

# Safety Annex for Architecture Analysis and Design Language

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**Abstract.** This paper describes a new methodology with tool support for model based safety analysis. It is implemented as a new Safety Annex for the Architecture Analysis and Design Language (AADL). The Safety Annex provides the ability to describe faults and faulty component behaviors in AADL models. In contrast to previous AADL-based approaches, the Safety Annex leverages a formal description of the nominal system behavior to propagate faults in the system. This approach ensures consistency with the rest of the system development process and simplifies the work of safety engineers. The language for describing faults is extensible and allows safety engineers to weave various types of faults into the nominal system model. The Safety Annex supports the injection of faults into component level outputs, and the resulting behavior of the system can be analyzed using model checking through the Assume-Guarantee Reasoning Environment (AGREE).

**Keywords:** Model-based systems engineering, fault analysis, safety engineering

## 1 Introduction

System safety analysis techniques are well-established and are a required activity in the development of safety-critical systems. While model based development methods are widely used in the aerospace industry, these methods are only recently being applied to system safety analysis. Model-based systems engineering (MBSE) methods and tools based on formal methods now permit system-level requirements to be specified and analyzed early in the development process [5, 6]. These tools can also be used to perform safety analysis based on the system architecture and initial functional decomposition. Design models can be integrated into the safety analysis process to help guarantee accurate and consistent results. This integration is especially important as the amount of safety-critical hardware and software in various domains has drastically increased due to the demand for greater autonomy, capability, and connectedness.

We have developed a Safety Annex for the Architecture Analysis and Design Language (AADL) [9] that provides the ability to reason about faults and faulty component

behaviors in AADL models. In the Safety Annex approach, we use formal assume-guarantee contracts to define the behavior of system components. The nominal model is then verified using the Assume Guarantee Reasoning Environment (AGREE) [6]. The Safety Annex provides a way to weave faults into the nominal system model and analyze the behavior of the system in the presence of faults. The Safety Annex also provides a library of common fault node definitions that is customizable to the needs of system and safety engineers.

There are other tools purpose-built for safety analysis, including AltaRica [14], smartIFlow [12] and xSAP [4]. These notations are separate from the system development model. Other tools extend existing system models, such as HiP-HOPS [1] and the AADL Error Model Annex, Version 2 (EMV2) [7]. EMV2 uses enumeration of faults in each component and explicit propagation of faulty behavior to perform error analysis. The required propagation relationships must be manually added to the system model and can become complex, leading to mistakes in the analysis.

In contrast, the Safety Annex supports model checking and quantitative reasoning by attaching behavioral faults to components and then using the normal behavioral propagation and proof mechanisms built into the AGREE AADL annex. This allows users to reason about the evolution of faults over time, and produce counterexamples demonstrating how component faults lead to failures. Our approach adapts the work of Joshi et. al in [13] to the AADL modeling language. More information on the approach is available in [15], and the tool and relevant documentation can be found at: <https://github.com/loonwerks/AMASE/>.

## **2 Preliminaries**

In this section, we describe the overall process flow and introduce some terminology associated with the certification context. We also describe the system architecture modeling and verification environment that we are using.

### **2.1 Safety Process Overview**

One of our goals is to transition the tools we have developed into use by the safety engineers who perform safety assessment on the avionics products. Therefore, we need to understand how the tools and the models will fit into the safety assessment and certification process.

### **2.2 Architecture Description and Design Contracts**

Pull AADL/AGREE background from previous papers to support points in the safety process

## **3 Detailed Approach**

### **3.1 Functionality**

In this section, we describe the main features and functionality of the Safety Annex. An AADL model of the nominal system behavior specifies the hardware and software com-

ponents of the system and their interconnections. This nominal model is then annotated with assume-guarantee contracts using the AGREE annex [6] for AADL. The nominal model requirements are verified using compositional verification techniques based on  $k$ -induction model checking [10].

Once the nominal model behavior is defined and verified, the Safety Annex can be used to specify possible faulty behaviors for each component. The faults are defined on each of the relevant components using a customizable library of fault nodes and the faults are assigned a probability of occurrence. A probability threshold is also defined at the system level. This extended model can be analyzed to verify the behavior of the system in the presence of faults. Verification of the nominal model with or without the fault model is controlled through the safety analysis option during AGREE verification.

To illustrate the syntax of the Safety Annex, we use an example based on the Wheel Brake System (WBS) described in [2] and used in our previous work [15]. The fault library contains commonly used fault node definitions. An example of a fault node is shown below:

```
node fail_to(val_in: real, alt_val: real, trigger: bool) returns (val_out: real);
let
  val_out = if (trigger) then alt_val else val_in;
tel;
```

The *fail\_to* node provides a way to inject a faulty input value. When the *trigger* condition is satisfied, the nominal component output value is overridden by the *fail\_to* failure value. In the WBS, the pump component generates an expected amount of pressure to a hydraulic line. Declaration of a non-deterministic fault in the pump component is shown below:

```
annex safety {**
  fault "In pump: pressure_output failed to non-deterministic value.": faults.fail_to {
    eq alt_value :real;
    inputs: val_in <- pressure_output,
            alt_val <- alt_value;
    outputs: pressure_output.val <- val_out ;
    probability: 1.0E-5 ;
    duration: permanent;
  }
**};
```

The *fault statement* consists of a unique description string, the fault node definition name, and a series of *fault subcomponent* statements.

**Inputs** in a fault statement are the parameters of the fault node definition. In the example above, *val\_in* and *alt\_val* are the two input parameters of the fault node. These are linked to the output from the Pump component (*pressure\_output.val*), and *alt\_value*, a nondeterministic value defined within the Safety Annex. When the analysis is run, these values are passed into the fault node definition.

**Outputs** of the fault definition correspond to the outputs of the fault node. The fault output statement links the component output (*pressure\_output.val*) with the fault node output (*val\_out*). If the fault is triggered, the nominal value of *pressure\_output.val* is overridden by the failure value output by the fault node. Faulty outputs can take deterministic or non-deterministic values.

**Probability** (optional) describes the probability of a fault occurrence.

**Duration** describes the duration of the fault; currently the Safety Annex supports transient and permanent faults.

To-add: HW fault and dependent fault:

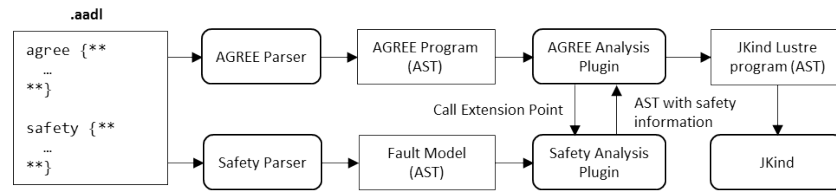
Fault propagations: it resides outside of fault statements and inside safety annex. This is because fault propagations are typically tied to the way components are connected or bound together; this information may not be available when faults are being specified for individual components. In addition, having fault propagations specified outside of a components fault statements makes it easier to reuse the component in different systems.

HW fault: no behavioral aspects.

Dependent fault modeling: With the introduction of dependent faults, active faults are divided into two categories: independently active (activated by its own triggering event) and dependently active (activated when the faults they depend on become active). Top level fault hypothesis (e.g., maximum number of faults that can be active, or probability threshold for occurrences of faults) now applies to independently active faults. Faulty behaviors augment nominal behaviors whenever their corresponding faults are active (either independently active or dependently active).

### 3.2 Architecture and Implementation


The architecture of the Safety Annex is shown in Figure 1. It is written in Java as a plug-in for the OSATE AADL toolset, which is built on Eclipse. It is not designed as a stand-alone extension of the language, but works with behavioral contracts specified AGREE AADL annex and associated tools [6]. AGREE allows *assume-guarantee* behavioral contracts to be added to AADL components. The language used for contract specification is based on the Lustre dataflow language [11]. AGREE improves scalability of formal verification to large systems by decomposing the analysis of a complex system architecture into a collection of smaller verification tasks that correspond to the structure of the architecture.



**Fig. 1.** Safety Annex Plug-in Architecture

AGREE contracts are used to define the nominal behaviors of system components as *guarantees* that hold when *assumptions* about the values the component's environment are met. The Safety Annex extends these contracts to allow faults to modify the behavior of component inputs and outputs. To support these extensions, AGREE implements an

Eclipse extension point interface that allows other plug-ins to modify the generated abstract syntax tree (AST) prior to its submission to the solver. If the Safety Annex is enabled, these faults are added to the AGREE contract and, when triggered, override the nominal guarantees provided by the component. An example of a portion of an initial AGREE node and its extended contract is shown in Figure 2. The `__fault` variables and declarations are added to allow the contract to override the nominal behavioral constraints (provided by guarantees) on outputs. In the Lustre language, assertions are constraints that are assumed to hold in the transition system.



```

agree node green_pump(
  time : real
) returns (
  pressure_output : common__pressure__i
);
let
  guarantees {
    "Pump always outputs something" :
      (pressure_output.val > 0.0)
  }
tel;

agree node green_pump(
  time : real;
  __fault__nominal__pressure_output : common__pressure__i;
  fault_trigger_green_pump_fault_22 : bool;
  green_pump_fault_22_alt_value : real
) returns (
  pressure_output : common__pressure__i
);
var
  green_pump_fault_22_node_val_out : common__pressure__i;
let
  assertions {
    (green_pump_fault_22_node_val_out = pressure_output)
  }
  guarantees {
    "Pump always outputs something" :
      (__fault__nominal__pressure_output.val > 0.0)
  }
  green_pump_fault_22_node_val_out =
    faults_fail_to(
      __fault__nominal__pressure_output,
      green_pump_fault_22_alt_value,
      fault_trigger_green_pump_fault_22);
tel;

```

**Fig. 2.** Nominal AGREE node and its extension with faults

An annotation in the AADL model determines the fault hypothesis. This may specify either a maximum number of faults that can be active at any point in execution (typically one or two), or that only faults whose probability of simultaneous occurrence is above some probability threshold should be considered. In the former case, we assert that the sum of the true *fault\_trigger* variables is below some integer threshold. In the latter, we determine all combinations of faults whose probabilities are above the specified probability threshold, and describe this as a proposition over *fault\_trigger* variables.

Once augmented with fault information, the AGREE model follows the standard AGREE translation path to the model checker JKind [10], an infinite-state model checker for safety properties. The augmentation includes traceability information so that when counterexamples are displayed to users, the active faults for each component are visualized.

## 4 Case Studies

To evaluate the effectiveness of the Safety Annex, we updated the WBS model [15] to specify faulty component behaviors. The components' nominal and faulty behaviors are modeled separately. At the top-level AADL component, the fault hypothesis was

specified as the maximum number of faults that can be active at any time. The AGREE contracts at the top-level component were verified using AGREE, with the “Perform Safety Analysis” option selected. This signals the tool to weave the nominal and faulty behaviors into one augmented AGREE model before feeding to the model checker.

In this example, the top level contract “Pedal pressed and no skid implies brake pressure applied” was verified in the presence of at most one fault active during execution. However, it was shown to be invalid when more than one fault was allowed. The counterexample indicated that both Selector’s outputs failed to non-deterministic values due to the faults introduced.

We also applied the Safety Annex to the Quad-Redundant Flight Control System (QFCS) model [3]. We introduced faulty behaviors to see the response of the system to several faults, and to evaluate fault mitigation logic in the model. The QFCS system-level properties failed when unhandled faulty behaviors were introduced.

We also used the Safety Annex to explore more complicated faults at the system level on a simplified QFCS model with cross-channel communication between its Flight Control Computers.

- Byzantine faults [8] were simulated by creating one-to-one connections from the source to multiple observers so that disagreements could be introduced by injecting faults on individual outputs. A system-level property failed due to the fault on the baseline model, but did not fail on the model with Byzantine fault handling protocol added. Using the Safety Annex like this can test a system’s vulnerability to Byzantine faults and verify mitigation mechanisms.
- Dependent faults in hardware were simulated by injecting faults to hardware components (physical layer) to affect their data outputs (logical layer), and consequently failing the software components bound to the hardware components. The relationship between the hardware components’ outputs and the software components’ inputs were specified in AGREE as part of the system’s nominal behavior.

#### **4.1 Safety Process Related Case Study**

value added to traditional safety assessment

#### **4.2 Scalability Related Case Study**

### **5 Related Work**

Address NFM review comments

### **6 Conclusion**

We have developed an extension to the AADL language with tool support for formal analysis of system safety properties in the presence of faults. Faulty behavior is specified as an extension of the nominal model, allowing safety analysis and system implementation to be driven from a single common model. This new Safety Annex leverages

the AADL structural model and nominal behavioral specification (using the AGREE annex) to propagate faulty component behaviors without the need to add separate propagation specifications to the model. Next steps will include extensions to automate injection of Byzantine faults as well as the ability to specify dependent faults.

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