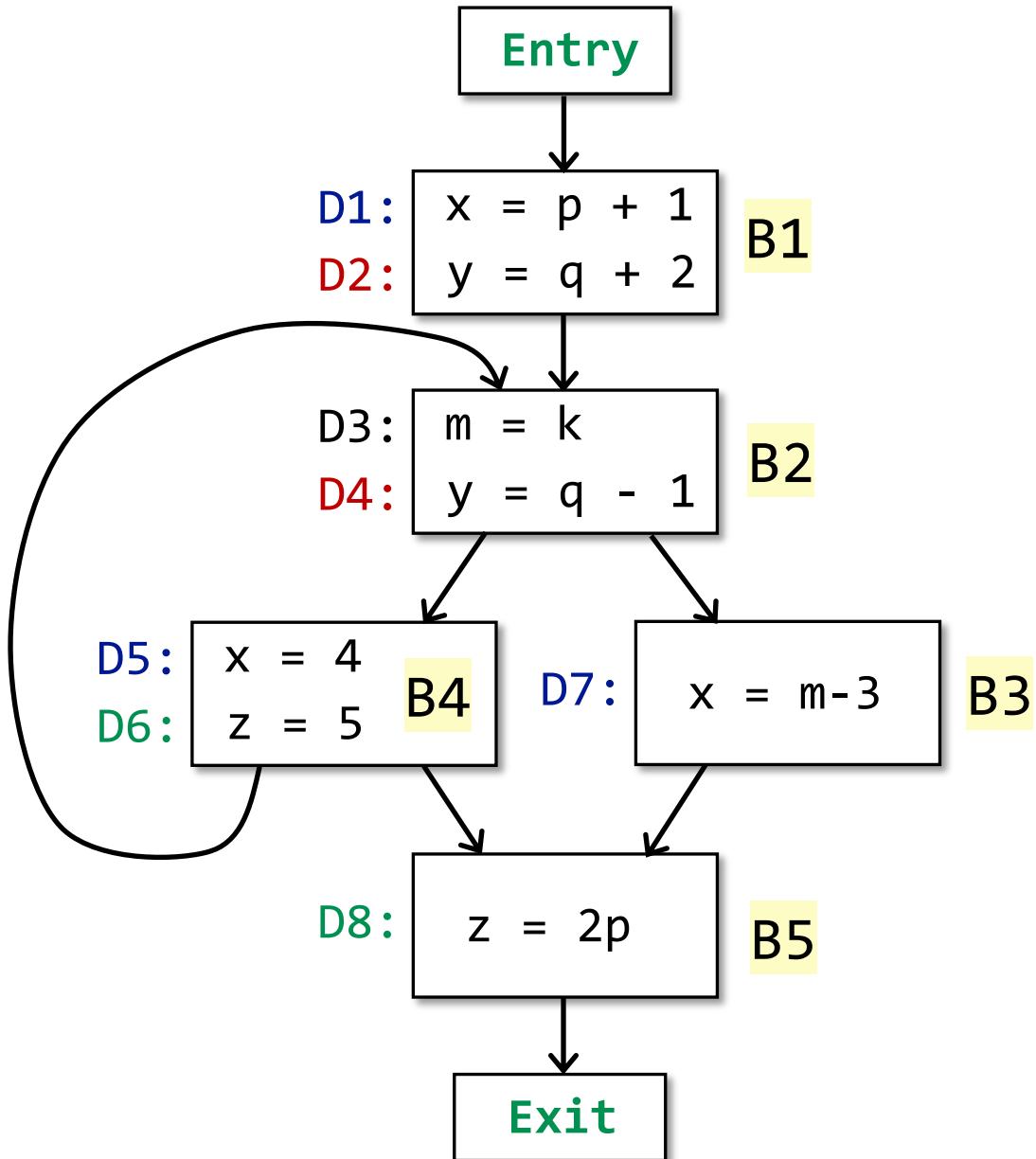


```

D1
D2
do {
    D3
    D4
    if(...) {
        D5
        D6
    } else {
        D7
        break
    }
} while(...)

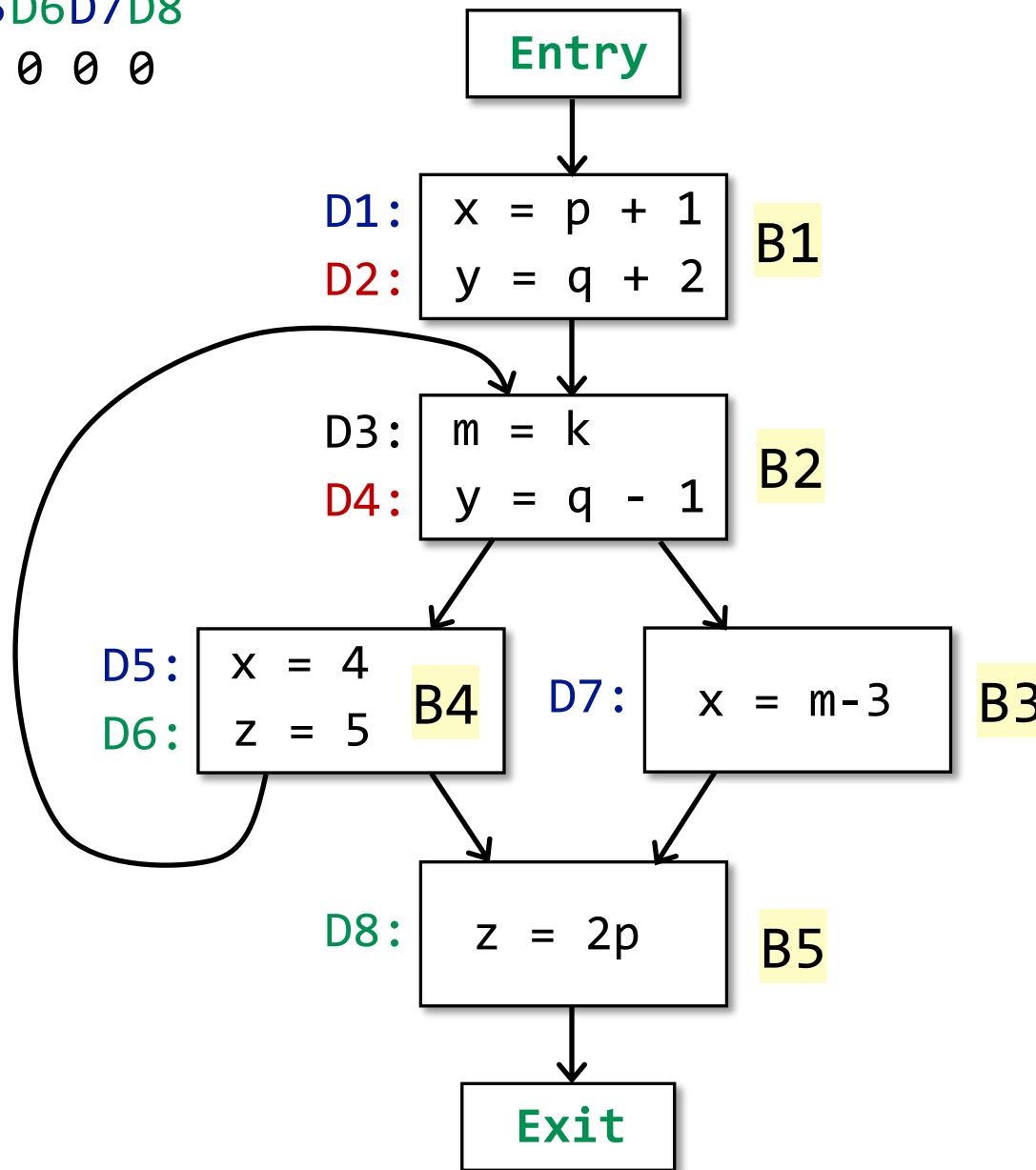
D8

```



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0 0



Algorithm of Reaching Definitions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

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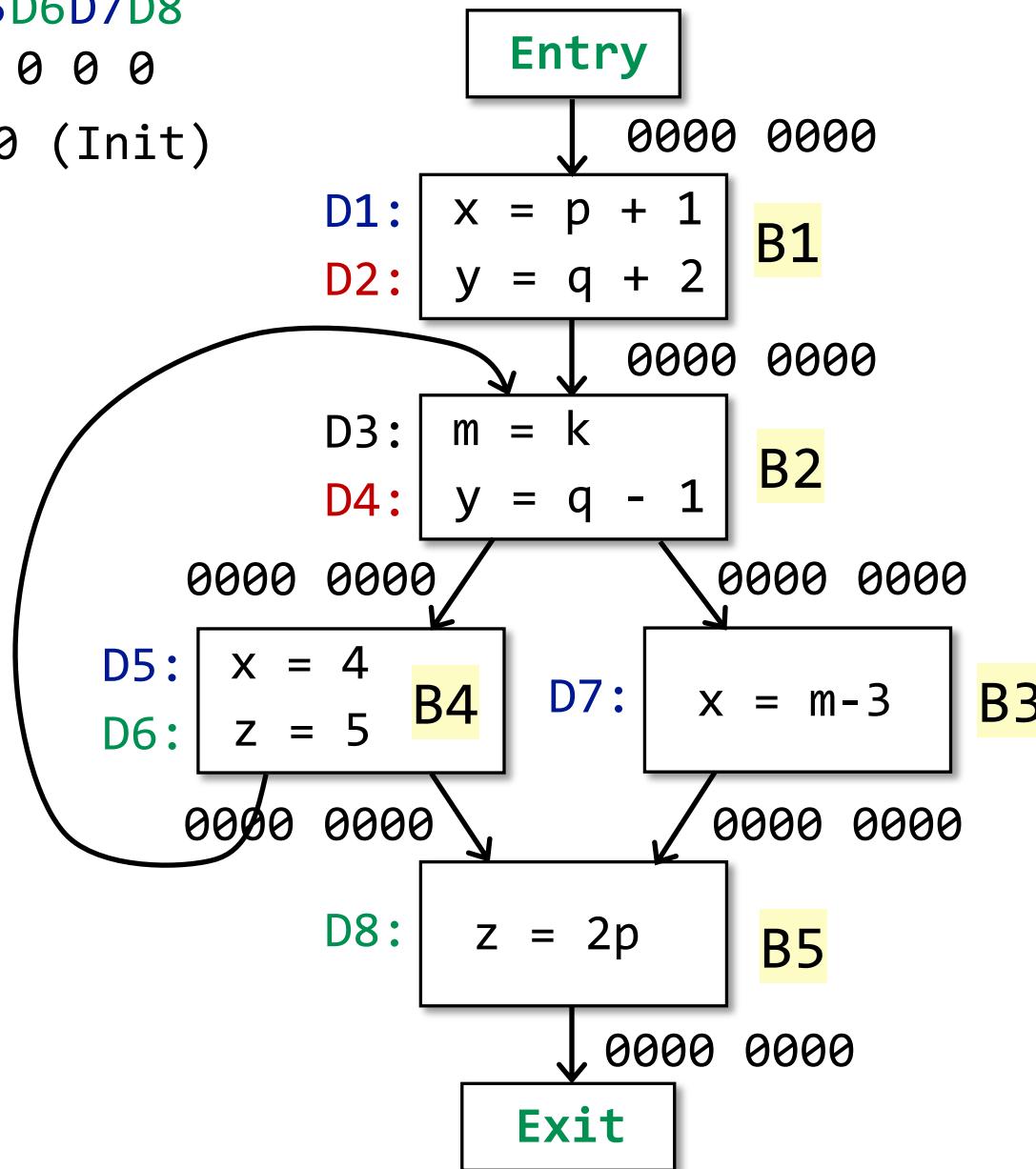
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    for (each basic block  $B \setminus entry$ ) {  
        IN[B] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P];$   
        OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
    }
```



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)



Algorithm of Reaching Definitions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[B] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P];$   
            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }
```

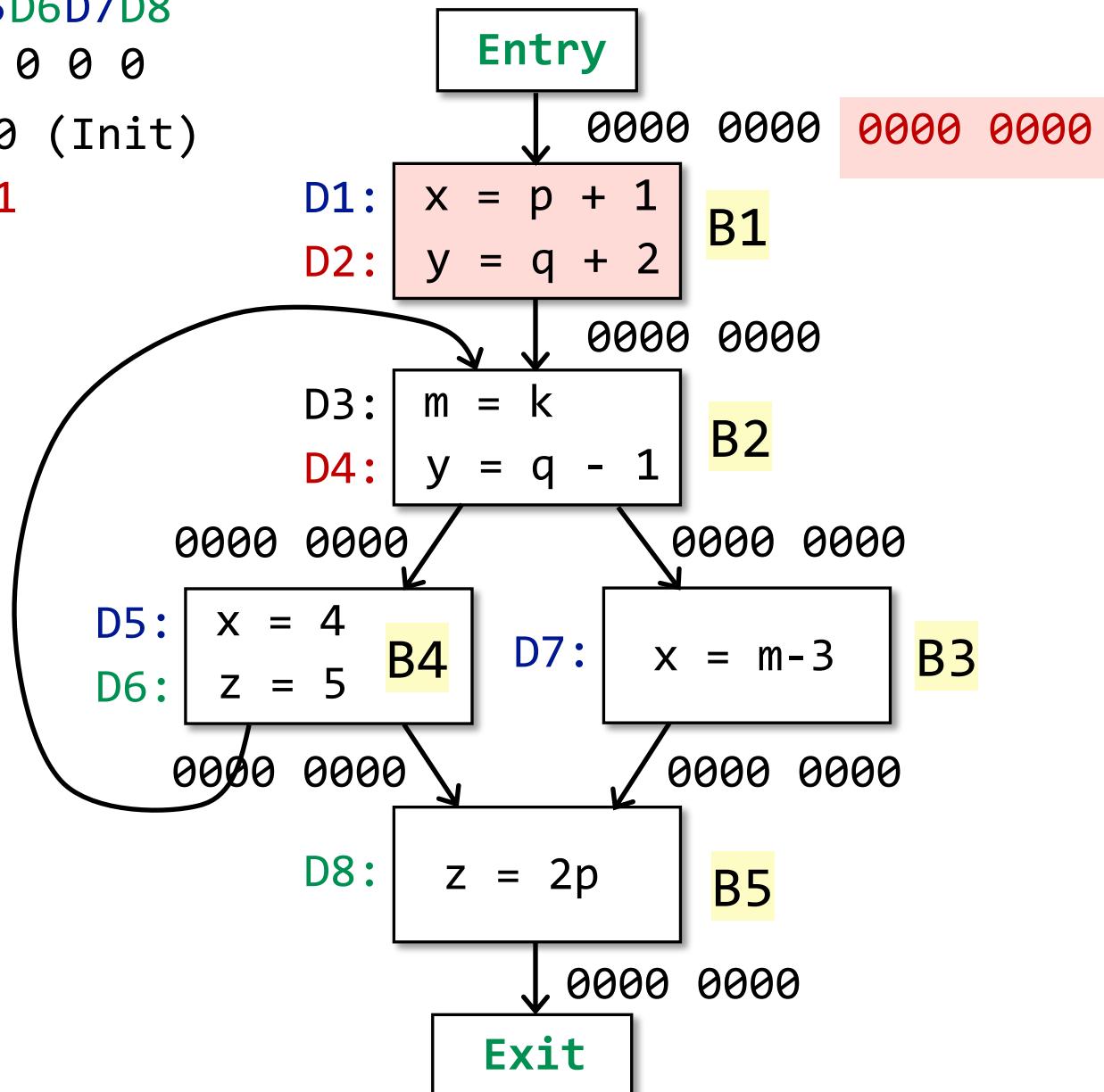


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1



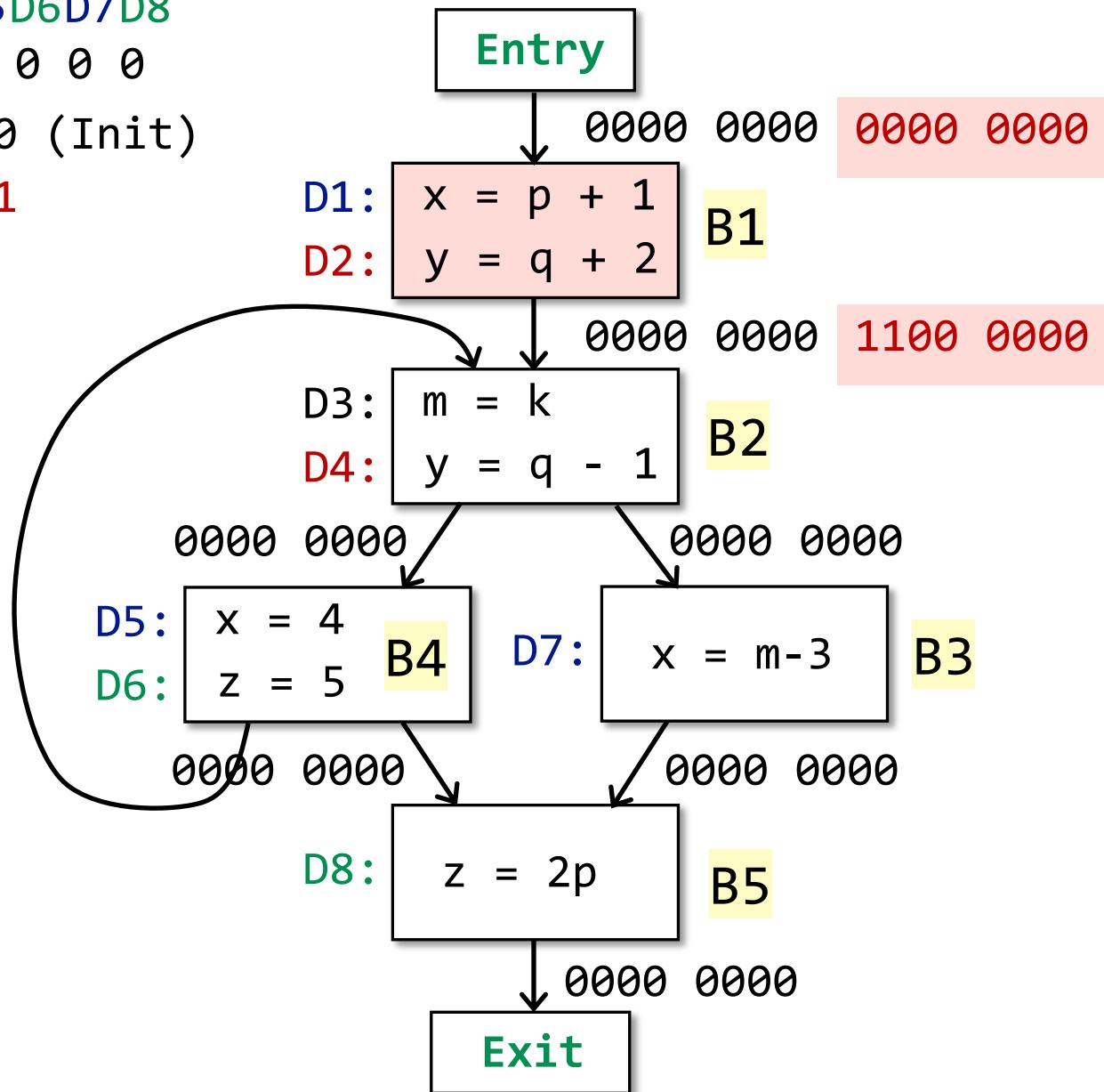
$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B);$ @ Nanjing University

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

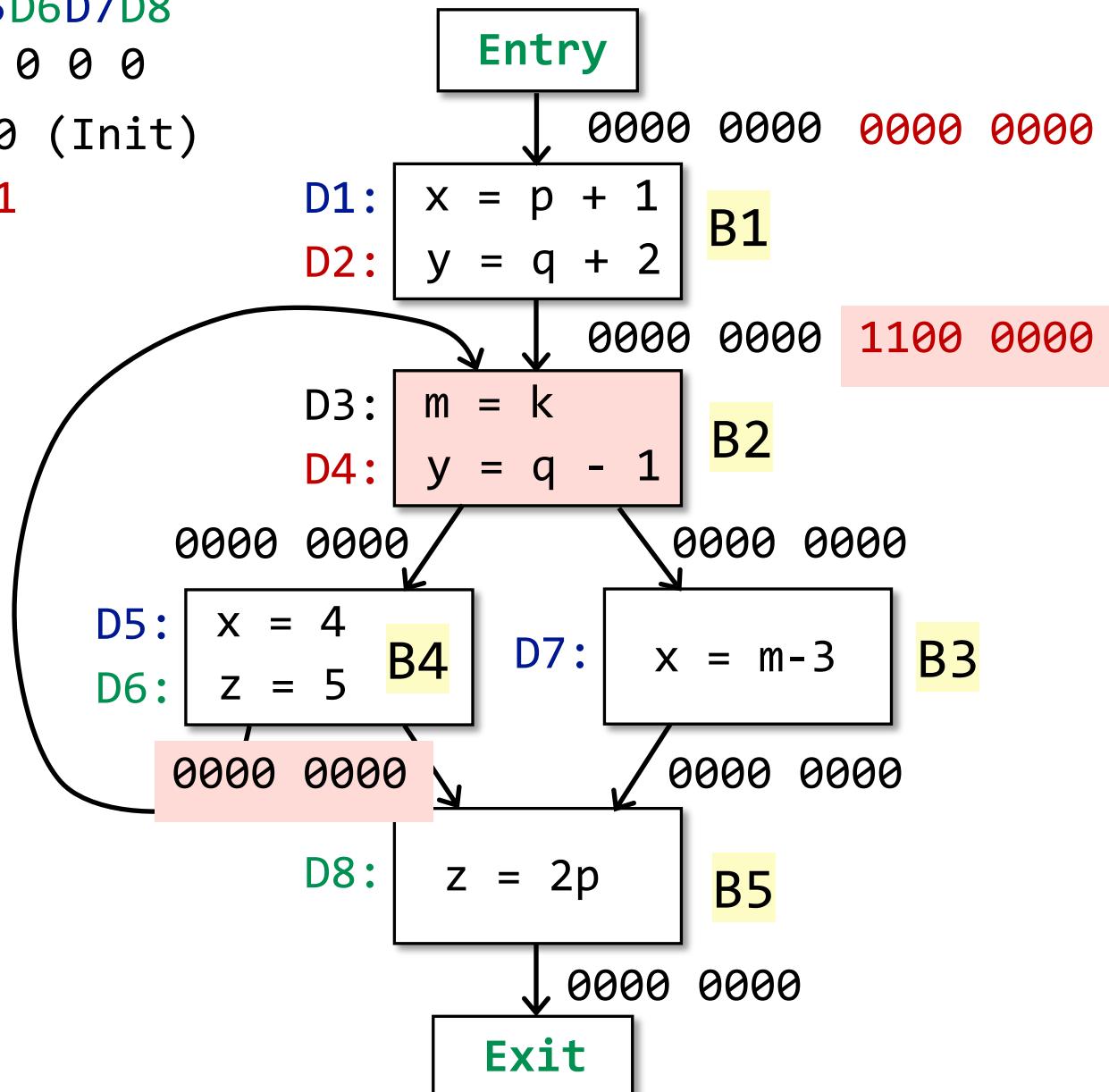


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1



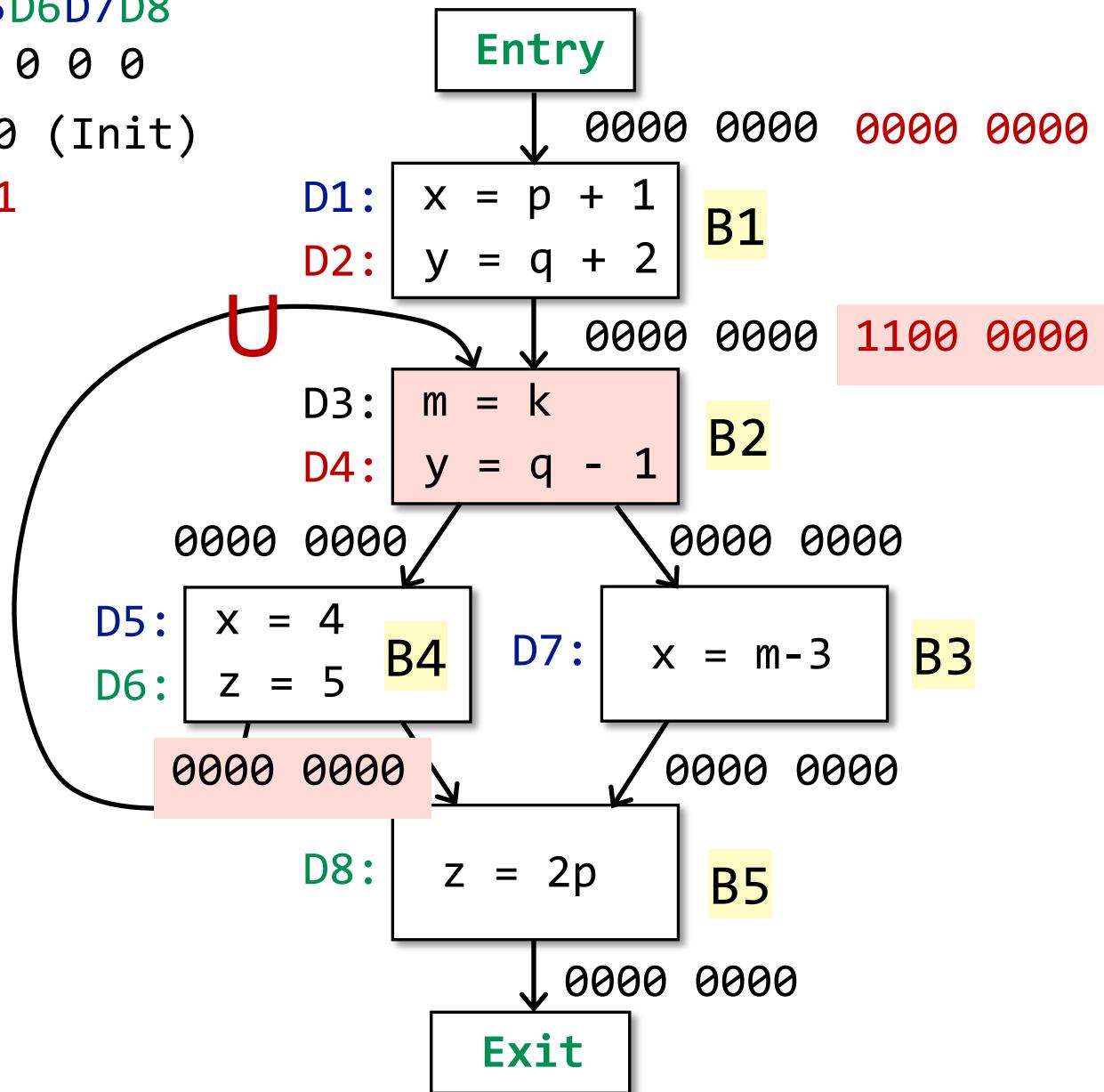
$$\text{IN}[B] = \bigcup_{P \text{ a predecessor of } B} \text{OUT}[P];$$

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

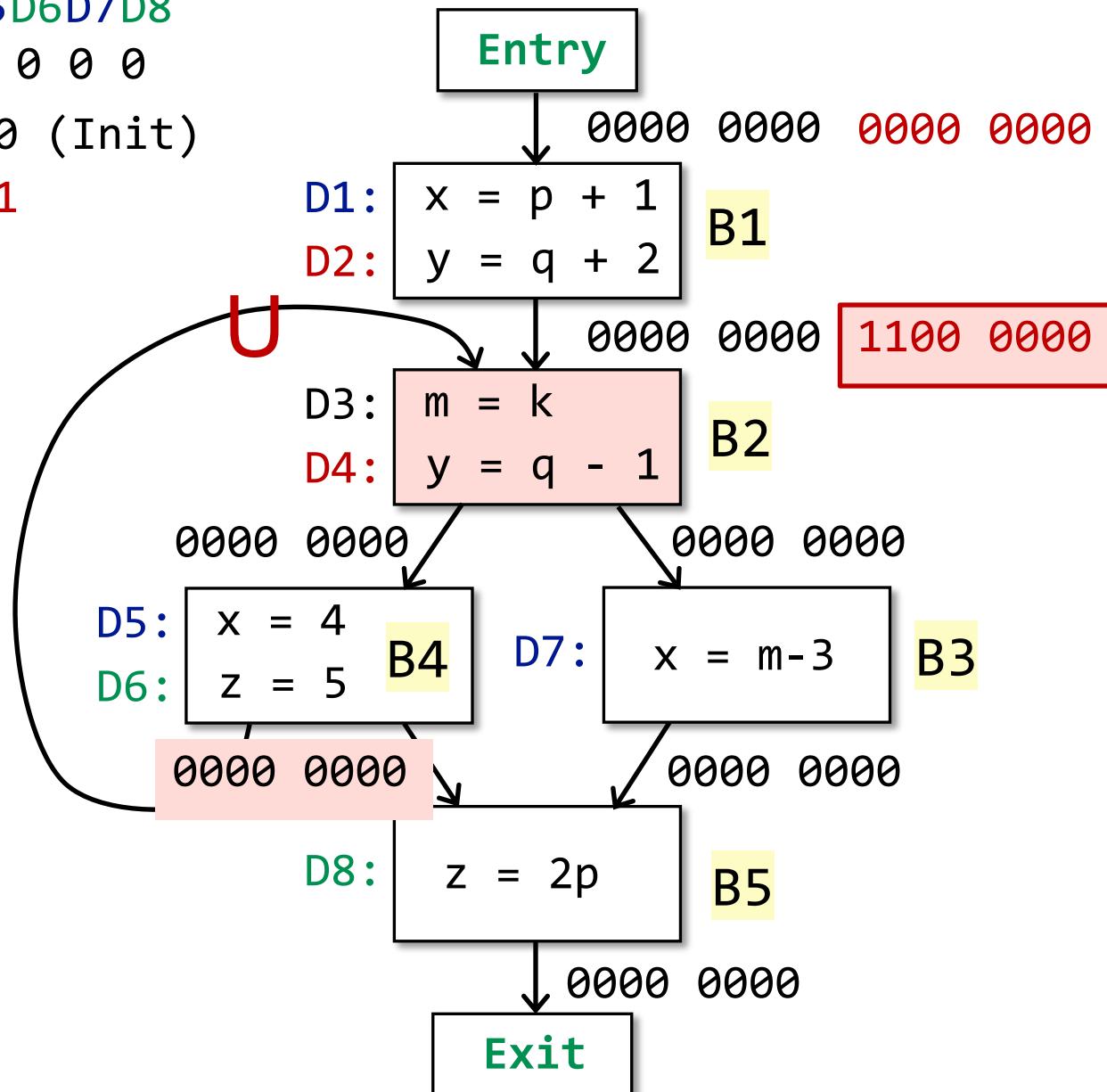


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

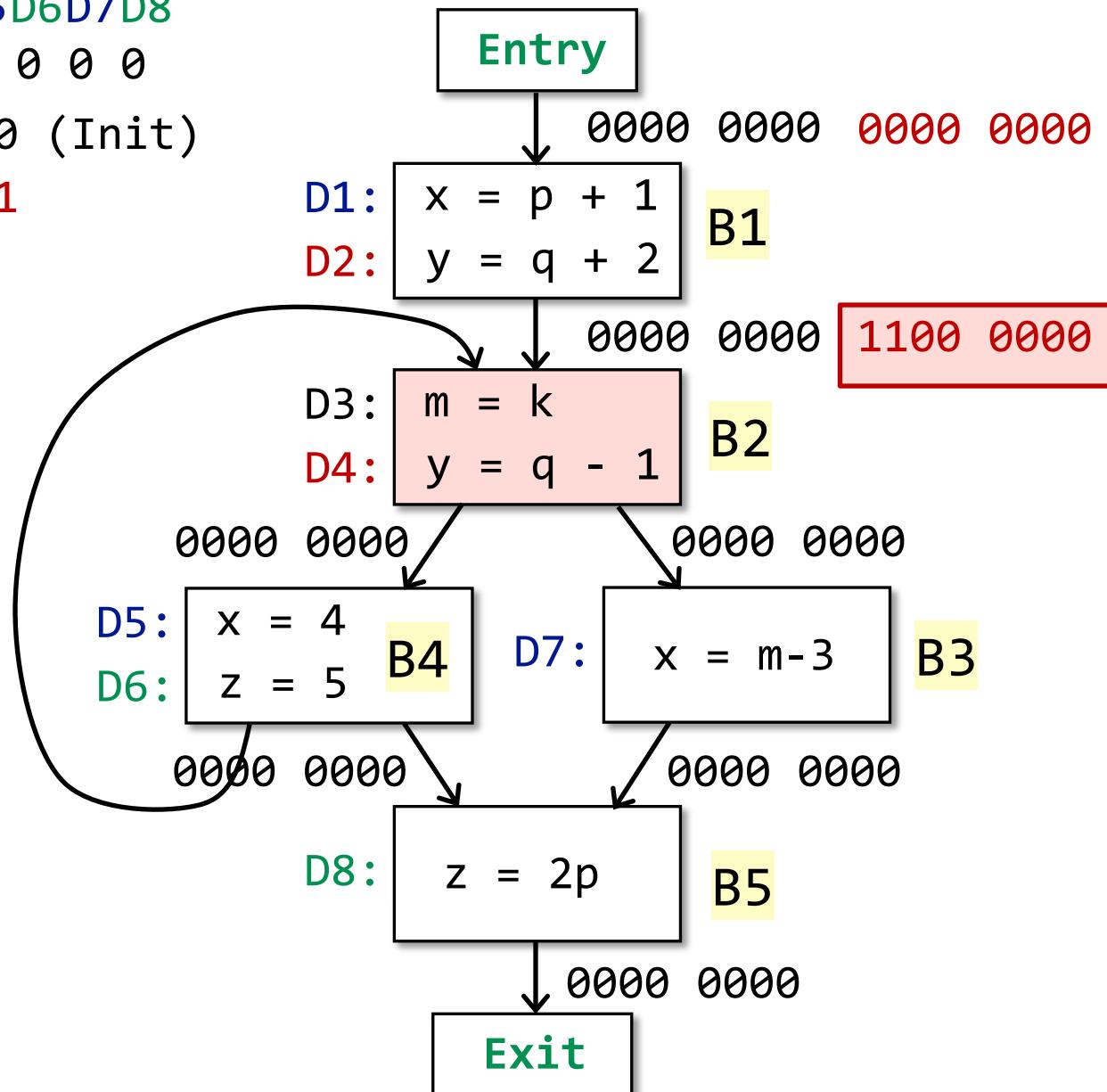


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

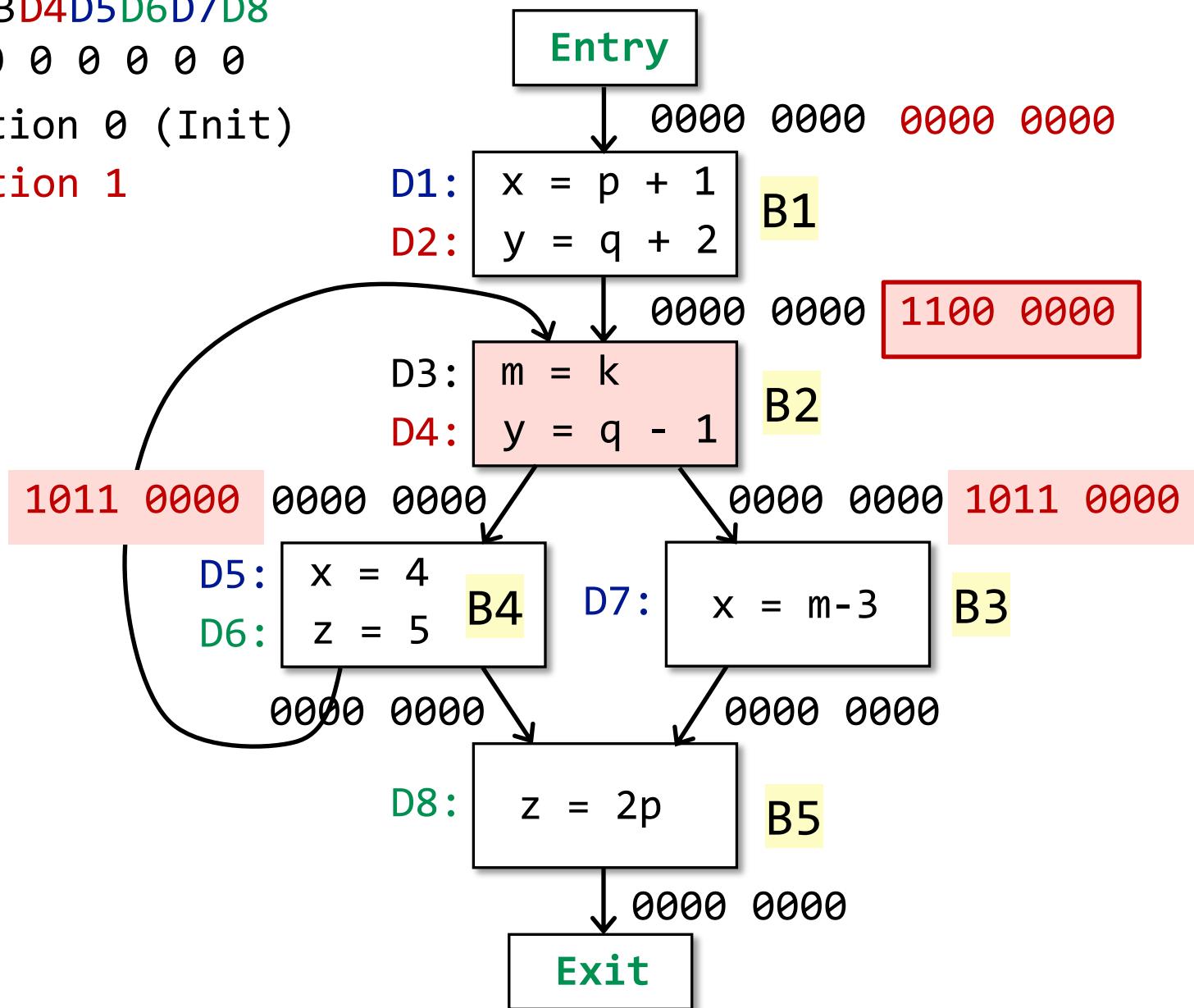


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

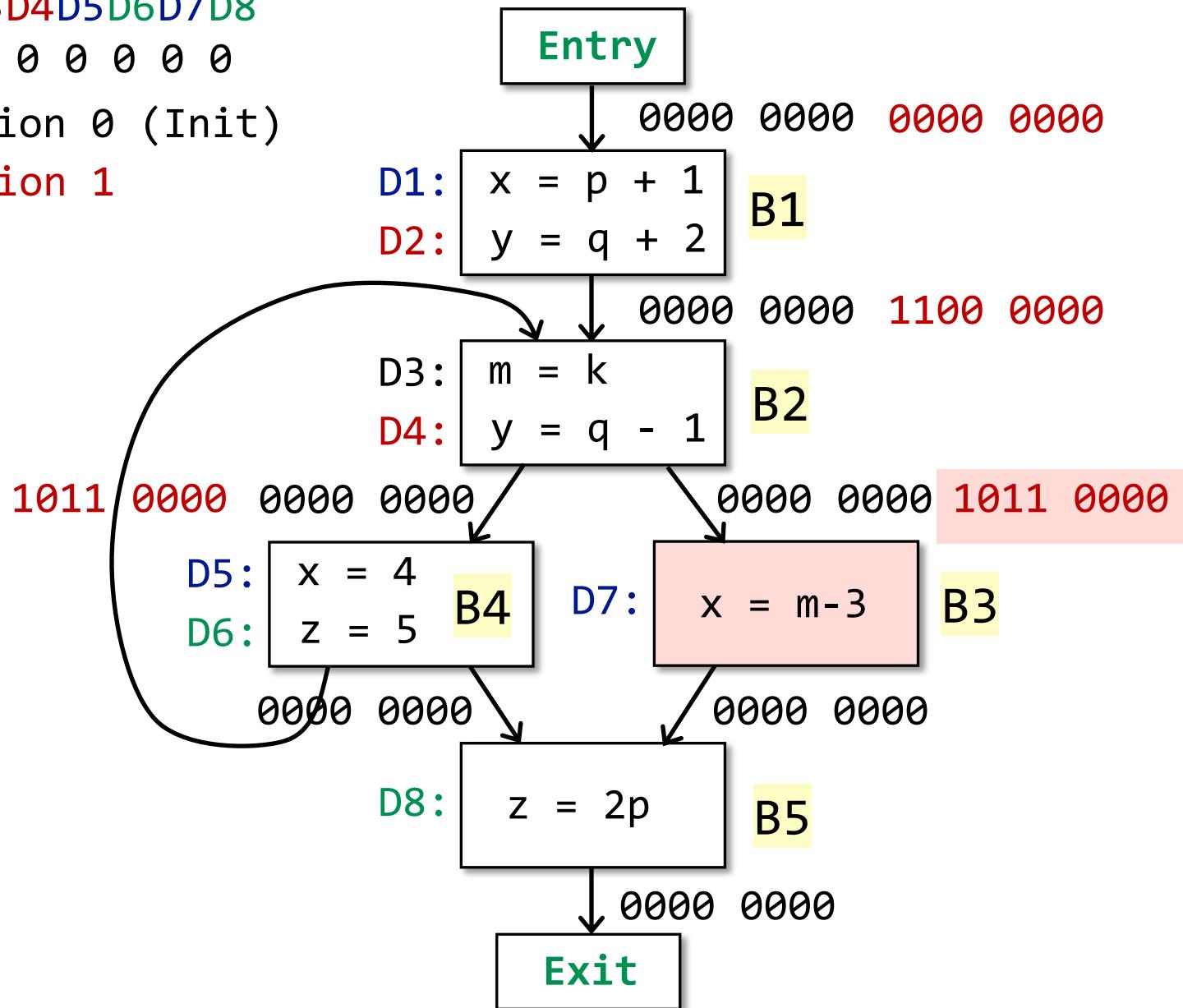


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

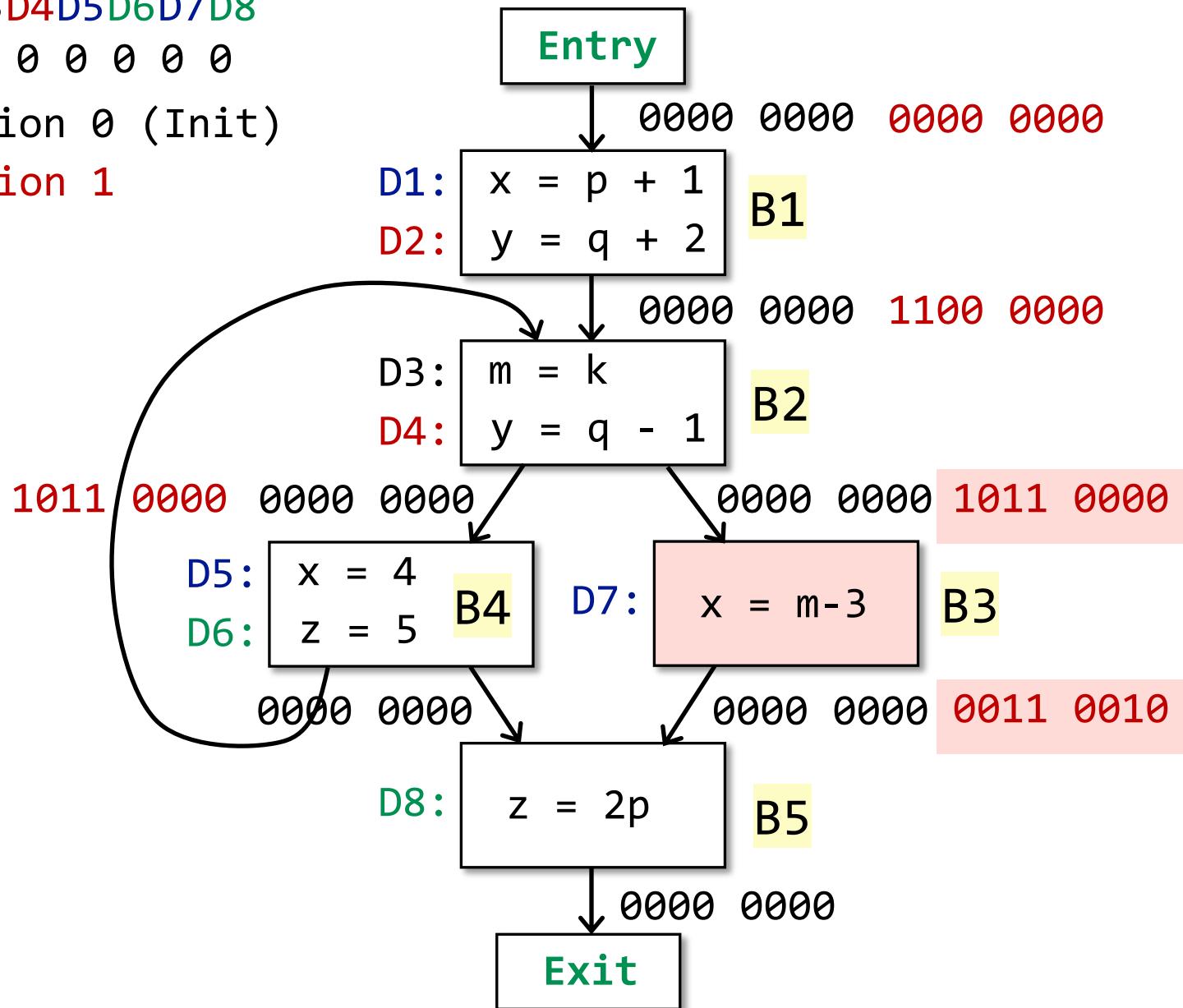


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

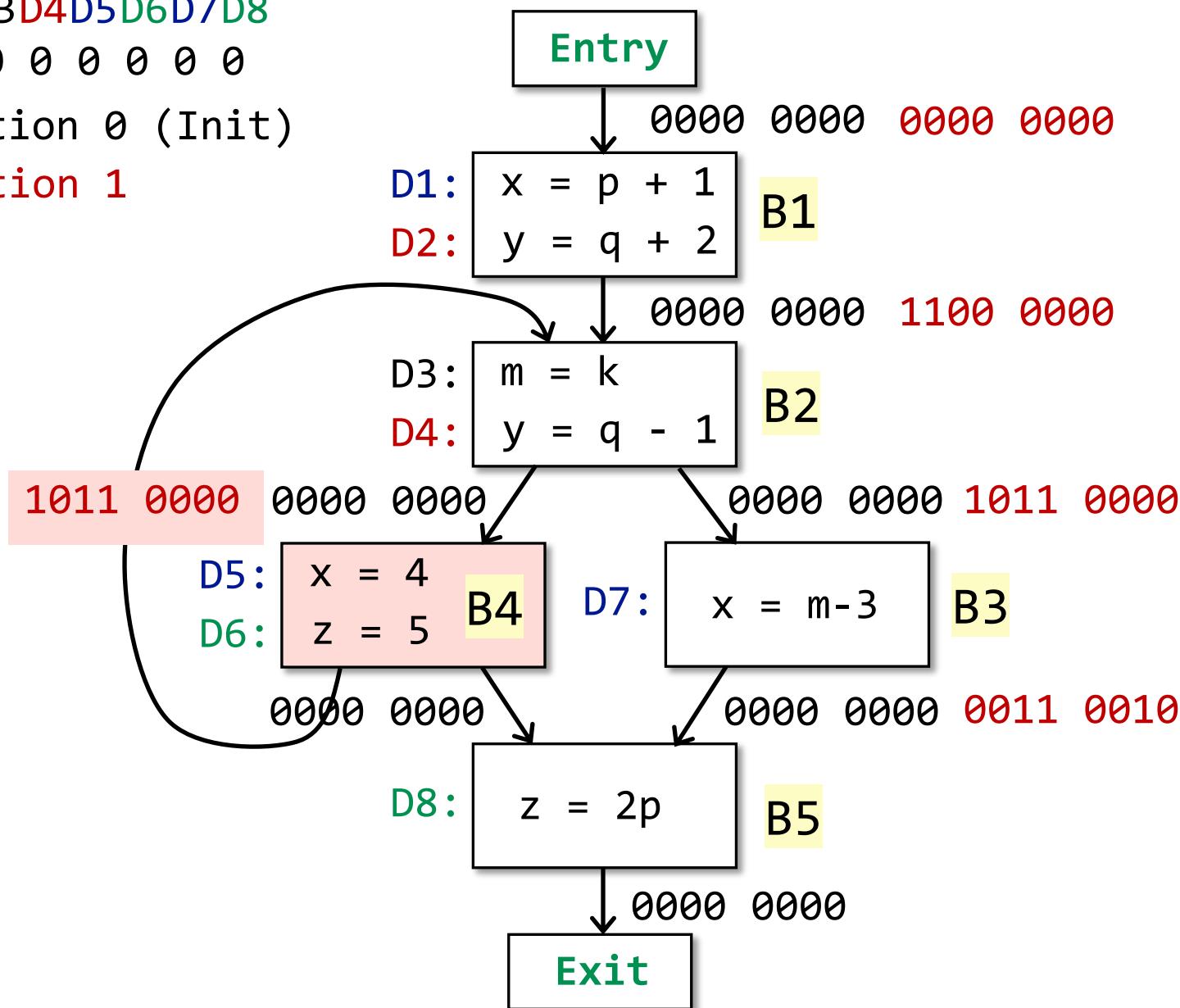


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

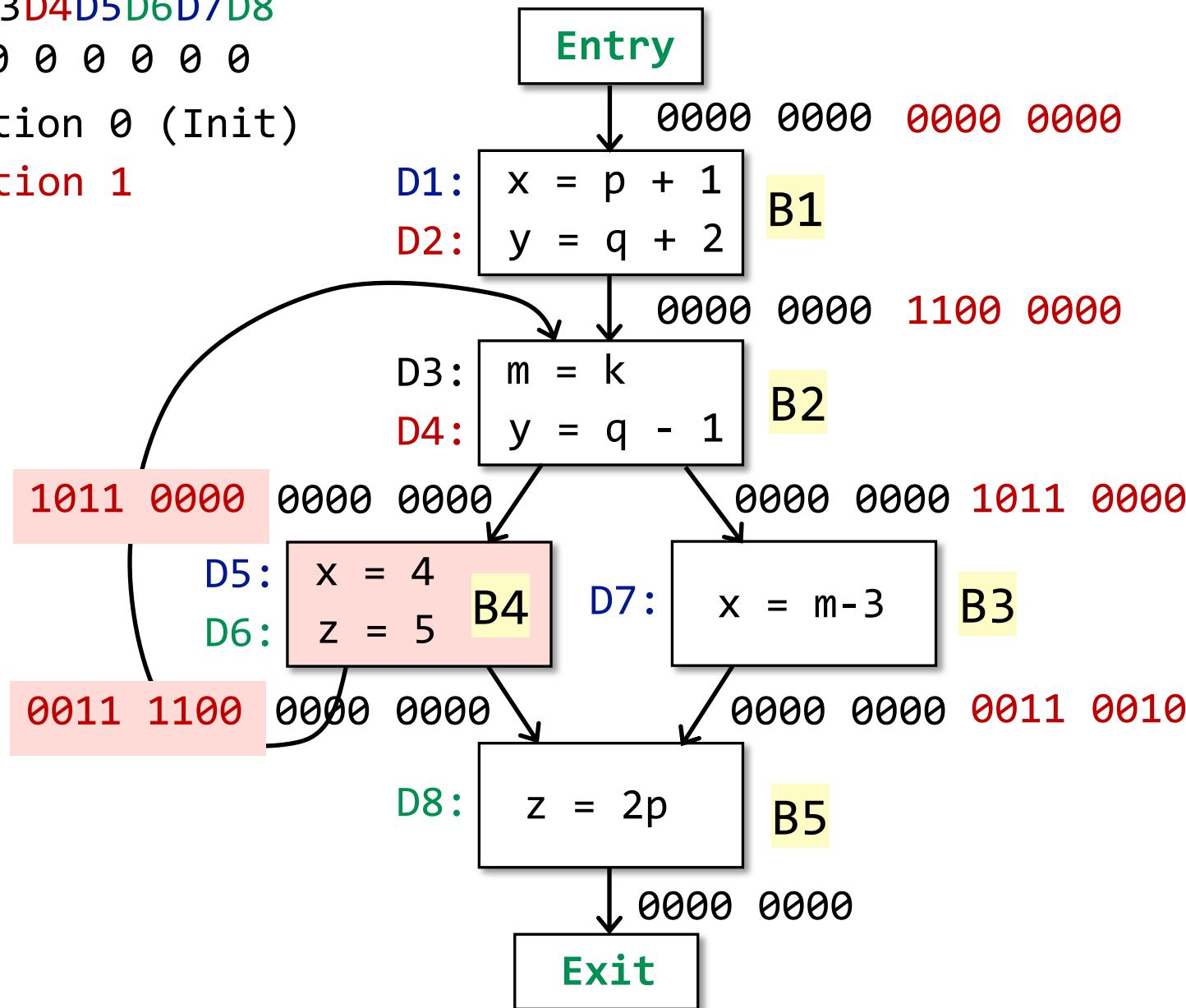


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

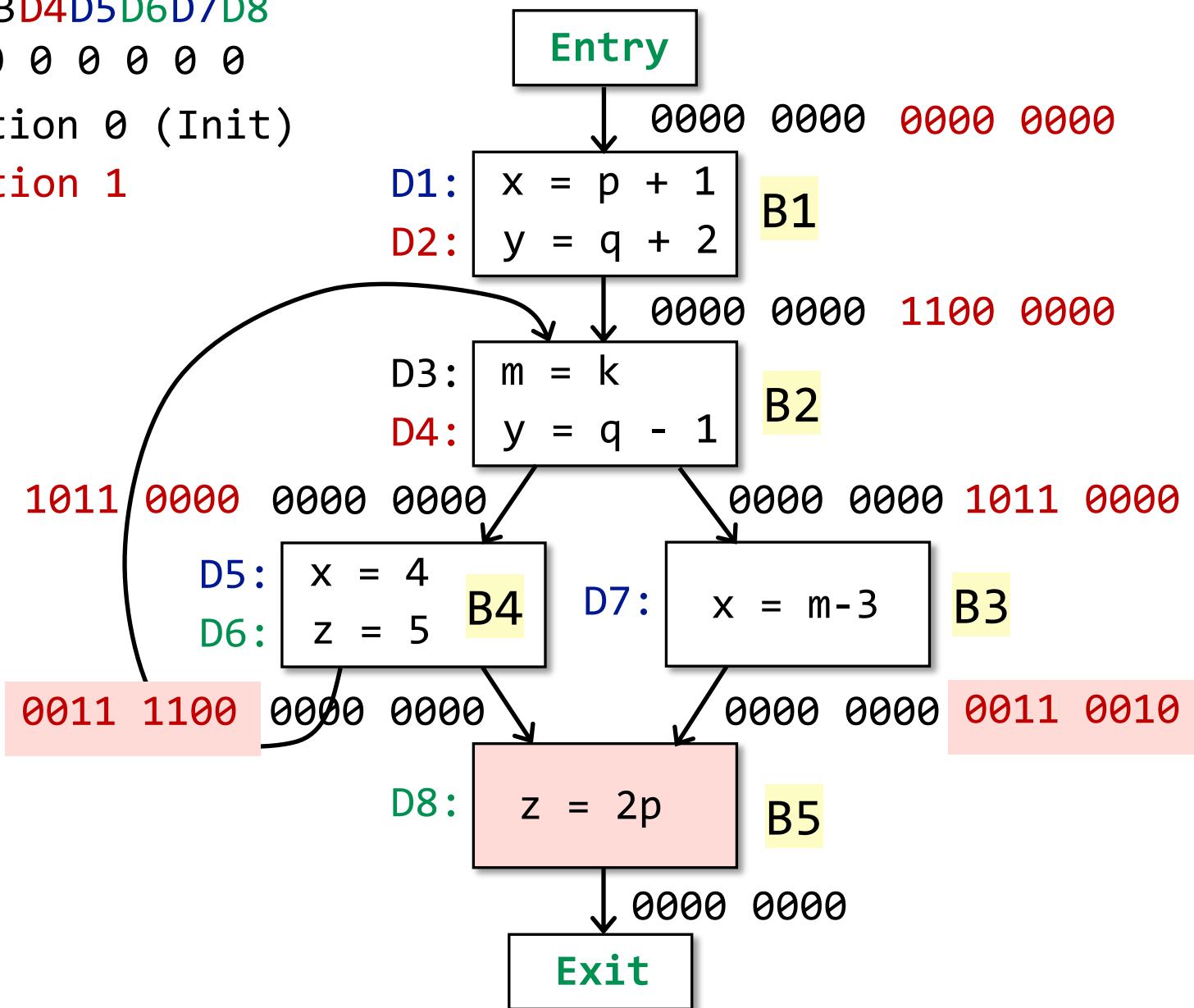


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

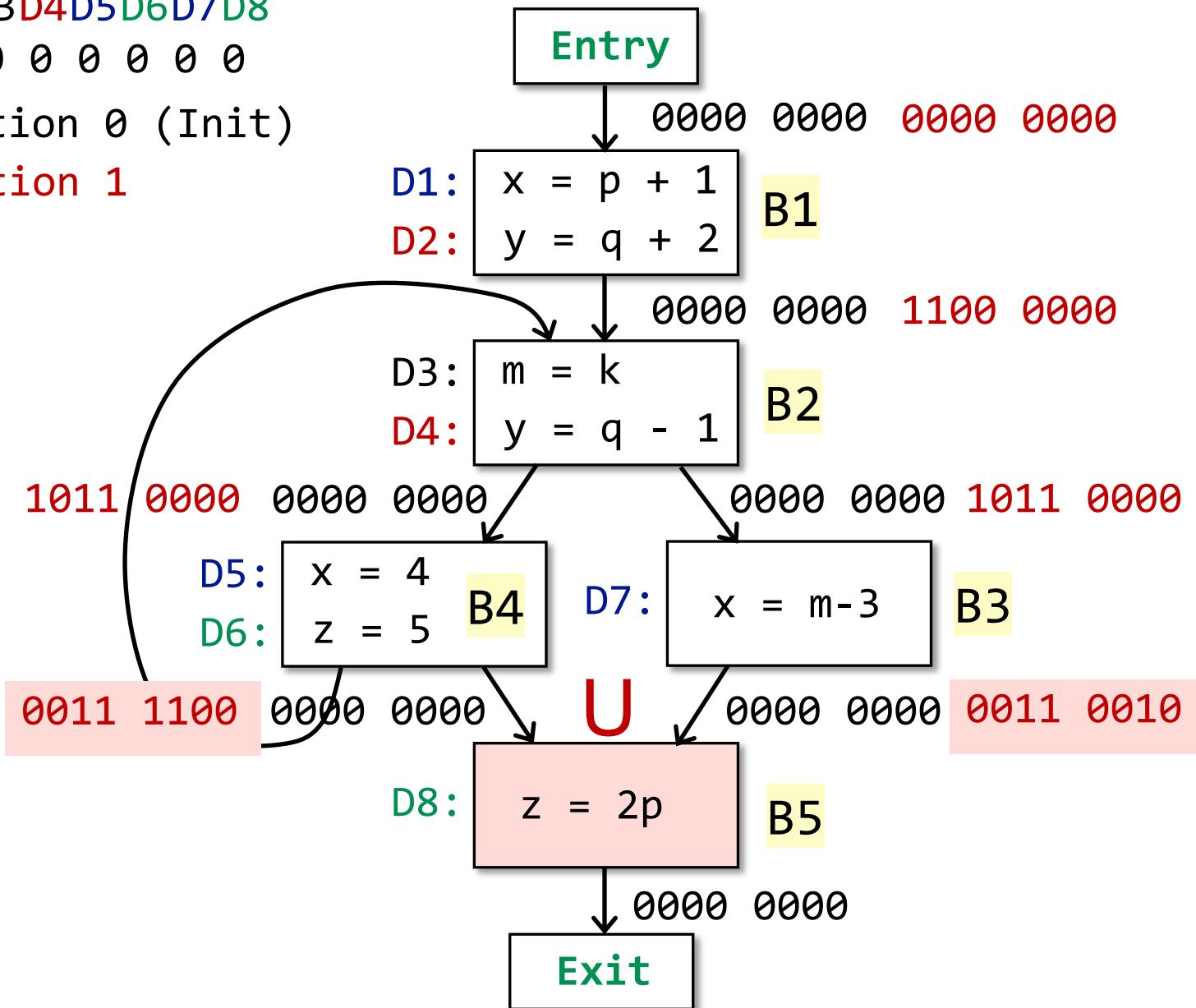


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

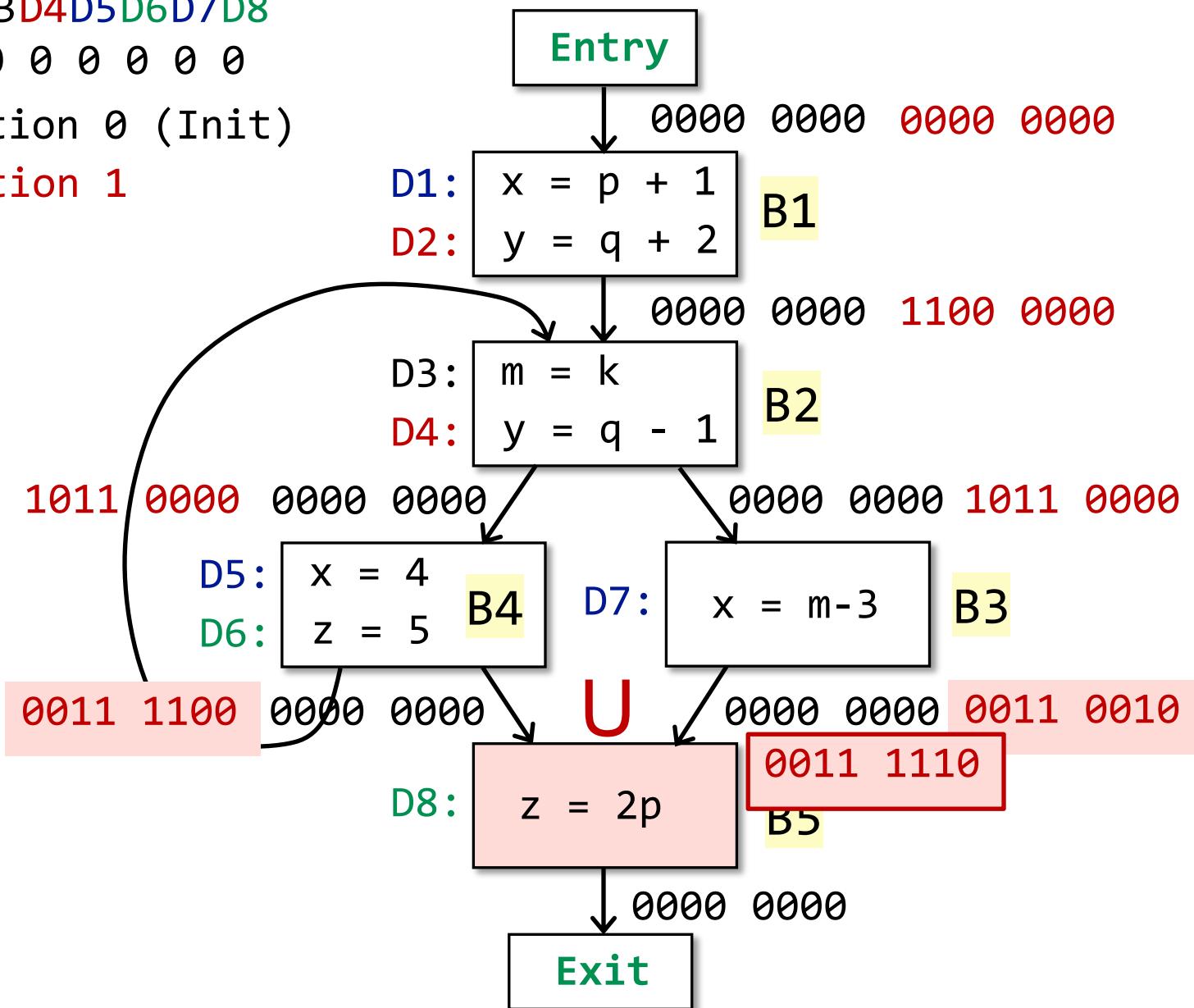


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

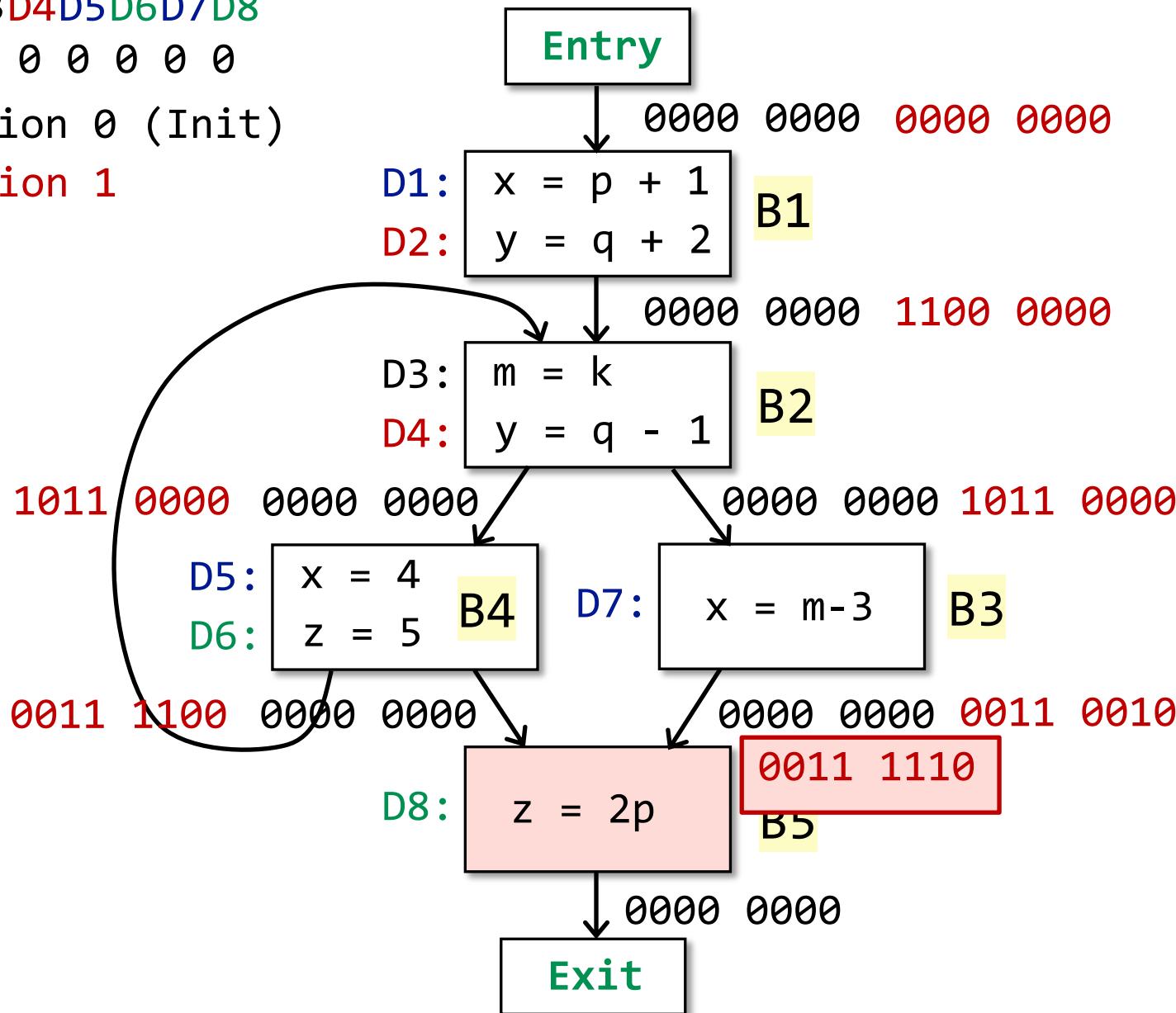


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

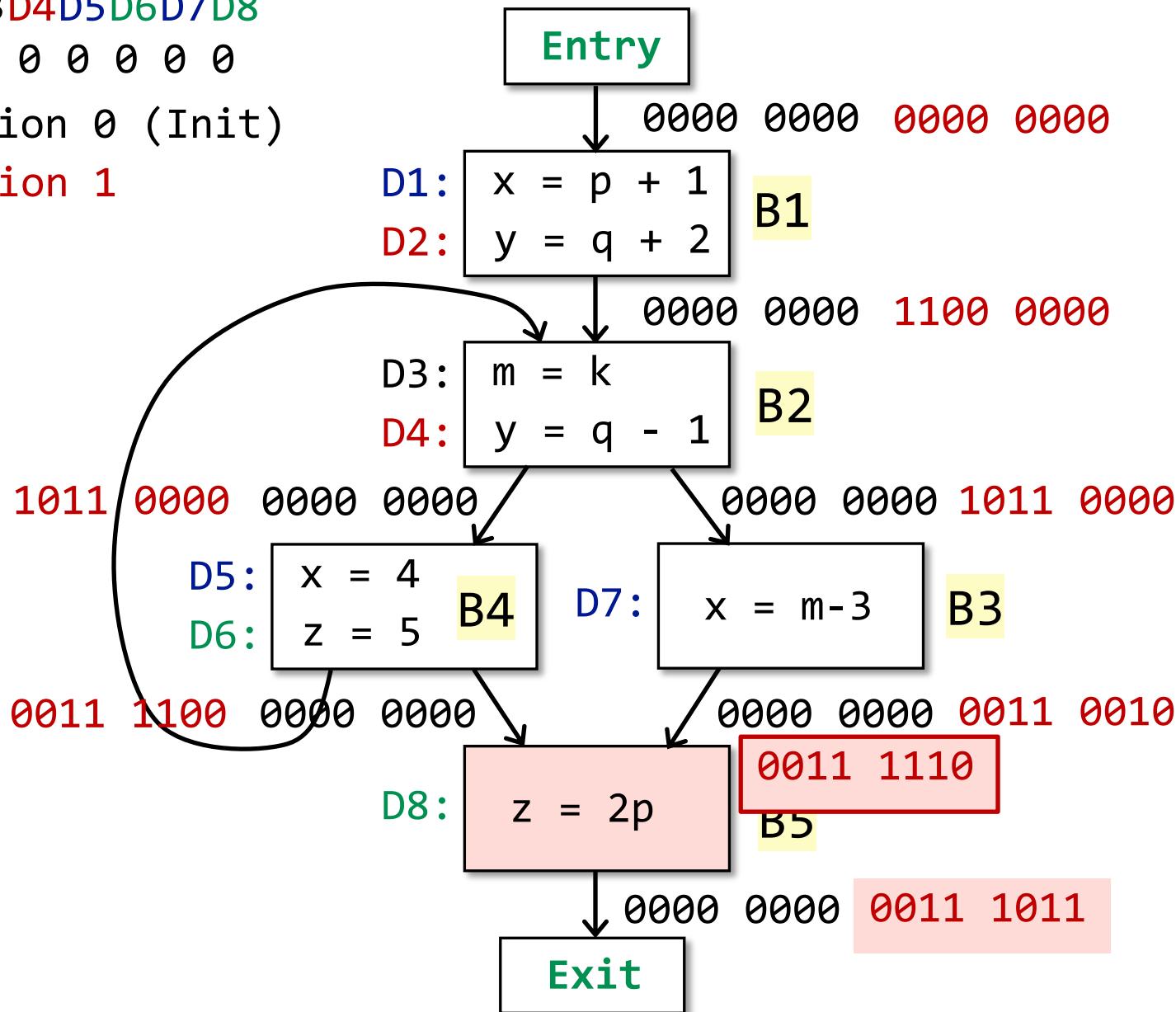


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

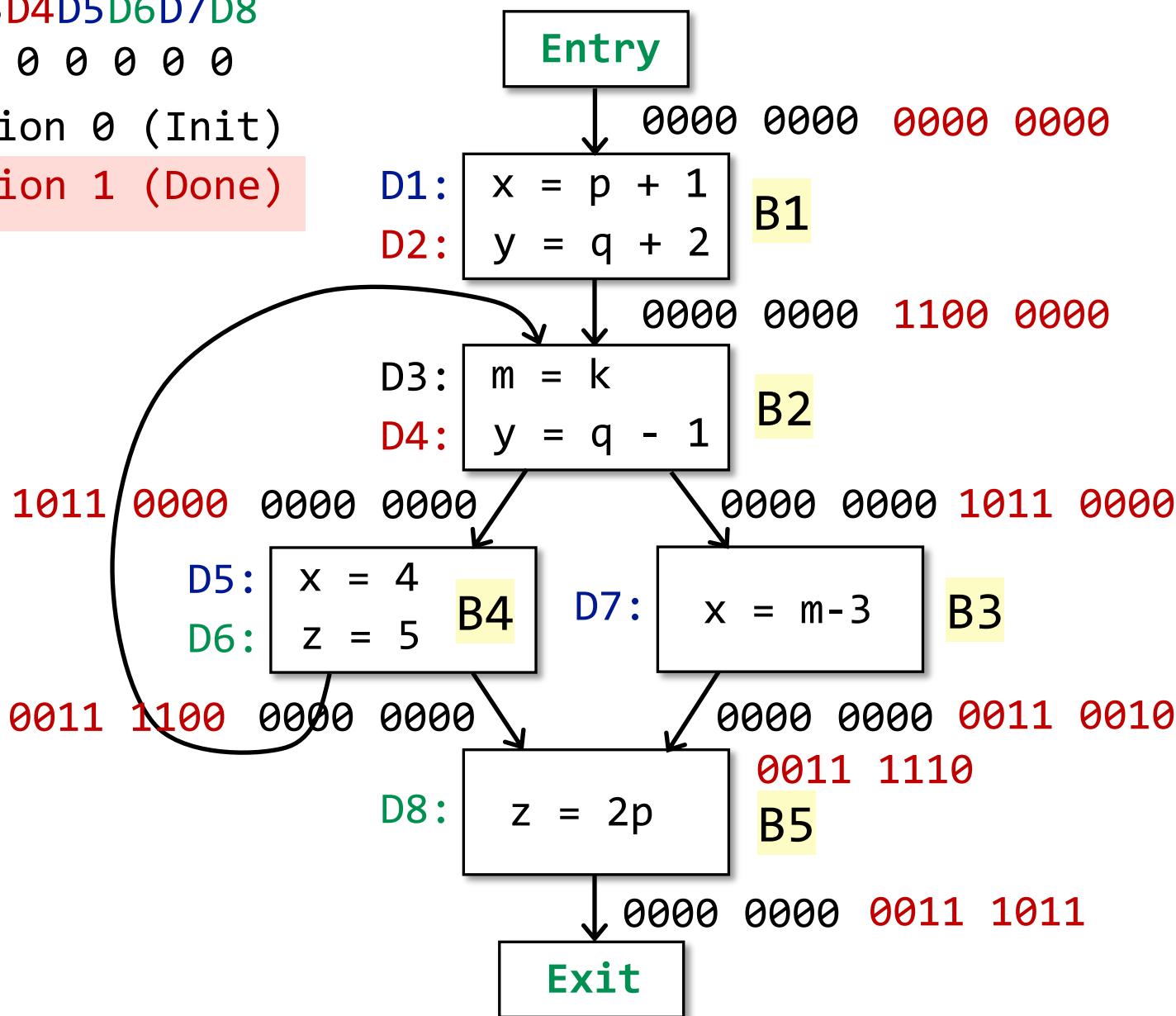


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)



Algorithm of Reaching Definitions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

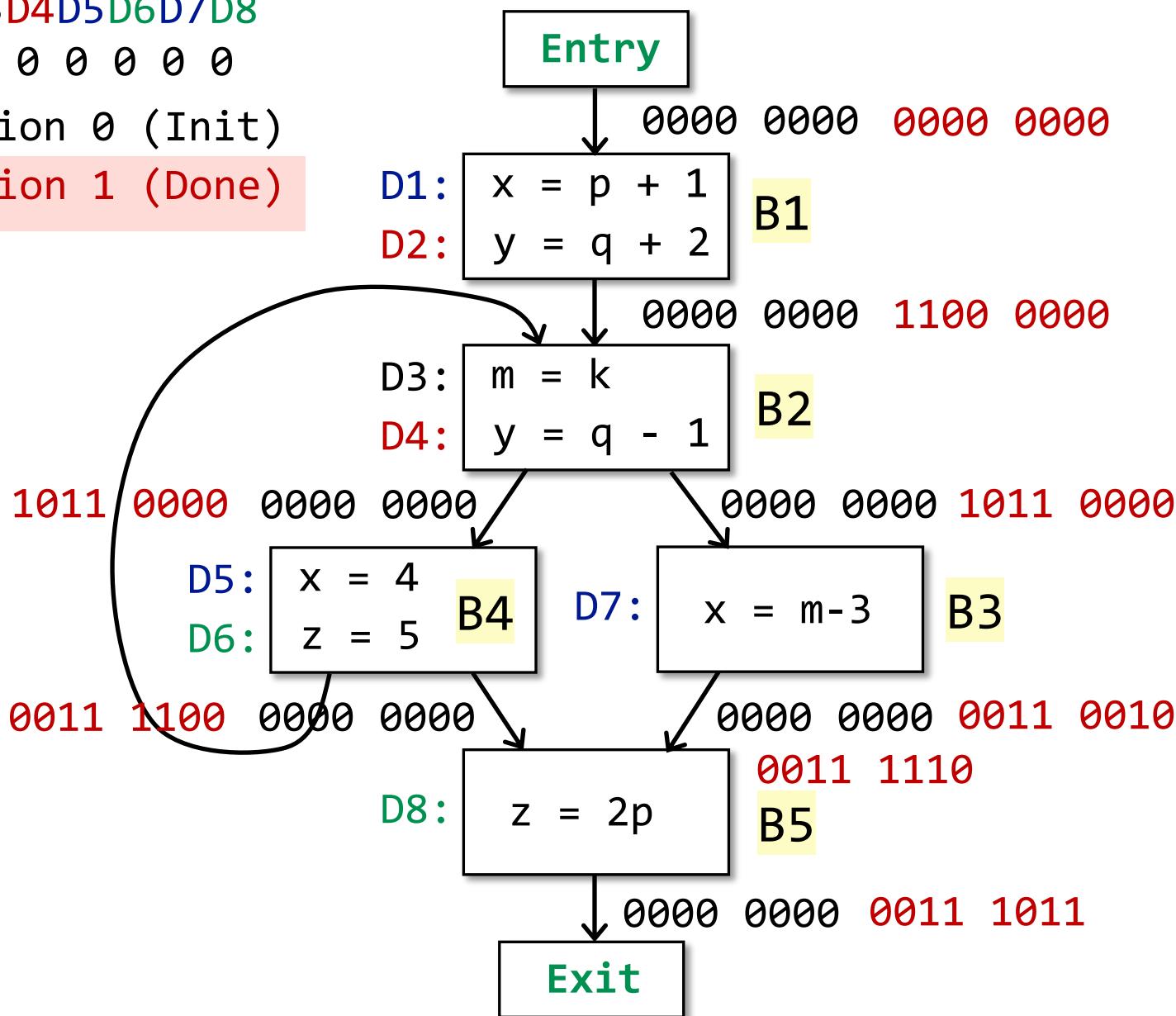
```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[ $B$ ] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[ $B$ ] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P]$ ;  
            OUT[ $B$ ] =  $gen_B \cup (IN[B] - kill_B)$ ;  
        }  
    }
```

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

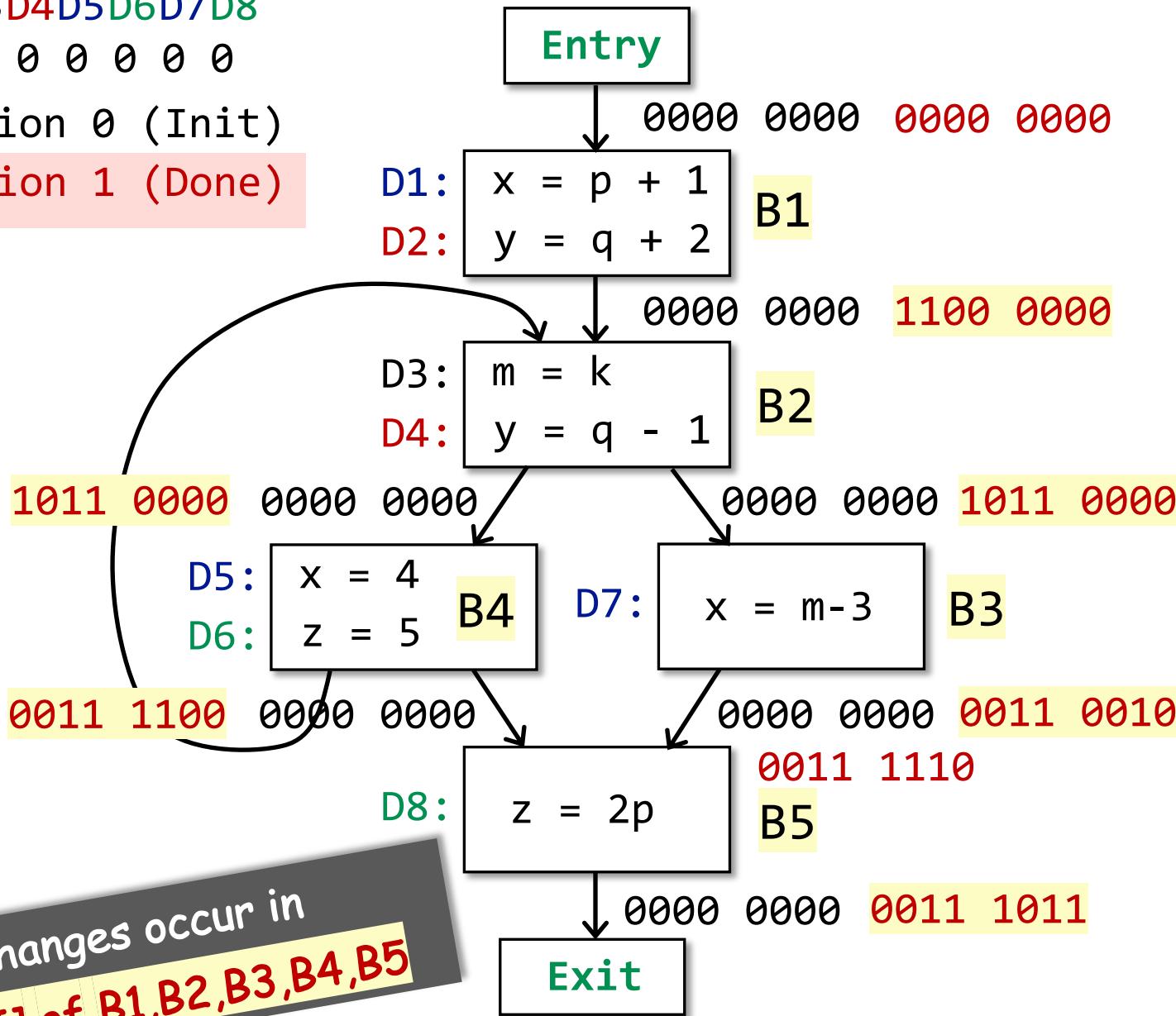


D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)



Changes occur in
OUT[] of B1,B2,B3,B4,B5

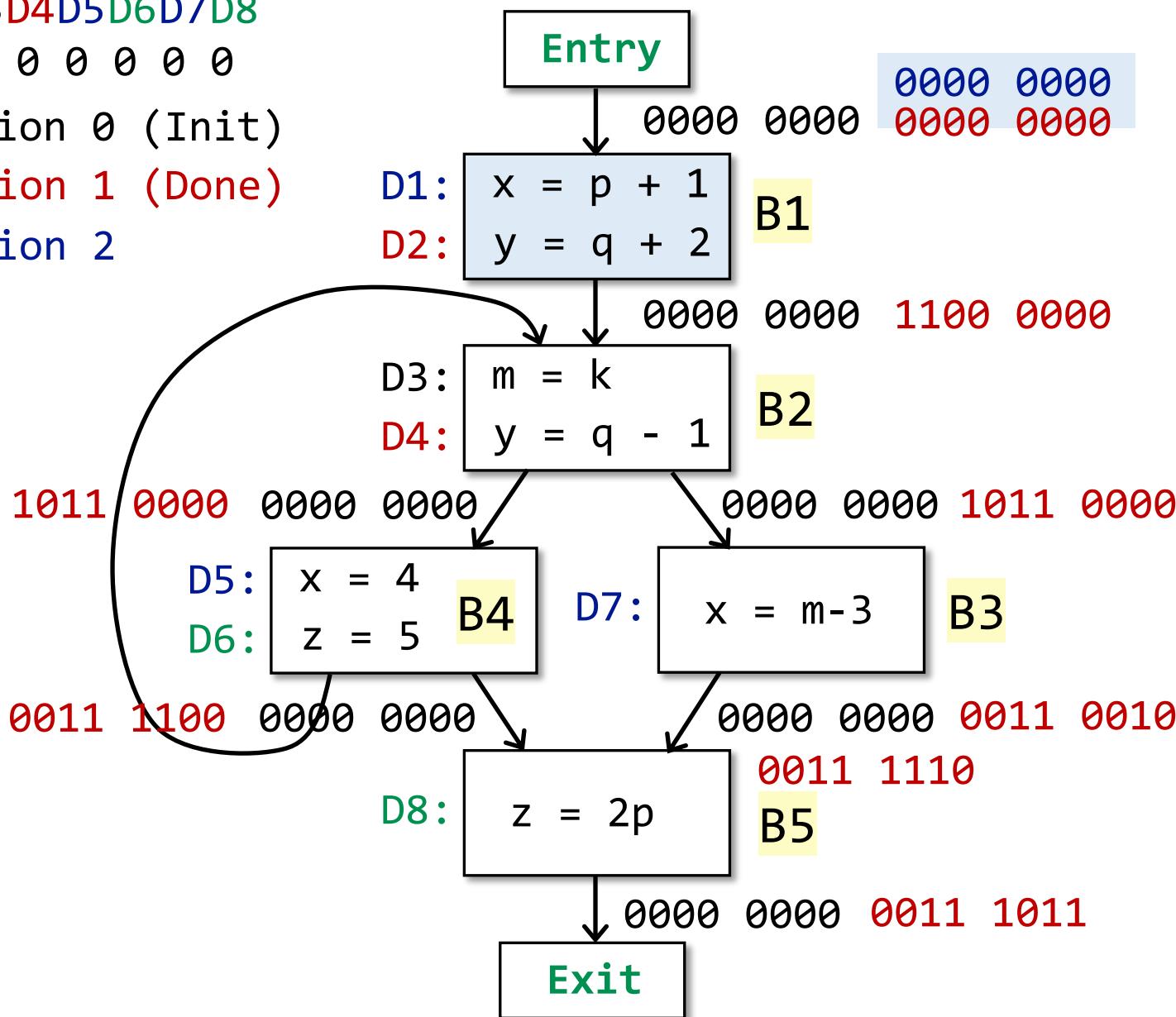
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



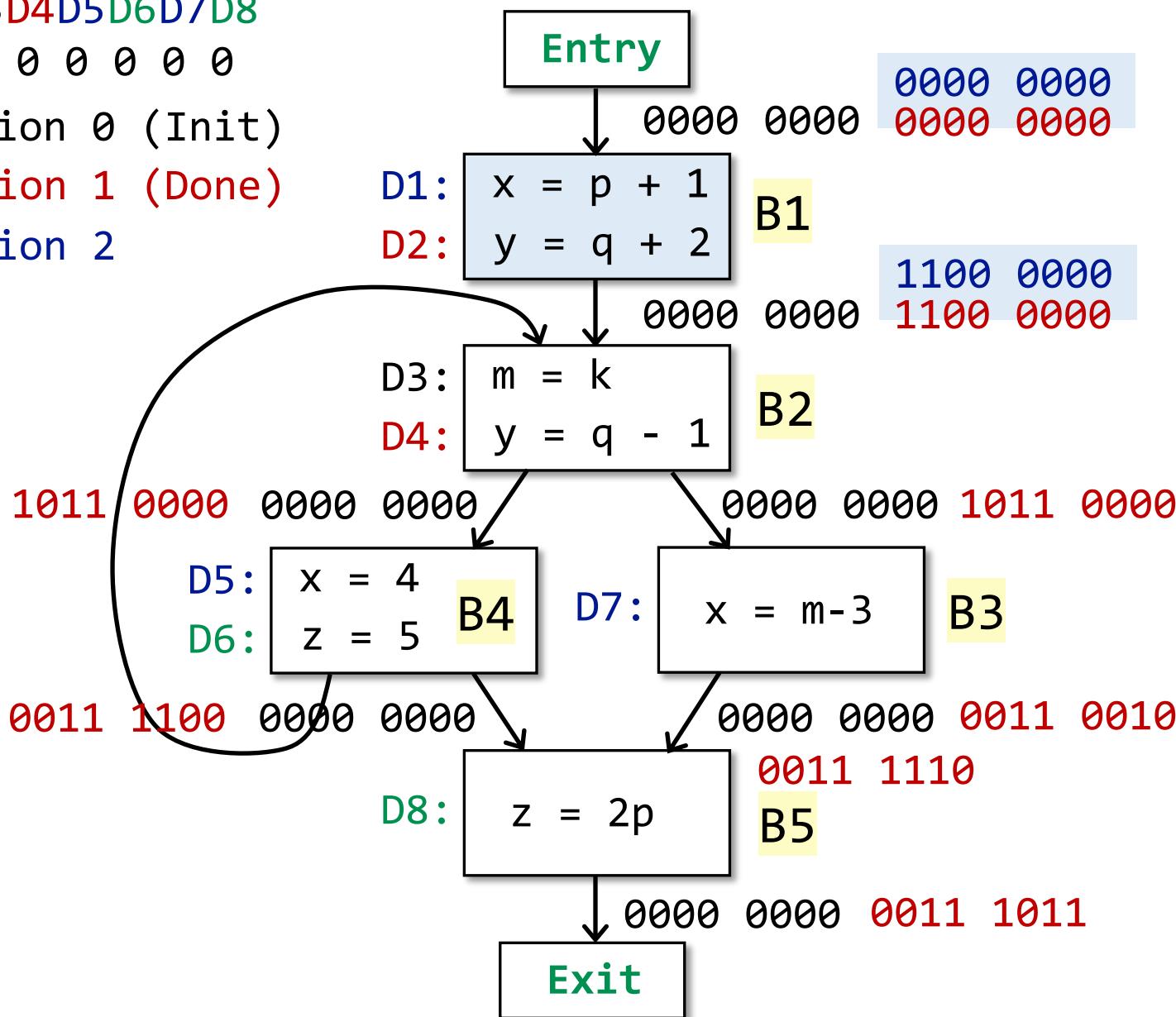
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



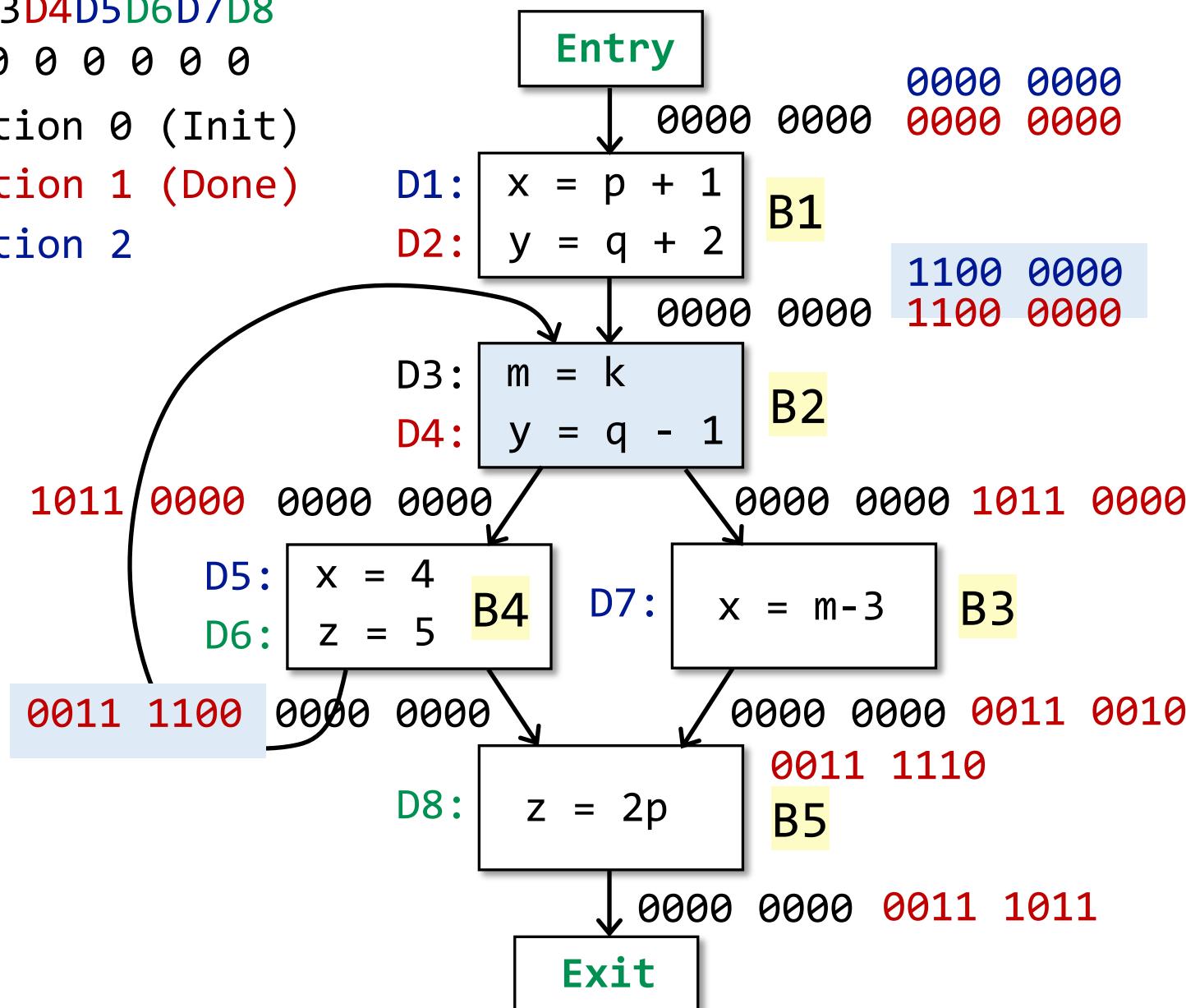
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



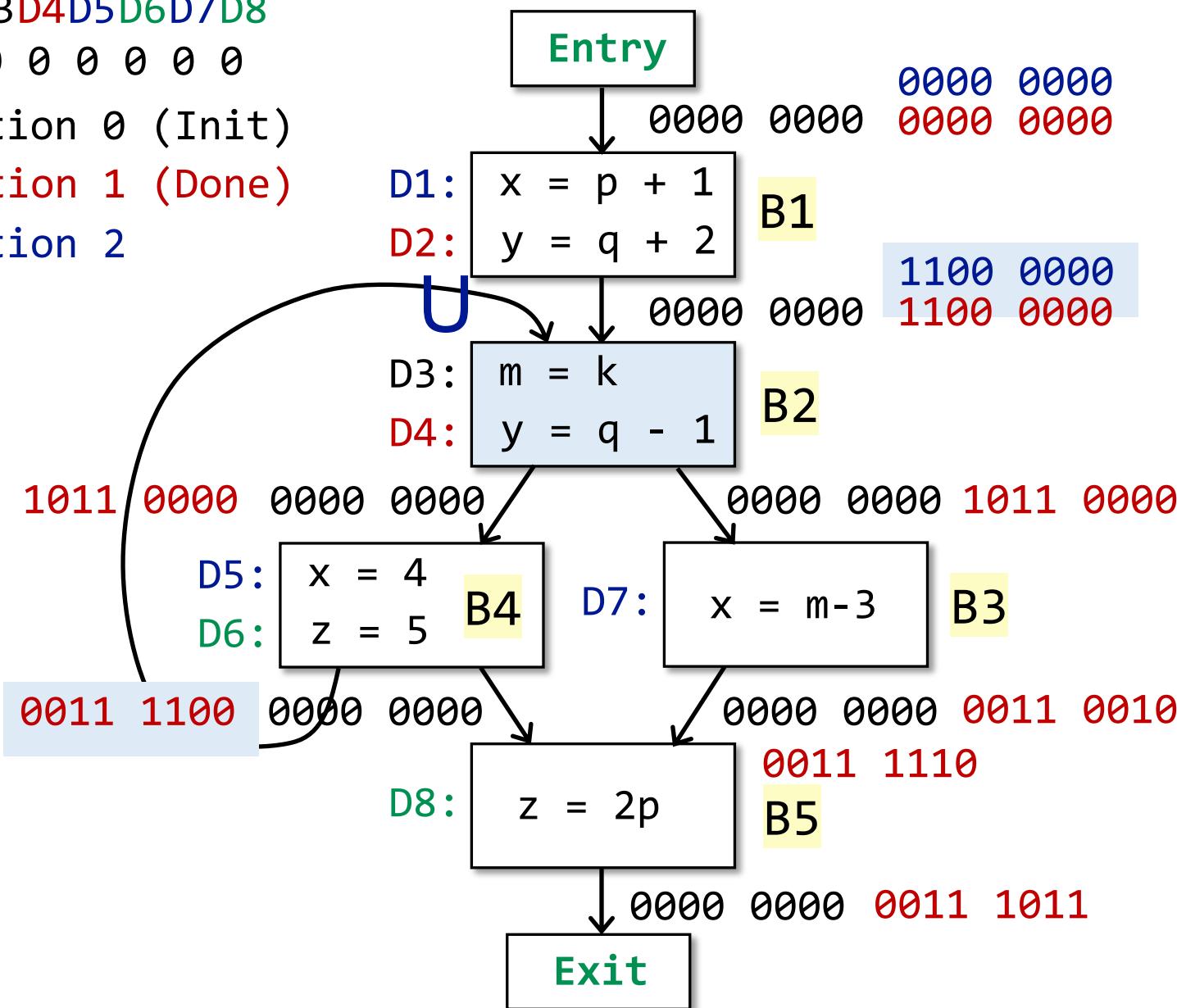
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



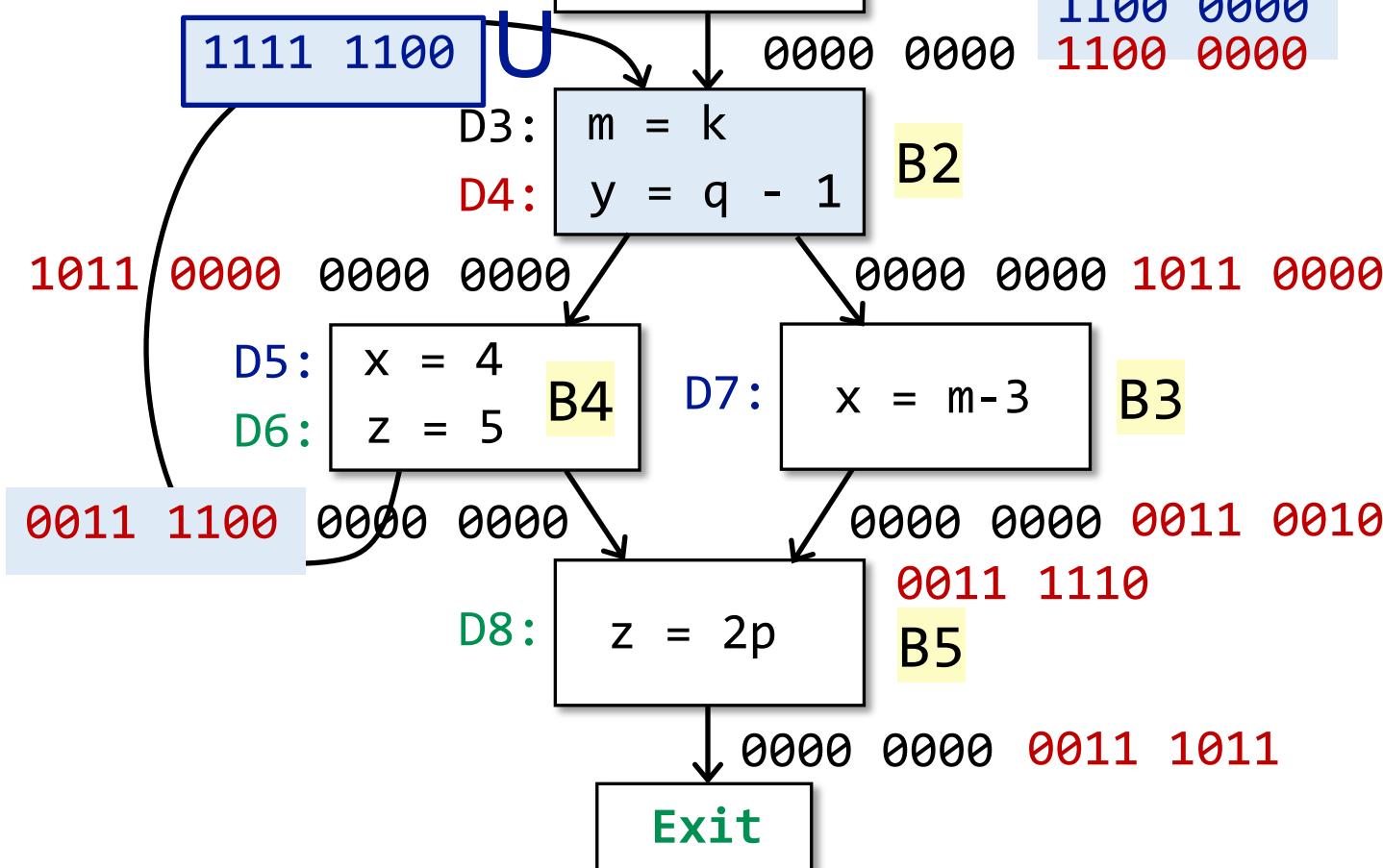
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



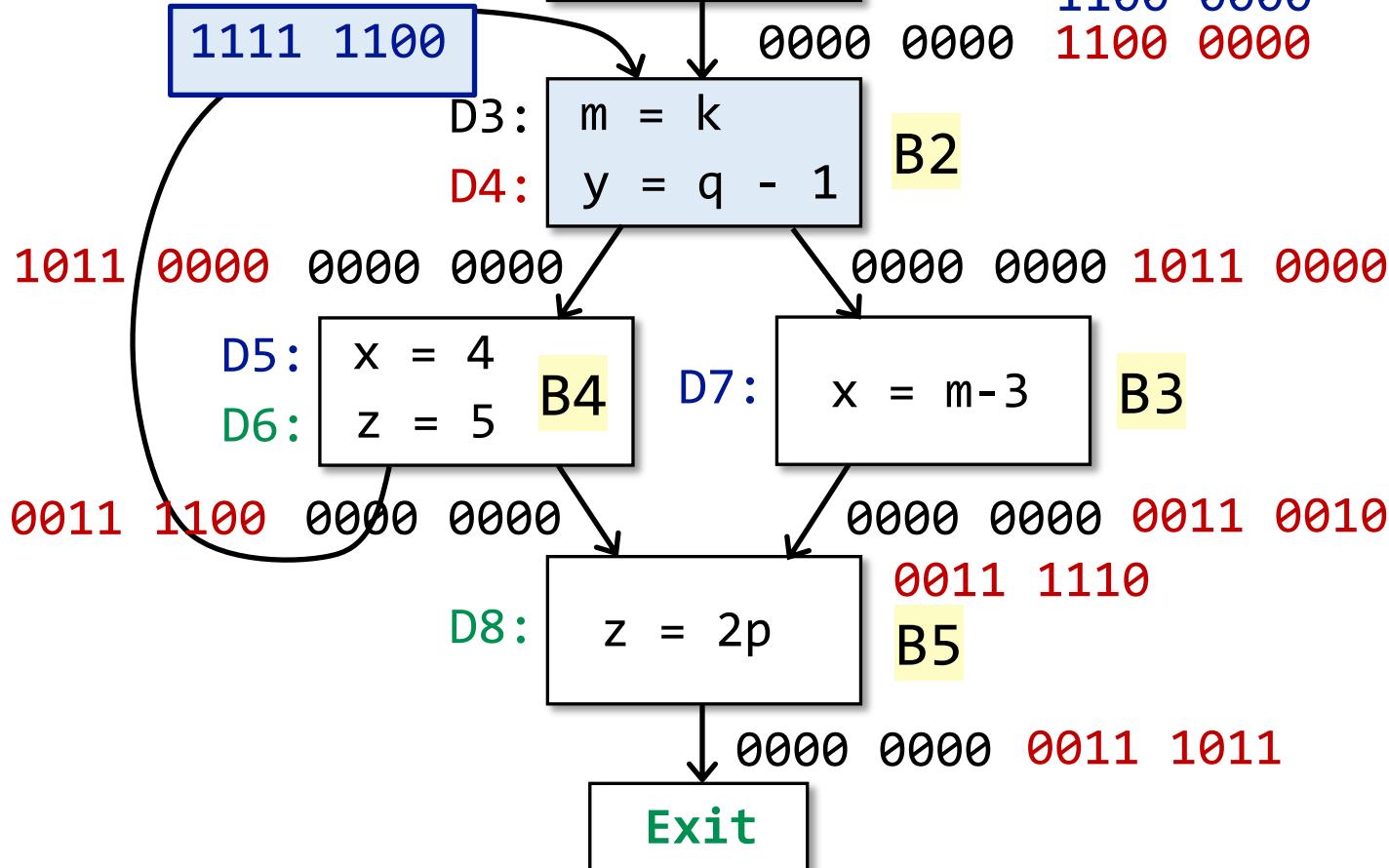
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



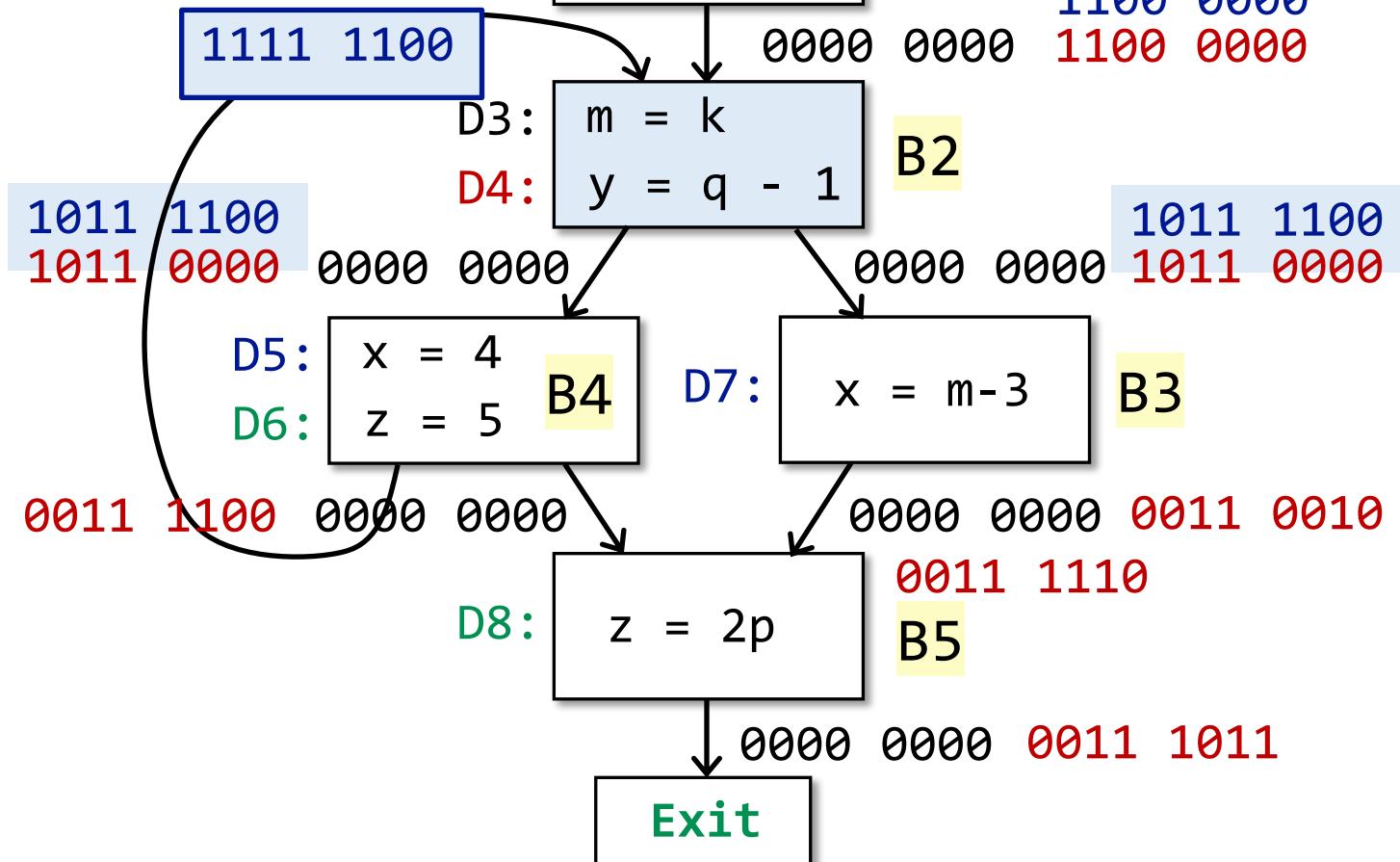
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

D5: $x = 4$
D6: $z = 5$ B4

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D7: $x = m - 3$

1011 1100
1011 0000

B3

D8: $z = 2p$

0011 1110
B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

D5: $x = 4$
D6: $z = 5$ B4

0011 1100 0000 0000

D8: $z = 2p$ B5

0011 1110

Exit

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

1100 0000
1100 0000

D3: $m = k$
D4: $y = q - 1$

B2

1011 1100
1011 0000

D7: $x = m - 3$

B3

0011 0110
0011 0010

0000 0000

0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

0011 1100
0000 0000

D8: $z = 2p$

0011 1110

B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000 0000 0000

0000 0000 0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

Exit

1100 0000
1100 0000

1011 1100
1011 0000

0011 0110
0011 0010

0011 1110
B5

0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

1011 1100
1011 0000

0011 0110
0011 0010

D8: $z = 2p$

0011 1110
B5

Exit

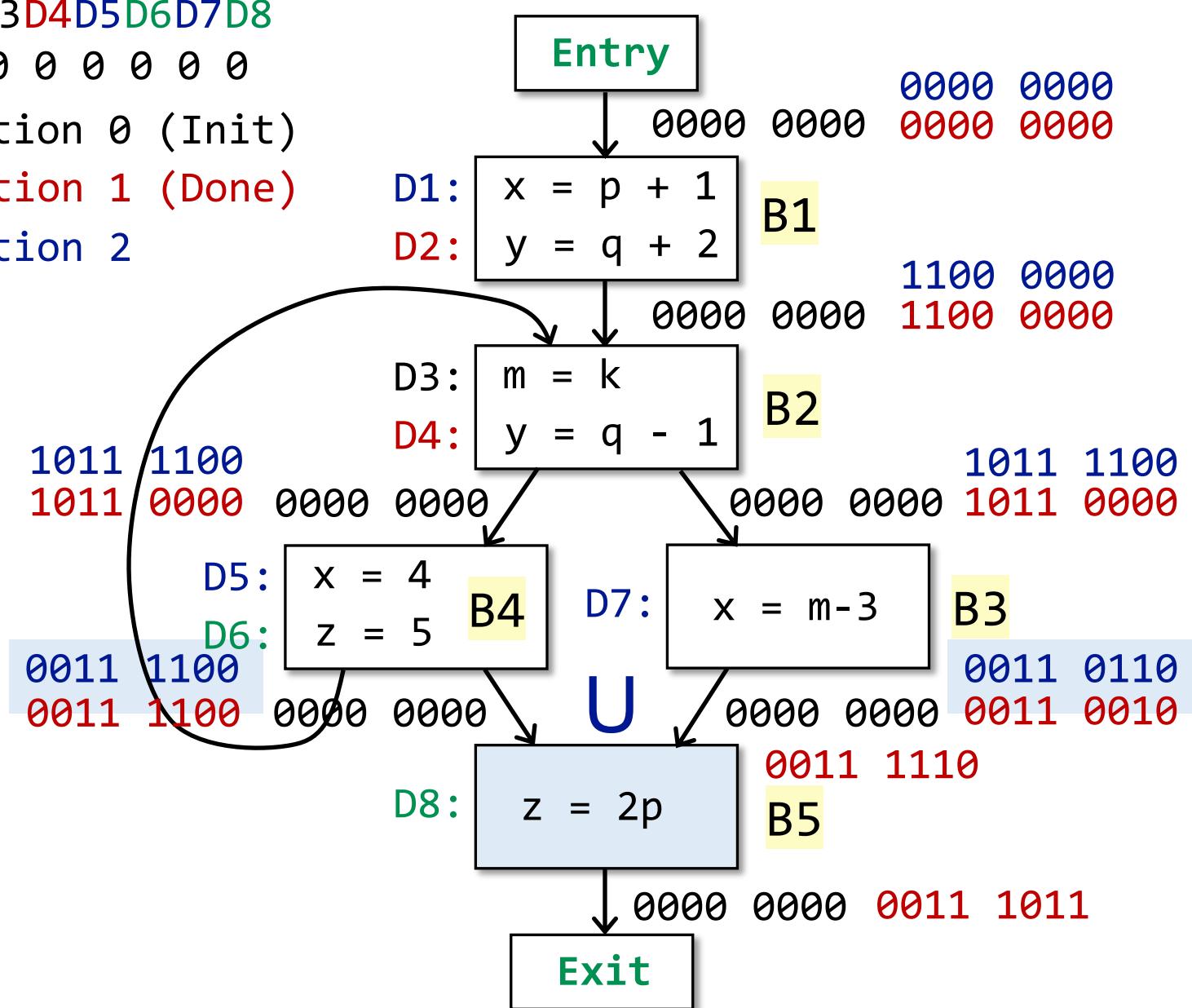
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

1011 1100
1011 0000

B3

D8: $z = 2p$

B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

B5

Exit

1100 0000
1100 0000

1011 1100
1011 0000

0011 0110
0011 0010

0011 1110 0011 1110

0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

B5

Exit

0000 0000
0000 0000

1100 0000
1100 0000

1011 1100
1011 0000

0011 0110
0011 0010

0011 1110 0011 1110

0011 1011
0011 1011

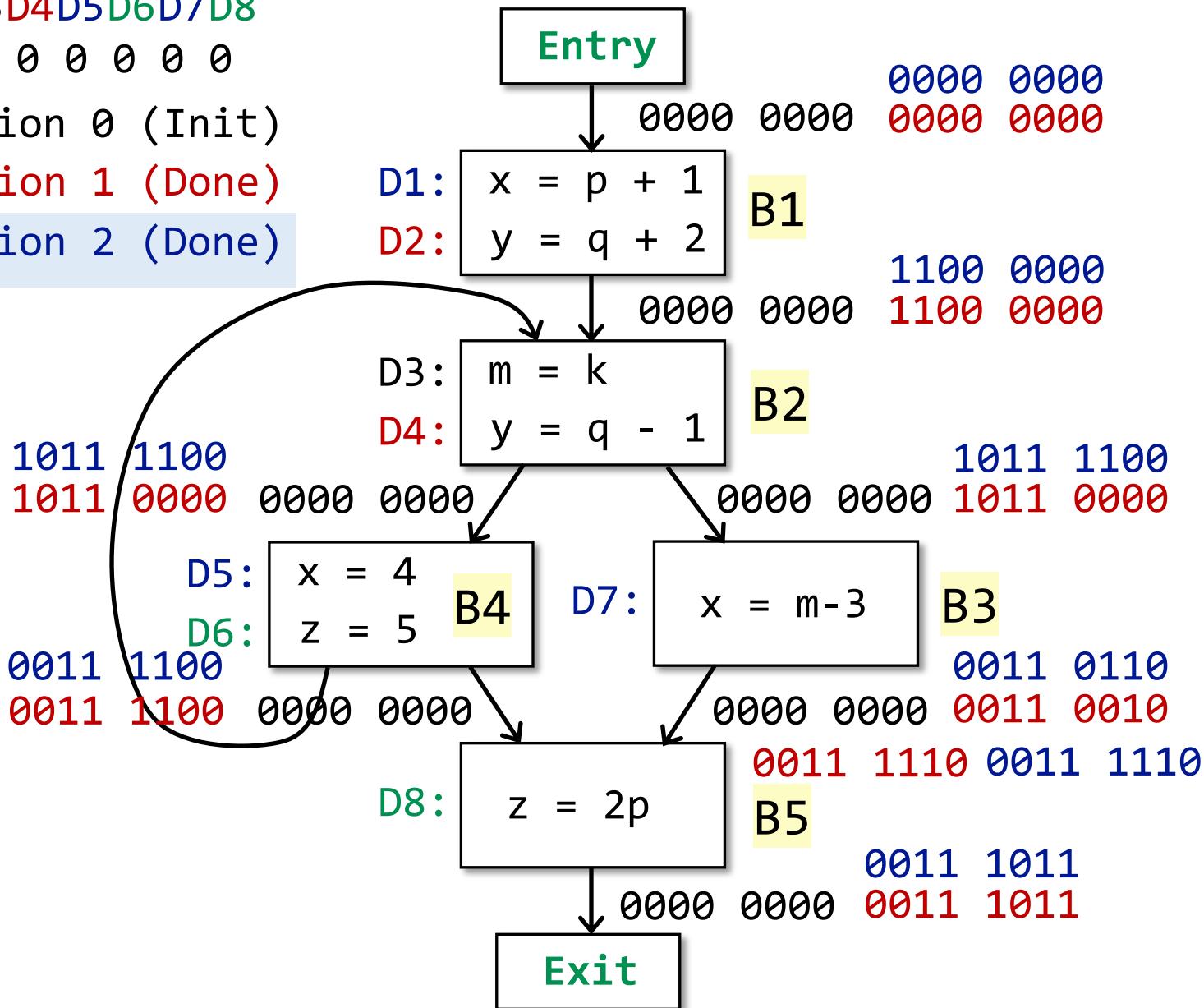
D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

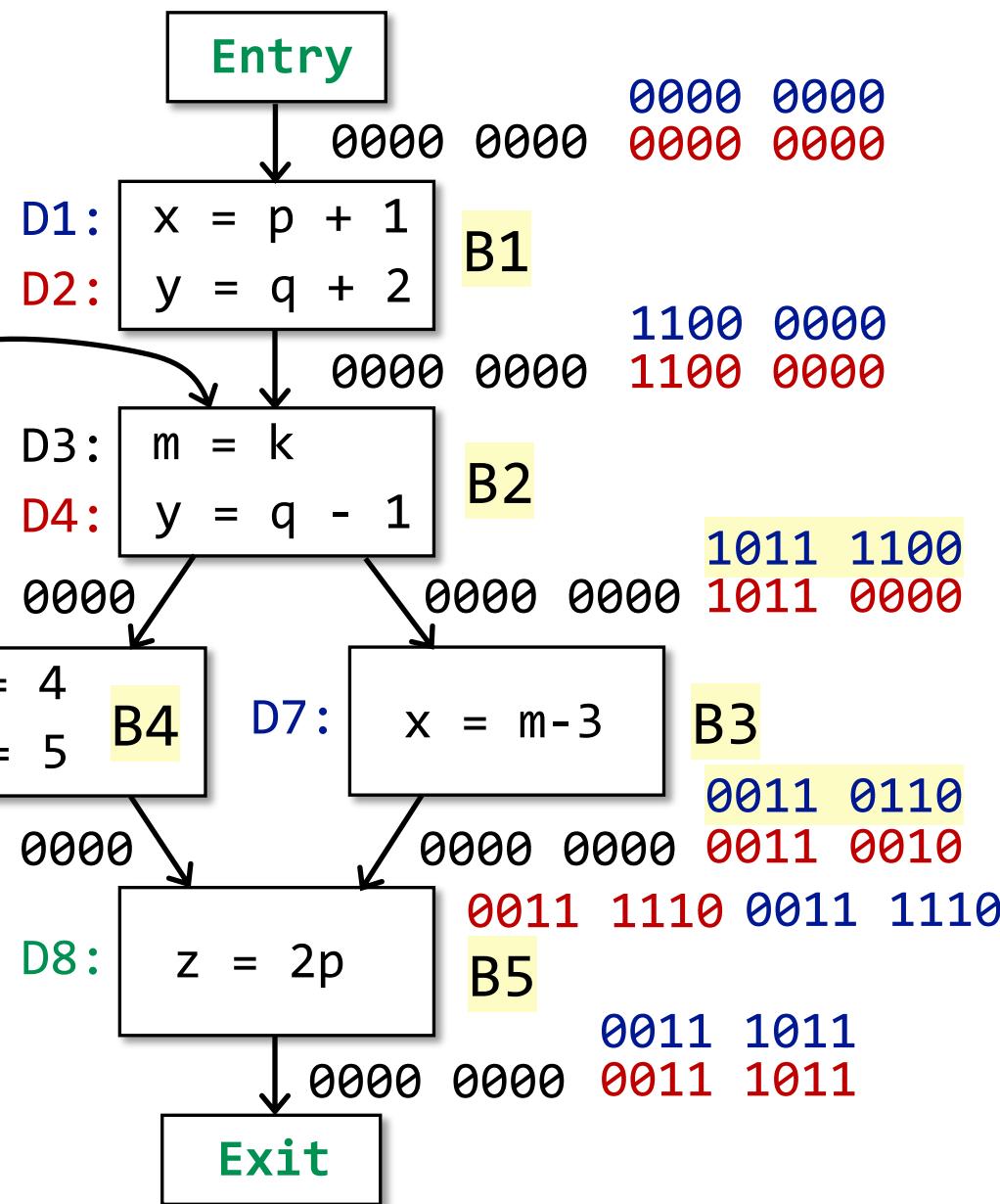
1011 1100
1011 0000

0011 1100
0011 1100

Changes occur in
OUT[] of B2,B3

Entry

Exit



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000	0000
0000	0000
0000	0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

0011 1110 0011 1110
0011 0110
0011 0010

0011 1011
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000	0000
0000	0000
0000	0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

1100	0000
1100	0000
1100	0000

D3: $m = k$
D4: $y = q - 1$

B2

1011	1100
1011	0000

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

0011	0110
0011	0010

D8: $z = 2p$

B5

0011	1110	0011	1110
0011	1011	0011	1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000 0000
0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

1100 0000
1100 0000
1100 0000

D3: $m = k$
D4: $y = q - 1$

B2

1011 1100
1011 0000

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

0011 0110
0011 0010

D8: $z = 2p$

B5

0011 1110 0011 1110
0011 1011
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

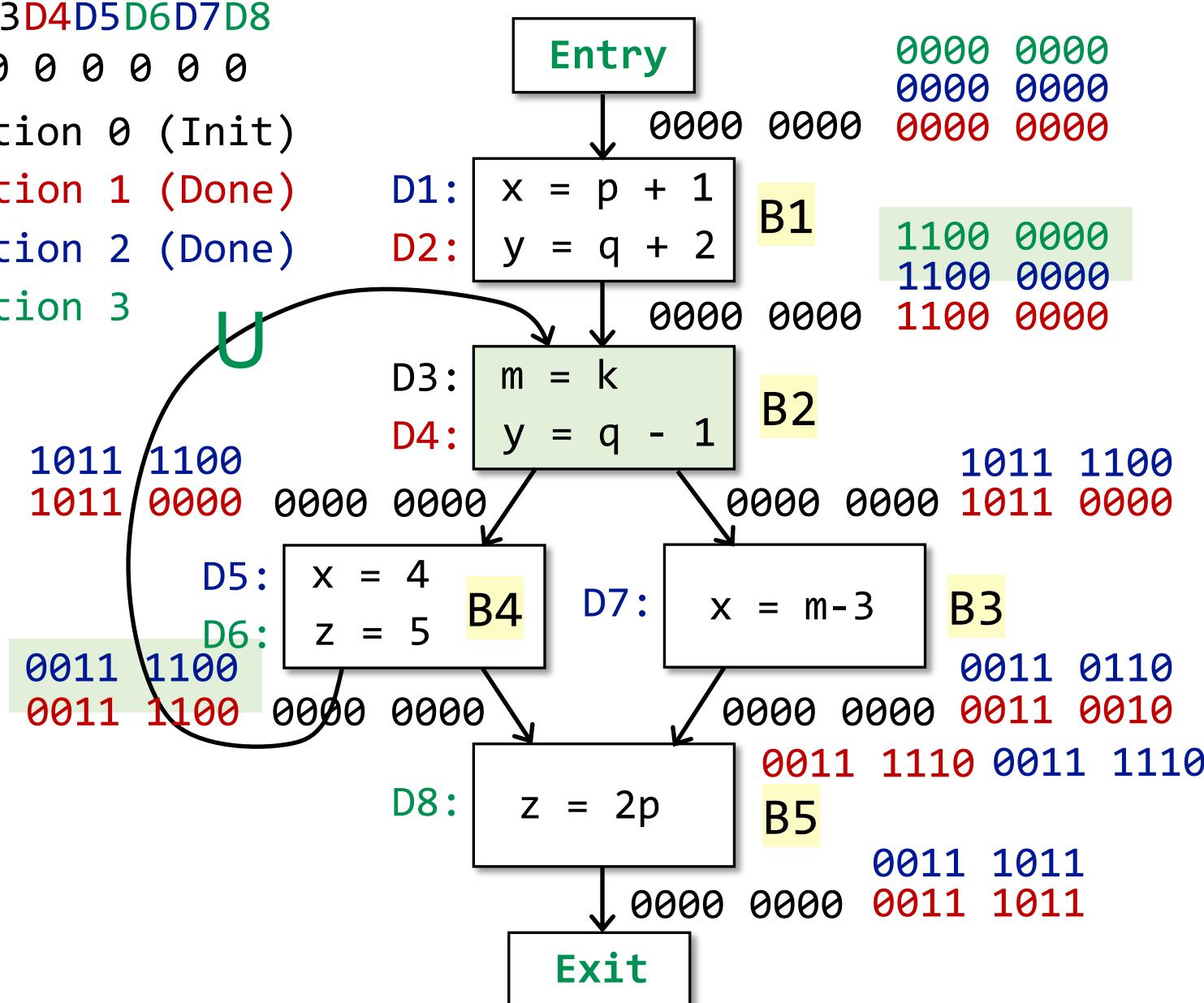
0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 0000

0011 1100
0011 1100

U

Entry

0000 0000

0000 0000
0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

1100 0000
1100 0000
1100 0000

D3: $m = k$
D4: $y = q - 1$

B2

1011 1100
1011 0000

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

0011 0110
0011 0010

D8: $z = 2p$

B5

0011 1110 0011 1110
0011 1011
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 0000

0011 1100
0011 1100

1111 1100

0000 0000

0000 0000
0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D8: $z = 2p$

B3

0011 1110 0011 1110
0011 0010

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

1111 1100

D3: $m = k$

D4: $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5: $x = 4$

D6: $z = 5$

B4

D7: $x = m - 3$

B3

0011 0110

0011 0010

0011 1110 0011 1110

D8: $z = 2p$

0011 1011

0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100

1111 1100

0000 0000

0000	0000
0000	0000
0000	0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

1100	0000
1100	0000
1100	0000

D3: $m = k$
D4: $y = q - 1$

B2

1011	1100
1011	1100
1011	0000

D5: $x = 4$
D6: $z = 5$

B4

0000 0000

D7: $x = m - 3$

0011	0110
0011	0010

D8: $z = 2p$

B5

0000 0000

0011	1110	0011	1110
0011	1011	0011	1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100

1111 1100

0000 0000

0000 0000

0000	0000
0000	0000
0000	0000

1100	0000
1100	0000
1100	0000

1011	1100
1011	1100
1011	0000

0011	0110
0011	0110
0011	0010

0011	1011
0011	1011

Entry

Entry

0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

B5

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100

1111 1100

0000 0000

0000 0000

0000 0000
0000 0000
0000 0000

1100 0000
1100 0000
1100 0000

1011 1100
1011 1100
1011 0000

0011 0110
0011 0110
0011 0010

0011 1110 0011 1110
B5

0011 1011
0011 1011

Exit

Entry

0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

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D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100
0011 1100

0011

0011

0011

Entry

0000 0000

0000 0000
0000 0000
0000 0000

D1: $x = p + 1$
D2: $y = q + 2$

B1

1100 0000
1100 0000
1100 0000

D3: $m = k$
D4: $y = q - 1$

B2

1011 1100
1011 1100
1011 0000

D5: $x = 4$
D6: $z = 5$

B4

D7:

$x = m - 3$

B3

0011 0110

0011 0110

0011 0010

D8: $z = 2p$

B5

0011 1110 0011 1110
0011 1011
0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3: $m = k$

D4: $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5: $x = 4$

D6: $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7: $x = m - 3$

B3

0011 1110

0011 0011

0011 1110

D8: $z = 2p$

B5

0011 1011

0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3: $m = k$

D4: $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5: $x = 4$

D6: $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7: $x = m - 3$

B3

0011 0110

0011 0010

0011 1110

D8: $z = 2p$

B5

0011 1011

0011 1011

U

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3: $m = k$

D4: $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5: $x = 4$

D6: $z = 5$

B4

B3

0011 0110

0011 0110

0011 0010

D7: $x = m - 3$

U

B5

0011 1110 0011 1110

0011 1011

0011 1011

D8: $z = 2p$

Exit

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100
0011 1100

0011 1110

0011 0110

0011 0110

0011 0010

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100
0011 1100

D1: $x = p + 1$
D2: $y = q + 2$

D3: $m = k$
D4: $y = q - 1$

D5: $x = 4$
D6: $z = 5$

D7: $x = m - 3$

D8: $z = 2p$

Entry

Exit

0000 0000
0000 0000
0000 0000

1100 0000
1100 0000
1100 0000

1011 1100
1011 1100
1011 0000

0011 0110
0011 0110
0011 0010

0011 1110 0011 1110
B5

0011 1011
0011 1011

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100
0011 1100

1111 1100

D1: $x = p + 1$
D2: $y = q + 2$

B1

D3: $m = k$
D4: $y = q - 1$

B2

D5: $x = 4$
D6: $z = 5$

B4

D7: $x = m - 3$

B3

D8: $z = 2p$

B5

0000 0000
0000 0000
0000 0000

1100 0000
1100 0000
1100 0000

1011 1100
1011 1100
1011 0000

0011 0110
0011 0110
0011 0010

0011 1110 0011 1110

0011 1011
0011 1011

0011 1011

Exit

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3: $m = k$

D4: $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5: $x = 4$

D6: $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7: $x = m - 3$

B3

0011 0110

0011 0010

D8: $z = 2p$

B5

0011 1011

0011 1011

0011 1011

Exit

1011 1100
1011 1100
1011 0000

0011 1100
0011 1100
0011 1100

0011 1110

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

0000 0000

0000 0000

0000 0000

0000 0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100 0000

1100 0000

1100 0000

D3: $m = k$

D4: $y = q - 1$

B2

1011 1100

1011 1100

1011 0000

D5: $x = 4$

D6: $z = 5$

B4

0011 0110

0011 0110

0011 0010

D7: $x = m - 3$

B3

0011 0110

0011 0010

0011 1110

D8: $z = 2p$

B5

0011 1011

0011 1011

0011 1011

Exit

No changes occur
in any OUT[]

D1 D2 D3 D4 D5 D6 D7 D8

0 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

0000 0000

0000	0000
0000	0000
0000	0000

D1: $x = p + 1$

D2: $y = q + 2$

B1

1100	0000
1100	0000
1100	0000

D3: $m = k$

D4: $y = q - 1$

B2

1011	1100
1011	1100
1011	0000

D5: $x = 4$

D6: $z = 5$

B4

D7:

$x = m - 3$

B3

0011	0110
0011	0110
0011	0010

D8: $z = 2p$

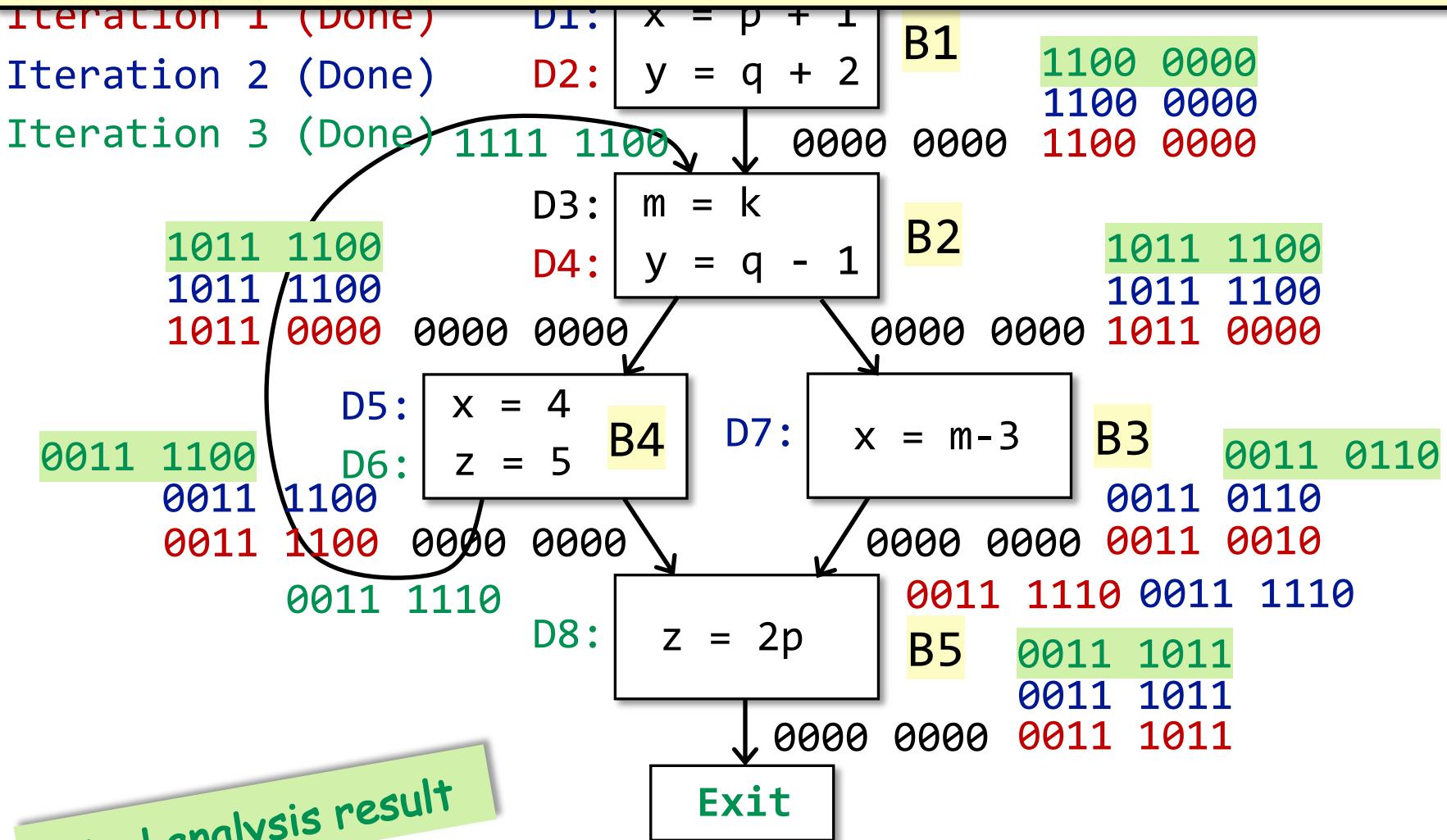
0011 1110 0011 1110

0011	1011
0011	1011
0011	1011

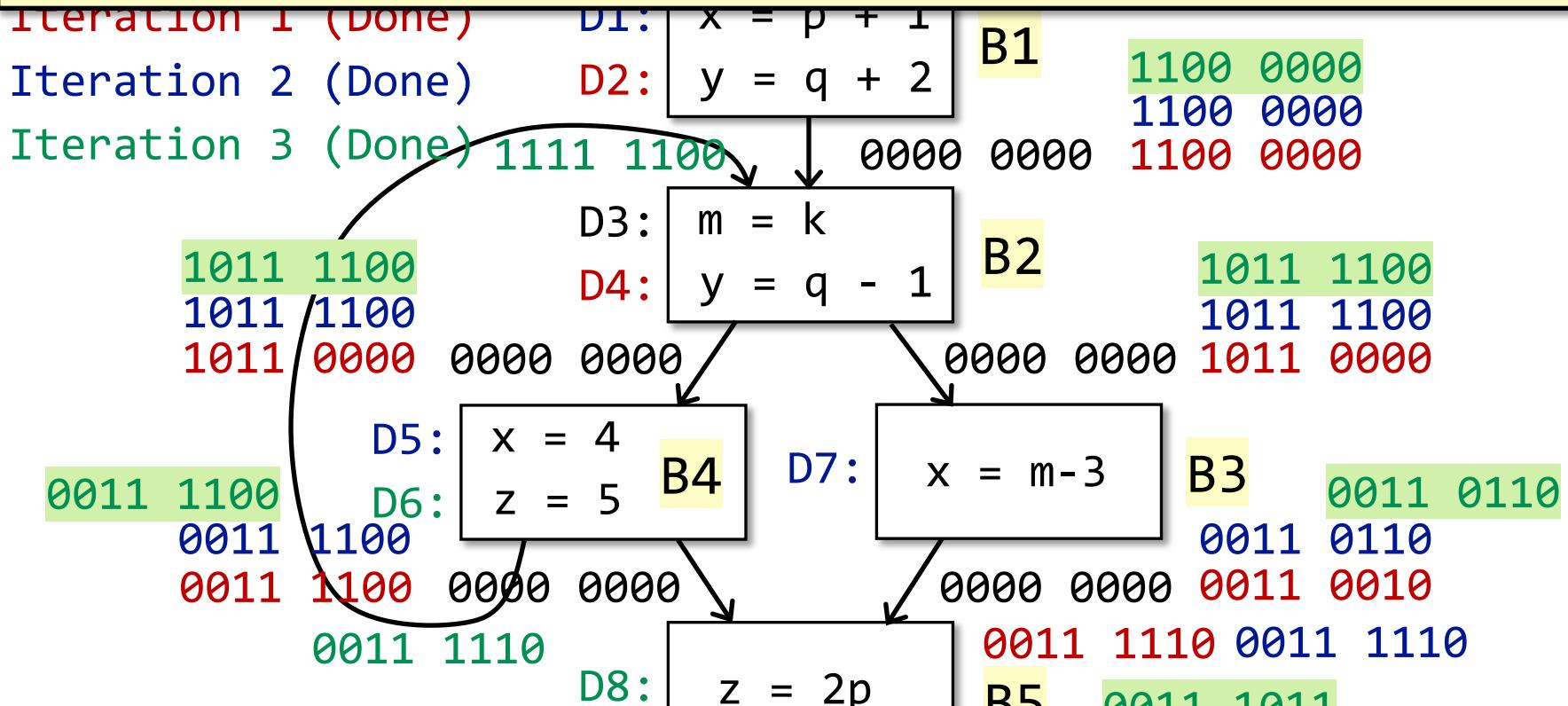
Exit

Final analysis result

In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.



In each data-flow analysis application, we associate with every program point a **data-flow value** that represents an *abstraction* of the set of all possible **program states** that can be observed for that point.



Data-flow analysis is to find a solution to a set of *safe-approximation-directed constraints* on the $\text{IN}[s]$'s and $\text{OUT}[s]$'s, for all statements s .

- constraints based on semantics of statements (*transfer functions*)
- constraints based on the flows of control

Algorithm of Reaching Definitions Analysis

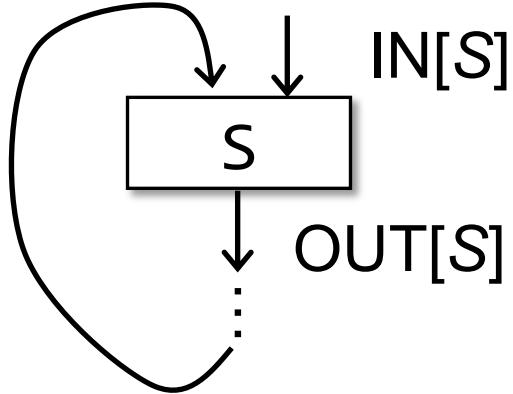
INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

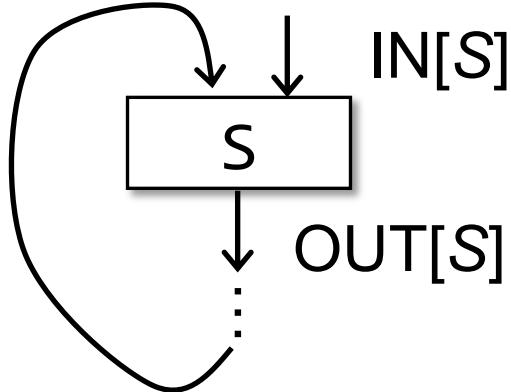
METHOD:

```
OUT[entry] = Ø;
for (each basic block B) {
    OUT[B] = Ø;
    while (changes to any OUT occur)
        for (each basic block B\entry) {
            IN[B] =  $\bigcup_{P \text{ a predecessor of } B}$  OUT[P];
            OUT[B] =  $gen_B \cup (IN[B] - kill_B)$ ;
        }
}
```

Why this iterative algorithm can finally stop?

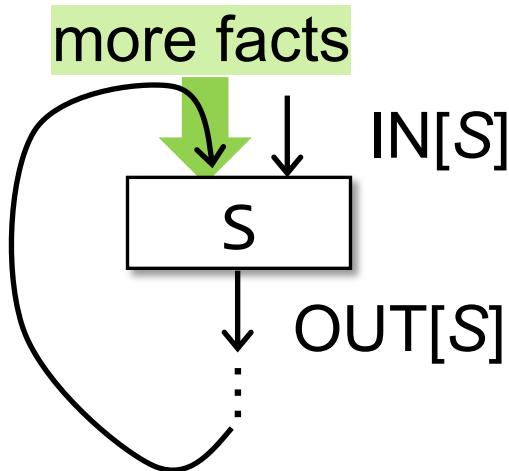


$$\text{OUT}[S] = \text{gen}_S \cup (\text{IN}[S] - \text{kill}_S);$$



$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

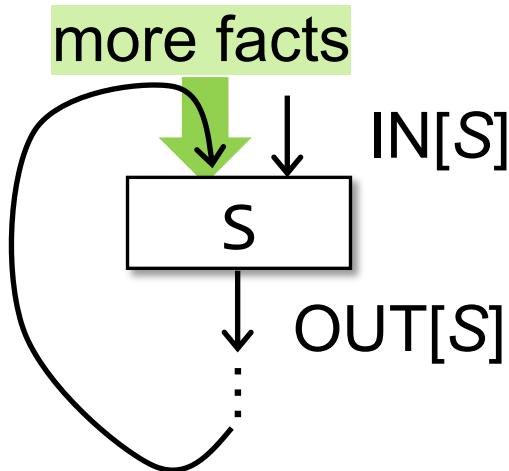
- gen_S and $kill_S$ remain unchanged



more facts

$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

- gen_S and $kill_S$ remain unchanged
- When more facts flow in $IN[S]$, the “more facts” either

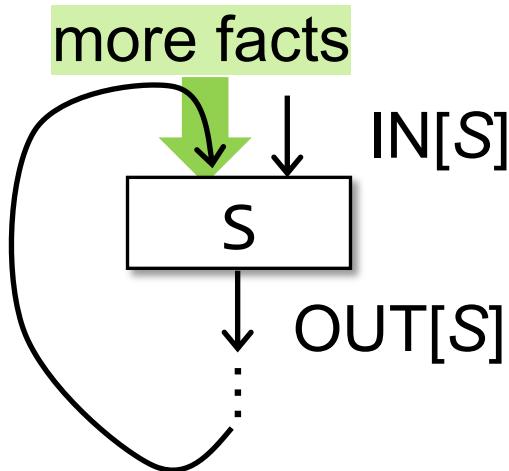


$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

$\underbrace{}_{\text{survivors}}$

more facts

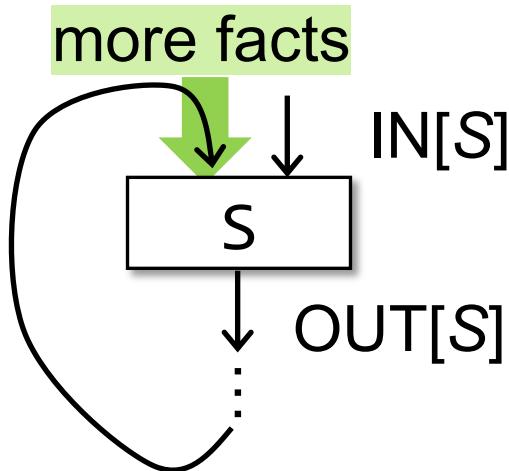
- gen_S and $kill_S$ remain unchanged
- When **more facts** flow in $IN[S]$, the “**more facts**” either
 - is killed, or
 - flows to $OUT[S]$ (**survivor_S**)



$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

survivors_S

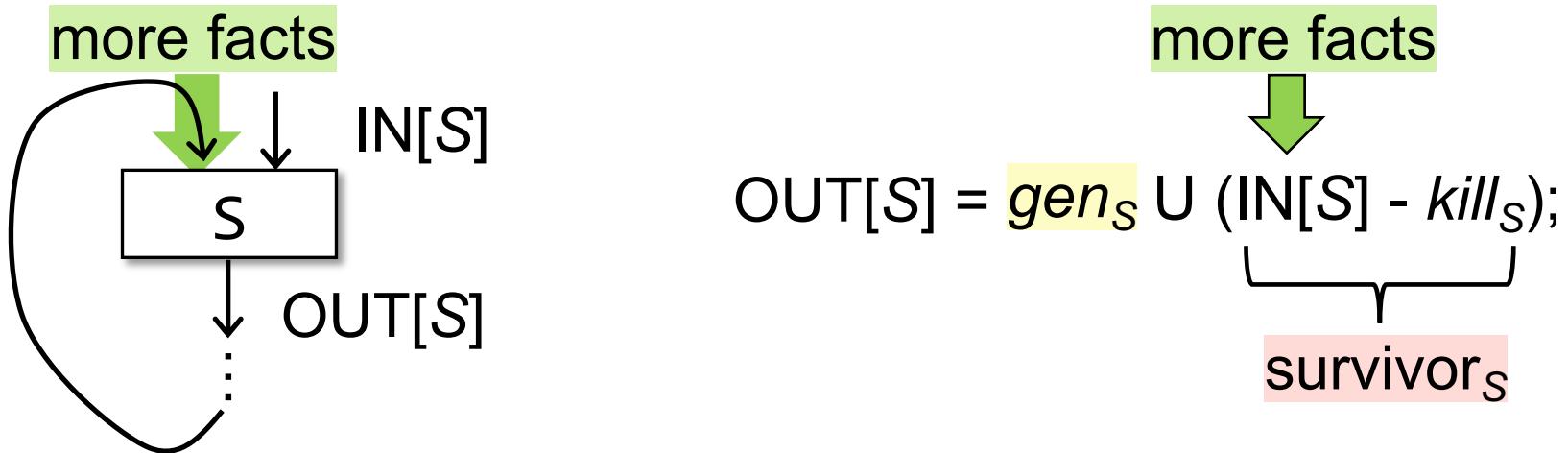
- gen_S and $kill_S$ remain unchanged
- When "more facts" flow in $IN[S]$, the "more facts" either
 - is killed, or
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- When a fact is added to $OUT[S]$, through either gen_S , or $survivors_S$, it stays there forever



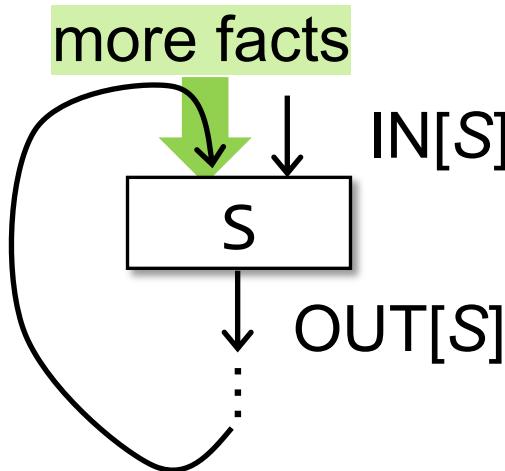
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survivors

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- Thus $OUT[S]$ never shrinks (e.g., $0 \rightarrow 1$, or $1 \rightarrow 1$)



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- Thus $OUT[S]$ never shrinks (e.g., $0 \rightarrow 1$, or $1 \rightarrow 1$)
- As the set of facts is finite (e.g., all definitions in the program),



$$OUT[S] = gen_S \cup (IN[S] - kill_S);$$

survivors

- gen_S and $kill_S$ remain unchanged
- When more facts flow in $IN[S]$, the “more facts” either
 - is killed, or
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- When a fact is added to $OUT[S]$, through either gen_S , or $survivor_S$, it stays there forever
- Thus $OUT[S]$ never shrinks (e.g., $0 \rightarrow 1$, or $1 \rightarrow 1$)
- As the set of facts is finite (e.g., all definitions in the program), there must exist a pass of iteration during which nothing is added to any OUT , and then the algorithm terminates

Algorithm of Reaching Definitions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ ) {  
    OUT[ $B$ ] = Ø;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[ $B$ ] =  $\bigcup_{P \text{ a predecessor of } B} OUT[P]$ ;  
            OUT[ $B$ ] =  $gen_B \cup (IN[B] - kill_B)$ ;  
        }  
}
```

*Safe to terminate
by this condition?*

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        }  
}
```

*Safe to terminate
by this condition?*

IN's will not change if
OUT's do not change

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*Safe to terminate
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if IN's do not change

$$OUT[B] = gen_B \cup (IN[B] - kill_B);$$

Reach a **fixed point**
Also related with **monotonicity**
(next lectures)

Data Flow Analysis Applications

(I) Reaching Definitions Analysis

(II) Live Variables Analysis

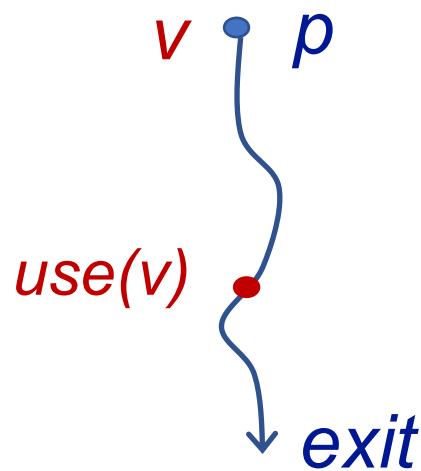
(III) Available Expressions Analysis

Live Variables Analysis

Live variables analysis tells whether the value of **variable v** at **program point p** could be used along some path in CFG starting at p. If so, v is live at p; otherwise, v is dead at p.

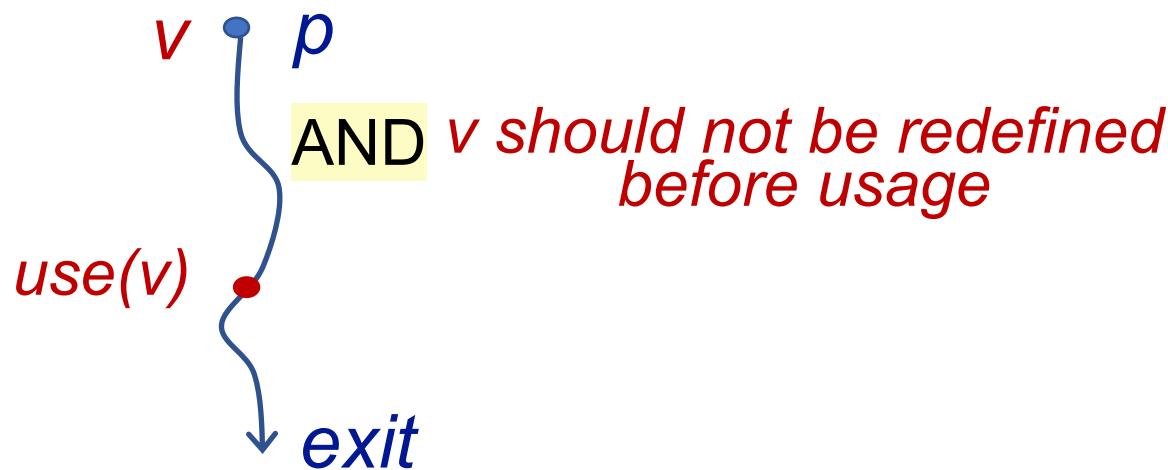
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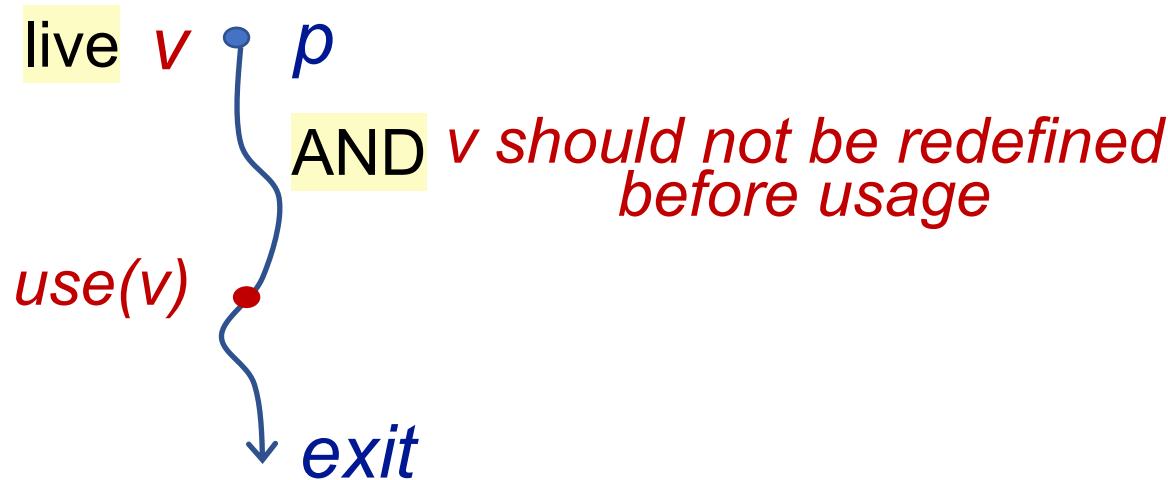
Live Variables Analysis

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Live Variables Analysis

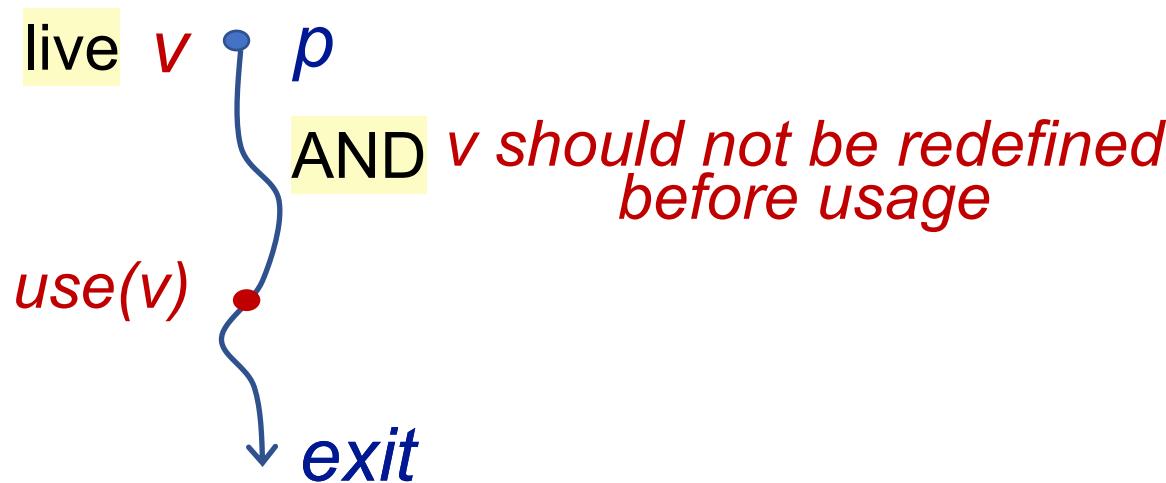
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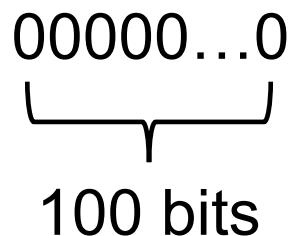
- Information of live variables can be used for register allocations. e.g., at some point all registers are full and we need to use one, then we should favor using a register with a dead value.



Understanding Live Variables Analysis

- Data Flow Values/Facts
 - All the variables in a program
 - Can be represented by bit vectors
- e.g., V1, V2, V3, V4, ..., V100 (100 variables)

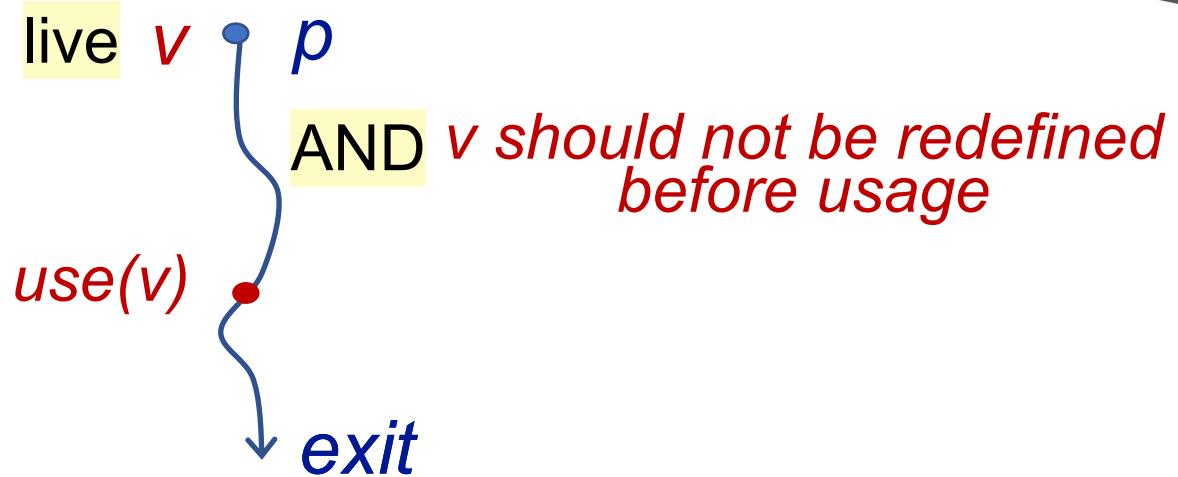
Abstraction



Bit i from the left represents variable Vi

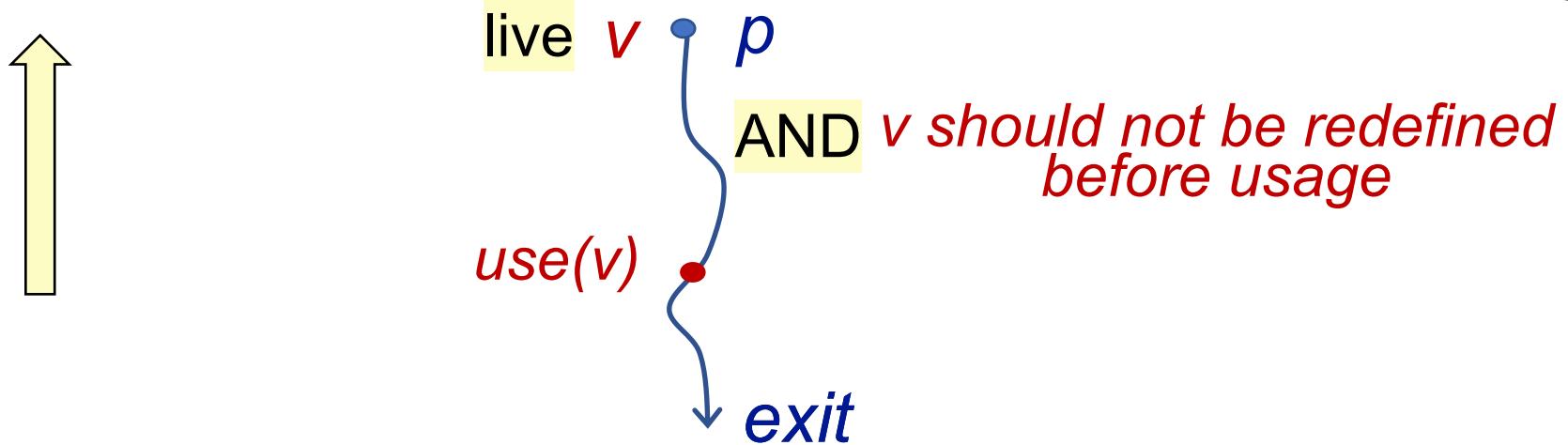
Understanding Live Variables Analysis

Safe-approximation



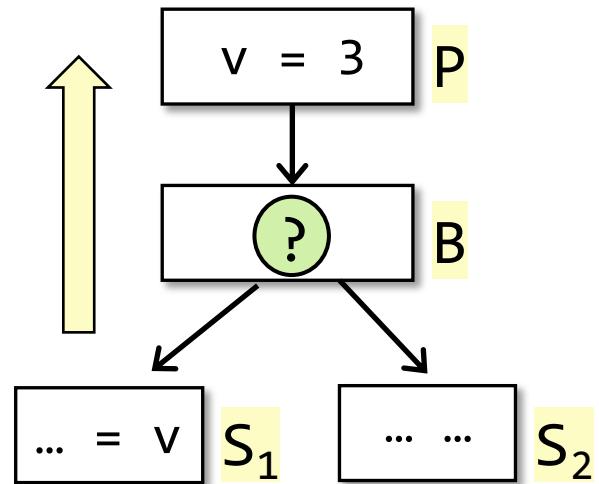
Understanding Live Variables Analysis

Safe-approximation

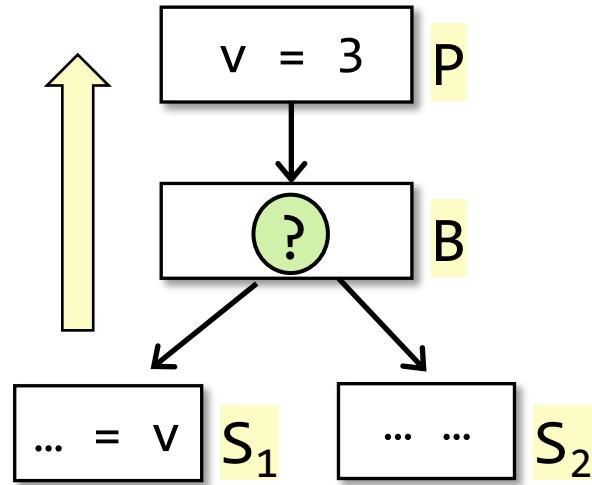


Understanding Live Variables Analysis

Safe-approximation

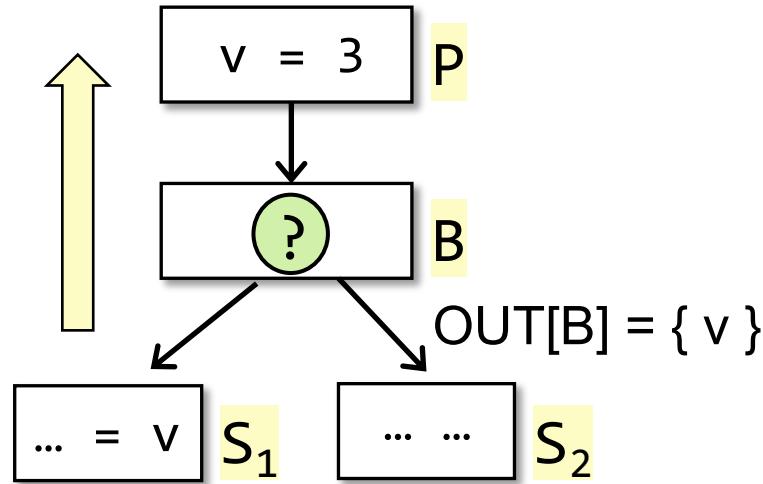


Understanding Live Variables Analysis



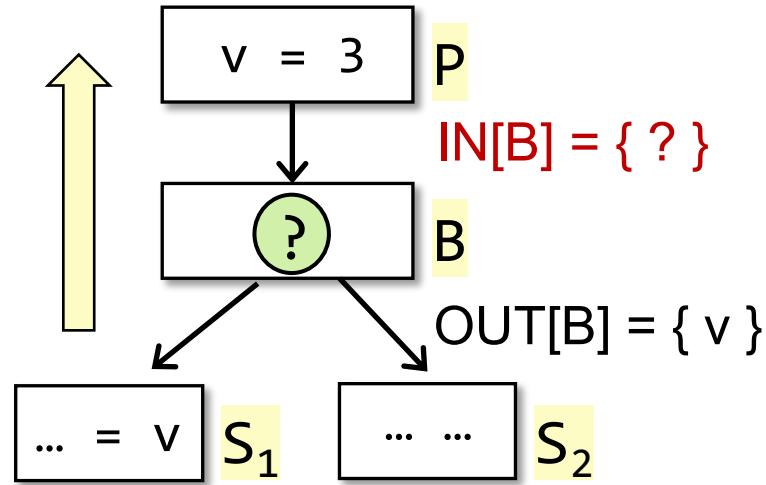
$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[S]$$

Understanding Live Variables Analysis



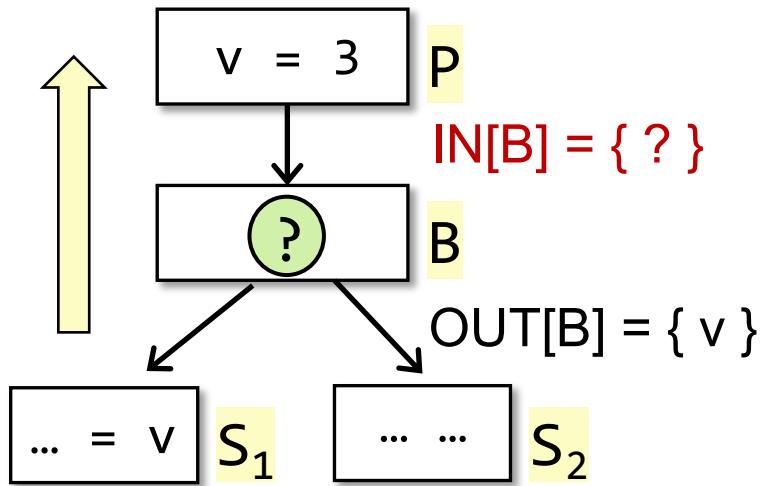
$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis

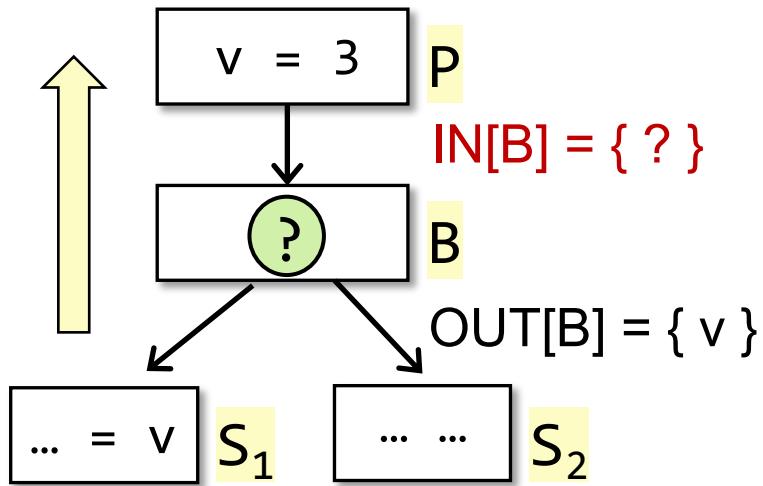


Tip: determine **whether** the variable **v** in some register **R** is **live**, or should we delete the value 3 of **v** in **R**, at the point of **IN[B]**?

Yes: **IN[B] = { v }** No: **IN[B] = { }**

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



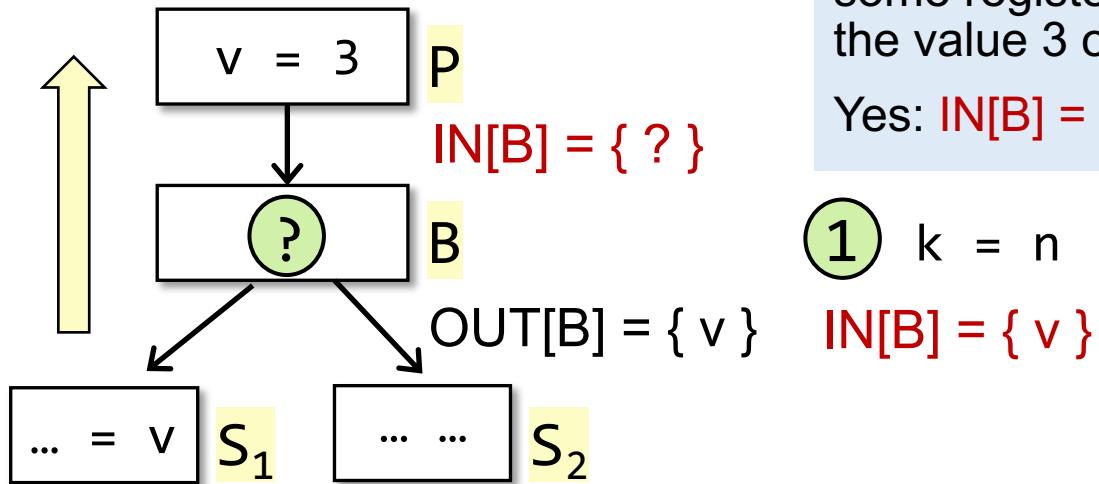
Tip: determine **whether** the variable v in some register R is **live**, or should we delete the value 3 of v in R , at the point of $IN[B]$?

Yes: $IN[B] = \{ v \}$ No: $IN[B] = \{ \}$

1 $k = n$

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



Tip: determine **whether** the variable v in some register R is **live**, or should we delete the value 3 of v in R , at the point of $\text{IN}[B]$?

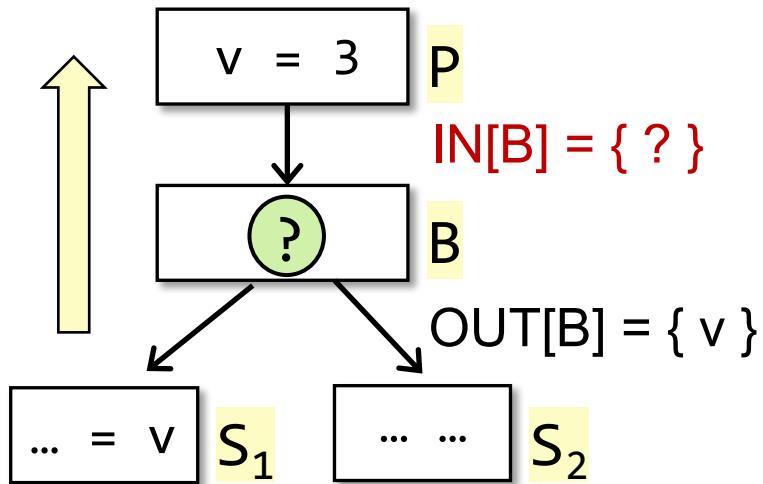
Yes: $\text{IN}[B] = \{ v \}$ No: $\text{IN}[B] = \{ \}$

$$1 \quad k = n$$

$$\text{IN}[B] = \{ v \}$$

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Understanding Live Variables Analysis



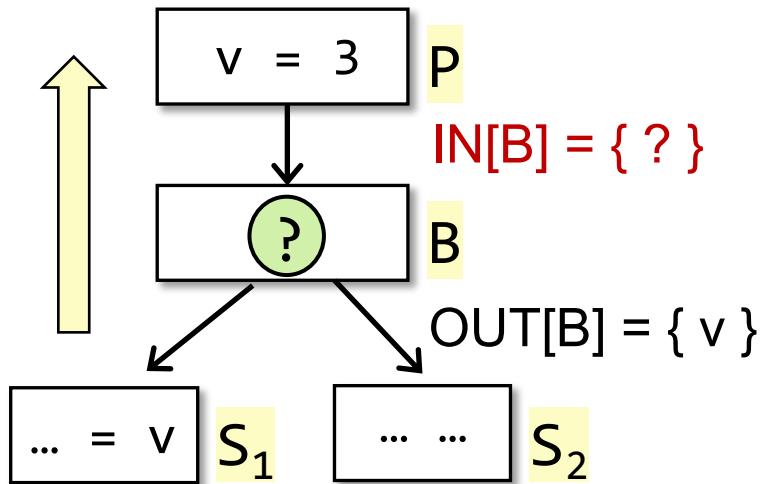
Tip: determine **whether** the variable v in some register R is **live**, or should we delete the value 3 of v in R, at the point of IN[B]?

Yes: IN[B] = { v } No: IN[B] = { }

$$\begin{array}{ll} 1 \quad k = n & 2 \quad k = v \\ \text{IN[B]} = \{ v \} & \end{array}$$

$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[S]$$

Understanding Live Variables Analysis



Tip: determine whether the variable v in some register R is live, or should we delete the value 3 of v in R , at the point of $IN[B]$?

Yes: $IN[B] = \{ v \}$ No: $IN[B] = \{ \}$

$$1 \quad k = n$$

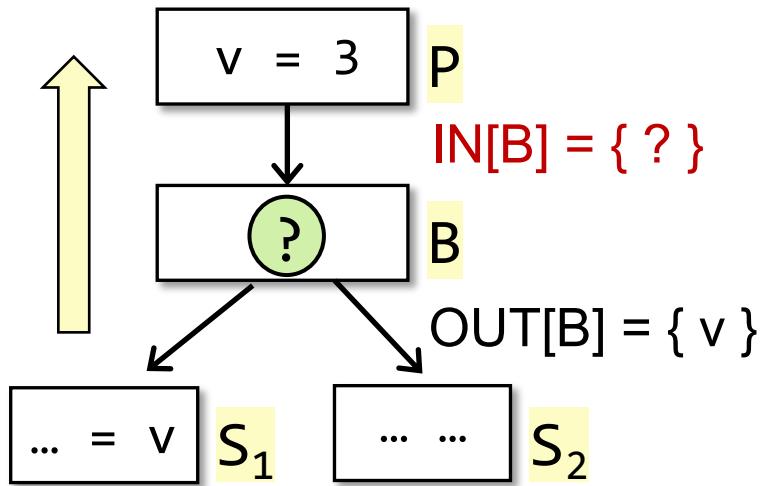
$$IN[B] = \{ v \}$$

$$2 \quad k = v$$

$$IN[B] = \{ v \}$$

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



Tip: determine whether the variable v in some register R is live, or should we delete the value 3 of v in R , at the point of $IN[B]$?

Yes: $IN[B] = \{ v \}$ No: $IN[B] = \{ \}$

1

$k = n$

$IN[B] = \{ v \}$

2

$k = v$

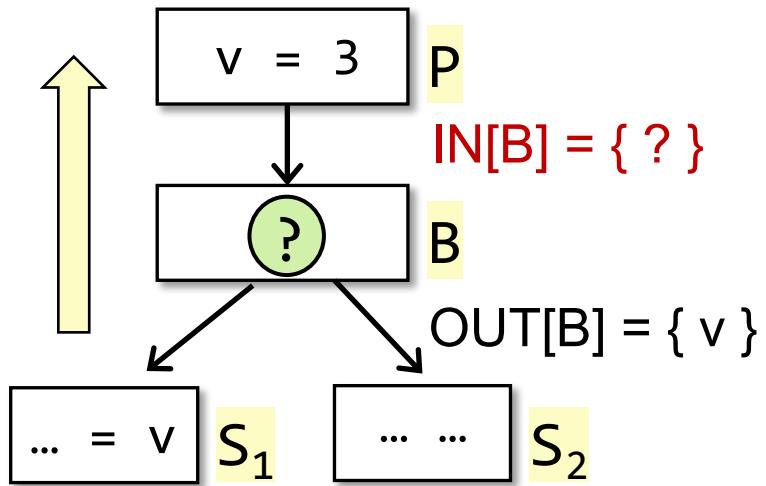
$IN[B] = \{ v \}$

3

$v = 2$

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



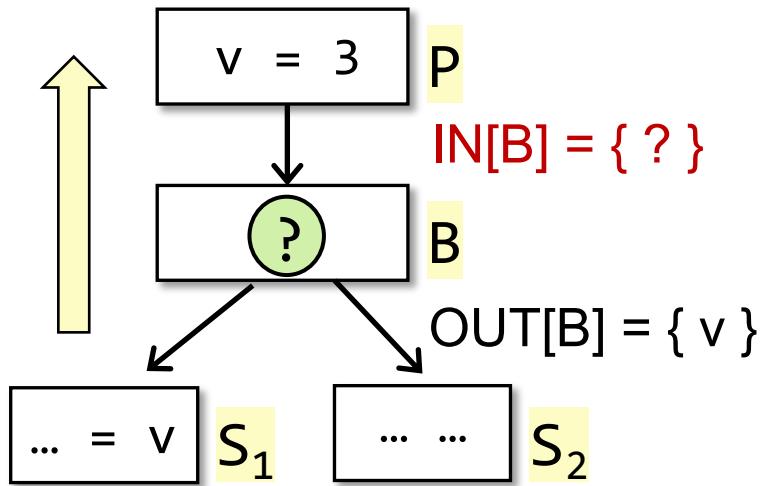
Tip: determine whether the variable v in some register R is live, or should we delete the value 3 of v in R , at the point of $IN[B]$?

Yes: $IN[B] = \{ v \}$ No: $IN[B] = \{ \}$

- (1) $k = n$ $IN[B] = \{ v \}$
- (2) $k = v$ $IN[B] = \{ v \}$
- (3) $v = 2$ $IN[B] = \{ \}$

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



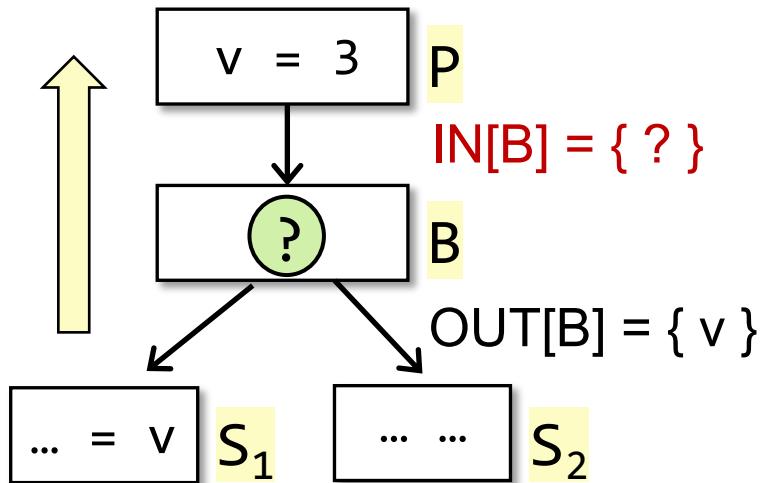
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Yes: $IN[B] = \{ v \}$ No: $IN[B] = \{ \}$

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- (2) $k = v$ $IN[B] = \{ v \}$
- (3) $v = 2$ $IN[B] = \{ \}$
- (4) $v = v - 1$

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



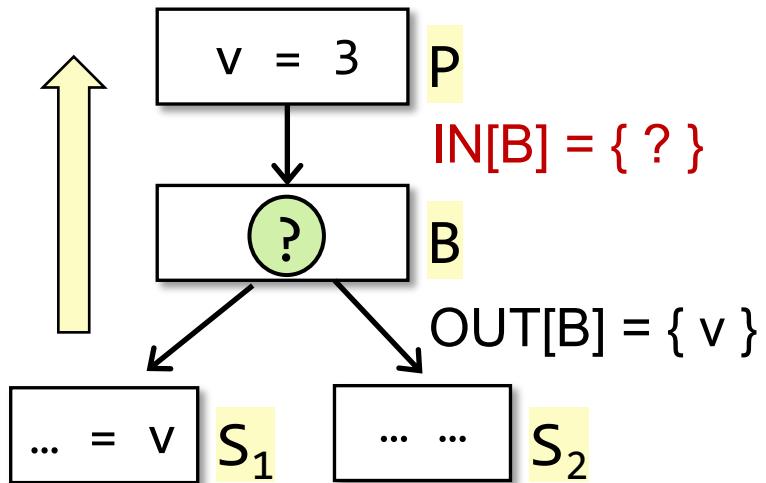
Tip: determine **whether** the variable **v** in some register **R** is live, or should we delete the value 3 of **v** in **R**, at the point of **IN[B]**?

Yes: **IN[B] = { v }** No: **IN[B] = { }**

- | | | | | | |
|----------|----------------------|----------|----------------------|----------|--------------------|
| 1 | $k = n$ | 2 | $k = v$ | 3 | $v = 2$ |
| | IN[B] = { v } | | IN[B] = { v } | | IN[B] = { } |
| 4 | $v = v - 1$ | | | | |
| | | | IN[B] = { v } | | |

$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[S]$$

Understanding Live Variables Analysis



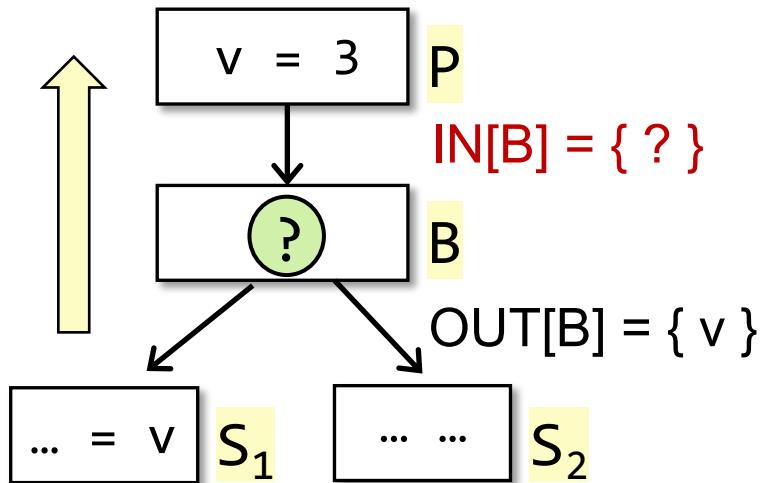
Tip: determine **whether** the variable v in some register R is **live**, or should we delete the value 3 of v in R , at the point of $\text{IN}[B]$?

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- | | | |
|--------------------------|--------------------------|------------------------|
| $1 \quad k = n$ | $2 \quad k = v$ | $3 \quad v = 2$ |
| $\text{IN}[B] = \{ v \}$ | $\text{IN}[B] = \{ v \}$ | $\text{IN}[B] = \{ \}$ |
| $4 \quad v = v-1$ | $5 \quad v = 2$ | |
| $k = v$ | | |
| $\text{IN}[B] = \{ v \}$ | | |

$$\text{OUT}[B] = \bigcup_{S \text{ a successor of } B} \text{IN}[S]$$

Understanding Live Variables Analysis



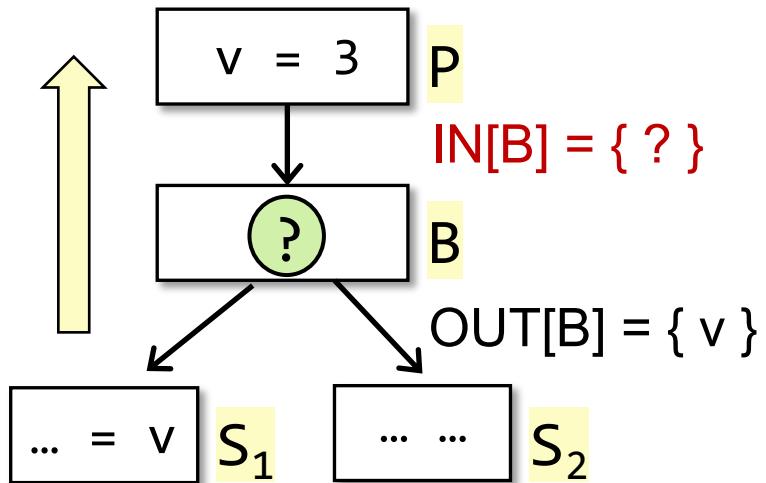
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- | | | | | | |
|----------|-------------------|----------|--------------------|----------|-----------------|
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| | $IN[B] = \{ v \}$ | | $IN[B] = \{ v \}$ | | $IN[B] = \{ \}$ |
| 4 | $v = v-1$ | 5 | $v = 2$
$k = v$ | | |
| | $IN[B] = \{ v \}$ | | $IN[B] = \{ \}$ | | |

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



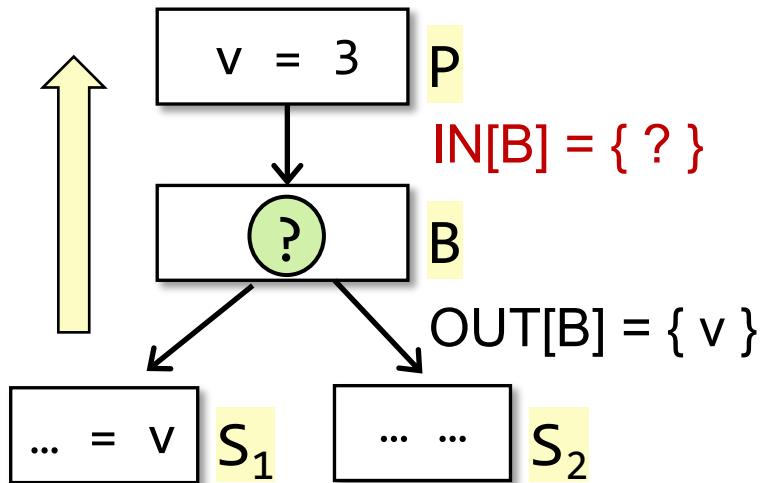
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- | | | | | | |
|---|-------------------|---|--------------------|---|--------------------|
| 1 | $k = n$ | 2 | $k = v$ | 3 | $v = 2$ |
| | $IN[B] = \{ v \}$ | | $IN[B] = \{ v \}$ | | $IN[B] = \{ \}$ |
| 4 | $v = v-1$ | 5 | $v = 2$
$k = v$ | 6 | $k = v$
$v = 2$ |
| | $IN[B] = \{ v \}$ | | $IN[B] = \{ \}$ | | |

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



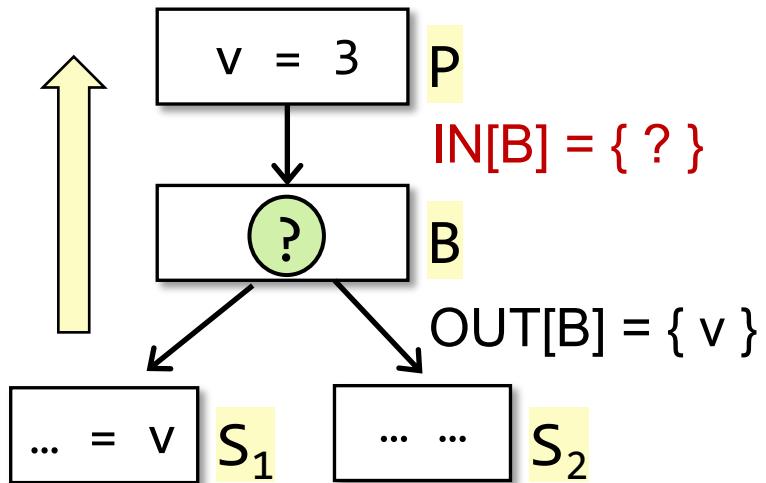
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- | | | | | | |
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| 4 | $v = v-1$ | 5 | $v = 2$
$k = v$ | 6 | $k = v$
$v = 2$ |
| | $IN[B] = \{ v \}$ | | $IN[B] = \{ \}$ | | $IN[B] = \{ v \}$ |

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

Understanding Live Variables Analysis



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- | | | |
|-------------------|----------------------------|----------------------------|
| $1 \quad k = n$ | $2 \quad k = v$ | $3 \quad v = 2$ |
| $IN[B] = \{ v \}$ | $IN[B] = \{ v \}$ | $IN[B] = \{ \}$ |
| $4 \quad v = v-1$ | $5 \quad v = 2$
$k = v$ | $6 \quad k = v$
$v = 2$ |
| $IN[B] = \{ v \}$ | $IN[B] = \{ \}$ | $IN[B] = \{ v \}$ |

$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

$$IN[B] = \text{use}_B \cup (OUT[B] - \text{def}_B)$$

Algorithm of Live Variables Analysis

INPUT: CFG (def_B and use_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
IN[exit] = ∅;  
for (each basic block  $B \setminus exit$ )  
    IN[ $B$ ] = ∅;  
    while (changes to any IN occur)  
        for (each basic block  $B \setminus exit$ ) {  
            OUT[ $B$ ] =  $\bigcup_{S \text{ a successor of } B} IN[S]$ ;  
            IN[ $B$ ] =  $use_B \cup (OUT[B] - def_B)$ ;  
        }  
    }
```

Algorithm of Live Variables Analysis

INPUT: CFG (def_B and use_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
IN[exit] = Ø;  
for (each basic block  $B \setminus exit$ )  
    IN[ $B$ ] = Ø;  
    while (changes to any IN occur)  
        for (each basic block  $B \setminus exit$ ) {  
            OUT[ $B$ ] =  $\bigcup_{S \text{ a successor of } B} IN[S]$ ;  
            IN[ $B$ ] =  $use_B \cup (OUT[B] - def_B)$ ;  
        }
```

Algorithm of Live Variables Analysis

INPUT: CFG (def_B and use_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
IN[exit] = Ø;  
for (each basic block  $B \setminus exit$ )  
    IN[B] = Ø;  
    while (changes to any IN occur)  
        for (each basic block  $B \setminus exit$ ) {  
            OUT[B] =  $\bigcup_{S \text{ a successor of } B} IN[S];$   
            IN[B] =  $use_B \cup (OUT[B] - def_B);$   
        }
```

Algorithm of Live Variables Analysis

INPUT: CFG (def_B and use_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
IN[exit] = ∅;  
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    IN[ $B$ ] = ∅;  
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            OUT[ $B$ ] =  $\bigcup_{S \text{ a successor of } B} IN[S]$ ;  
            IN[ $B$ ] =  $use_B \cup (OUT[B] - def_B)$ ;  
        }
```

Algorithm of Live Variables Analysis

INPUT: CFG (def_B and use_B computed for each basic block B)

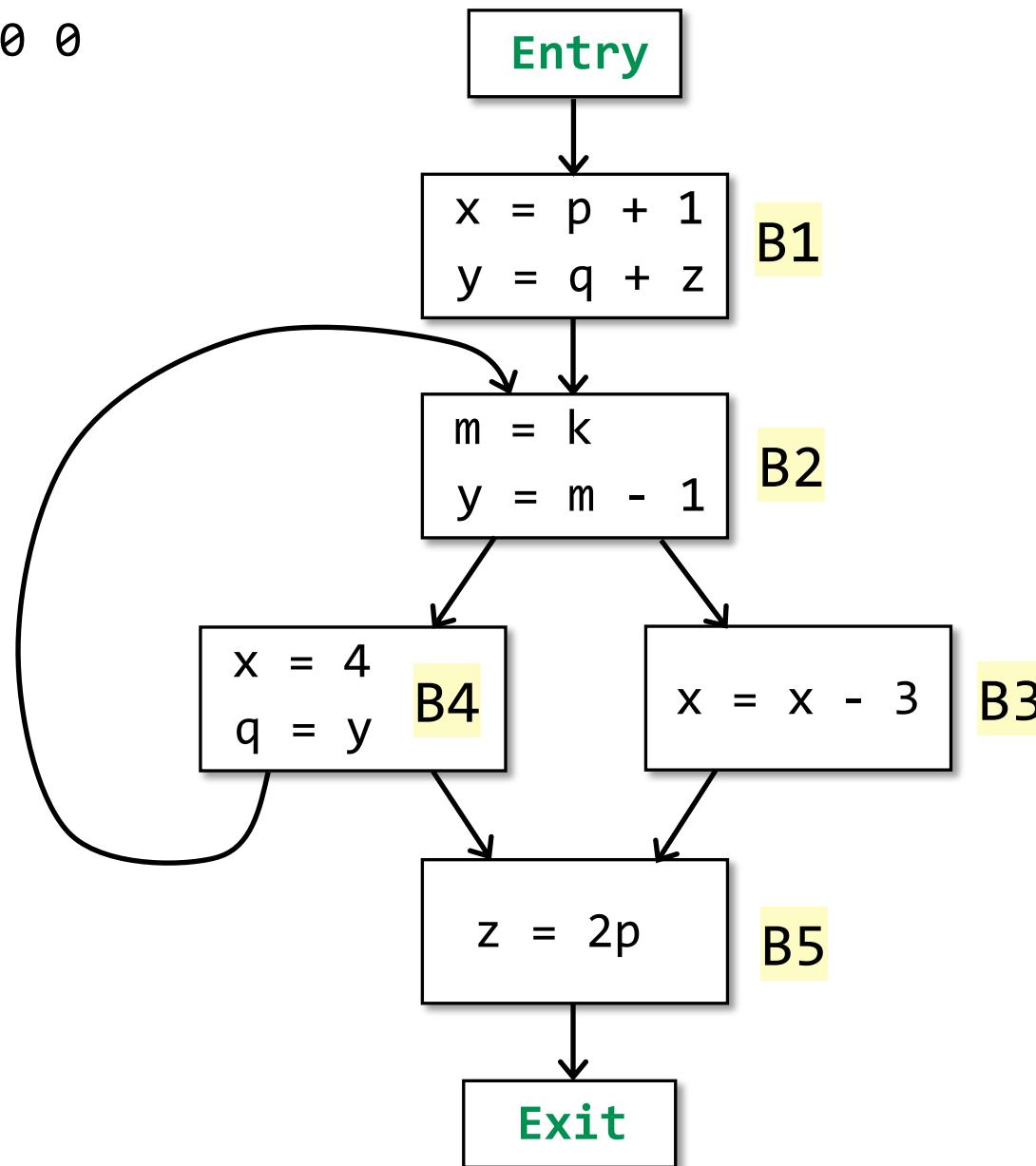
OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
IN[exit] = ∅;  
for (each basic block  $B \setminus exit$ )  
    IN[ $B$ ] = ∅;  
    while (changes to any IN occur)  
        for (each basic block  $B \setminus exit$ ) {  
            OUT[ $B$ ] =  $\bigcup_{S \text{ a successor of } B} IN[S]$ ;  
            IN[ $B$ ] =  $use_B \cup (OUT[B] - def_B)$ ;  
        }
```

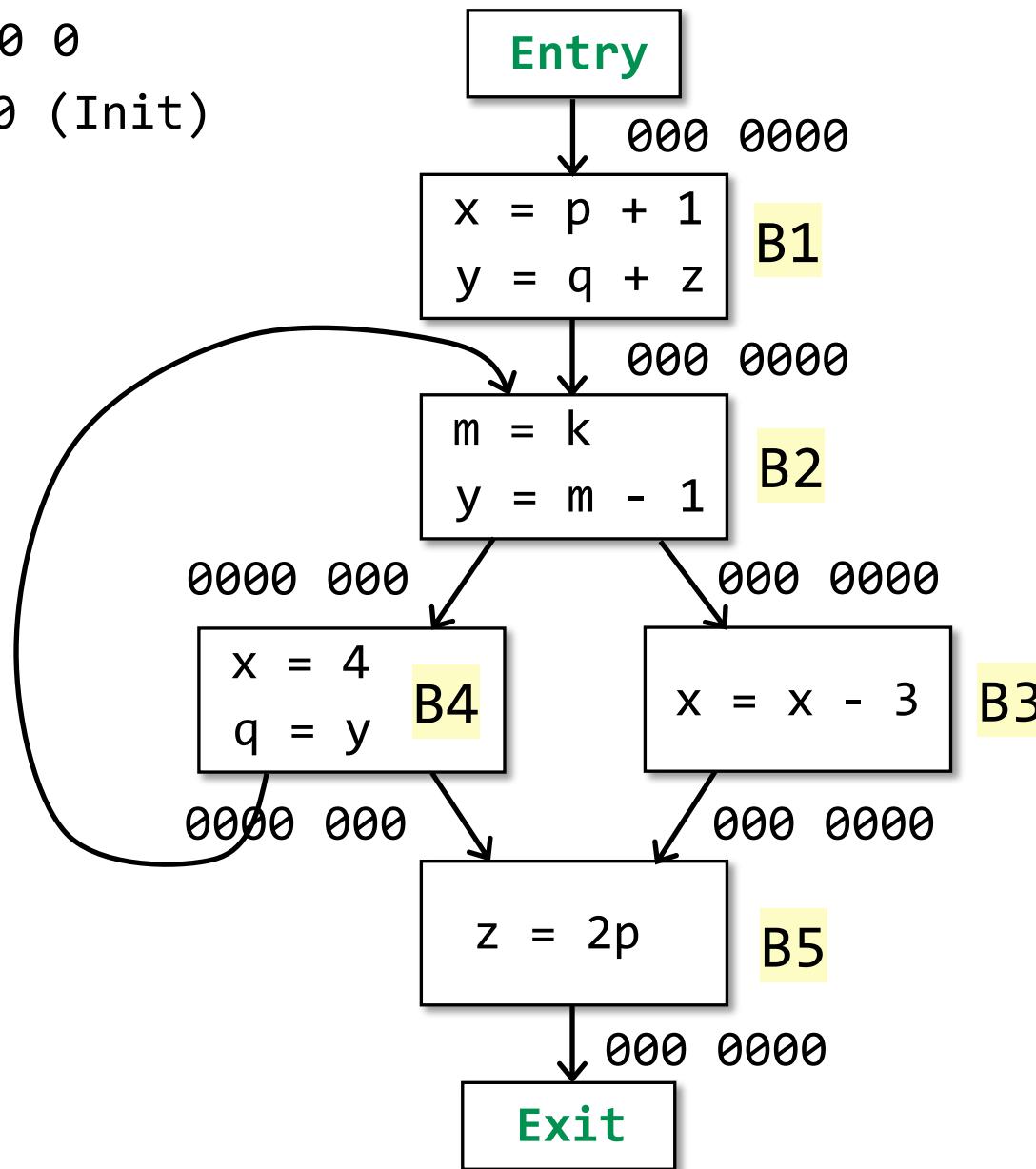


x	y	z	p	q	m	k
0	0	0	0	0	0	0



x y z p q m k
0 0 0 0 0 0 0

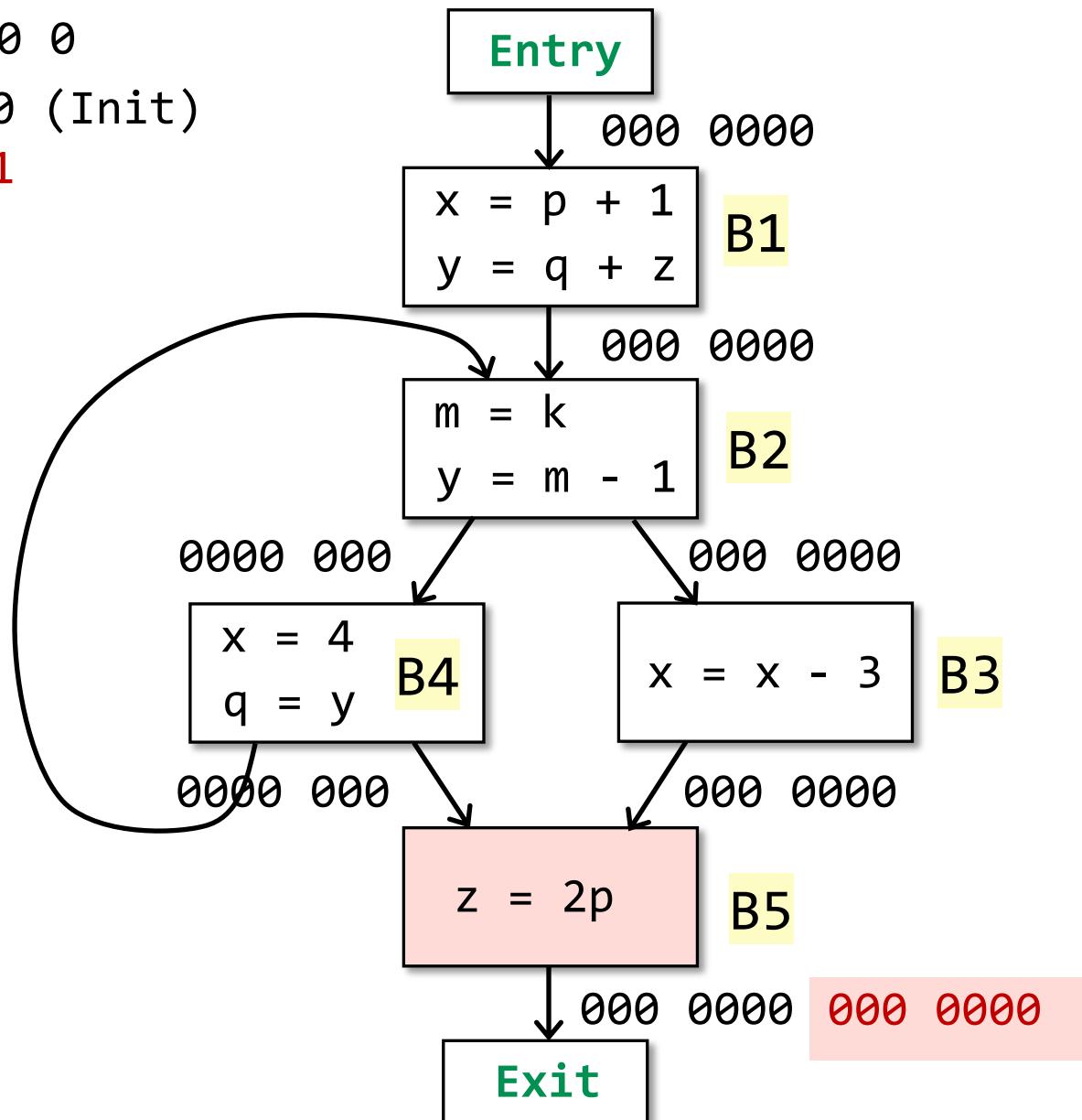
Iteration 0 (Init)



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

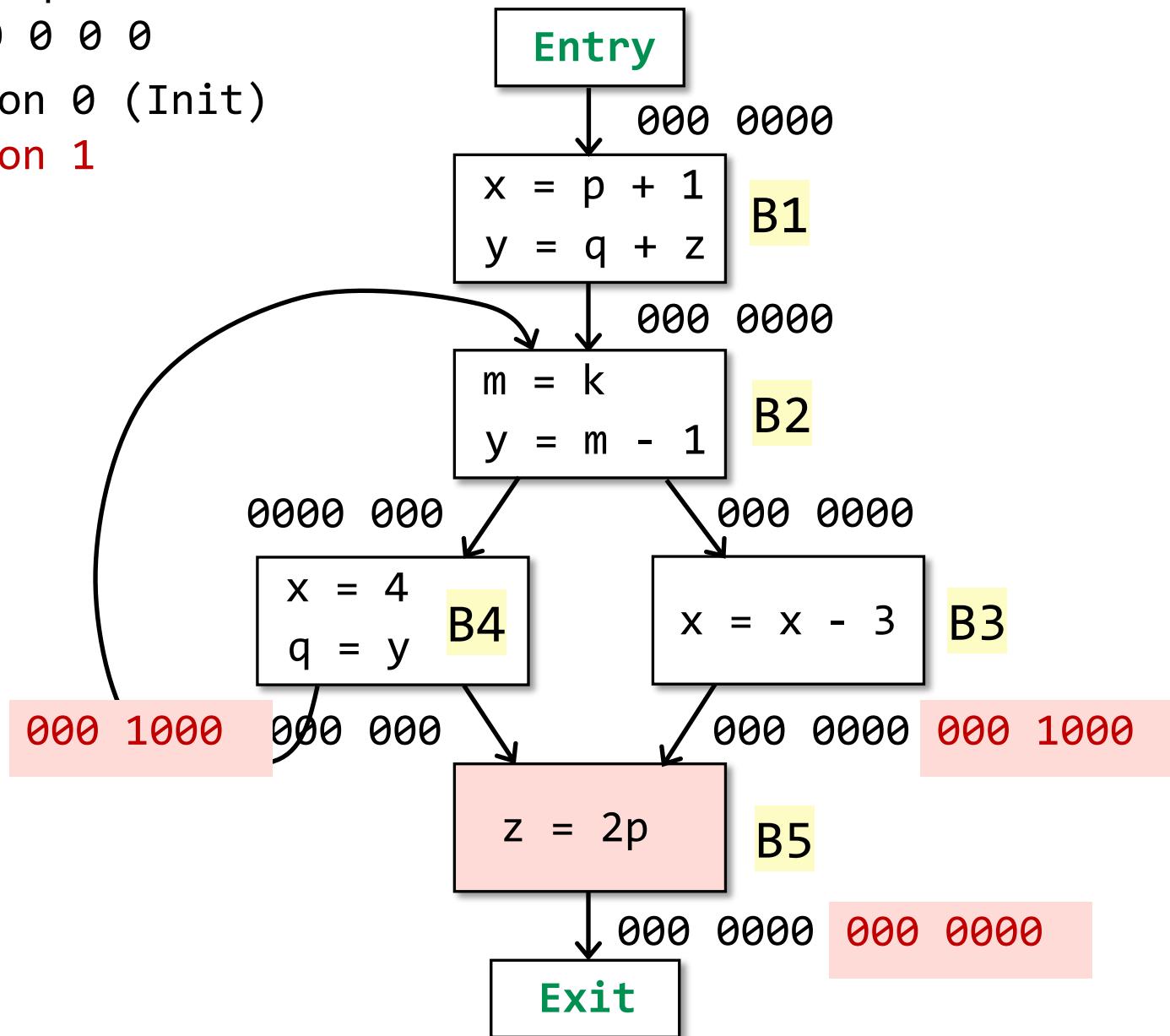


$$IN[B] = use_B \cup (OUT[B] - def_B)$$

x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

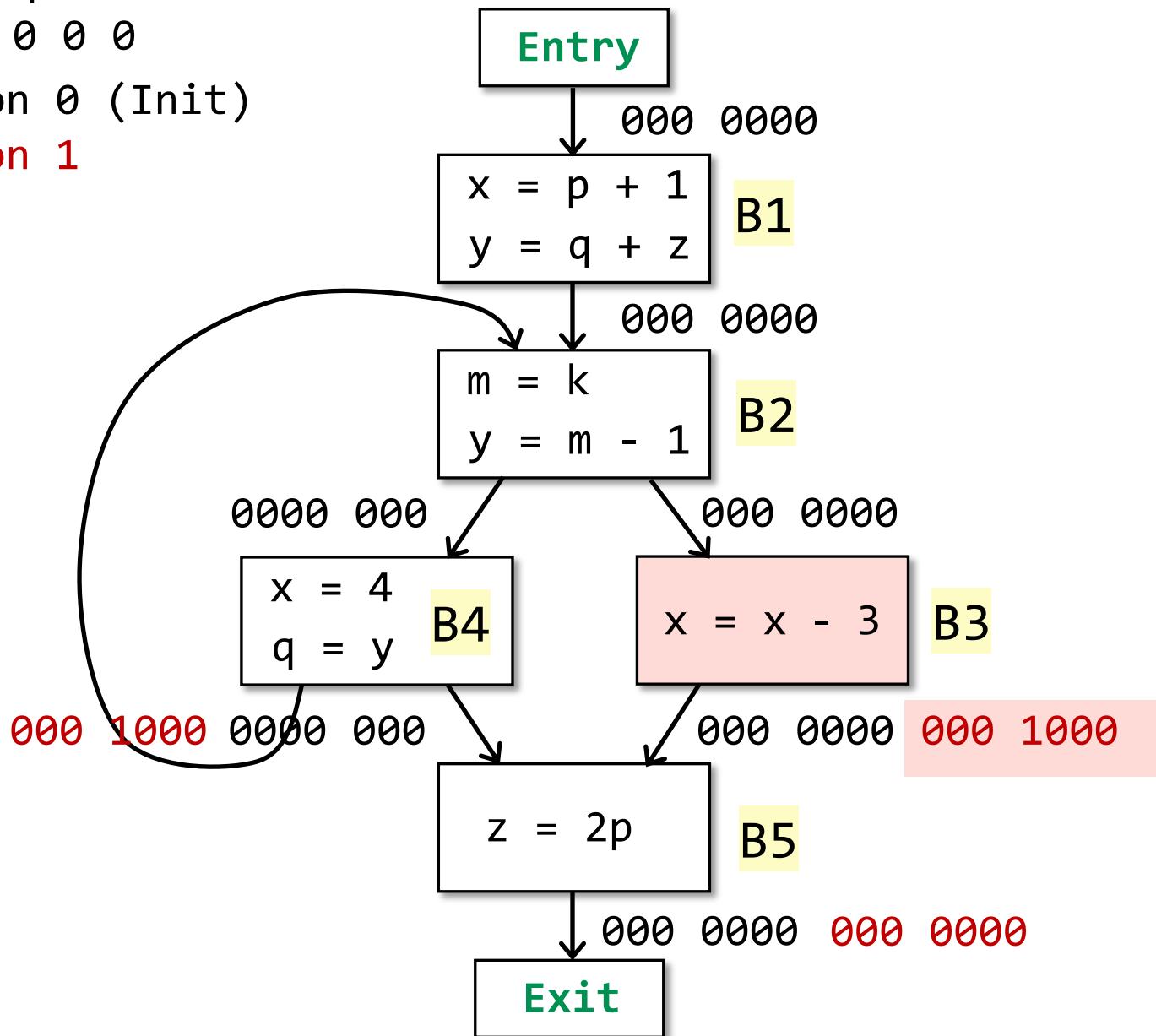
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

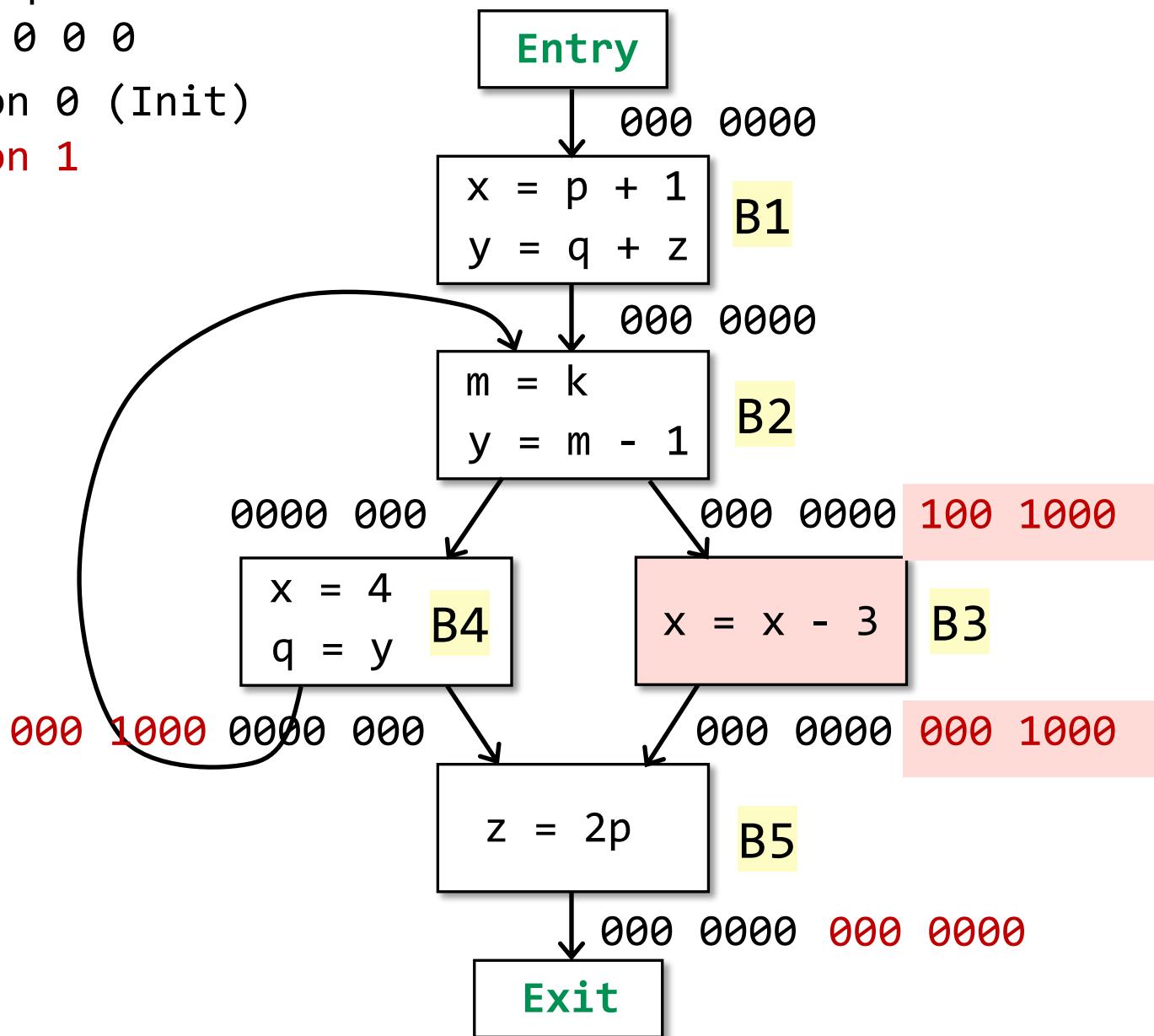
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

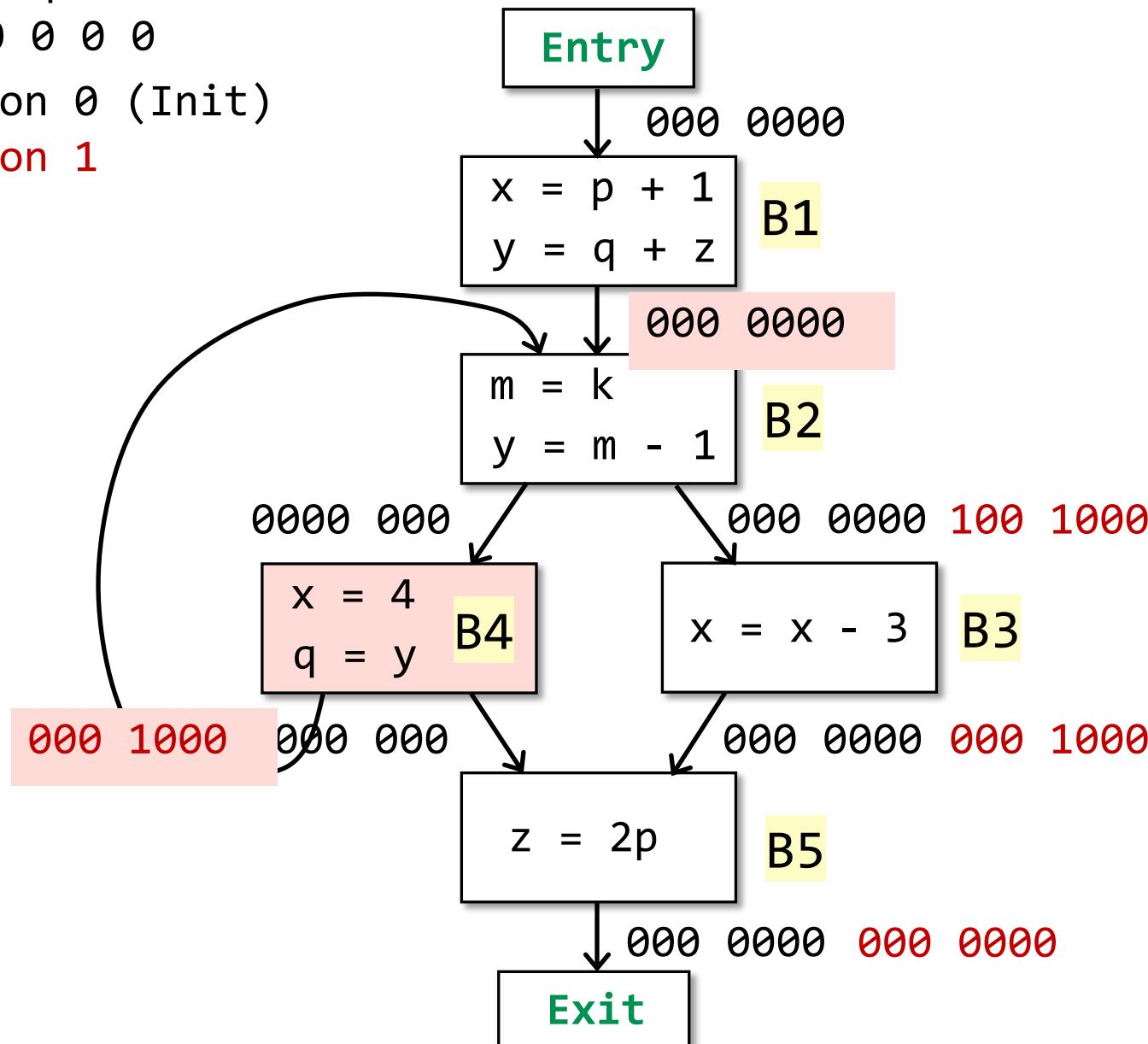
Iteration 1



x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1

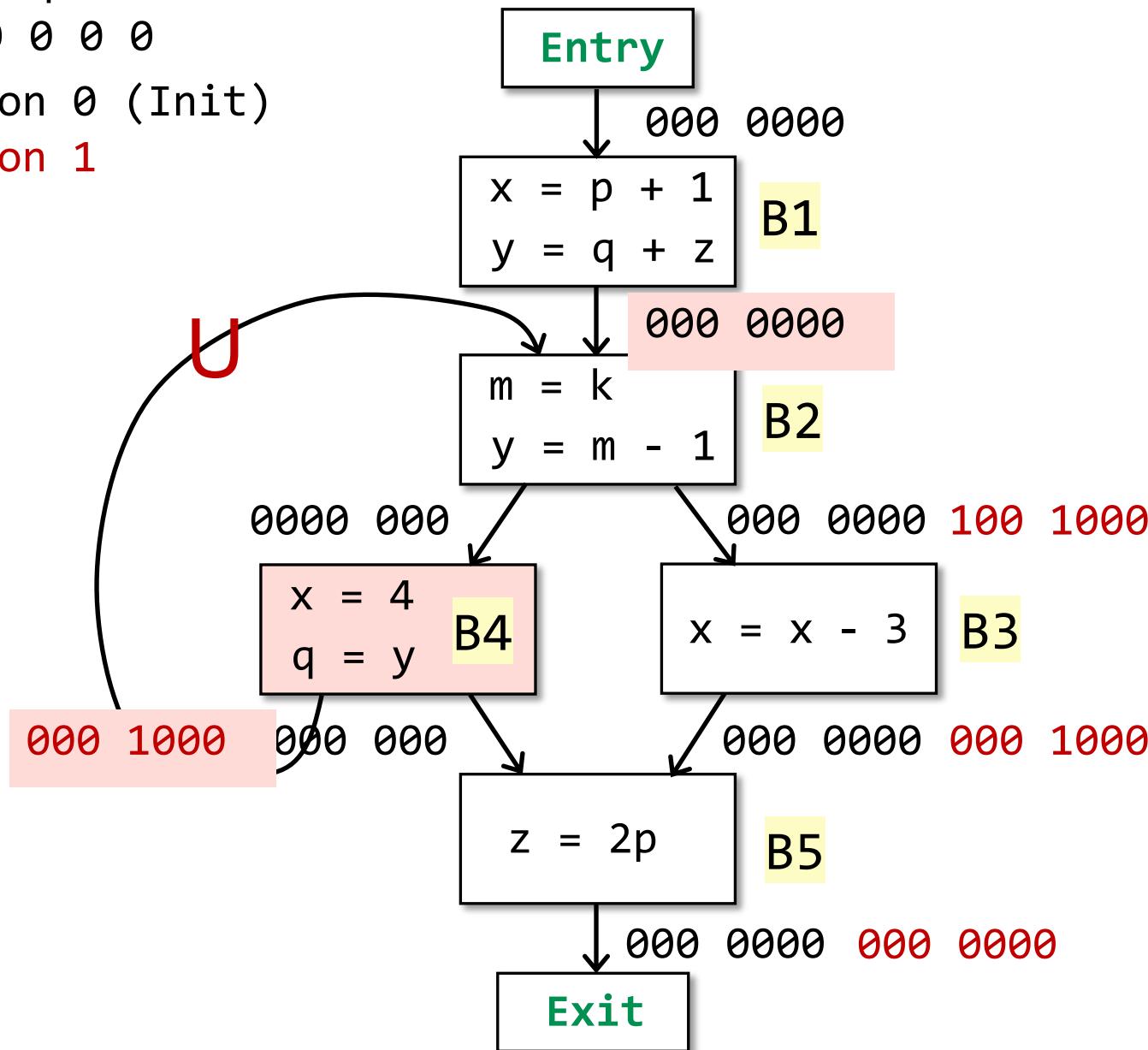


$$OUT[B] = \bigcup_{S \text{ a successor of } B} IN[S]$$

x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

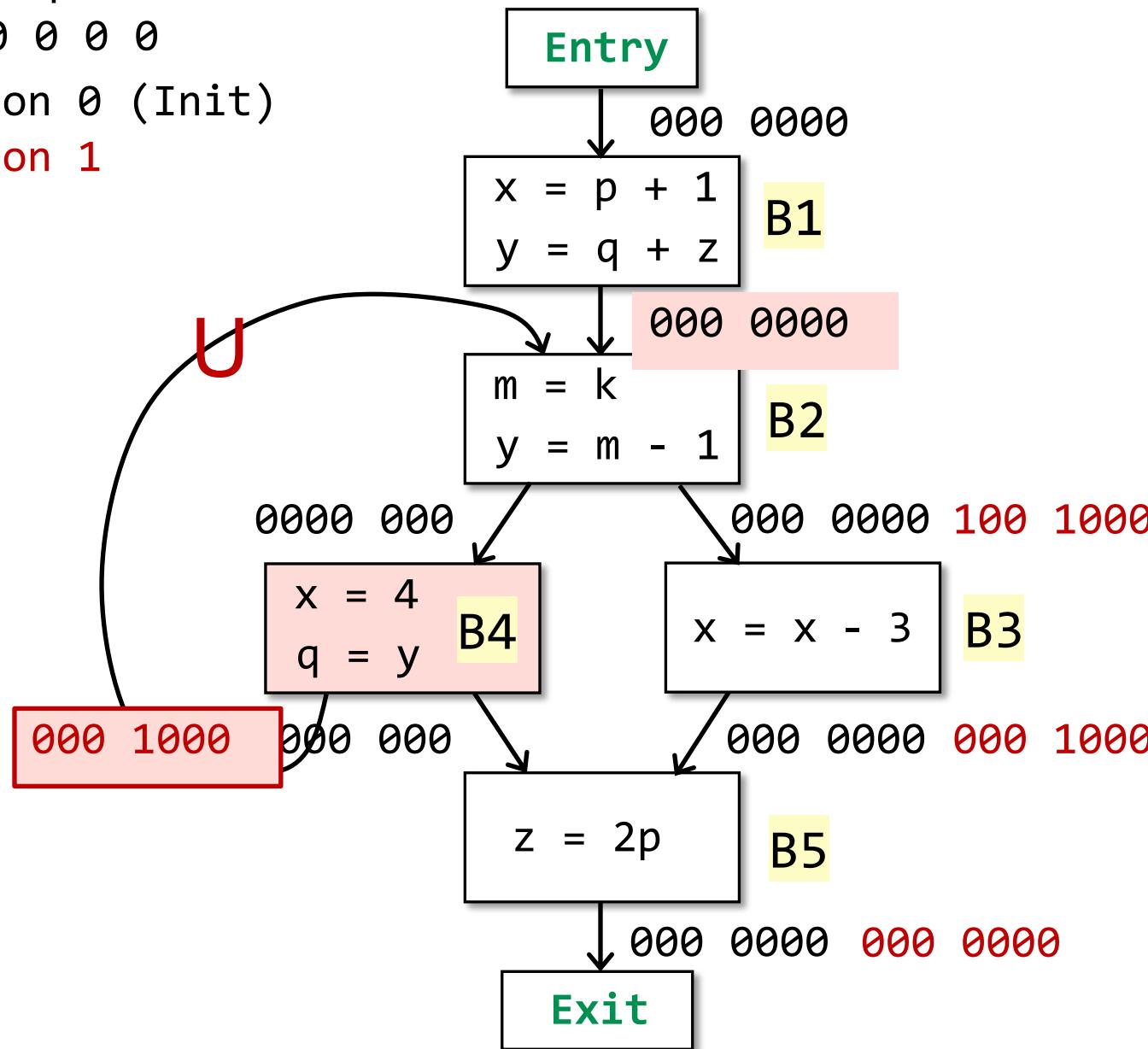
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

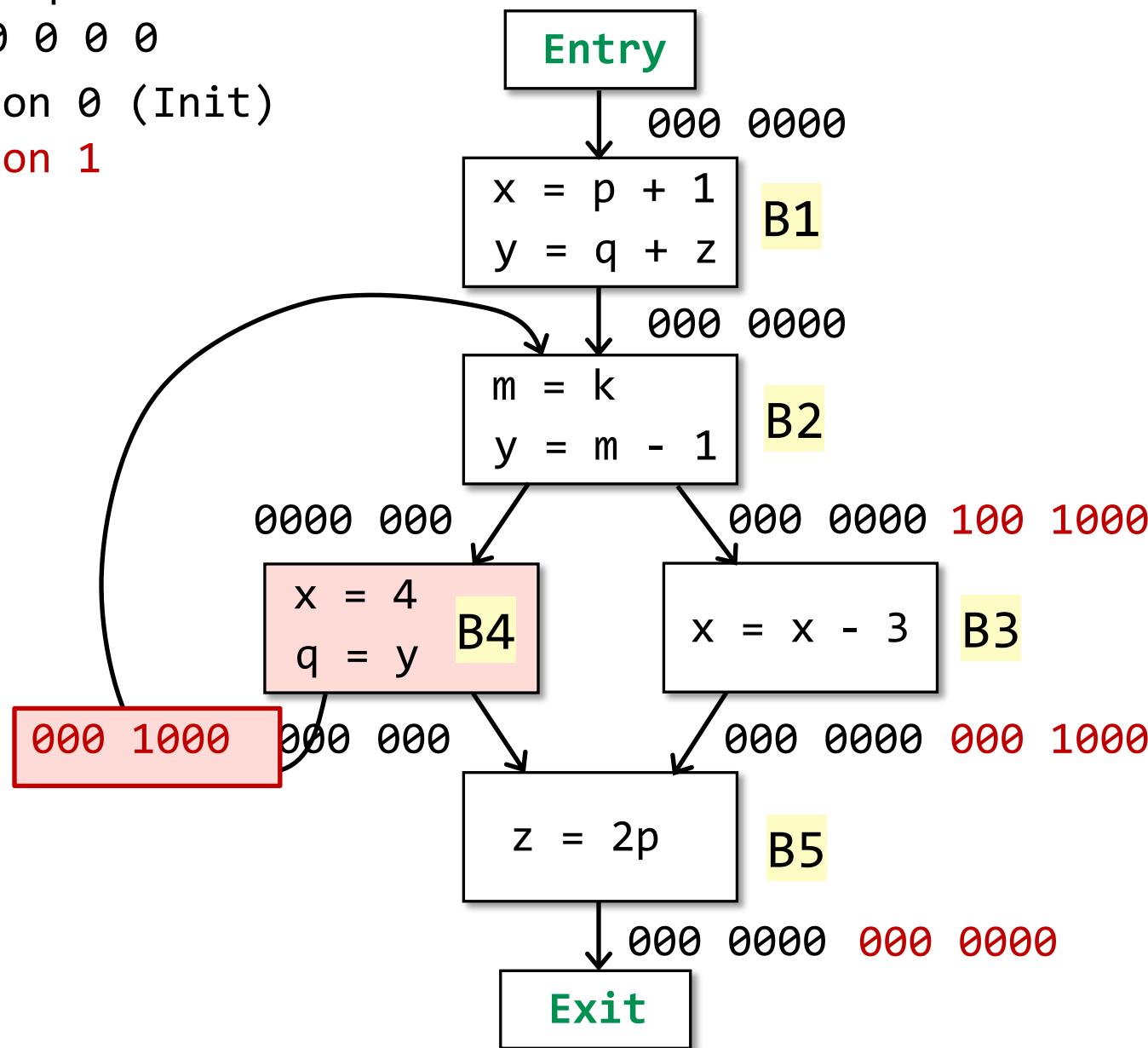
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

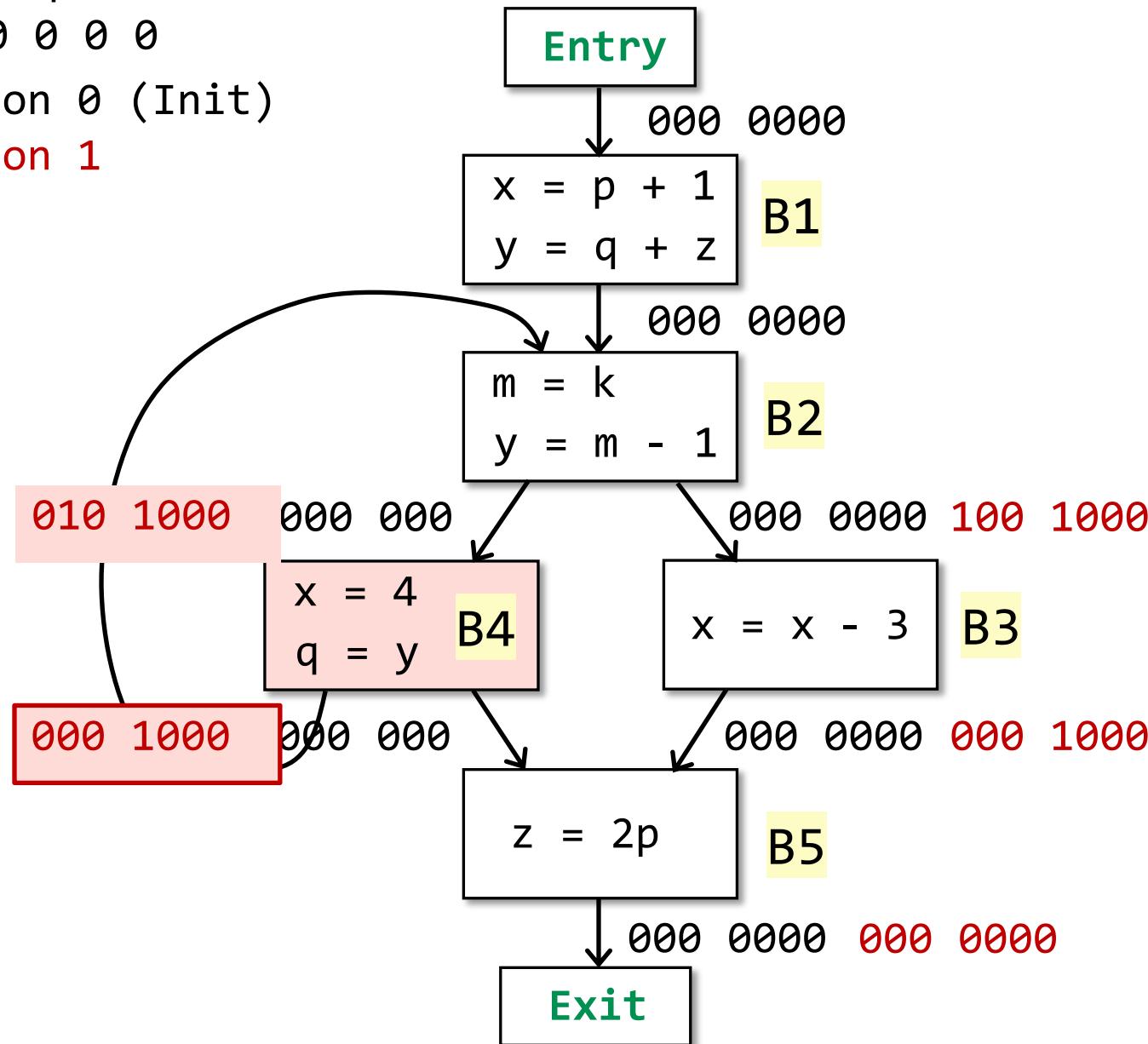
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

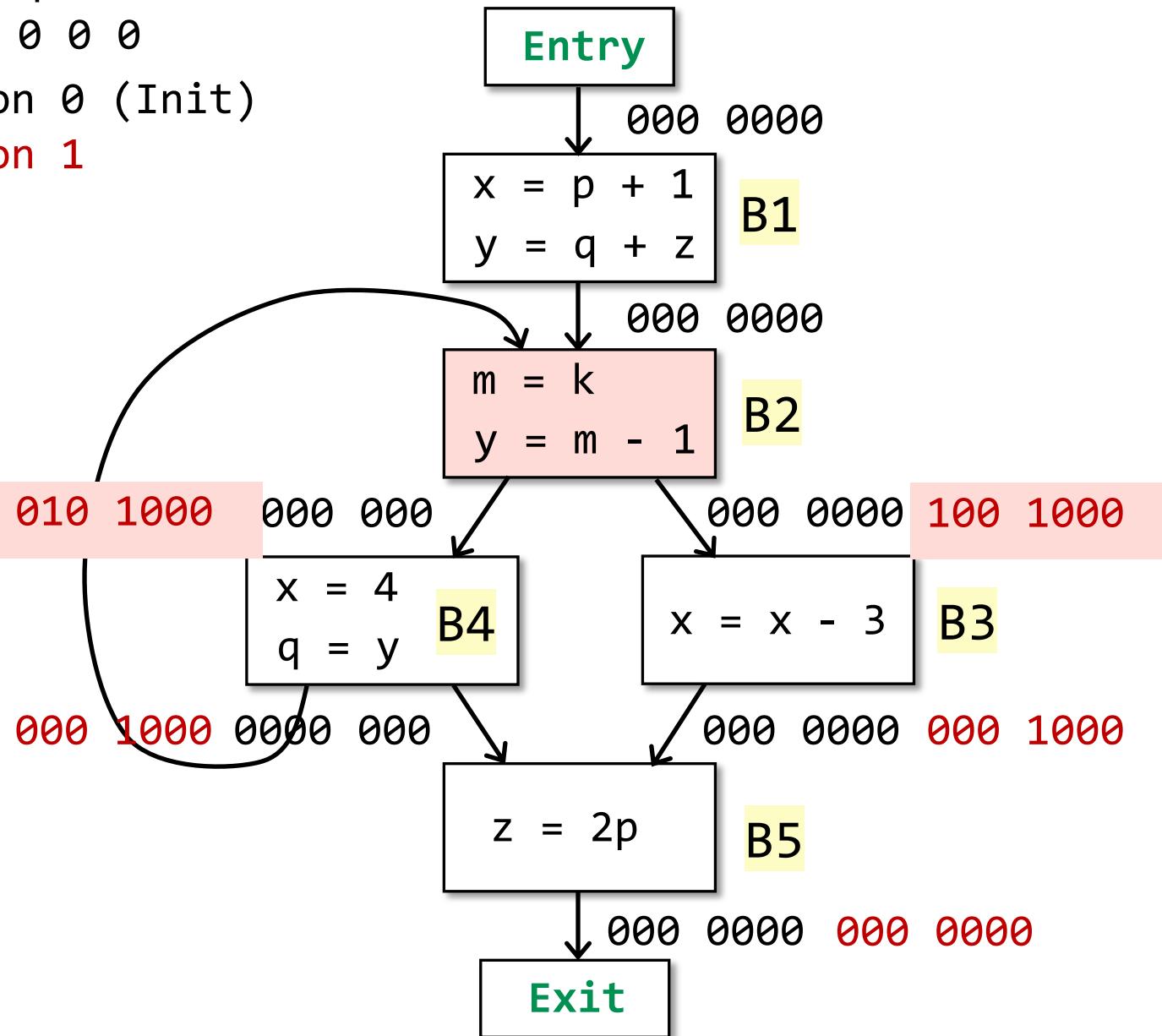
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

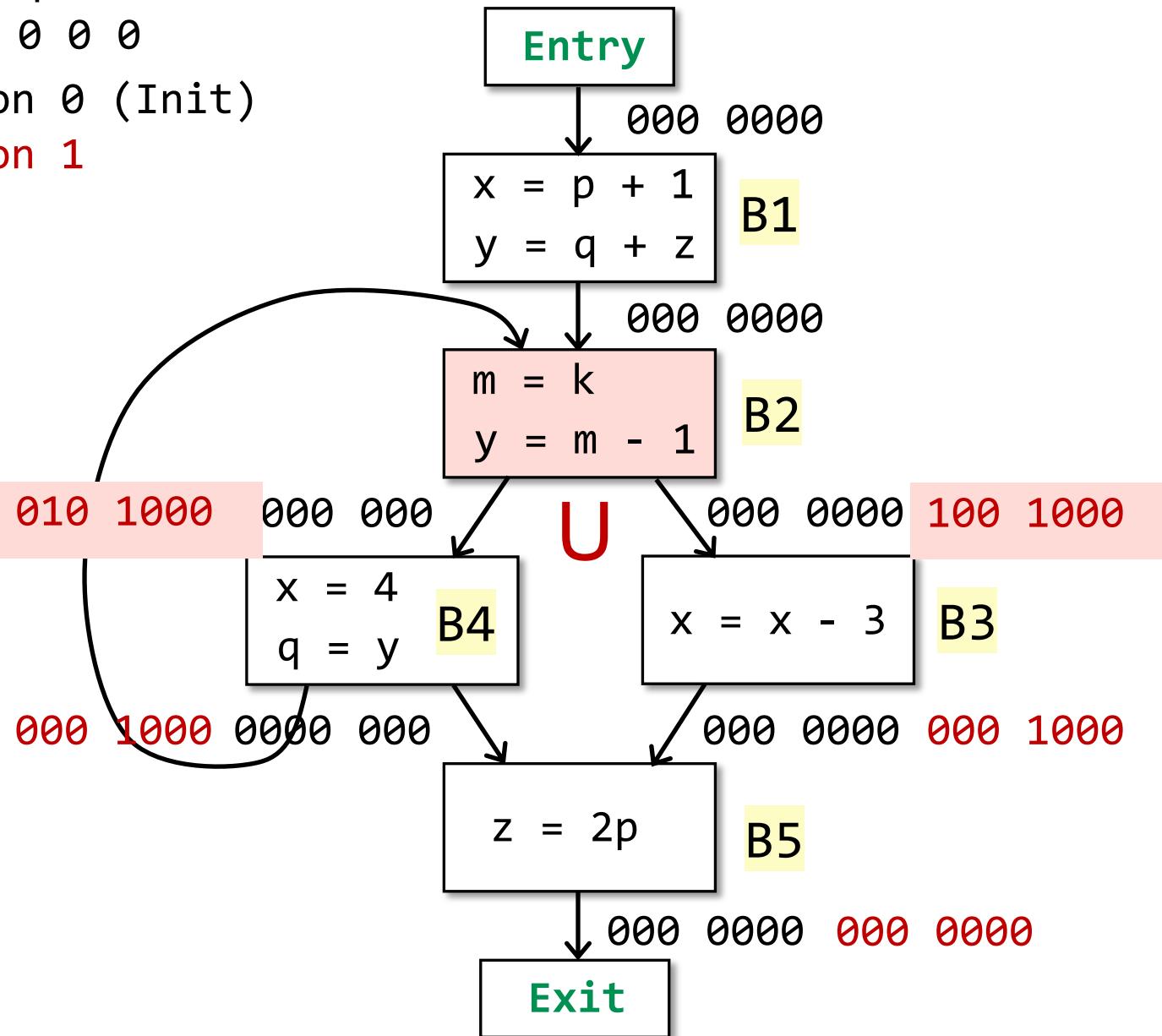
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

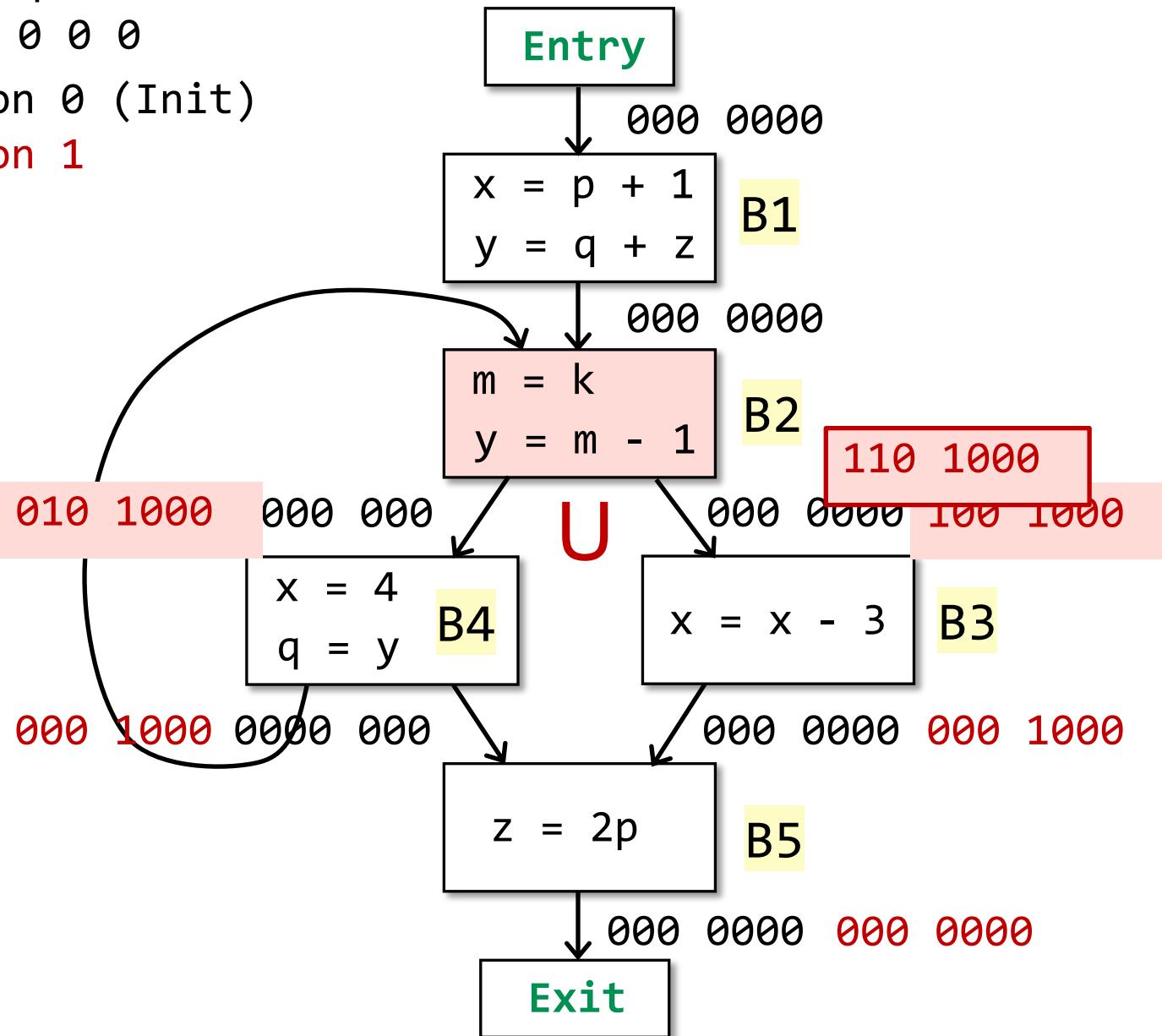
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

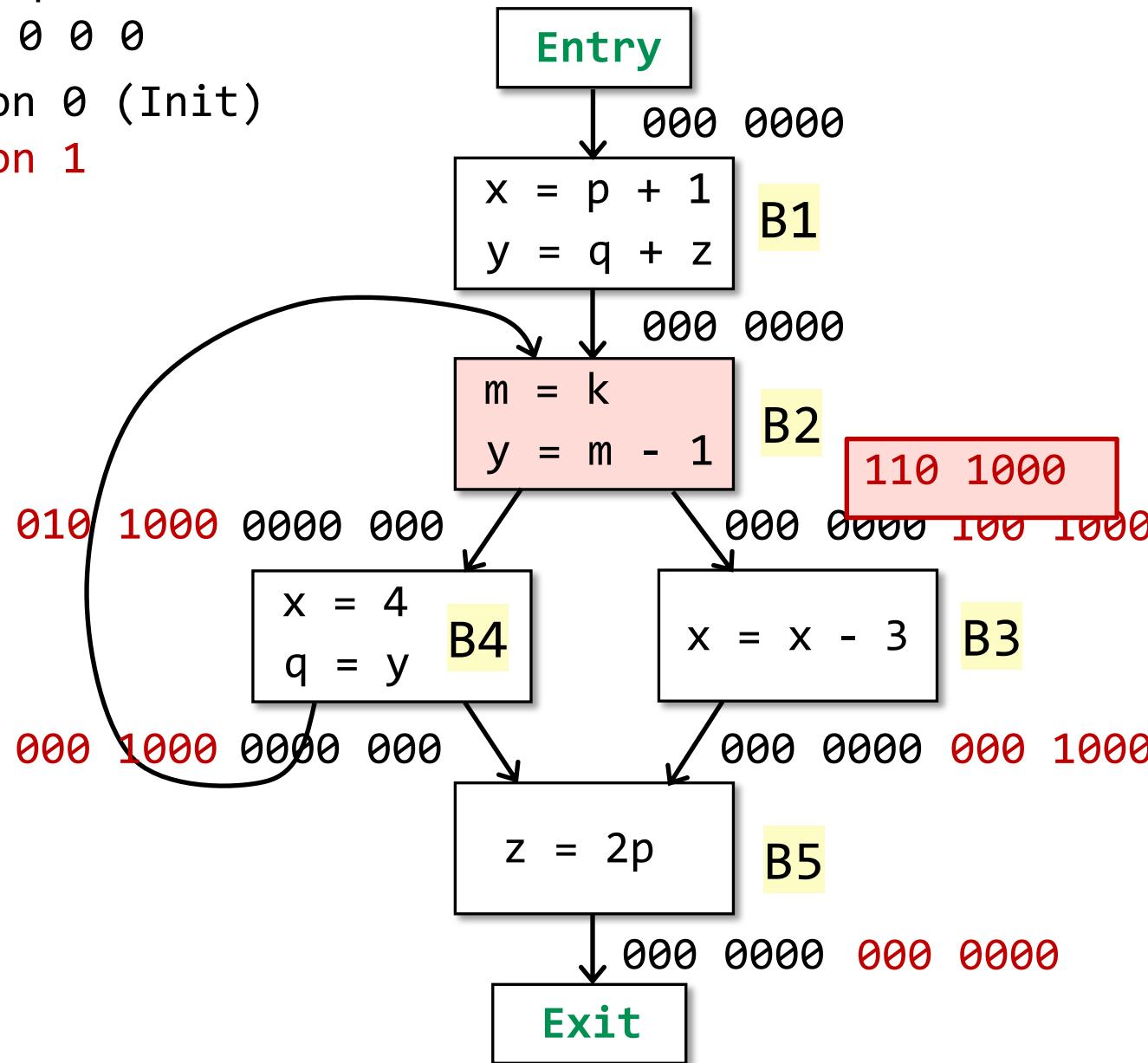
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

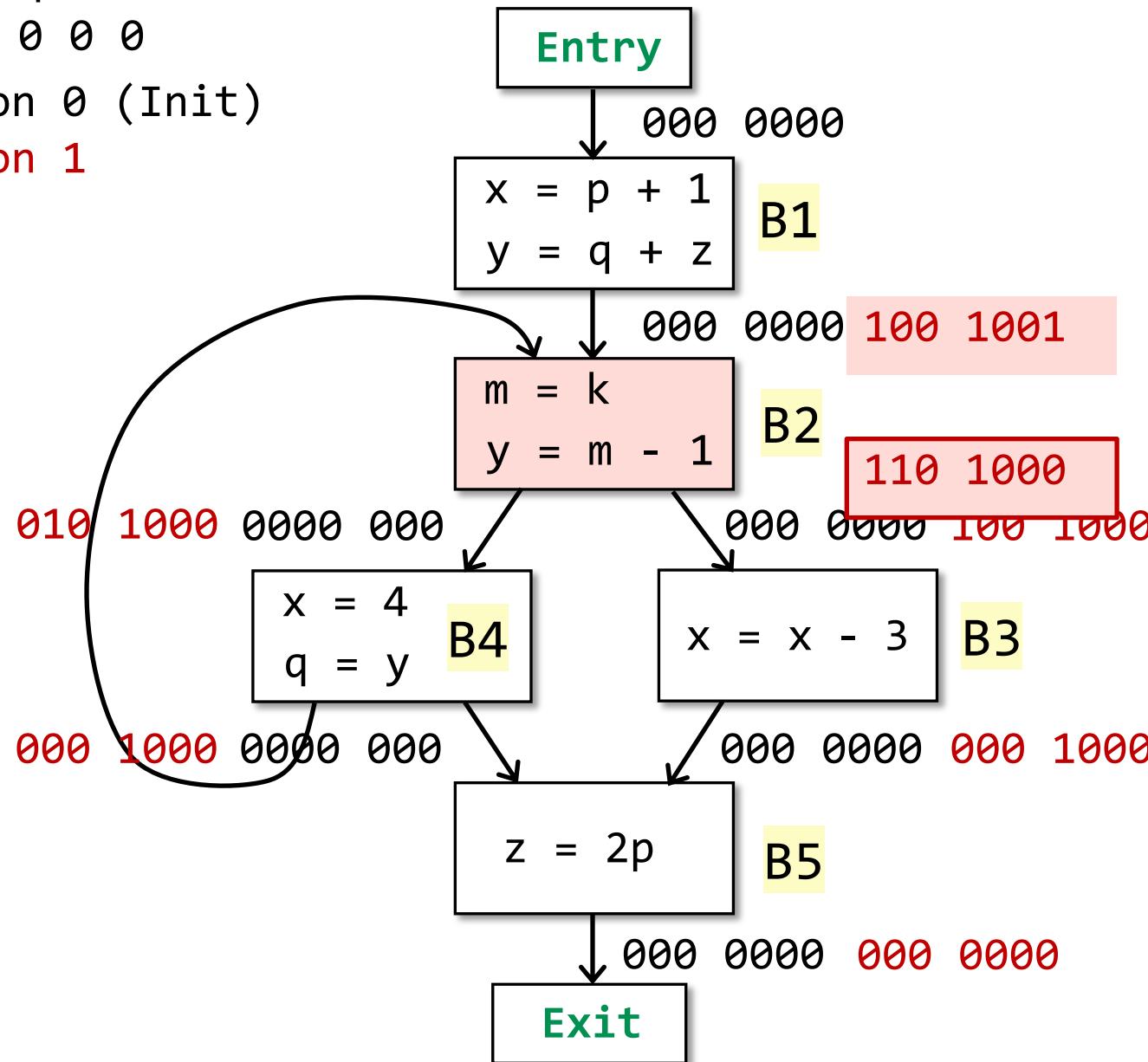
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

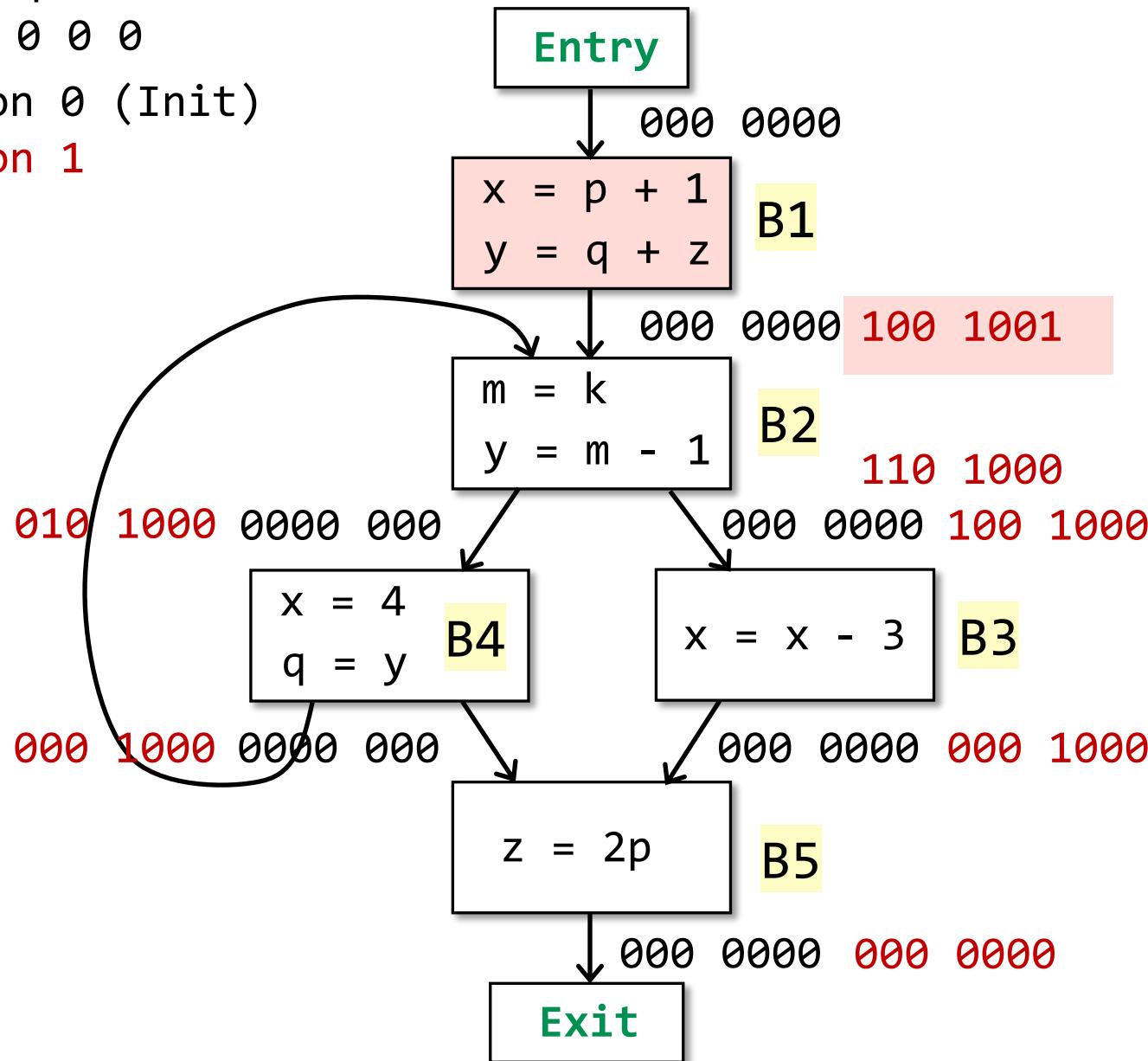
Iteration 1



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

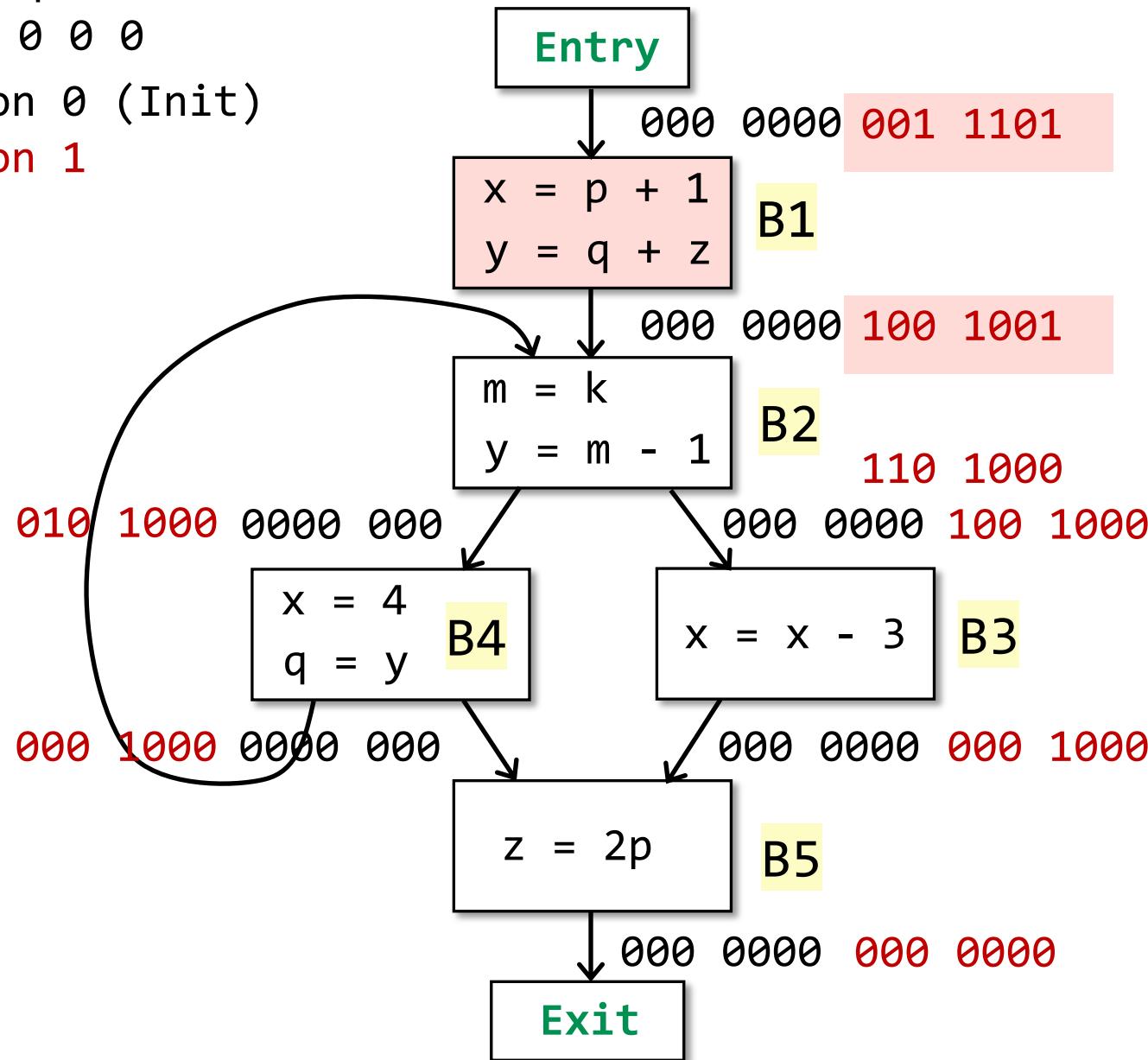
Iteration 1



x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

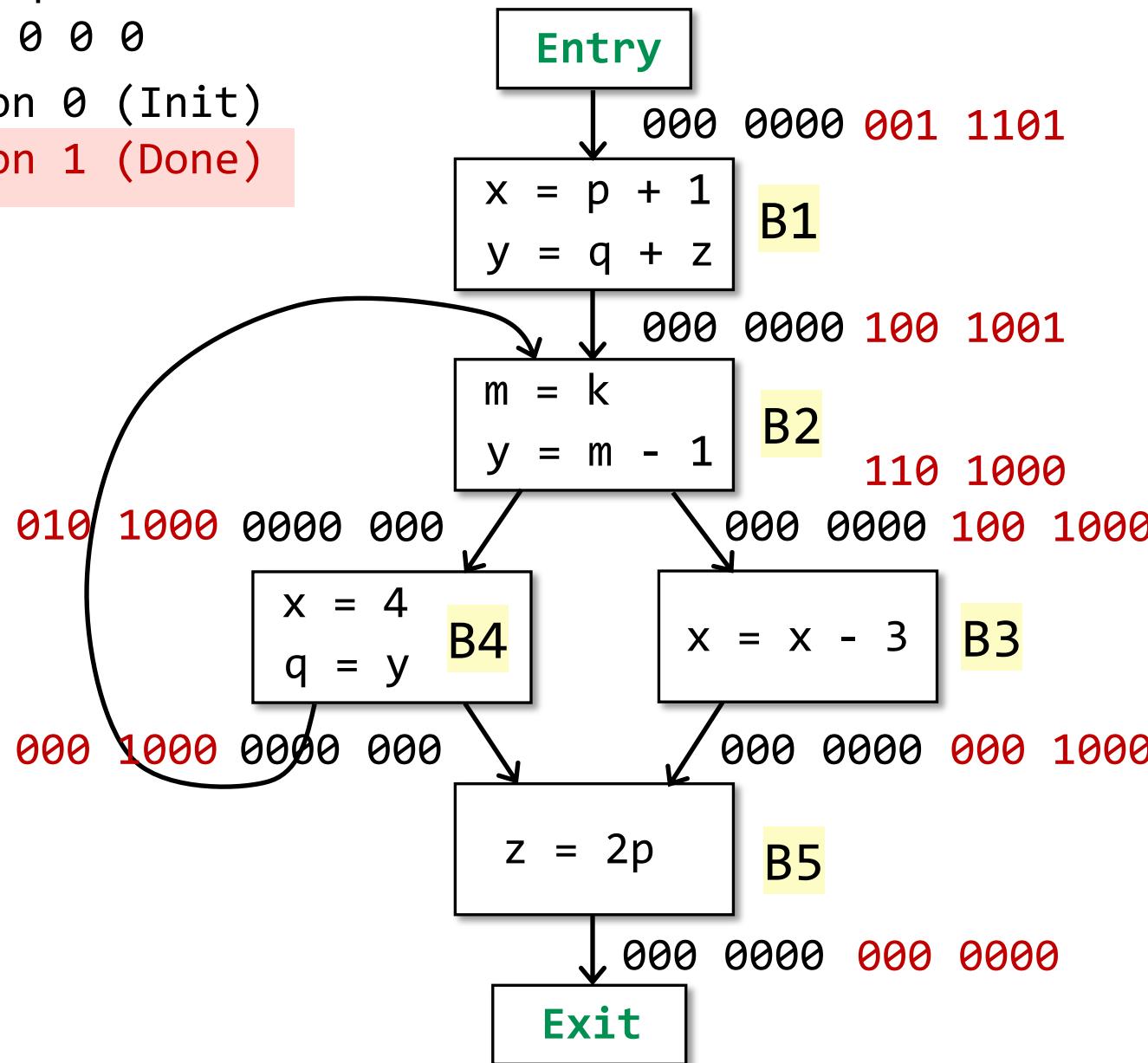
Iteration 1



x	y	z	p	q	m	k
0	0	0	0	0	0	0

Iteration 0 (Init)

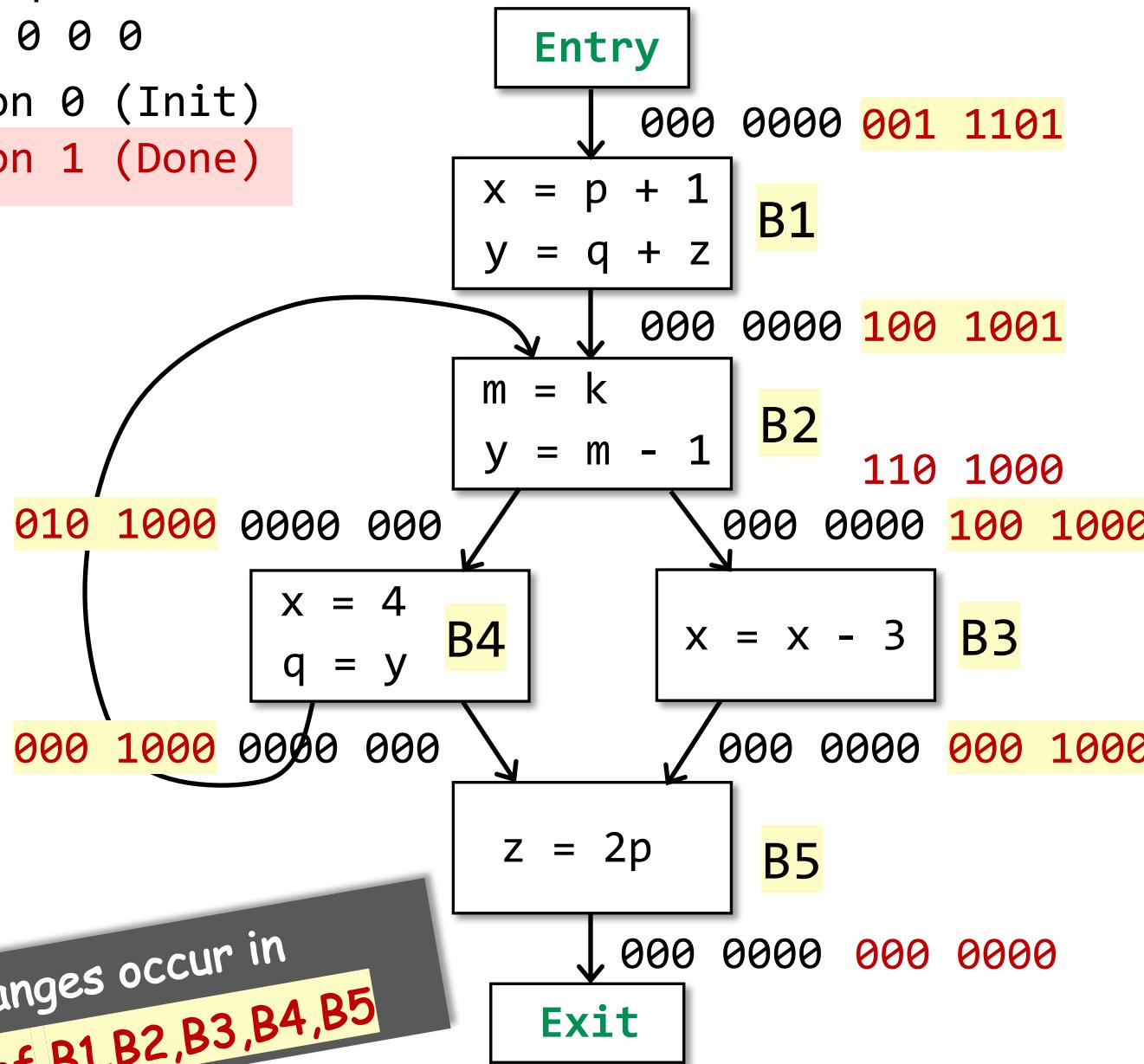
Iteration 1 (Done)



x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

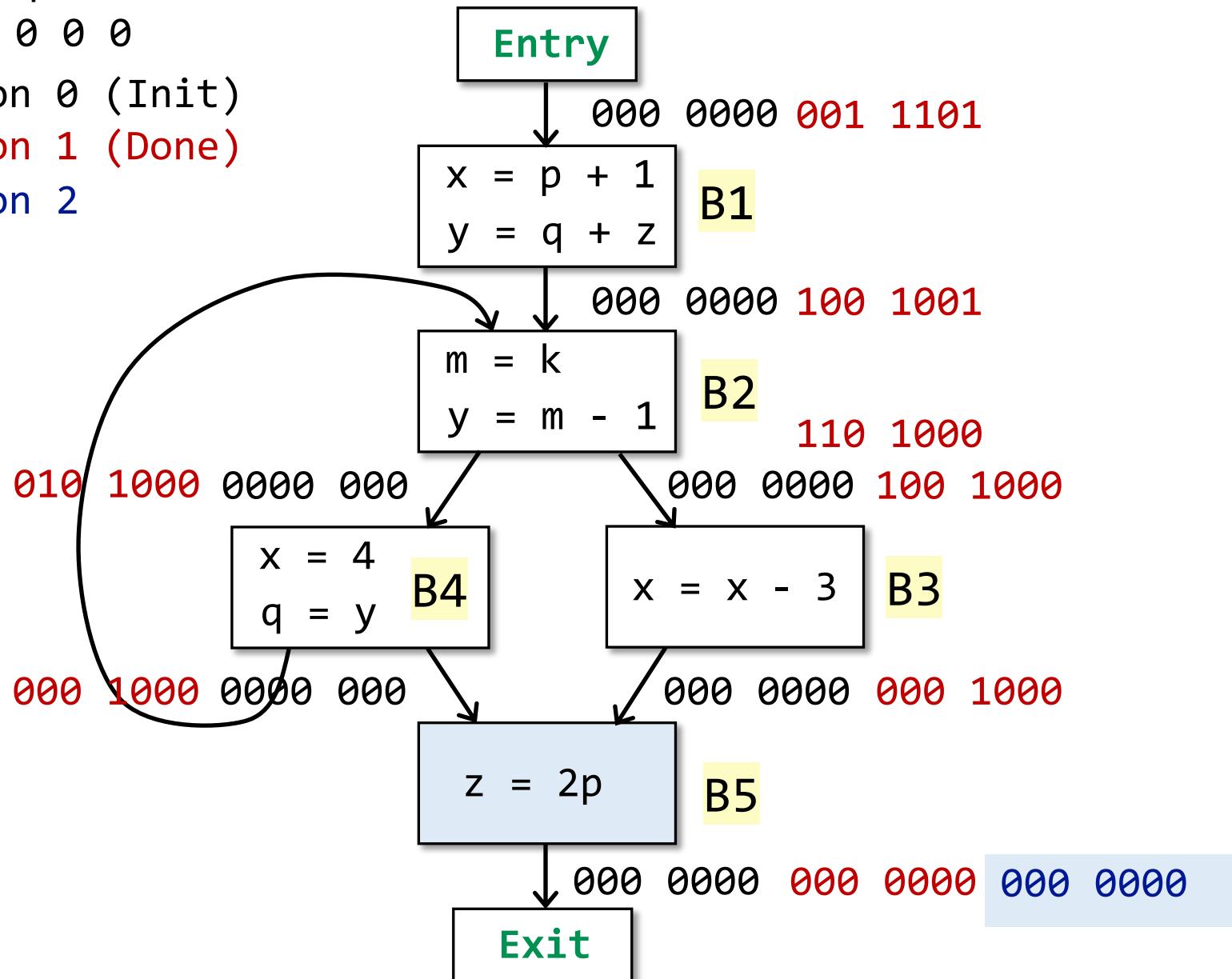


x	y	z	p	q	m	k
0	0	0	0	0	0	0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

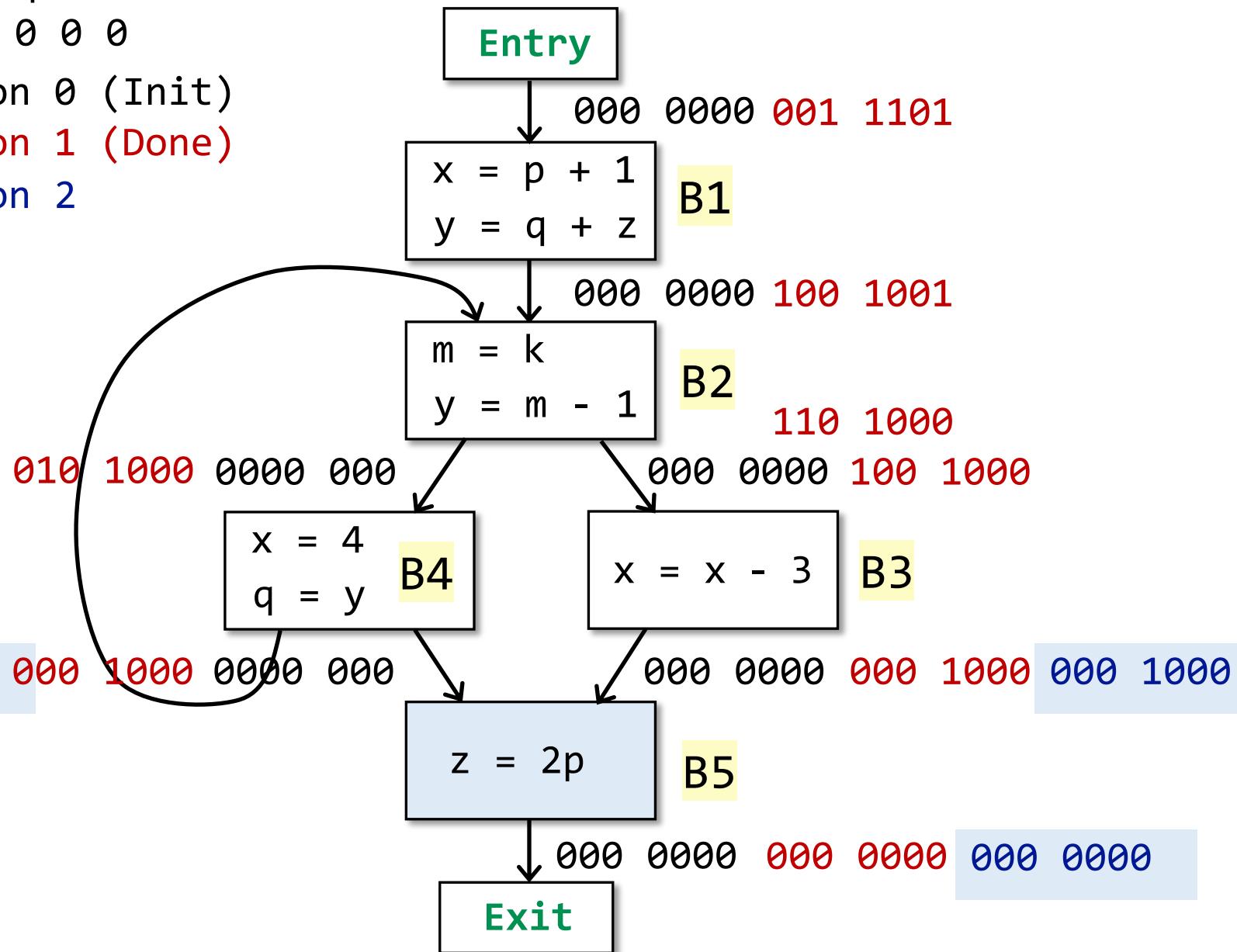


x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

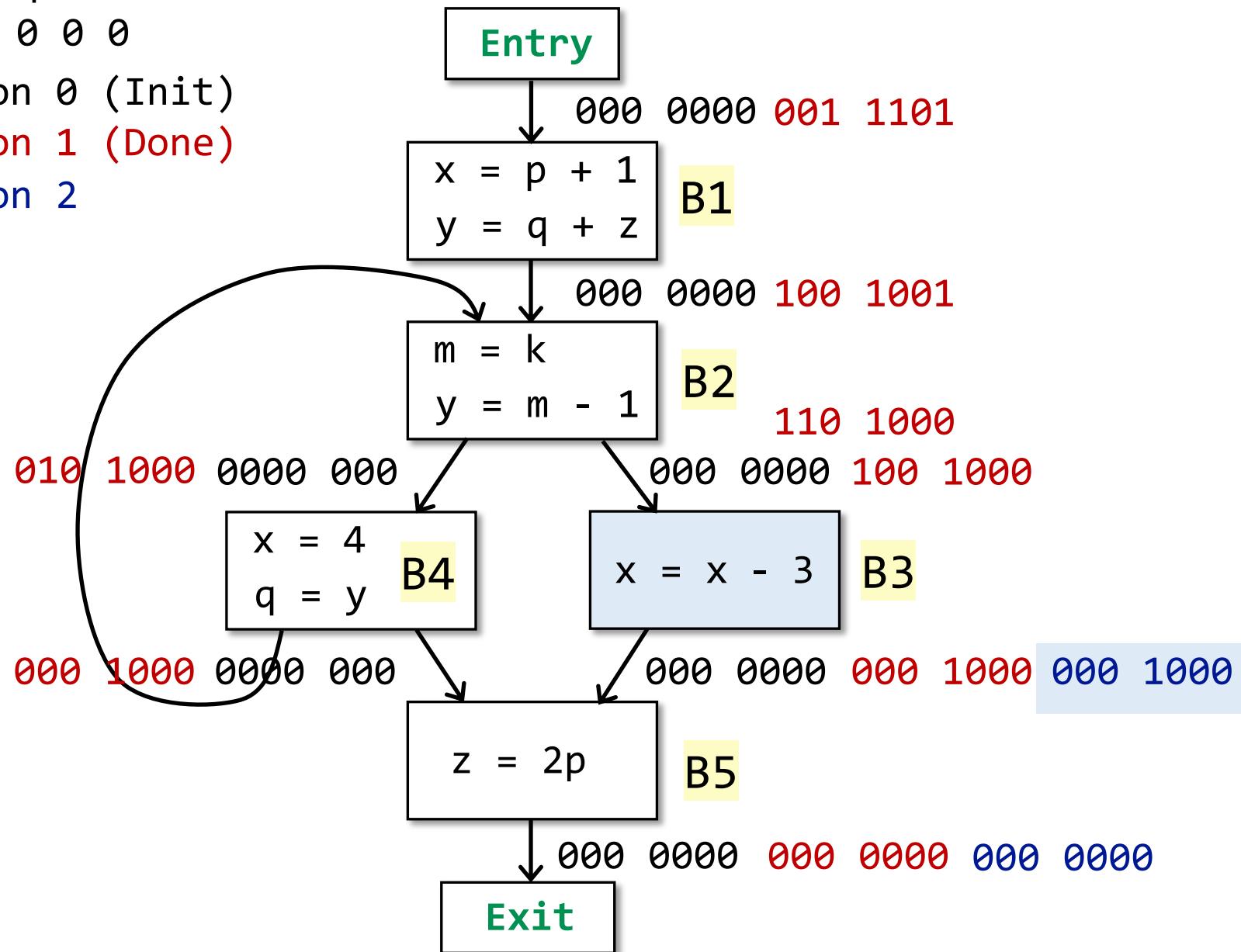


x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

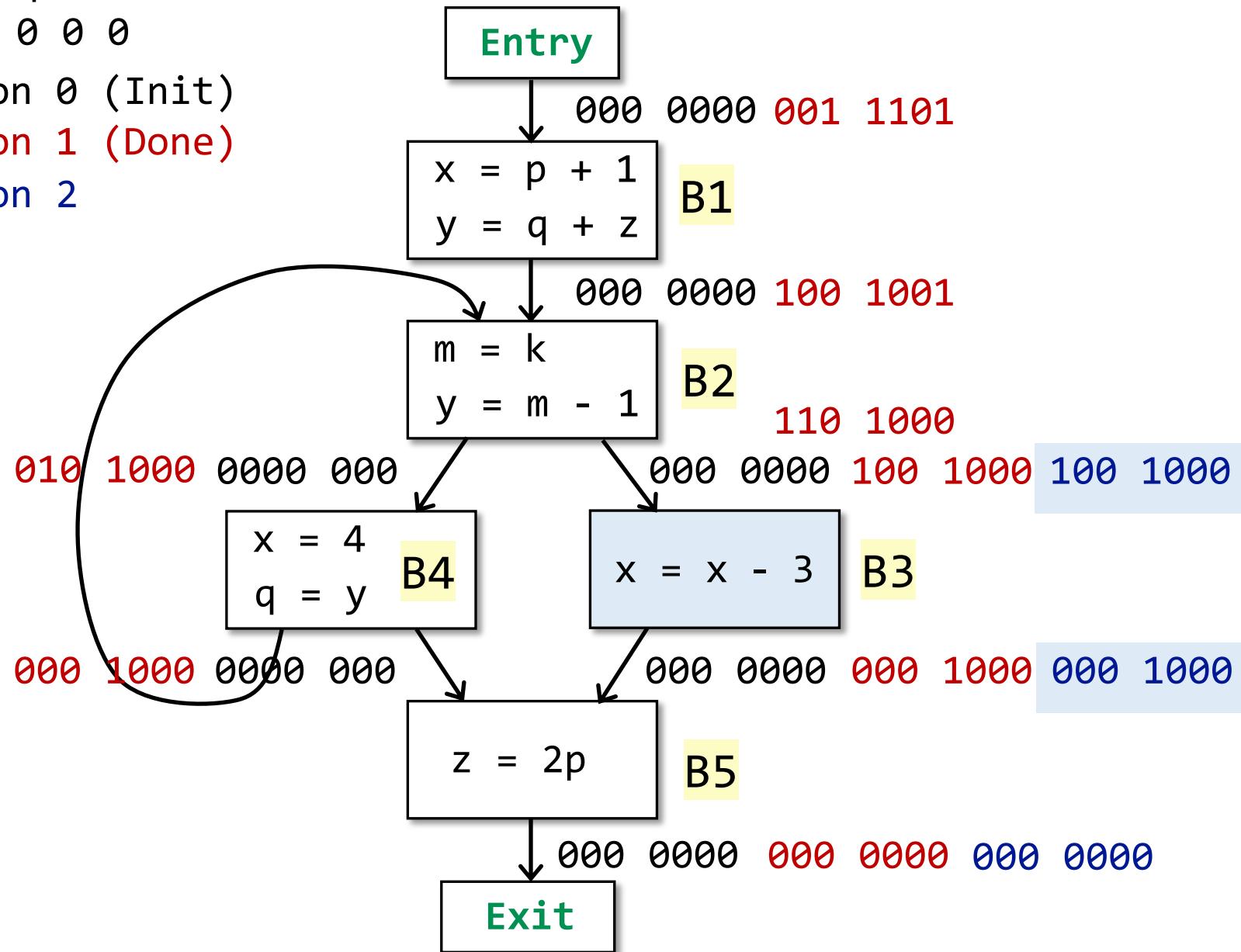


x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

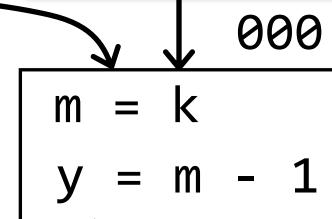
Iteration 2

Entry



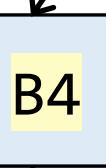
x = p + 1
y = q + z

B1



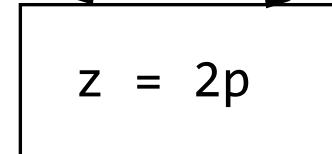
m = k
y = m - 1

B2



x = 4
q = y

B4



z = 2p

B5

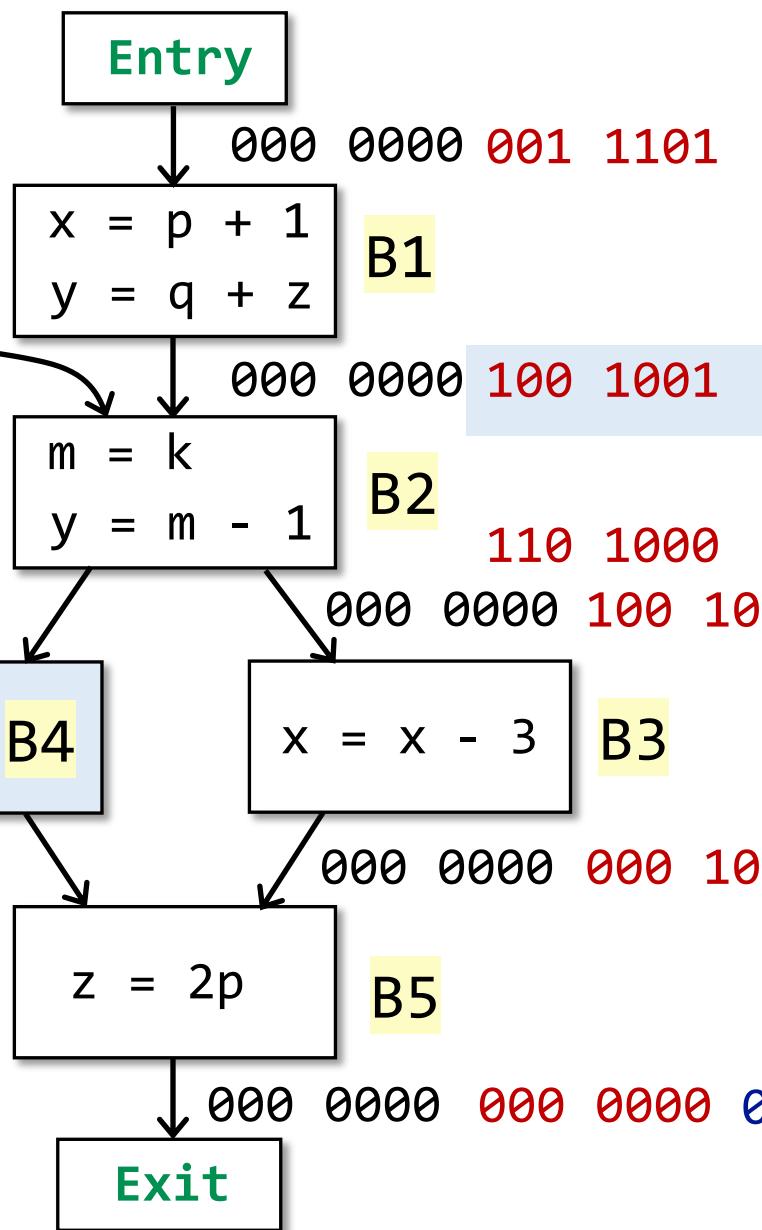
x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry

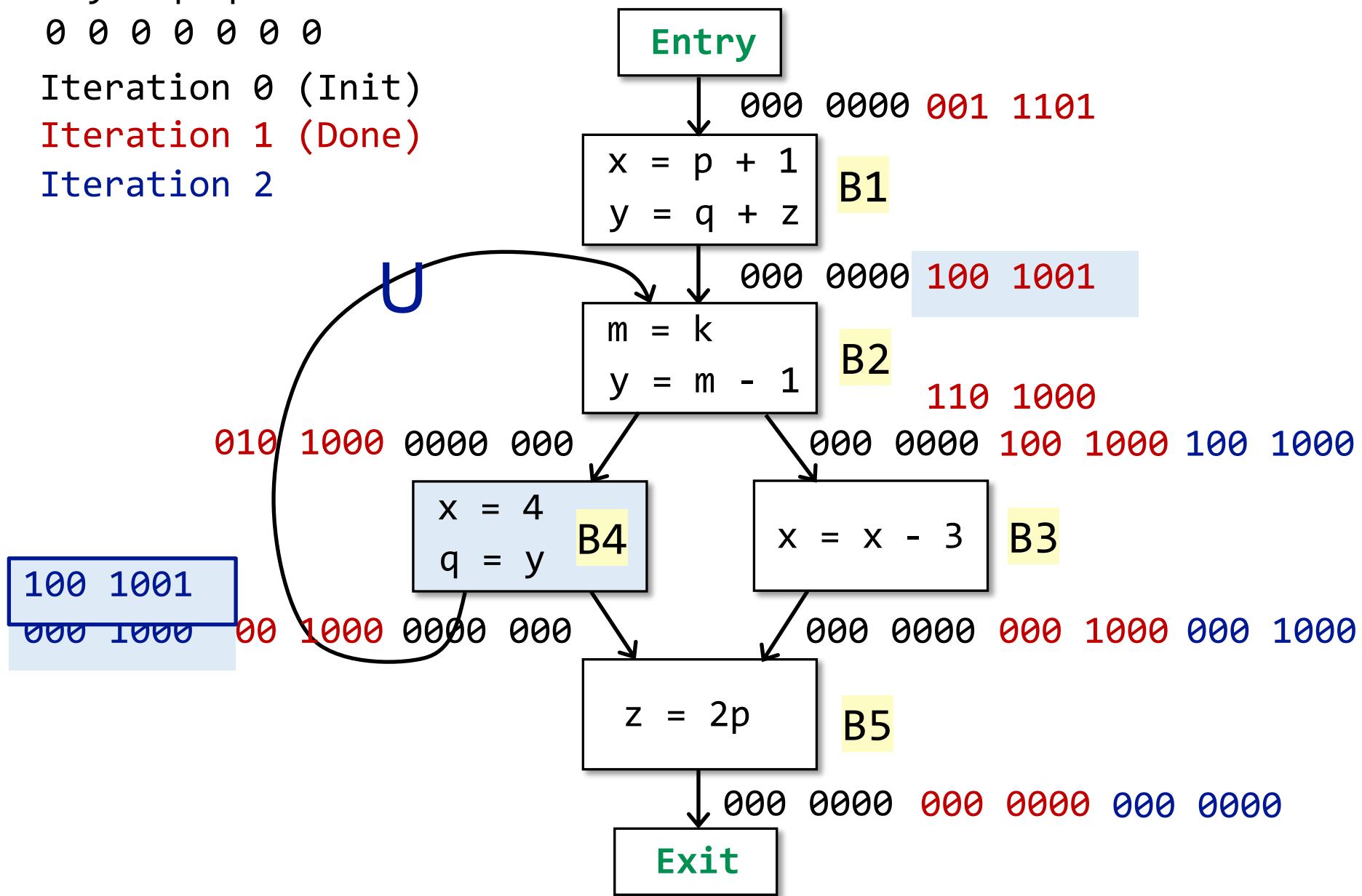


x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2



x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



010 1000 0000 000
 100 1001
 000 1000 000 1000

000 0000 001 1101
 B1

000 0000 100 1001
 B2

110 1000

000 0000 100 1000 100 1000
 B3

000 1000 000 1000

000 0000 000 0000 000 0000
 B5

000 0000 000 0000 000 0000
 Exit

x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

010 1001 10 1000 0000 000
100 1001
000 1000 000 1000 000 000

Entry

000 0000 001 1101

x = p + 1
y = q + z

B1

m = k
y = m - 1

B2

000 0000 100 1001

110 1000

000 0000 100 1000 100 1000

x = 4
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

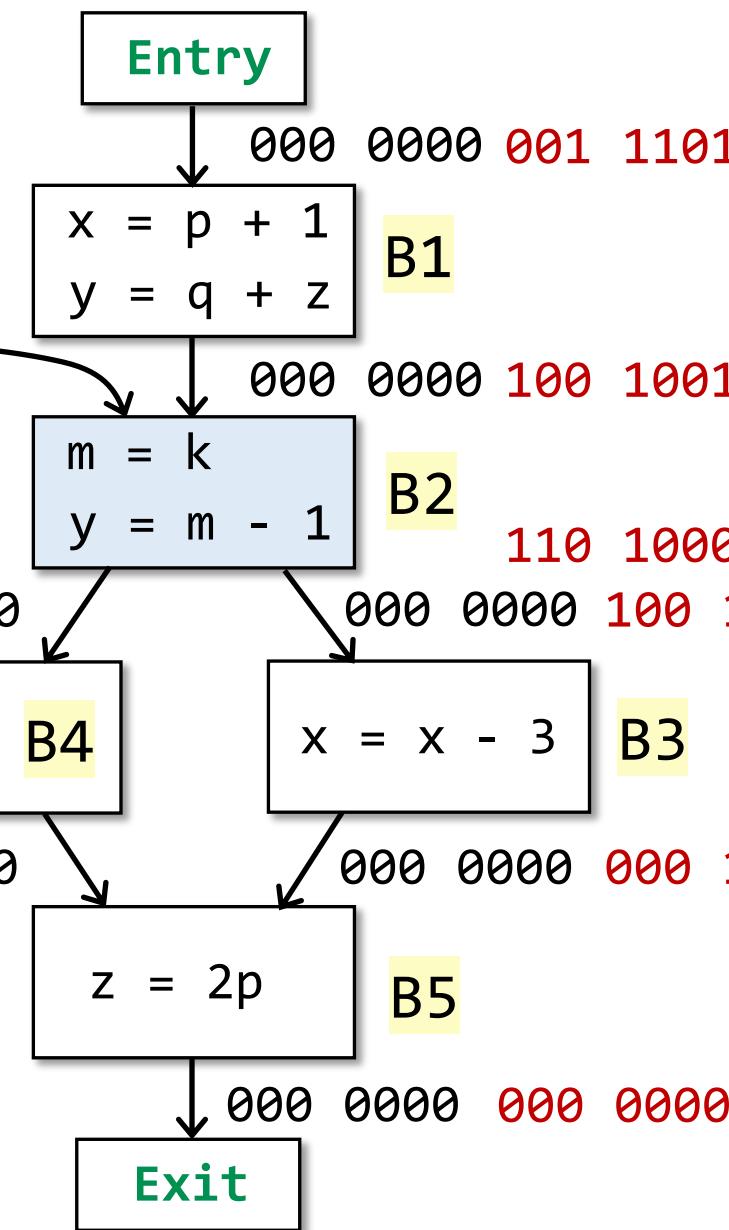
x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



010 1001	10 1000	0000 000	B1	000 0000 100 1001	110 1000	100 1000	100 1000
100 1001	000 1000	000 1000	B2	000 0000 100 1000	100 1000	100 1000	100 1000
000 1000	1000 0000 000	000 0000 000	B3	000 0000 000 1000	000 1000 000 1000	000 1000 000 1000	000 1000 000 1000

x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry

$x = p + 1$
 $y = q + z$

B1

$m = k$
 $y = m - 1$

B2

$x = 4$
 $q = y$

B4

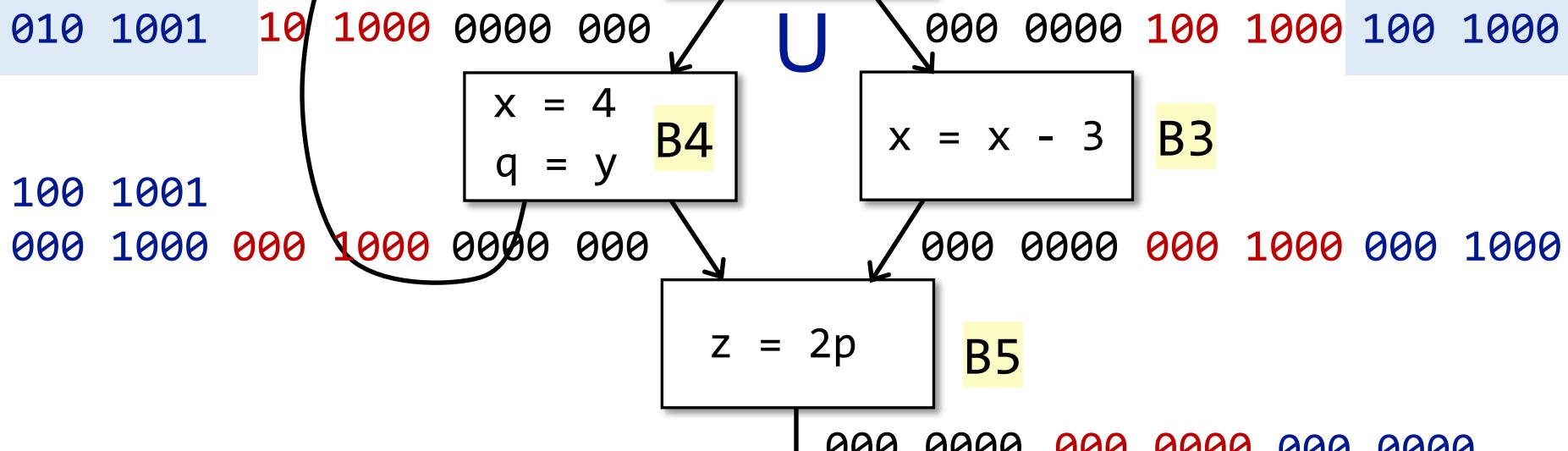
$x = x - 3$

B3

$z = 2p$

B5

Exit



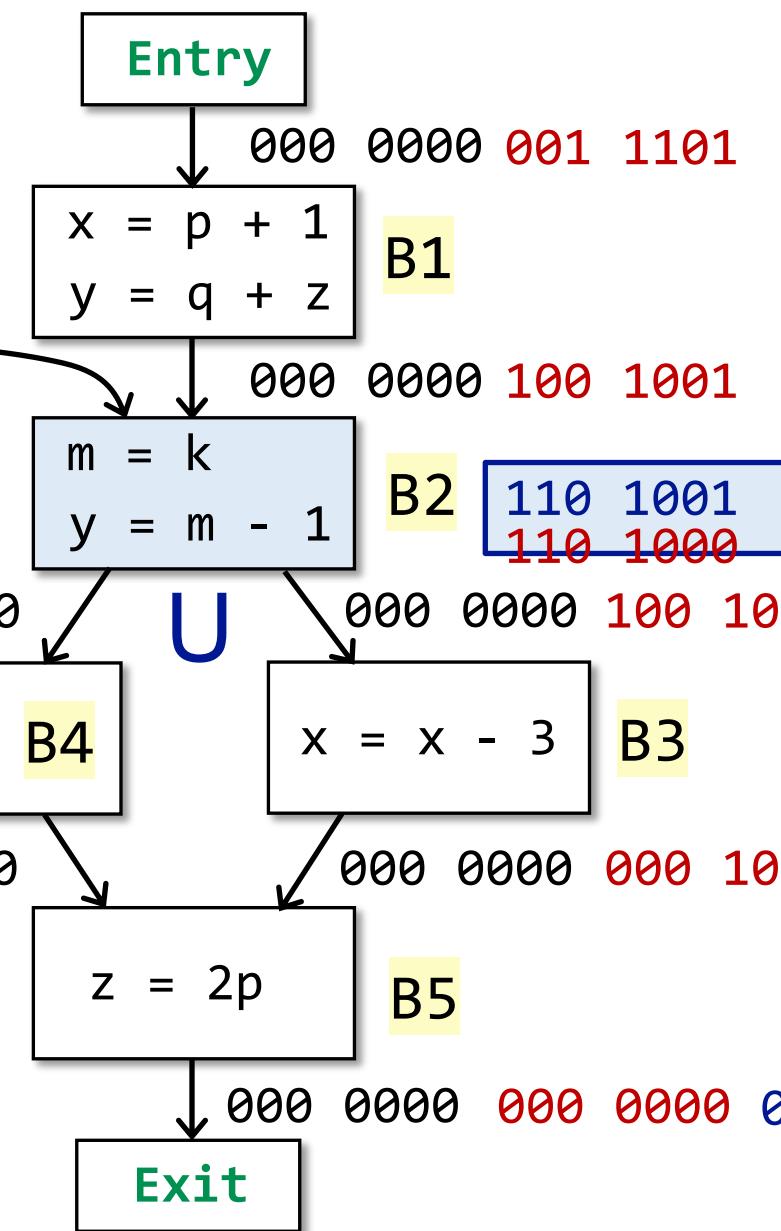
x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

Entry



010 1001	10 1000	0000 000											
100 1001													
000 1000	000 1000	0000 000											

x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

010 1001 010 1000 0000 000
 100 1001
 000 1000 000 1000 000 000

Entry

000 0000 001 1101

x = p + 1
 y = q + z

B1

m = k
 y = m - 1

B2

110 1001
 110 1000

x = 4
 q = y

B4

x = x - 3

B3

z = 2p

B5

Yue Li @ Nanjing University

x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

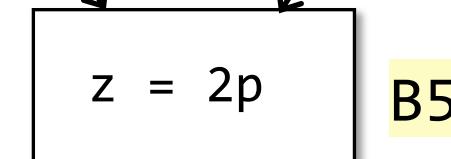
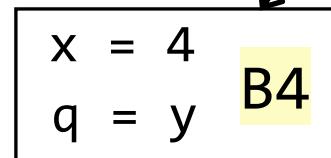
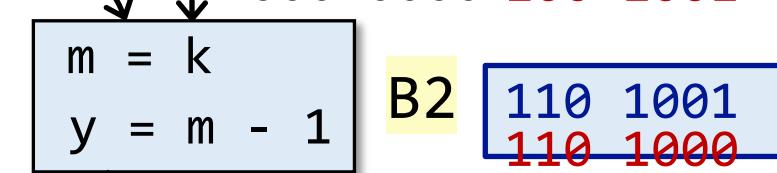
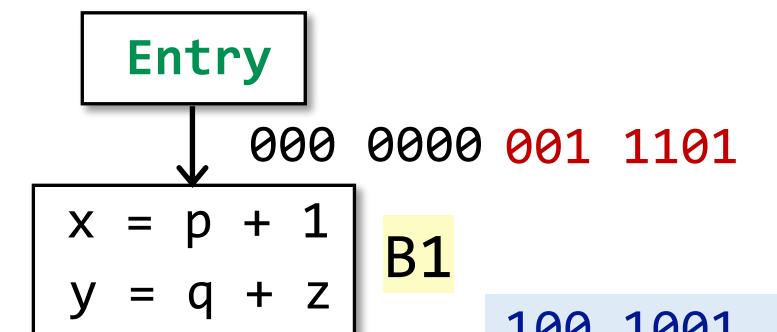
010 1001 010 1000 0000 000

100 1001

000 1000 000 1000 000 000 1000



Entry



Exit

x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

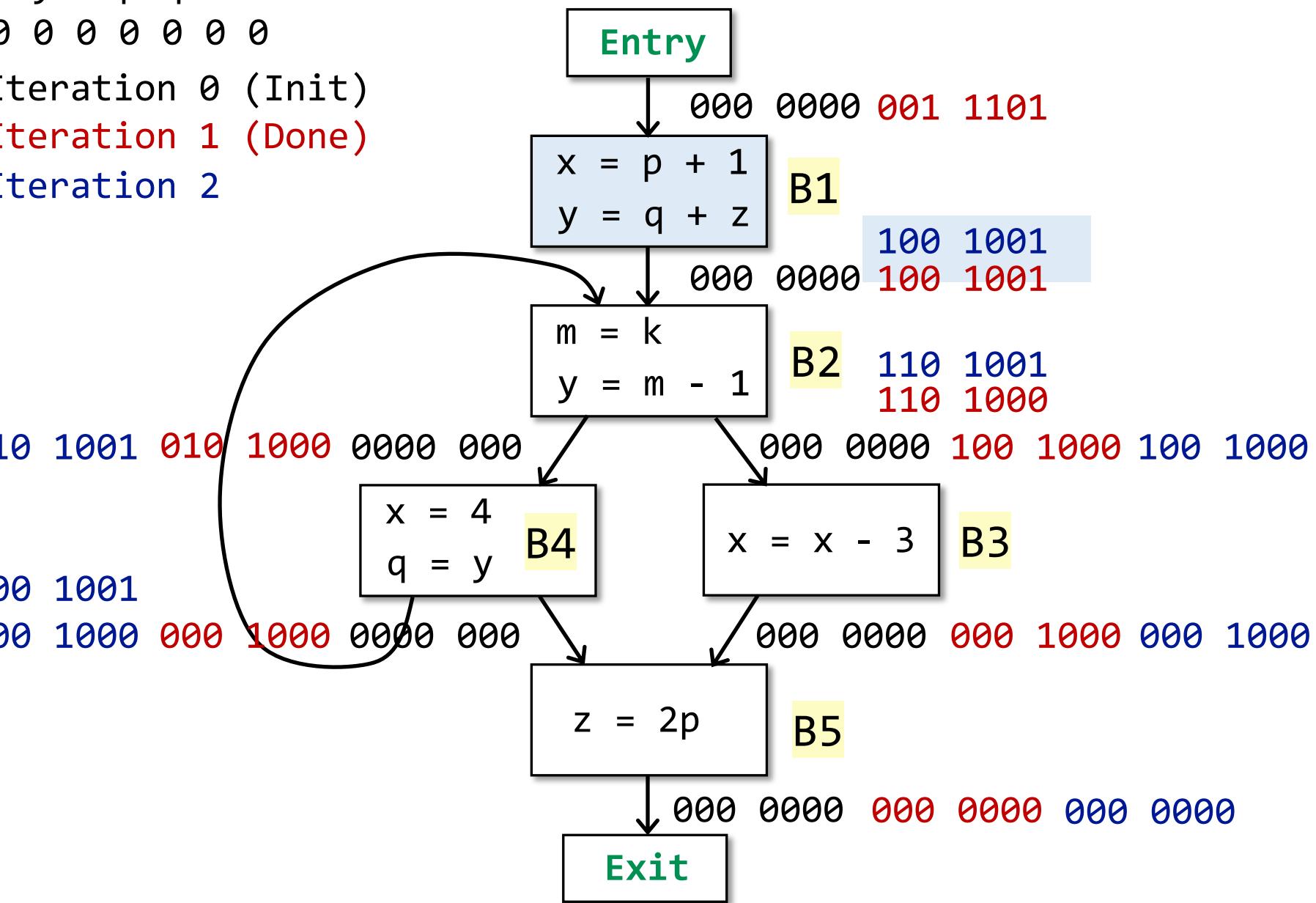
Iteration 1 (Done)

Iteration 2

010 1001 010 1000 0000 000

100 1001

000 1000 000 1000 0000 000



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

010 1001 010 1000 0000 000

100 1001
000 1000 000 1000 000 000

Entry

$x = p + 1$
 $y = q + z$

001 1101
001 1101

$m = k$
 $y = m - 1$

100 1001
100 1001

$x = 4$
 $q = y$

110 1001
110 1000

B1

B2

B3

$x = x - 3$
 $z = 2p$

100 1000 100 1000 100 1000

B4

B5

Exit

x y z p q m k
 0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Entry

$x = p + 1$
 $y = q + z$

B1

$m = k$
 $y = m - 1$

B2

$x = 4$
 $q = y$

B4

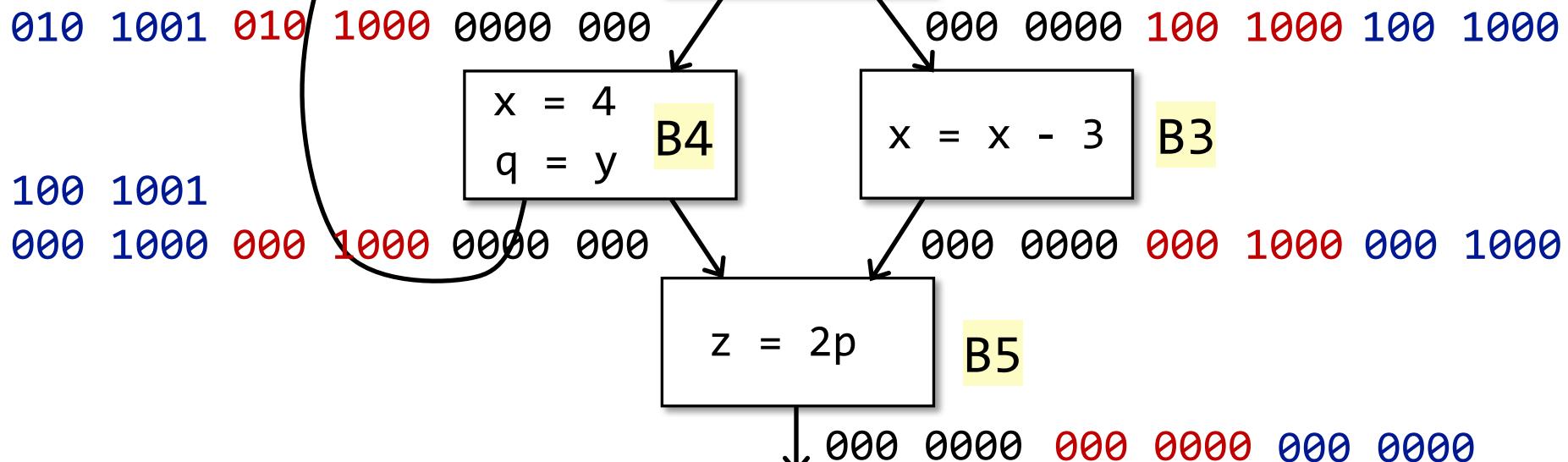
$x = x - 3$

B3

$z = 2p$

B5

Exit



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Entry

$x = p + 1$
 $y = q + z$

B1

$m = k$
 $y = m - 1$

B2

$x = 4$
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

Changes occur in
IN[] of B4

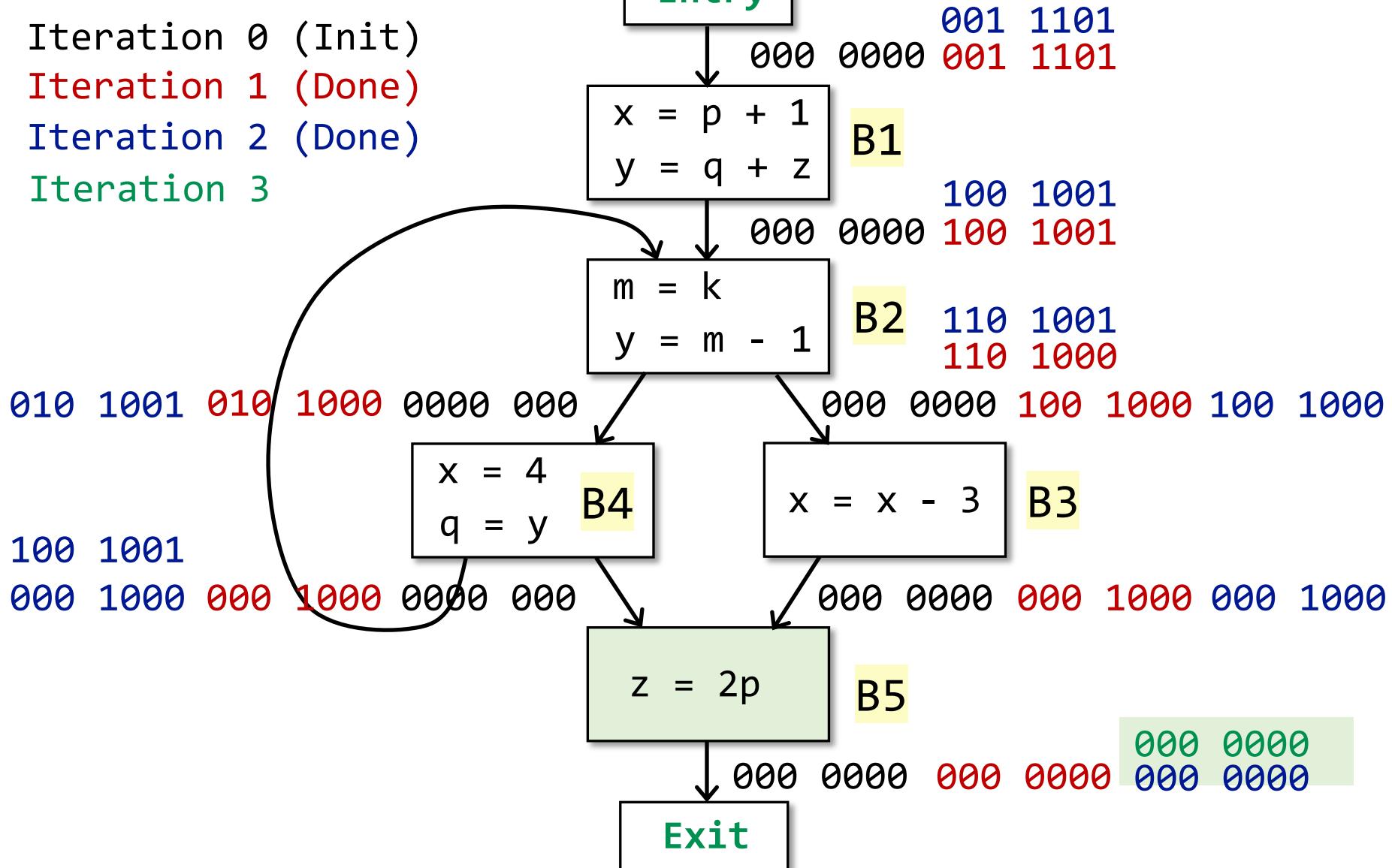
x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001 010 1000 0000 000
100 1001
000 1000 000 1000 0000 000
000 1000

Entry

000 0000

001 1101
001 1101

x = p + 1
y = q + z

B1

100 1001
100 1001

m = k
y = m - 1

B2

110 1001
110 1000

x = 4
q = y

B4

x = x - 3

B3

000 1000
000 1000

z = 2p

B5

000 0000
000 0000

Exit

x y z p q m k
0 0 0 0 0 0 0

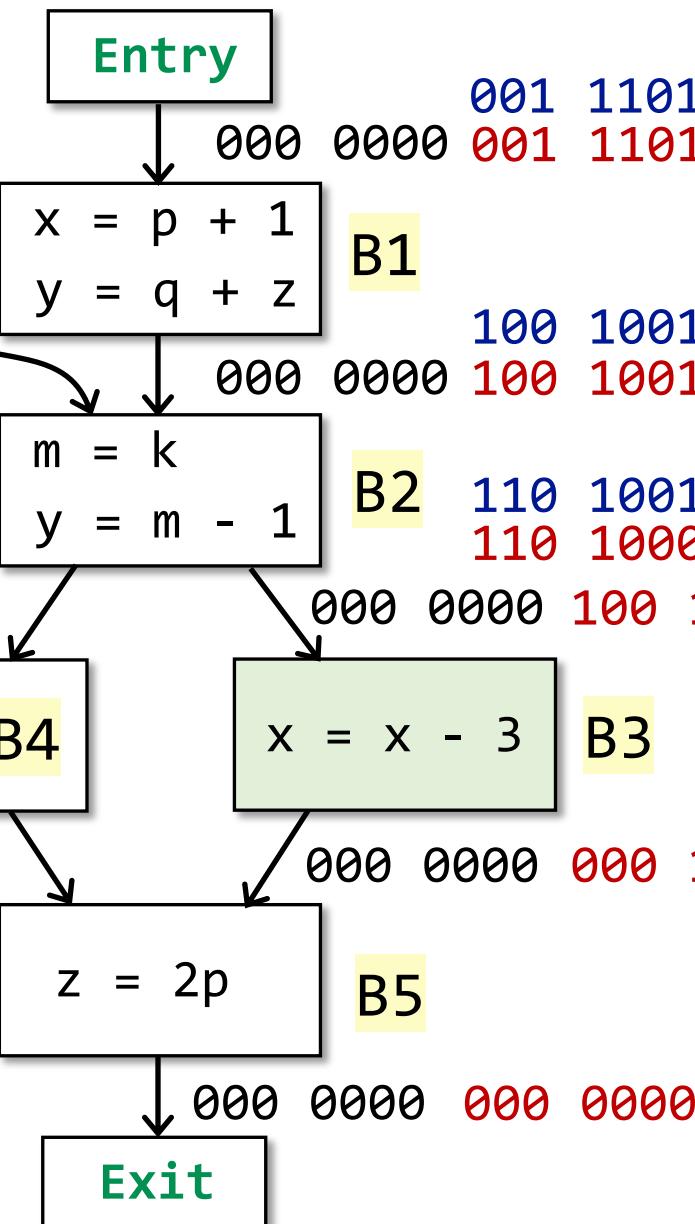
Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001 010 1000 0000 000
100 1001
000 1000 000 1000 0000 000
000 1000



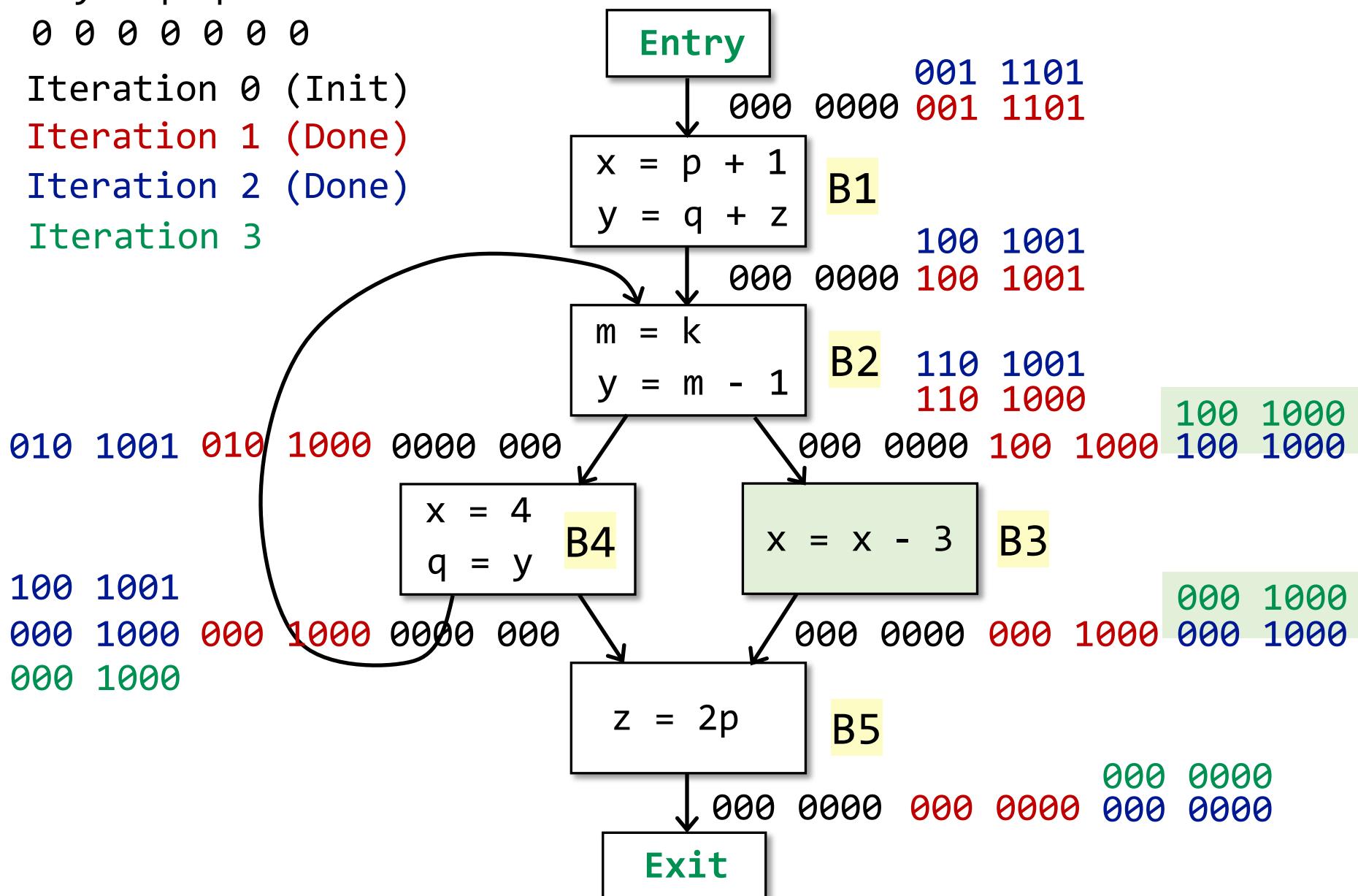
x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



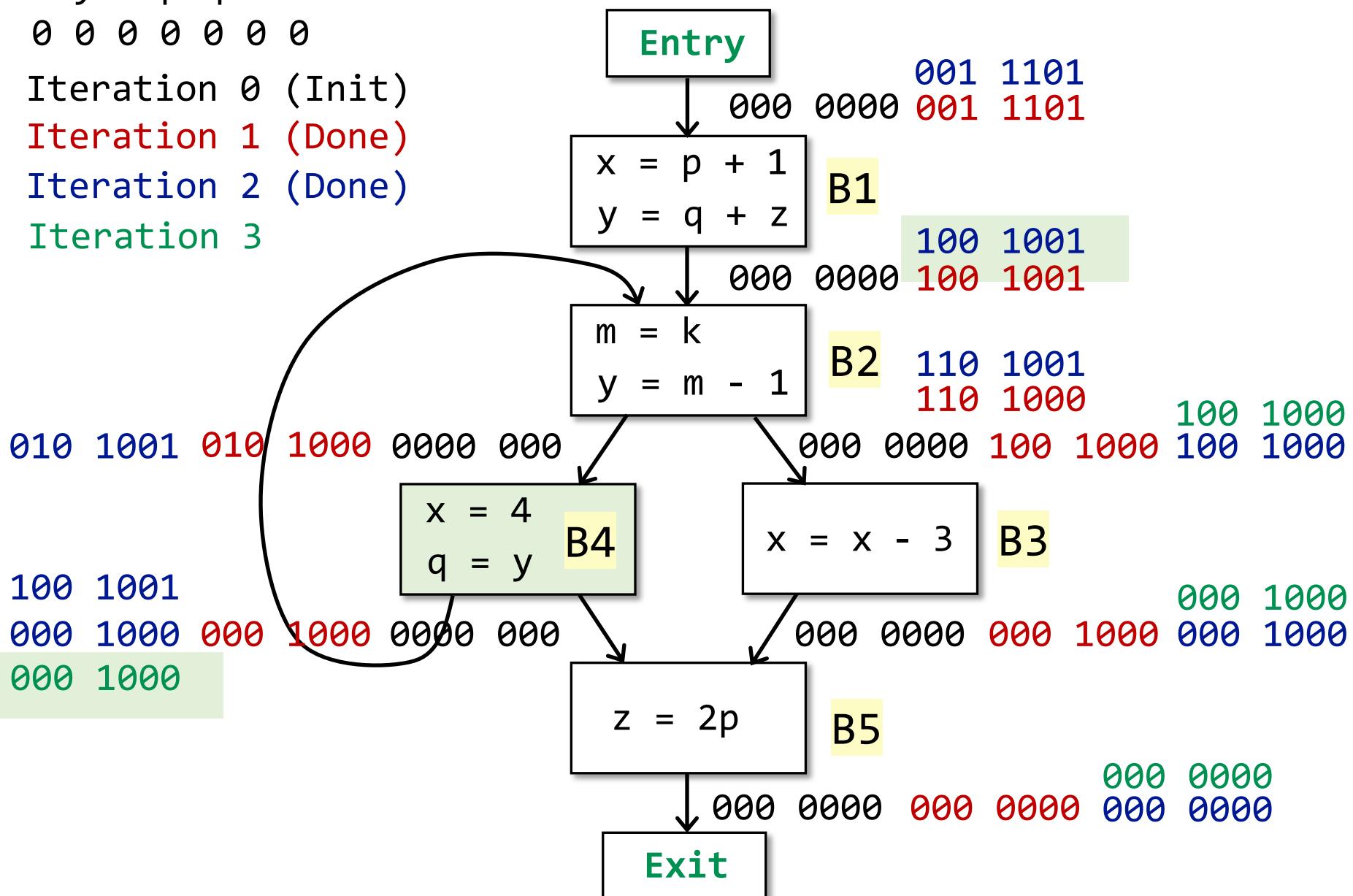
x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



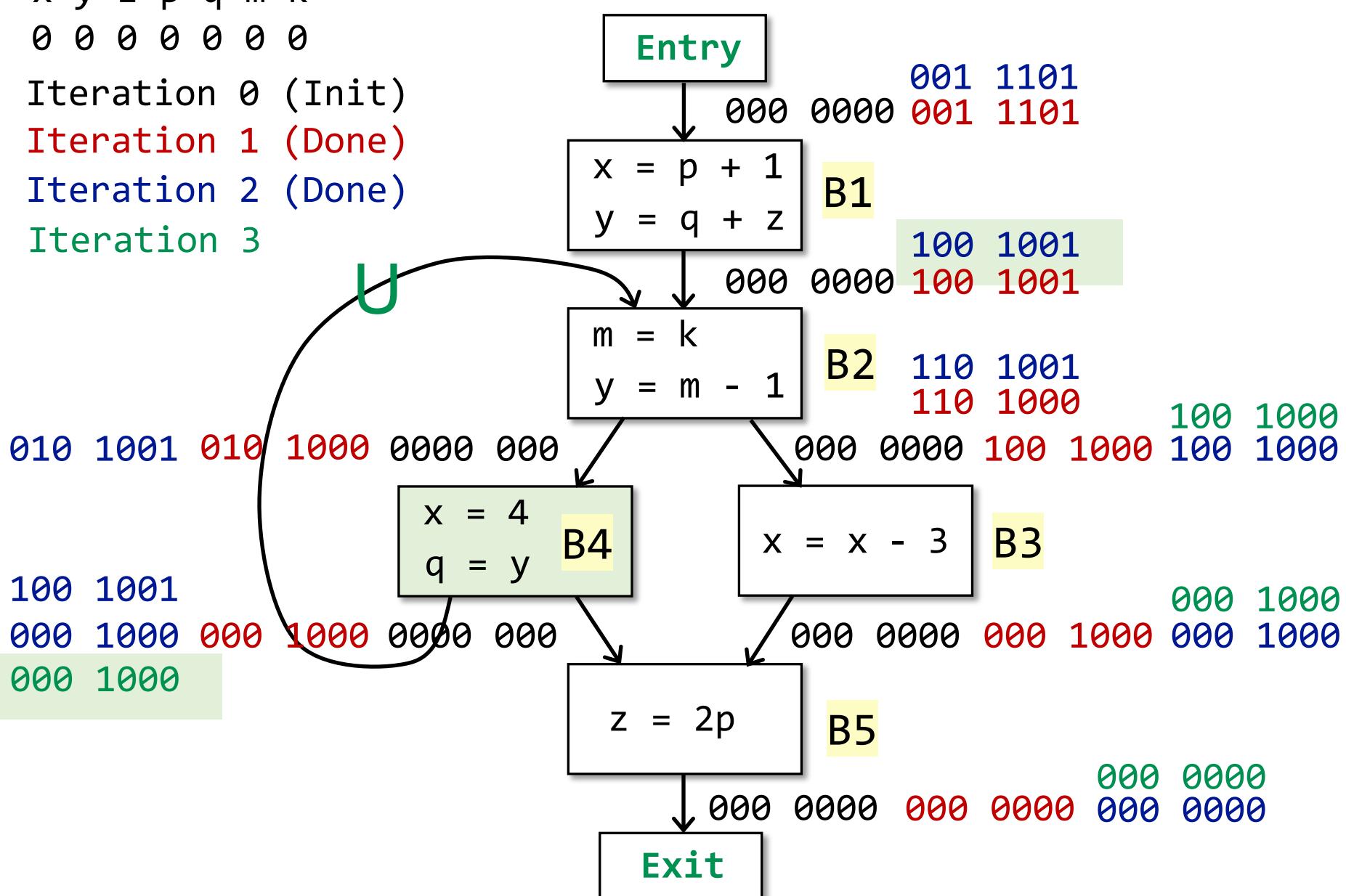
x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



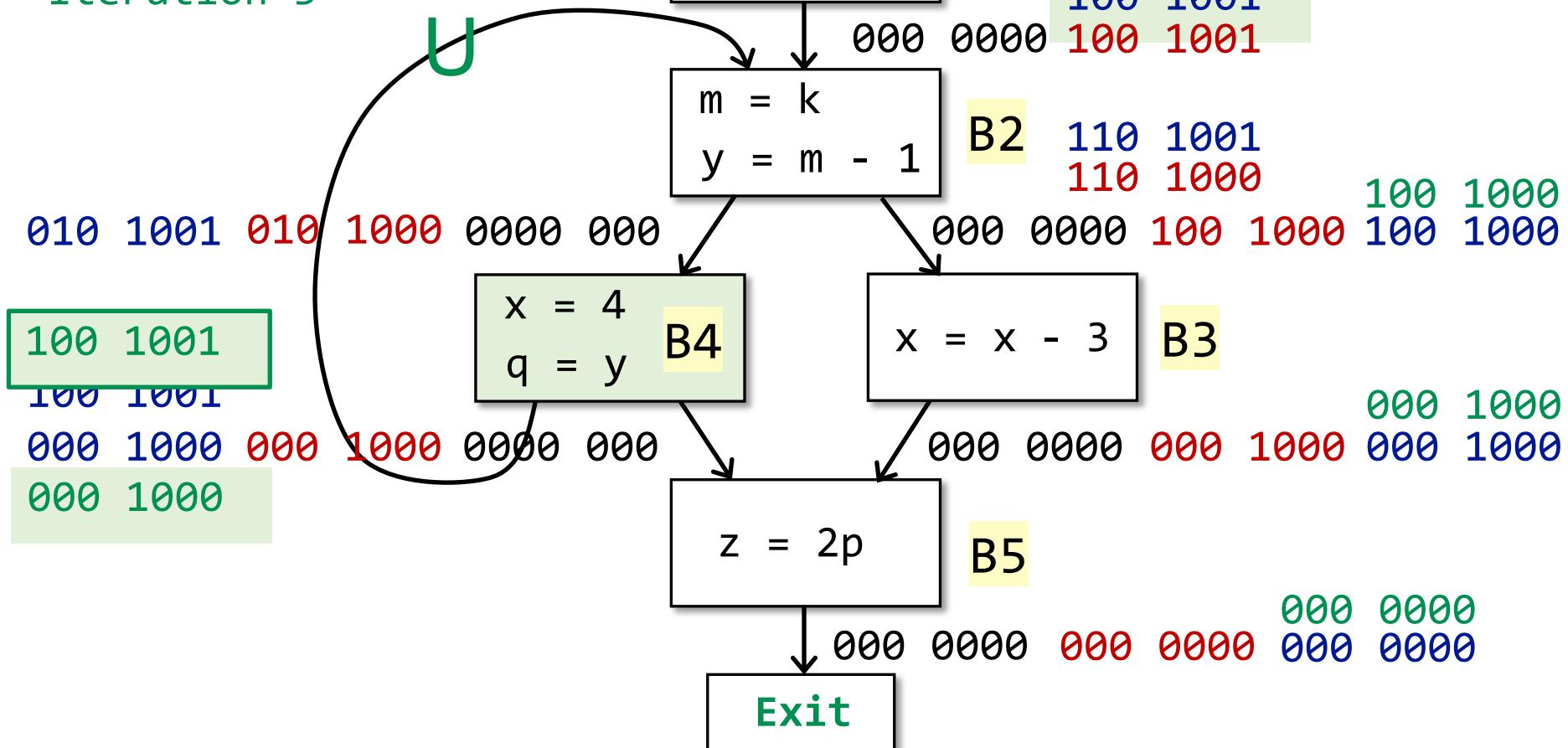
x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3



x y z p q m k
0 0 0 0 0 0 0

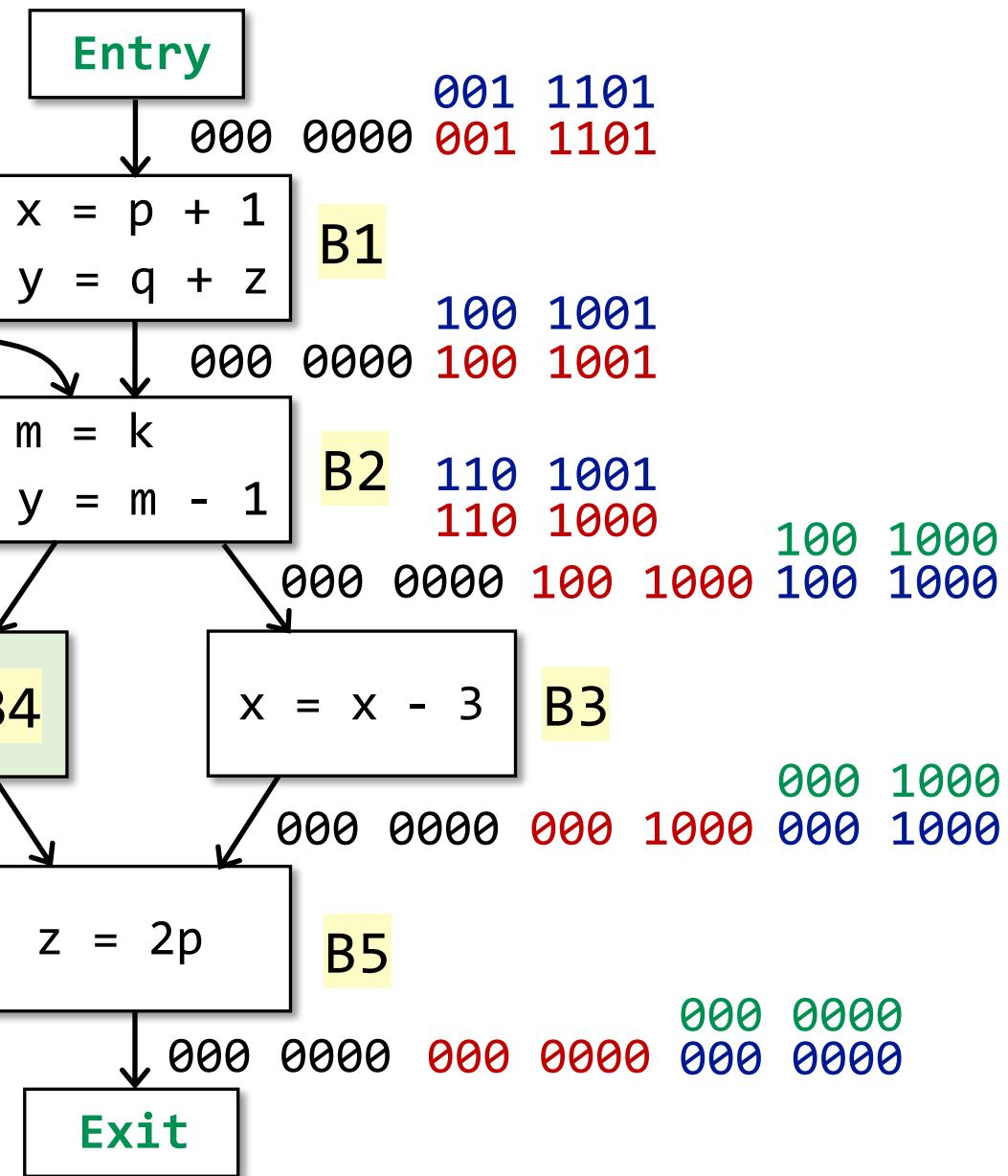
Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001 010	1000 0000 000	000 1000 000	000 1000 000	000 1000 000	000 1000 000	000 1000 000
100 1001		1000 0000 000	1000 0000 000	1000 0000 000	1000 0000 000	1000 0000 000
100 1001		000 1000 000	000 1000 000	000 1000 000	000 1000 000	000 1000 000
000 1000		000 1000 000	000 1000 000	000 1000 000	000 1000 000	000 1000 000



x y z p q m k
0 0 0 0 0 0 0

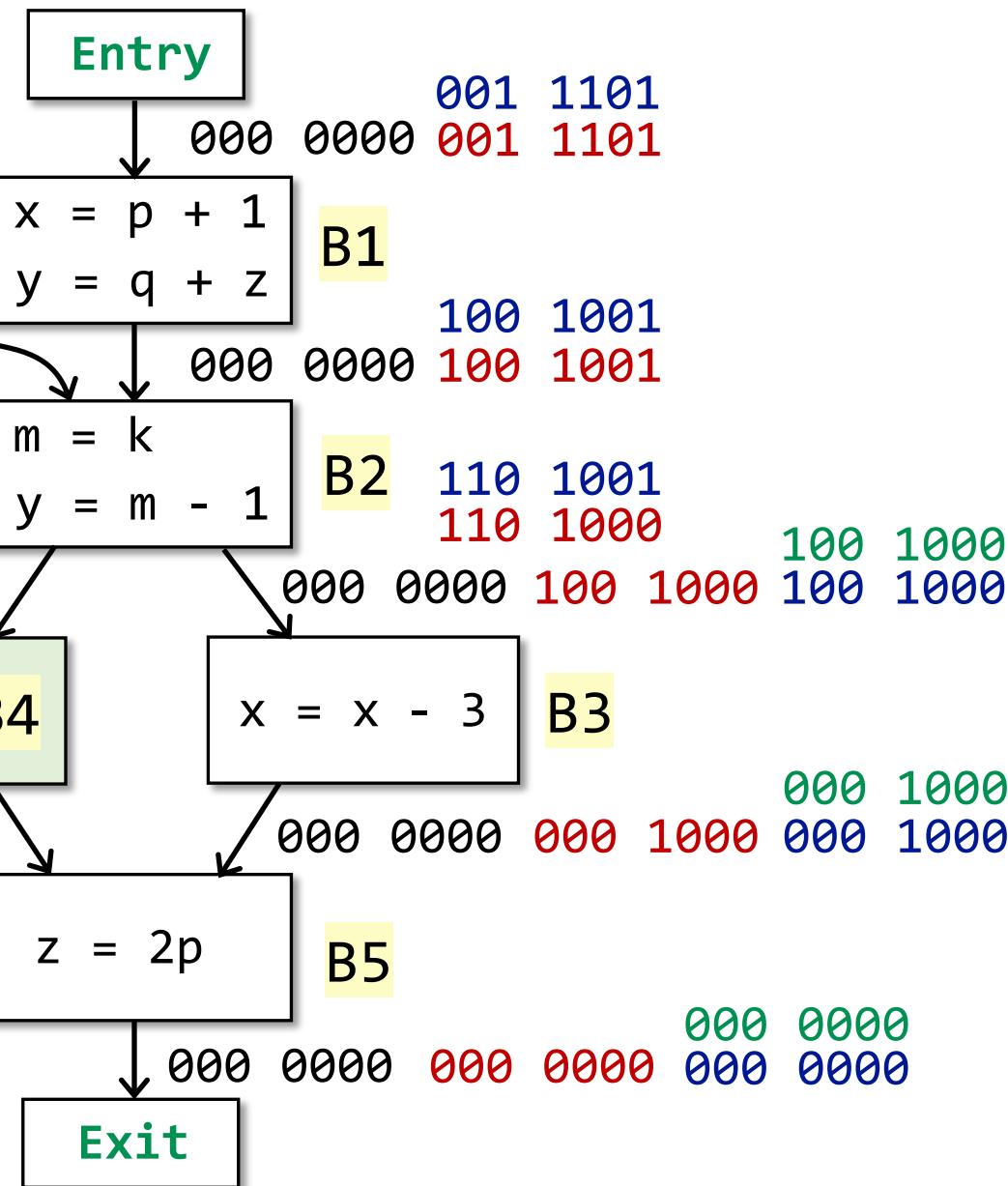
Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010	1001
010	1001
010	1001
100	1001
100	1001
000	1000
000	1000



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

Entry

$x = p + 1$
 $y = q + z$

B1

$m = k$
 $y = m - 1$

B2

$x = 4$
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

010 1001

010 1001

010 1000

100 1001

100 1001

000 1000

000 1000

001 1101
001 1101

100 1001
100 1001

110 1001
110 1000

100 1000
100 1000

000 1000
000 1000

000 0000
000 0000

x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

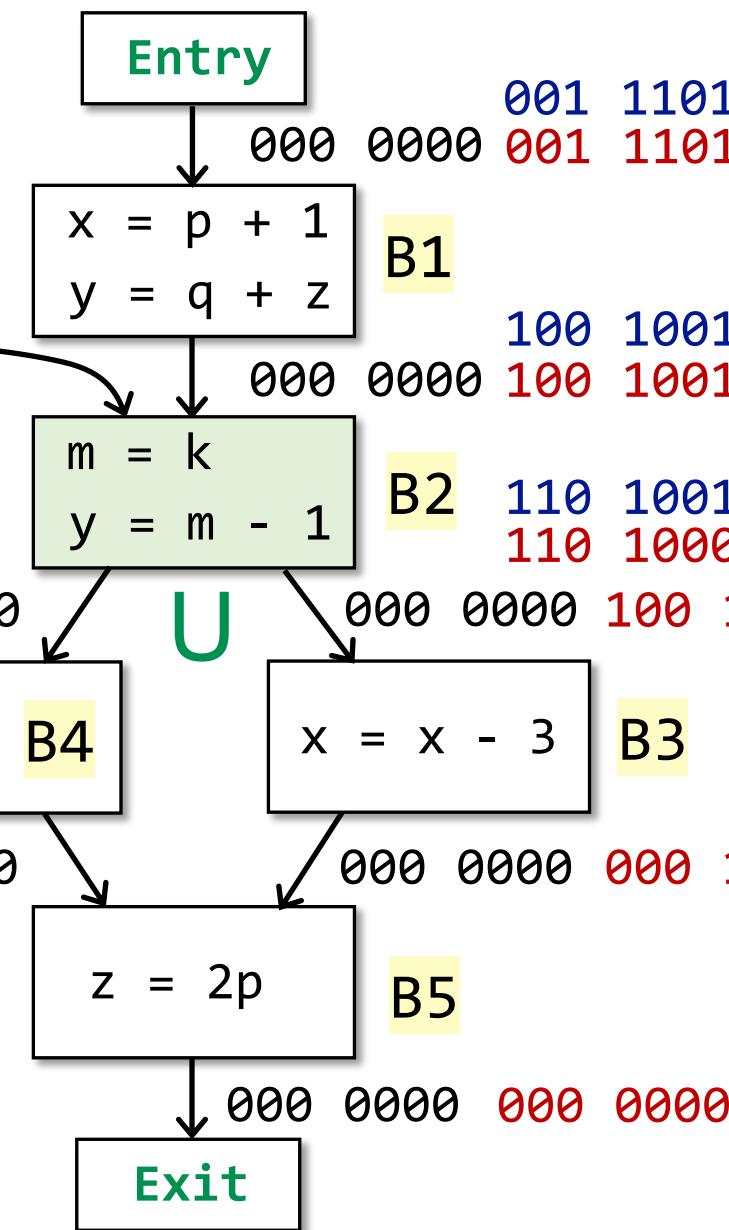
Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001	010 1001	000 0000	000 0000	001 1101	001 1101
010 1001	010 1001	010 1000	000 0000	100 1001	100 1001
100 1001	100 1001	1000 0000	000 0000	110 1001	110 1000
100 1001	100 1001	000 1000	000 0000	100 1000	100 1000
000 1000	000 1000	000 1000	000 0000	000 1000	000 1000
000 1000	000 1000	000 1000	000 0000	000 0000	000 0000

Entry



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

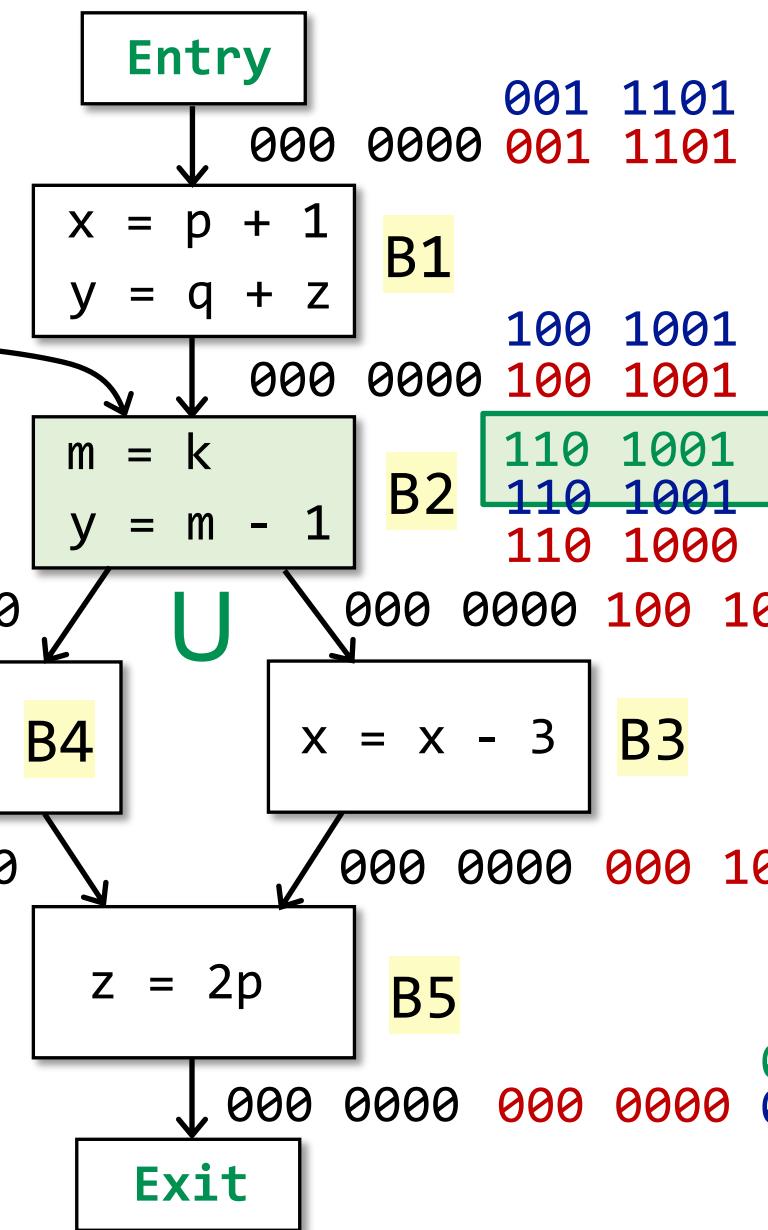
Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001	010 1001	010 1000	0000 000	000 000	000 000	000 000	000 000	000 000	000 000
010 1001	010 1001	010 1000	0000 000	000 000	000 000	000 000	000 000	000 000	000 000
100 1001	100 1001	100 1000	000 000	000 000	000 000	000 000	000 000	000 000	000 000
100 1001	100 1001	100 1000	000 000	000 000	000 000	000 000	000 000	000 000	000 000
000 1000	000 1000	000 1000	000 000	000 000	000 000	000 000	000 000	000 000	000 000
000 1000	000 1000	000 1000	000 000	000 000	000 000	000 000	000 000	000 000	000 000

Entry



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

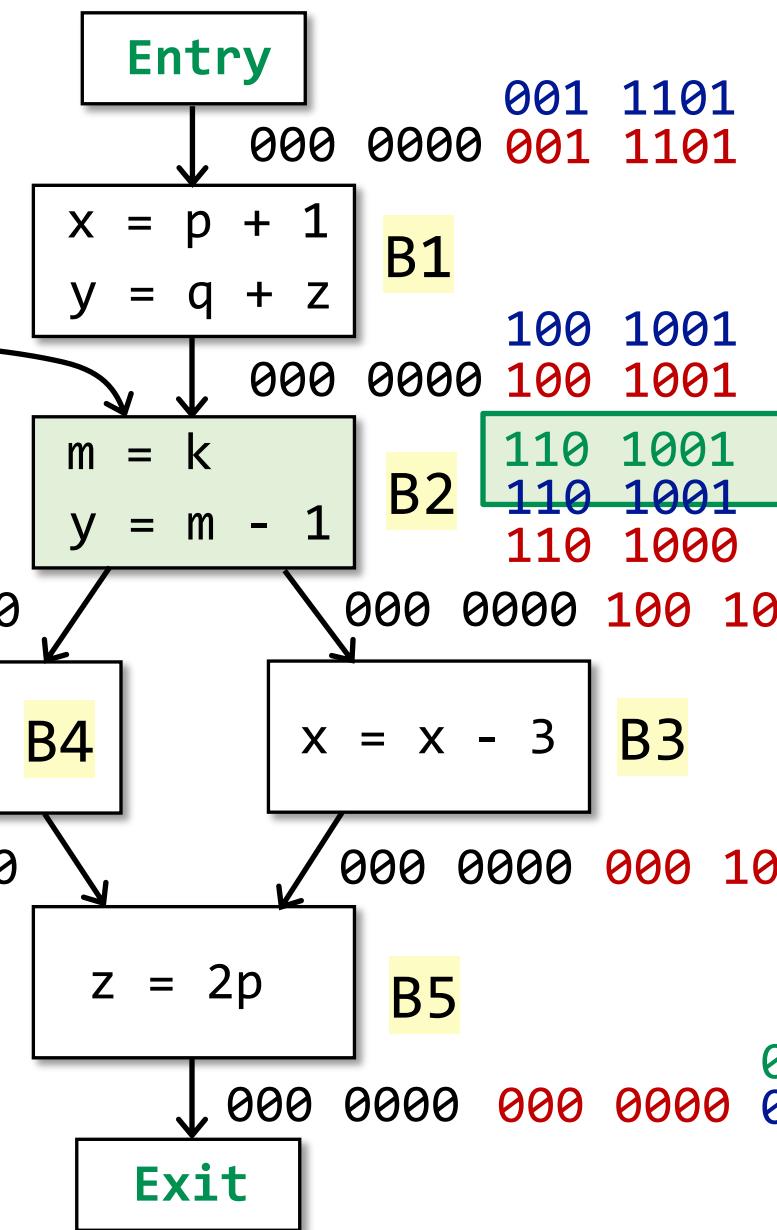
Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001
010 1001 010 1000 0000 000
100 1001
100 1001
000 1000 000 1000 0000 000
000 1000

Entry



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001

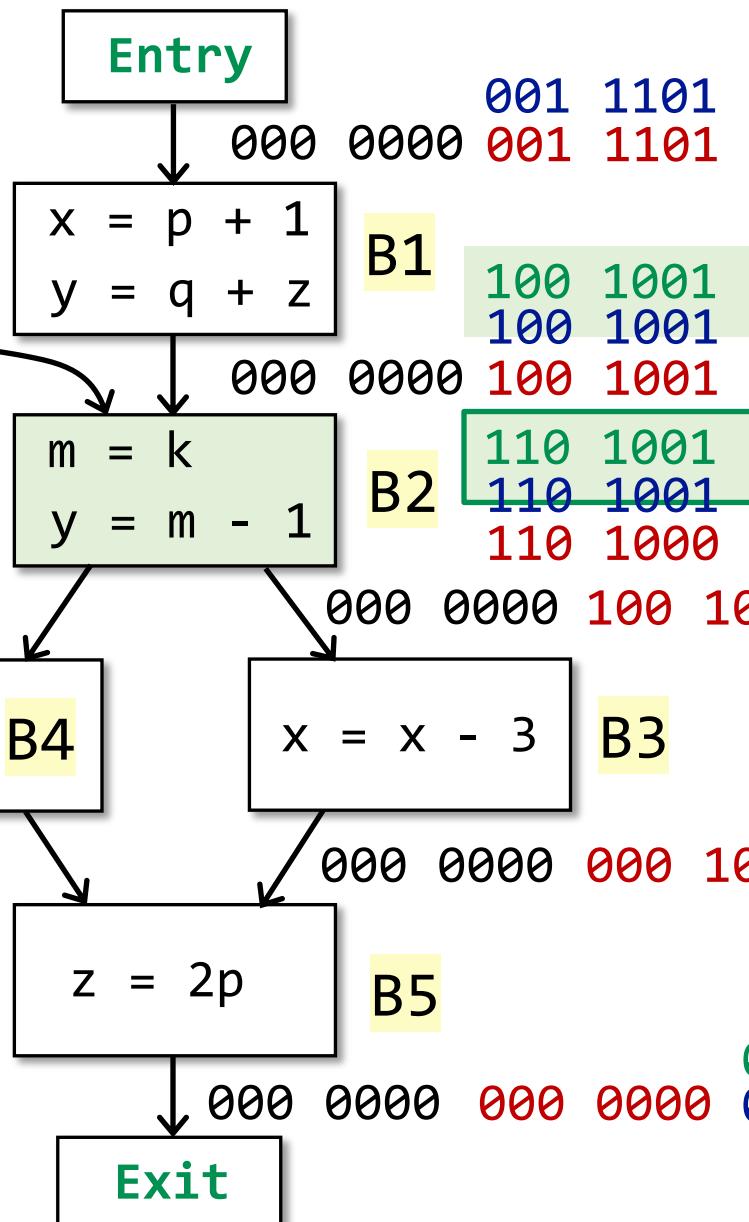
010 1001 010 1000 0000 000

100 1001

100 1001

000 1000 000 1000 000 000

000 1000



x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001

010 1001 010 1000 0000 000

100 1001

100 1001

000 1000

000 1000

Entry

$x = p + 1$
 $y = q + z$

B1

001 1101
001 1101

100 1001
100 1001
100 1001

$m = k$
 $y = m - 1$

B2

110 1001
110 1001
110 1000

$x = 4$
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

x	y	z	p	q	m	k
0	0	0	0	0	0	0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3

010 1001

010 1001

010 1000 0000 000

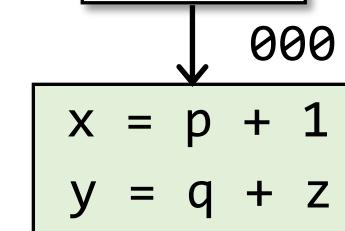
100 1001

100 1001

000 1000

000 1000

Entry



B1

B2

B3

B4

Exit

001	1101
001	1101
001	1101

100	1001
100	1001
100	1001

110	1001
110	1001
110	1000

100	1000
100	1000

000	1000
000	1000

000	0000
000	0000

x y z p q m k

0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

Entry

$x = p + 1$

$y = q + z$

B1

$m = k$

$y = m - 1$

B2

$x = 4$

$q = y$

B4

$x = x - 3$

B3

$z = 2p$

Exit

001 1101

001 1101

001 1101

100 1001

100 1001

100 1001

110 1001

110 1001

110 1000

100 1000

100 1000

100 1000

000 1000

000 1000

000 1000

000 0000

000 0000

000 0000

x y z p q m k

0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

010 1001

010 1001

010 1000

0000 000

100 1001

100 1001

000 1000

000 1000

0000 000

000 1000

Entry

$x = p + 1$
 $y = q + z$

B1

$m = k$
 $y = m - 1$

B2

$x = 4$
 $q = y$

B4

$x = x - 3$

B3

$z = 2p$

B5

Exit

001 1101

001 1101

001 1101

100 1001

100 1001

100 1001

110 1001

110 1001

110 1000

100 1000

100 1000

000 1000

000 1000

000 0000

000 0000

000 0000

No changes occur
in any IN[]

x y z p q m k
0 0 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

Iteration 3 (Done)

010 1001

010 1001

100 1001

100 1001

000 1000

000 1000

010 1000

1000 0000 000

000 1000 000

1000 0000 000

000 0000 000

000 0000 000

000 0000 000

000 0000 000

001 1101
001 1101
001 1101

100 1001
100 1001
100 1001

110 1001
110 1001
110 1000

100 1000
100 1000

000 1000
000 1000

000 0000
000 0000

Entry

000 0000

x = p + 1
y = q + z

B1

000 0000

m = k
y = m - 1

B2

x = 4
q = y

B4

x = x - 3

B3

z = 2p

B5

Exit

Final analysis result

Data Flow Analysis Applications

(I) Reaching Definitions Analysis

(II) Live Variables Analysis

(III) Available Expressions Analysis

Available Expressions Analysis

An expression $x \ op \ y$ is available at program point p if (1) **all** paths from the entry to p **must** pass through the evaluation of $x \ op \ y$, and (2) after the last evaluation of $x \ op \ y$, there is no redefinition of x or y

Available Expressions Analysis

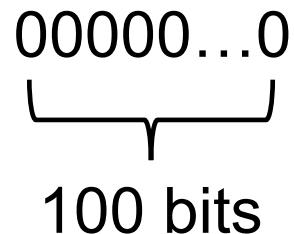
An expression $x \ op \ y$ is available at program point p if (1) **all** paths from the entry to p **must** pass through the evaluation of $x \ op \ y$, and (2) after the last evaluation of $x \ op \ y$, there is no redefinition of x or y

- This definition means at program p, we can replace expression $x \ op \ y$ by the result of its last evaluation
- The information of available expressions can be used for detecting global common subexpressions.

Understanding Available Expressions Analysis

- Data Flow Values/Facts
 - All the expressions in a program
 - Can be represented by bit vectors
- e.g., E1, E2, E3, E4, ..., E100 (100 expressions)

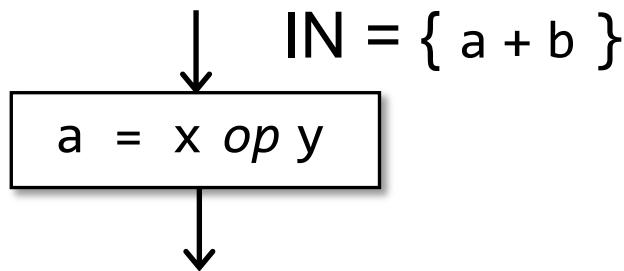
Abstraction



Bit i from the left represents expression E_i

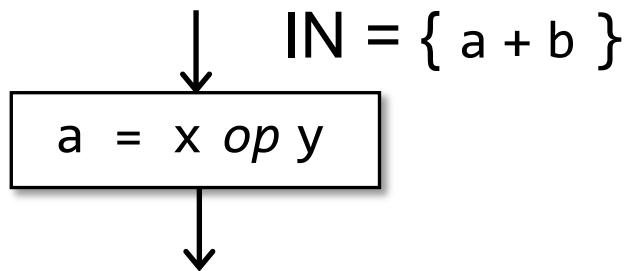
Understanding Available Expressions Analysis

Safe-approximation



Understanding Available Expressions Analysis

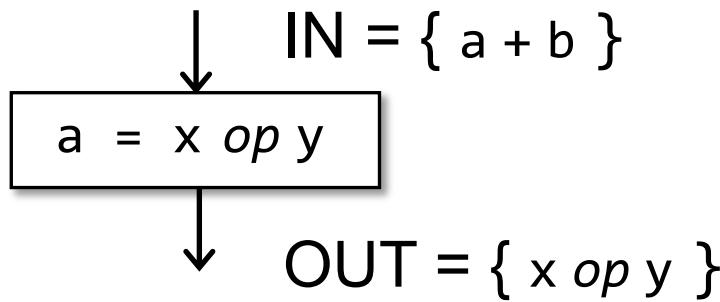
Safe-approximation



- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

Understanding Available Expressions Analysis

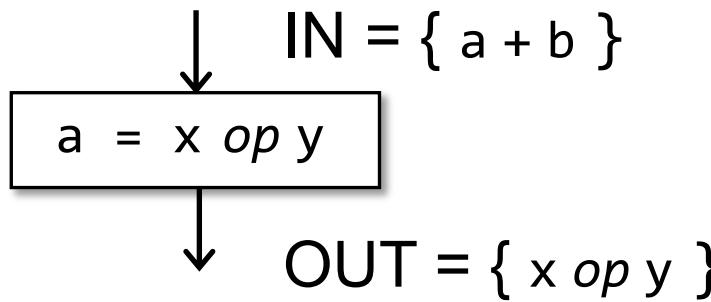
Safe-approximation



- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

Understanding Available Expressions Analysis

Safe-approximation



- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

Understanding Available Expressions Analysis

Safe-approximation

$$\downarrow \quad \text{IN} = \{ a + b \}$$

$$a = x \ op \ y$$

$$\downarrow \quad \text{OUT} = \{ x \ op \ y \ }$$

- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$

$$a = e^{16} * x$$

$$\begin{aligned} x &= \dots \\ b &= e^{16} * x \end{aligned}$$

$$c = e^{16} * x$$

Understanding Available Expressions Analysis

Safe-approximation

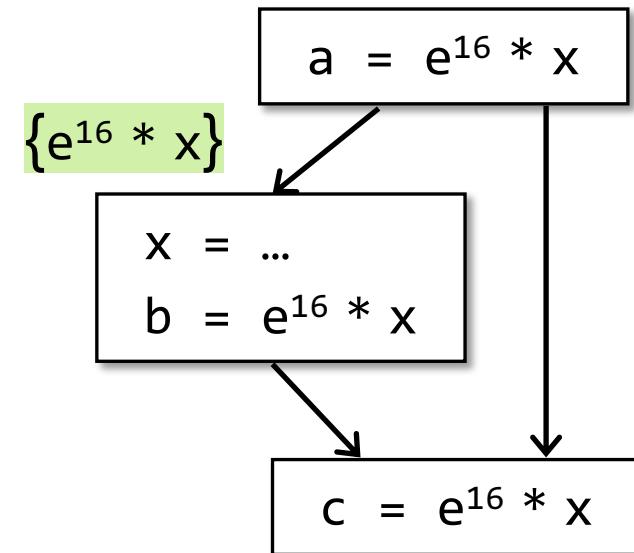
$$\downarrow \quad \text{IN} = \{ a + b \}$$

$$a = x \ op \ y$$

$$\downarrow \quad \text{OUT} = \{ x \ op \ y \ }$$

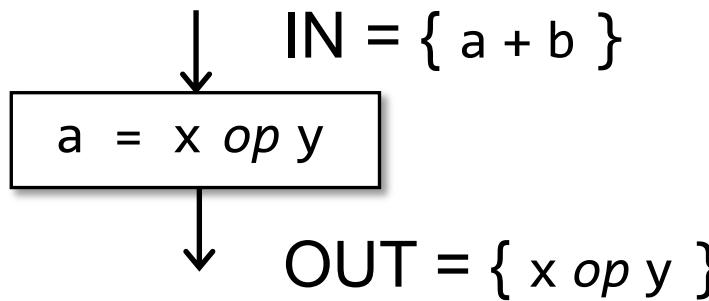
- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



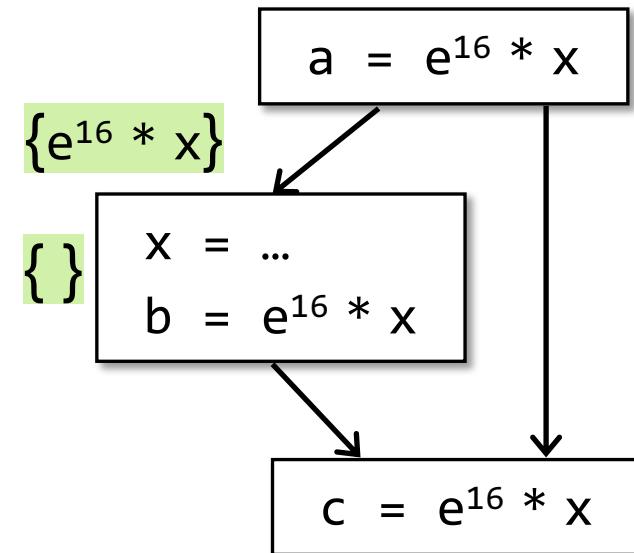
Understanding Available Expressions Analysis

Safe-approximation



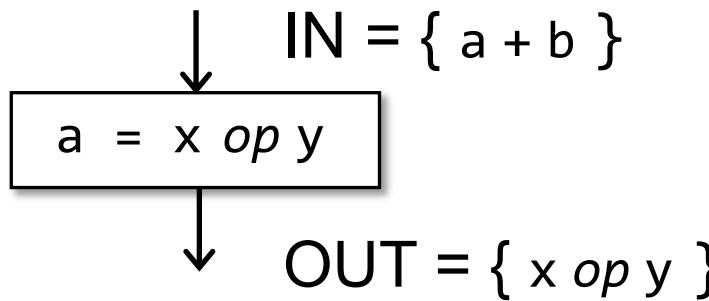
- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$



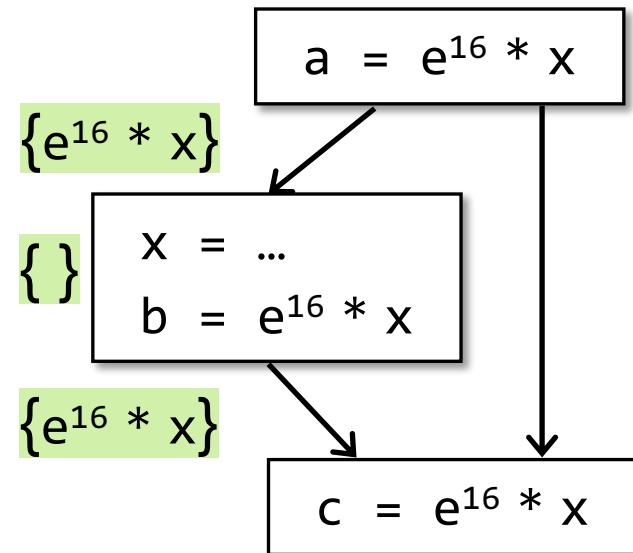
Understanding Available Expressions Analysis

Safe-approximation



- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$



Understanding Available Expressions Analysis

Safe-approximation

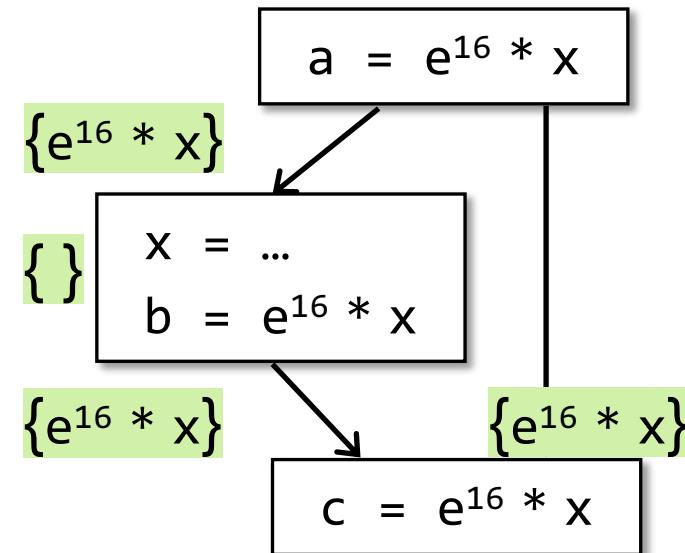
$$\downarrow \quad \text{IN} = \{ a + b \}$$

$$a = x \ op \ y$$

$$\downarrow \quad \text{OUT} = \{ x \ op \ y \ }$$

- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



Understanding Available Expressions Analysis

Safe-approximation

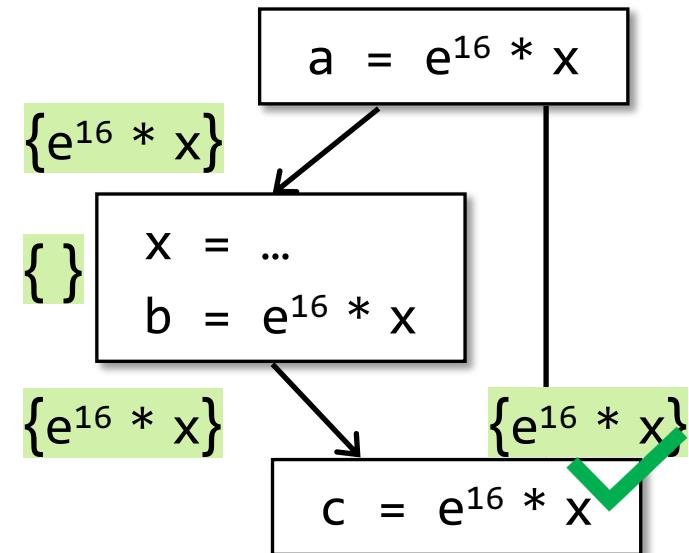
$$\downarrow \quad \text{IN} = \{ a + b \}$$

$$a = x \ op \ y$$

$$\downarrow \quad \text{OUT} = \{ x \ op \ y \ }$$

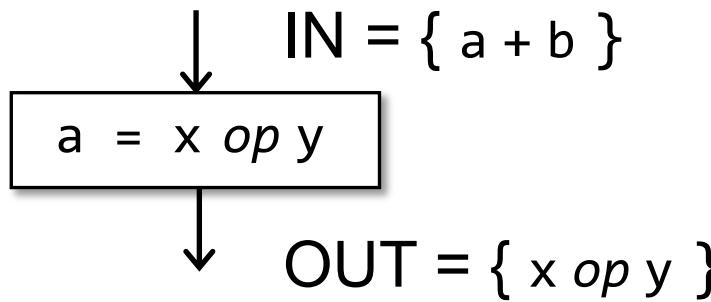
- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$\text{OUT}[B] = \text{gen}_B \cup (\text{IN}[B] - \text{kill}_B)$$



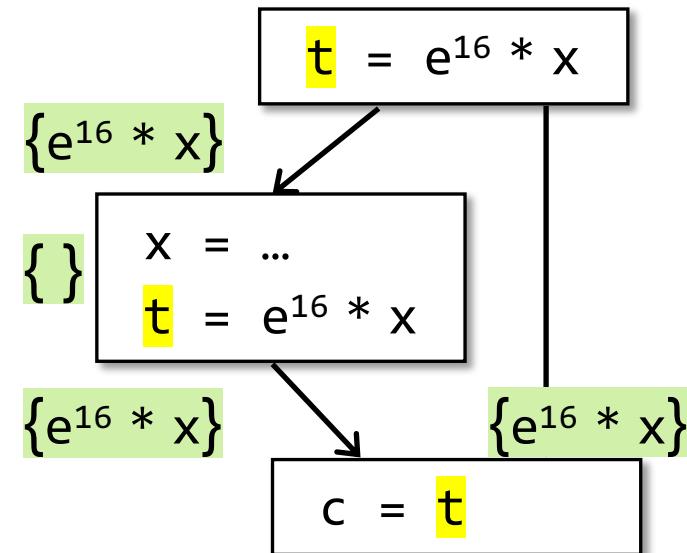
Understanding Available Expressions Analysis

Safe-approximation



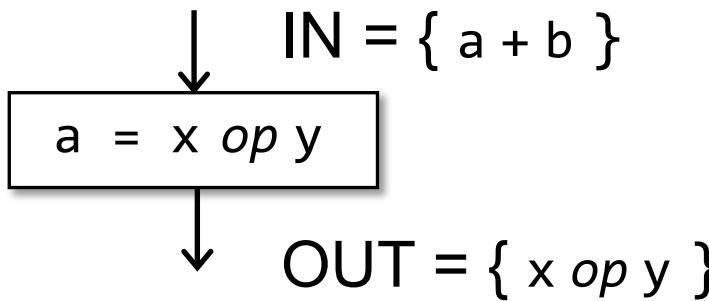
- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$



Understanding Available Expressions Analysis

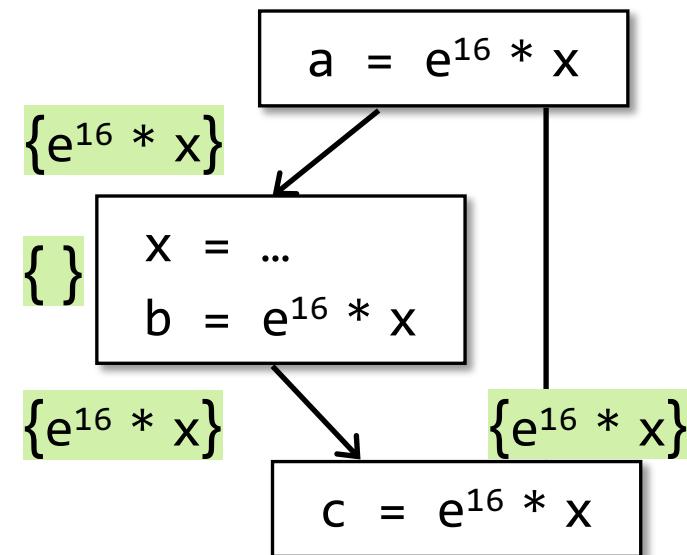
Safe-approximation



- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

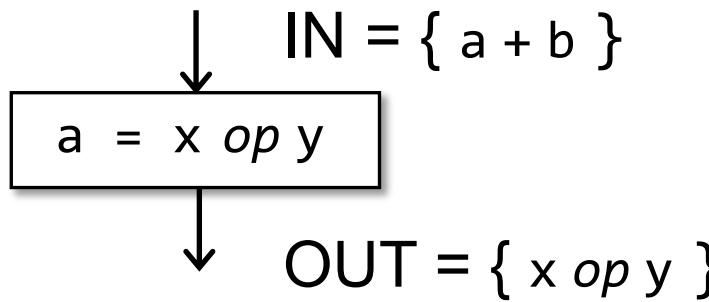
$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

$$IN[B] = \bigcap_{P \text{ a predecessor of } B} OUT[P]$$



Understanding Available Expressions Analysis

Safe-approximation

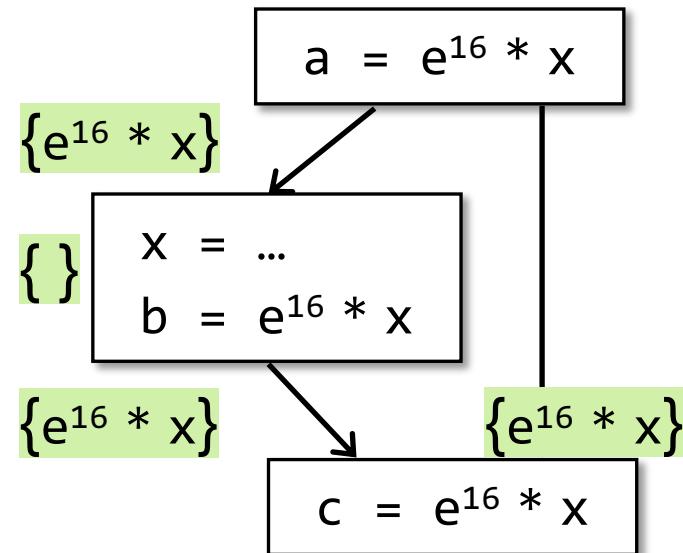


- Add to OUT the expression $x \ op \ y$ (gen)
- Delete from IN any expression involving variable a (kill)

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

$$IN[B] = \bigcap_{P \text{ a predecessor of } B} OUT[P]$$

All paths from entry to point p must pass through the evaluation of $x \ op \ y$



Understanding Available Expressions Analysis

$$\downarrow \quad \text{IN} = \{ a + b \}$$

Safe-approximation

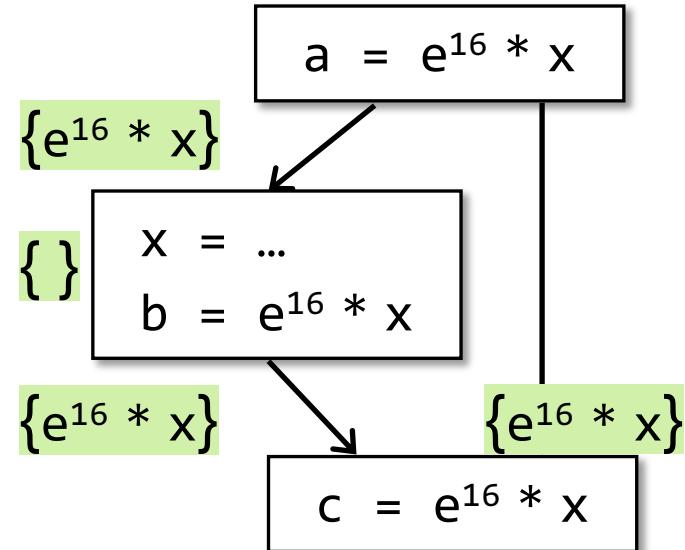
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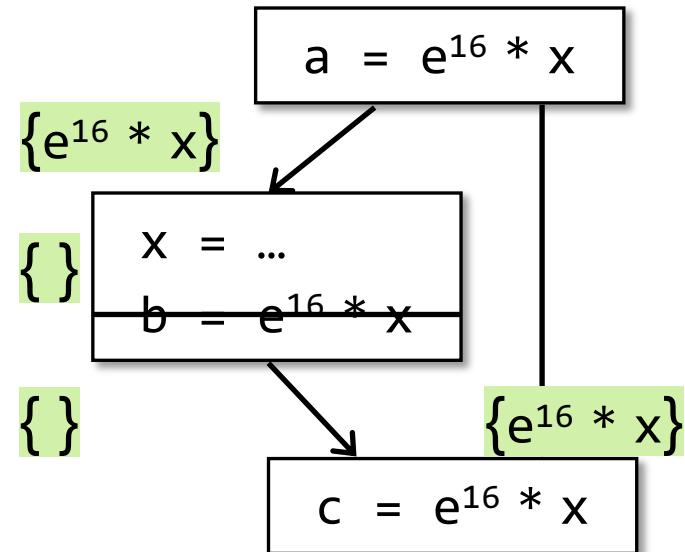
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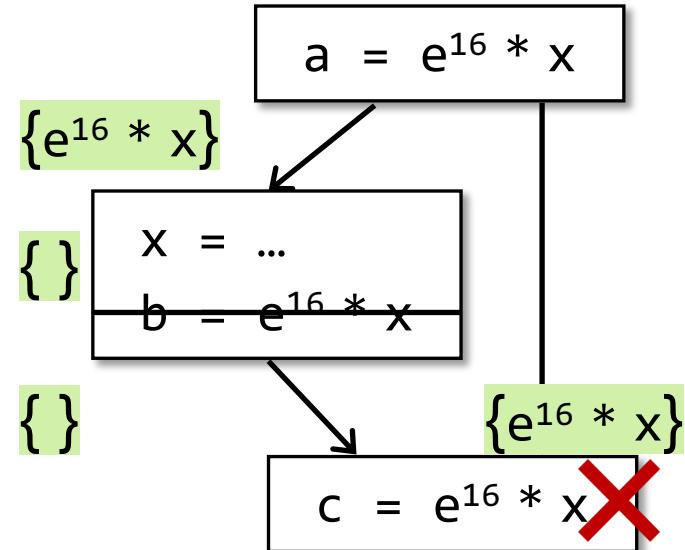
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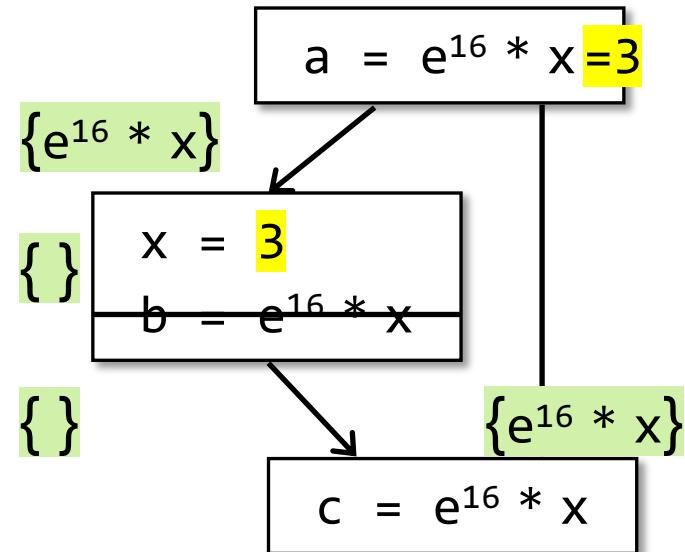
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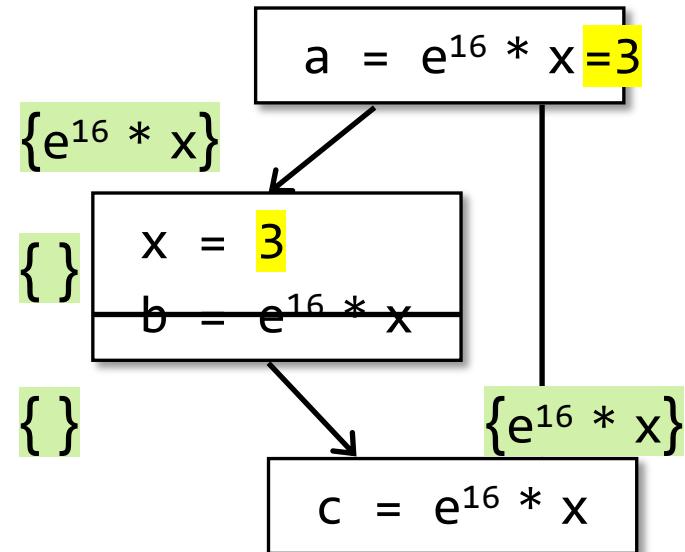
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Algorithm of Available Expressions Analysis

INPUT: CFG ($kill_B$ and gen_B computed for each basic block B)

OUTPUT: $IN[B]$ and $OUT[B]$ for each basic block B

METHOD:

```
OUT[entry] = Ø;  
for (each basic block  $B \setminus entry$ )  
    OUT[B] = U;  
    while (changes to any OUT occur)  
        for (each basic block  $B \setminus entry$ ) {  
            IN[B] =  $\bigcap_{P \text{ a predecessor of } B} OUT[P];$   
            OUT[B] =  $gen_B \cup (IN[B] - kill_B);$   
        }  
    }
```

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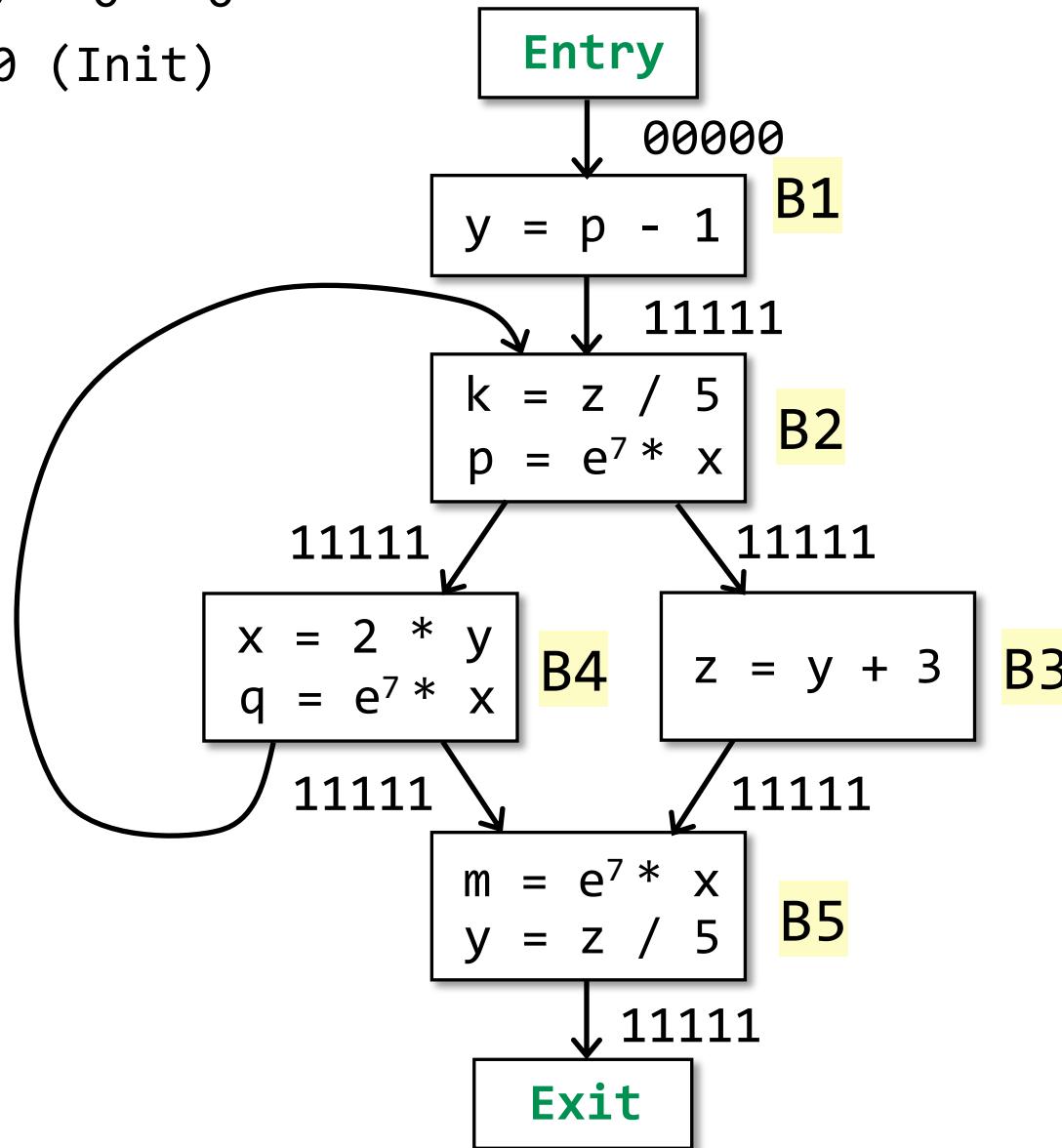
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```



p-1 z/5 2*y e⁷*x y+3
0 0 0 0 0

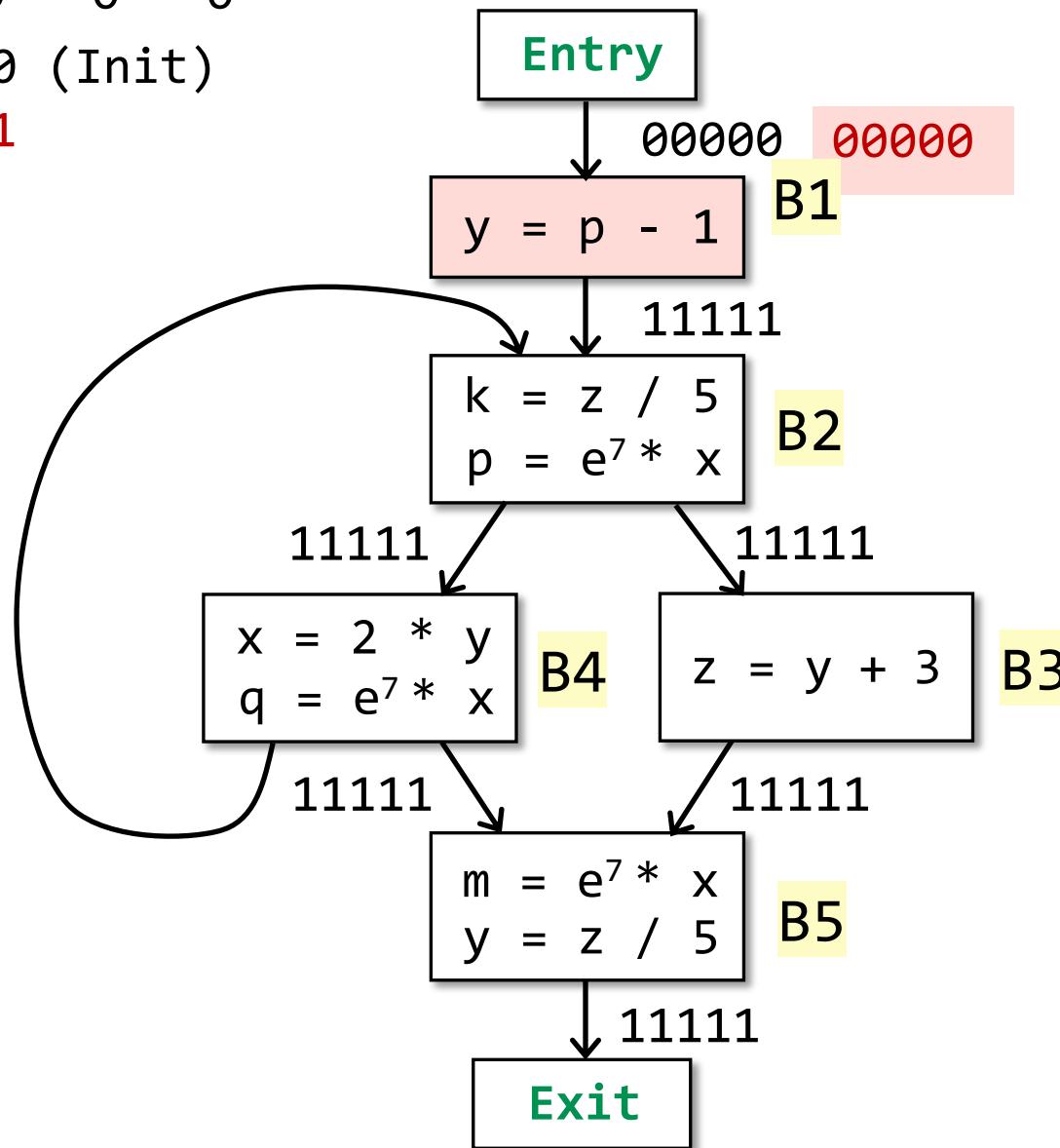
Iteration 0 (Init)



p-1 z/5 2*y e⁷*x y+3
0 0 0 0 0

Iteration 0 (Init)

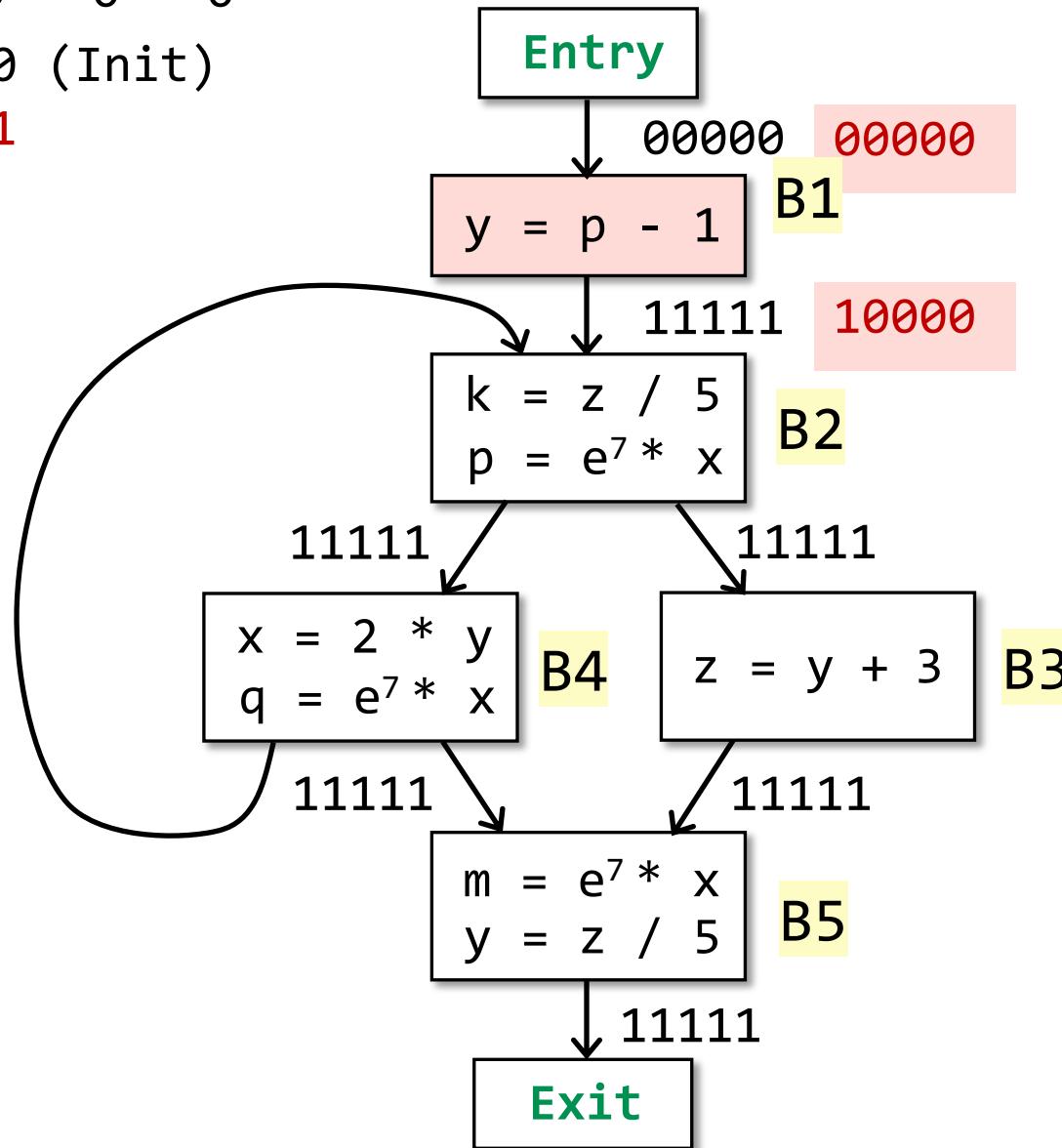
Iteration 1



$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
0 0 0 0 0

Iteration 0 (Init)

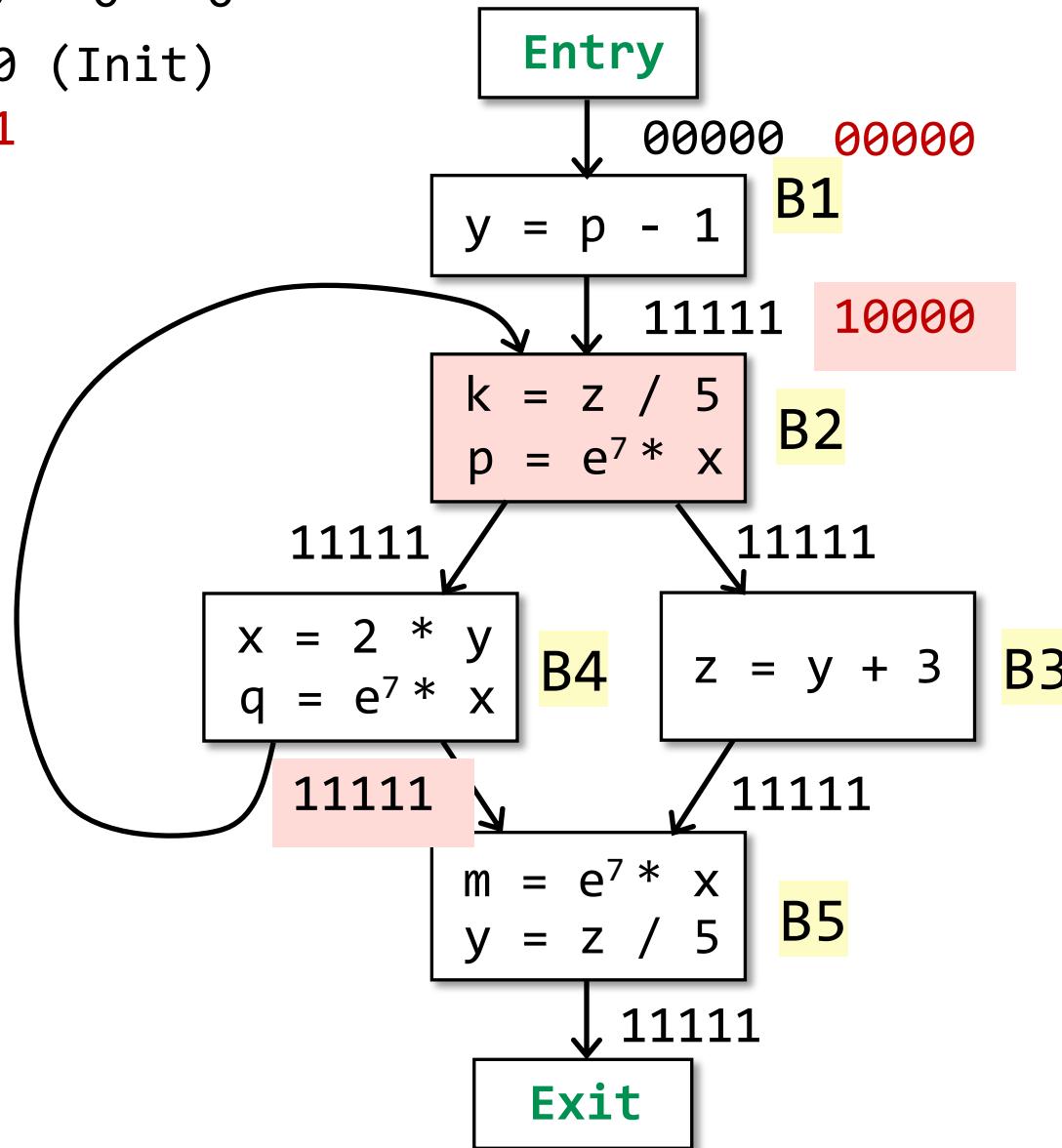
Iteration 1



p-1 z/5 2*y e⁷*x y+3
0 0 0 0 0

Iteration 0 (Init)

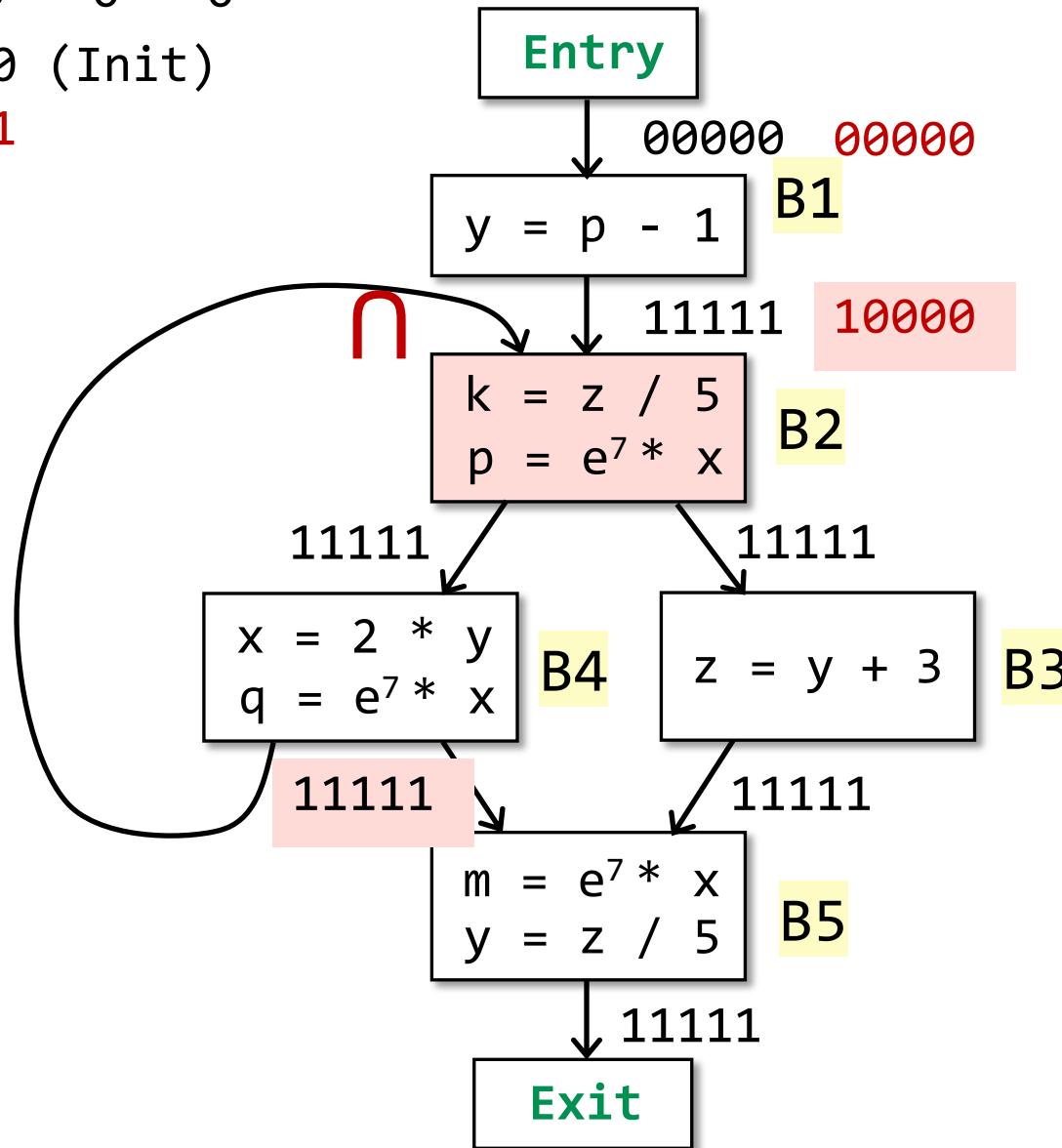
Iteration 1



p-1 z/5 2*y e⁷*x y+3
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Iteration 0 (Init)

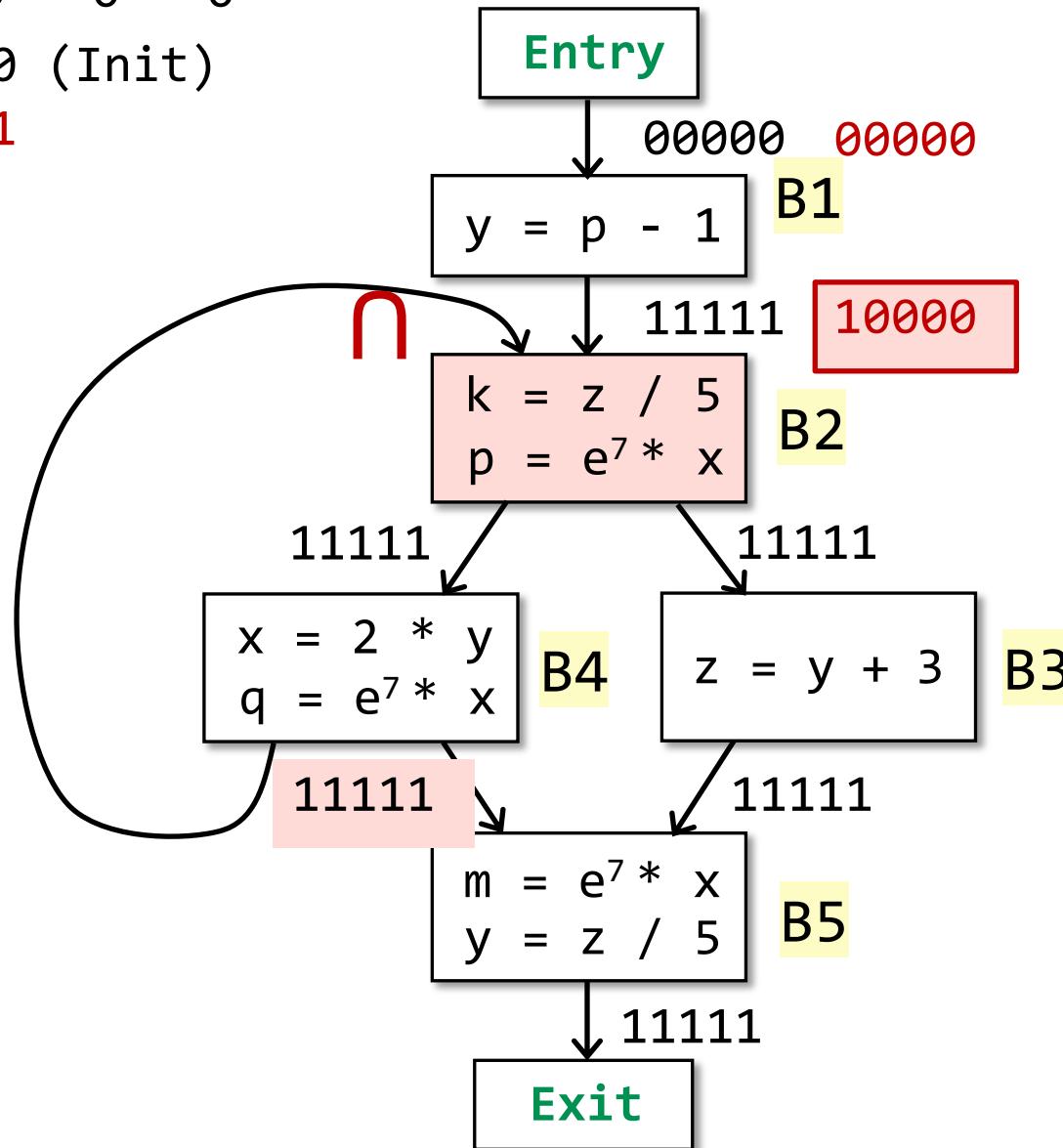
Iteration 1



$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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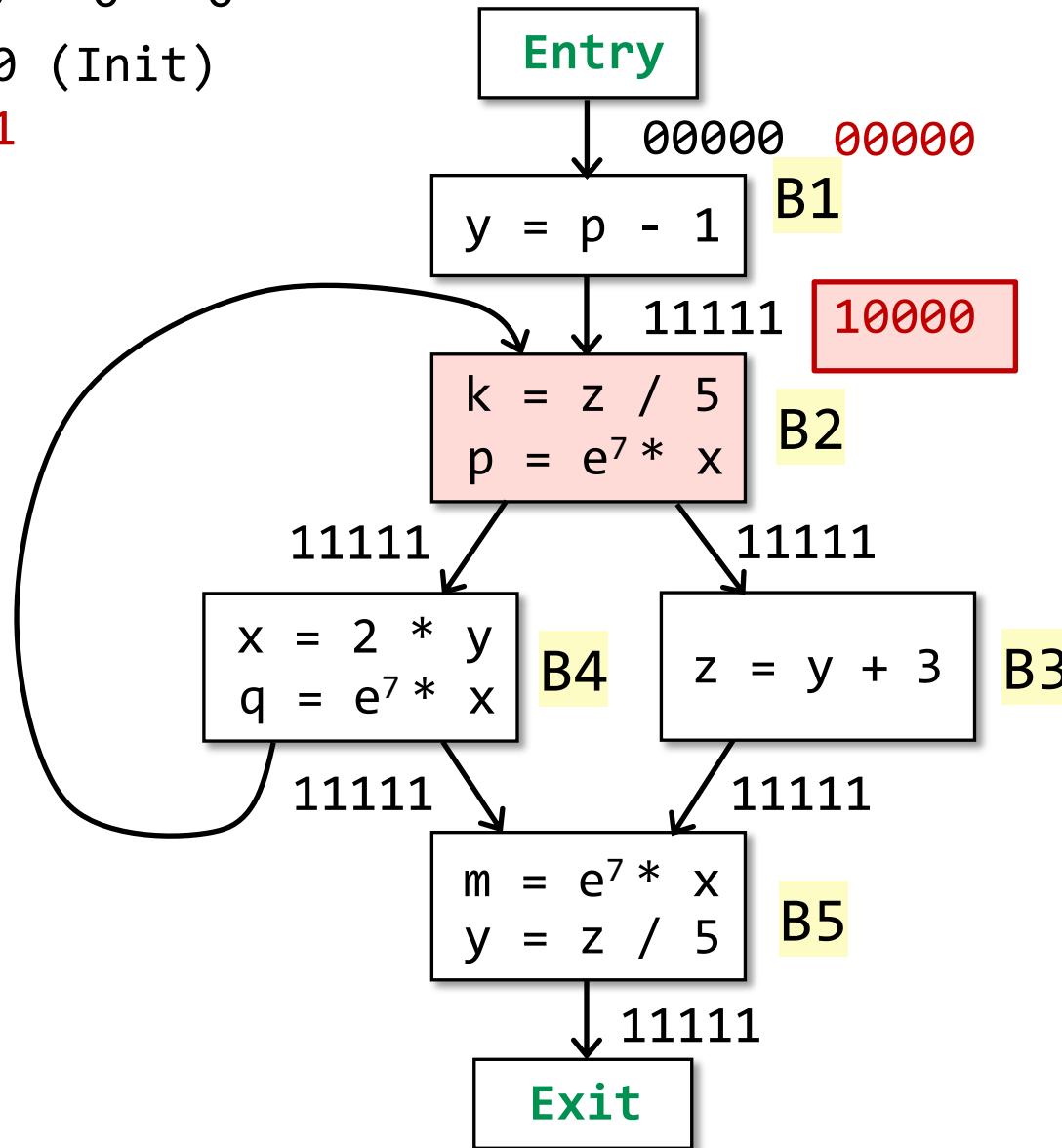
Iteration 1



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Iteration 0 (Init)

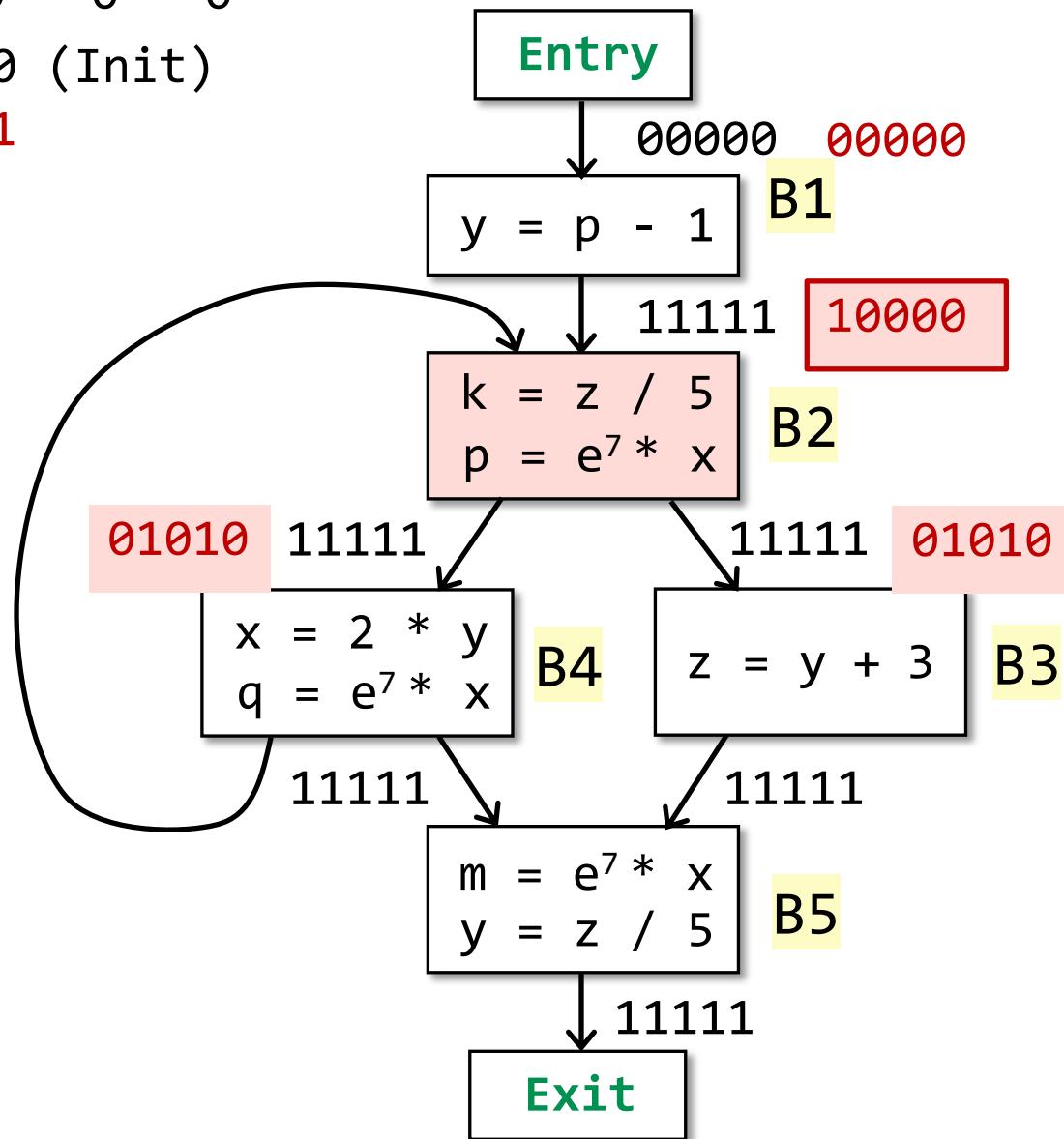
Iteration 1



$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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Iteration 0 (Init)

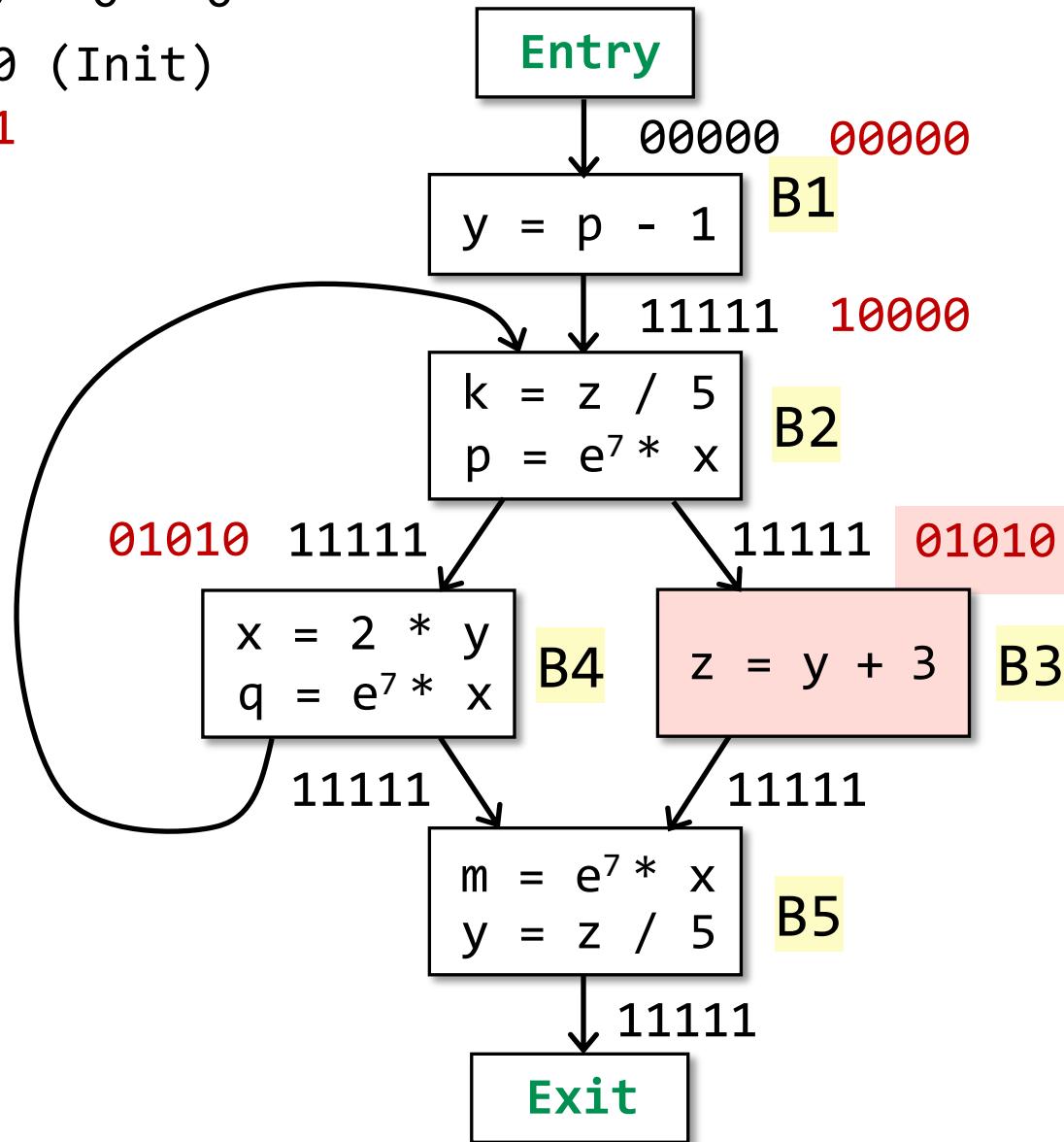
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Iteration 0 (Init)

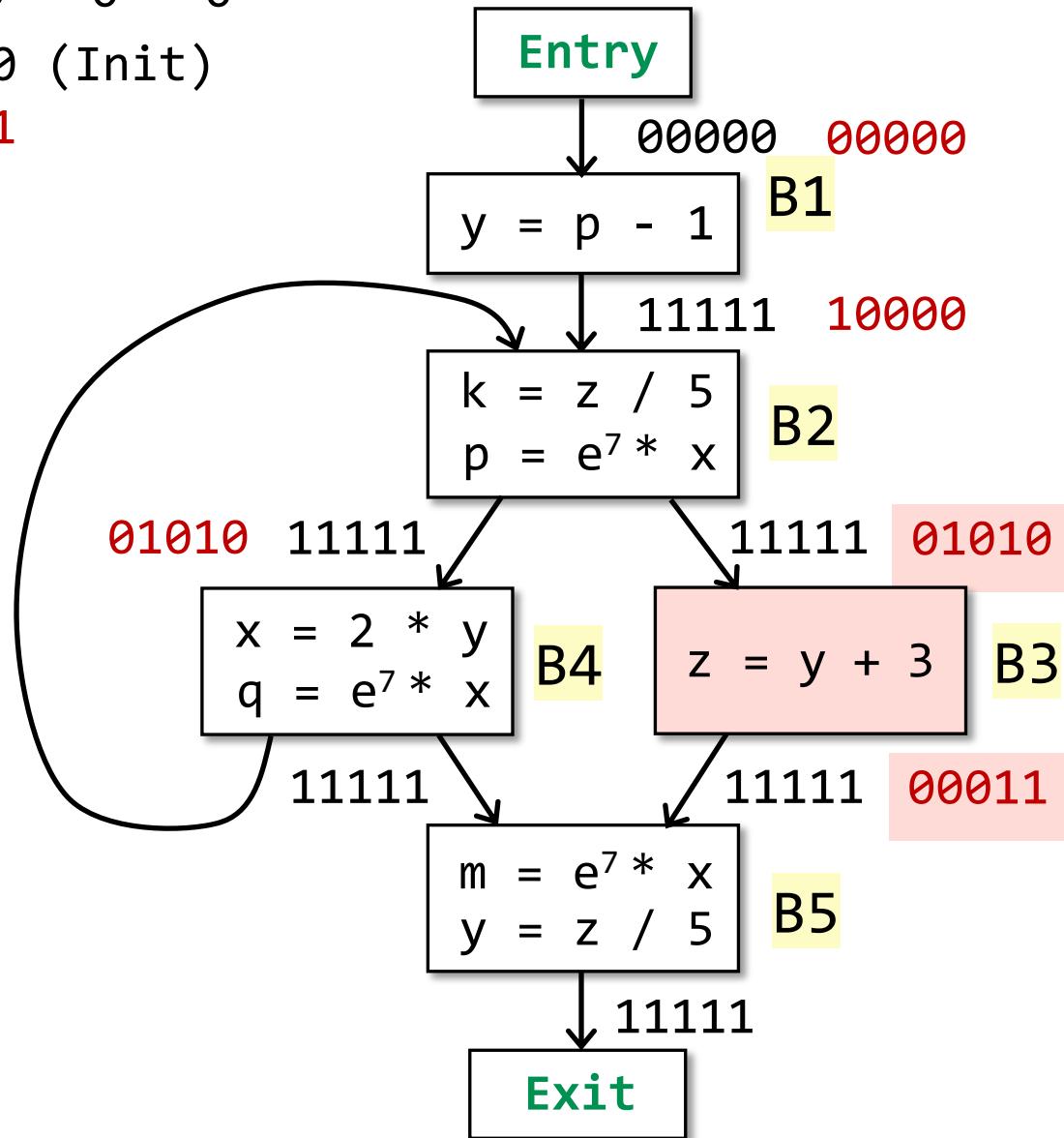
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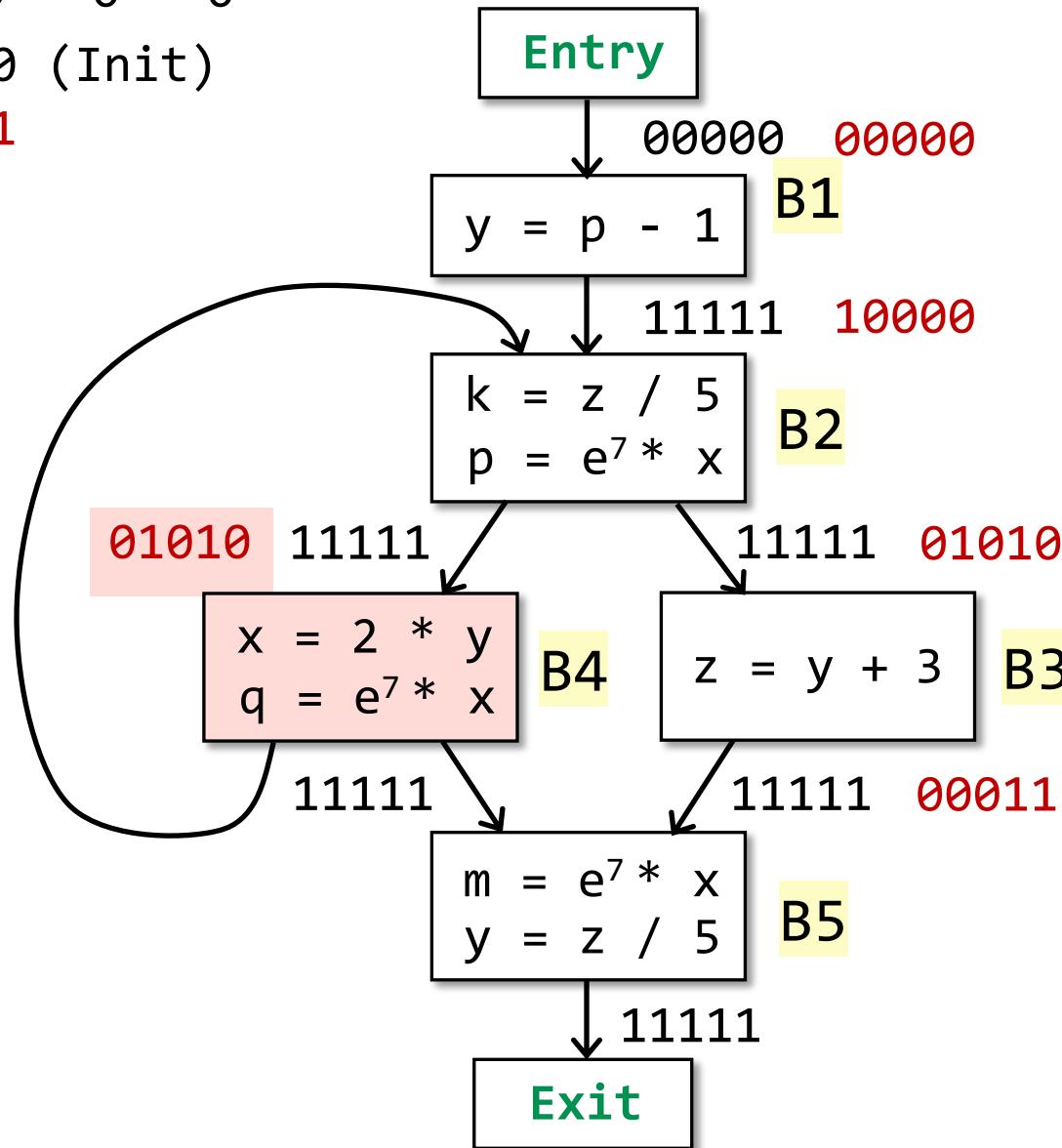
Iteration 1



$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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Iteration 0 (Init)

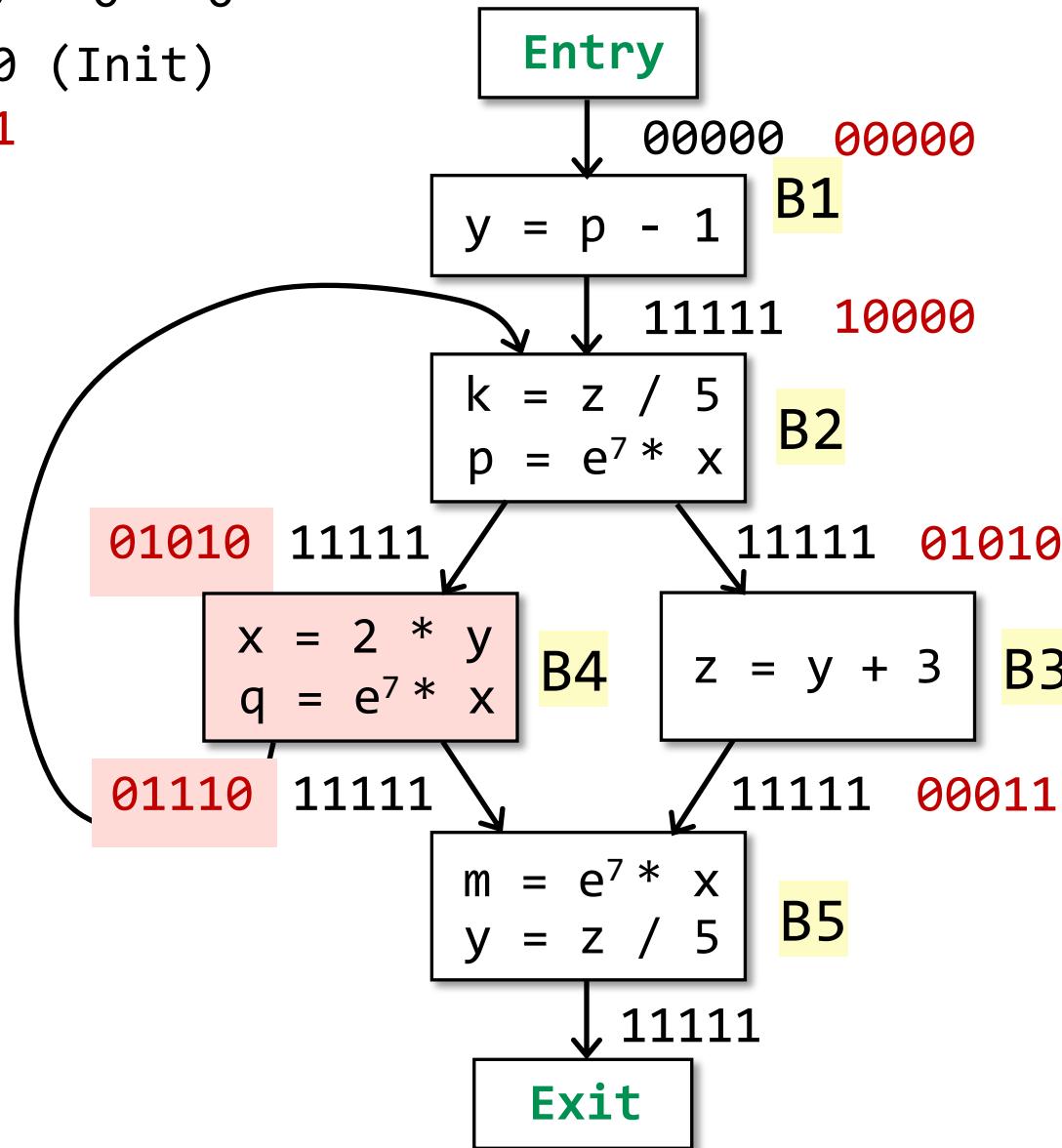
Iteration 1



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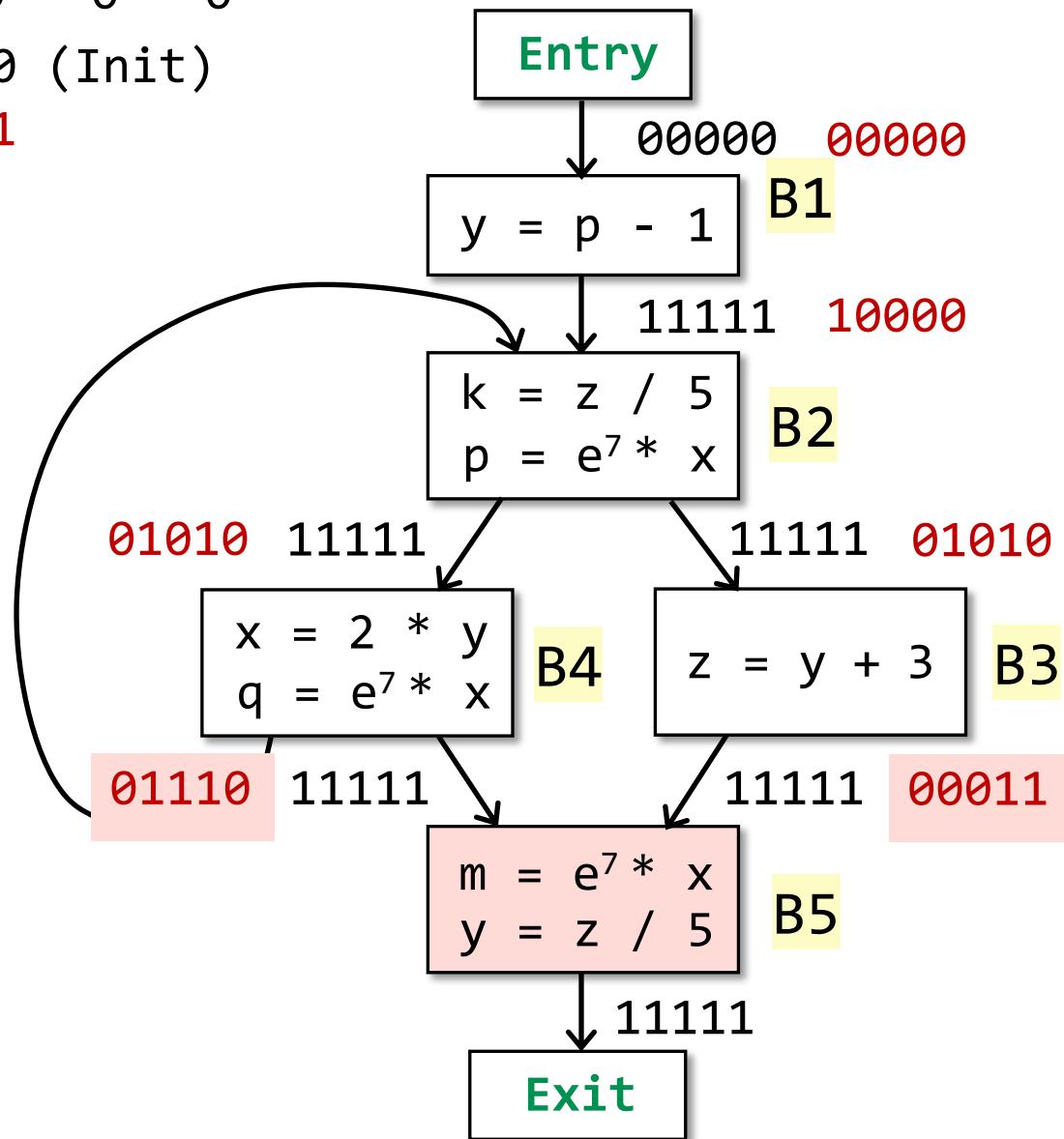
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Iteration 0 (Init)

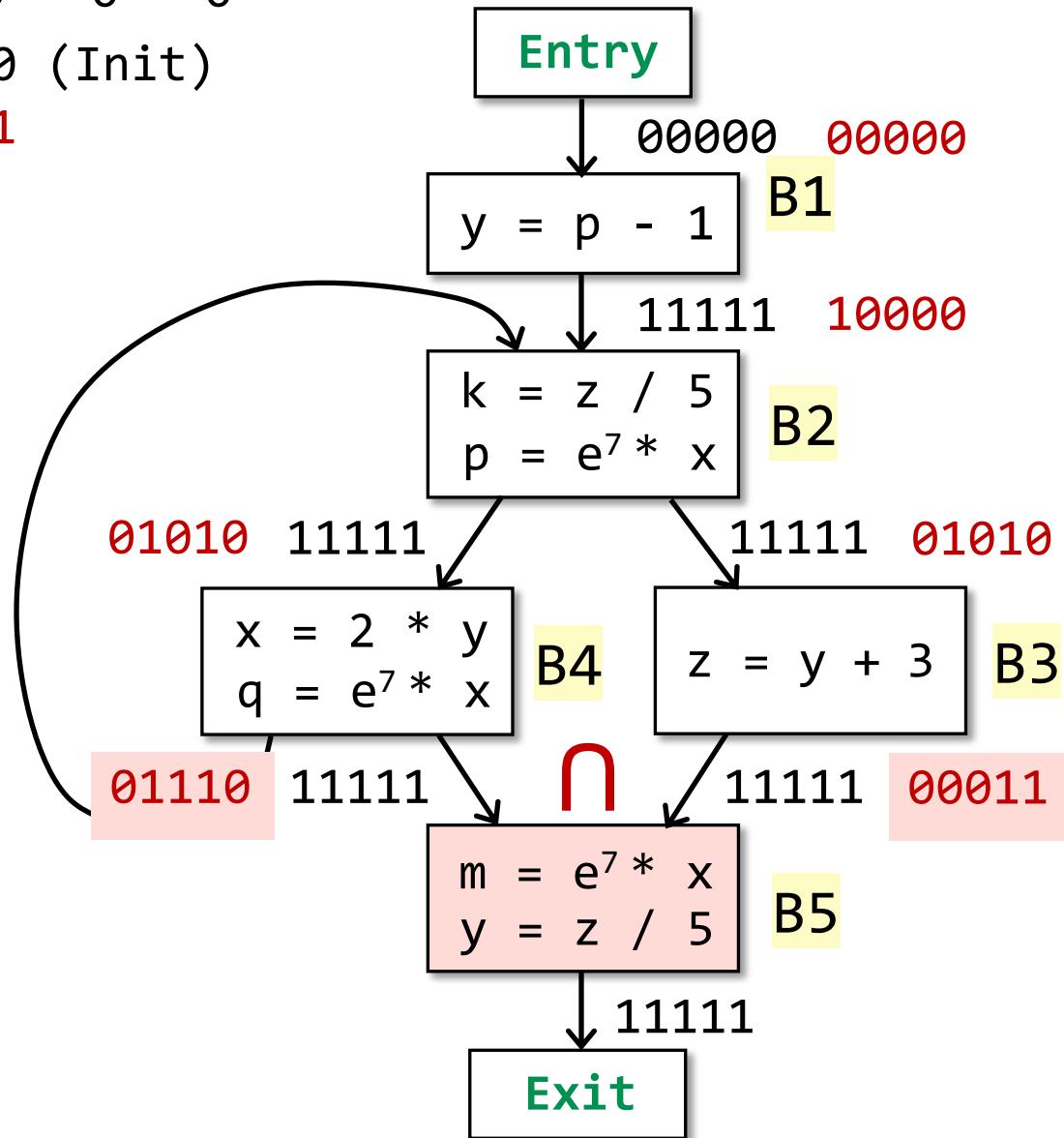
Iteration 1



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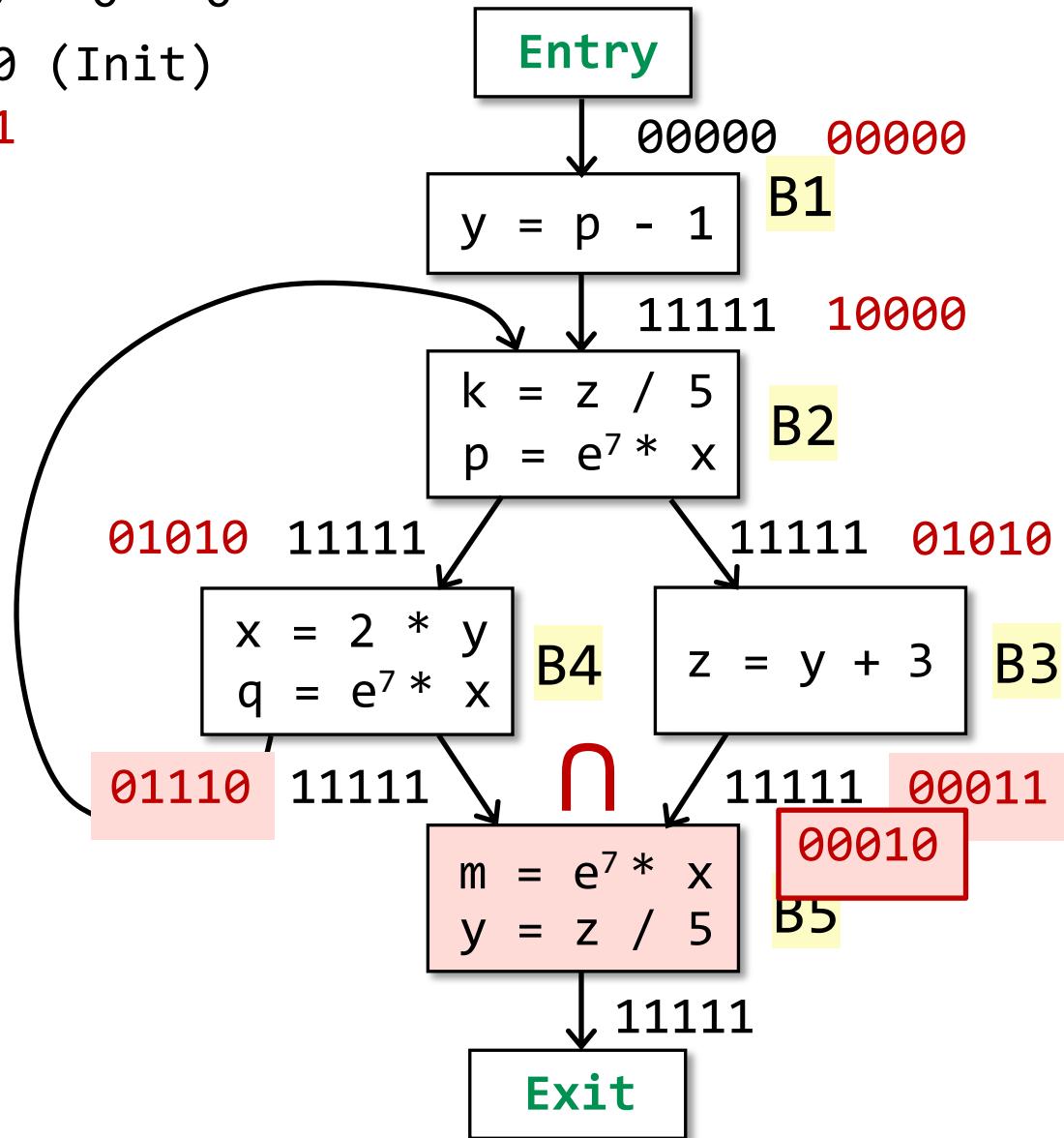
Iteration 1



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Iteration 0 (Init)

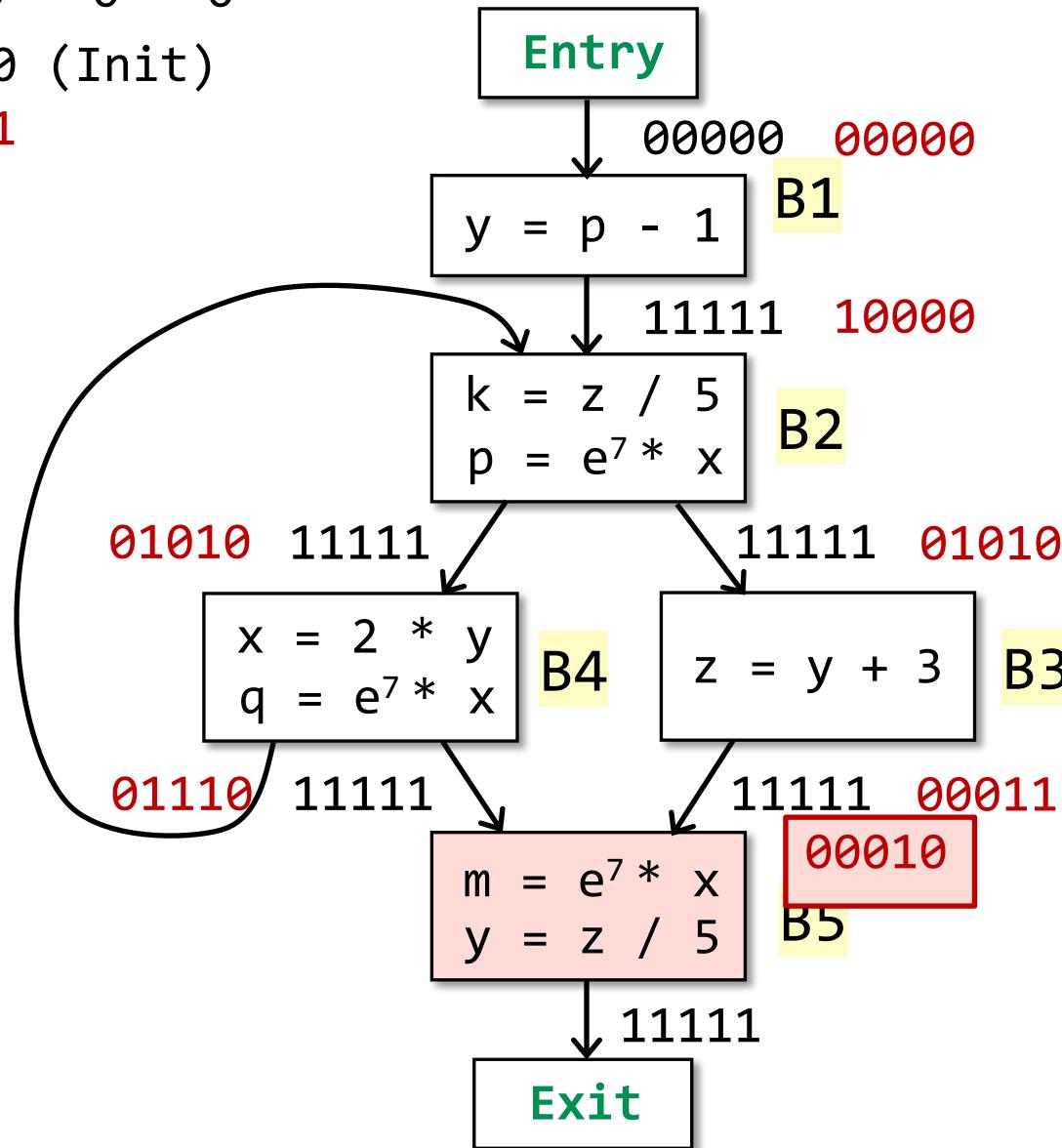
Iteration 1



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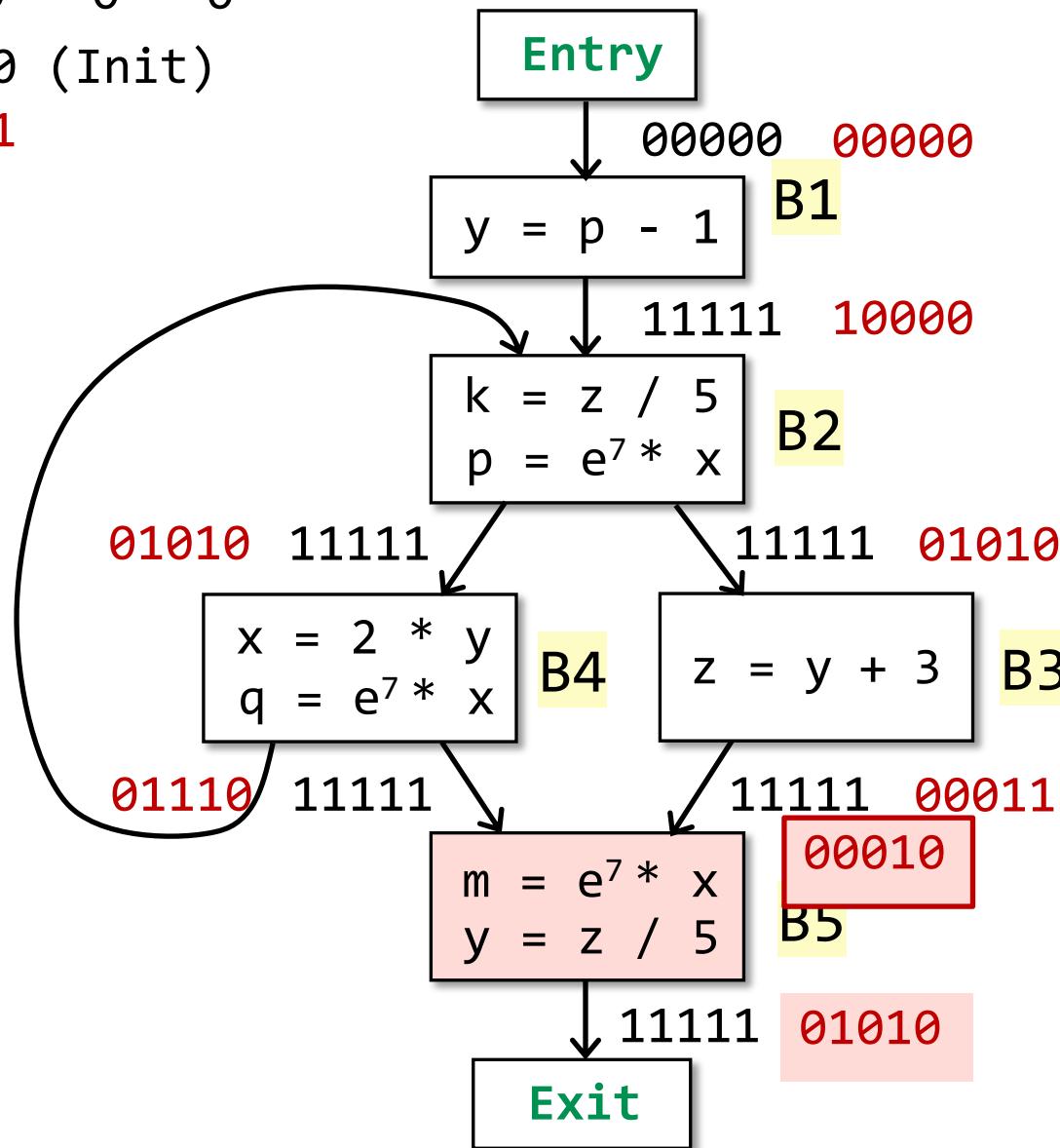
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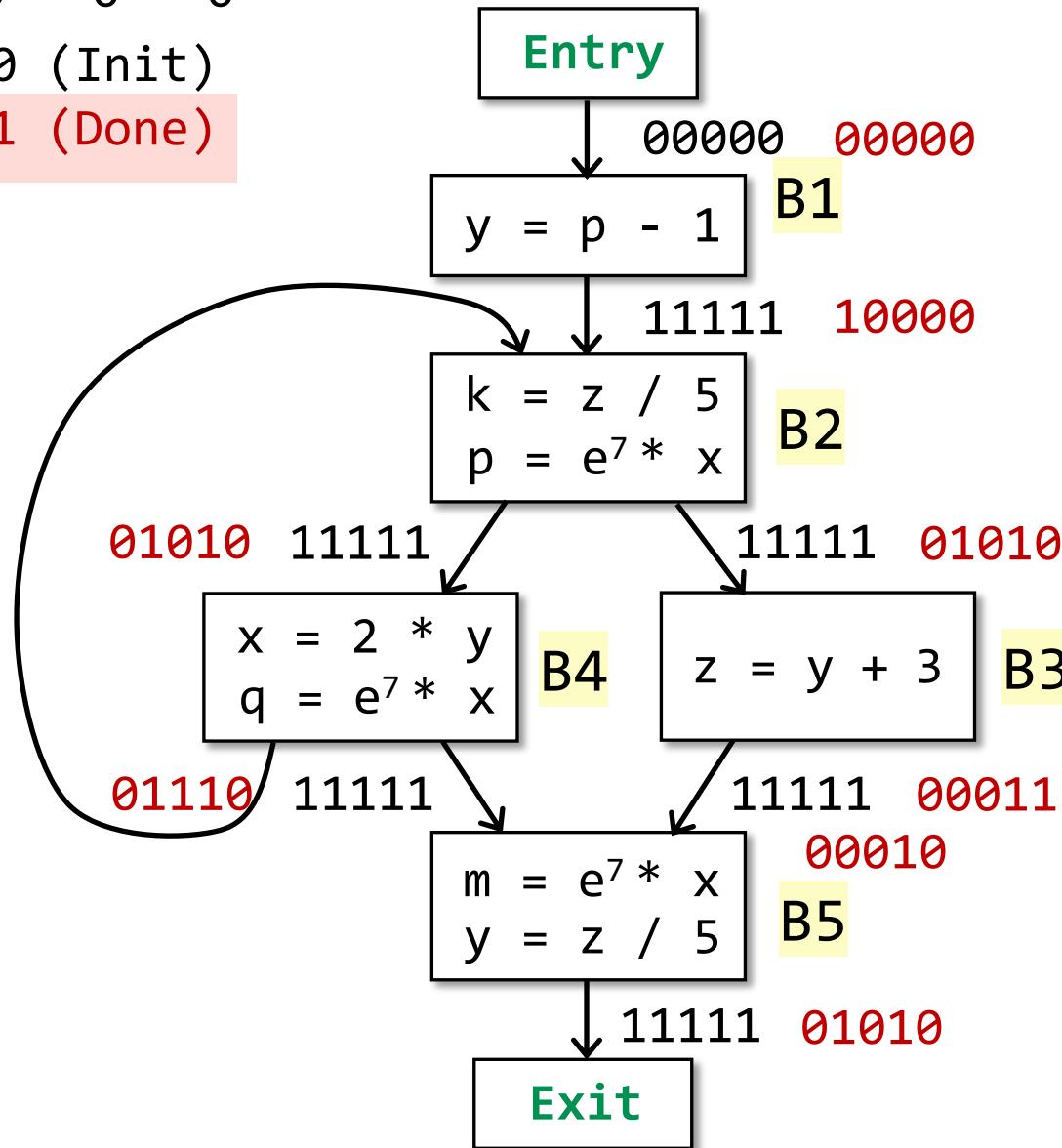
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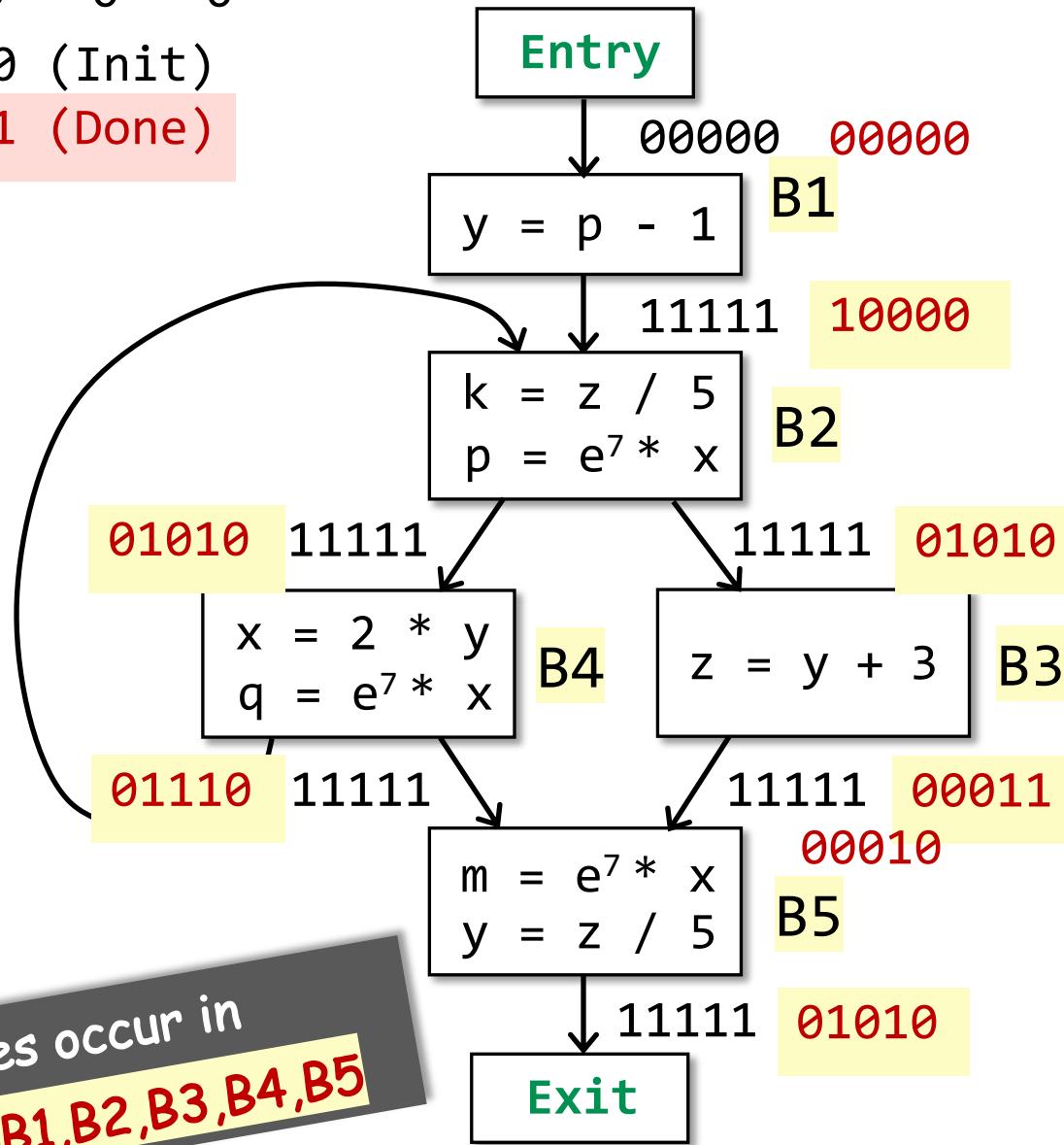
Iteration 1 (Done)



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Iteration 0 (Init)

Iteration 1 (Done)



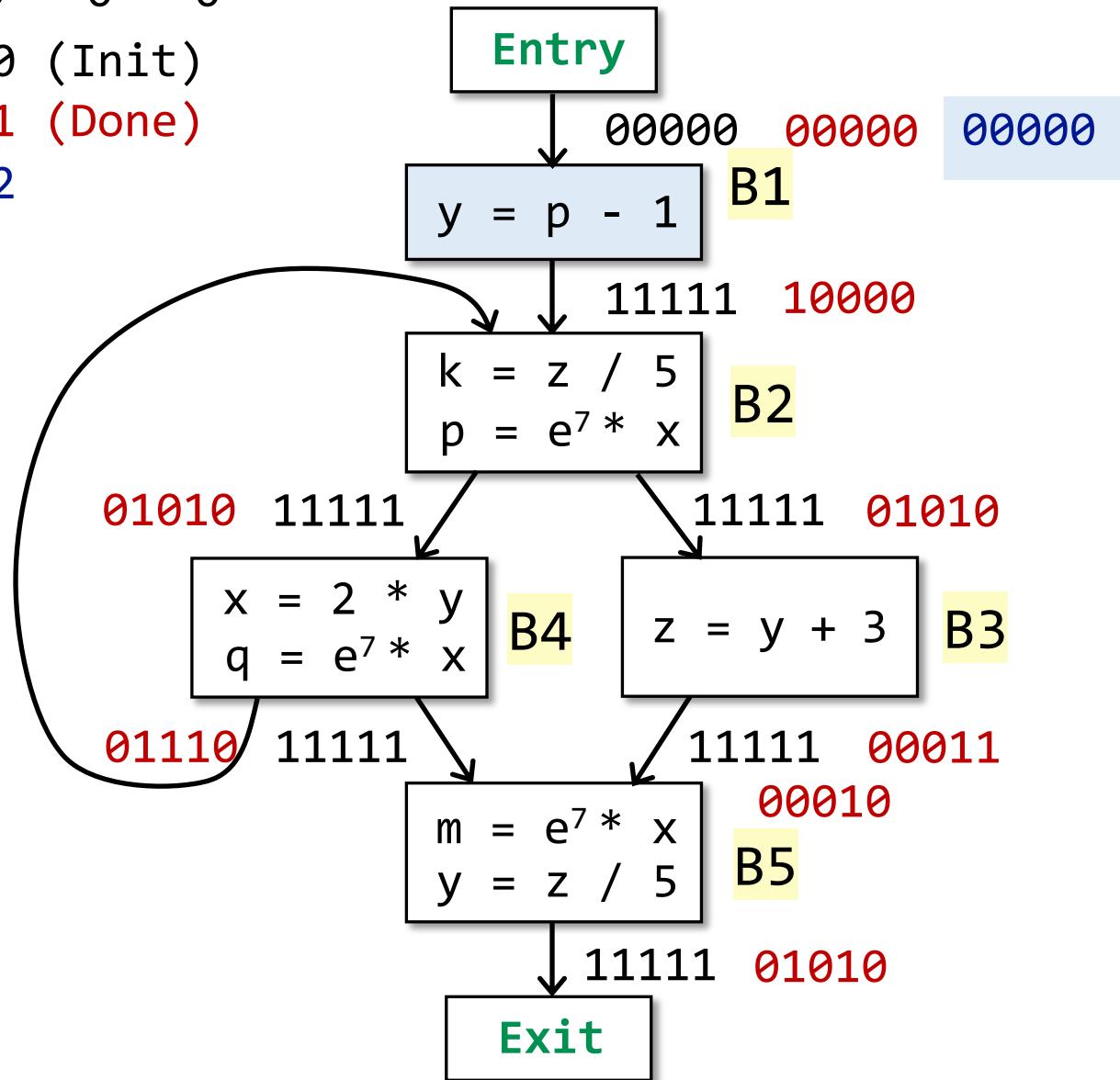
Changes occur in
 OUT[] of B1,B2,B3,B4,B5

$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

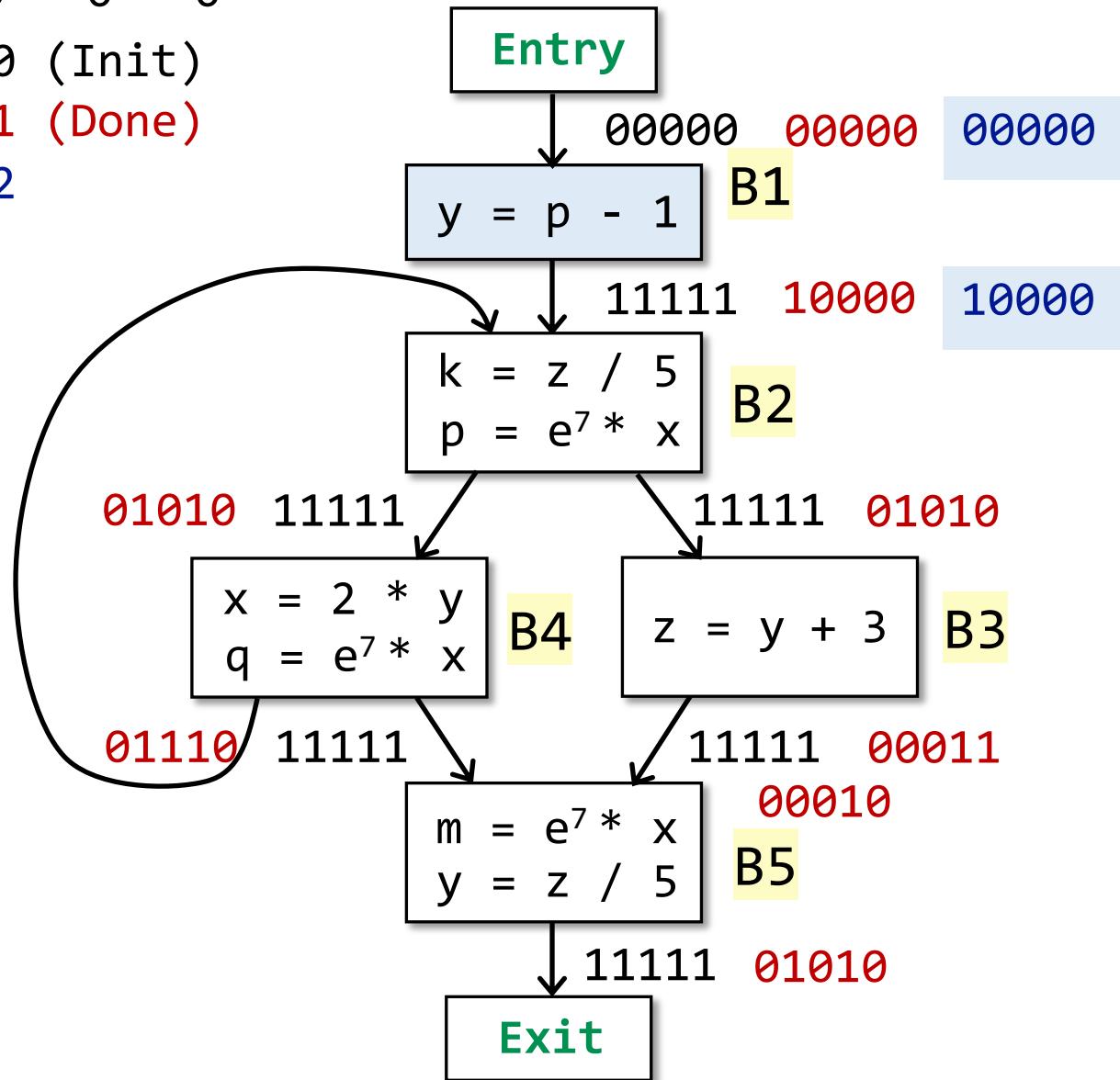


$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2

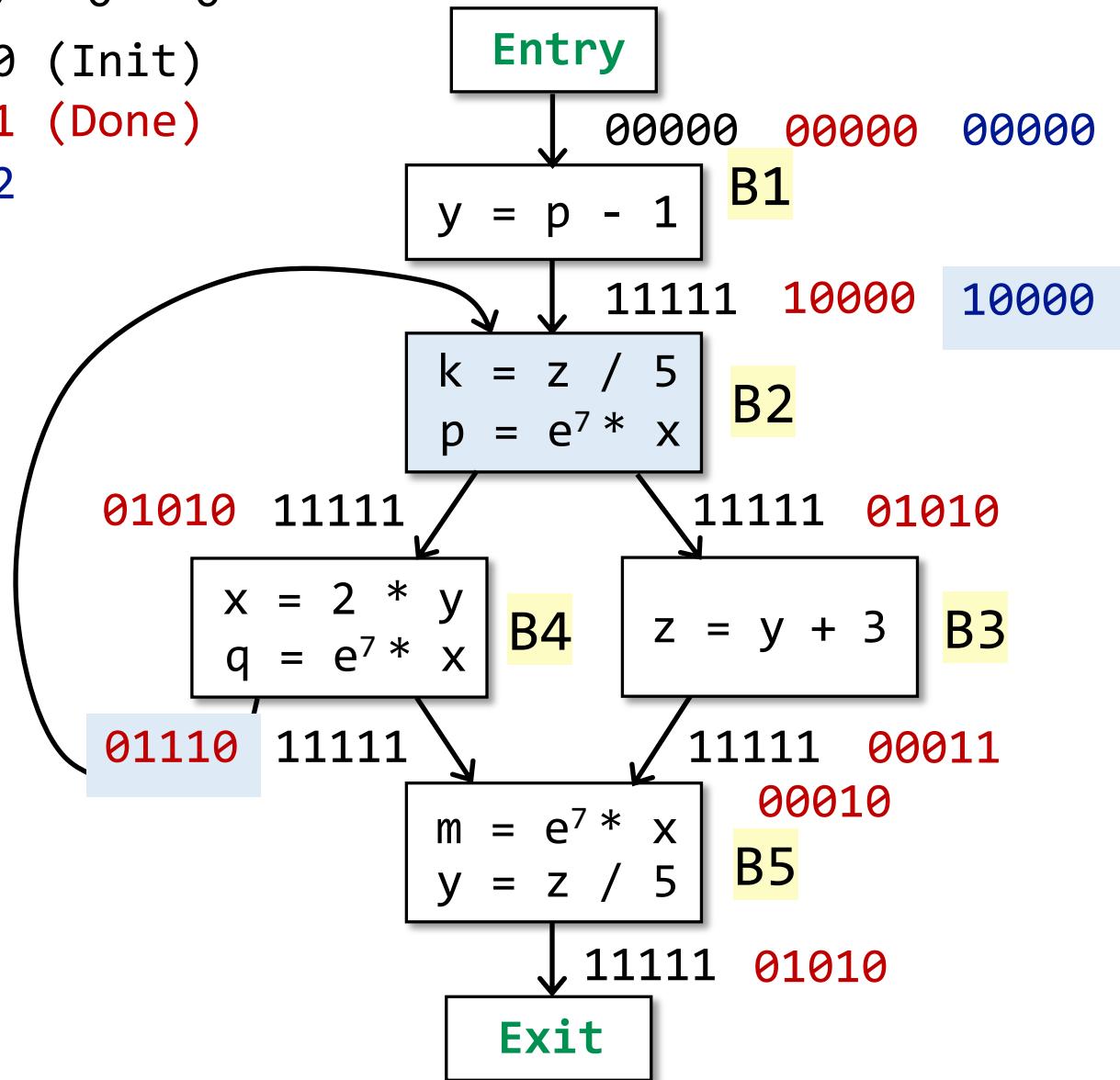


$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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Iteration 0 (Init)

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Iteration 2

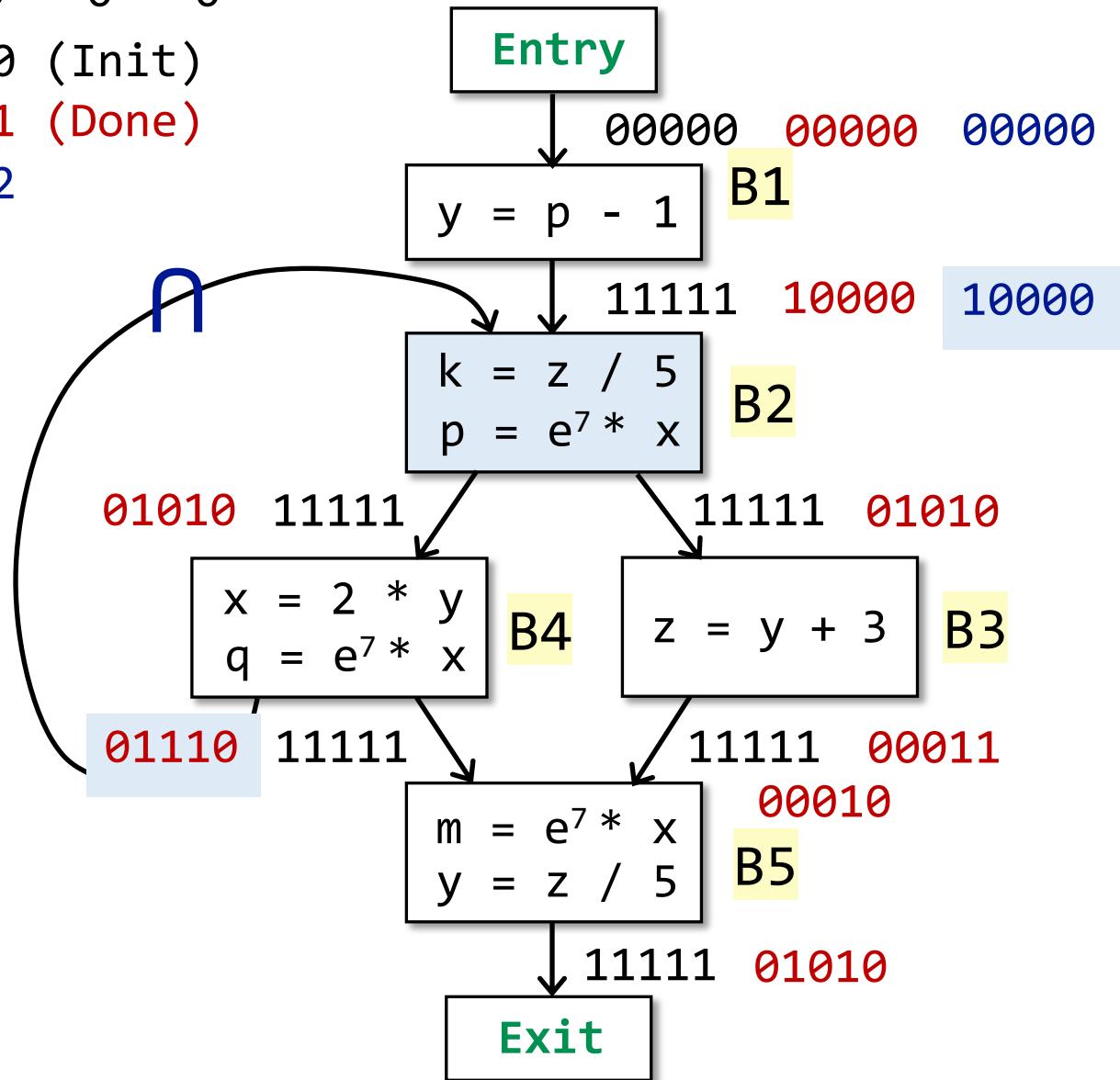


$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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Iteration 2

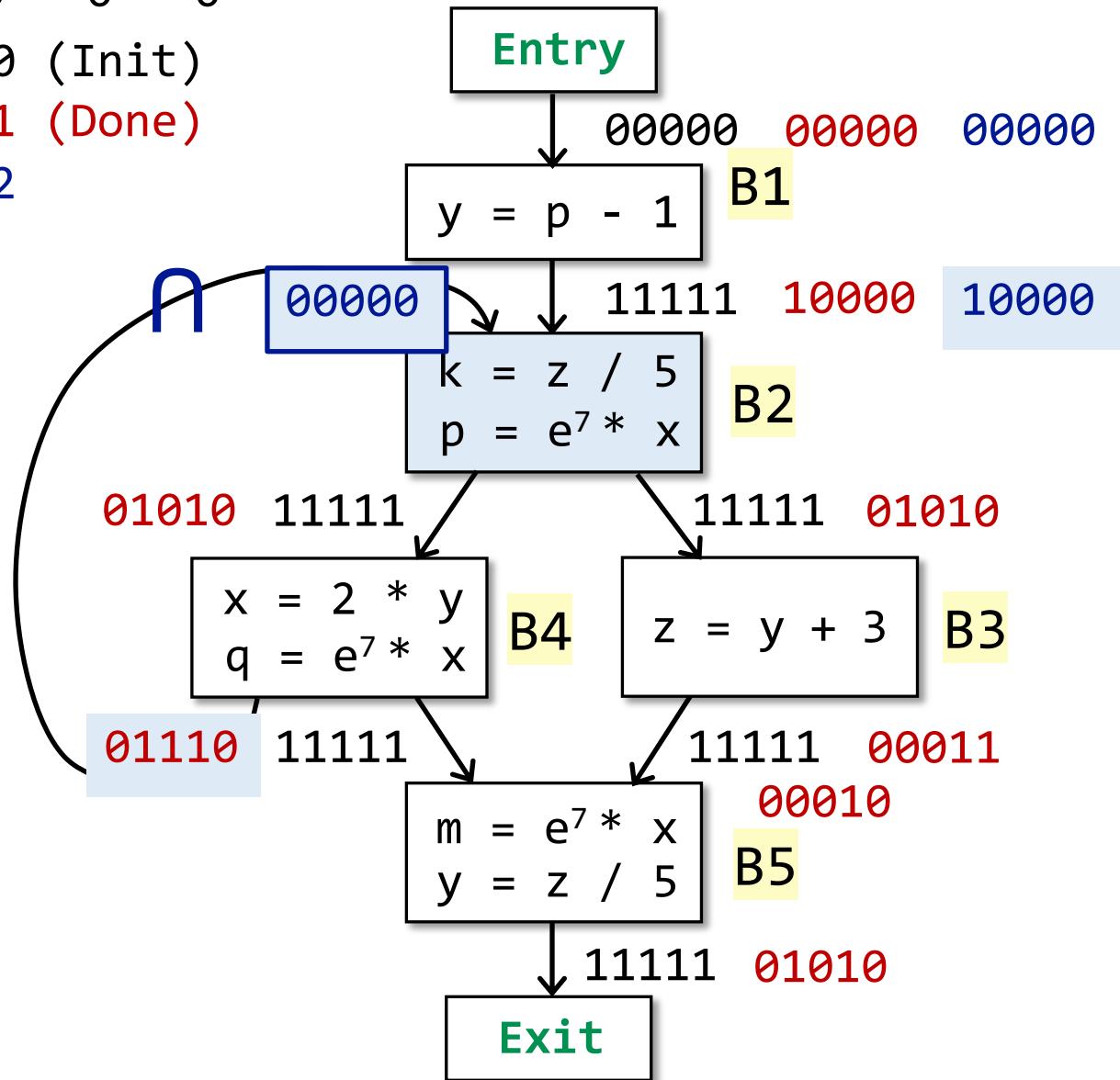


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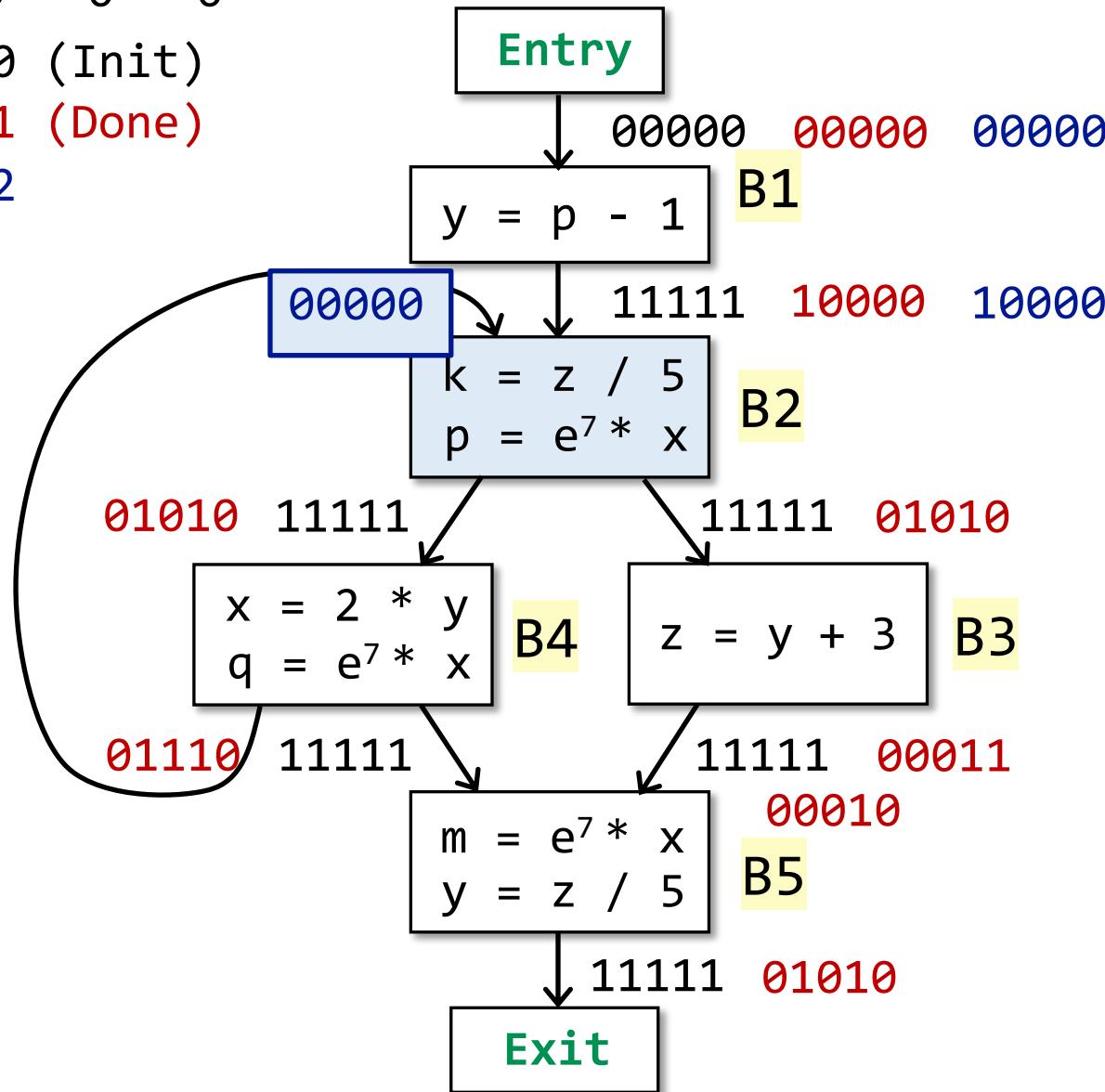


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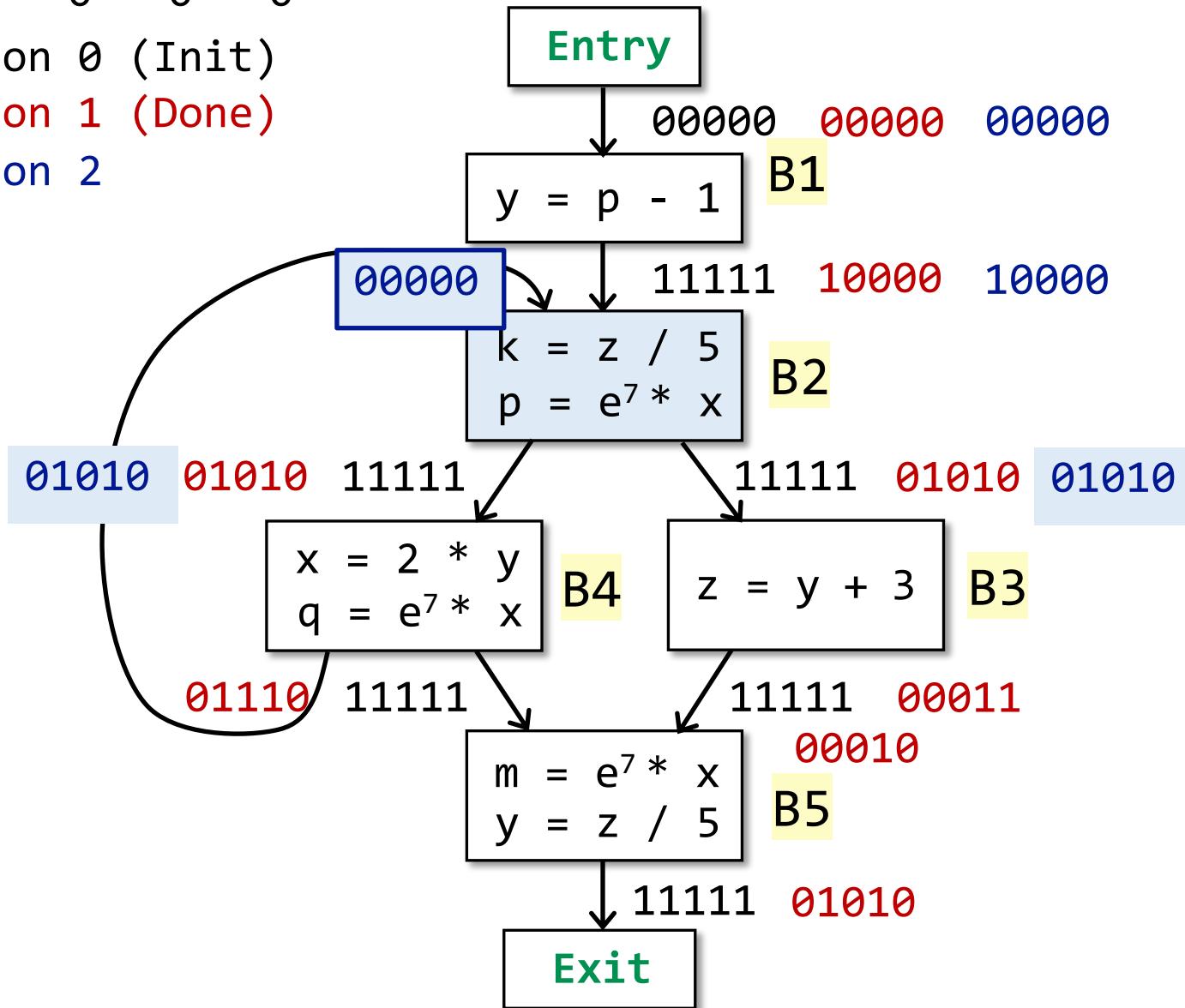


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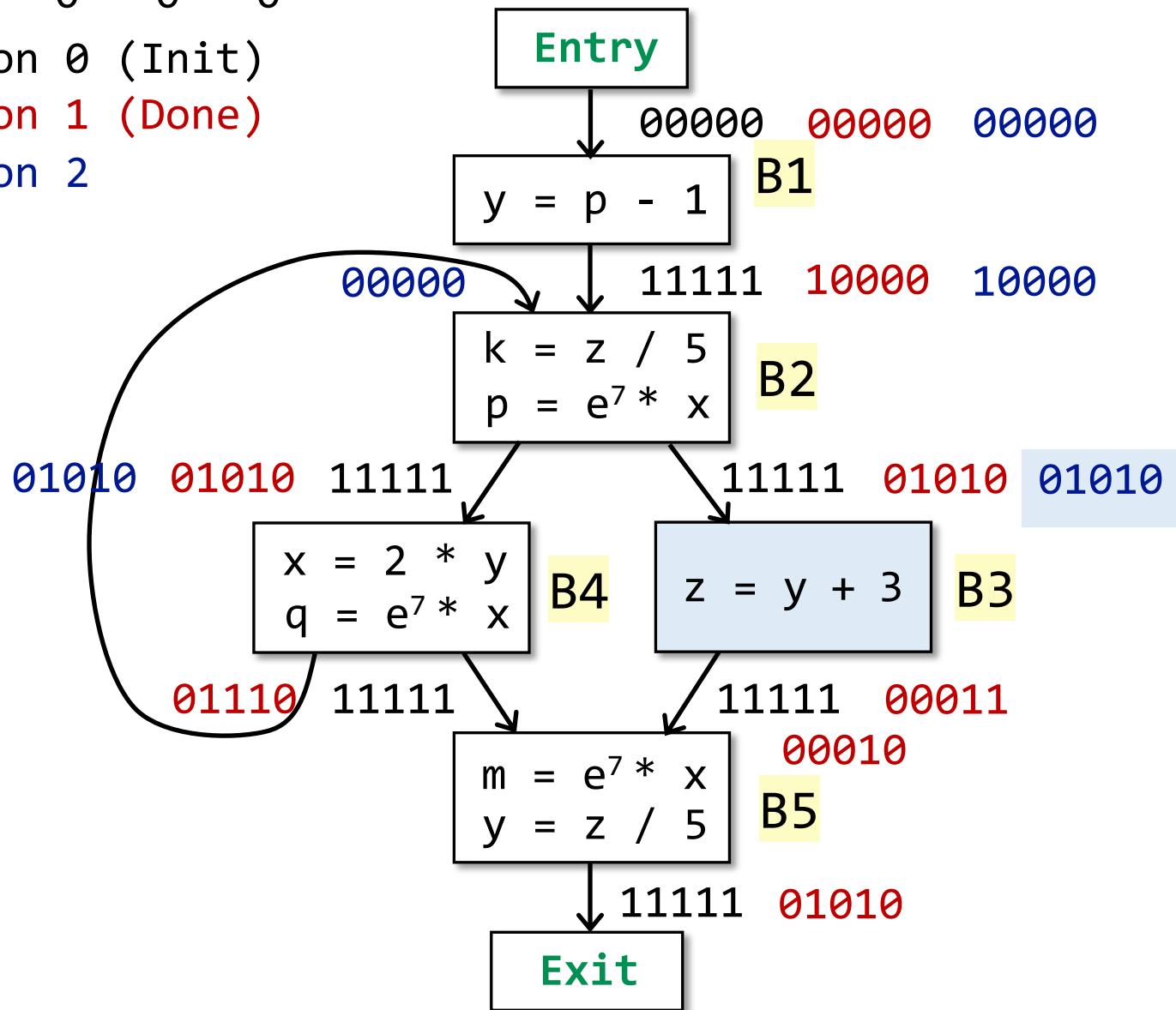


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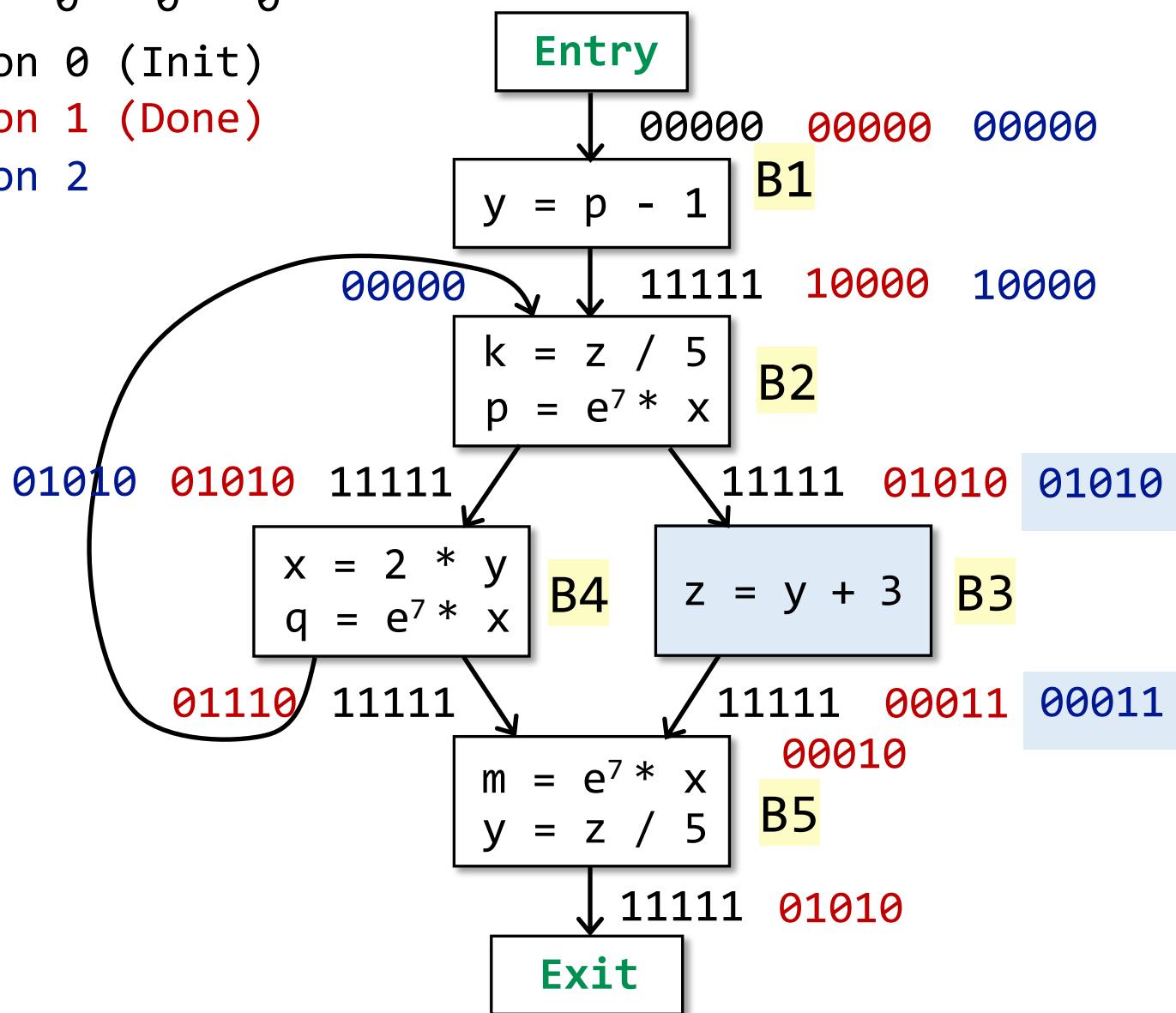


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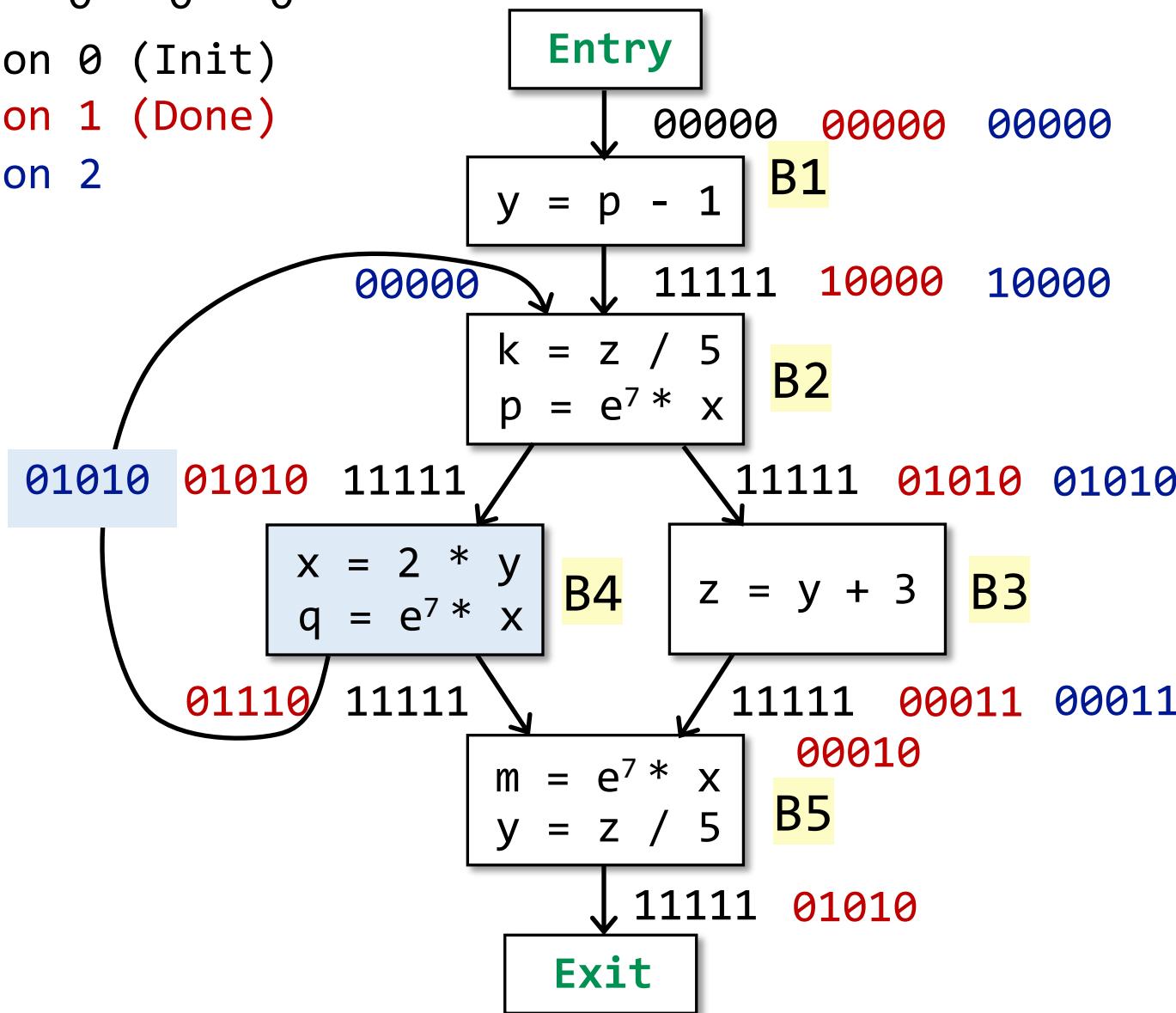


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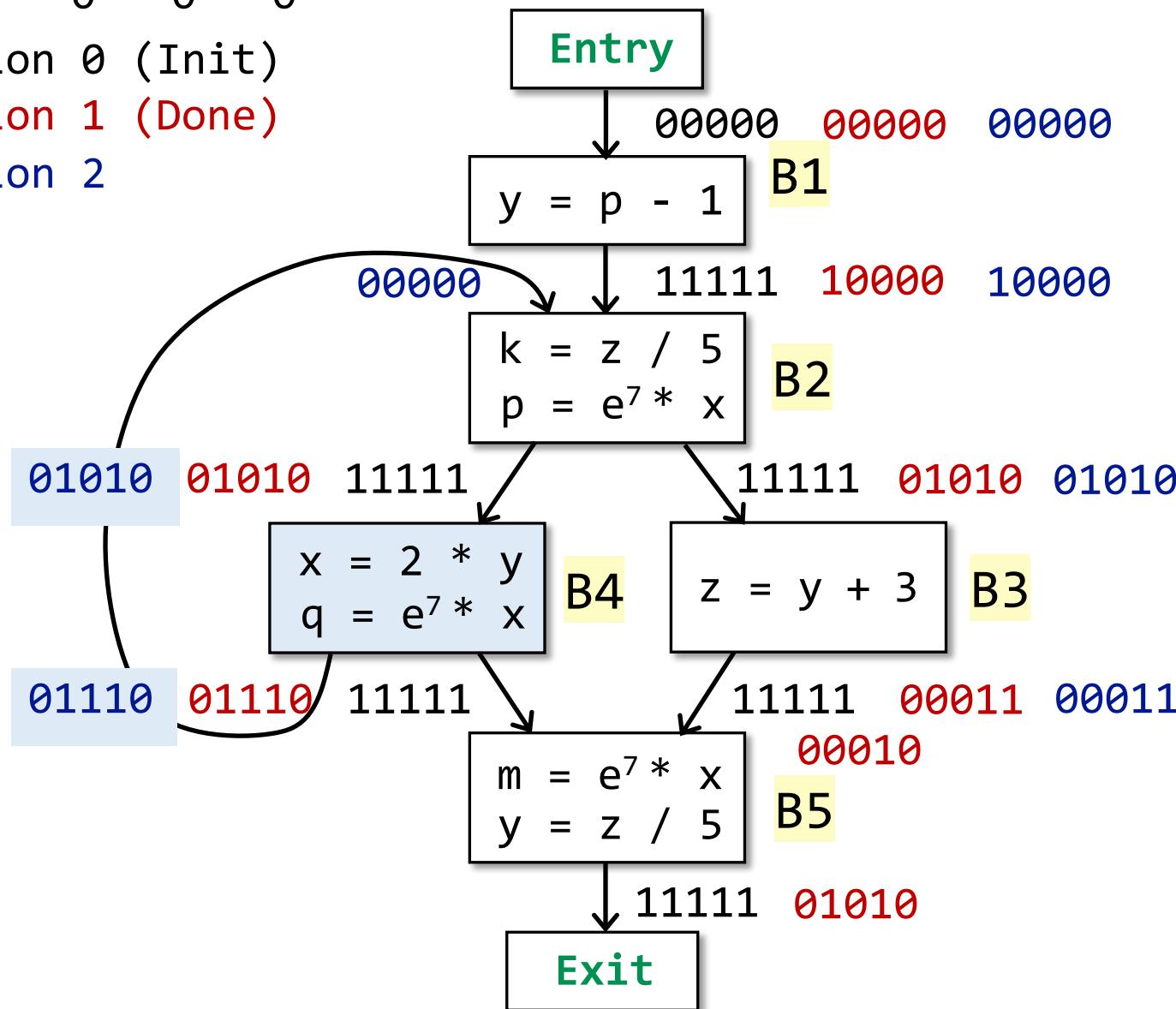


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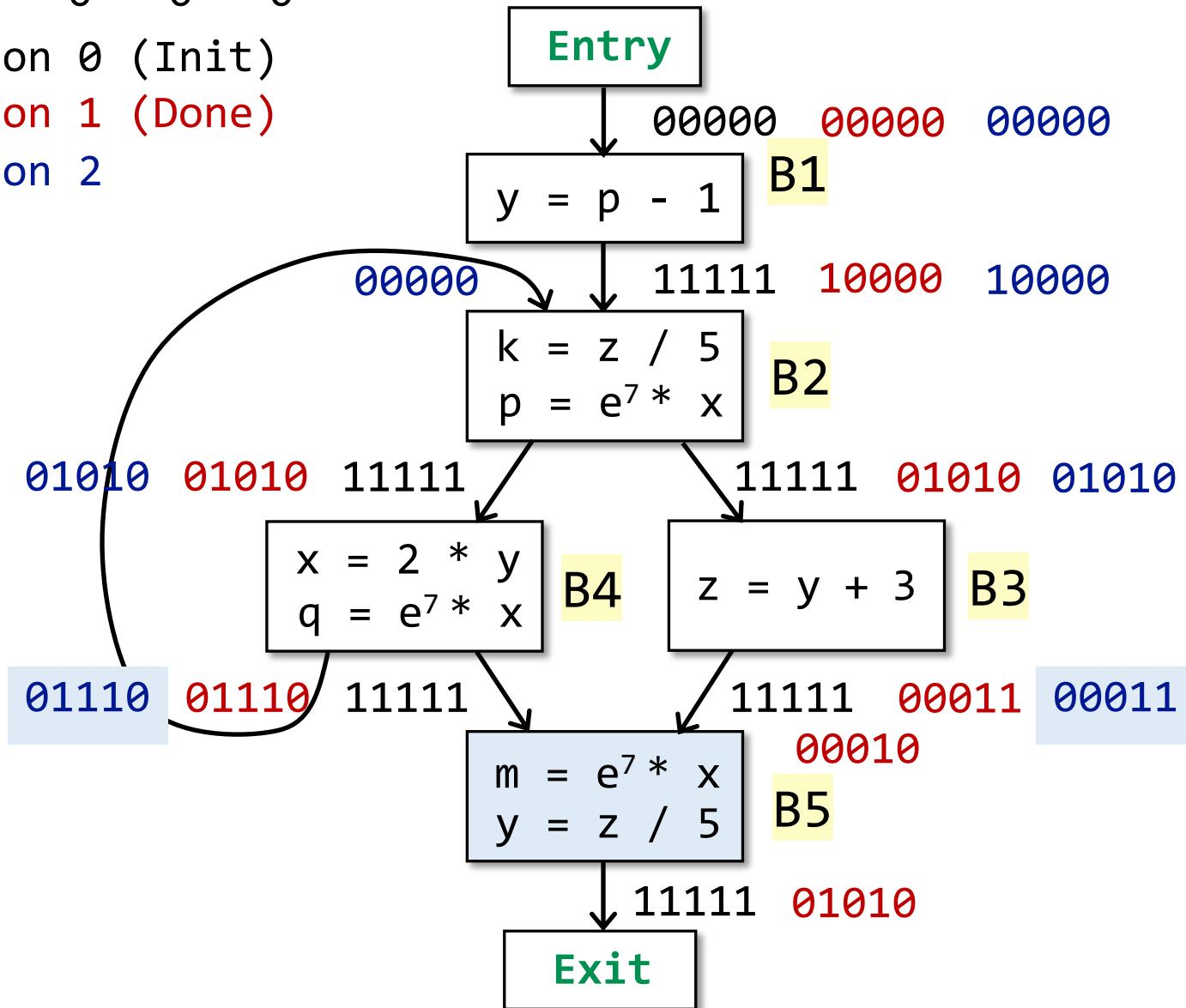


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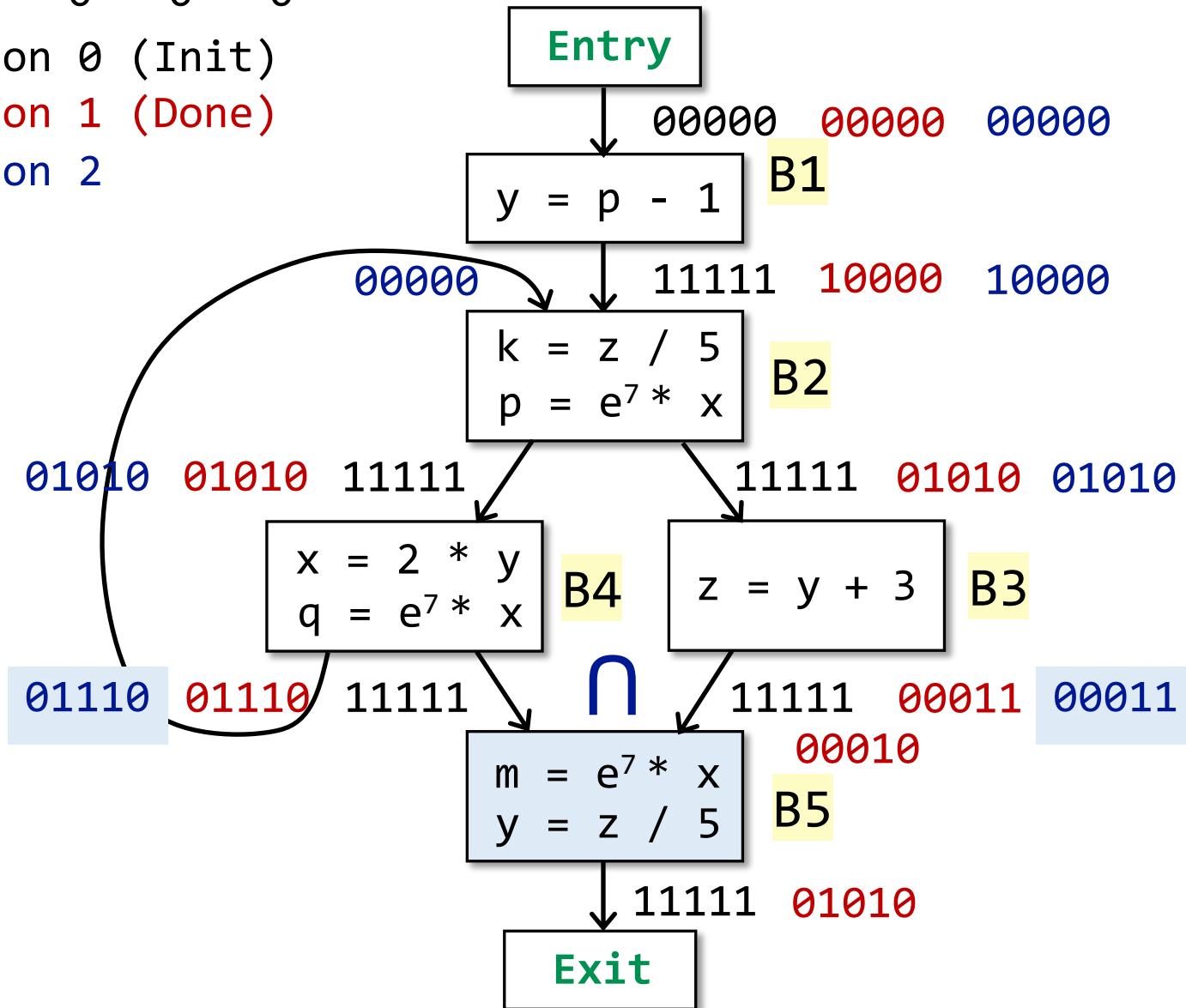


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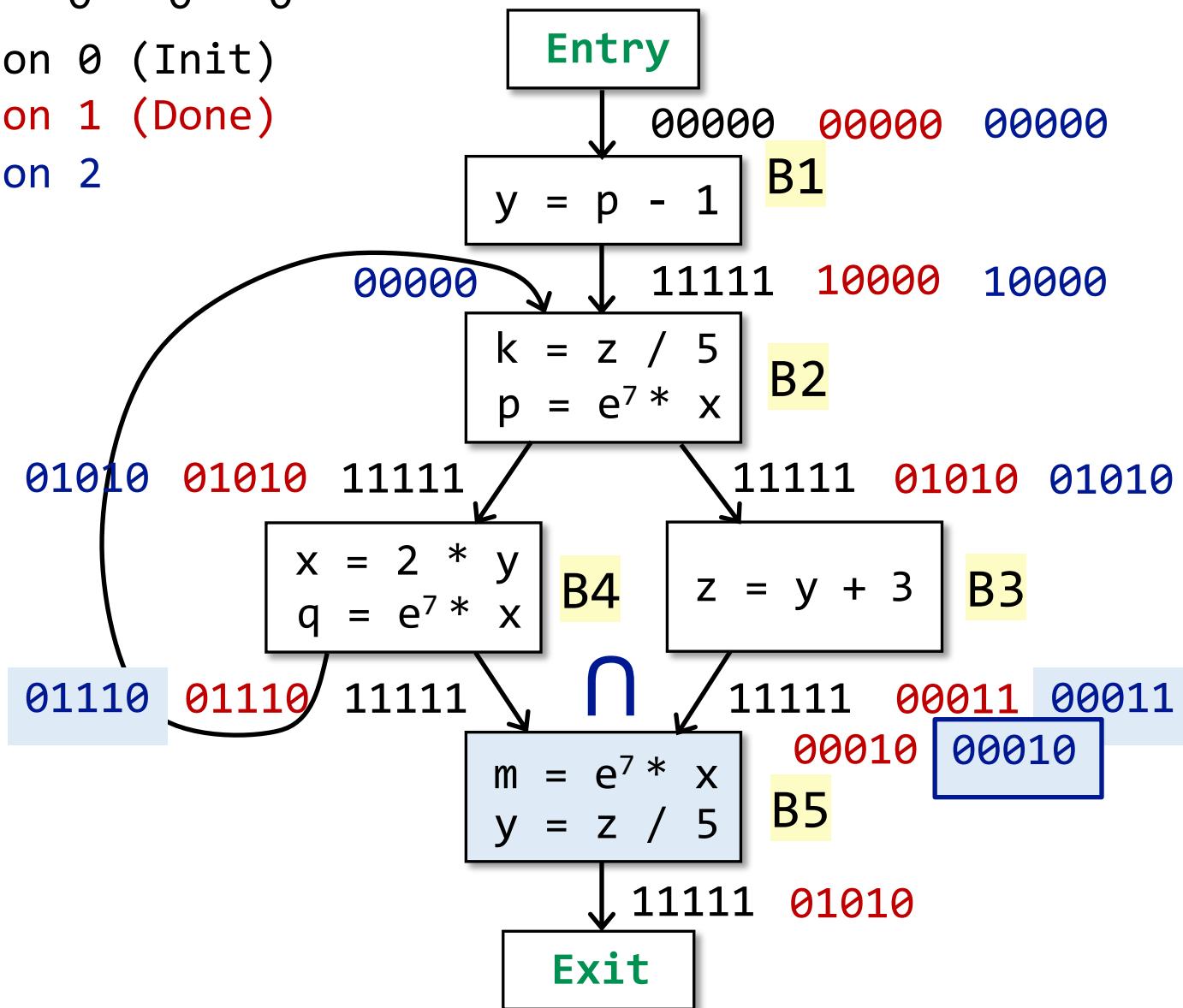


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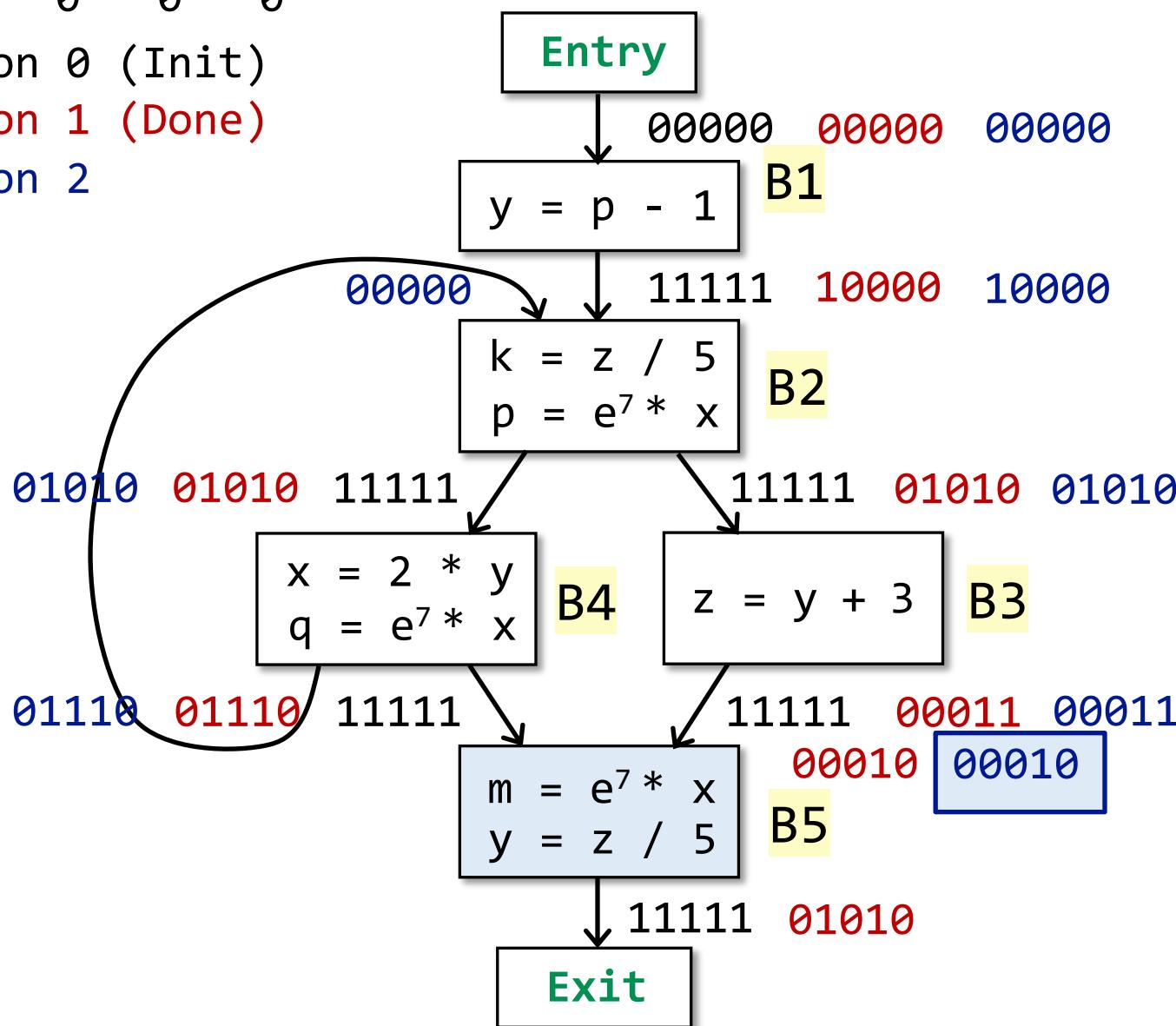


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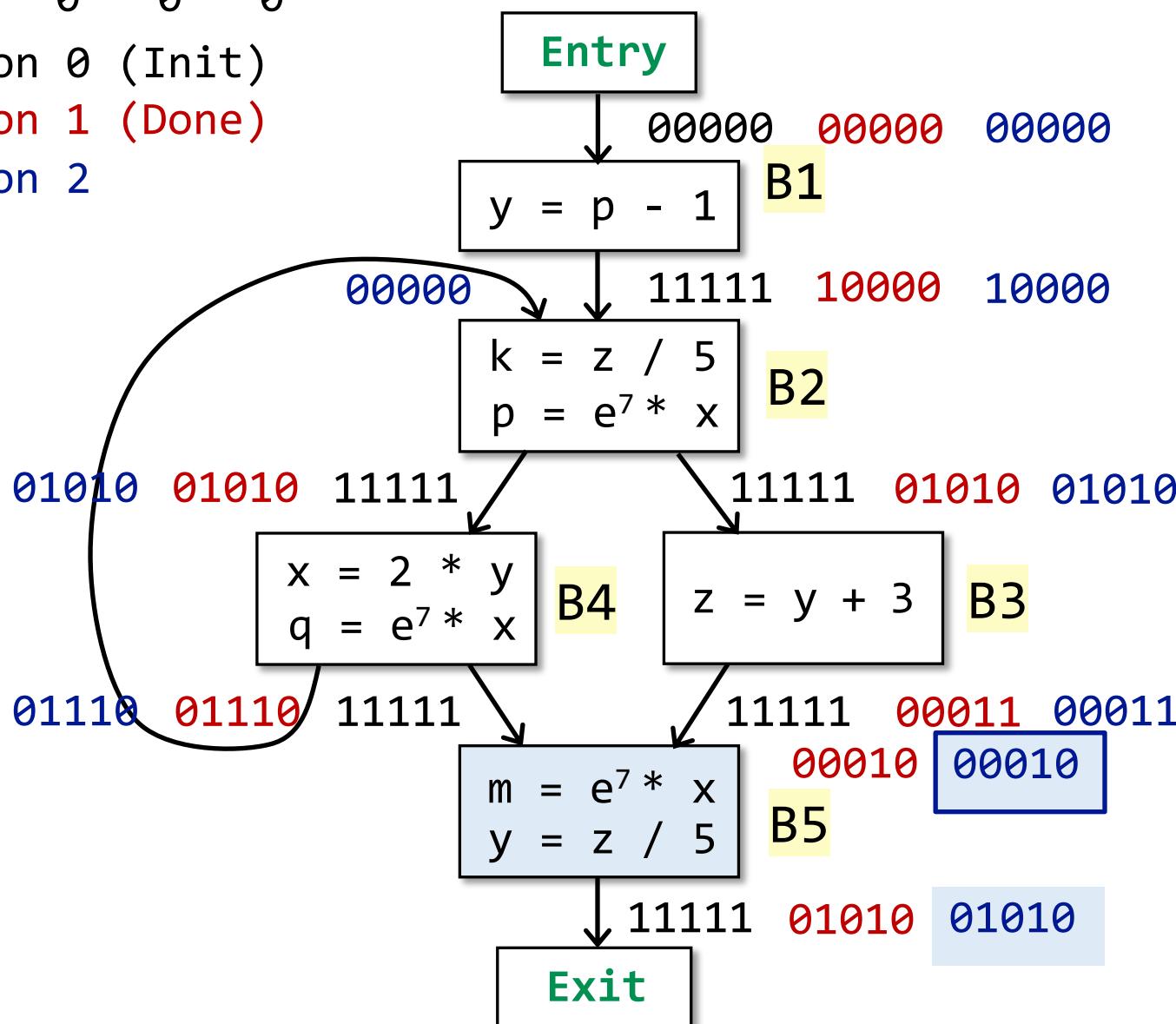


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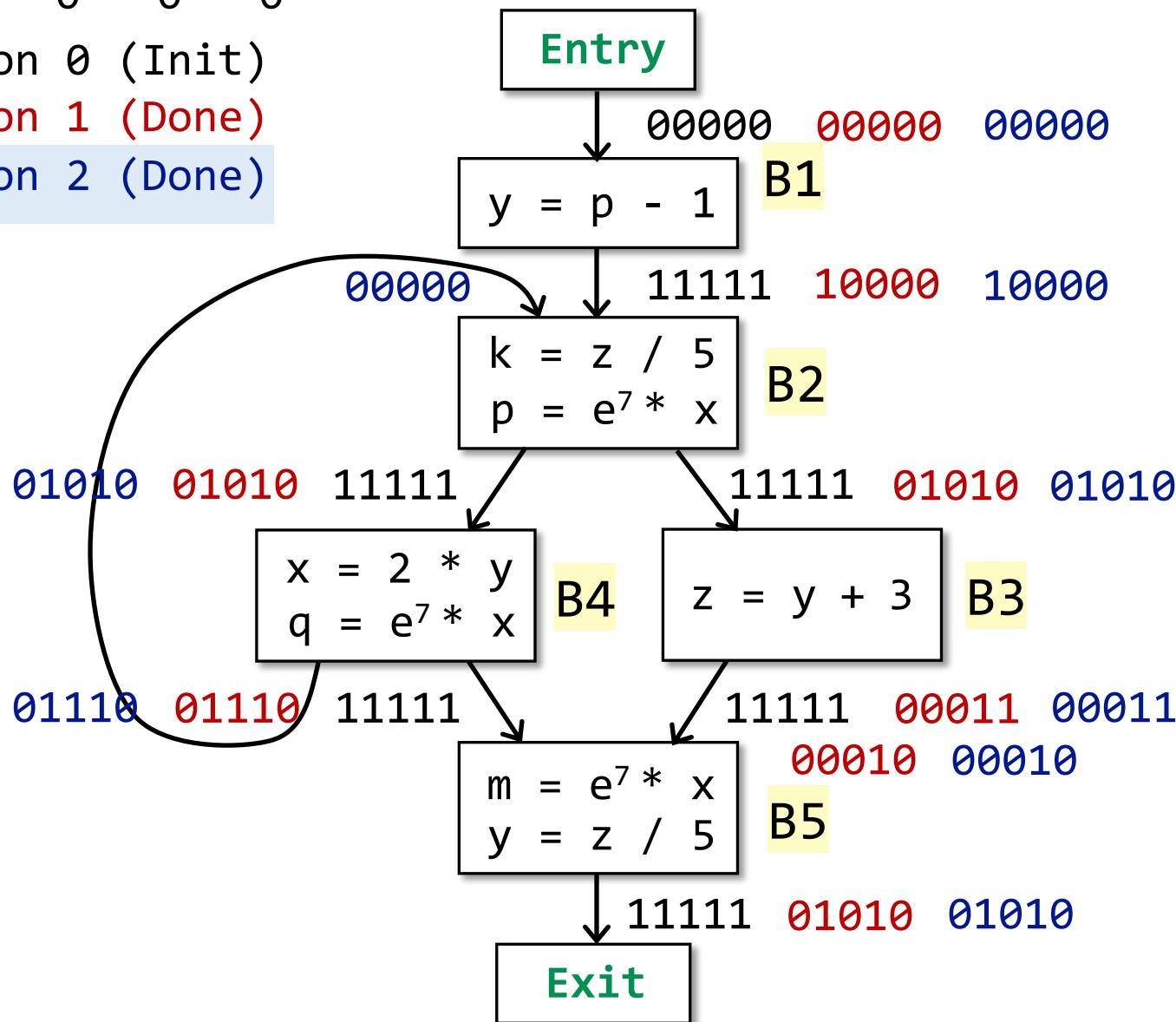


$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
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Iteration 0 (Init)

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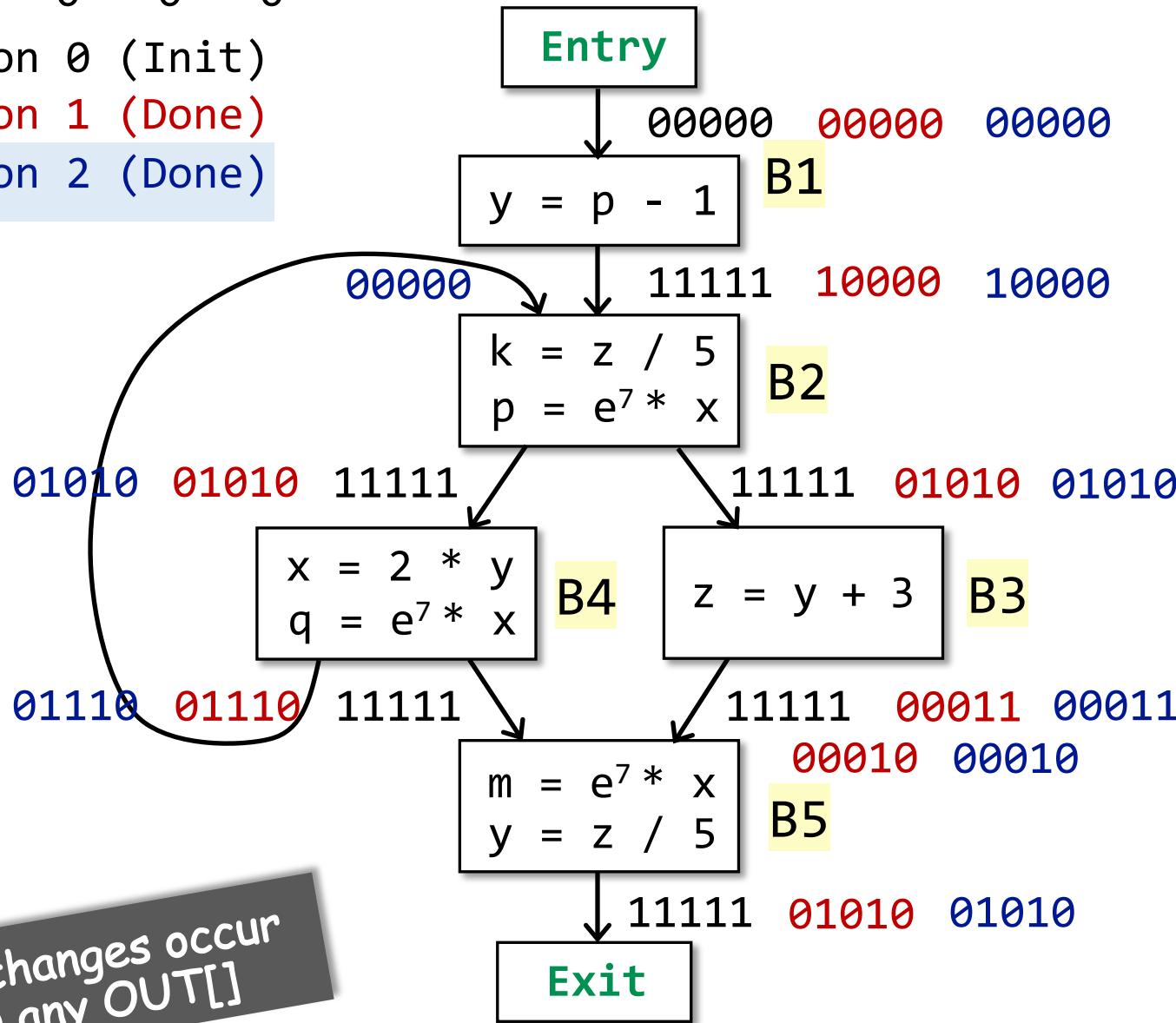


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Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)

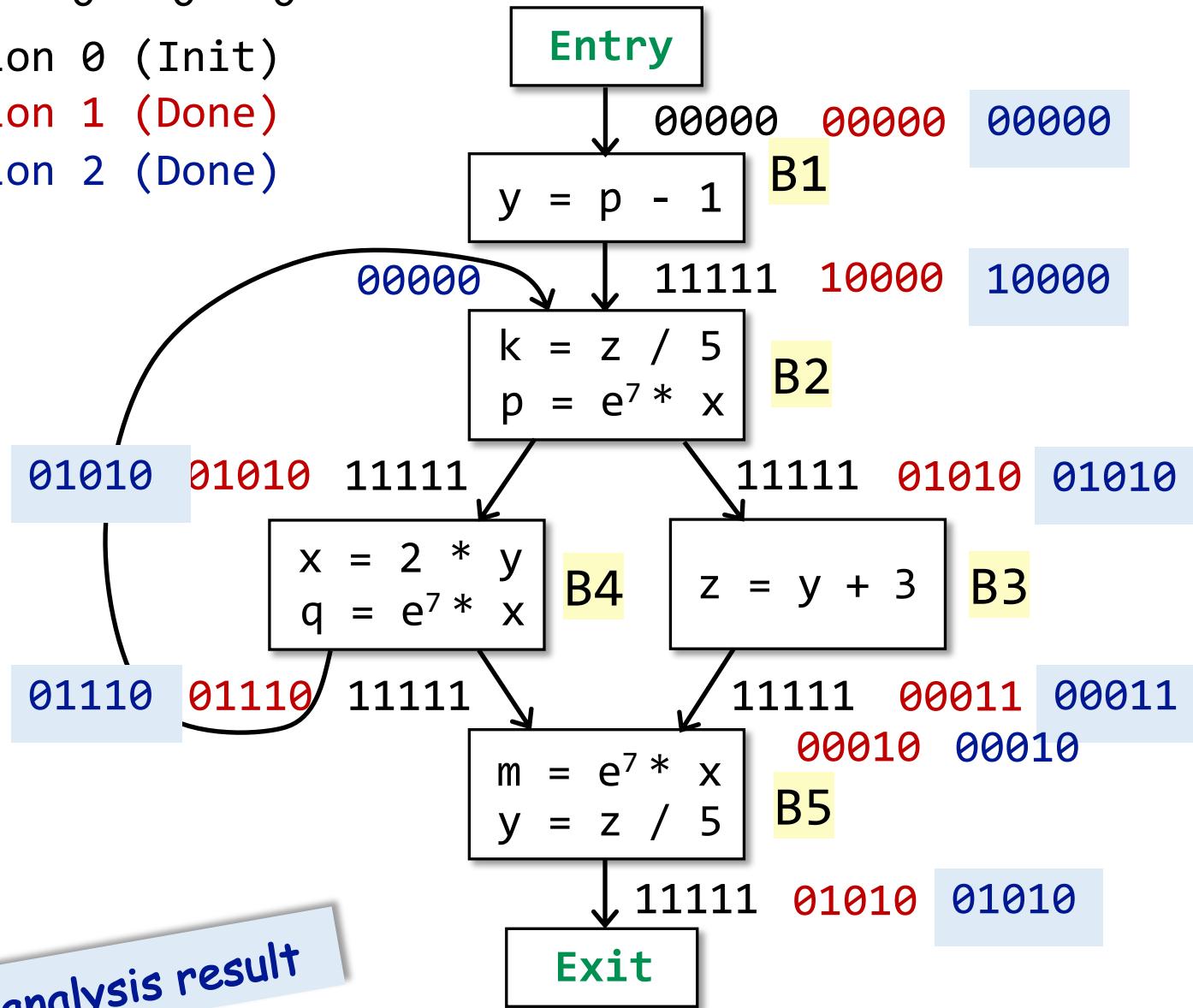


$p - 1$ $z / 5$ $2 * y$ $e^7 * x$ $y + 3$
 0 0 0 0 0

Iteration 0 (Init)

Iteration 1 (Done)

Iteration 2 (Done)



Analysis Comparison

	Reaching Definitions	Live Variables	Available Expressions
Domain			
Direction			
May/Must			
Boundary			
Initialization			
Transfer function			
Meet			

Analysis Comparison

	Reaching Definitions	Live Variables	Available Expressions
Domain	Set of definitions	Set of variables	Set of expressions
Direction			
May/Must			
Boundary			
Initialization			
Transfer function			
Meet			

Analysis Comparison

	Reaching Definitions	Live Variables	Available Expressions
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Direction	Forwards	Backwards	Forwards
May/Must			
Boundary			
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Transfer function			
Meet			

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May/Must	May	May	Must
Boundary			
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May/Must	May	May	Must
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Initialization			
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Transfer function			
Meet			

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Boundary	$\text{OUT}[\text{entry}] = \emptyset$	$\text{IN}[\text{exit}] = \emptyset$	$\text{OUT}[\text{entry}] = \emptyset$
Initialization	$\text{OUT}[B] = \emptyset$	$\text{IN}[B] = \emptyset$	$\text{OUT}[B] = U$
Transfer function		$\text{OUT}/\text{IN} = \text{gen } U \ (\text{IN}/\text{OUT} - \text{kill})$	
Meet			

Analysis Comparison

	Reaching Definitions	Live Variables	Available Expressions
Domain	Set of definitions	Set of variables	Set of expressions
Direction	Forwards	Backwards	Forwards
May/Must	May	May	Must
Boundary	$\text{OUT}[\text{entry}] = \emptyset$	$\text{IN}[\text{exit}] = \emptyset$	$\text{OUT}[\text{entry}] = \emptyset$
Initialization	$\text{OUT}[B] = \emptyset$	$\text{IN}[B] = \emptyset$	$\text{OUT}[B] = U$
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Meet	U	U	\cap

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Meet	U	U	\cap

Analysis Comparison

According to the meaning of the analysis
We'll draw a theoretical framework to systematically explain them in next lectures

	Reaching Definitions	Live Variables	Available Expressions
Domain	Set of definitions	Set of variables	Set of expressions
Direction	Forwards	Backwards	Forwards
May/Must	May	May	Must
Boundary	$\text{OUT}[\text{entry}] = \emptyset$	$\text{IN}[\text{exit}] = \emptyset$	$\text{OUT}[\text{entry}] = \emptyset$
Initialization	$\text{OUT}[B] = \emptyset$	$\text{IN}[B] = \emptyset$	$\text{OUT}[B] = U$
Transfer function		$\text{OUT}/\text{IN} = \text{gen } U \ (\text{IN}/\text{OUT} - \text{kill})$	
Meet	U	U	n

summary

1. Overview of Data Flow Analysis
2. Preliminaries of Data Flow Analysis
3. Reaching Definitions Analysis
4. Live Variables Analysis
5. Available Expressions Analysis

The X You Need To Understand in This Lecture

- Understand the three data flow analyses:
 - reaching definitions
 - live variables
 - available expressions
- Can tell the differences and similarities of the three data flow analyses
- Understand the iterative algorithm and can tell why it is able to terminate

注意注意！
划重点了！



Assignment One:
Live variable analysis and iterative solver

Evaluation Criteria

- Coding Assignments 40% <http://tai-e.pascal-lab.net>
- Final Exam 60%

OJ注意事项：

- 用学校邮箱注册
- 填写 Student ID
- Nickname 用真实姓名
- 抄袭：第一次发现，该次作业0分
 第二次发现，课程0分
- 不交：得5%诚信分
- 迟交：扣得分 $5\% \times D$, D为迟交天数
- 每帮助找出1个新的有效test case, 实验分数+1