

TOPIC

**1.PROPERTIES OF WELL FORMED
FORMULA**

2.CONVERSION TO CluASAL FORM

3.INFERENCE RULES



PROPERTIES OF WELL FORMED FORMULA

1. Propositional logic uses a symbolic “language” to represent the logical structure, or form, of a compound proposition.
2. Has rules of syntax- grammatical rules for putting symbols together in the right way.
3. Any expression that obeys the syntactic rules of propositional logic is called a WELL FORMED FORMULA, or WFF



Definition

A well formed formula is any formula that is capable of being generated by some combination of seven information rules.

- ❖ Every propositional letter of PL (e.g. A, B, C) is a WFF.
- ❖ If P is a WFF, then $\neg(P)$ is a WFF.
- ❖ If P and Q are WFF, then $(P \wedge Q)$ is a WFF.
- ❖ If P and Q are WFF, then $(P \vee Q)$ is a WFF.
- ❖ If P and Q are WFF, then $(P \supset Q)$ is a WFF.
- ❖ If P and Q are WFF, then $(P \leftrightarrow Q)$ is a WFF.
- ❖ Nothing else is a wff except what can be formed by repeated application of 1-6.



Example(WFF)

- P
- $(P \rightarrow Q)$
- $\neg((P \rightarrow \neg(Q)))$
- $\neg(\neg(P \rightarrow (Q)))$

Example(not WFF)

- $P \wedge$
- PQ
- $\neg \neg(Q))$
- $\wedge \wedge P \neg \neg Q \rightarrow S$



Three type of WFF: Atomic WFFs

Definition

An Atomic wff in PL is a well formed formula in PL consisting only of a single propositional letter.

Example:

$>P$

$>Q$

$>A$

$>B$



Complex WFFs

A complex wff in PL is a well formed formula in PL that contains at least one propositional letter and a truth-functional operator.

Example:

$\neg(P)$

$(P \wedge Q)$

$(P \supset (A))$



Literal Wff

A literal wff in PL is a well formed formula in PL that consist of an atomic wff P or a negated atomic wff $\neg(P)$.

Example:

$\Box p$

$\Box \neg(P)$

$\Box Q$

$\Box \neg(Q)$



Conversion to clausal form /normal form

Firstly :

1. PL does not permit us to make generalized statement about classes of similar object.
2. These are serious limitation when reasoning about real world entity.
3. E.g. statement= \Rightarrow it should be possible to conclude that john must take the pascal course.

ALL STUDENT IN COMPUTER SCIENCE MUST TAKE PASCAL

JOHN IS A COMPUTER SCIENCE MAJOR



SYNTAX OF FOPL

1. CONNECTIVES: These are 5 connective symbols \Rightarrow (not or
negation \neg)
 \wedge (AND or
conjunction \wedge)
 \vee (OR or
disjunction \vee)
 \Rightarrow (implication \Rightarrow)
 \Leftrightarrow (equivalence or if and
only if \Leftrightarrow)

2. QUANTIFIERS: 2 quantifier symbols are
 \exists (existential quantifier)
 \forall (universal quantifier)

3. CONSTANT: constant are fixed value terms that belong to a



4. VARIABLE: variable are terms that assume different value over a given domain.

5. FUNCTION: function symbol denote relations defined on a domain D .

6. PREDICATE: predicate symbols denote relations or functional mapping from the elements of a domains D to the value true and false.

capital letter and capitalized word such as $P, Q, R, \text{EQUAL}, \text{MARRIED}$ are used to represent predicate.



Symbolic formation

E1: All employee earning \$1400 or more per year pay taxes.

E2: Some employee are sick today.

E3: no employee earns more than the president.

Predicate and functions:

$E(x)$ for x is an employee.

$P(x)$ for x is president.

$i(x)$ for income of x . (lower case denotes a function).

$GE(u,v)$ for u is greater than or equal to v .



Using abbreviation E1, E2, E3

E1': 

E2': 

E3': 

E1' E2' E3' (woof)



Clausal form

- 1.Clause:** as the disjunction of a number of literals.
- 2.Ground clause:** is one in which no variable occur in the expression.
- 3.Horn clause:** is a clause with a most one positive literal.



How to convert PL to

Conjunctive normal form

$(A \vee B) \wedge (A \vee \neg B)$

Disjunctive normal form

$(A \wedge B) \vee (A \wedge \neg B)$

1. Remove \neg
2. Remove \wedge
3. Move negation(\neg)
inward/ apply
demorgan's law
4. Use distributive law



Remove 

A

▲A▶B

Remove 

A

(A)◀(B

Convert in CNF 

 A B

 A▶B

▲ A▶B C

▲ A)◀▲B▶C

 A◀▲B C

 A▶C  B▶C



CONVERT INTO DNF

$\neg(A \vee B) \vee (\neg S \vee T)$

$(\neg A \vee \neg B) \vee (\neg S \vee T)$

$(\neg A \vee \neg B) \vee (\neg S) \vee (\neg A \vee \neg B) \vee T$



INFERENCE RULES

INFERENCE: Deriving conclusion from evidences

RULES OF INFERENCE: Template for constructing valid argument

Types

1. Modus ponens

~~P~~ ~~Q~~

P

Q



2. MODUS TOLLEN:

$P \supset Q$

$\neg Q$

$\neg P$

$[(P \supset Q) \wedge \neg Q] \supset \neg P$

3. HYPOTHETICAL SYLLOGISM

$P \supset Q$

$Q \supset R$

$P \supset R$

$[(P \supset Q) \wedge (Q \supset R)] \supset P \supset R$



4. Disjunctive syllogism

$P \vee Q$

$\neg P$

Q

$[(P \vee Q) \wedge \neg P] \rightarrow Q$

5. ADDITION

P

$P \vee Q$

$P \rightarrow (P \vee Q)$



6. RESOLUTION

$P \supset Q$

$P \supset R$

$Q \supset R$

$[(P \supset Q) \wedge (\neg P \supset R)] \equiv (Q \supset R)$

