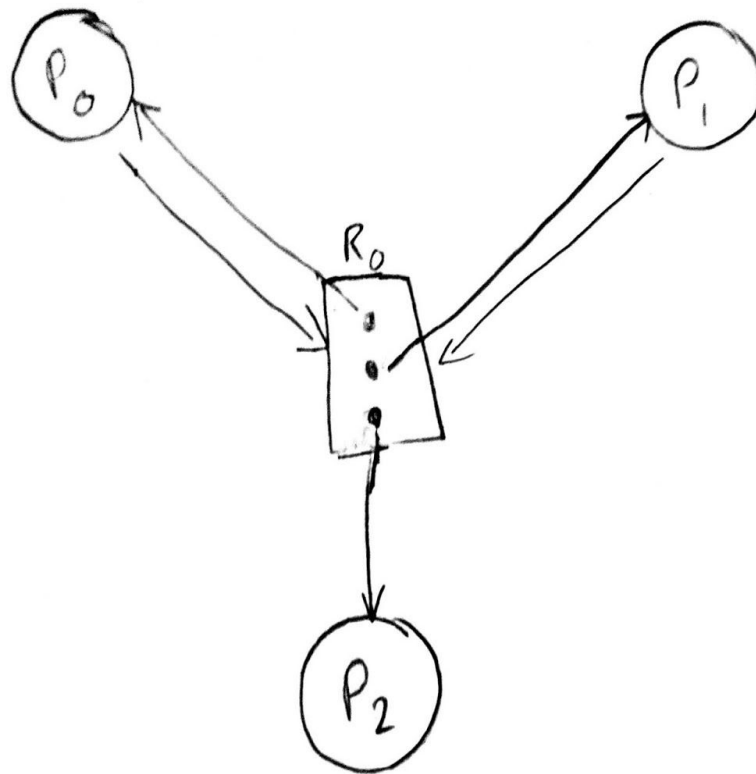


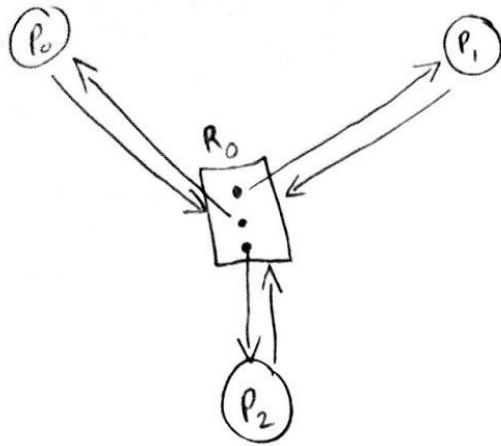
Q1

Unsafe state that finishes.



Order of execution:  $P_2, P_1, P_0$

Unsafe state that leads to deadlock.



Order of execution: There is a deadlock immediately since all three processes are waiting on each other to release a resource that they won't be able to get since  $R_0$  has been already fully utilized.  
Sequence:

$P_0$  requests  $R_0$  - blocked

$P_1$  requests  $R_0$  - blocked

$P_2$  requests  $R_0$  - blocked

Q2

<u>Available:</u>		<u>Need:</u>
00112	P3 releases	✓ P <sub>0</sub> 01002
+ 11110		✓ P <sub>1</sub> 02100
11222	P0 releases	✓ P <sub>2</sub> 10300
+ 10211		✓ P <sub>3</sub> 00111
21433	P <sub>2</sub> releases	
+ 11010		
32443		P <sub>3</sub> , P <sub>0</sub> , P <sub>2</sub> , P <sub>1</sub>
+		

The lowest value  $x$  can be is 1 such that the system is in a safe state. What I did was apply the banker's algorithm by first calculating the need matrix as shown above. Then every time a process finished I updated the available vector by adding all the resources a process releases. By doing the algorithm with  $x=1$ , the safe execution is P<sub>3</sub>, P<sub>0</sub>, P<sub>2</sub>, and finally P<sub>1</sub>.